Digital Forensics
Reverse Engineering
Lecture 6

Deciphering Code Structures

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Reverse Engineering – Deciphering Code Structures

- Goal:
  - Learn about the most common logical and control flow constructs used in high-level languages
  - How they are implemented in assembly languages including loops, and conditional blocks
- It helps to know about the types of typical assembly language sequences usually observed in assembly language code

- An example of a conditional codes
  - "je" will cause a jumpt to "SomePlace" if "EAX" equals to 7

```
cmp
eax, 7
je
SomePlace
```

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- Control flow and program layout
  - Branches are the most popular tools for implementing logic of the program
  - Most common programming constructs
  - Identifying branches and understanding their meaning and purpose is important in reverse engineering

- Deciphering functions
  - Basic building block in a program
  - On I-32 bits processors function are called using the "CALL" instruction
    - It stores the current instruction pointer in the stack and jumps to the function address
    - It helps to distinguish it from unconditional jumps.

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- Internal functions
  - They are called from the same binary executable that contains their implementation
  - Compilers generate an internal function call sequence
  - They embed the function's address into the code (easy to detect)
  - E.g., "Call CodeSectionAddress"

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- Imported functions
  - When a module is making a call into a function implemented in another binary executable
  - During the compilation process the compiler has no idea where the imported function can be found
    - Unable to embed the function's address into the code
  - "Imported Address Table" and "Import Directory" are usually used to implement calling imported functions
    - The Import Directory is used in runtime for resolving the function's name with a matching function in the target executable
    - IAT stores the actual address of the target function
    - The caller then loads the function's pointer from the IAT and calls it
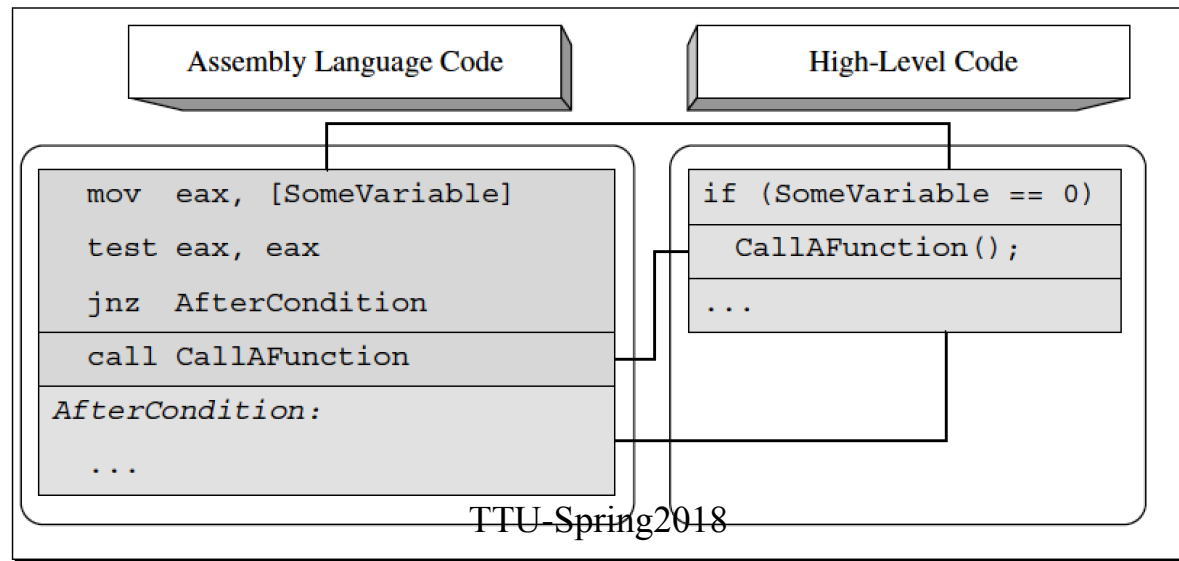
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- Imported functions
    - An example of a typical imported function call
        - "call   DWORD   PTR [IAT_Pointer]
        - "DWORD   PTR" tells the CPU to jump not to the address of IAT_Pointer but to the address that is pointed to by IAT_Pointer
    - Detecting imported calls is easy because of these types of calls

