

Quality test for Relational DBs

Reading: Elmasri & Navathe, Fundamentals of Database Systems, Chapter 14.





Recap: Normal Forms

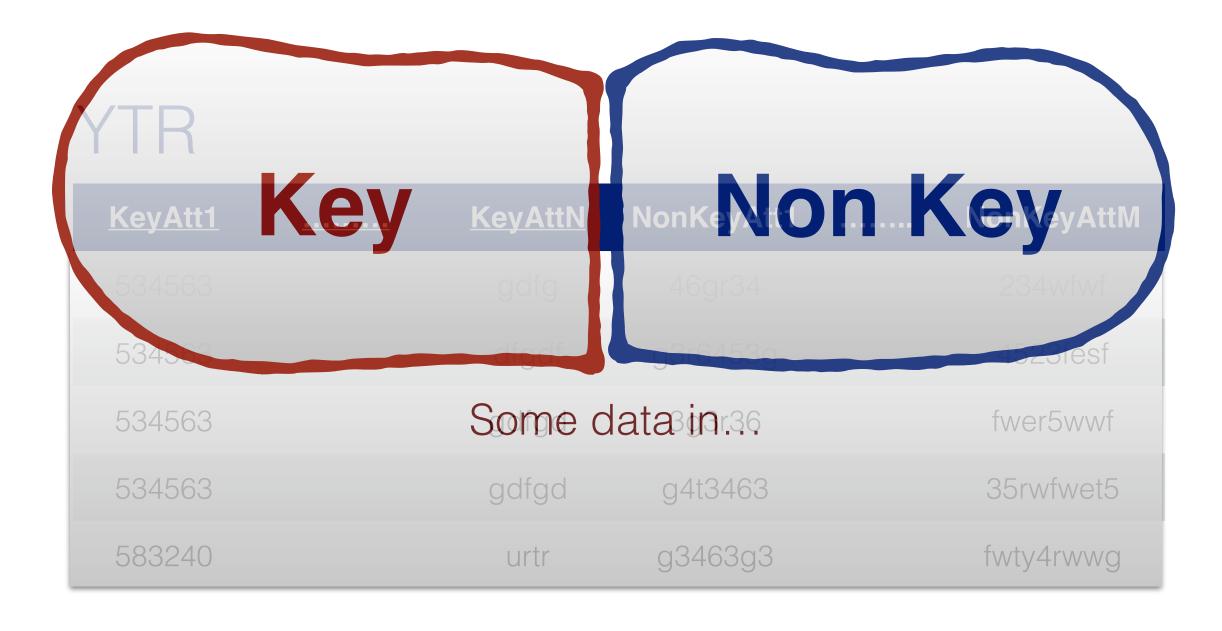
- Based on functional dependencies:
 - 1st Normal Form (1NF)
 - 2nd Normal Form (2NF)
 - 3rd Normal Form (3NF)
 - Boyce-Codd Normal Form (BCNF)
- Based on "multivalued" and "join" dependencies:
 - 4th and 5th Normal Forms (4NF, 5NF)
- Each NF extends the previous (so a relation in 3NF is also in 2NF and 1NF)



A graphic guide to FD based Normal Forms

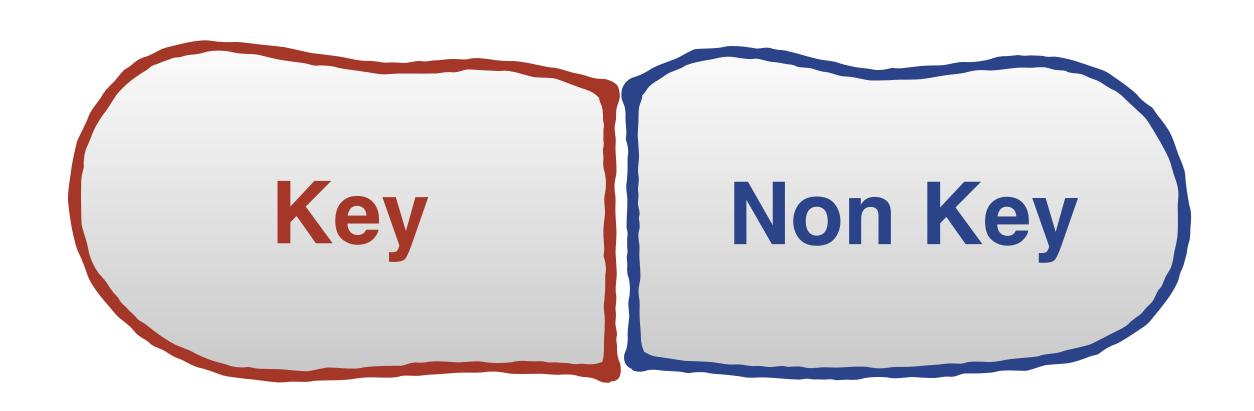
consider Your Typical Relation:

