CHAPTER 1 The Foundations: Logic and Proofs

SECTION 1.1 Propositional Logic

Manipulating propositions and constructing truth tables are straightforward. A truth table is constructed by finding the truth values of compound propositions from the inside out; see the solution to Exercise 27, for instance. This exercise set also introduces fuzzy logic.

- Propositions must have clearly defined truth values, so a proposition must be a declarative sentence with no free variables.
 - a) This is a true proposition.
 - b) This is a false proposition (Tallahassee is the capital).
 - c) This is a true proposition.
 - d) This is a false proposition.
 - e) This is not a proposition (it contains a variable; the truth value depends on the value assigned to x).
 - f) This is not a proposition, since it does not assert anything.
- 3. a) Today is not Thursday.
 - b) There is pollution in New Jersey.
 - c) $2 + 1 \neq 3$
 - d) It is not the case that the summer in Maine is hot and sunny. In other words, the summer in Maine is not hot and sunny, which means that it is not hot or it is not sunny. It is not correct to negate this by saying "The summer in Maine is not hot and not sunny." [For this part (and in a similar vein for part (b)) we need to assume that there are well-defined notions of hot and snnny; otherwise this would not be a proposition because of not having a definite truth value.]
- 5. This is pretty straightforward, using the normal words for the logical operators.
 - a) Sharks have not been spotted near the shore.
 - b) Swimming at the New Jersey shore is allowed, and sharks have been spotted near the shore.
 - c) Swimming at the New Jersey shore is not allowed, or sharks have been spotted near the shore.
 - d) If swimming at the New Jersey shore is allowed, then sharks have not been spotted near the shore.
 - e) If sharks have not been spotted near the shore, then swimming at the New Jersey shore is allowed.

 f) If swimming at the New Jersey shore is not allowed, then sharks have not been spotted near the shore.
 - g) Swimming at the New Jersey shore is allowed if and only if sharks have not been spotted near the shore.
 - h) Swimming at the New Jersey shore is not allowed, and either swimming at the New Jersey shore is allowed or sharks have not been spotted near the shore. Note that we were able to incorporate the parentheses by using the word "either" in the second half of the sentence.
- 7. a) Here we have the conjunction $p \wedge q$.
 - b) Here we have a conjunction of p with the negation of q, namely $p \land \neg q$. Note that "but" logically means the same thing as "and."