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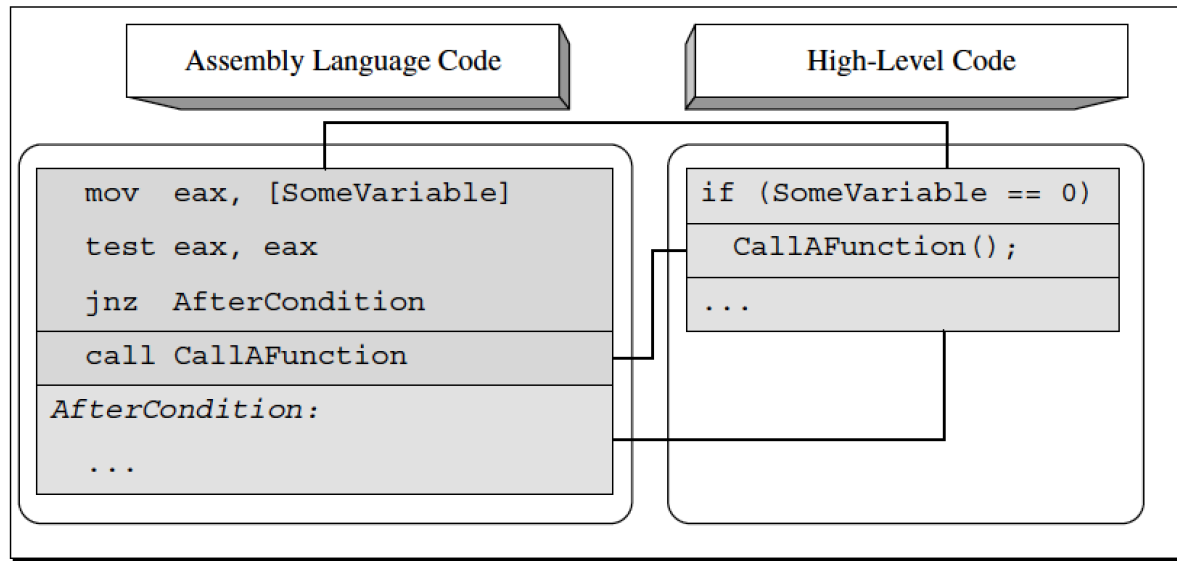
```
if (SomeVariable == 0)
        CallAFunction();
```

- Single-Branch conditionals

  - Requires a logical check to determine whether "SomeVariable" contains "0" or not

  - Skip the conditional block by performing a conditional jump if "SomeVariable" is nonzero

  - Use a simple TEST to perform a simple zero check for EAX.

    - A bitwise AND operation on EAX and setting flags to reflect the result

    - An effective way : TEST sets the zero flag (ZF) with respect to the result of the bitwise AND operation

Assembly Language Code	High-Level Code
`mov    eax, [SomeVariable]`	`if (SomeVariable == 0)`
`test eax, eax`	`    CallAFunction();`
`jnz    AfterCondition`	`...`
`call CallAFunction`	
`AfterCondition:`	
`  ...`	

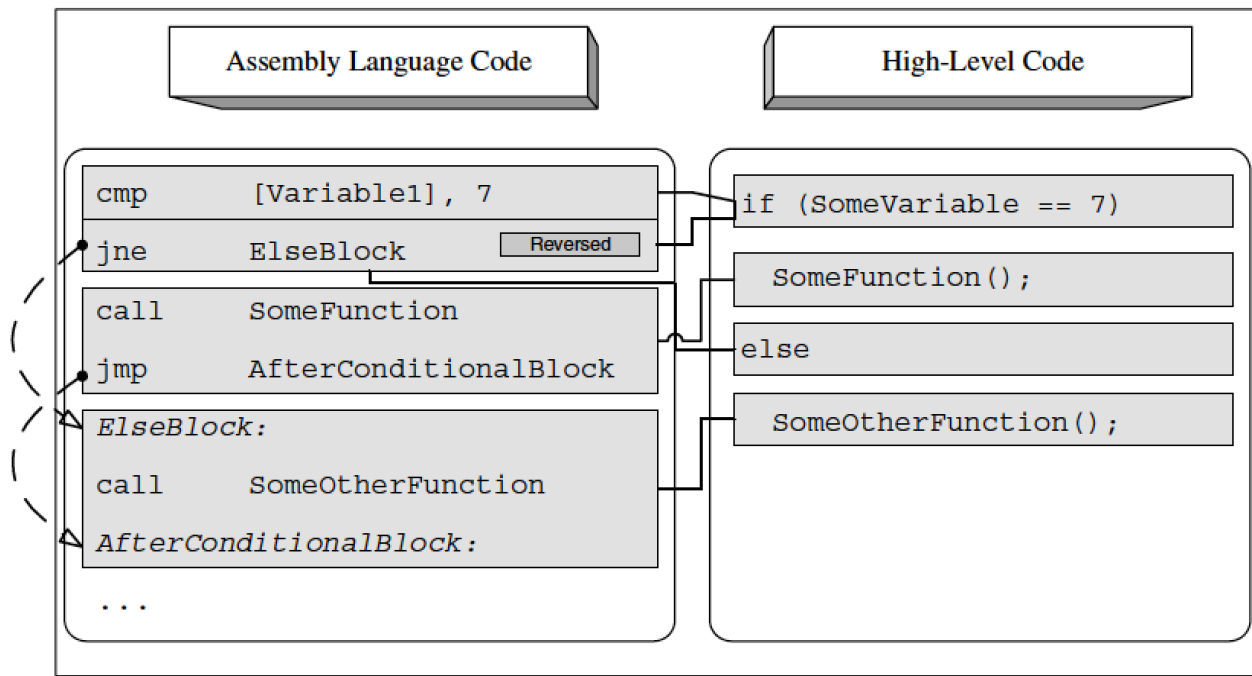# Reverse Engineering – Deciphering Code Structures

```
if (SomeVariable == 0)
    CallAFunction();
```