YTR = {KeyAtt1, KeyAtt2,...KeyAttN, NonKeyAtt1,....NonKeyAttM}





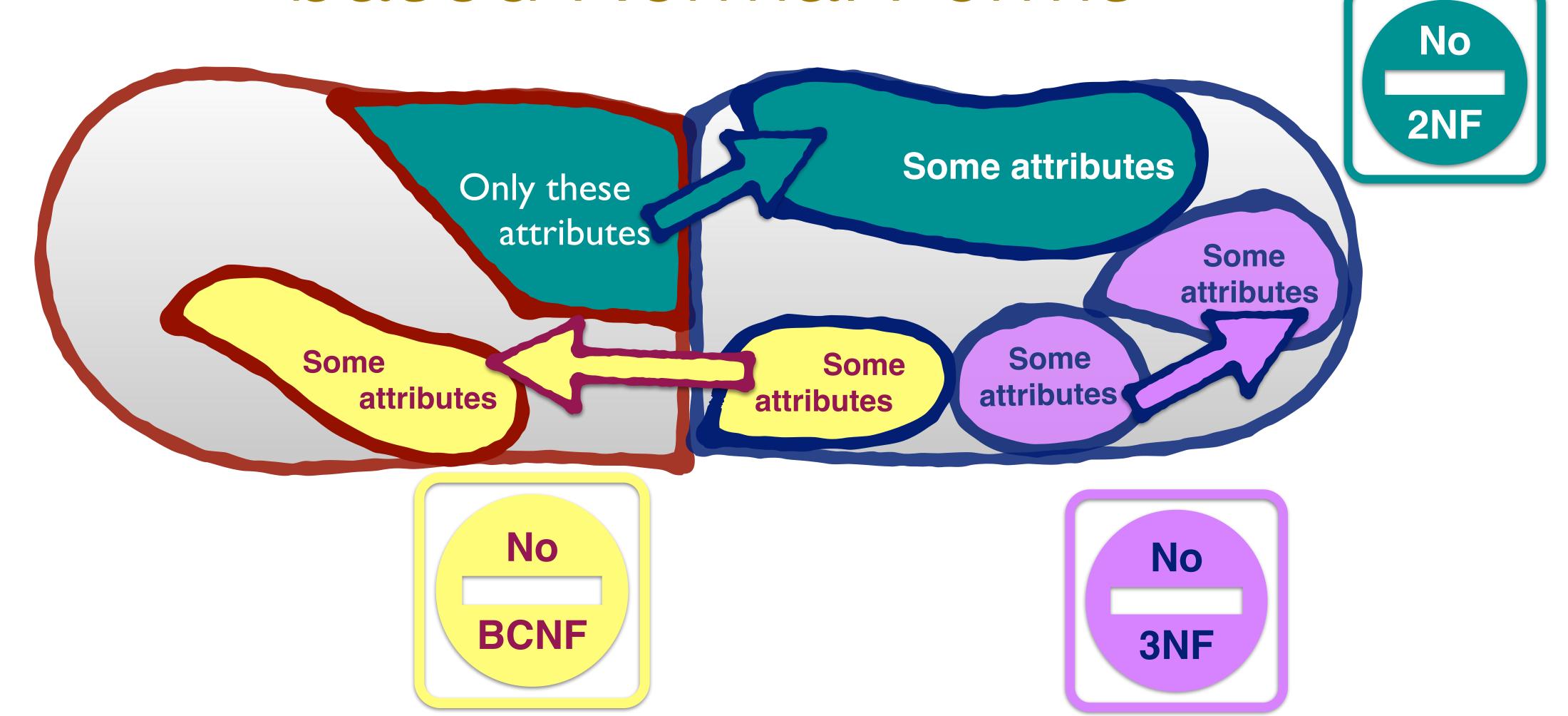
A graphic guide to FD based Normal Forms





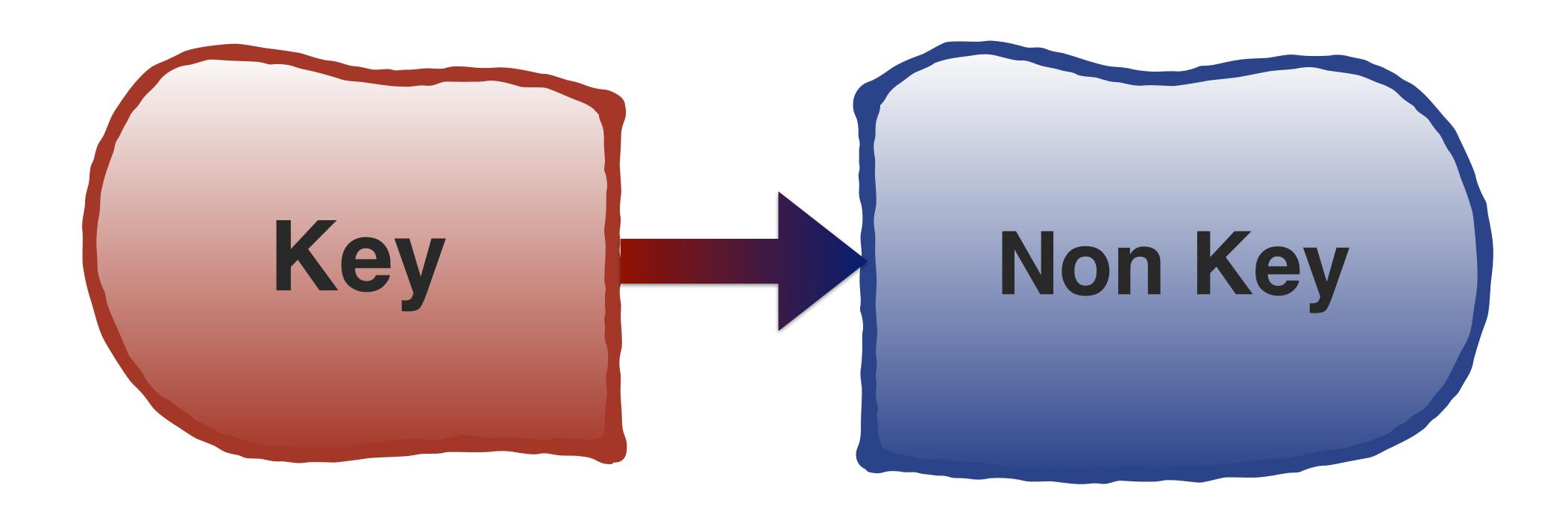


A graphic guide to FD based Normal Forms





BCNF (and 3NF, and 2NF)



- 1. The whole key is needed to determine ALL non key attributes.
- 2. There are **NO OTHER dependencies**.