- c) Again this is a conjunction: $\neg p \land \neg q$.
- d) Here we have a disjunction, $p \vee q$. Note that \vee is the inclusive or, so the "(or both)" part of the English sentence is automatically included.
- e) This sentence is a conditional statement, $p \to q$.
- f) This is a conjunction of propositions, both of which are compound: $(p \lor q) \land (p \to \neg q)$.
- g) This is the biconditional $p \leftrightarrow q$.
- **9.** a) This is just the negation of p, so we write $\neg p$.
 - b) This is a conjunction ("but" means "and"): $p \wedge \neg q$.
 - c) The position of the word "if" tells us which is the antecedent and which is the consequence: $p \to q$.
 - d) $\neg p \rightarrow \neg q$
 - e) The sufficient condition is the antecedent: $p \rightarrow q$.
 - f) $q \wedge \neg p$
 - g) "Whenever" means "if": $q \rightarrow p$.
- 11. a) "But" is a logical synonym for "and" (although it often suggests that the second part of the sentence is likely to be unexpected). So this is $r \wedge \neg p$.
 - b) Because of the agreement about precedence, we do not need parentheses in this expression: $\neg p \land q \land r$.
 - c) The outermost structure here is the conditional statement, and the conclusion part of the conditional statement is itself a biconditional: $r \to (q \leftrightarrow \neg p)$.
 - d) This is similar to part (b): $\neg q \land \neg p \land r$.
 - e) This one is a little tricky. The statement that the condition is necessary is a conditional statement in one direction, and the statement that this condition is not sufficient is the negation of the conditional statement in the other direction. Thus we have the structure (safe \rightarrow conditions) $\land \neg$ (conditions \rightarrow safe). Fleshing this out gives our answer: $(q \rightarrow (\neg r \land \neg p)) \land \neg ((\neg r \land \neg p) \rightarrow q)$. There are some logically equivalent correct answers as well.
 - f) We just need to remember that "whenever" means "if" in logic: $(p \wedge r) \rightarrow \neg q$.
- 13. In each case, we simply need to determine the truth value of the hypothesis and the conclusion, and then use the definition of the truth value of the conditional statement. The conditional statement is true in every case except when the hypothesis (the "if" part) is true and the conclusion (the "then" part) is false.
 - a) Since the hypothesis is true and the conclusion is false, this conditional statement is false.
 - b) Since the hypothesis is false and the conclusion is true, this conditional statement is true.
 - c) Since the hypothesis is false and the conclusion is false, this conditional statement is true. Note that the conditional statement is false in both part (b) and part (c); as long as the hypothesis is false, we need look no further to conclude that the conditional statement is true.
 - d) Since the hypothesis is false, this conditional statement is true.
- 15. a) Presumably the diner gets to choose only one of these beverages, so this is an exclusive or.
 - b) This is probably meant to be inclusive, so that long passwords with many digits are acceptable.
 - c) This is surely meant to be inclusive. If a student has had both of the prerequisites, so much the better.
 - d) At first glance one might argue that no one would pay with both currencies simultaneously, so it would seem reasonable to call this an exclusive *or*. There certainly could be cases, however, in which the patron would pay a portion of the bill in dollars and the remainder in euros. Therefore, an inclusive *or* seems better.
- 17. a) If this is an inclusive or, then it is allowable to take discrete mathematics if you have had calculus or computer science or both. If this is an exclusive or, then a person who had both courses would not be allowed

to take discrete mathematics—only someone who had taken exactly one of the prerequisites would be allowed in. Clearly the former interpretation is intended; if anything, the person who has had both calculus and computer science is even better prepared for discrete mathematics.

- b) If this is an inclusive or, then you can take the rebate, or you can sign up for the low-interest loan, or you can demand both of these incentives. If this is an exclusive or, then you will receive one of the incentives but not both. Since both of these deals are expensive for the dealer or manufacturer, surely the exclusive or was intended.
- c) If this is an inclusive or, you can order two items from column A (and none from B), or three items from column B (and none from A), or five items (two from A and three from B). If this is an exclusive or, which it surely is here, then you get your choice of the two A items or the three B items, but not both.
- d) If this is an inclusive or, then lots of snow, or extreme cold, or a combination of the two will close school. If this is an exclusive or, then one form of bad weather would close school but if both of them happened then school would meet. This latter interpretation is clearly absurd, so the inclusive or is intended.
- 19. a) If the wind blows from the northeast, then it snows. ["Whenever" means "if."]
 - b) If it stays warm for a week, then the apple trees will bloom. [Sometimes word order is flexible in English, as it is here. Other times it is not—"The man bit the dog" does not have the same meaning as "The dog bit the man."]
 - c) If the Pistons win the championship, then they beat the Lakers.
 - d) If you get to the top of Long's Peak, then you must have walked eight miles. [The necessary condition is the conclusion.]
 - e) If you are world famous, then you will get tenure as a professor. [The sufficient condition is the antecedent.]
 - f) If you drive more than 400 miles, then you will need to buy gasoline. [The word "then" is sometimes omitted in English sentences, but it is still understood.]
 - g) If your guarantee is good, then you must have bought your CD player less than 90 days ago. [Note that "only if" does not mean "if"; the clause following the "only if" is the conclusion, not the antecedent.]
 - h) If the water is not too cold, then Jan will go swimming. [Note that "unless" really means "if not." It also can be taken to mean "or."]
- 21. In each case there will be two statements. It is being asserted that the first one holds true if and only if the second one does. The order doesn't matter, but often one order is more colloquial English.
 - a) You buy an ice cream cone if and only if it is hot outside.
 - b) You win the contest if and only if you hold the only winning ticket.
 - c) You get promoted if and only if you have connections.
 - d) Your mind will decay if and only if you watch television.
 - e) The train runs late if and only if it is a day I take the train.
- 23. Many forms of the answers for this exercise are possible.
 - a) One form of the converse that reads well in English is "I will ski tomorrow only if it snows today." We could state the contrapositive as "If I don't ski tomorrow, then it will not have snowed today." The inverse is "If it does not snow today, then I will not ski tomorrow."
 - b) The proposition as stated can be rendered "If there is going to be a quiz, then I will come to class." The converse is "If I come to class, then there will be a quiz." (Or, perhaps even better, "I come to class only if there's going to be a quiz.") The contrapositive is "If I don't come to class, then there won't be a quiz." The inverse is "If there is not going to be a quiz, then I don't come to class."
 - c) There is a variable ("a positive integer") in this sentence, so technically it is not a proposition. Nevertheless, we can treat sentences such as this in the same way we treat propositions. Its converse is "A positive integer