- **Single-Branch conditionals**
  - Note. The condition is reversed.
    - In the source code: the program checks whether "SomeVariable" equals to "zero"
    - The compiled version is reversing the condition: The conditional instruction checks whether "SomeVariable" is nonzero

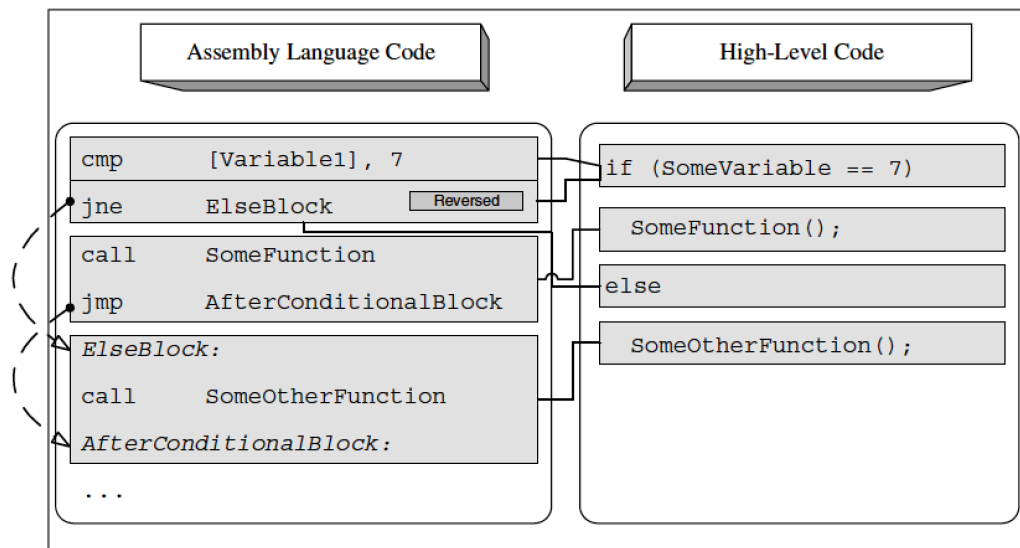| Assembly Language Code	High-Level Code
`mov  eax, [SomeVariable]`	`if (SomeVariable == 0)`
`test eax, eax`	`    CallAFunction();`
`jnz  AfterCondition`	`...`
`call CallAFunction`	
`AfterCondition:`	
`  ...`	

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- Two-way conditionals
 - Typically implemented in high-level languages using (if-else)

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- Two-way conditionals
  - How compiler implements two-way conditionals?
    - The conditional branch points to the "else" block and not to the code follows the conditional statement
    - The primary conditional block is placed right after the conditional jump
    - The conditional block always ends with an unconditional jump that essentially skips the "else" block
    - The "else" block is placed at the end of the conditional block, right after that conditional jump.
    - The unconditional "JMP" jumps after the functional call