Normal Forms: not enough

- When considered in isolation, NFs do not guarantee a good design
- Important additional properties to check:
 - Dependency preservation property: Each functional dependency is represented in some relation after normalisation.
 Desirable but can be sacrificed for other factors.
 - Non-additive join property: No spurious tuples are generated after normalisation. Extremely critical.
- Algorithms can be used to guarantee that a set of relations satisfies either or both of the properties above



REMEMBER OUR **PROBLEMATIC** TABLE?

StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein

Tutor

Einstein

Mozart

Darwin

Bohr

Department

Physics

Music

Biology

Physics

StudentID	Department
123	Physics
123	Music
456	Biology
789	Physics
999	Physics

StudentID

COMP

Students

StudentID	Department	Tutor
123	Physics	Einstein
123	Music	Mozart
456	Biology	Darwin
789	Physics	Bohr
999	Physics	Einstein

let's test them all

Tutor	Department
Einstein	Physics
Mozart	Music
Darwin	Biology
Bohr	Physics

StudentID	Tutor
999	Physics
789	Physics
456	Biology
123	Music
123	Physics

Department

this was non BCNF 3 alternative solutions

StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein





Dependency preservation

Students

StudentID	Department	Tutor
123	Physics	Einstein
123	Music	Mozart
456	Biology	Darwin
789	Physics	Bohr
999	Physics	Einstein

Department

Physics

- Functional Dependencies:
 - 1. {StudentID, Department}→ Tutor
 - 2. Tutor → Department

Einstein	Physics
Mozart	Music
Darwin	Biology

Bohr

Tutor

StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein

- This set of tables is BCNF
- However: Functional Dependency 1. is **not** preserved
- There is no way we can tell it holds



Non-additive join

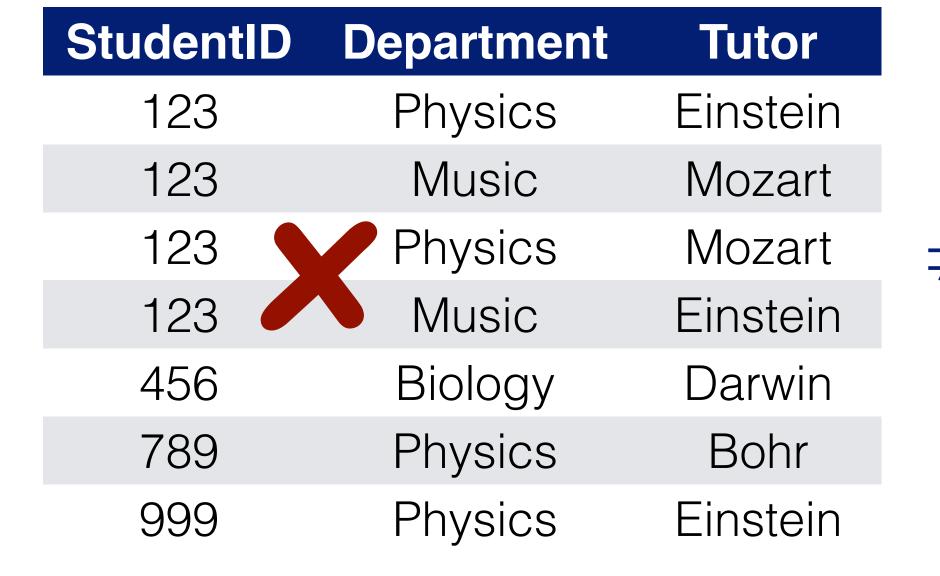
1.

StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein

StudentID	Department
123	Physics
123	Music
456	Biology
789	Physics
999	Physics

This BCNF decomposition does not have the **non-additive join property** as spurious tuples can be produced

R1 join R2 =



Students

StudentID	Department	Tutor
123	Physics	Einstein
123	Music	Mozart
456	Biology	Darwin
789	Physics	Bohr
999	Physics	Einstein



Testing for non-additive join

- Problem: relying on finding examples is not the best strategy
- We need a more systematic procedure
- "Tableaux test for non-additive join"
 - starting from the "universal relation" and the set of FDs
 - gives an algorithm for easy checking
 - if the test succeeds, we can be sure ALL instances of the database will never generate spurious tuples
 - if the test fails, it also gives the counterexample we need

