Artificial Intelligence CSE 4617

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What is Search for?

- Assumptions about the world
 - Single agent → No adverseries
 - Deterministic actions
 - Fully observed state
 - Discrete state space

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 - Paths have various costs, depths
 - Heuristics give problem-specific guidance

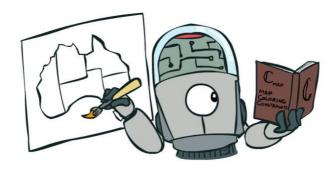


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 - Paths have various costs, depths
 - Heuristics give problem-specific guidance
- Identification: assignments to variables
 - The goal itself is important, not the path
 - All paths at the same depth (for some formulations)
 - CSPs are a specialized class of identification problems







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- Simple example of a formal representation language
- Allows useful general-purpose algorithms with more power than standard search algorithms

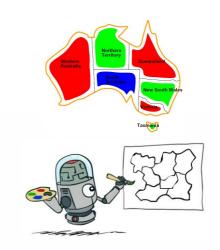




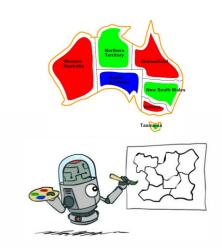
CSP Examples



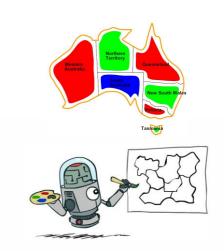
Variables:



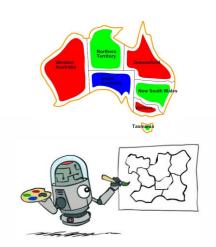
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 - *WA*, *NT*, *Q*, *NSW*, *V*, *SA*, *T*



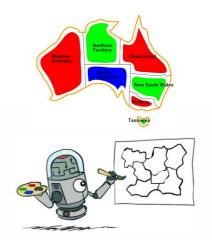
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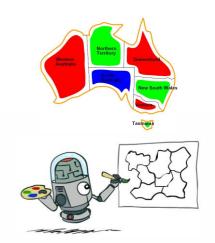
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- Constraints: adjacent regions must have different colors

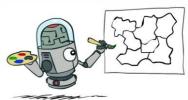


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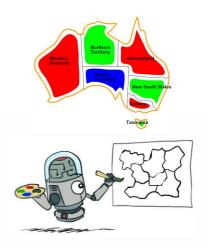


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 (WA, NT) ∈ {(red, green), (red, blue), . . . }

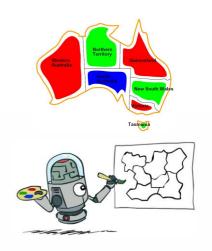


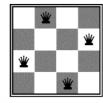


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 - Implicit: $WA \neq NT$
 - Explicit: (WA, NT) ∈ {(red, green), (red, blue),...}
- Solutions are assignments satisfying all constraints
 - $\{WA = \text{red}, NT = \text{green}, Q = \text{red}, NSW = \text{green}, V = \text{red}, SA = \text{blue}, T = \text{green}\}$

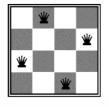






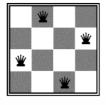


■ Formulation 1:



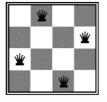


- Formulation 1:
 - Variables: X_{ij}



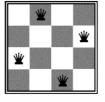


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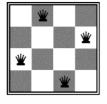


