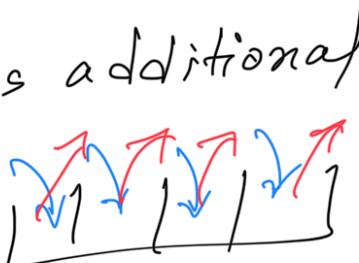
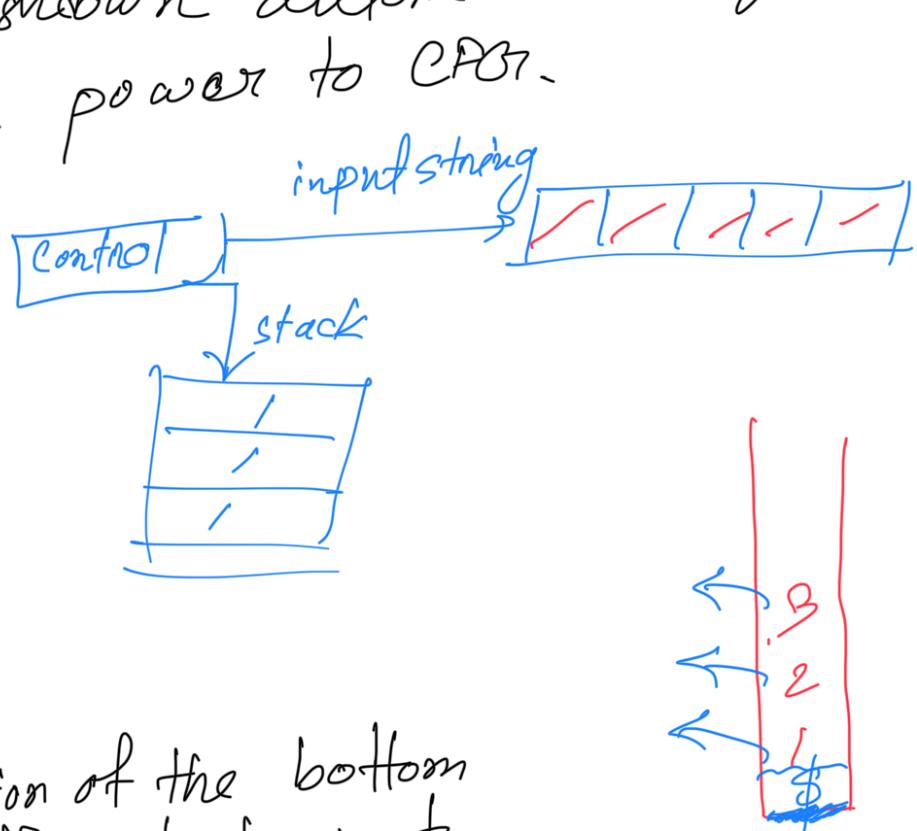


Pushdown Automata

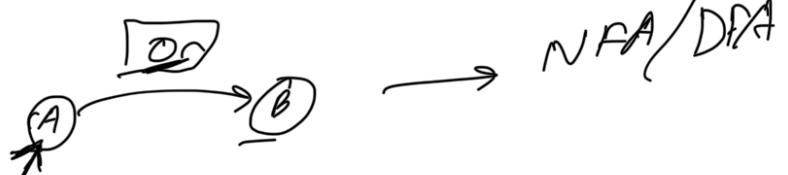
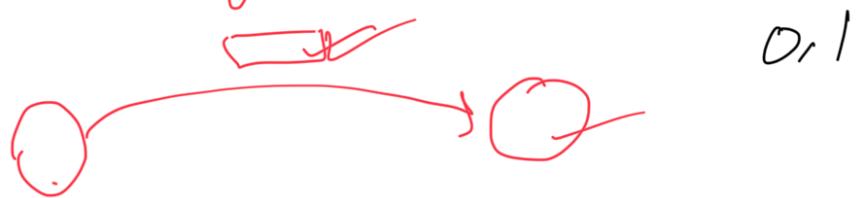
- * This is like NFA but has additional component → **Stack.**
 - provides additional memory beyond the finite memory of automata.
 - Allows to recognize some non-regular language.
- * Pushdown automata is equivalent in power to CPT.



Indication of the bottom of the stack $\Rightarrow \emptyset$

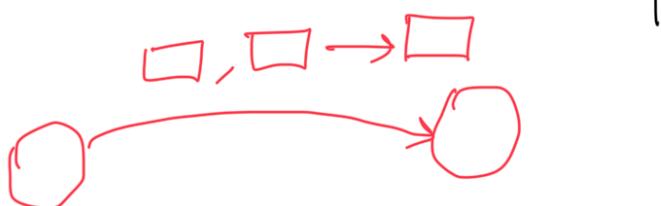
Write symbol on the stack \Rightarrow push

Read symbol from the stack \rightarrow pop

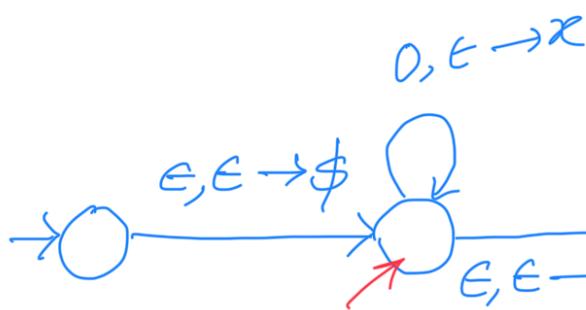
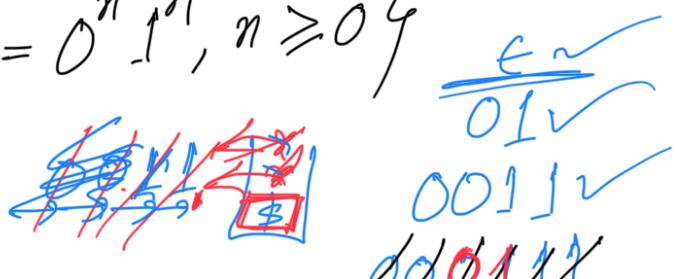


format

input symbol, pop symbol \longrightarrow push symbol.