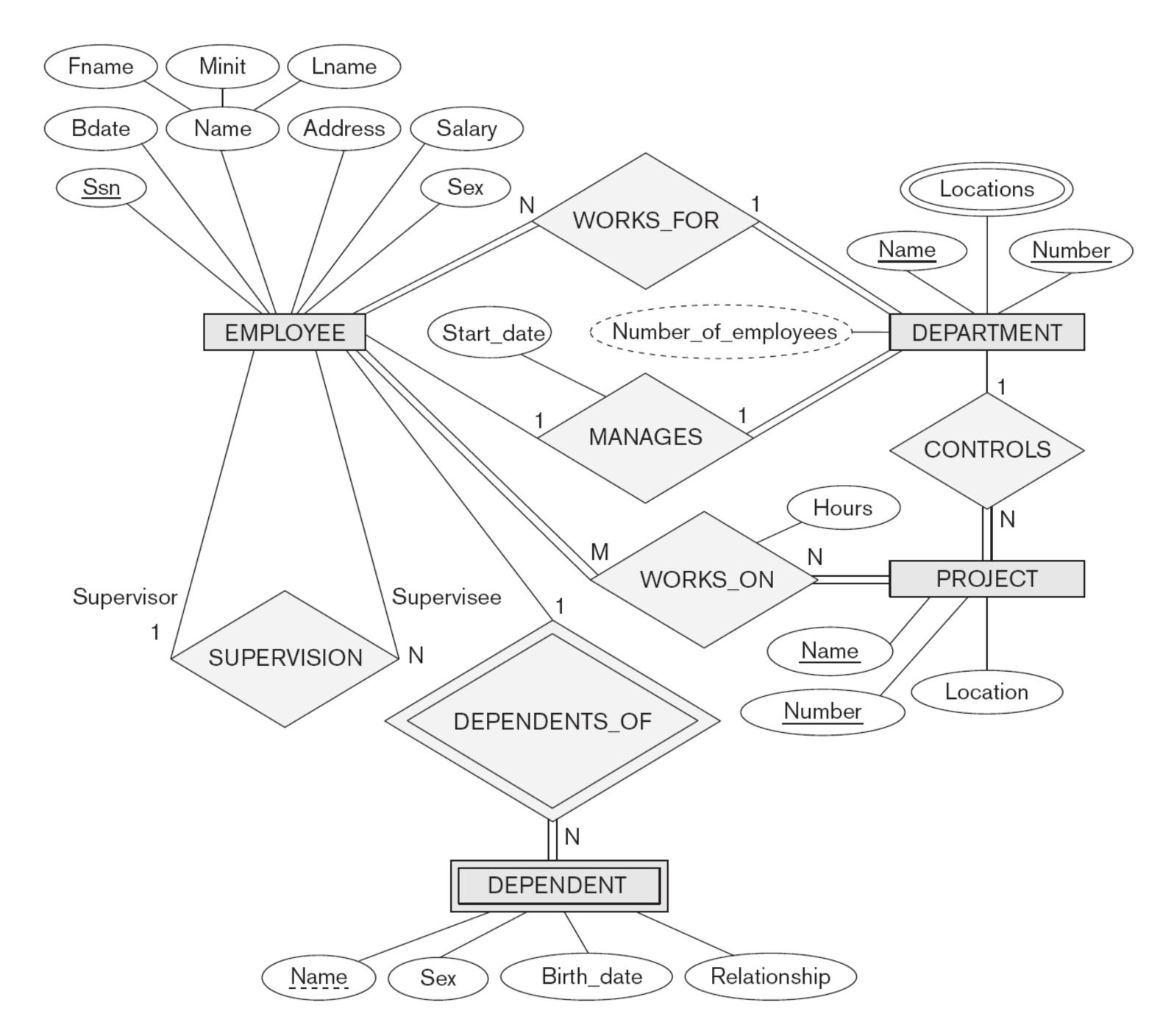




Tableaux test

- Principle: simulate what would happen if a report was created that joined *all tables* in the database
- One would assume that only very few individuals are such that their relationships can be traced through all tables
- If a join over all tables produced a potentially infinite amount of such individuals, we would be in the presence of spurious records
- We would like to test this, without having to populate the database with 1000s of records in the hope of finding a good example by chance

In the company example, this would mean to search for all employees **who** work for a department, **and** manage a department, **and** work on a project, **which is** controlled by a department, **and** have a supervisor, **and** have a dependent.







We can simulate this by building a table containing **ALL** attributes of the database (the "Universal Relation") and **one row for each** table (relation) of the DB, and filling it in with generic values, arbitrary, **but constrained by the functional dependencies**

	Attribute I	Attribute 2	• • •	• • •	• • •	Attribute N
Relation I						
Relation 2						
Relation M						



Let's test our example

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Tutor	Department
Einstein	Physics
Mozart	Music
Darwin	Biology
Bohr	Physics

StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein

- The Universal relation is {StudentID, Tutor, Department}
- The Relations are R1={Tutor, Department} and R2={StudentID,Tutor}
- The only functional dependency that we can represent here is
 Tutor → Department

- R1={Tutor, Department} and R2={StudentID, Tutor}
- FD: Tutor → Department

1. Prepare the table:

	StudentID	Tutor	Department
Rı			
R ₂			

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R1={Tutor, Department} and R2={StudentID, Tutor}

FD: Tutor → Department

2. Fill each row with generic values, ensuring consistency:

We use each relation in turn and add values to "activated" attributes in each row

	StudentID	Tutor	Department
Rı		Value I	Value 2
R_2			

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- R1={Tutor, Department} and
- R2={StudentID,Tutor}
 - FD: Tutor → Department

2. Fill each row with generic values, ensuring consistency:

We use each relation in turn and add values to "activated" attributes in each row

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3	Value I	

- R1={Tutor, Department} and R2={StudentID, Tutor}
- FD: Tutor → Department

2. Fill each row with generic values, ensuring consistency:

We use each relation in turn and add values to "activated" attributes in each row

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3	Value I	

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- R1={Tutor, Department} and R2={StudentID, Tutor}
- FD: Tutor → Department

"All tuples with same Tutor must have the same Department"

3. Apply the Functional Dependencies in turn

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3	Value I	Value 2

- R1={Tutor, Department} and R2={StudentID, Tutor}
- FD: Tutor → Department

If we manage to complete a row of values, we pass the test!

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3	Value I	Value 2

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- R1={Tutor, Department} and R2={StudentID, Tutor}
- FD: Tutor → Department

If we manage to complete a row of values, we pass the test!