is a prime if it has no divisors other than 1 and itself." (Note that this can be false, since the number 1 satisfies the hypothesis but not the conclusion.) The contrapositive of the original proposition is "If a positive integer has a divisor other than 1 and itself, then it is not prime." (We are simplifying a bit here, replacing "does not have no divisors" by "has a divisor." Note that this is always true, assuming that we are talking about positive divisors.) The inverse is "If a positive integer is not prime, then it has a divisor other than 1 and itself."

- **25.** A truth table will need 2^n rows if there are n variables.
 - a) $2^1 = 2$
- **b)** $2^4 = 16$
- c) $2^6 = 64$
- **d**) $2^4 = 16$
- 27. To construct the truth table for a compound proposition, we work from the inside out. In each case, we will show the intermediate steps. In part (d), for example, we first construct the truth table for $p \vee q$, then the truth table for $p \wedge q$, and finally combine them to get the truth table for $(p \vee q) \rightarrow (p \wedge q)$. For parts (a) and (b) we have the following table (column three for part (a), column four for part (b)).

p	$\neg p$	$p \land \neg p$	$p \lor \neg p$
T	\mathbf{F}	F	T
\mathbf{F}	${ m T}$	\mathbf{F}	T

For part (c) we have the following table.

\underline{p}	q	$\neg q$	$p \lor \neg q$	$(p \lor \neg q) \rightarrow q$
T	Τ	\mathbf{F}	T	T
T	F	T	T	F
\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	T
\mathbf{F}	\mathbf{F}	${f T}$	${ m T}$	\mathbf{F}

For part (d) we have the following table.

p	q	$\underline{p \lor q}$	$p \wedge q$	$(p \lor q) \to (p \land q)$
Ţ	Τ	T	${ m T}$	${f T}$
\mathbf{T}	F	\mathbf{T}	\mathbf{F}	F
\mathbf{F}	T	T	\mathbf{F}	F
\mathbf{F}	F	F	\mathbf{F}	T

For part (e) we have the following table. This time we have omitted the column explicitly showing the negations of p and q. Note that this true proposition is telling us that a conditional statement and its contrapositive always have the same truth value.

p	q	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$(p \to q) \leftrightarrow (\neg q \to \neg p)$
T	${\rm T}$	${ m T}$	${f T}$	${f T}$
T	\mathbf{F}	\mathbf{F}	F	${f T}$
\mathbf{F}	\mathbf{T}	${f T}$	${ m T}$	${f T}$
F	\mathbf{F}	\mathbf{T}	${f T}$	${ m T}$

For part (f) we have the following table. The fact that this proposition is not always true tells us that knowing a conditional statement in one direction does not tell us that the conditional statement is true in the other direction.

