

# Linked Stacks and Queues

# 4

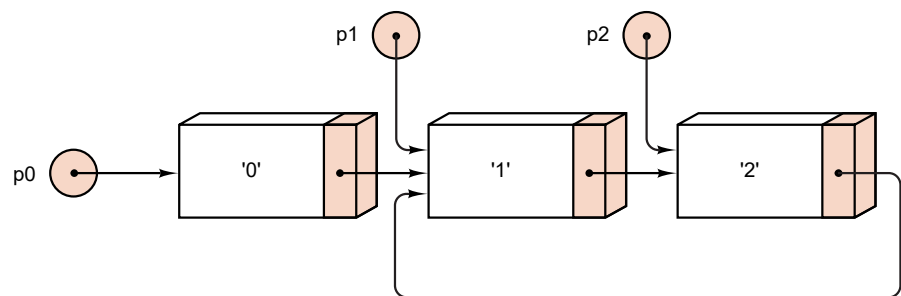
## 4.1 POINTERS AND LINKED STRUCTURES

### Exercises 4.1

**E1.** Draw a diagram to illustrate the configuration of linked nodes that is created by the following statements.

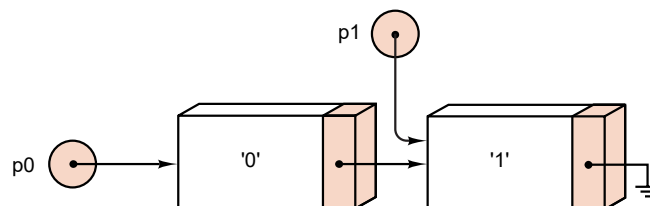
```
Node *p0 = new Node('0');  
Node *p1 = p0->next = new Node('1');  
Node *p2 = p1->next = new Node('2', p1);
```

*Answer*



**E2.** Write the C++ statements that are needed to create the linked configuration of nodes shown in each of the following diagrams. For each part, embed these statements as part of a program that prints the contents of each node (both data and next), thereby demonstrating that the nodes have been correctly linked.

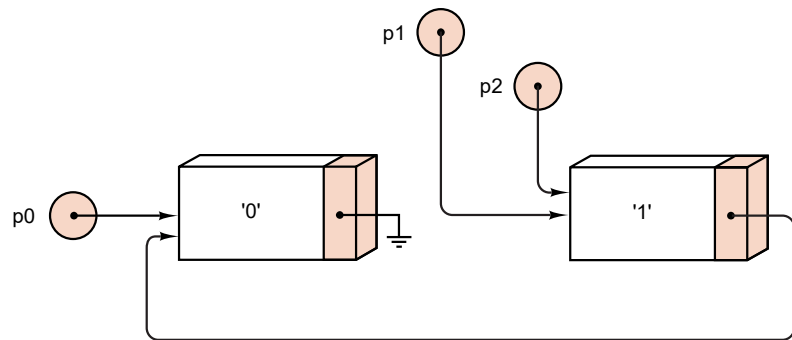
(a)



*Answer*

```
Node *p0 = new Node('0');  
Node *p1 = p0->next = new Node('1');
```

(b)

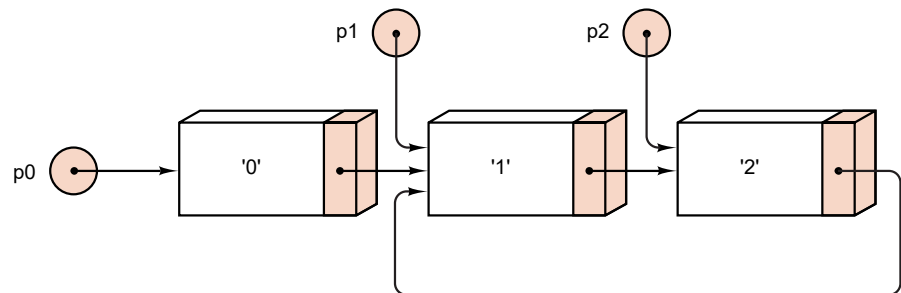


**Answer**

```

Node *p0 = new Node('0');
Node *p1 = new Node('1', p0);
Node *p2 = p1;
  
```

(c)



**Answer**

```

Node *p0 = new Node('0');
Node *p1 = p0->next = new Node('1');
Node *p2 = p1->next = new Node('2', p1);
  
```

## 4.2 LINKED STACKS

### Exercises 4.2

**E1.** Explain why we cannot use the following implementation for the method `push` in our linked Stack.

```

Error_code Stack::push(Stack_entry item)
{
    Node new_top(item, top_node);
    top_node = new_top;
    return success;
}
  
```

**Answer** There is a type incompatibility in the assignment statement `top_node = new_top`. Moreover, even if we corrected this incompatibility with the assignment `top_node = &new_top`, the statically acquired node would disappear at the end of the function, and the remaining nodes of the Stack would become garbage.

**E2.** Consider a linked stack that includes a method `size`. This method `size` requires a loop that moves through the entire stack to count the entries, since the number of entries in the stack is not kept as a separate member in the stack record.

```

typedef int Stack_entry;
#include "stack.h"
#include "stack.c"
int ackermann(int m, int n)
/* Pre:  Integers n and k are both at least 0.
   Post: The Ackermann function A(m, n) is returned. */
{
    Stack s;
    s.push(m);
    s.push(n);
    while (1) {
        s.top(n); s.pop();
        if (s.top(m) == underflow) return n; s.pop();
        if (m == 0) s.push(n + 1);
        else if (n == 0) {
            s.push(m - 1);
            s.push(1);
        }
        else {
            s.push(m - 1);
            s.push(m);
            s.push(n - 1);
        }
    }
}

```

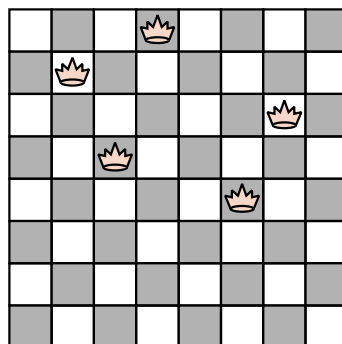
## 5.3 BACKTRACKING: POSTPONING THE WORK

### Exercises 5.3

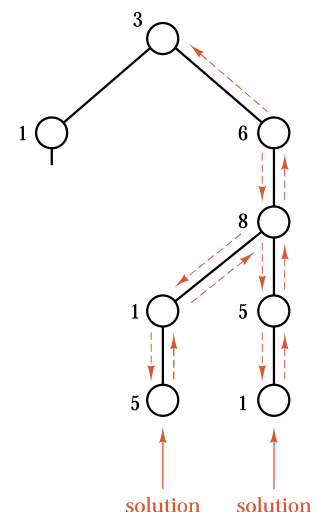
E1. What is the maximum depth of recursion in the function `solve_from`?

**Answer** The maximum depth of recursion occurs when the solution for the eight-queens problem has been found. Hence the maximum depth of recursion is eight.

E2. Starting with the following partial configuration of five queens on the board, construct the recursion tree of all situations that the function `solve_from` will consider in trying to add the remaining three queens. Stop drawing the tree at the point where the function will backtrack and remove one of the original five queens.



**Answer**



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