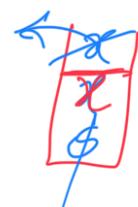


$$\textcircled{1} \quad L = \{ \omega \in \{0, 1\}^*: \omega = 0^n 1^n, n \geq 0 \}$$



1, x \rightarrow E

$\epsilon, \$ \rightarrow E$



000111✓

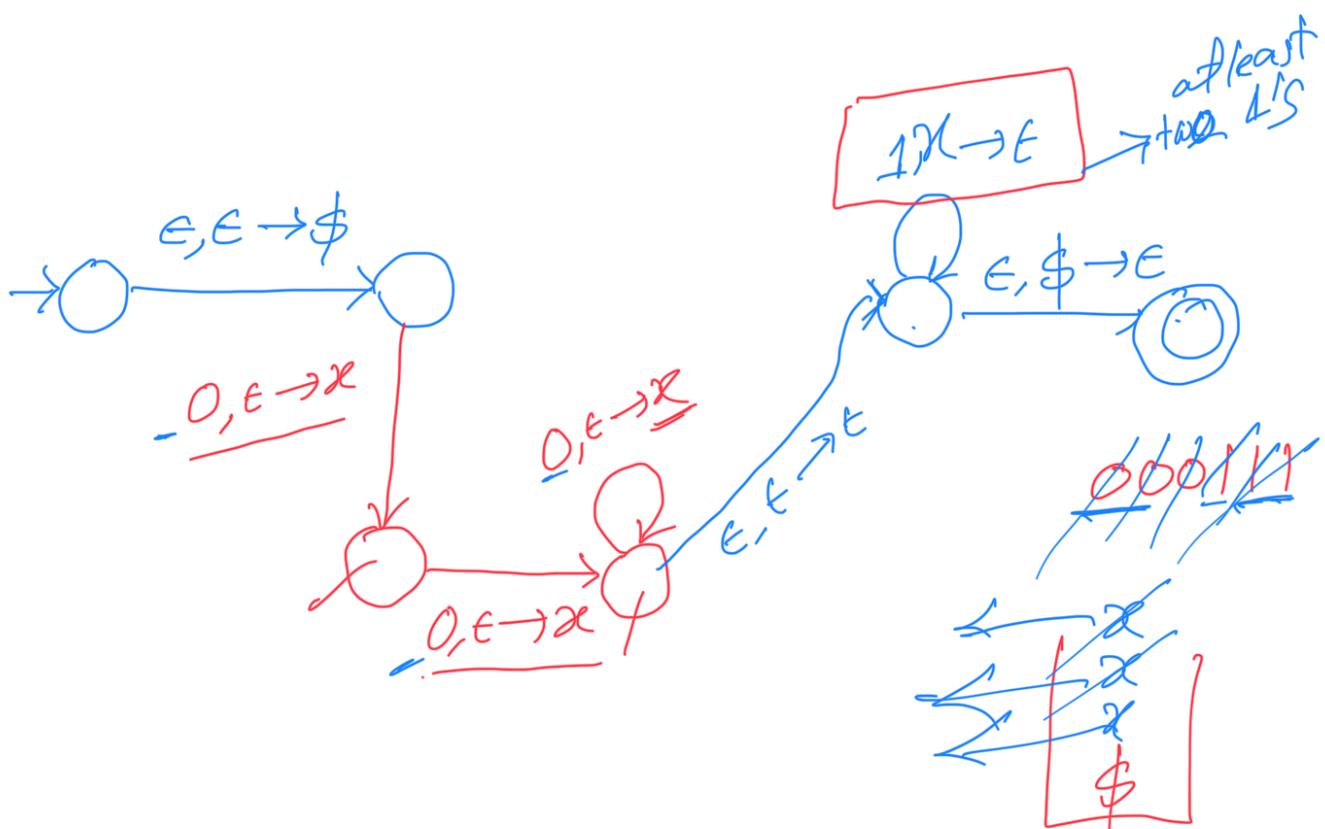
~~00110X~~

四

11

$$\textcircled{11} \quad L = \left\{ \omega \in S_0, 1^{\rho^*} : \omega = O^n - 1^n, n \geq \underline{\underline{2}} \right\}$$