p	\underline{q}	$p \rightarrow q$	$q \rightarrow p$	$(p \to q) \to (q \to p)$
\mathbf{T}	\mathbf{T}	${ m T}$	${ m T}$	T
${ m T}$	\mathbf{F}	\mathbf{F}	${f T}$	T
\mathbf{F}	\mathbf{T}	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	F	T	${f T}$	T

29. To construct the truth table for a compound proposition, we work from the inside out. In each case, we will show the intermediate steps. In part (a), for example, we first construct the truth table for $p \lor q$, then the

truth table for $p \oplus q$, and finally combine them to get the truth table for $(p \vee q) \to (p \oplus q)$. For parts (a), (b), and (c) we have the following table (column five for part (a), column seven for part (b), column eight for part (c)).

p	q	$p \lor q$	$p \oplus q$	$\underline{(p\vee q)\ \rightarrow\ (p\oplus q)}$	$p \wedge q$	$(p \oplus q) \to (p \land q)$	$(p \vee q) \oplus (p \wedge q)$
Τ	Τ	Т	F	F	${ m T}$	T	\mathbf{F}
Τ	F	T	T	T	\mathbf{F}	F	T
\mathbf{F}	Τ	Т	T	T	F	F	Τ
F	F	F	F	T	F	T	F

For part (d) we have the following table.

p	q	$\neg p$	$p \leftrightarrow q$	$\neg p \leftrightarrow q$	$\underline{(p \leftrightarrow q) \oplus (\neg p \leftrightarrow q)}$
Τ	T	\mathbf{F}	${f T}$	F	${ m T}$
Τ	F	F	F	T	${ m T}$
F	Τ	Τ	\mathbf{F}	${ m T}$	T
F	F	T	${f T}$	F	T

For part (e) we need eight rows in our truth table, because we have three variables.

p	q	7.	$\underline{\neg p}$	$\neg r$	$p \leftrightarrow q$	$\neg p \leftrightarrow \neg r$	$(p \leftrightarrow q) \oplus (\neg p \leftrightarrow \neg r)$
\mathbf{T}	T	Τ	\mathbf{F}	\mathbf{F}	${ m T}$	${ m T}$	\mathbf{F}
T	Τ	F	F	${ m T}$	T	F	${f T}$
T	F	Τ	F	\mathbf{F}	\mathbf{F}	${ m T}$	${f T}$
T	F	\mathbf{F}	F	\mathbf{T}	F	F	\mathbf{F}
F	Τ	Τ	T	\mathbf{F}	F	\mathbf{F}	F
F	T	F	T	T	F	T	${f T}$
\mathbf{F}	\mathbf{F}	Τ	T	F	${f T}$	F	T
F	\mathbf{F}	\mathbf{F}	${ m T}$	T	${ m T}$	T	\mathbf{F}

For part (f) we have the following table.

\underline{p}	q	$\neg q$	$p \oplus q$	$p \oplus \neg q$	$(p \oplus q) \to (p \oplus \neg q)$
\mathbf{T}	Τ	F	F	\mathbf{T}	T
Τ	F	\mathbf{T}	\mathbf{T}	F	F
F	\mathbf{T}	F	${ m T}$	F	F
F	F	\mathbf{T}	F	${ m T}$	Т

31. The techniques are the same as in Exercises 27-30. For parts (a) and (b) we have the following table (column four for part (a), column six for part (b)).

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p	\underline{q}	$\overline{\neg q}$	$p \rightarrow \neg q$	$\neg p$	$\neg p \leftrightarrow q$
	Τ	\mathbf{F}	F	F	F
Τ	F	${ m T}$	T	\mathbf{F}	T
F	Τ	\mathbf{F}	T	${f T}$	T
F	\mathbf{F}	${ m T}$	T	T	\mathbf{F}

For parts (c) and (d) we have the following table (columns six and seven, respectively).

