

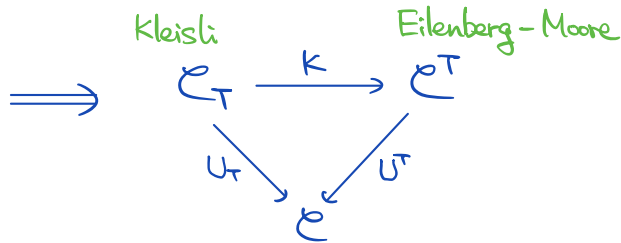
# Monad theory through enrichment

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(Joint work with John Bourke)

$\mathcal{C}$ : category

$T: \mathcal{C} \rightarrow \mathcal{C}$ : monad



$\mathcal{C}_T \cong$  full subcat. of  $\mathcal{C}^T$  consisting of the free  $T$ -algebras  $\left(TA, \begin{smallmatrix} T^2A \\ \downarrow \mu_A \\ TA \end{smallmatrix}\right)$ .

Every  $T$ -algebra  $\left(A, \begin{smallmatrix} TA \\ \downarrow \alpha \\ A \end{smallmatrix}\right)$  is the canonical coequalizer of free  $T$ -algebras:

$$\left(T^2A, \begin{smallmatrix} T^3A \\ \downarrow \mu_{TA} \\ T^2A \end{smallmatrix}\right) \begin{array}{c} \xrightarrow{\mu_A} \\ \xrightarrow{T\alpha} \end{array} \left(TA, \begin{smallmatrix} T^2A \\ \downarrow \mu_A \\ TA \end{smallmatrix}\right) \xrightarrow{\alpha} \left(A, \begin{smallmatrix} TA \\ \downarrow \alpha \\ A \end{smallmatrix}\right) : \text{coeq. in } \mathcal{C}^T.$$

**Q1.** Is  $\mathcal{C}^T$  a "(co)completion" of  $\mathcal{C}_T$  in a suitable sense?

**A1.** When e.g.  $\mathcal{C} = \mathbf{Set}$ ,  $\mathcal{C}^T$  is the **exact completion** of the weakly left exact category  $\mathcal{C}_T$ . [Vitale, Left covering functors, 1994]

What about general  $\mathcal{C}$ ?

In general,  $\mathcal{C}^T$  might not be exact, might not have (reflexive) coequalizers, ... . In fact,  $\mathcal{C}^T$  could be any category!

However,  $\mathcal{E}^T$  has certain (co)completeness **relative to**  $\mathcal{C}$ .

- $\mathcal{E}^T$  **lifts** limits in  $\mathcal{C}$ :

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{X} & \mathcal{E}^T \\ & & \downarrow U^T \\ & & \mathcal{C} \end{array} \quad \begin{array}{l} \text{If } \lim U^T X \text{ exists in } \mathcal{C}, \text{ then} \\ \lim X \text{ exists and } U^T(\lim X) \cong \lim U^T X. \\ (U^T \text{ creates limits.}) \end{array}$$

- $\mathcal{E}^T$  lifts colimits in  $\mathcal{C}$  preserved by  $T$ .
- In particular,  $\mathcal{E}^T$  lifts **absolute colimits** in  $\mathcal{C}$ .

The coequalizer

$$\left( T^2 A, \begin{array}{c} T^3 A \\ \downarrow \mu_{TA} \\ T^2 A \end{array} \right) \begin{array}{c} \xrightarrow{\mu_A} \\ \xrightarrow{T\alpha} \end{array} \left( T A, \begin{array}{c} T^2 A \\ \downarrow \mu_A \\ T A \end{array} \right) \xrightarrow{\alpha} \left( A, \begin{array}{c} T A \\ \downarrow \alpha \\ A \end{array} \right)$$

is a lifting of an absolute (in fact, split) coequalizer in  $\mathcal{C}$ :

$$\begin{array}{ccccc} T^2 A & \begin{array}{c} \xrightarrow{\mu_A} \\ \xrightarrow{T\alpha} \end{array} & T A & \xrightarrow{\alpha} & A \\ & \searrow \eta_{TA} & \nwarrow \eta_A & & \end{array}$$

$\Rightarrow$  Instead of the **ordinary category theory**, let us work in the **category theory over**  $\mathcal{C}$ .

Eilenberg-Moore

Kleisli

Q2. Is  $\begin{array}{c} \mathcal{E}^T \\ \downarrow U^T \\ \mathcal{C} \end{array}$  a (co)completion of  $\begin{array}{c} \mathcal{C}_T \\ \downarrow U_T \\ \mathcal{C} \end{array}$  in a suitable sense?