~~0111~~ ✓
~~0111~~ ✓
000111✓



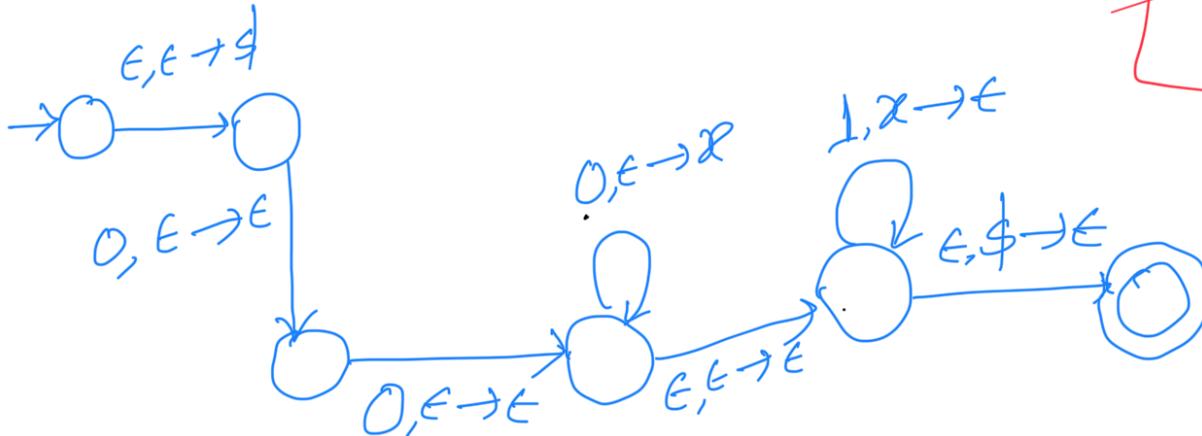
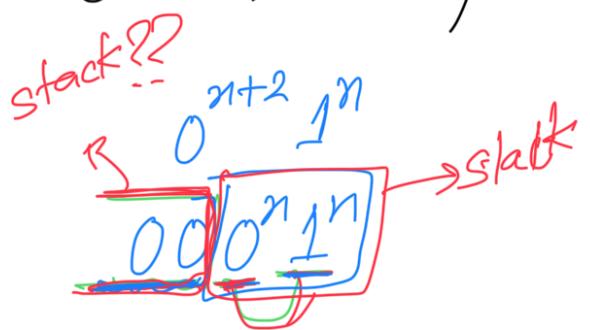
$$\begin{array}{c} \swarrow \searrow \\ x \quad x \\ \nearrow \searrow \\ -x \quad y \quad z \end{array}$$

