

MCMC

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# Chapter 1

## Introduction

In the ever-evolving landscape of mathematical and statistical research and application, the integration of simulation techniques has emerged as a powerful tool to unravel complex phenomena, validate theoretical frameworks, and facilitate a deeper understanding of intricate mathematical structures. Simulation techniques are usually used when any problem is heard to calculate with the help of traditional analytical method like simulation of a Solitaire Card game, to analysis this traditionally we have to deal with the number like  $52! = 8.06 \times 10^{67}$  this number is beyond astronomically large so, dealing with this number is very hard even for a high performant computer, to deal with this problem Polish-American mathematician and nuclear physicist Stanistaw Ulam with the help of Von Neumann developpe Monte Carlo Simulation.

## Chapter 2

# Generating Random Variables

### 2.1 Generating Discrete Random Variables

Main component of a simulation study is the ability to generate random number, where a random number represents the value of random variable uniform distribution on  $(0, 1)$ .

#### 2.1.1 Pseudorandom Number Generation

Random numbers were originally either manually or mechanically generated, by using spinning wheels or dice rolling or card shuffling but the modern approach is to use a computer to successively generate pseudorandom numbers.

One of the common approaches to generate pseudorandom numbers starts with an initial value  $x_0$ , called seed, and then recursively computes successive values  $x_n, n \geq 1$ , by letting

$$x_n = ax_{n-1} \text{ modulo } m \quad (2.1)$$

where  $a$  and  $m$  are given positive integers, and where the equation (2.1) means that  $ax_{n-1}$  is divided by  $m$  and remainder is taken as the value of  $x_n$ . Thus, each value of  $x_n$  is either  $0, 1, \dots, m-1$  and the quantity  $x_n/m$  is pseudorandom number and follows an approximation to the value of a uniform  $(0, 1)$  random variable.

The approach specified by equation (2.1) to generate random numbers is called the Multiplicative Congruential Method.

Another method is

$$x_n = (ax_{n-1} + c) \text{ modulo } m$$

this method is known as *Mixed Congruential Generators* or *Linear congruential Generations (LCGs)* where  $c$  is a non-negative integer.

#### 2.1.2 The Inverse Transform Method

Suppose we want to generate the value of a discrete random variable  $X$  having probability mass function

$$P(X = x_i) = p_i, \quad i = 0, 1, \dots, \quad \sum_i p_i = 1$$

To do this, we generate a random number from a uniform distribution  $(0, 1)$   $U$ , and set

$$X = \begin{cases} x_0 & \text{if } U < p_0 \\ x_1 & \text{if } p_0 \leq U \leq p_0 + p_1 \\ \vdots & \\ x_j & \text{if } \sum_{i=0}^{j-1} p_i \leq U \leq \sum_{i=0}^j p_i \\ \vdots & \end{cases}$$

Since, for  $0 < a < b < 1$ ,  $P(a \leq U < b) = b - a$ , we have,

$$P(X = x_j) = P\left(\sum_{i=0}^{j-1} p_i \leq U < \sum_{i=0}^j p_i\right) = p_j.$$

So,  $X$  has the desired distribution.

*Example 2.1.1* (Bernoulli Distribution). Let,  $X \sim Ber(p)$  where  $p$  is success probability i.e.  $P(X = 0) = 1 - p$  and  $P(X = 1) = p$  and  $0 \leq p \leq 1$ . Then, to generate  $X$  we first generate  $U \sim U[0, 1]$  then, we set

$$X = \begin{cases} 1, & \text{if } U \leq p \\ 0, & \text{if } U > p \end{cases}$$

Hence,  $X$  follows Bernoulli Distribution with the parameter  $p$ .

**Algorithm for Inverse Transform Algorithm for Generating Bernoulli Distribution:**

STEP 1: Generate a random variable  $U \sim U[0, 1]$ .

STEP 2: If  $U \leq p$  set  $X = 1$  or set  $X = 0$ .

STEP 3: Go to STEP 1.