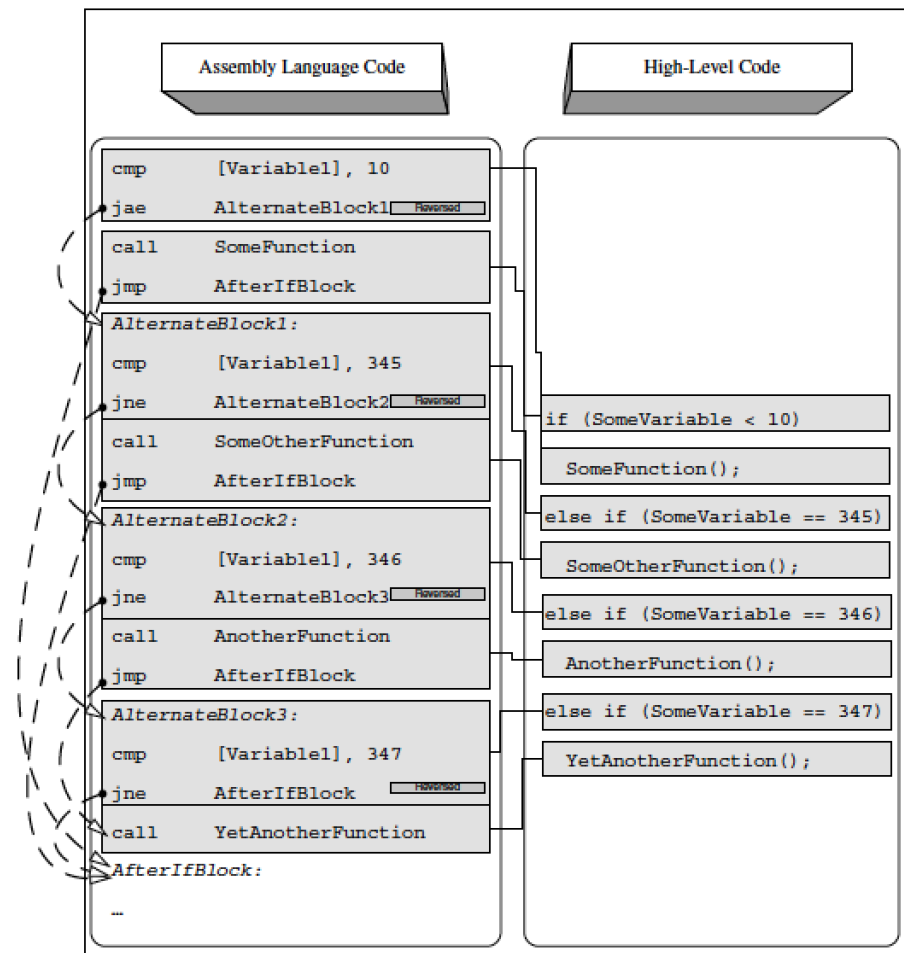| Assembly Language Code	High-Level Code
`cmp     [Variable1], 7`	`if (SomeVariable == 7)`
`jne     ElseBlock` *Reversed*	
`call    SomeFunction`	`    SomeFunction();`
`jmp     AfterConditionalBlock`	`else`
`ElseBlock:`	`    SomeOtherFunction();`
`call    SomeOtherFunction`	
`AfterConditionalBlock:`	
`...`	

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- Two-way conditionals – A more complex case
  - A separate conditional statement is used for each of the two code blocks
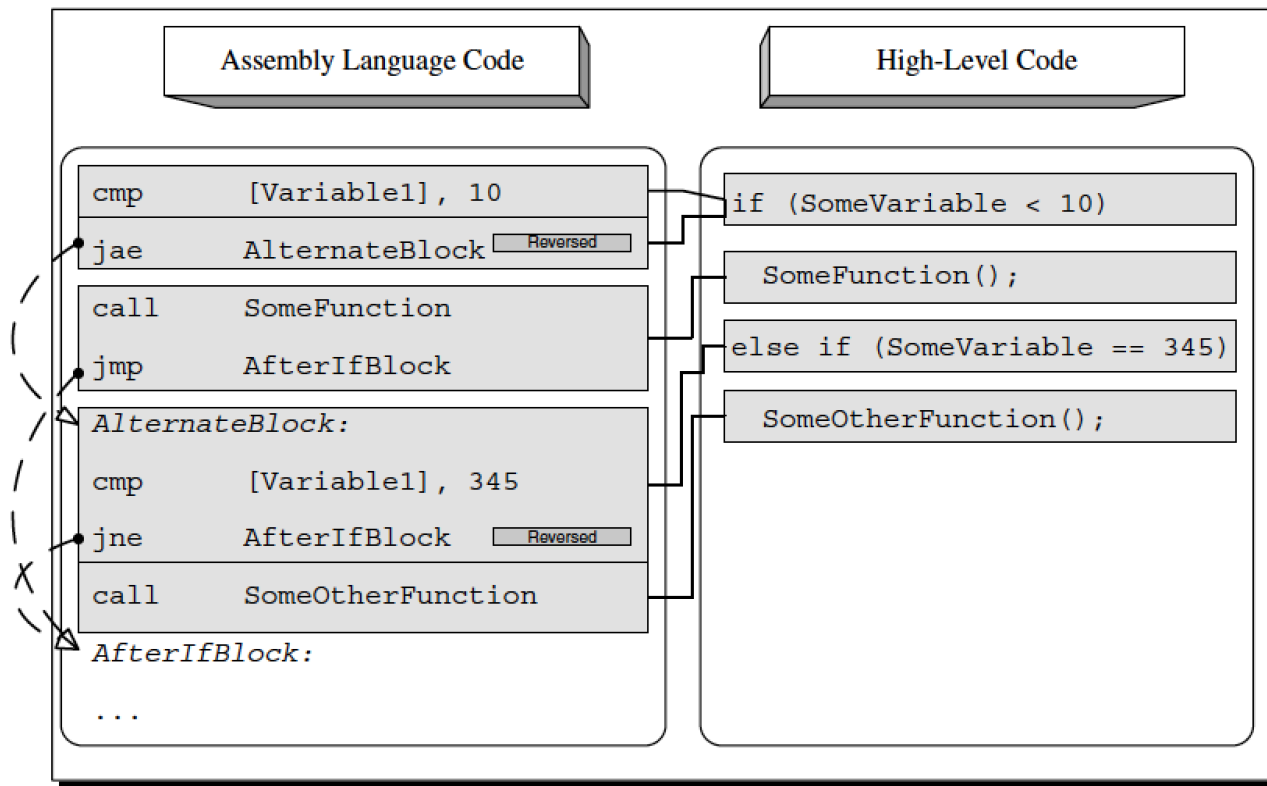
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- Multiple-alternative conditionals
  - Maybe implemented using "switch"
  - The compiler adds additional "alternate blocks" that consists of:
  1. One or more logical checks
  2. The actual conditional code block
  3. The final JMP that skips to the end of The entire block.