100

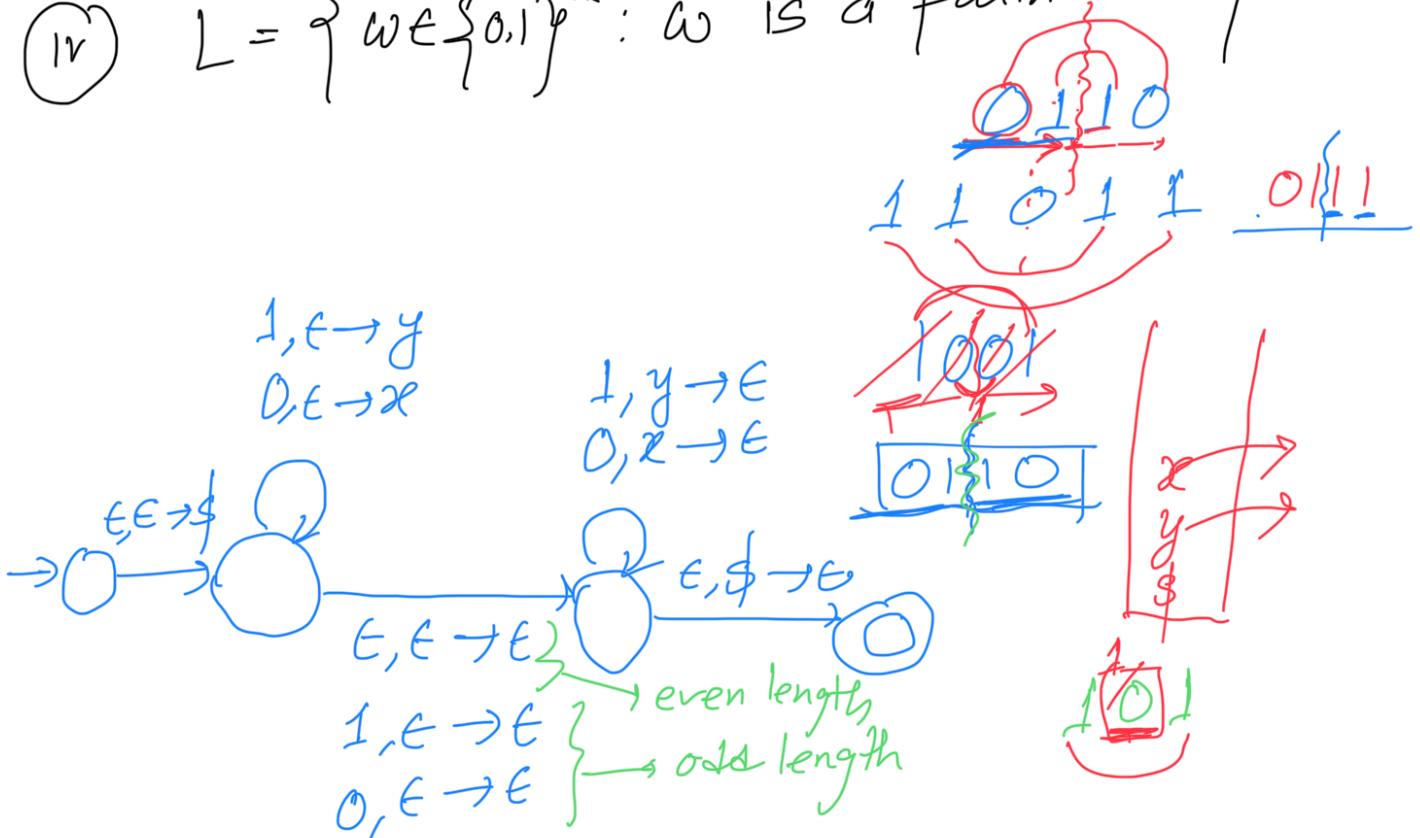
$$C_{n+2} = C_n + 2^n \quad n \geq 0$$

(11) $L = \{w \in \{0,1\}^*: w = u^+ + \dots + u^-\}$

$$\begin{array}{c} 00 \\ \hline 0001 \\ \hline 000011 \end{array}$$

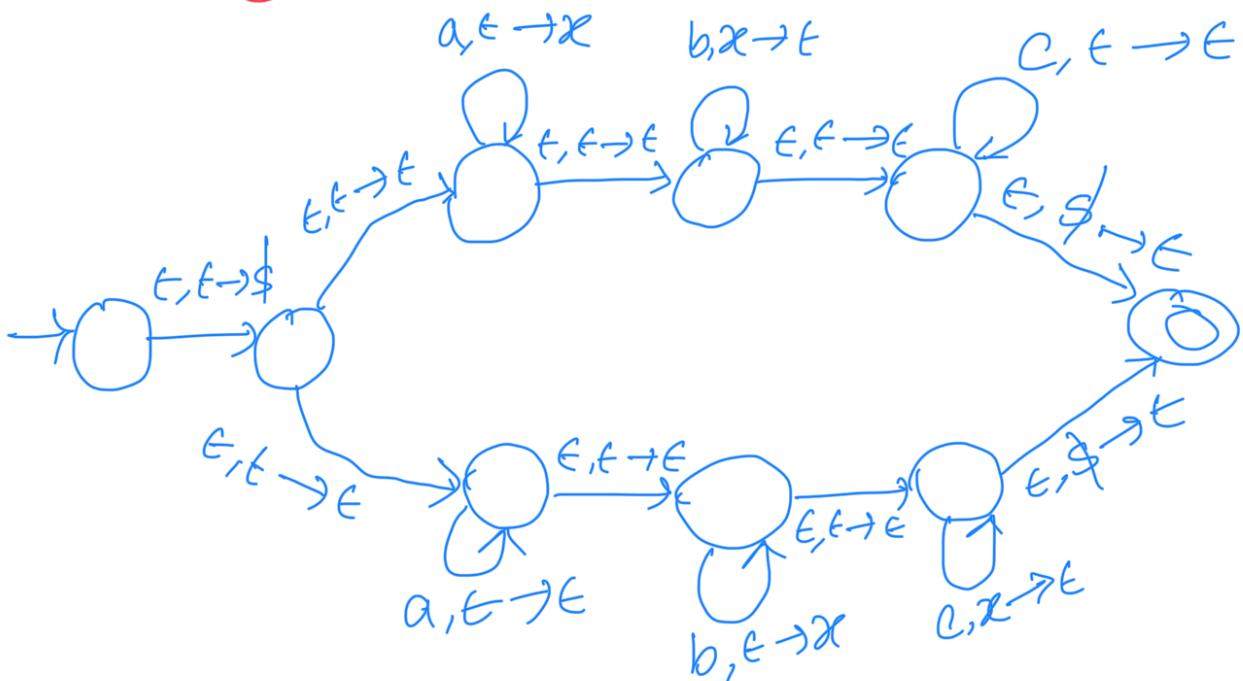


(12) $L = \{w \in \{0,1\}^*: w \text{ is a palindrome}\}$

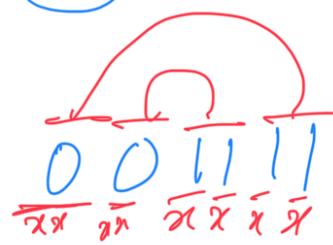
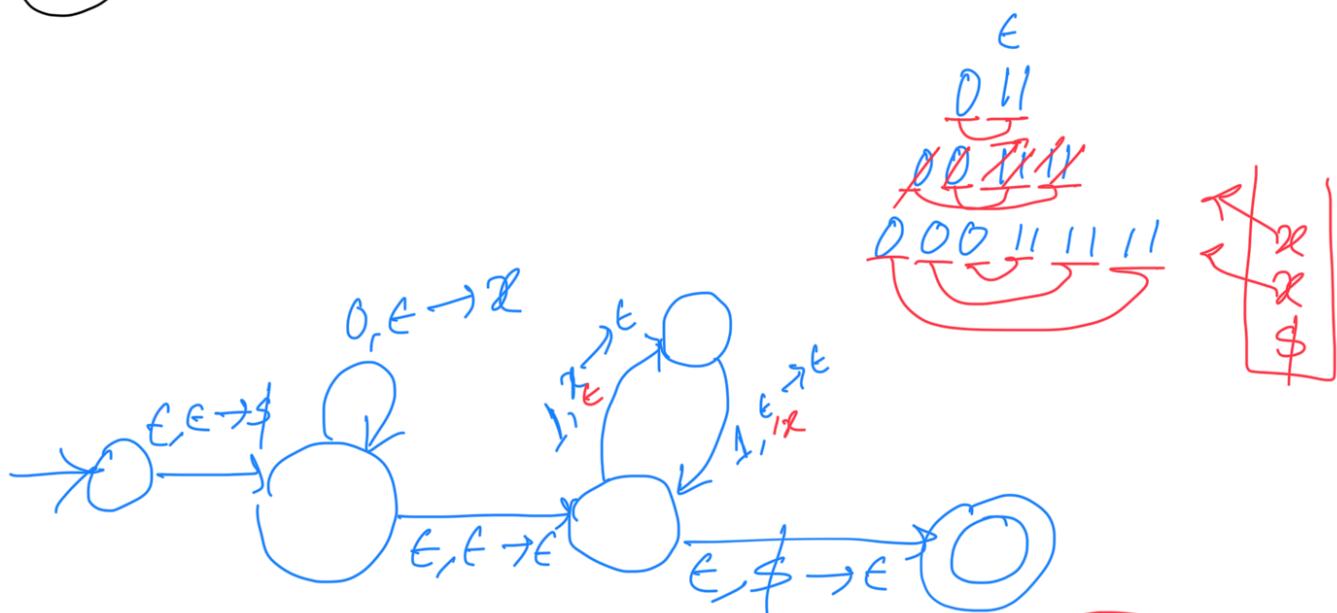