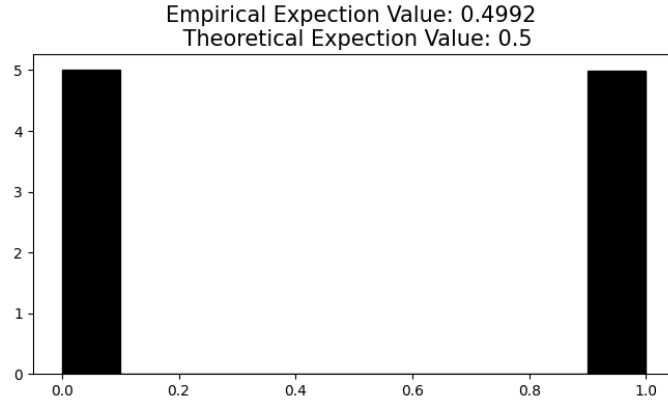


Figure 2.1: Inverse Transform method for generating Bernoulli random numbers with  $p = 0.5$

*Example 2.1.2* (Binomial Distribution). Let,  $X \sim Bin(n, p)$  then,  $X$  has probability mass function

$$f(r) = P(X = r) = \binom{n}{r} p^r (1 - p)^{n-r}, \quad i = 1, 2, \dots$$

The generation of  $X \sim Bin(n, p)$  by Inverse Transform Algorithm can be tedious. We can use the relation between Binomial and Bernoulli distribution. If  $x_i \sim Ber(p), \forall i = 1, 2, \dots, n$  then,  $\sum_{i=1}^n x_i \sim Bin(n, p)$ .

Hence, by generating  $x_i$   $n$  independent random variable from Bernoulli distribution and summing them we get binomial distribution

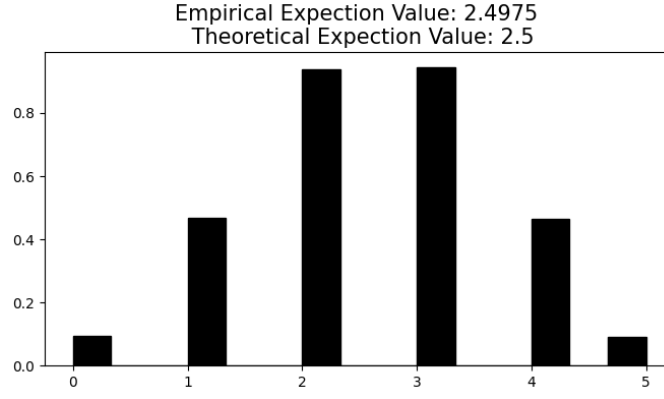


Figure 2.2: Generating binomial random numbers with  $n = 5$  and  $p = 0.5$

## 2.2 Generating Continuous Random Variables

### 2.2.1 The Inverse Transform Algorithm

To generate Continuous random variables The Inverse Transform Algorithm is very important method. It is based on a following theorem.

**Theorem 2.2.1.** *Let  $U$  be a uniform  $(0, 1)$  random variable. For any continuous distribution function  $F$  the random variable  $X$  defined by*

$$X = F^{-1}(U)$$

*has distribution  $F$ .*

*Proof.* Let,  $F_X$  denote the distribution function of  $X = F^{-1}(U)$ . Then,

$$\begin{aligned} F_X(x) &= P(X \leq x) \\ &= P(F^{-1}(U) \leq x) \end{aligned}$$

Since,  $F$  is a cumulative distribution function it follows that  $F(x)$  is monotonic increasing function of  $x$  and range of  $F(x)$  is  $(0, 1)$ . Then,

$$\begin{aligned} F_X(x) &= P(F(F^{-1}(U)) \leq F(x)) \\ &= P(U \leq F(x)) \\ &= F(x) \text{ since } U \sim U(0, 1) \end{aligned}$$

□

The above theory tells us we can generate a random variable  $X$  from the continuous distribution function  $F$  by generating a random number  $U \sim U(0, 1)$  and setting  $X = F^{-1}(U)$ .