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- Compound conditions with logical operator **&&**
  - In assembly language:

```
cmp                    [Variable1], 100
jne                    AfterCondition
cmp                    [Variable2], 50
jne                    AfterCondition
ret
AfterCondition:
...
```
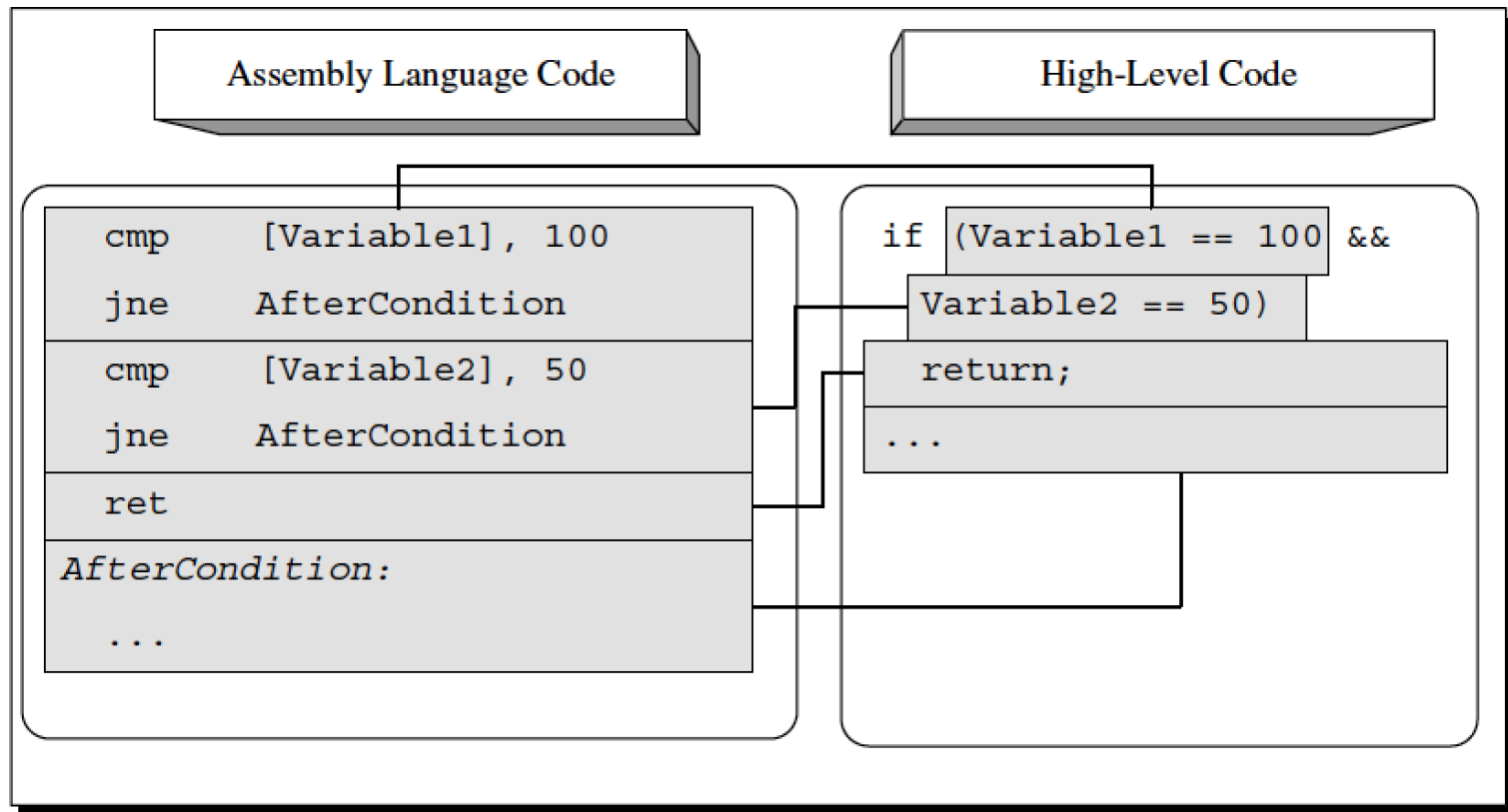
  - The equivalent high-level language

```
if (Variable1 == 100 && Variable2 == 50)
                                        return;
```