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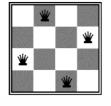




$$\forall ij, k(X_{ij}, X_{ik}) \in \{(0,0), (0,1), (1,0)\}$$

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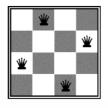




```
\forall ij, k (X_{ij}, X_{ik}) \in \{(0,0), (0,1), (1,0)\}
\forall ij, k (X_{ij}, X_{kj}) \in \{(0,0), (0,1), (1,0)\}
\forall ij, k (X_{ij}, X_{i+k,j+k}) \in \{(0,0), (0,1), (1,0)\}
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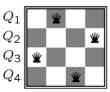
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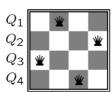
$$\forall ij,k(X_{ij},X_{i+k,j-k}) \in \{(0,0),(0,1),(1,0)\}$$

$$\sum_{i,j} X_{ij} = N$$

- Formulation 2:
 - Variables: Q_k
 - Domains: $\{1, 2, 3, ..., N\}$
 - Constraints:



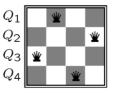
- Formulation 2:
 - Variables: Q_k
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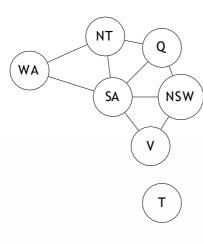
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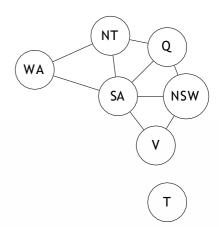


Constraint Graphs



Constraint Graphs

- Binary CSP: each constraint relates (at most) two variables
- Binary constraint graph: nodes are variables, arcs show constraints
- General-purpose CSP algorithms use the graph structure to speed up search. E.g., Tasmania is an independent subproblem!



Example: Cryptarithmetic

Variables

_ Domains:

Constraints:

T W O + T W O F O U R



Example: Cryptarithmetic

- Variables
 - F, T, U, W, R, O, X_1 , X_2 , X_3
- Domains:
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Example: Cryptarithmetic

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Example: Cryptarithmetic

- Variables
 - F, T, U, W, R, O, X_1 , X_2 , X_3
- Domains:
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- Constraints: alldiff(F, T, U, W, R, O)





Example: Cryptarithmetic

- Variables
 - $F, T, U, W, R, O, X_1, X_2, X_3$
- Domains:
 - {0,1,2,3,4,5,6,7,8,9}
- Constraints:

alldiff(
$$F$$
, T , U , W , R , O)
 $O + O = R + 10 \times X_1$

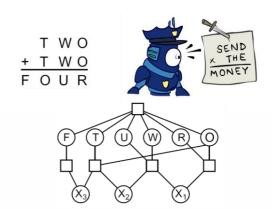




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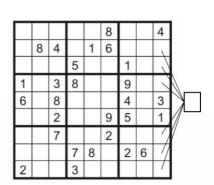
- Variables
 - Each (open) square
- Domains
 - {1,2,...,9}
- Constraints

1	3	5 8			9		
6	8				4		3
3000	2			9	5		1
	7			2			
		7	8		2	6	
2		3					

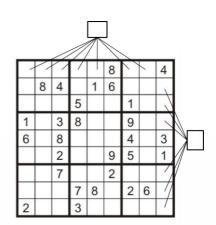
- Variables
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 - Unary constraints for given values

1 6	8	3 8 2	5 8	1	8 6	1 9 4 5		3
		7			2			1
			7	8		2	6	
2			3					

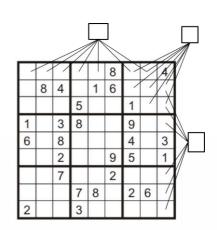
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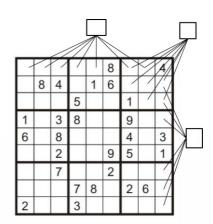
