

Consider a list **A** of length n that is already sorted, and set $low = 0$ and $high = n-1$ in the quick sort call. The sublist would then be divided into 2 sublists of lengths 0 and $n-1$, and the $(n-1)$ -length sublist will be divided similarly (into sublists of lengths 0 and $n-2$), and this will continue recursively until we get to 2 sublists of length 0 and 1, i.e. there will be $n-1$ recursive calls to `recQuickSort` within the main call in `quickSort`.

This means that *partition* will then have to be run on n sublists of decreasing (by 1) length. Note that since $low = 0$ and $high = n-1$, the left marker is shifted all the way to the right of the list, so the inner while loop performing this shift takes $O(n-i)$ time for the i th recursive call and the inner while loop for the shifting of right marker only takes $O(1)$ time as the right marker will then be smaller than the left one. The outer while loop will terminate after this, and hence for the i th recursive call to `recQuickSort`, *partition* will take $O(n - i)$ time and so `recQuickSort` will also take $O(n - i)$ time.

Hence, since we have a total of n calls to `recQuickSort` in the worst case, we get a worst case time for quick sort of:

$$\sum_{i=1}^n (n - i) = n(n) - \sum_{i=1}^n i = n^2 - n = O(n^2)$$