

## Homework #2

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## Question 5:

Solve the following questions from the Discrete Math zyBook.

### Part A.

#### Exercise 1.12.2:

##### Section B:

1.  $\neg q$  *Hypothesis*
2.  $\neg q \vee \neg r$  *Addition (1)*
3.  $\neg(q \wedge r)$  *De Morgan (2)*
4.  $p \rightarrow (q \wedge r)$  *Hypothesis*
5.  $\neg p$  *Modus Tollens (3, 4)*

##### Section E:

1.  $p \vee q$  *Hypothesis*
2.  $\neg q$  *Hypothesis*
3.  $p$  *Disjunctive Syllogism (1, 2)*
4.  $\neg p \vee r$  *Hypothesis*
5.  $r$  *Disjunctive Syllogism (3, 4)*

#### Exercise 1.12.3:

##### Section C:

1.  $p \vee q$  *Hypothesis*
2.  $\neg p$  *Hypothesis*
3.  $\neg p \wedge (p \vee q)$  *Conjunction (1, 2)*
4.  $(\neg p \wedge p) \vee (\neg p \wedge q)$  *Distributive (3)*
5.  $F \vee (\neg p \wedge q)$  *Complement (4)*
6.  $\neg p \wedge q$  *Identity (5)*
7.  $q$  *Simplification (6)*

### Exercise 1.12.5:

For this we can make three statements, as in the examples from sections (a) and (b).

j: I will get a job

c: I will buy a new car

h: I will buy a new house

#### Section C:

The form of the argument is:

$$\frac{(c \wedge h) \rightarrow j \quad \neg j}{\therefore \neg c}$$

This argument is invalid, if we do the truth table we can see that the fourth line invalidates it:

$c$	$h$	$j$	$(c \wedge h) \rightarrow j$	$\neg j$	$\neg c$
$T$	$T$	$T$	$F$	$F$	$F$
$T$	$T$	$F$	$T$	$T$	$F$
$T$	$F$	$T$	$T$	$F$	$F$
$T$	$F$	$F$	$T$	$T$	$F$
$F$	$T$	$T$	$T$	$F$	$T$
$F$	$T$	$F$	$T$	$T$	$T$
$F$	$F$	$T$	$T$	$F$	$T$
$F$	$F$	$F$	$T$	$T$	$T$

#### Section D:

The form of the argument is:

$$\frac{(c \wedge h) \rightarrow j \quad \neg j \quad h}{\therefore \neg c}$$

This argument is valid because of  $c = j = F$  and  $h = T$ , all three hypothesis are true and the conclusion is true.

## Part B.

### Exercise 1.13.3:

#### Section B:

The form of the argument is:

$$\frac{\begin{array}{c} \exists_x(P(x) \vee Q(x)) \\ \exists_x \neg Q(x) \end{array}}{\therefore \exists_x P(x)}$$

We can make a truth table for this and check:

$P(x)$	$Q(x)$
$F$	$F$
$F$	$T$

We can see that the if  $P(x)$  is false, and  $Q(x)$  is true, then  $\exists_x(P(x) \vee Q(x))$  is true and we also see if both  $P(x)$  and  $Q(x)$  are both false, then  $\exists_x \neg Q(x)$  is true. However the statement  $\exists_x P(x)$  is false in both, and therefore this argument is invalid.

### Exercise 1.13.5:

The domain for each problem is the set of students in a class.

#### Section D:

The two equations for this:

$M(x)$ :  $x$  missed class.

$D(x)$ :  $x$  got a detention.

We can set up equations for this as:

$$\frac{\begin{array}{c} \forall_x(M(x) \rightarrow D(x)) \\ \neg M(Penelope) \end{array}}{\therefore \neg D(Penelope)}$$

This argument is invalid when  $M(Penelope) = T$  and  $D(Penelope) = F$ , as the first hypothesis is thus false, but the conclusion is true due to universal instantiation.