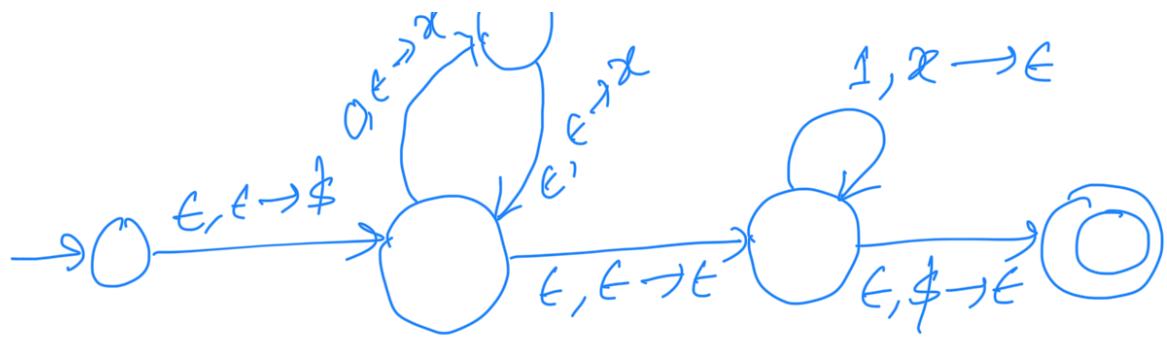
(13) $L = \{w \in \{a,b,c\}^*: w = a^i b^j c^k, \text{ where } i=j \text{ or } j=k\}$

$\cup \quad \cap \quad \sim \quad - \quad |$
 $= \underline{\underline{a^ib^jc^k}} \quad / \quad \underline{\underline{a^ib^jc^k}}$
 ① ②