This DB is safe from producing spurious records



	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3	Value I	Value 2



Let's test another example

2

Tutor	Department	StudentID	Department
Einstein	Physics	123	Physics
Mozart	Music	123	Music
Darwin	Biology	456	Biology
Bohr	Physics	789	Physics
		999	Physics

- The Universal relation is {StudentID, Tutor, Department}
- The Relations are R1={Tutor, Department} and R2={StudentID,Department}
- The functional dependency that we can represent here is only:
 Tutor → Department

- R1={Tutor, Department} and R2={StudentID,Department}
- FD: Tutor → Department

1. Prepare the table:

	StudentID	Tutor	Department
Rı			
R ₂			

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R1={Tutor, Department} and R2={StudentID,Department}

FD: Tutor → Department

2. Fill each row with generic values, ensuring consistency:

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂			



- R1={Tutor, Department} and
 R2={StudentID,Department}
 - FD: Tutor → Department
 - 2. Fill each row with generic values, ensuring consistency:

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2

- R1={Tutor, Department} and R2={StudentID,Department}
- FD: Tutor → Department
- 2. Fill each row with generic values, ensuring consistency:

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2



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- R1={Tutor, Department} and R2={StudentID,Department}
- FD: Tutor → Department

"All tuples with same Tutor must have the same Department"

3. Apply the Functional Dependencies in turn

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2



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- R1={Tutor, Department} and R2={StudentID,Department}
- FD: Tutor Department

There are no 2 tuples where we can apply this FD

3. Apply the Functional Dependencies in turn

	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2



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- R1={Tutor, Department} and R2={StudentID,Department}
- FD: Tutor Department

There are no 2 tuples where we can apply this FD

If we manage to complete a row of values, we pass the test!

This DB could produce spurious records



	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2



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	StudentID	Tutor	Department
Rı		Value I	Value 2
R ₂	Value 3		Value 2



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	StudentID	Tutor	Department
Rı		Value I	CompSci
R ₂	Value 3		CompSci



COMP
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	StudentID	Tutor	Department
Rı		Turing	CompSci
R ₂	Value 3	Hopper	CompSci





StudentID	Tutor	Department
Stl	Turing	CompSci
St2	Hopper	CompSci

R1={Tutor, Department}

Tutor	Department
Turing	CompSci
Hopper	CompSci

R2={StudentID,Department}

StudentID	Department
Stl	CompSci
St2	CompSci





StudentID	Department
Stl	CompSci
St2	CompSci



Tutor	Department
Turing	CompSci
Hopper	CompSci

StudentID	Department	Tutor
Stl	CompSci	Turing
St2	Corpoci	Turing
Stl	Compsci	Hopper
St2	CompSci	Hopper

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StudentID	Tutor	Department
Stl	Turing	CompSci
St2	Hopper	CompSci



What about solution 1?

1.

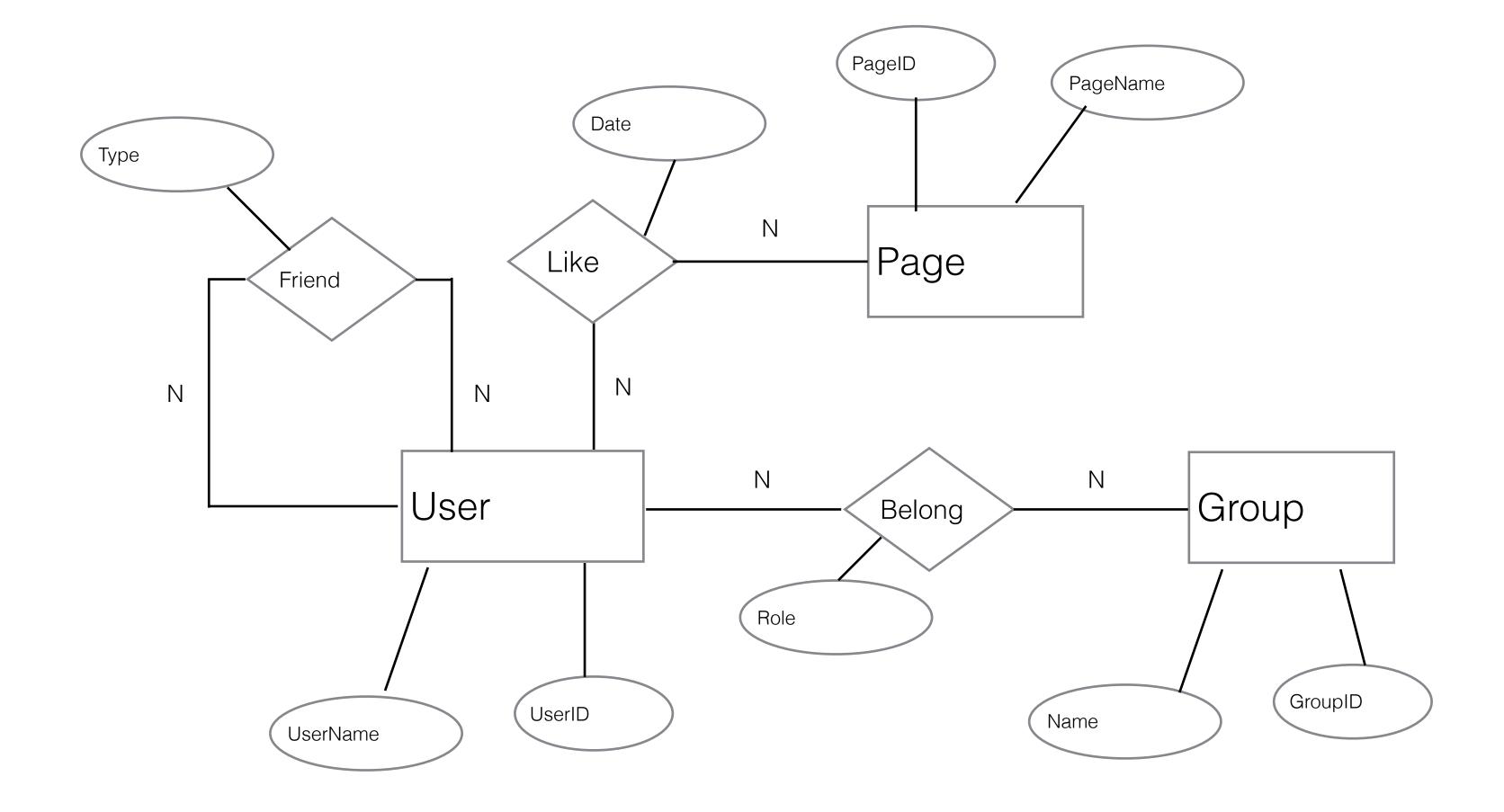
StudentID	Tutor
123	Einstein
123	Mozart
456	Darwin
789	Bohr
999	Einstein