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- Compound conditions with logical operator &&

Assembly Language Code	High-Level Code
| ```
cmp     [Variable1], 100

jne     AfterCondition

cmp     [Variable2], 50

jne     AfterCondition

ret

AfterCondition:

    ...
``` | ```
if (Variable1 == 100 &&

    Variable2 == 50)

    return;

...
``` |
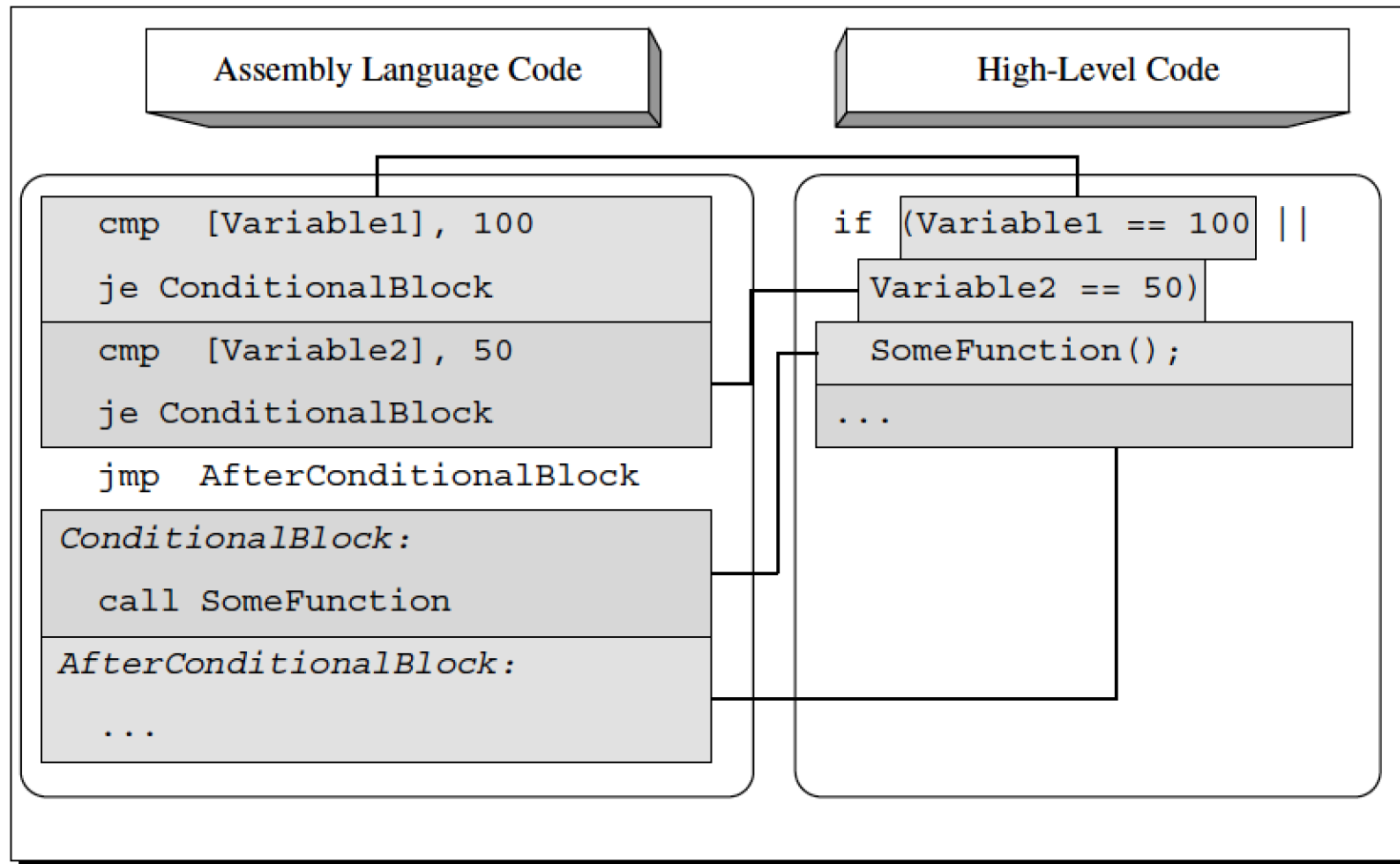
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- Compound conditions with logical operator ||
  - The assembly language version

```
cmp                     [Variable1], 100
je                      ConditionalBlock
cmp                     [Variable2], 50
je                      ConditionalBlock
jmp                     AfterConditionalBlock
ConditionalBlock:
call      SomeFunction
AfterConditionalBlock:
...
```

