

# KING FAHD UNIVERSITY OF PETROLEUM & MINERALS

## COLLEGE OF COMPUTER SCIENCES & ENGINEERING

### COE 301 Computer Organization

### Project -Term 232

#### Pipelined Processor Design

##### Project Objectives:

- Designing a Pipelined 16-bit MIPS-like Processor
- Using Logisim Simulator to model and test the processor
- Teamwork Practice

##### Instruction Set Architecture

In this project, students should design a simple 16-bit MIPS-like processor with eight 16-bit general purpose registers: R0 through R7. R0 is hardwired to zero and cannot be written to. There is also one special-purpose 12-bit register, which is the program counter (PC). All instructions are 16 bits, and there are three instruction formats: R-type, I-type, and J-type as shown below:

##### R-type format

4-bit opcode (Op), 3-bit register numbers (Rs, Rt, and Rd), and 3-bit function field (funct)

Op <sup>4</sup>	Rd <sup>3</sup>	Rs <sup>3</sup>	Rt <sup>3</sup>	funct <sup>3</sup>
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##### I-type format

4-bit opcode (Op), 3-bit register numbers (Rs and Rt), and 6-bit immediate constant

Op <sup>4</sup>	Rd <sup>3</sup>	Rs <sup>3</sup>	Immediate <sup>6</sup>
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##### J-type format

4-bit opcode (Op) and 12-bit immediate constant

Op <sup>4</sup>	Immediate <sup>12</sup>
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For **R-type** instructions, Rs and Rt specify the two source register numbers, and Rd specifies the destination register number. The function field can specify at most eight functions for a given opcode. Opcodes 0000 and 0001 are reserved for R-type instructions.

For **I-type** instructions, Rs specifies a source register number, and Rt can be a second source or a destination register number. The immediate constant is only 6 bits because of the fixed-size nature of the instruction. The 6-bit immediate constant is assumed to be sign-extended for arithmetic instructions and zero extended for logical instructions.

For **J-type**, a 12-bit immediate constant is used for J (jump), JAL (jump-and-link), and LUI (load upper immediate) instructions.

## Instruction Encoding

Sixteen R-type instructions, eleven I-type instructions, and three J-type instructions are defined. These instructions, their meaning, and their encoding are shown below:

Instr.	Meaning		Encoding			
OR	$\text{Reg(Rd)} = \text{Reg(Rs)} \mid \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 000
NAND	$\text{Reg(Rd)} = \sim(\text{Reg(Rs)} \& \text{Reg(Rt)})$	Op = 0000	Rd	Rs	Rt	f = 001
XOR	$\text{Reg(Rd)} = \text{Reg(Rs)} \wedge \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 010
ADD	$\text{Reg(Rd)} = \text{Reg(Rs)} + \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 100
SUB	$\text{Reg(Rd)} = \text{Reg(Rs)} - \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 101
SGE	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ signed } \geq \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 110
SGEU	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ unsigned } \geq \text{Reg(Rt)}$	Op = 0000	Rd	Rs	Rt	f = 111
SLL	$\text{Reg(Rd)} = \text{Reg(Rs)} \ll \text{Reg(Rt)}$	Op = 0001	Rd	Rs	Rt	f = 000
SRL	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ zero } \gg \text{Reg(Rt)}$	Op = 0001	Rd	Rs	Rt	f = 001
SRA	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ sign } \gg \text{Reg(Rt)}$	Op = 0001	Rd	Rs	Rt	f = 010
ROR	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ ROR by Reg(Rt)}$	Op = 0001	Rd	Rs	Rt	f = 100
MOVZ	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ if } (\text{Reg(Rt)} == 0)$	Op = 0001	Rd	Rs	Rt	f = 101
MOVN	$\text{Reg(Rd)} = \text{Reg(Rs)} \text{ if } (\text{Reg(Rt)} != 0)$	Op = 0001	Rd	Rs	Rt	f = 111
JR	PC = lower 12 bits of Reg(Rs)	Op = 0001	000	Rs	000	f = 110
JALR	Rd=PC+1, PC=Rs	Op = 0001	Rd	Rs	000	f = 111
ANDI	$\text{Reg(Rt)} = \text{Reg(Rs)} \& \text{zero\_ext(im}_6)$	Op = 0010	Rd	Rs	Immediate <sup>6</sup>	
ORI	$\text{Reg(Rt)} = \text{Reg(Rs)} \mid \text{zero\_ext(im}_6)$	Op = 0011	Rd	Rs	Immediate <sup>6</sup>	
ADDI	$\text{Reg(Rt)} = \text{Reg(Rs)} + \text{sign\_ext(im}_6)$	Op = 0100	Rd	Rs	Immediate <sup>6</sup>	
SW	$\text{Mem}(\text{Reg(Rs)} + \text{sign\_ext(im}_6)) = \text{Reg(Rt)}$	Op = 0101	Rd	Rs	Immediate <sup>6</sup>	
LW	$\text{Reg(Rt)} = \text{Mem}(\text{Reg(Rs)} + \text{sign\_ext(im}_6))$	Op = 0110	Rd	Rs	Immediate <sup>6</sup>	
BNE	Branch if $(\text{Reg(Rs)} \neq \text{Reg(Rt)})$	Op = 0111	Rd	Rs	Immediate <sup>6</sup>	
BEQ	Branch if $(\text{Reg(Rs)} == \text{Reg(Rt)})$	Op = 1000	Rd	Rs	Immediate <sup>6</sup>	
BGT	Branch if $(\text{Reg(Rs)} \text{ signed } > \text{Reg(Rt)})$	Op = 1001	Rd	Rs	Immediate <sup>6</sup>	
BGE	Branch if $(\text{Reg(Rs)} \text{ signed } \geq \text{Reg(Rt)})$	Op = 1010	Rd	Rs	Immediate <sup>6</sup>	
BLT	Branch if $(\text{Reg(Rs)} \text{ signed } < \text{Reg(Rt)})$	Op = 1011	Rd	Rs	Immediate <sup>6</sup>	
BLE	Branch if $(\text{Reg(Rs)} \text{ signed } \geq \text{Reg(Rt)})$	Op = 1100	Rd	Rs	Immediate <sup>6</sup>	
LUI	$\text{R1} = \text{Immediate}^{12} \ll 4$	Op = 1101	Immediate <sup>12</sup>			
J	$\text{PC} = \text{Immediate}^{12}$	Op = 1110	Immediate <sup>12</sup>			
JAL	$\text{R7} = \text{PC} + 1, \text{PC} = \text{Immediate}^{12}$	Op = 1111	Immediate <sup>12</sup>			