#### Section E:

We can set up equations for this as, we also define  $A(x)$  in addition to above to indicate if  $x$  got an A.

$$\frac{\begin{array}{c} \forall_x(M(x) \vee D(x)) \rightarrow \neg A \\ Penelope \text{ is a student.} \\ A(p) \end{array}}{\therefore \neg D(Penelope)}$$

This is valid and we can use the rules of inference to prove it:

- |     |   |                                      |
|-----|---|--------------------------------------|
| 1.  | $\forall_x(M(x) \vee D(x)) \rightarrow \neg A$  | <i>Hypothesis</i>                    |
| 2.  | $Penelope \text{ is a student.}$                | <i>Hypothesis</i>                    |
| 3.  | $M(p) \vee D(p) \rightarrow \neg A(p)$          | <i>Universal Instillation</i> (1, 2) |
| 4.  | $(\neg((M(p) \vee D(p))) \vee \neg A(p))$       | <i>Conditional Identity</i> (3)      |
| 5.  | $(\neg(M(p) \wedge \neg(D(p))) \vee \neg A(p))$ | <i>De Morgan</i> (4)                 |
| 6.  | $\neg A(p) \vee (\neg M(p) \wedge \neg D(P))$   | <i>Commutative Law</i> (5)           |
| 7.  | $A(p) \rightarrow (\neg M(p) \wedge \neg D(p))$ | <i>Conditional Identity</i> (6)      |
| 8.  | $A(p)$  | <i>Hypothesis</i>                    |
| 9.  | $\neg M(p) \wedge \neg D(p)$                    | <i>Modus Ponens</i> (7, 8)           |
| 10. | $\neg D(p) \wedge \neg M(p)$                    | <i>Commutative</i> (9)               |
| 11. | $\neg D(p)$                                     | <i>Simplification</i> (10)           |

## Question 6:

Solve the following questions from the Discrete Math zyBook.

### Exercise 2.2.1:

C.

If  $x$  is a real number and  $x \leq 3$ , then  $12 - 7x + x^2 \geq 0$ .

$x \leq 3$	<i>Hypothesis</i>
$x - x \leq 3 - x = 0 \leq 3 - x$	<i>Subtract <math>x</math></i>
$0 \leq 3 - x$ and $0 \leq 4 - x$	<i>If left one is true, so is right.</i>
$(3 - x)(4 - x) \geq 0$	<i>Expand this out to get the following.</i>
$12 - 7x + x^2 \geq 0$	

D.

*The product of two odd integers is an odd integer.*

Assume that both  $x$  and  $y$  are both odd integers. Because  $x$  and  $y$  are both odd, then having two integers  $a$  and  $b$  then  $x = 2a + 1$  and  $y = 2b + 1$ . Then we can see that:

$$\begin{aligned} xy &= (2a + 1)(2b + 1) \\ xy &= 2(2ab + a + b) + 1 \end{aligned}$$

Because any number divided by 2 is even, then  $2(2ab + a + b)$  is even. So adding 1 to this must make  $xy$  odd.

## Question 7:

Solve the following questions from the Discrete Math zyBook.

### Exercise 2.3.1:

**D.**

*For every integer  $n$ , if  $n^2 - 2n + 7$  is even, then  $n$  is odd.*

Assuming that  $n$  is even, we can represent every even  $n$  as  $2k$ , where  $k$  is any arbitrary integer. Then it follows that:

$$\begin{aligned}n^2 - 2n + 7 &= (2k)^2 - 2(2k) + 7 \\n^2 - 2n + 7 &= 2(2k^2 - 2k) + 7\end{aligned}$$

Because  $k$  is an integer then  $2(2k^2 - 2k)$  is an even integer. Therefore adding 7 to that makes it an odd integer. Hence it proves the theorem, when  $n$  is even, the equation is odd, so the contrapositive must be true, in that the equation is odd when the  $n$  is odd.