... but what exactly is "category theory over  $\mathcal{C}$ "?

We take the following view:

(weighted) (co)limits, free (co)completion, ...

category theory over  $\mathcal{C} = \underline{\text{enriched category theory}}$   
over a bicategory  $\Sigma \text{Set} \downarrow \mathcal{C}$ .

[Fujii and Lack, The oplax limit of an enriched category, 2024]

$$\text{Cat}(\mathcal{C}) \cong (\Sigma \text{Set} \downarrow \mathcal{C})\text{-Cat.}$$

$\uparrow$                        $\uparrow$   
 2-categories

Bicategory-enriched category theory [Betti, Carboni, Street, Walters, ...]

$\mathcal{B}$ : bicategory

A  $\mathcal{B}$ -category  $A$  consists of:

- a set  $\text{ob}(A)$  of objects equipped with a function

$$\begin{array}{ccc} \text{ob}(A) & \xrightarrow{1-1} & \text{ob}(B) \\ \downarrow & & \downarrow \\ x & \longmapsto & |x| : \text{the extent of } x. \end{array}$$

- $\forall x, y \in \text{ob}(A)$ . a 1-cell  $|x| \xrightarrow{A(x,y)} |y|$  in  $\mathcal{B}$ .  
 $\uparrow$  the hom 1-cell from  $x$  to  $y$ .

- $\forall x \in \text{ob}(A)$ . a 2-cell  $|x| \begin{array}{c} \xrightarrow{I_x} \\ \Downarrow j_x \\ \xrightarrow{A(x,x)} \end{array} |x|$  in  $\mathcal{B}$ .

- $\forall x, y, z \in \text{ob}(A)$ . a 2-cell  $|x| \begin{array}{ccc} \xrightarrow{A(x,y)} & |y| & \xrightarrow{A(y,z)} \\ & \Downarrow m_{xyz} & \\ |x| & & |z| \end{array} \xrightarrow{A(x,z)} |z|$  in  $\mathcal{B}$ .

satisfying the unit and associativity axioms.

For any bicategory  $B$  and any  $B$ -category  $A$ ,  
there exists a bicategory  $B \downarrow A$  with

$$(B\text{-Cat})/A \cong (B \downarrow A)\text{-Cat}.$$

[Fujii and Lack, The oplax limit of an enriched category, 2024].

Monoidal cat.  $(\text{Set}, 1, *) = 1\text{-obj. bicategory } \Sigma\text{Set}.$



In particular, for any  $(\text{Set-})$  category  $\mathcal{C}$ ,  $\exists$  bicategory  $\Sigma\text{Set} \downarrow \mathcal{C}$

with

$$\text{Cat}/\mathcal{C} \cong (\Sigma\text{Set} \downarrow \mathcal{C})\text{-Cat}.$$

$\Sigma\text{Set} \downarrow \mathcal{C}$  : bicategory s.t.

-  $\text{ob}(\Sigma\text{Set} \downarrow \mathcal{C}) = \text{ob}(\mathcal{C})$  ←  $\Sigma\text{Set} \downarrow \mathcal{C}$  : monoidal cat only when  $\mathcal{C} = \text{monoid } M$ .

( $\Rightarrow \Sigma\text{Set} \downarrow \mathcal{C} = \text{Set}/M$  with the canonical monoidal str.)

-  $\forall A, B \in \text{ob}(\mathcal{C}) = \text{ob}(\Sigma\text{Set} \downarrow \mathcal{C}).$

$$(\Sigma\text{Set} \downarrow \mathcal{C})(A, B) = \text{Set}/\mathcal{C}(A, B).$$

So a  $(\Sigma\text{Set} \downarrow \mathcal{C})$ -category  $\mathbb{E}$  consists of

- a set  $\text{ob}(\mathbb{E})$  t.w. a function  $\text{ob}(\mathbb{E}) \xrightarrow{1-1} \text{ob}(\Sigma\text{Set} \downarrow \mathcal{C}) = \text{ob}(\mathcal{C})$

