

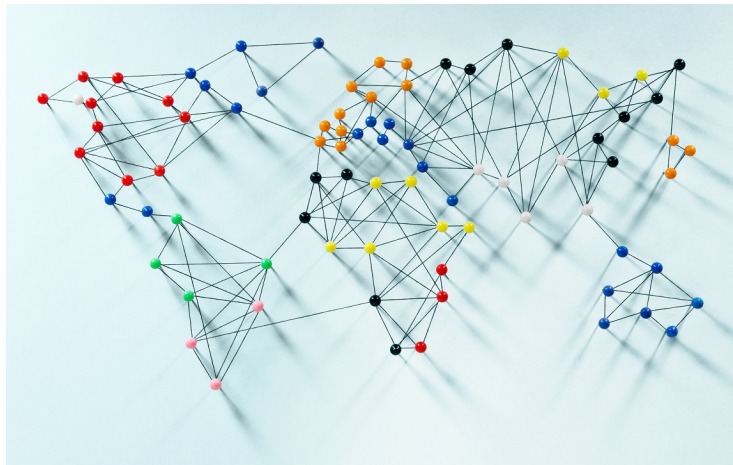
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يُونَيْتِيْ اِسْلَامْ اِنْتَارَا بَغْسِيَا مِلْدِسِيَا

Garden of Knowledge and Virtue

ECIE 3101 and ECIE 3312

MST Visualizer

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Objective

The objective of this project was to create a Graphical User Interface (GUI) that visualizes the application of two different algorithms – Kruskal’s algorithm and Prim’s algorithm – in finding the Minimum Spanning Tree (MST) of a graph. The primary goal was to make the understanding of these algorithms easier and interactive, aiming to promote better learning.

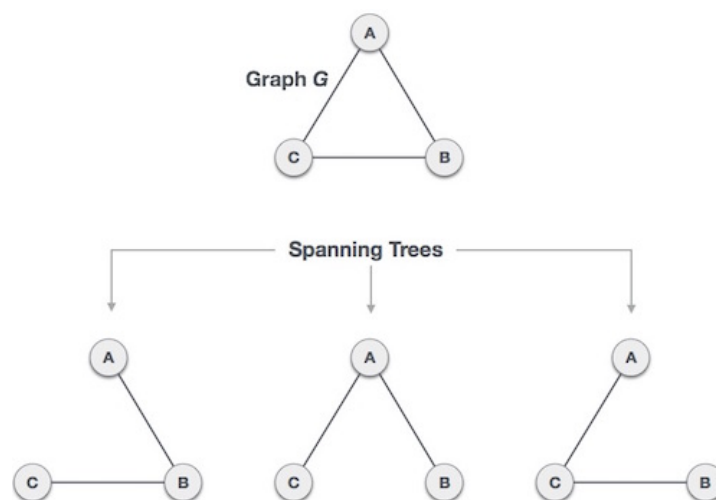
The project has been implemented in Python, using libraries such as NetworkX for handling and operating on the graph, and Tkinter for creating the GUI. An executable file was created to make the program accessible to users.

Introduction

Spanning Tree and Minimum Spanning Tree

A Spanning Tree is a subset of Graph G that has the least amount of connections feasible connecting every vertex. As a result, a spanning tree has no cycles and cannot be detached.

This concept leads us to the conclusion that any connected and undirected Graph G contains at least one spanning tree. A disconnected graph does not have any spanning tree, as it cannot be spanned to all its vertices.



A complete undirected graph can have a maximum n^{n-2} number of spanning trees, where n is the number of nodes. In the above addressed example, n is 3, hence $3^{3-2} = 3$ spanning trees are possible.

A Minimum Spanning Tree (MST) is a subset of the edges of a connected, edge-weighted undirected graph that connects all the vertices together, without any cycles, and with the minimum possible total edge weight. The MST has direct applications in network design, including telephone networks and computer networks.

Kruskal's Algorithm

Kruskal's algorithm is a minimum-spanning-tree algorithm that finds an edge of the least possible weight that connects any two trees in the forest. It is a greedy algorithm in graph theory, as it finds a minimum spanning tree for a connected weighted graph, adding increasing cost arcs at each step. We begin with the edges that have the lowest weight and keep adding edges until we accomplish our goal. The following are the stages for applying Kruskal's algorithm:

1. Sort all of the edges from light to heavy.
2. Add the edge to the spanning tree that has the lowest weight. Reject this edge if adding it resulted in a loop.
3. Keep adding the edges until all vertices have been reached.