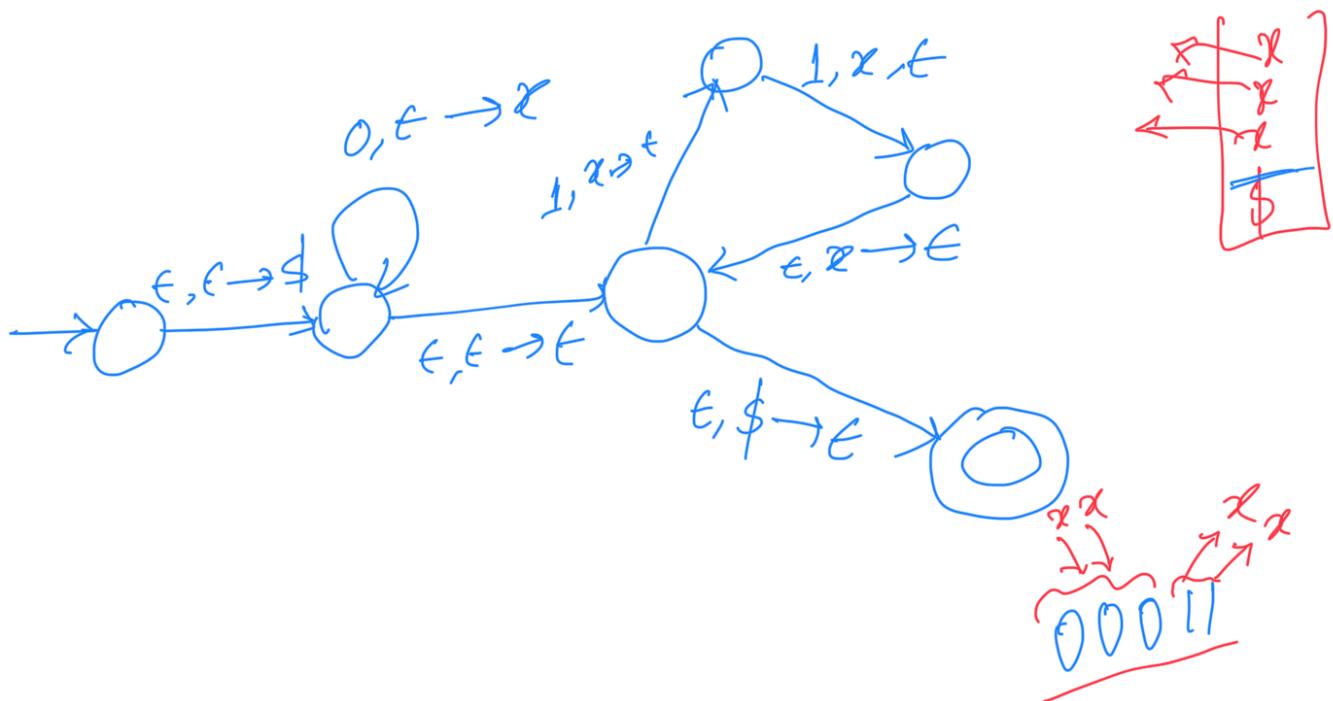


VI $L = \{w \in \{0,1\}^*: w = 0^n 1^m 2^n, n \geq 0\}$

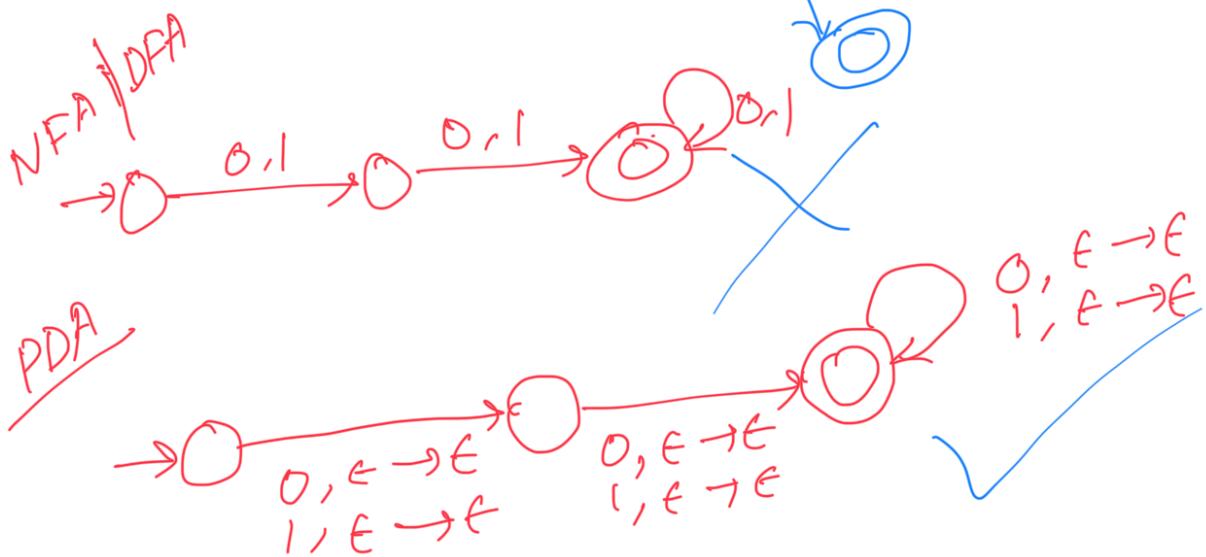
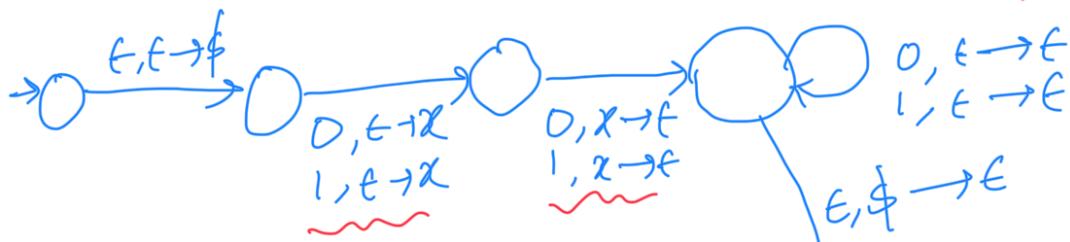




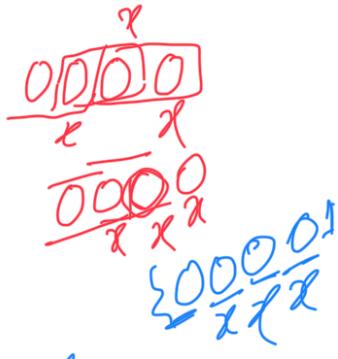
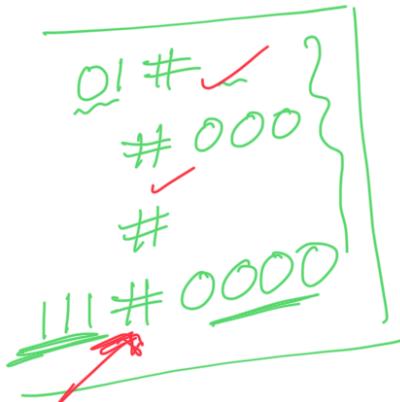
VII) $L = \{w \in \{0,1\}^*: w = 0^{3n}1^{2n}, n \geq 0\}$



VIII) $L = \{w \in \{0,1\}^*: \text{The length of } w \text{ is at least } \frac{\text{two}}{\text{one}}\}$