*Example 2.2.1 (Exponential Distribution).* Suppose we want to generate a random variable  $x \sim \text{Exp}(\lambda)$ , then its probability density function is

$$f(x) = \lambda e^{-\lambda x}.$$

Hence, The cumulative distribution function is,

$$F(x) = 1 - e^{-\lambda x}$$

if we let  $x = F^{-1}(u)$ , then,

$$\begin{aligned} u &= F(x) = 1 - e^{-\lambda x} \\ 1 - u &= e^{-\lambda x} \\ x &= -\frac{\ln(1 - u)}{\lambda} \end{aligned}$$

Hence, we can generate an exponential random variable with parameter 1 by generating a uniform  $(0, 1)$  random number  $U$  and then setting

$$X = F^{-1}(U) = -\frac{\ln(1 - U)}{\lambda}.$$

We see that if  $U \sim U(0, 1)$  then also  $1 - U \sim U(0, 1)$  thus  $\ln(1 - U)$  has the same distribution as  $\ln(U)$  so,

$$X = F^{-1}(U) = -\frac{\ln(U)}{\lambda}.$$

will also work. If we use second expression then the algorithm will take less computing power hence less time.

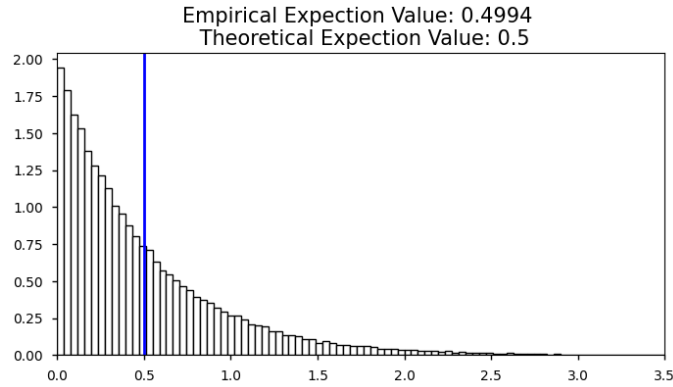


Figure 2.3: Inverse Transform method for generating  $Exp(2)$

*Example 2.2.2* (Gamma Distribution). Let  $X \sim G(n, \lambda)$  Then, its probability mass function is given by,

$$f(x) = \frac{1}{\Gamma(n)} \lambda^n x^{n-1} e^{-\lambda x}$$

We know if  $X_i \sim Exp(\lambda) \forall i = 1, 2, \dots, n$  then  $Y = \sum_i X_i \sim G(n, \lambda)$ . As,

$$\begin{aligned} M_Y(t) &= E[e^{tY}] = E\left[e^{\sum_{i=1}^n X_i t}\right] = E\left[\prod_{i=1}^n e^{X_i t}\right] \\ &= \prod_{i=1}^n E[e^{X_i t}] \quad \text{As all } X_i \text{ are independent} \\ &= \prod_{i=1}^n \frac{\lambda}{\lambda - t} = \left(\frac{\lambda}{\lambda - t}\right)^n \end{aligned}$$

Then, Generating  $n$  number of  $X_i \sim Exp(\lambda)$  and summing them we can easily generate a random variable which follows gamma distribution



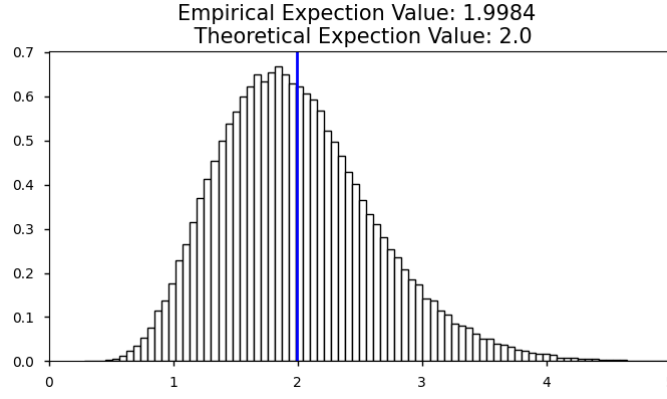


Figure 2.4:  $G(10, 5)$  generated by summing of  $Exp(5)$

### 2.2.2 Accept - Reject Method

The accept-reject method is useful when it is difficult to directly simulate  $f(x)$  but we can generate another density  $g(x)$  such that  $f(x)/g(x)$  is uniformly bounded and it is much easier to simulate  $g(x)$ . We simulate  $X$  from  $g$ , and retain it or toss it according to a probability proportional to  $f(x)/g(x)$ . Because an  $X$  value is either retained or discarded, depending on whether it passes the admission rule, the method is called the accept-reject method. The density  $g(x)$  is called the envelope density.