StudentID	Department
123	Physics
123	Music
456	Biology
789	Physics
999	Physics

- The Universal relation is {StudentID, Tutor, Department}
- The Relations are R1={StudentID, Tutor} and R2={StudentID,Department}
- The functional dependencies that we can represent here are...
 - NONE!
- The decomposition fails the test immediately



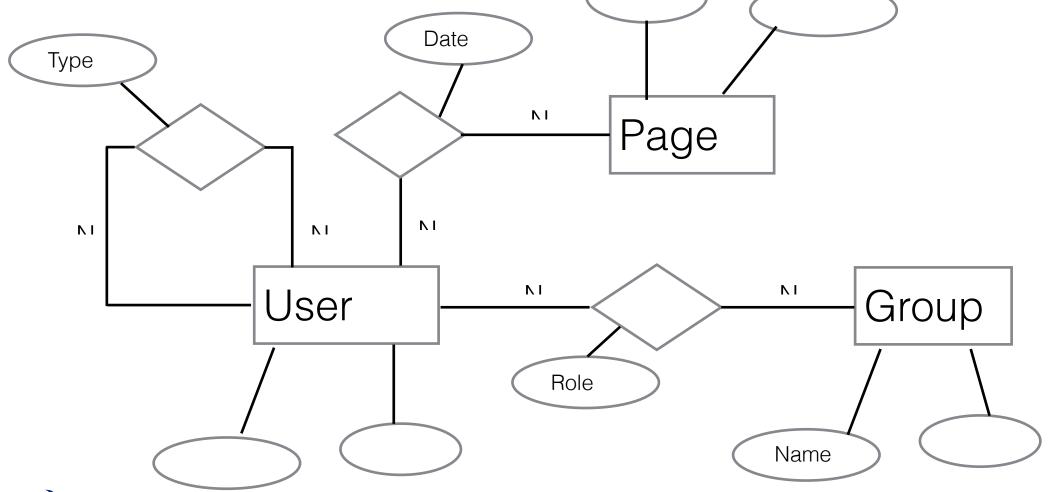
No design is waterproof!



A simplified social network!







Relations:

- R1={<u>UserID</u>, Username}
- R2={GroupID, GroupName}
- R3={<u>PageID</u>, PageName}
- R4={UserID,PageID,Date} (user likes page)
- R5={<u>UserID,GroupID</u>,Role} (user belongs to group)
- R6={<u>UserID,FriendID</u>,Type} (*user friends user*)

Role

User

Group

Name

Functional dependencies

- UserID → Username
- GroupID → GroupName
- PageID → PageName
- {UserID,PageID} → Date
- {UserID, GroupID} → Role
- {UserID,FriendID} → Type

- BCNF set of tables
- and we also can preserve all dependencies



Construct the rows, one for each relation:

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I				Value5		Value7			
R5	Value I		Value3					Value8		
R6	Value I								Value9	Value 10



UserID → Username

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I				Value5		Value7			
R5	Value I		Value3					Value8		
R6	Value I								Value9	Value I 0



UserID → Username

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5		Value7			
R5	Value I	Value2	Value3					Value8		
R6	Value I	Value2							Value9	Value I 0



• GroupID → GroupName

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5		Value7			
R5	Value I	Value2	Value3					Value8		
R6	Value I	Value2							Value9	Value I 0



• GroupID → GroupName

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5		Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



PageID → PageName

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5		Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



PageID → PageName

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value 10



• {UserID,PageID} → Date

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2			> <					
R2			Value3	Value4						
R3	X				Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



• {UserID,GroupID} → Role

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2	>							
R2	>		Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



• {UserID,FriendID} → Type

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7		>	
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



This DB could produce spurious records!!!