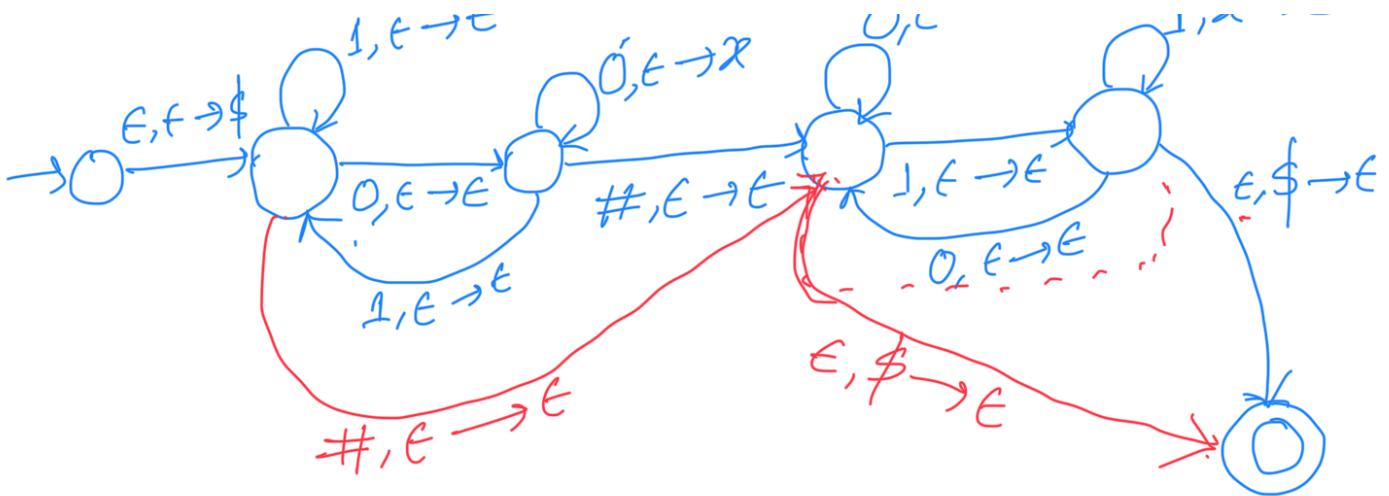


IX) $L = \{ w_1 \# w_2 \mid \text{The number of } '00' \text{ in } w_1 \text{ is the same as the number of } '11' \text{ in } w_2 \}$



001 # 10011
000 # 0110011,
110 # 000

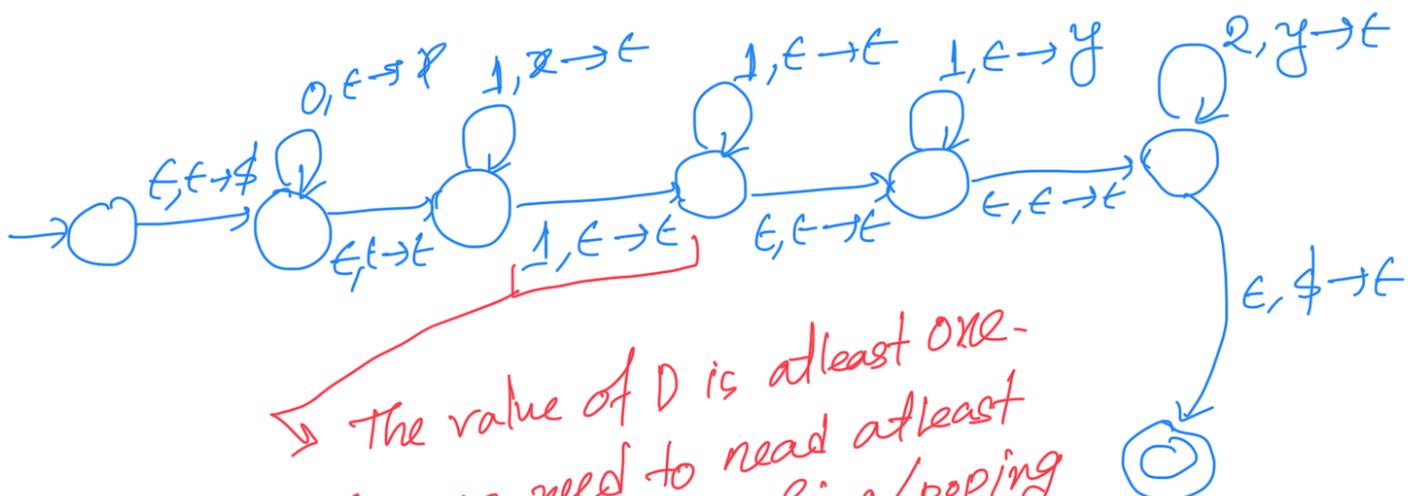
$\sim \leftarrow \epsilon$ $1 \rightarrow \epsilon$



(X) $L = \{w \in \{0, 1, 2\}^*: w = 0^i 1^j 2^k, j \geq i+k, i, k \geq 0\}$



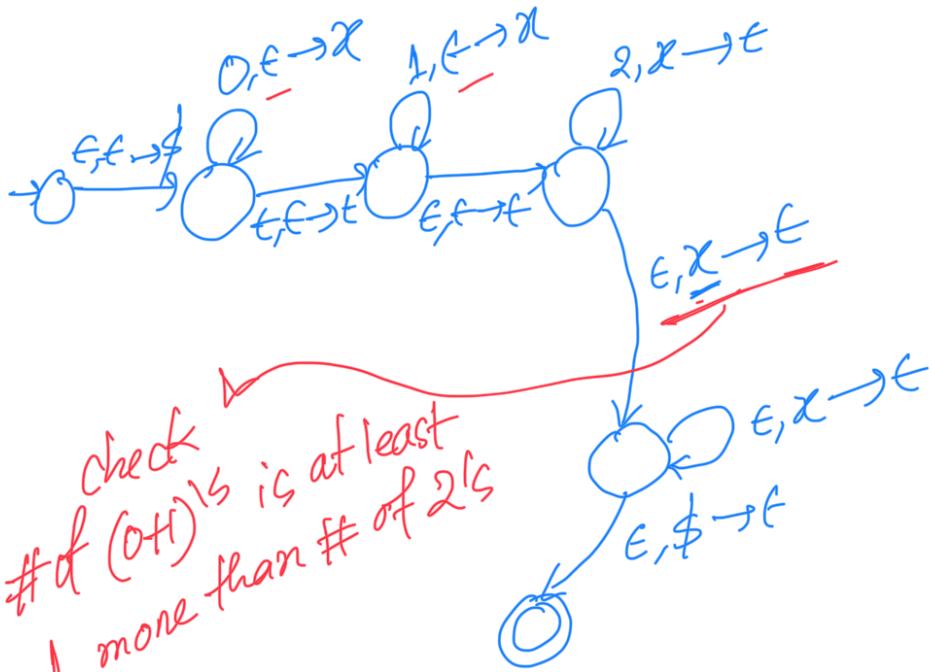
$$j = i + k + D \rightarrow 1 + \infty$$



