

Talker: 張志謙

PQC Final Report

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Definition:

要計算兩個多項式 $A(x), B(x)$ ，分別將其分成 k 部分（每部分長度固定），再對這些部分執行運算。隨著 k 的增長，可以組合許多乘法子運算，從而降低算法的整體複雜度，然後再次使用 *Toom - k* 算法遞歸計算乘法子運算，依此類推。

流程:

1. 拆分&求值 \Rightarrow Linear Transformation
2. 點乘
3. 插值 \Rightarrow Linear Transformation

拆分&求值:

Let $A(x) = a_{n-1}x^{n-1} + \dots, a_0$, $B(x) = b_{n-1}x^{n-1} + \dots, b_0$, and we set $y = x^{\frac{n}{k}}$,

$$\alpha_{k-1} = a_{n-1}x^{\frac{n}{k}-1} + \dots + a_{n-\frac{n}{k}}, \alpha_{k-2} = a_{n-1-\frac{n}{k}}x^{\frac{n}{k}-1} + \dots + a_{n-\frac{2n}{k}}, \dots$$

$$\beta_{k-1} = b_{n-1}x^{\frac{n}{k}-1} + \dots + b_{n-\frac{n}{k}}, \dots$$

$$\text{then } \begin{cases} A(y) = \alpha_{k-1}y^{k-1} + \dots \alpha_0 \\ B(y) = \beta_{k-1}y^{k-1} + \dots + \beta_0 \end{cases}$$

拆分&求值:

choose $2k - 1$ points p_0, \dots, p_{2k-2} on $A(y)$

$$\begin{bmatrix} A(p_0) \\ A(p_1) \\ \vdots \\ A(p_{2k-2}) \end{bmatrix} = \begin{bmatrix} (p_0)^0 & (p_0)^1 & \dots & (p_0)^{k-1} \\ (p_1)^0 & (p_1)^1 & \dots & (p_1)^{k-1} \\ \dots & \dots & \dots & \dots \\ (p_{2k-2})^0 & (p_{2k-2})^1 & \dots & (p_{2k-2})^{k-1} \end{bmatrix} \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \vdots \\ \alpha_{k-1} \end{bmatrix}$$

Toom-cook

點乘: Let $C(x) = A(x) \times B(x)$, then

$$\begin{bmatrix} C(p_0) \\ C(p_1) \\ \cdot \\ \cdot \\ C(p_{2k-2}) \end{bmatrix} = \begin{bmatrix} A(p_0)B(p_0) \\ A(p_1)B(p_1) \\ \cdot \\ \cdot \\ A(p_{2k-2})B(p_{2k-2}) \end{bmatrix}$$

插值:

Use $C(p_0), \dots, C(p_{2k-2})$ to recover $C(y) = c_{2k-2}y^{2k-2} + \dots, c_0$

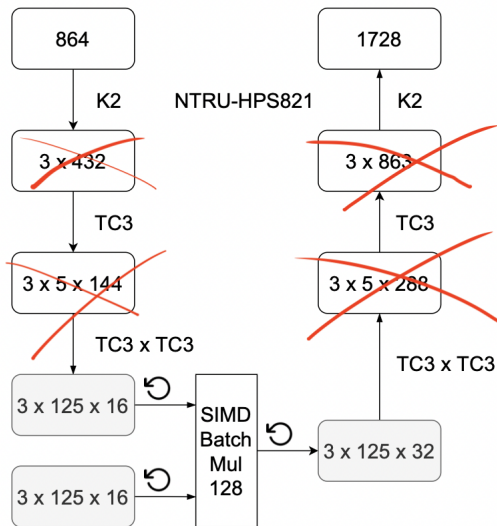
$$\begin{bmatrix} C(p_0) \\ C(p_1) \\ \vdots \\ C(p_{2k-2}) \end{bmatrix} = \begin{bmatrix} (p_0)^0 & (p_0)^1 & \dots & (p_0)^{k-1} \\ (p_1)^0 & (p_1)^1 & \dots & (p_1)^{k-1} \\ \dots & \dots & \dots & \dots \\ (p_{2k-2})^0 & (p_{2k-2})^1 & \dots & (p_{2k-2})^{k-1} \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_{2k-2} \end{bmatrix}$$

插值:

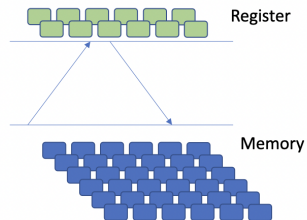
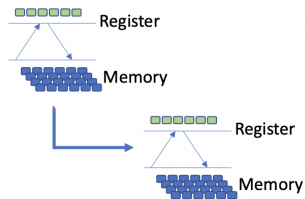
Use $C(p_0), \dots, C(p_{2k-2})$ to recover $C(y) = c_{2k-2}y^{2k-2} + \dots, c_0$

$$\begin{bmatrix} c_0 \\ c_1 \\ \vdots \\ c_{2k-2} \end{bmatrix} = \begin{bmatrix} (p_0)^0 & (p_0)^1 & \dots & (p_0)^{k-1} \\ (p_1)^0 & (p_1)^1 & \dots & (p_1)^{k-1} \\ \dots & \dots & \dots & \dots \\ (p_{2k-2})^0 & (p_{2k-2})^1 & \dots & (p_{2k-2})^{k-1} \end{bmatrix}^{-1} \begin{bmatrix} C(p_0) \\ C(p_1) \\ \vdots \\ C(p_{2k-2}) \end{bmatrix}$$

Improvement



Improvement



Improvement

Consider all column vectors in

$$\begin{bmatrix} (p_0)^0 & (p_0)^1 & \dots & (p_0)^{k-1} \\ (p_1)^0 & (p_1)^1 & \dots & (p_1)^{k-1} \\ \dots & \dots & \dots & \dots \\ (p_{2k-2})^0 & (p_{2k-2})^1 & \dots & (p_{2k-2})^{k-1} \end{bmatrix} :$$

if $\forall i, j$ satisfy $p_i \neq p_j$, then these vectors are independent.

Hence, Toom-Cook multiplication and interpolation are fundamentally linear maps of the form:

$$\begin{cases} \text{TC}_k : R_{\frac{n}{k}-1}^k(x) \rightarrow R_{\frac{n}{k}-1}^{2k-1}(x) \\ \text{TC}_k^{-1} : R_{\frac{2n}{k}-2}^{2k-1}(x) \rightarrow R_{\frac{2n}{k}-2}^k(x) \end{cases}$$

Improvement

- ▶ $\hat{TC}(A(x)) = TC_{k_\eta}(TC_{k_{\eta-1}}(\dots TC_{k_1}(A(x))))$
- ▶ $\theta = \hat{TC}(A(x)) * \hat{TC}(B(x))$
- ▶ $C(x) = TC_{k_1}^{-1}(TC_{k_2}^{-1}(\dots TC_{k_\eta}^{-1}(\theta)))$

Improvement

Definition:

$$O_n = \begin{bmatrix} 0 & \dots & 0 \\ \cdot & \dots & \cdot \\ \cdot & \dots & \cdot \\ 0 & \dots & 0 \end{bmatrix}_{n \times n}, I_n = \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \cdot & \dots & \dots & \cdot \\ \cdot & \dots & \dots & \cdot \\ 0 & \dots & \dots & 1 \end{bmatrix}_{n \times n},$$

$$T_n = \begin{bmatrix} I_n & \\ I_n & I_n \\ & I_n \end{bmatrix}_{3n \times 2n}$$

Improvement

$$T_{\frac{n}{2}} = \begin{bmatrix} I_{\frac{n}{2}} & \\ I_{\frac{n}{2}} & I_{\frac{n}{2}} \\ & I_{\frac{n}{2}} \end{bmatrix}_{\frac{3}{2}n \times n}$$

$$\Rightarrow \begin{bmatrix} T_{\frac{n}{4}} & & \\ & T_{\frac{n}{4}} & \\ & & T_{\frac{n}{4}} \end{bmatrix}_{\frac{3}{2}n \times n} T_{\frac{n}{2}} =$$

$$\begin{bmatrix} I_{\frac{n}{4}} & & & & & & \\ I_{\frac{n}{4}} & I_{\frac{n}{4}} & & & & & \\ & I_{\frac{n}{4}} & & & & & \\ & & I_{\frac{n}{4}} & & & & \\ & & I_{\frac{n}{4}} & I_{\frac{n}{4}} & & & \\ & & & I_{\frac{n}{4}} & & & \\ & & & & I_{\frac{n}{4}} & & \\ & & & & & I_{\frac{n}{4}} & \\ & & & & & I_{\frac{n}{4}} & I_{\frac{n}{4}} \\ & & & & & I_{\frac{n}{4}} & \\ & & & & & & I_{\frac{n}{4}} \end{bmatrix}_{\frac{9}{4}n \times \frac{3}{2}n}$$

$$\begin{bmatrix} I_{\frac{n}{2}} & \\ I_{\frac{n}{2}} & I_{\frac{n}{2}} \\ & I_{\frac{n}{2}} \end{bmatrix}_{\frac{3}{2}n \times n}$$

Reference

Time-memory trade-off in Toom-Cook multiplication: an application to module-lattice based cryptography