Prim's Algorithm

Prim's algorithm is also a minimum-spanning-tree algorithm that finds an edge of the least possible weight that connects any two trees in the forest. It differs from Kruskal's algorithm, as it starts with an arbitrary node and grows the MST one edge at a time. Starting with one vertex, we continue to add edges with the lowest weight until we achieve our objective. Implementing Prim's algorithm involves the following steps:

1. Set up a random vertex to serve as the minimal spanning tree's first node.
2. Find the smallest edge between the tree's existing vertices and all other edges, then add it to the tree.
3. Repeat step 2 as necessary until a minimal spanning tree is formed.

Methodology

Our MST Visualization Tool was designed and developed leveraging the Python programming language, specifically employing the NetworkX and Tkinter libraries to fulfill the requirements of the project.

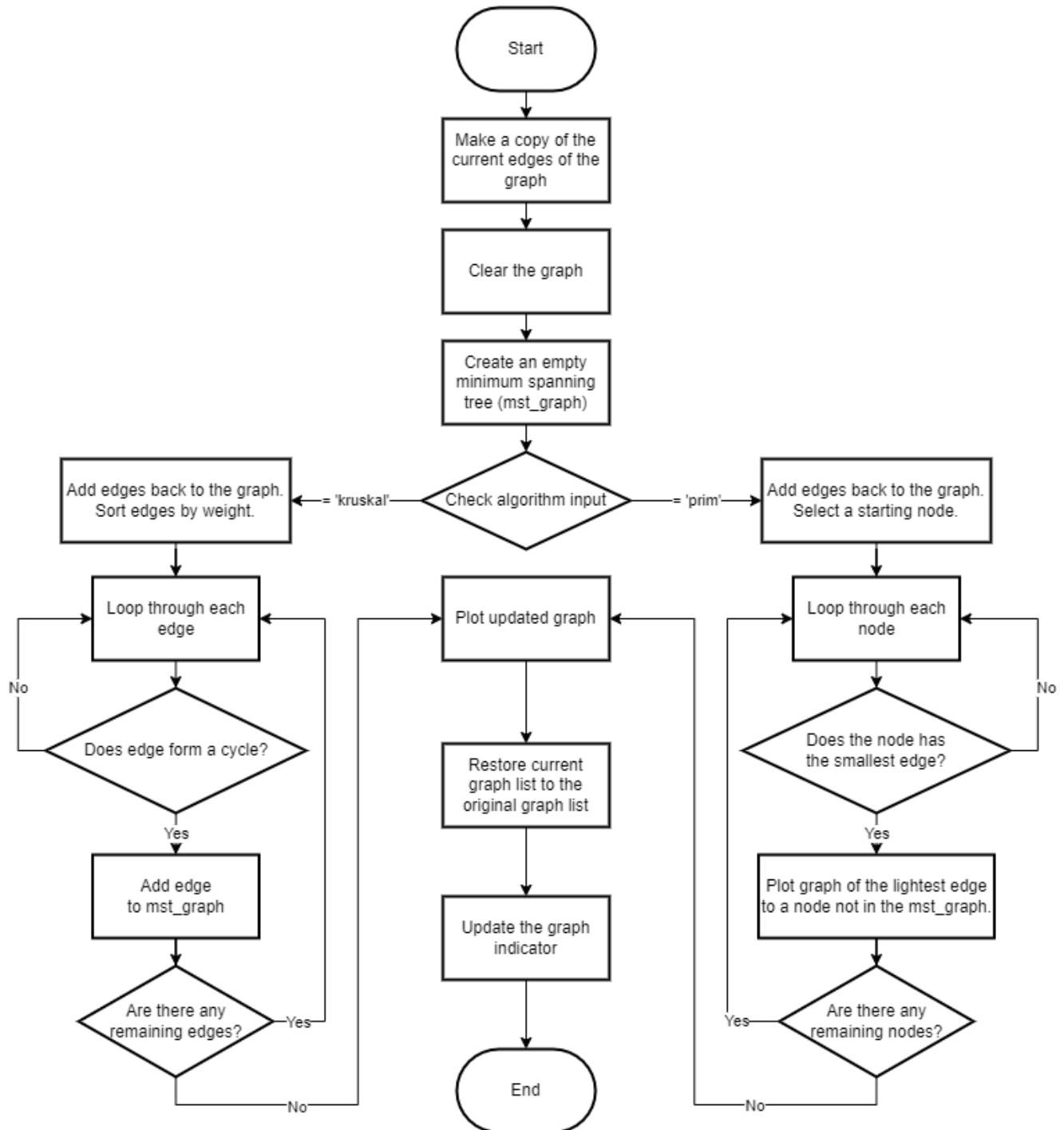
1. **NetworkX** is a robust library used to create, manipulate, and study the structure, dynamics, and functions of complex networks. In our tool, it plays a central role in performing all graph operations. This includes creating and managing graphs, adding nodes and edges, and implementing both Kruskal's and Prim's algorithms for Minimum Spanning Tree (MST) calculation.
2. **Tkinter**, on the other hand, is the standard Python interface to the Tk GUI toolkit. It offers a potent object-oriented interface to GUIs. We utilize Tkinter to design an intuitive and user-friendly interface where users can input graph data and observe the MST computation process and results.