**F.**

*For every non-zero real number  $x$ , if  $x$  is irrational, then  $\frac{1}{x}$  is also irrational.*

Assume that  $\frac{1}{x}$  is rational and  $x$  is nonzero. There must then be integers  $a$  and  $b$  where  $\frac{1}{x} = \frac{a}{b}$ . Then cross multiplying we have:

$$\begin{aligned}b &= ax \\x &= \frac{b}{a}\end{aligned}$$

$x$  is in the form of a rational number. Therefore the theorem is true due to the contrapositive, since  $x$  is irrational, the reciprocal is also irrational.

**G.**

*For every pair of real numbers  $x$  and  $y$ , if  $x^3 + xy^2 \leq x^2y + y^3$ , then  $x \leq y$ .*

We assume that  $x > y$ , therefore  $x^3 > x^2y$  and  $xy^2 > y^3$ . If we add these two statements up, we obtain that  $x^3 + xy^2 > x^2y + y^3$ . Which proves the theorem by being the contrapositive.

**L.**

*For every pair of real numbers  $x$  and  $y$ , if  $x + y > 20$ , then  $x > 10$  or  $y > 10$ .*

Assume that when  $x \leq 10$  and  $y \leq 10$  then  $x + y \leq 20$ . Then we assume that  $x = y = 10$ . Thus it follows that:

$$\begin{aligned}x + y &\leq 20 \\10 + 10 &\leq 20 \\20 &\leq 20\end{aligned}$$

Which ends up being true, therefore the theorem is true due to the contrapositive.

## Question 8:

Solve the following questions from the Discrete Math zyBook.

### Exercise 2.4.1:

**C. *The average of three real numbers is greater than or equal to at least one of the numbers.***

Proving this by contradiction, we can assume that there are three real numbers,  $x$ ,  $y$  and  $z$  such that the average of the them is less than each of the 3 numbers. So we have:  $(x + y + z)/3 < x$ ,  $(x + y + z)/3 < y$ ,  $(x + y + z)/3 < z$ . Adding up the three equations together, we have:  $(x + y + z) < x + y + z$ , which cannot be true, since a number cannot be less than itself, therefore the theorem must be true.

**E. *There is no smallest integer.***

Proving this by contradiction, let us assume there is a smallest integer and we will call it  $n$ . If  $n$  is the smallest integer, then nothing can be smaller/lesser than it. If we subtract 1 from an integer, we also get an integer. Then we get that there must exist an integer  $n - 1$  and we know that  $n - 1 < n$ , therefore we know that there is no integer  $n$  smaller than itself less one. Therefore the theorem must be true.

## Question 9:

Solve the following questions from the Discrete Math zyBook.

### Exercise 2.5.1:

#### Section C:

If integers  $x$  and  $y$  have the same parity, then  $x + y$  is even. The parity of a number tells whether the number is odd or even. If  $x$  and  $y$  have the same parity, they are either both even or both odd.

#### Case 1: Even Integers

If we assume both  $x$  and  $y$  are both even, then integers  $a$  and  $b$  exist such that  $x = 2a$  and  $y = 2b$ . Therefore:

$$\begin{aligned}x + y &= 2a + 2b \\x + y &= 2(a + b)\end{aligned}$$

Because all integers when multiplied by 2 are even, then  $x + y$  is also even, since that is equal to  $2(a + b)$ .

#### Case 2: Odd Integers

If we assume both  $x$  and  $y$  are both odd, then integers  $a$  and  $b$  exist such that  $x = 2a + 1$  and  $y = 2b + 1$ . Therefore:

$$\begin{aligned}x + y &= 2a + 2b + 1 + 1 \\x + y &= 2a + 2b + 2 \\x + y &= 2(a + b + 1)\end{aligned}$$

Because all integers when multiplied by 2 are even, then  $x + y$  is also even, since that is equal to  $2(a + b + 1)$ .