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 - Can also have a bunch of pairwise inequalities



Varieties of CSPs and Constraints



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- Discrete Variables
 - Finite domains
 - Size d means $O(d^n)$ complete assignments



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Continuous Variables

• E.g., start/end times for Hubble Telescope observations



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 SA ≠ WA
- Higher-order constraints involve 3 or more variables, e.g., cryptarithmetic column constraints

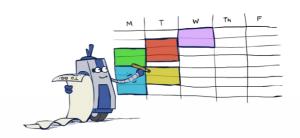


- Unary constraints involve a single variable (equivalent to reducing domains), e.g.: $SA \neq green$
- Binary constraints involve pairs of variables, e.g.: $SA \neq WA$
- Higher-order constraints involve 3 or more variables, e.g., cryptarithmetic column constraints
- Preferences (soft constraints)
 - E.g., red is better than green
 - Often representable by a cost for each variable assignment
 - Gives constrained optimization problems
 - (We'll ignore these until we get to Bayes' nets)



Real-World CSPs

- Scheduling problem
- Timetabling problem
- Assignment problem
- Hardware configuration
- Transportation scheduling
- Factory scheduling Circuit
- layout
- Fault diagnosis
- ... lotsmore!
- Many real-world problems involve real-valued variables...



Solving CSPs



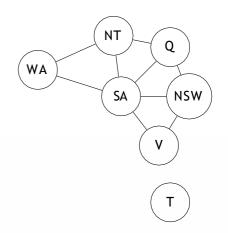
Standard Search Formulation

- States defined by the values assigned so far (partial assignments)
 - Initial state: the empty assignment, {}
 - Successor function: assign a value to an unassigned variable
 - Goal test: the current assignment is complete and satisfies all constraints



Search Methods

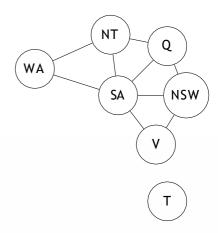
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Website: <u>simple-naive</u>

Search Methods

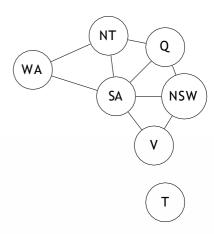
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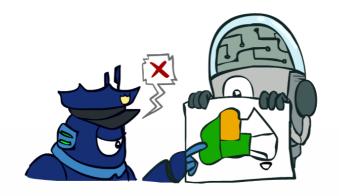
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Search Methods

- What would BFS do?
- What would DFS do?
- What problems does naïve search have?



Website: simple -naive



■ Backtracking search is the basic uninformed algorithm for solving CSPs

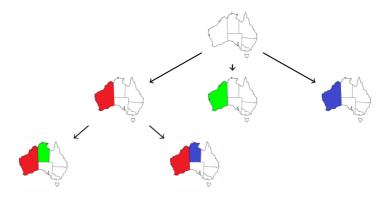
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 - i.e., [WA = red then NT = green] same as [NT = green then WA = red]
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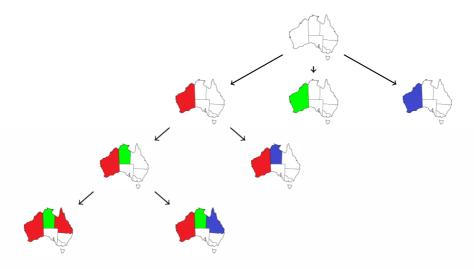
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 - "Incremental goal test"
- Depth-first search with these two improvements is called backtracking search