The process of our tool unfolds through the following steps:

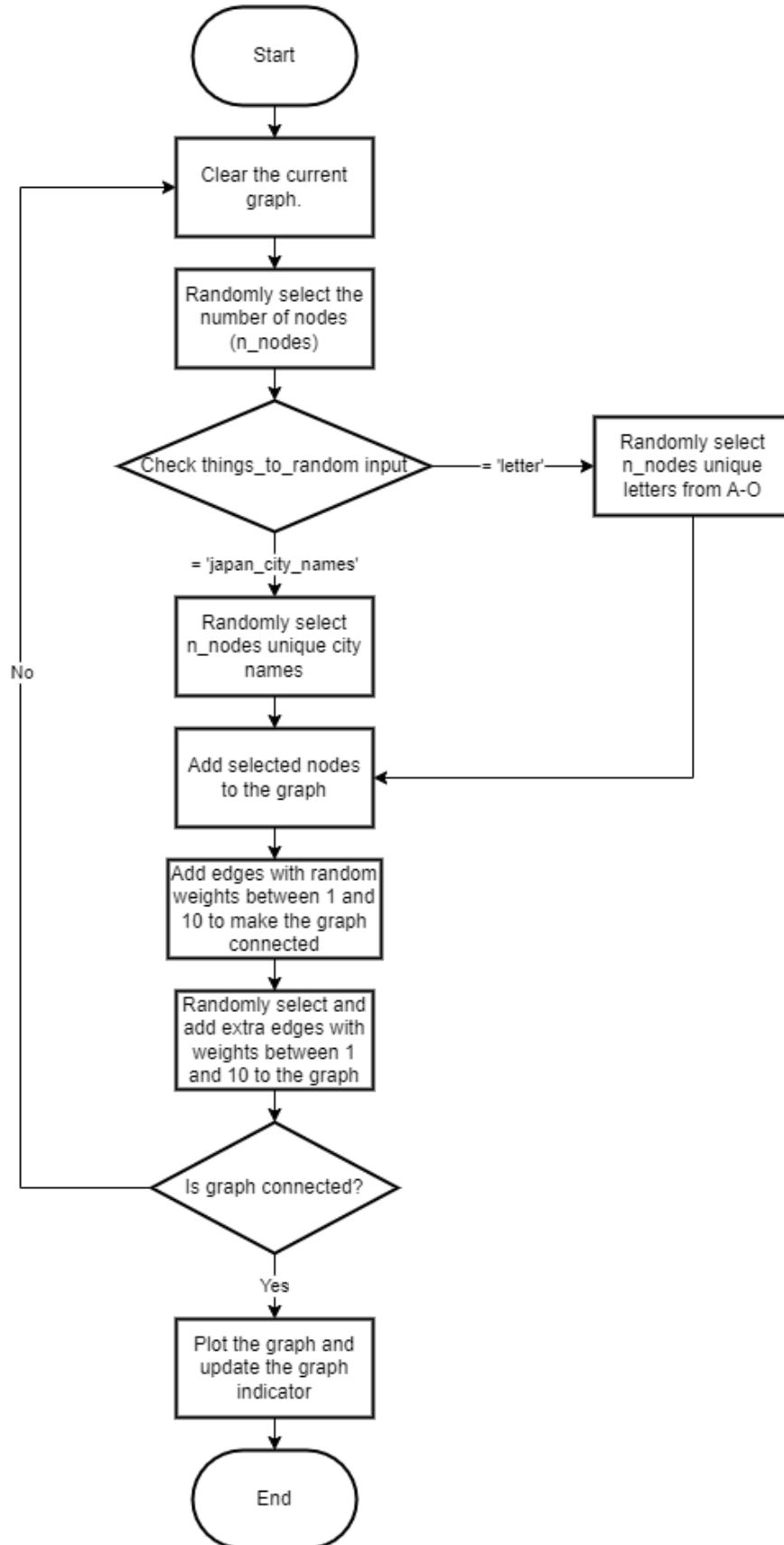
1. The tool first generates a GUI window using Tkinter. This window allows users to enter their graph details or opt for a randomly generated graph.
2. If the user chooses to manually input the graph, they input edges (node1, node2, weight) through the GUI interface. These edges are subsequently added to the graph structure maintained by NetworkX.
3. If the user opts for a randomly generated graph, the tool will randomly select nodes and edge weights and create a fully connected graph.
4. The user then selects the MST algorithm (either Kruskal or Prim) they wish to use for visualization purposes.
5. The selected algorithm is executed on the graph, and the process of constructing the MST is dynamically visualized on the GUI.
6. Upon completion, the resulting MST is displayed. The user can choose to add more edges, generate a new random graph, or run a different algorithm if they wish.