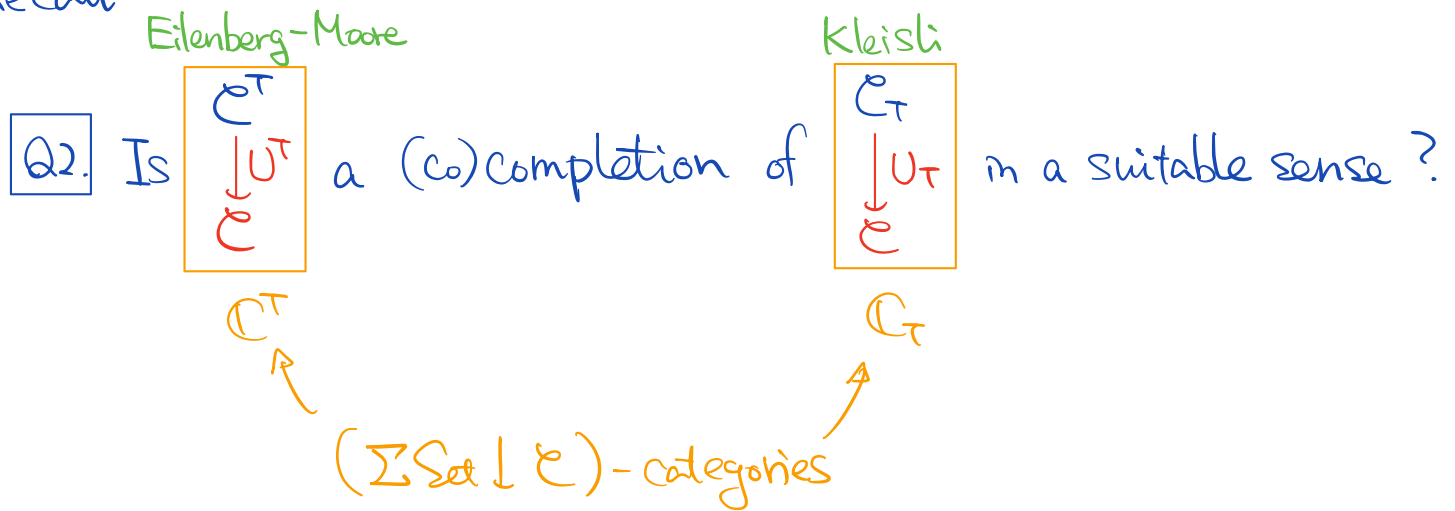
$$\begin{array}{ccc} \downarrow & & \downarrow \\ \chi & \xrightarrow{\quad\quad\quad} & \cup \chi \end{array}$$

-  $\forall x, y \in \text{ob}(\mathbb{E}),$  an object  $\mathbb{E}(x, y) = \begin{pmatrix} \mathbb{E}(x, y) \\ \downarrow \cup x, y \\ \mathcal{C}(\cup x, \cup y) \end{pmatrix} \in \text{Set}/\mathcal{C}(\cup x, \cup y)$

⋮

$\Rightarrow$  A category  $\mathbb{E}$  equipped with a functor  $\begin{array}{c} \mathbb{E} \\ \downarrow \cup \\ \mathcal{C} \end{array}$ .

Recall



A2. [Bourke and F.]  $\mathcal{C}^T$  is the free completion of  $\mathcal{C}_T$  under lifts of absolute colimits in  $\mathcal{C}$ .

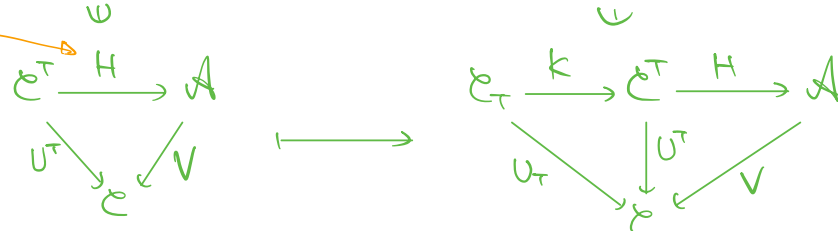
=  $\Phi$ -colimits for a suitable class  $\Phi$  of weights in  $(\Sigma\text{Set} \downarrow \mathcal{C})$ -enriched category theory.

In concrete terms, this means:

- $\begin{array}{c} \mathcal{C}^T \\ \downarrow U^T \\ \mathcal{C} \end{array}$  lifts absolute colimits in  $\mathcal{C}$ 
  - $\forall \mathcal{D} \xrightarrow{W} \text{Set}. \quad \forall \mathcal{D} \xrightarrow{X} \mathcal{C}^T.$
  - $\forall$  weighted colimit  $W * U_T X$  in  $\mathcal{C}$  which is an absolute colimit, a weighted colimit  $W * X$  exists in  $\mathcal{C}^T$  and satisfies  $U^T(W * X) = W * U_T X$ .
- $\forall \begin{array}{c} A \\ \downarrow V \\ \mathcal{C} \end{array}$  which lifts absolute colimits in  $\mathcal{C}$ ,