The value of D is at least one.
So we need to read at least one 1 without pushing/popping anything to/from the stack.

XI

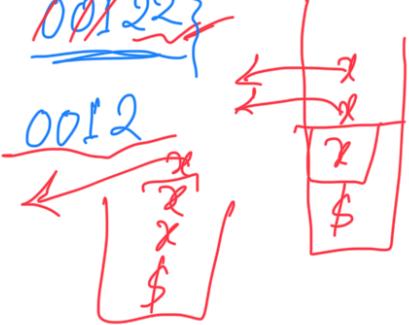
$$L = \{ \omega \in \{0, 1, 2\}^*, \omega = 0^i 1^j 2^k, i+j > k \}$$



$$i+j = k+d \rightarrow 1 \rightarrow \infty$$

$$k = i+j-d$$

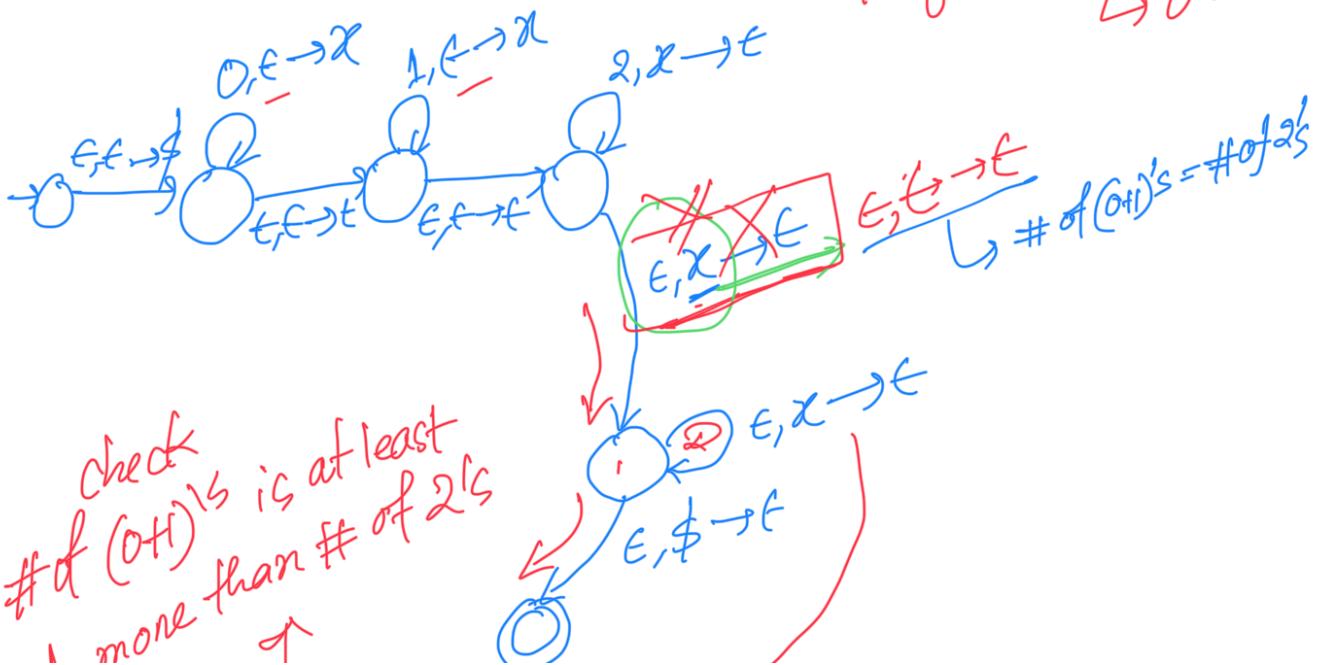
$$\frac{i}{0} \frac{j}{1} \frac{k}{2} \{$$



XII

$$L = \{ \omega \in \{0, 1, 2\}^*, \omega = 0^i 1^j 2^k, i+j \geq k \}$$

$$i+j = k+d \rightarrow 0 \rightarrow \infty$$





Practice

- ① $L = \{ w \in \{a,b\}^*: w = a^i b^j, \text{ where } i > j, j \geq 0 \}$
- ② $L = \{ w \in \{0,1,2\}^*: w = 0^i 1^j 2^k, \text{ where } i+j+k = K, i, j, k \geq 0 \}$
- ③ $L = \{ w \in \{0,1\}^*: 0 \text{ and } 1 \text{ alternates in } w \}$
- ④ $L = \{ w \in \{a,b\}^*: \text{the count of "a" in } w \text{ is a multiple of 3} \}$