\boldsymbol{p}	q	$p \rightarrow q$	$\underline{\neg p}$	$\neg p \rightarrow q$	$(p \to q) \lor (\neg p \to q)$	$(p \to q) \land (\neg p \to q)$
${f T}$	\mathbf{T}	T	F	${ m T}$	T	T
${f T}$	\mathbf{F}	\mathbf{F}	F	T	T	\mathbf{F}
\mathbf{F}	Τ	T	Τ	${f T}$	${f T}$	${ m T}$
\mathbf{F}	\mathbf{F}	${ m T}$	Τ	\mathbf{F}	${f T}$	F

For parts (e) and (f) we have the following table (this time we have not explicitly shown the columns for negation). Column five shows the answer for part (e), and column seven shows the answer for part (f).

p	q	$p \leftrightarrow q$	$\neg p \leftrightarrow q$	$(p \leftrightarrow q) \vee (\neg p \leftrightarrow q)$	$\neg p \leftrightarrow \neg q$	$(\neg p \leftrightarrow \neg q) \leftrightarrow (p \leftrightarrow q)$
\mathbf{T}	Τ	${f T}$	\mathbf{F}	${f T}$	${ m T}$	T
T	F	F	${ m T}$	T	F	T
F	T	F	${ m T}$	T	\mathbf{F}	T
F	F	T	\mathbf{F}	T	\mathbf{T}	T

33. The techniques are the same as in Exercises 27-32, except that there are now three variables and therefore eight rows. For part (a), we have

\underline{p}	\underline{q}	7	$\neg q$	$\neg q \lor r$	$p \to (\neg q \lor r)$
Τ	Τ	Τ	F	T	T
T	Τ	F	\mathbf{F}	\mathbf{F}	F
\mathbf{T}	F	\mathbf{T}	${ m T}$	${ m T}$	T
\mathbf{T}	F	F	T	T	${ m T}$
F	Τ	Τ	F	${ m T}$	${ m T}$
F	T	F	F	F	T
F	F	T	\mathbf{T}	${ m T}$	${ m T}$
F	\mathbf{F}	F	T	T	${ m T}$

For part (b), we have

p	q	r	$\neg p$	$q \rightarrow r$	$\neg p \rightarrow (q \rightarrow r)$
\mathbf{T}	Τ	T	F	T	T
${ m T}$	T	F	F	F	${ m T}$
T	F	Τ	F	${f T}$	T
${\rm T}$	F	F	F	${ m T}$	${ m T}$
F	T	T	${ m T}$	${ m T}$	${ m T}$
F	\mathbf{T}	F	T	\mathbf{F}	F
F	F	T	\mathbf{T}	T	T
F	F	F	T	${ m T}$	${ m T}$

Parts (c) and (d) we can combine into a single table.

\underline{p}	$q_{_}$	\underline{r}	$p \longrightarrow q$	$\underline{\neg p}$	$\underline{\neg p \rightarrow r}$	$(p \to q) \lor (\neg p \to r)$	$(p \to q) \land (\neg p \to r)$
\mathbf{T}	T	T	${f T}$	F	T	${f T}$	T
Τ	Т	F	${ m T}$	F	T	${ m T}$	T
Τ	F	\mathbf{T}	\mathbf{F}	F	T	T	F
T	F	F	\mathbf{F}	\mathbf{F}	${f T}$	${f T}$	F
F	T	T	T	\mathbf{T}	T	T	T
F	T	F	${ m T}$	T	F	${ m T}$	F
F	F	\mathbf{T}	${f T}$	T	T	T	T
E	F	F	т	Υ	F	Т	F

For part (e) we have

p q	\underline{r}	$p \leftrightarrow q$	$\neg q$	$\neg q \leftrightarrow r$	$\underline{(p \leftrightarrow q) \lor (\neg q \leftrightarrow r)}$
TT	T	T	F	F	${ m T}$
T T	F	${f T}$	\mathbf{F}	${ m T}$	T
TF	T	\mathbf{F}	${ m T}$	${ m T}$	${f T}$
T F	F	\mathbf{F}	${f T}$	\mathbf{F}	\mathbf{F}
F T	T	F	\mathbf{F}	F	F
F T	F	F	F	${ m T}$	T
\mathbf{F} \mathbf{F}	T	T	T	T	T
\mathbf{F} \mathbf{F}	F	${ m T}$	${f T}$	F	${f T}$

