

Lax comma categories: descent and exponentiability

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109th Peripatetic Seminar on Sheaves and Logic

17 November 2024

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- objects: morphisms $f: B \rightarrow X$ of \mathbb{A} with B in \mathbb{B} ,
- morphisms $f \rightarrow g$: 2-cells θ of \mathbb{A} of the form

$$\begin{array}{ccc} A & \xrightarrow{h} & B \\ & \searrow f \quad \xRightarrow{\theta} \quad \swarrow g & \\ & X & \end{array}$$

where h is a morphism in \mathbb{B} .

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Thus, $\mathbf{Set} \Downarrow \mathcal{X} = \mathbf{Fam}(\mathcal{X})$.

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that is, $\alpha(a) \leq \beta(f(a))$ for all $a \in A$.

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that is, we have a fibration $\mathbb{A} \Downarrow X \rightarrow \mathbb{A}$.

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Let \mathcal{X} be a infinitary distributive category. The following are equivalent:

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- $\mathbf{Fam}(\mathcal{X})$ is cartesian closed.

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If \mathcal{X} is cartesian closed, then so is $\mathbf{Cat} \Downarrow \mathcal{X}$.

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We say p is an *effective descent morphism* (descent morphism) if p^* is monadic (premonadic).

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Theorem (Lucatelli Nunes, P. 2024)

If X satisfies mild conditions, then

$$\mathbb{A} \Downarrow X \rightarrow \mathbb{A}$$

preserves effective descent morphisms.

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- $f: (\alpha(a))_{a \leq a'} \rightarrow (\beta(b))_{b \leq b'}$ is a descent morphism in $\mathbf{Fam}(X)$.

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- Characterization of effective descent morphisms in $\mathbf{Cat} \downarrow X$.

Dank wel!