For the four shift and rotate instructions (SLL, SRL, SRA, ROR), the least significant 4 bits of register  $R_t$  specify the shift or rotate amount. The Load Upper Immediate (LUI) is a J-type instruction, it loads the 12-bit immediate constant into the upper 12 bits of register  $R_1$ . The LUI can be combined with ORI (or ADDI) to load any 16-bit constant into register  $R_1$ . The JAL, JALR and JR instructions are used for procedure calls and returns. Although the instruction set is reduced, it is still rich enough to write useful programs.

## Memory

Your processor will have separate instruction and data memories with  $2^{12} = 4096$  words each. Each word is 16 bits or 2 bytes. Memory is *word addressable*. Only words (not bytes) can be read from or written to memory, and each address is a word address. This will simplify the processor implementation. The PC contains a word address (not a byte address). Therefore, it is sufficient to increment the PC by 1 (rather than 2) to point to the next instruction in memory. Also, the Load and Store instructions can only load and store words. There is no instruction to load or store a byte in memory.

## Addressing Modes

For branch instructions (BEQ, BNE, BLT, BLE, BGT and BGE), PC-relative addressing mode is used:  $PC = PC + \text{sign-extend}(\text{immediate})^6$ . For jump instructions (J and JAL), absolute addressing is used:  $PC = \text{Immediate}^{12}$ . For LW and SW instructions, base-displacement addressing mode is used. The base address in register  $R_s$  contains the base address and the memory address is computed by adding the base address with the offset.

## Program Execution

The program once loaded in the code memory will start at address 0 in the instruction memory. The data segment will be loaded and will start also at address 0 in the data memory. You may also have a stack segment if you want to support procedures. The stack segment can occupy the upper part of the data memory and can grow backwards towards lower addresses. The stack segment can be implemented completely in software. To terminate the execution of a program, the last instruction in the program can jump or branch to itself indefinitely.

## Build a Single-Cycle Processor

Start by building the datapath and control of a single-cycle processor and ensure its correctness. You should have sufficient test cases that ensure the correct execution of ALL instructions in the instruction set. You should also write test cases that show the correct execution of complete programs. To verify the correctness of your design, show the values of all registers ( $R_1$  to  $R_7$ ) at the top-level of your design. Provide output pins for registers  $R_1$  through  $R_7$  and make their values visible at the top level of your design to simplify testing and verification.