The method proceeds as follows,

STEP 1: Find a density  $g$  and a finite constant  $c$  such that  $\frac{f(x)}{g(x)} \leq c \forall x$ .

STEP 2: Generate  $X \sim g$ .

STEP 3: Generate  $U \sim U(0, 1)$ , independent of  $X$ .

STEP 4: Retain this generated value  $X$  if  $U \leq \frac{f(x)}{cg(x)}$ .

STEP 5: Repeat the same until the required number of  $n$  values of  $X$  has been obtained.

The following theorem supports the method.

**Theorem 2.2.2.** Let  $X \sim g$ , and  $U$ , independent of, be a distributed as  $U[0, 1]$ . Then the conditional density of  $X$  given that  $U \leq \frac{f(X)}{cg(X)}$  is  $f$ .

*Proof.* Denote the CDF of  $f$  by  $F$ . Then,

$$\begin{aligned} P\left(X \leq x | U \leq \frac{f(X)}{cg(X)}\right) &= \frac{P\left(X \leq x, U \leq \frac{f(X)}{cg(X)}\right)}{P\left(U \leq \frac{f(X)}{cg(X)}\right)} \\ &= \frac{\int_{-\infty}^x \int_0^{\frac{f(t)}{cg(t)}} g(t) du dt}{\int_{-\infty}^{\infty} \int_0^{\frac{f(t)}{cg(t)}} g(t) du dt} \\ &= \frac{\int_{-\infty}^x f(t) dt}{\int_{-\infty}^{\infty} f(t) dt} = \frac{F(x)}{1} = F(x). \end{aligned}$$

□

*Example 2.2.3* (Generating a Normal Random Variable). To generate a standard normal variable  $Z$  i.e.  $Z \sim N(0, 1)$ , note first that the absolute value of  $Z$  has probability density function

$$f(x) = \frac{2}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} \quad 0 \leq x < \infty. \quad (2.2)$$

Then, we can choose  $g$  as the exponential density function with mean 1 i.e.

$$g(x) = e^{-x} \quad 0 \leq x \leq \infty$$

Now,

$$\frac{f(x)}{g(x)} = \sqrt{\frac{2}{\pi}} e^{x - \frac{x^2}{2}}$$

and so the maximum value of  $f(x)/g(x)$  occurs at the value of  $x$  that maximize  $x - x^2/2$  hence  $x = 1$  so we take

$$c = \max_x \frac{f(x)}{g(x)} = \frac{f(1)}{g(1)} = \sqrt{\frac{2e}{\pi}}.$$

Now,

$$\frac{f(x)}{cg(x)} = \exp\left(x - \frac{x^2}{2} - \frac{1}{2}\right) = \exp\left(\frac{-(x-1)^2}{2}\right)$$

Then, it follows that we can generate the absolute value of a standard normal random variable as follows:

STEP 1: Generate  $X \sim \text{Exp}(1)$ .

STEP 2: Generate  $U \sim U(0, 1)$ , independent of  $X$ .

STEP 3: If  $U \leq \exp(-(X-1)^2/2)$ , retain  $X$ , Otherwise, return to Step 1.

Once, we have simulated a random variable  $X$  having density function as in Equation (2.2) we can obtain a standard normal  $Z$  by letting  $Z$  be equally likely to be either  $X$  or  $-X$ . In Step 3, the value  $X$  is accepted if  $U \leq \exp(-(X-1)^2/2)$ , which is equivalent to  $-\ln U \geq (X-1)^2/2$ . However, in Example (2.2.1) we have seen that  $-\ln U \sim \text{Exp}(1)$  When  $U \sim U(0, 1)$ .

So, summing up, we can generate the standard normal random variable  $Z$  as follows:

STEP 1: Generate independent  $X_1, X_2 \sim \text{Exp}(1)$

STEP 2: If  $X_2 \geq (X_1 - 1)^2/2$  retain  $X_1$ . Otherwise, return to Step 1.