Finally, for part (f) we have

\underline{p}	q	<u>r</u>	$\underline{\neg p}$	$\neg q$	$\neg p \leftrightarrow \neg q$	$q \leftrightarrow r$	$(\neg p \leftrightarrow \neg q) \leftrightarrow (q \leftrightarrow r)$
\mathbf{T}	Τ	T	\mathbf{F}	\mathbf{F}	${f T}$	${f T}$	${ m T}$
\mathbf{T}	\mathbf{T}	F	\mathbf{F}	F	Ť	F	\mathbf{F}
\mathbf{T}	\mathbf{F}	\mathbf{T}	\mathbf{F}	${f T}$	\mathbf{F}	\mathbf{F}	T
T	\mathbf{F}	F	\mathbf{F}	${ m T}$	\mathbf{F}	${ m T}$	\mathbf{F}
\mathbf{F}	Τ	Τ	${ m T}$	F	\mathbf{F}	T	\mathbf{F}
F	Τ	F	T	\mathbf{F}	\mathbf{F}	\mathbf{F}	${f T}$
F	F	Τ	${ m T}$	${ m T}$	${f T}$	\mathbf{F}	F
\mathbf{F}	\mathbf{F}	F	${ m T}$	${ m T}$	T	${ m T}$	${ m T}$

35. This time the truth table needs $2^4 = 16$ rows. Note the systematic order in which we list the possibilities.

$p_{\underline{}}$	q	$r_{\underline{}}$	s	$p \leftrightarrow q$	$r \leftrightarrow s$	$(p \leftrightarrow q) \leftrightarrow (r \leftrightarrow s)$
Τ	\mathbf{T}	Τ	\mathbf{T}	${f T}$	T	${ m T}$
Τ	\mathbf{T}	\mathbf{T}	F	T	F	\mathbf{F}
Τ	Τ	\mathbf{F}	Τ	T	\mathbf{F}	F
${\rm T}$	\mathbf{T}	\mathbf{F}	F	${ m T}$	${ m T}$	${ m T}$
Τ	F	\mathbf{T}	\mathbf{T}	F	Ť	F
Τ	\mathbf{F}	Τ	\mathbf{F}	F	F	${ m T}$
\mathbf{T}	\mathbf{F}	F	Τ	F	F	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	F	\mathbf{F}	T	\mathbf{F}
F	T	\mathbf{T}	Τ	F	${ m T}$	F
F	Τ	Τ	F	F	F	Ť
\mathbf{F}	\mathbf{T}	\mathbf{F}	Τ	F	F	T
\mathbf{F}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}	${ m T}$	\mathbf{F}
\mathbf{F}	\mathbf{F}	Τ	Τ	${ m T}$	Ť	${ m T}$
\mathbf{F}	\mathbf{F}	Τ	F	T	F	F
F	\mathbf{F}	\mathbf{F}	Τ	${ m T}$	\mathbf{F}	\mathbf{F}
F	\mathbf{F}	F	F	${f T}$	${ m T}$	${ m T}$

- **37.** a) bitwise $OR = 111\ 1111$; bitwise $AND = 000\ 0000$; bitwise $XOR = 111\ 1111$
 - b) bitwise $OR = 1111\ 1010$; bitwise $AND = 1010\ 0000$; bitwise $XOR = 0101\ 1010$
 - c) bitwise $OR = 10\ 0111\ 1001$; bitwise $AND = 00\ 0100\ 0000$; bitwise $XOR = 10\ 0011\ 1001$
 - d) bitwise $QR = 11 \ 1111 \ 1111$; bitwise $AND = 00 \ 0000 \ 0000$; bitwise $XOR = 11 \ 1111 \ 1111$
- **39.** For "Fred is not happy," the truth value is 1 0.8 = 0.2.