$$\Phi\text{-Cocts} \left( \begin{array}{c} \mathcal{C}^T \\ \downarrow U^T \\ \mathcal{C} \end{array}, \begin{array}{c} A \\ \downarrow V \\ \mathcal{C} \end{array} \right) \simeq \text{Cat}/\mathcal{C} \left( \begin{array}{c} \mathcal{C}_T \\ \downarrow U_T \\ \mathcal{C} \end{array}, \begin{array}{c} A \\ \downarrow V \\ \mathcal{C} \end{array} \right).$$

preserving lifts of absolute colimits in  $\mathcal{C}$



## Cor. (Monadicity theorem)

$\mathcal{A}$

$\downarrow V : \text{functor}$

$\mathcal{C}$

Suppose

- $V$  has a left adjoint
- $(A, V) : \text{lifts absolute colimits in } \mathcal{C}$
- $V$  is conservative.  $\Phi\text{-cocomplete}$

Then  $\mathcal{A} \simeq \mathcal{C}^T$  where  $T$  is the monad on  $\mathcal{C}$  induced by  $V$  and its left adjoint.

### Proof

One can show that under these assumptions

$$\begin{array}{ccc} \mathcal{C}_T & \xrightarrow{J} & \mathcal{A} \\ & \searrow U_T & \swarrow V \\ & \mathcal{C} & \end{array}$$

is a free completion under  $\Phi$ -colimits.

Hence the claim follows from the uniqueness up to equivalence of free completions. □

We can also describe the free completion under  $\Phi$ -colimits.

Given  $\begin{array}{c} \mathcal{A} \\ \downarrow V \\ \mathcal{C} \end{array}$ , its free completion  $\begin{array}{c} \Phi(\mathcal{A}) \\ \downarrow \Phi(V) \\ \mathcal{C} \end{array}$  under  $\Phi$ -colimits is

given as follows. [Betti, Cocompleteness over coverings, 1985]

obj.  $(\mathcal{A}^{\text{op}} \xrightarrow{W} \text{Set}, C \in \mathcal{C}, W \xrightarrow{\alpha} \mathcal{C}(V-, C))$  s.t.  
 $\alpha$  exhibits  $C$  as the weighted colimit  $W * V$  in  $\mathcal{C}$ , and  
 moreover  $W * V$  is an absolute colimit.

"1-step completion" works since  $\Phi$  is a saturated class of weights.

Prop. [Bourke and F.]

$\begin{array}{c} \mathcal{A} \\ \downarrow V \\ \mathcal{C} \end{array}$

Suppose  $V$  has a left adjoint  $G$ . Then in

$$\begin{array}{ccccc} \mathcal{A}^{\text{op}} & \xrightarrow{W} & \text{Set} & & \mathcal{A}^{\text{op}} & \xrightarrow{W} & \text{Set} \\ V^{\text{op}} \downarrow & \Downarrow \alpha & \downarrow \text{id} & \xleftrightarrow{\text{mate}} & G^{\text{op}} \uparrow & \Downarrow \beta & \uparrow \text{id} \\ \mathcal{C}^{\text{op}} & \xrightarrow{\mathcal{C}(G-, C)} & \text{Set} & & \mathcal{C}^{\text{op}} & \xrightarrow{\mathcal{C}(G-, C)} & \text{Set} \end{array}$$

we have  $(W, C, \alpha) \in \Phi(\mathcal{A}) \iff \beta : \text{iso.}$

Cor. [Bourke and F.]

$$\begin{array}{c} \mathcal{A} \\ G \uparrow \downarrow V \\ \mathcal{C} \end{array} \Rightarrow \begin{array}{ccc} \Phi(\mathcal{A}) & \longrightarrow & [\mathcal{A}^{\text{op}}, \text{Set}] \\ \Phi(V) \downarrow & \Downarrow \cong & \downarrow [G^{\text{op}}, \text{Set}] \\ \mathcal{C} & \xrightarrow{\gamma} & [\mathcal{C}^{\text{op}}, \text{Set}] \end{array} \begin{array}{l} : \text{bipullback} \\ \text{(in fact, iso-comma obj.)} \end{array}$$

$\gamma$  Yoneda embedding

Cor. (Linton's theorem)

$$\begin{array}{c} \mathcal{C} : \text{category} \\ T : \text{monad on } \mathcal{C} \end{array} \Rightarrow \begin{array}{ccc} \mathcal{C}^T & \longrightarrow & [\mathcal{C}_T^{\text{op}}, \text{Set}] \\ U^T \downarrow & & \downarrow [F_T^{\text{op}}, \text{Set}] \\ \mathcal{C} & \xrightarrow{\gamma} & [\mathcal{C}^{\text{op}}, \text{Set}] \end{array} \begin{array}{l} : \text{bipullback} \\ \text{(in fact, pullback)} \end{array}$$