STEP 3: Generate  $U \sim U(0, 1)$  and set,

$$Z = \begin{cases} X_1 & \text{if } U \leq \frac{1}{2}, \\ -X_1 & \text{if } U > \frac{1}{2}. \end{cases}$$

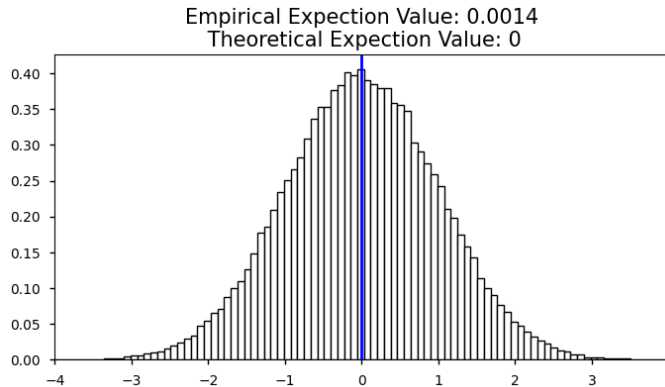


Figure 2.5: Generating  $N(0, 1)$  with Accept - Reject method

If we want to generate normal random variable to have mean  $\mu$  and variance  $\sigma^2$ , just take  $\mu + \sigma Z$ .

*Example 2.2.4 (Generating Beta Distribution).* If  $\alpha$  and  $\beta$  are both greater than 1, then Beta density is uniformly bounded and its maximum attained at  $\frac{\alpha-1}{\alpha+\beta-2}$ . As a result the  $U[0, 1]$  density can be

served as an envelope density for generating such Beta distribution by using accept-reject method. Precisely, generate  $U, X \sim U[0, 1]$  (independently), and retain the value if  $U \leq \frac{f(X)}{\sup_x f(X)}$ , where,

$$f(X) = \frac{\Gamma(\alpha + \beta)}{\Gamma(\alpha)\Gamma(\beta)} x^{\alpha-1} (1-x)^{\beta-1}, 0 < x < 1.$$

Because

$$\sup_x f(X) = f\left(\frac{\alpha-1}{\alpha+\beta-2}\right) = \frac{\Gamma(\alpha+\beta)}{\Gamma(\alpha)\Gamma(\beta)} \frac{(\alpha-1)^{\alpha-1}(\beta-1)^{\beta-1}}{(\alpha+\beta-2)^{\alpha+\beta-2}}$$

The algorithm finally works out as follows:

STEP 1: Generate independent  $U, X \sim U[0, 1]$ .

STEP 2: Retain the value  $X$  if,

$$U \leq \frac{X^{\alpha-1}(1-X)^{\beta-1}(\alpha+\beta-2)^{\alpha+\beta-2}}{(\alpha-1)^{\alpha-1}(\beta-1)^{\beta-1}}.$$

Otherwise, return to STEP 1.

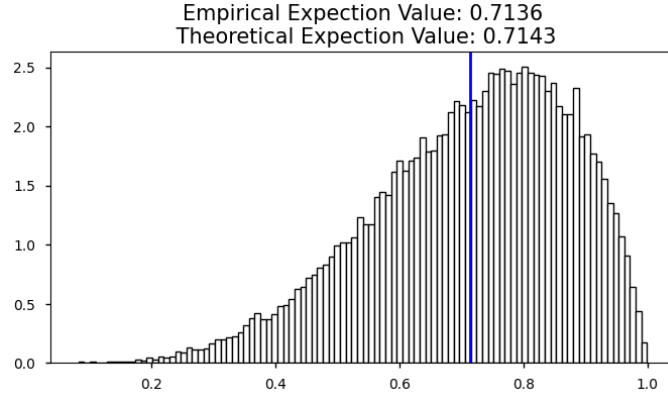


Figure 2.6: Generating Beta(5,2) with accept-reject method

An issue about an accept-reject method is the acceptance rate. Our goal make it as large as possible to increase the efficiency of the method. This can be achieved by choosing  $c$  to be smallest possible number, described in the result bellow.