For "John is not bappy," the truth value is 1 - 0.4 = 0.6.

41. For "Fred is happy, or John is happy," the truth value is max(0.8, 0.4) = 0.8.

For "Fred is not happy, or John is not happy," the truth value is $\max(0.2, 0.6) = 0.6$ (using the result of Exercise 39).

- 43. One great problem-solving strategy to try with problems like this, when the parameter is large (100 statements here) is to lower the parameter. Look at a simpler problem, with just two or three statements, and see if you can figure out what's going on. That was the approach used to discover the solution presented here.
 - a) Some number of these statements are true, so in fact exactly one of the statements must be true and the other 99 of them must be false. That is what the 99th statement is saying, so it is true and the rest are false.
 - b) The 100th statement cannot be true, since it is asserting that all the statements are false. Therefore it must be false. That makes the first statement true. Now if the 99th statement were true, then we would conclude that statements 2 through 100 were false, which contradicts the truth of statement 99. So statement 99 must be false. That means that statement 2 is true. We continue in this way and conclude that statements 1 through 50 are all true and statements 51 through 100 are all false.

- c) If there are an odd number of statements, then we'd run into a contradiction when we got to the middle. If there were just three statements, for example, then statement 3 would have to be false, making statement 1 true, and now the truth of statement 2 would imply its falsity and its falsity would imply its truth. Therefore this situation cannot occur with three (or any odd number of) statements. It is a logical paradox, showing that in fact these are not statements after all.
- 45. There are many correct answers to this problem, but all involve some sort of double layering, or combining a question about the kind of person being addressed with a question about the information being sought. One solution is to ask this question: "If I were to ask you whether the right branch leads to the ruins, would you say 'yes'?" If the villager is a truth-teller, then of course he will reply "yes" if and only if the right branch leads to the ruins. Now let us see what the liar says. If the right branch leads to the ruins, then he would say "no" if asked whether the right branch leads to the ruins. Therefore, the truthful answer to your convoluted question is "no." Since he always lies, he will reply "yes." On the other hand, if the right branch does not lead to the ruins, then he would say "yes" if asked whether the right branch leads to the ruins; and so the truthful answer to your question is "yes"; therefore he will reply "no." Note that in both cases, he gives the same answer to your question as the truth-teller; namely, he says "yes" if and only if the right branch leads to the ruins. A more detailed discussion can be found in Martin Gardner's Scientific American Book of Mathematical Puzzles and Diversions (Simon and Schuster, 1959), p. 25; reprinted as Hexaflexagons and Other Mathematical Diversions: The First Scientific American Book of Puzzles and Games (University of Chicago Press, 1988).
- 47. a) Since "whenever" means "if," we have $q \to p$.
 - **b)** Since "but" means "and," we have $q \land \neg p$.
 - c) This sentence is saying the same thing as the sentence in part (a), so the answer is the same: $q \to p$.
 - d) Again, we recall that "when" means "if" in logic: $\neg q \rightarrow \neg p$.
- **49.** Let m, n, k, and i represent the propositions "The system is in multiuser state," "The system is operating normally," "The kernel is functioning," and "The system is in interrupt mode," respectively. Then we want to make the following expressions simultaneously true by our choice of truth values for m, n, k, and i:

$$m \leftrightarrow n$$
, $n \rightarrow k$, $\neg k \lor i$, $\neg m \rightarrow i$, $\neg i$

In order for this to happen, clearly i must be false. In order for $\neg m \to i$ to be true when i is false, the hypothesis $\neg m$ must be false, so m must be true. Since we want $m \leftrightarrow n$ to be true, this implies that n must also be true. Since we want $n \to k$ to be true, we must therefore have k true. But now if k is true and i is false, then the third specification, $\neg k \lor i$ is false. Therefore we conclude that this system is not consistent.