Flowchart

1. Kruskal and Prim's Algorithm



2. Generating random graph



Pseudocode

1. Kruskal and Prim's Algorithm

function calculate_mst(*algorithm*):

store a copy **of** the graph's edges in a variable "*edges_copy*"

clear the current graph

create an empty graph "*mst_graph*"

if the *algorithm* **is** "*kruskal*":

add the edges **from** "*edges_copy*" back to the graph

sort the edges **by** weight **in** ascending **order**

for each edge **in** the sorted edges:

if adding the edge **to** "*mst_graph*" does **not** form a cycle:

add the edge **to** "*mst_graph*"

update **and** display the "*mst_graph*"

pause **for** a moment before proceeding **to** the **next** edge

else if the *algorithm* **is** "*prim*":

add the edges **from** "*edges_copy*" back **to** the graph

choose a starting node, typically the first node **in** the graph

add the starting node **to** "*mst_graph*"

while "*mst_graph*" has fewer nodes than the graph:

mst_graph find the minimum-weight edge that connects a node **in** "*mst_graph*" **and** a node **not in**

add the found edge **to** "*mst_graph*"

update **and** display the "*mst_graph*"

pause **for** a moment before proceeding **to** the **next** edge

clear the graph and **add** the edges **from** "*edges_copy*" back **to** the graph

update the graph connectivity indicator

2. Generating random graph

function random_graph(*things_to_random*):

clear the **current** graph

generate a random number **of** nodes **between** a minimum **and** maximum limit

if *things_to_random* **is** 'letter':

generate a list **of unique** random letters **as** node names

else if *things_to_random* **is** 'japan_city_names':

generate a list **of unique** random Japanese city names **as** node names

until the graph **is** fully connected:

clear the graph

add the nodes **to** the graph

connect each node **to** the next **one with** a random weight

generate a random number **of** extra edges

for each extra edge:

choose two random nodes that **are not** already connected

connect the chosen nodes **with** a random weight

display the graph **and update** the graph connectivity **indicator**

Running Time

The time complexity of Kruskal's and Prim's algorithms differs based on the structures and characteristics of the graph they operate on.

Overall time complexity for Kruskal's and Prim's algorithms are $O(E \log E)$ and $O((E+V) \log V)$ respectively, where E is the number of edges and V is the number of vertices. Prim's algorithm performs well when it comes to dense graphs while Kruskal's algorithm shines when it comes to sparse graphs since it uses a simpler data structure.

Both Kruskal's and Prim's algorithms have a time complexity of $O(E \log E)$, where E is the number of edges. This is because both algorithms require sorting of edges, which take $O(E \log E)$ amount of time to be executed. However, in terms of efficiency, Kruskal's tends to be better for sparse graphs since it uses a simpler data structure, while Prim's is more efficient for dense graphs.

Conclusion

The MST Visualization tool serves as a compelling and interactive resource for individuals aiming to understand the operational intricacies of Kruskal's and Prim's algorithms in constructing a Minimum Spanning Tree. By providing a detailed visualization of each step in the process, the tool enables users to effectively grasp these foundational graph theory concepts in action.

Successfully meeting its predetermined objectives, the project delivers an engaging and dynamic learning platform. It not only fosters comprehension of key theoretical ideas but also promotes interactive learning, enriching the user's grasp of these fundamental graph algorithms. In essence, the MST Visualization tool represents a significant stride in the realm of algorithmic learning tools.

Executable file download link

Version 1.0:

https://drive.google.com/file/d/1R5rQd9D2-bp4U-uEc7dTsKH-wmf8-bzc/view?usp=drive_link

Version 1.1:

<https://drive.google.com/file/d/1MTNpcTC0UD0VMV09J79YMqJQ01pGAEsR/view?usp=sharing>

Version 2.0:

<https://drive.google.com/file/d/15u-Tg6aIrdUdc3qPF6z6Wr5KIO5RNPis/view?usp=sharing>