```
function BACKTRACKING-SEARCH(csp) returns a solution, or failure
  return RECURSIVE-BACKTRACKING(?, csp)
function RECURSIVE-BACKTRACKING(assignment, csp) returns a solution, or failure
  if assignment is complete then return assignment
  var \leftarrow SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUE(var, assignment, csp) do
     if value is consistent with assignment given CONSTRAINTS[csp] then
       add {var= value} to assignment
       result \leftarrow RECURSIVE-BACKTRACKING(assignment, csp)
       if result \( \pm \) failure then return result
       remove {var= value} from assignment
  return failure
```

■ Backtracking = DFS + variable-ordering + fail-on-violation

Website: simple -backtracking

Improving Backtracking

- General-purpose ideas give huge gains in speed
- Filtering: Can we detect inevitable failure early?
- Ordering:
 - Which variable should be assigned next?
 - In what order should its values be tried?
- Structure: Can we exploit the problem structure?

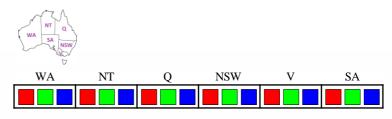


Filtering

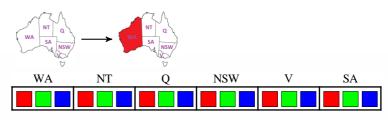


Filtering: Forward Checking

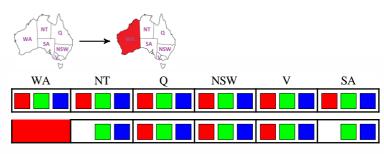
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- Forward checking: Cross off values that violate a constraint when added to the existing assignment



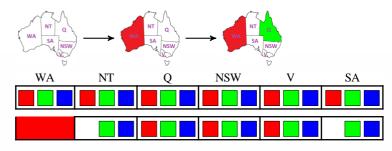
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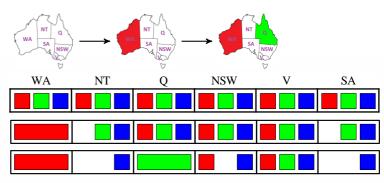


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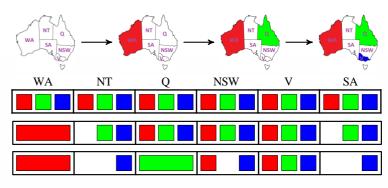


Website: <u>simple</u> - backtracking, forward

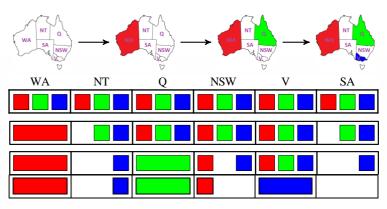
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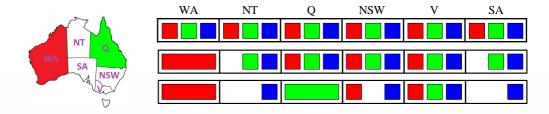
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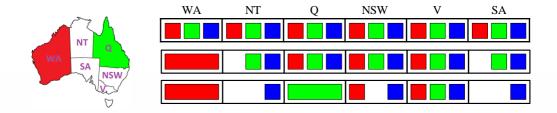
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Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:

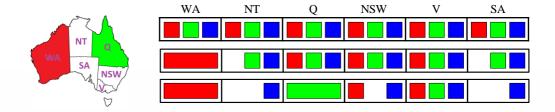


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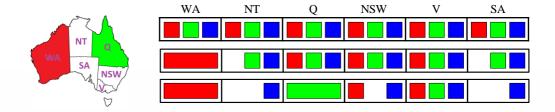
NT and SA cannot both be blue!

■ Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:

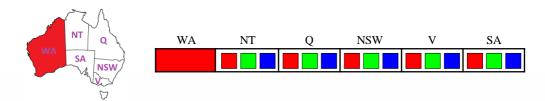


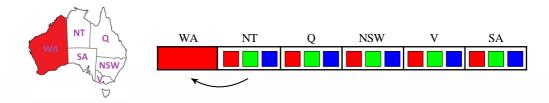
- NT and SA cannot both be blue!
- Why didn't we detect this yet?

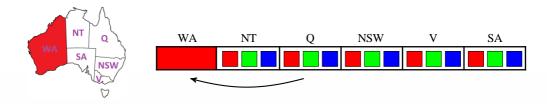