	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



The counterexample

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0



The counterexample

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	UI	John							••	••
R2			GI	Bikers					••	••
R3					PI	StarWars			••	••
R4	UI	John			PΙ	StarWars	1.4.19		••	••
R5	UI	John	GI	Bikers				Treasur er	••	
R6	UI	John								



The counterexample

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	UI	John	G2	Geeks	P2	Pizza	1.1.19	Admin	••	••
R2	U2	Mary	GI	Bikers	P3	Liverpool	1.2.19	Owner	••	••
R3	U3	Sarah	G3	Chefs	PI	StarWars	1.3.19	Fan	••	••
R4	UI	John	G4	Travel	PI	StarWars	1.4.19	Support er	••	••
R5	UI	John	GI	Bikers	P4	Van Gogh	1.5.19	Treasur er	••	
R6	UI	John	G5	MovieClub	P5	Java!	1.6.19	Chief		



	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	Frie	Туре
RI	UI	John	G2	Geeks	P2	Pizza	1.1.19	Admin	••	••
R2	U2	Mary	GI	Bikers	P3	Liverpool	1.2.19	Owner	••	••
R3	U3	Sarah	G3	Chefs	PI	StarWars	1.3.19	Fan	••	••
R4	UI	John	G4	Travel	PI	StarWars	1.4.19	Supporter	••	••
R5	UI	John	GI	Bikers	P4	Van Gogh	1.5.19	Treasurer	••	••
R6	UI	John	G5	MovieClub	P5	Java!	1.6.19	Chief	• • •	

UserID	Username
UI	John
U2	Mary
U3	Sarah

GroupID	GroupName
GI	Bikers
G2	Geeks
G3	Chefs
G4	Travel
G5	MovieClub

PageID	PageName
PI	StarWars
P2	Pizza
P3	Liverpool
P4	Van Gogh
P5	Java!

UserID	GroupID	Role
UI	G2	Admin
U2	G۱	Owner
U3	G3	Fan
UI	G4	Supporter
UI	GI	Treasurer
UI	G5	Chief

UserID	PageID	Date
UI	P2	1.1.19
U2	P3	1.2.19
U3	PI	1.3.19
UI	PI	1.4.19
UI	P4	1.5.19
UI	P5	1.6.19

- looks reasonable!
 - what you would expect
- BUT!





	UserID	Username	PageID	Date	PageName	GroupID	Role	GroupName
	UI	John	P2	1.1.19	Pizza	G2	ADMIN	Geeks
	UI	John	P2	1.1.19	Pizza	G4	SUPPORTER	Travel
	UI	John	P2	1.1.19	Pizza	GI	TREASURER	Bikers
	UI	John	P2	1.1.19	Pizza	G5	CHIEF	MovieClub
UserIE	UI	John	PI	1.4.19	StarWars	G2	ADMIN	Geeks
UI	UI	John	PI	1.4.19	StarWars	G4	SUPPORTER	Travel
U2	UI	John	PI	1.4.19	StarWars	GI	TREASURER	Bikers
	UI	John	PI	1.4.19	StarWars	G5	CHIEF	MovieClub
U3	UI	John	P4	1.5.19	VanGogh	G2	ADMIN	Geeks
	UI	John	P4	1.5.19	VanGogh	G4	SUPPORTER	Travel
	UI	John	P4	1.5.19	VanGogh	GI	TREASURER	Bikers
	UI	John	P4	1.5.19	VanGogh	G5	CHIEF	MovieClub
	UI	John	P5	1.6.19	Java	G2	ADMIN	Geeks
	UI	John	P5	1.6.19	Java	G4	SUPPORTER	Travel
	UI	John	P5	1.6.19	Java	GI	TREASURER	Bikers
	UI	John	P5	1.6.19	Java	G5	CHIEF	MovieClub
	U2	Mary	Р3	1.2.19	Liverpool	GI	OWNER	Bikers
	U3	Sarah	PI	1.3.19	StarWars	G3	FAN	Chefs



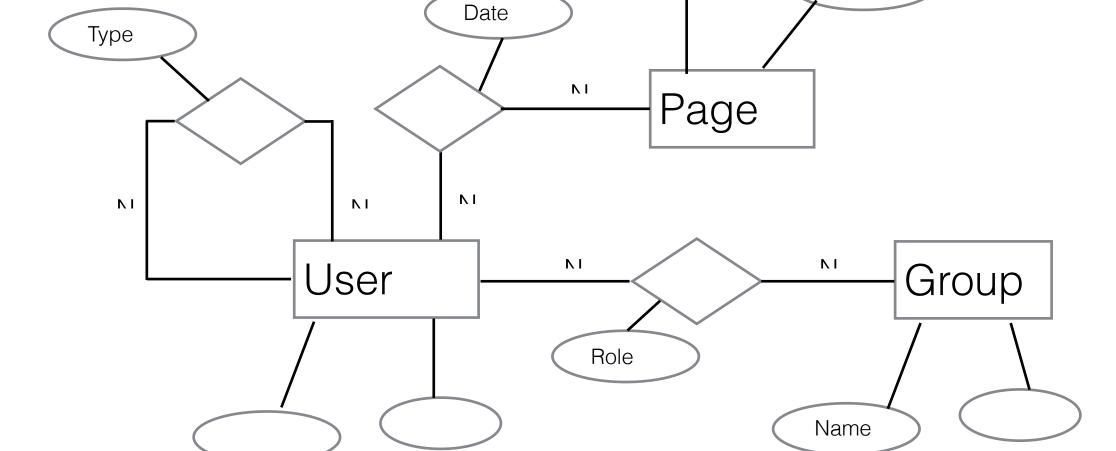


We are stuck then?