## Build a Pipelined Processor

Once you have succeeded in building a single-cycle processor and verified its correctness, design and implement a pipelined version of your design. Make a copy of your single-cycle design, then convert it and implement a pipelined datapath and its control logic. Add pipeline registers between stages. Design the control logic to detect data dependencies among instructions and implement the forwarding logic. You should handle properly the data hazards when there are data dependencies between instructions and stall the pipeline when needed. You should handle properly the control hazards of the branch and jump instructions.

## Design Alternatives

When designing the datapath and control unit, explore alternative design options and justify why a given design alternative is chosen. For example, when designing the control unit consider implementing it using a decoder and a set of OR/NOR gates vs. using a ROM to store the control signals vs. optimizing the equation of each control signal separately. When designing the ALU and the shifter unit, consider alternative designs and justify why a design alternative is chosen. The same should be applied for all design decisions in your CPU.

## Testing

To test the correct functionality of your designed CPU, you need to implement the following programs:

1. A program to count the number of ones in a register,
2. The **bubble sort procedure** given in class. Use this procedure to sort an array of 8 words of your choice,
3. A program that tests each of the remaining untested instructions to demonstrate their correct operation.

## WARNING

Although Logisim is stable, it might crash from time to time. Therefore, it is best to save your work often. Make several copies and versions of your design before making changes, in case you need to go back to an older version.

## Project Report

The report document must contain sections highlighting the following:

### Design and Implementation

- Specify clearly the design giving detailed description of the datapath, its components, control, and the implementation details (highlighting the design choices you made and why, and any notable features that your processor might have). Document clearly design alternatives explored and why a given design is selected.
- Provide drawings of the component circuits and the overall datapath.
- Provide a complete description of the control logic and the control signals. Provide a table giving the control signal values for each instruction. Provide the logic equations for each control signal.
- Provide a complete description of the forwarding logic, the cases that were handled, and the cases that stall the pipeline, and the logic that you have implemented to stall the pipeline.
- Provide list of sources for any parts of your design that are not entirely yours (if any).
- Carry out the design and implementation with the following aspects in mind: -
  - Correctness of the individual components
  - Correctness of the overall design when wiring the components together
  - Completeness: all instructions were implemented properly, detecting dependences and forwarding was handled properly, and stalling the pipeline was handled properly for all cases.

## 2 – Simulation and Testing

- Carry out the simulation of the processor developed using Logisim.
- Test each of the components individually and demonstrate its correct operation including the ALU and register file.
- Describe all the features of the simulator used for simulating your design with a clear emphasis on its advantages and limitations (if any) for simulating the design, list the known bugs or missing features (if any).
- Describe the test programs that you used to test your design with enough comments describing the program, its inputs, and its expected output. List all the instructions that were tested and work correctly. List all the instructions that do not run properly.
- Describe all the cases that you handled involving dependencies between instructions, forwarding cases, and cases that stall the pipeline.
- Also provide snapshots of the Simulator window with your test program loaded and showing the simulation output results.

## 3 – Teamwork

- This project is a teamwork project with a maximum of three students. Make sure to write the names of all the group members on the project report title page.
- Each group should assign a group leader that leads the conduction of the project, divide the project tasks among the team members.
- Project tasks should be divided among the group members so that each group member contributes equally in the project and everyone is involved in all the following activities:
  - Design and Implementation
  - Simulation and Testing
  - Design and results reporting
- Clearly show the work done by each group member using a chart and prepare an execution plan showing the time frame for completing the subtasks of the project. You can also mention how many meetings were conducted between the group members to discuss the design, implementation, and testing.
- Students who help other team members should mention that to earn credit for that.

## Submission Guidelines

All submissions will be done through Blackboard.

Attach one zip file containing all the design circuits, the test programs source code and binary instruction files that you have used to test your design, their test data, as well as the report document.

Submit also a hard copy of the report to the lab instructor.

## Grading policy:

The grade will be divided according to the following components:

- Correctness: whether your implementation is working.
- Completeness and testing: whether all instructions and cases have been implemented, handled, and tested properly.
- Participation and contribution to the project. It should be noted that being in a group that implemented the project correctly does not qualify you to get full mark. Your mark will depend on your contribution in the project.

- Report organization and clarity.

**Submission Deadlines:**

<b>Task</b>	<b>Deadline</b>
Single-cycle CPU Design Demo	Week 13
Pipelined CPU Design Demo	Week 15
Final Report Submission	Week 15