- 51. Let s be "The router can send packets to the edge system"; let a be "The router supports the new address space"; let r be "The latest software release is installed." Then we are told $s \to a$, $a \to r$, $r \to s$, and $\neg a$. Since a is false, the first conditional statement tells us that s must be false. From that we deduce from the third conditional statement that r must be false. If indeed all three propositions are false, then all four specifications are true, so they are consistent.
- 53. This is similar to Example 17, about universities in New Mexico. To search for beaches in New Jersey, we could enter NEW AND JERSEY AND BEACHES. If we enter (JERSEY AND BEACHES) NOT NEW, then we'll get websites about beaches on the isle of Jersey, except for sites that happen to use the word "new" in a different context (e.g., a recently opened beach there). If we were sure that the word "isle" was in the name of the location, then of course we could enter ISLE AND JERSEY AND BEACHES.

- 55. If A is a knight, then he is telling the truth, in which case B must be a knave. Since B said nothing, that is certainly possible. If A is a knave, then he is lying, which means that his statement that at least one of them is a knave is false; hence they are both knights. That is a contradiction. So we can conclude that A is a knight and B is a knave.
- 57. If A is a knight, then he is telling the truth, in which case B must be a knight as well, since A is not a knave. (If $p \lor q$ and $\neg p$ are both true, then q must be true.) Since B said nothing, that is certainly possible. If A is a knave, then his statement is patently true, but that is a contradiction to the behavior of knaves. So we can conclude that A is a knight and B is a knight.
- 59. If A is a knight, then he should be telling the truth, but he is asserting that he is a knave. So that cannot be. If A is a knave, then in order for his statement to be false, B must be a knight. So we can conclude that A is a knave and B is a knight.
- 61. Because of the first piece of information that Steve has, let's assume first that Fred is not the highest paid. Then Janice is. Therefore Janice is not the lowest paid, so by the second piece of information that Steve has, Maggie is the highest paid. But that is a contradiction. Therefore we know that Fred is the highest paid. Next let's assume that Janice is not the lowest paid. Then our second fact implies that Maggie is the highest paid. But that contradicts the fact that Fred is the highest paid. Therefore we know that Janice is the lowest paid. So it appears that the only hope of a consistent set of facts is to have Fred paid the most, Maggie next, and Janice the least. (We have just seen that any other assumption leads to a contradiction.) This assumption does not contradict either of our two facts, since in both cases, the hypothesis is false.
- 63. Let's use the letters B, C, G, and H for the statements that the butler, cook, gardner, and handyman are telling the truth, respectively. We can then write each fact as a true proposition: $B \to C$; $\neg(C \land G)$, which is equivalent to $\neg C \lor \neg G$ (see the discussion of De Morgan's law in Section 1.2); $\neg(\neg G \land \neg H)$, which is equivalent to $G \lor H$; and $H \to \neg C$. Suppose that B is true. Then it follows from the first of our propositions that C must also be true. This tells us (using the second proposition) that G must be false, whence the third proposition makes H true. But now the fourth proposition is violated. Therefore we conclude that B cannot be true. If fact, the argument we have just given also proves that C cannot be true. Therefore we know that the butler and the cook are lying. This much already makes the first, second, and fourth propositions true, regardless of the truth of G or H. Thus either the gardner or the handyman could be lying or telling the truth; all we know (from the third proposition) is that at least one of them is telling the truth.
- 65. This is often given as an exercise in constraint programming, and it is difficult to solve by hand. The following table shows a solution consistent with all the clues, with the houses listed from left to right. Reportedly the solution is unique.

NATIONAL	ľTY	Norwegian	Italian	Englishman	Spaniard	Japanese
COLOR		Yellow	Blue	Red	White	Green
PET		Fox	Horse	Snail	Dog	Zebra
JOB		Diplomat	Physician	Photographer	Violinist	Painter
DRINK		Water	Tea	Milk	Juice	Coffee

In this solution the Japanese man owns the zebra, and the Norwegian drinks water. The logical reasoning needed to solve the problem is rather extensive, and the reader is referred to the following website containing the solution to a similar problem: http://mathforum.org/dr.math/problems/joseph8.5.97.html.