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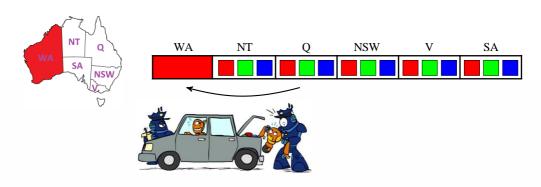


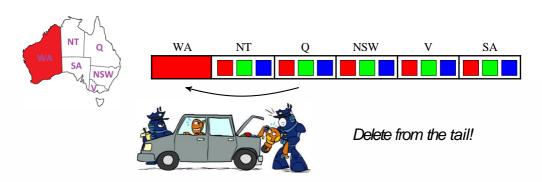
- NT and SA cannot both be blue!
- Why didn't we detect this yet?
- Constraint propagation: reason from constraint to constraint



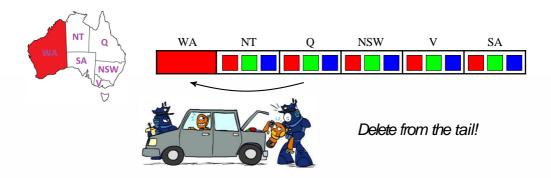






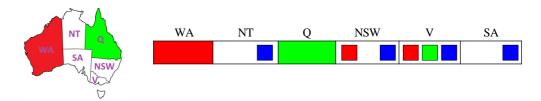


An arc $X \to Y$ is consistent iff for every x in the tail there is some y in the head which could be assigned without violating a constraint

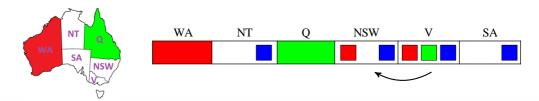


Forward checking: Enforcing consistency of arcs pointing to each new assignment

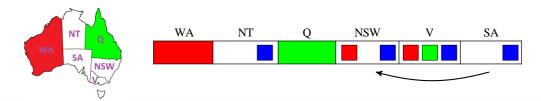
■ A simple form of propagation makes sure all arcs are consistent:



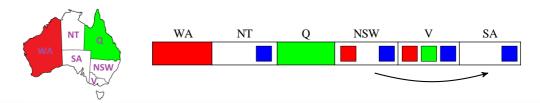
■ A simple form of propagation makes sure all arcs are consistent:



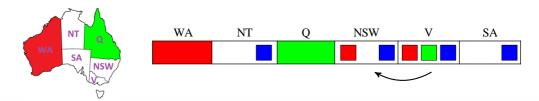
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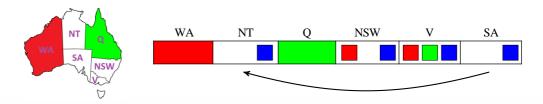
■ A simple form of propagation makes sure all arcs are consistent:



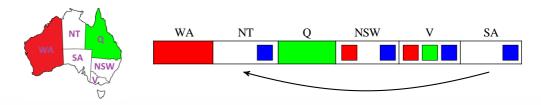
■ A simple form of propagation makes sure all arcs are consistent:



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■ A simple form of propagation makes sure all arcs are consistent:



- Important: If X loses a value, neighbors of X need to be rechecked!
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment
- What's the downside of enforcing arc consistency?

Enforcing Arc Consistency in a CSP

```
function AC-3(csp) returns the CSP, possibly with reduced domains
  inputs: CSP, a binary CSP with variables \{X_1, X_2, \dots, X_N\}
  local variables: queue, a queue of arcs, initially all the arcs in csp
  while queue is not empty do
     (X_i, X_i) \leftarrow \mathsf{REMOVE}\text{-}\mathsf{FIRST}(queue)
     if REMOVE-INCONSISTENT-VALUES (X_i, X_i) then
        for each X_k in NEIGHTBORS[X_i] do
           add (X_k, X_i) to queue
function REMOVE-INCONSISTENT-VALUES (X_i, X_i) returns true iff succeeds
  removed ← false
  for each x in DOMAIN[X_i] do
     if no value y in DOMAIN[X<sub>i</sub>] allows (x, y) to satisfy the constraint X_i \leftrightarrow X_i
        then delete x from DOMAIN[X_i]; removed \leftarrow true
  return removed
```

Applet: CSP - fiveQueens

Suggested Reading

Russell & Norvig: Chapter 6.1