- NO, there is a simple trick that can fix any database
- Always add a "bridge relation" to your set of tables
- This is a relation that contains the "superkey" of the whole database, i.e. the key of the Universal Relation
- This will capture the individuals that "tick all boxes"







Relations:

- R1={<u>UserID</u>, Username}
- R2={GroupID, GroupName}
- R3={PageID, PageName}
- R4={<u>UserID,PageID,Date</u>} (user likes page)
- R5={<u>UserID,GroupID</u>,Role} (user belongs to group)
- R6={<u>UserID,FriendID</u>,Type} (user friends user)
- R7={<u>UserID,GroupID,PageID,FriendID</u>} superkey



Updated Tableaux test

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I				Value3		Value7			
R5	Value I		Value3					Value8		
R6	Value I								Value9	Value I 0
R7	Value I		Value3		Value3				Value9	



Updated Tableaux test

The bridge R7 will always guarantee the test is passed

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	Value I	Value2								
R2			Value3	Value4						
R3					Value5	Value6				
R4	Value I	Value2			Value5	Value6	Value7			
R5	Value I	Value2	Value3	Value4				Value8		
R6	Value I	Value2							Value9	Value I 0
R7	Value I	Value2	Value3	Value4	Value5	Value6	Value7	Value8	Value9	Value I 0



Updated counterexample

	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	FriendID	Туре
RI	UI	John	G2	Geeks	P2	Pizza	1.1.19	Admin	••	••
R2	U2	Mary	GI	Bikers	P3	Liverpool	1.2.19	Owner	••	••
R3	U3	Sarah	G3	Chefs	ΡI	StarWars	1.3.19	Fan	••	••
R4	UI	John	G4	Travel	PΙ	StarWars	1.4.19	Support er	••	••
R5	UI	John	GI	Bikers	P4	Van Gogh	1.5.19	Treasur er	••	••
R6	UI	John	G5	MovieClub	P5	Java!	1.6.19	Chief		
R7	UI	John	GI	Bikers	PI	StarWars	1.4.19	Treasur er		



	UserID	Username	GroupID	GroupName	PageID	PageName	Date	Role	Frie	Туре
RI	UI	John	G2	Geeks	P2	Pizza	1.1.19	Admin	• •	••
R2	U2	Mary	GI	Bikers	P3	Liverpool	1.2.19	Owner	• •	••
R3	U3	Sarah	G3	Chefs	PI	StarWars	1.3.19	Fan	• •	••
R4	UI	John	G4	Travel	PI	StarWars	1.4.19	Supporter	• •	••
R5	UI	John	GI	Bikers	P4	Van Gogh	1.5.19	Treasurer	••	••
R6	UI	John	G5	MovieClub	P5	Java!	1.6.19	Chief	• • •	• • •
R7	UI	lohn	GI	Bikers	PI	StarWars	1.4.19	Treasurer		

UserID	Username
UI	John
U2	Mary
U3	Sarah

GroupID	GroupName
GI	Bikers
G2	Geeks
G3	Chefs
G4	Travel
G5	MovieClub

PageID	PageName
PI	StarWars
P2	Pizza
P3	Liverpool
P4	Van Gogh
P5	Java!

UserID	GroupID	Role
UI	G2	Admin
U2	GI	Owner
U3	G3	Fan
UI	G4	Supporter
UI	GI	Treasurer
UI	G5	Chief

UserID	PageID	Date
UI	P2	1.1.19
U2	P3	1.2.19
U3	PI	1.3.19
UI	PI	1.4.19
UI	P4	1.5.19
UI	P5	1.6.19

These are the users who both have a page and belong to a group, in all legal combinations

UserID	GroupID	PageID
UI	G2	P2
U2	GI	P3
U3	G3	PI
UI	G4	PI
UI	GI	P4
UI	G5	P5
UI	GI	PI



UserID	Username	PageID	Date	PageName	GroupID	Role	GroupName
UI	John	P2	1.1.19	Pizza	G2	ADMIN	Geeks
UI	John	PI	1.4.19	StarWars	G4	SUPPORTER	Travel
UI	John	PI	1.4.19	StarWars	GI	TREASURER	Bikers
UI	John	P4	1.5.19	VanGogh	GI	TREASURER	Bikers
UI	John	P5	1.6.19	Java	G5	CHIEF	MovieClub
U2	Mary	P3	1.2.19	Liverpool	GI	OWNER	Bikers
U3	Sarah	PI	1.3.19	StarWars	G3	FAN	Chefs
	UI UI UI UI UI	UI John UI John UI John UI John UI John UI Mary	UI John P2 UI John PI UI John PI UI John P4 UI John P5 U2 Mary P3	UI John P2 1.1.19 UI John PI 1.4.19 UI John PI 1.4.19 UI John P4 1.5.19 UI John P5 1.6.19 U2 Mary P3 1.2.19	UI John P2 1.1.19 Pizza UI John PI 1.4.19 StarWars UI John PI 1.4.19 StarWars UI John P4 1.5.19 VanGogh UI John P5 1.6.19 Java U2 Mary P3 1.2.19 Liverpool	UI John P2 I.I.19 Pizza G2 UI John PI I.4.19 StarWars G4 UI John PI I.4.19 StarWars GI UI John P4 I.5.19 VanGogh GI UI John P5 I.6.19 Java G5 U2 Mary P3 I.2.19 Liverpool GI	UI John P2 I.I.19 Pizza G2 ADMIN UI John P1 I.4.19 StarWars G4 SUPPORTER UI John P1 I.4.19 StarWars G1 TREASURER UI John P4 I.5.19 VanGogh GI TREASURER UI John P5 I.6.19 Java G5 CHIEF U2 Mary P3 I.2.19 Liverpool GI OWNER



Conclusions

- We have gone full cycle (design wise at least)
 - from requirements to user stories
 - from user stories to conceptual design, in (E)ER
 - from conceptual design to logical design (in the relational model)
- we can identify issues and critical points in each of the phases, and we can express all of the above in a formal report
- Ready to go back to a more general notion of quality in CS/IT