**Theorem 2.2.3** (Acceptance Rate). *For an accept-reject scheme, the probability that an  $X \sim g$  is acceded is  $\frac{1}{c}$ , and is maximized when  $c$  is chosen to be  $c = \sup_x \frac{f(x)}{g(x)}$ .*

*Proof.*

$$\begin{aligned} P\left(U \leq \frac{f(x)}{cg(x)}\right) &= \int_{-\infty}^{\infty} \int_0^{\frac{f(x)}{cg(x)}} g(t) du dt \\ &= \int_{-\infty}^{\infty} \frac{f(t)}{cg(t)} g(t) dt = \int_{-\infty}^{\infty} \frac{f(t)}{c} dt = \frac{1}{c}. \end{aligned}$$

Because any  $c$  that can be chosen must be at least as large as  $\sup_x \frac{f(x)}{g(x)}$ , obviously  $1/c$  is maximized by choosing  $c = \sup_x \frac{f(x)}{g(x)}$ .  $\square$

In the example (2.2.3) for  $N(0, 1)$  the acceptance rate is  $\sqrt{\frac{\pi}{2e}} = 0.7601$ . And for example (2.2.4) for Beta(5,2) the acceptance rate is 0.4069

## 2.3 The Polar Method for Generating Normal Random Variables

Let  $X$  and  $Y$  be independent standard normal random variable and let  $R$  and  $\theta$  denote the polar coordinates of vector  $(X, Y)$ . That is,

$$R^2 = X^2 + Y^2$$

$$\tan \theta = \frac{Y}{X}$$

Since  $X$  and  $Y$  are independent, their joint density is the product of their individual densities and thus given by

$$f(x, y) = \frac{1}{\sqrt{2\pi}} e^{-x^2/2} \frac{1}{\sqrt{2\pi}} e^{-y^2/2}$$

$$= \frac{1}{\sqrt{2\pi}} e^{-(x^2+y^2)/2}$$

To determine the joint density of  $R^2$  and  $\Theta$  - call it  $f(d, \theta)$  we make the change of variables

$$d = x^2 + y^2, \quad \theta = \tan^{-1} \left( \frac{y}{x} \right)$$

Then the joint density function of  $R^2$  and  $\Theta$  is,

$$f(d, \theta) = \frac{1}{2} \frac{1}{2\pi} e^{-d/2}, \quad 0 < d < \infty, 0 < \theta < 2\pi. \quad (2.3)$$

As  $f(d, \theta)$  is equal to product of the product of  $Exp(1/2)$  density and  $U(0, 2\pi)$ , it follows that,  $R^2$  and  $\Theta$  are independent, with  $R^2 \sim Exp(1/2)$  and  $\Theta \sim U(0, 2\pi)$

Hence to generate a pair of independent standard normal random variables  $X$  and  $Y$  by generating  $R^2$  and  $\Theta$  in polar coordinates and then transform back to rectangular coordinates. Hence the algorithm is:

STEP 1: Generate random number  $U_1, U_2 \sim U(0, 1)$ .

STEP 2:  $R^2 = -2 \ln U_1$  and  $\Theta = 2\pi U_2$ .

STEP 3: Now let,

$$X = R \cos \Theta = \sqrt{-2 \ln U_1} \cos(2\pi U_2)$$

$$Y = R \sin \Theta = \sqrt{-2 \ln U_1} \sin(2\pi U_2). \quad (2.4)$$

The transformation given by equations (2.4) are known as Box-Muller transformation.

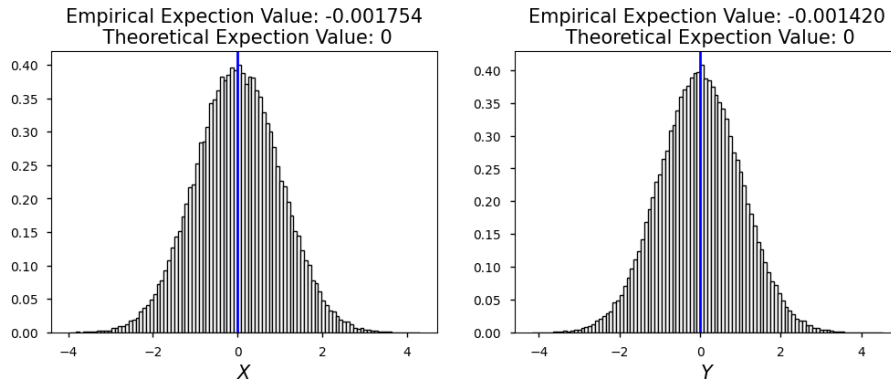


Figure 2.7: Generating independent  $X, Y \sim N(0, 1)$  with polar method