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- Compound conditions with logical operator ||

  –

Assembly Language Code	High-Level Code
| ```
cmp   [Variable1], 100
je ConditionalBlock
cmp   [Variable2], 50
je ConditionalBlock
jmp   AfterConditionalBlock
ConditionalBlock:
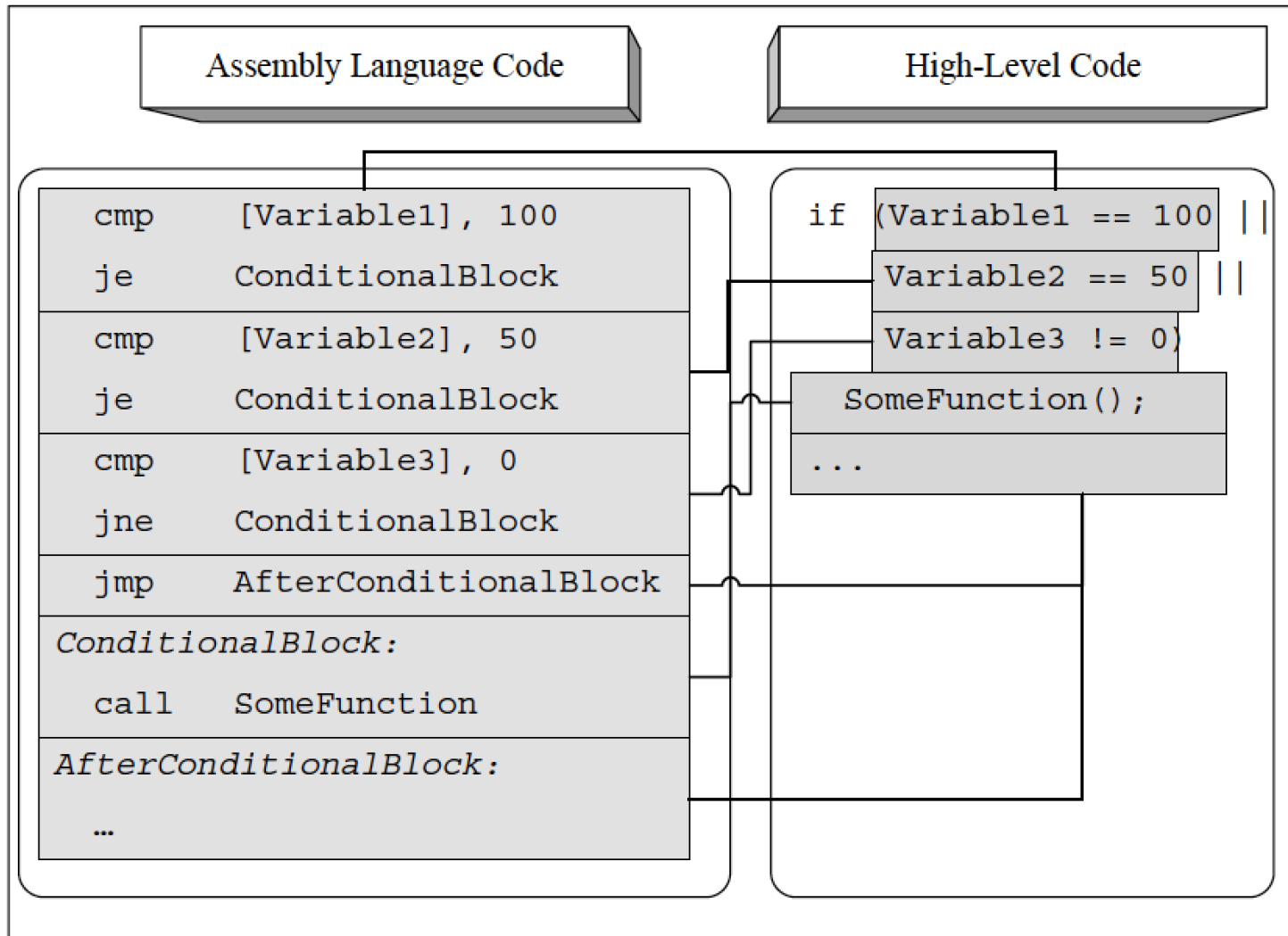  call SomeFunction
AfterConditionalBlock:
  ...
``` | ```
if (Variable1 == 100 ||
   Variable2 == 50)
  SomeFunction();
...
``` |
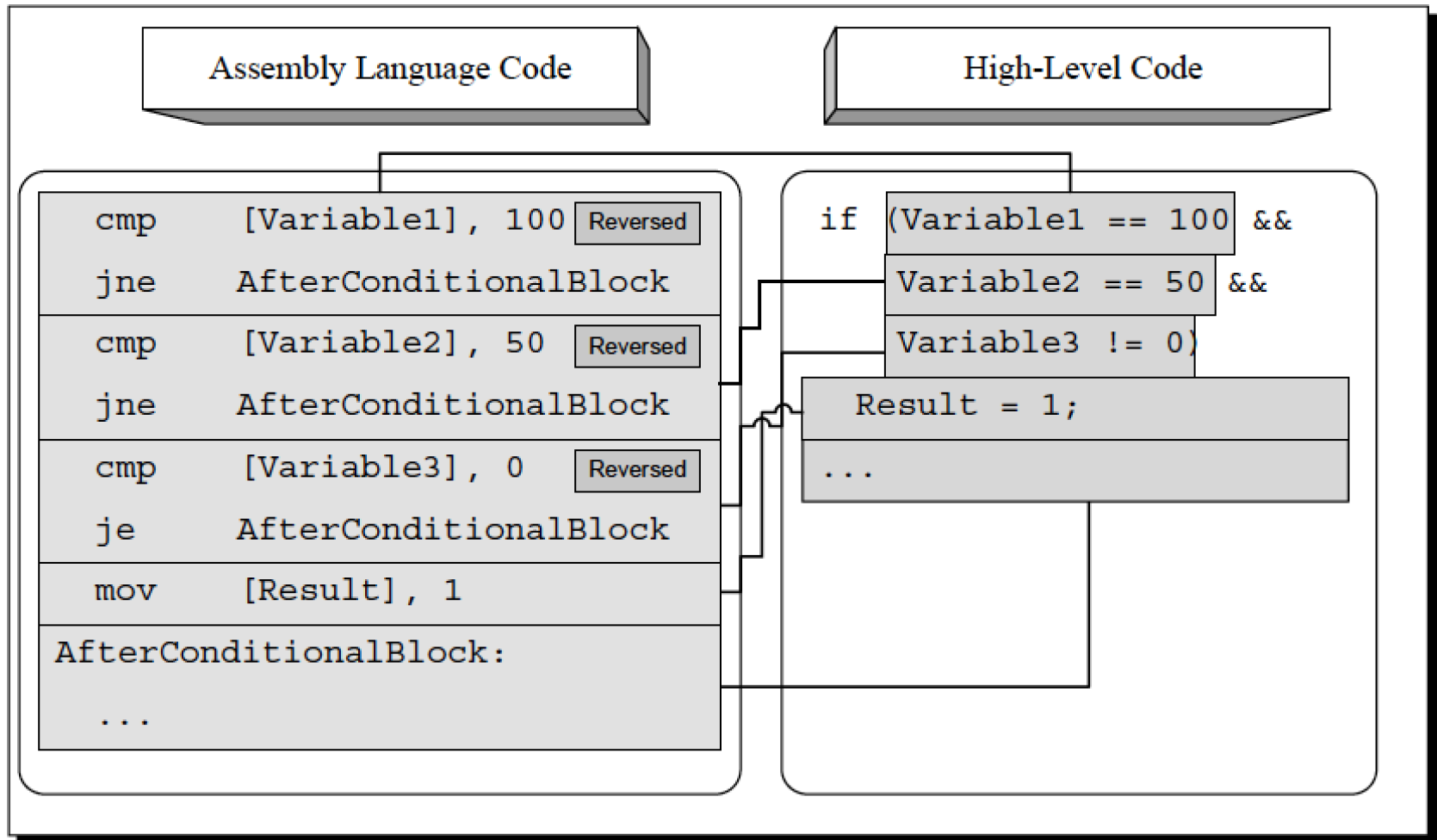
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- Compound conditions with three conditions combined using OR

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- Compound conditions with three conditions combined using AND

Assembly Language Code	High-Level Code

```
cmp      [Variable1], 100  [Reversed]
jne      AfterConditionalBlock
cmp      [Variable2], 50   [Reversed]
jne      AfterConditionalBlock
cmp      [Variable3], 0    [Reversed]
je       AfterConditionalBlock
mov      [Result], 1
AfterConditionalBlock:
   ...
```

```
if (Variable1 == 100 &&
    Variable2 == 50 &&
    Variable3 != 0)
    Result = 1;
    ...
```

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- Loops
  - A chunk of conditional code that is repeatedly executed until the condition is no longer satisfied

- Pretested loops
  - To present: a pretest loop the assembly language code must contain two jump instructions:
    - A conditional branch instruction in the beginning
    - An unconditional jump at the end that jumps back to the beginning of the loop

```
c = 0;
while (c < 1000)
{
    array[c] = c;
    c++;
}
```

```
mov          ecx, DWORD PTR [array]
xor          eax, eax
LoopStart:
mov          DWORD PTR [ecx+eax*4], eax
add          eax, 1
cmp          eax, 1000
jl           LoopStart
```

References

- Appendix A – Reversing the Secrets of Reverse Engineering (Eldad Eilam)