

Lab 5: Bubble Sort

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Download New .sof File for DE0

- A new LC3_data_path.sof is placed in the lab folder
- Reminder: address for LEDs is 0xFFFFD
- Encoding for HEX10 (address 0xFFFFF) and Hex32 (address 0xFFFFE) is already done by the Verilog code of the soft core LC3
- AND R0,R0,#0
- ADD R1,R0,#15
- STR R1,R0,#-1 ; show 0F on Hex10
- STR R1,R0,#-2 ; show 0F on Hex32

Check HEXs, LEDs with This Program

- .ORIG x3000
- AND R0,R0,#0
- LD R1,n255
- loop STR R1,R0,#-1 ; write to HEX10
- LD R2,time
- delay ADD R2,R2,#-1
- BRp delay
- ADD R1,R1,#-1
- BRzp loop
- LD R1,n255
- loop2 STR R1,R0,#-2 ; write to HEX32
- LD R2,time
- delay2 ADD R2,R2,#-1
- BRp delay2
- ADD R1,R1,#-1
- BRzp loop2
- LD R1,n255
- loop3 STR R1,R0,#-3 ; write to LEDs
- LD R2,time
- delay3 ADD R2,R2,#-1
- BRp delay3
- ADD R1,R1,#-1
- BRzp loop3
- bye BRnzp bye
- n255 .FILL x00FF
- time .FILL x0080
- .END

Task 1: Bubble Sort

- Write the code to bubble sort a list of 10 numbers
- At the end, all numbers should be sorted in increasing order
- Show the number of swaps on LED9_0
- Show the number of comparisons on Hex3_0
- Repeat experiment with 20 numbers and 30 numbers (or more if needed) with data of decreasing numbers
- Record #swaps and #comparisons for the three experiments in a table

bubble0.c

- `#include <stdio.h>`
- `void swap_next(int* array, int j)`
 - `{`
 - `int temp;`
 - `temp = array[j];`
 - `array[j] = array[j+1];`
 - `array[j+1] = temp;`
 - `}`
- `int main(void) {`
- `int size=10, i, j, temp;`
- `int nums[size];`
- `for (i = 0; i<size; i = i+1)`
- `nums[i] = size-i-1;`
- `for (i = size-1; i>0; i = i-1) {`
- `for (j = 0; j<i; j = j+1)`
- `if (nums[j]>nums[j+1])`
- `swap_next(nums, j);`
- `}`
- `}`

bubble1.c

- `int main(void) {`
- `int size=10, i, j, temp, c=0,`
`swaps=0;`
- `int nums[size];`
- `for (i = 0; i<size; i = i+1)`
- `nums[i] = size-i-1;`
- `for (i = size-1; i>0; i = i-1) {`
- `counter_j = i;`
- `for (j = 0; j<i; j = j+1) {`
- `c = c + 1;`
- `if (nums[j] > nums[j+1]) {`
- `swap_next(nums, j);`
- `swaps = swaps+1;`
- `}`
- `counter_j = counter_j-1;`
- `if (counter_j<0) exit j loop;`
- `}`
- `if (i=0) exit i loop;`
- `}`
- `}`

Getting from C to Assembly Language

- for (i = size-1; i>0; i = i-1) {
 - for (j = 0; j<i; j = j+1) {
 - if (nums[j] > nums[j+1]) {
 - swap_next(nums, j);
 - }
 - }
 - }
- for (i = size-1; i>0; i = i-1) {
 - for (k = i; k>0; k = k-1) {
 - j = size-1-k;
 - if (nums[j] > nums[j+1]) {
 - swap_next(nums, j);
 - }
 - }
 - }
 - ; LC-3: BRp inLoop
 - ; exit inLoop when k=0

Task 1: Example

- $i = 4$
- $j = 0, 1, 2, 3, 4$
- $a[j] = 4, 3, 2, 1, 0$
- At $j=0$, $a[0]=4 > a[1]=3$, swap to get $a[]: 3, 4, 2, 1, 0$
- At $j=1$, $a[1]=4 > a[2]=2$, swap to get $a[]: 3, 2, 4, 1, 0$
- At $j=2$, $a[2]=4 > a[3]=1$, swap to get $a[]: 3, 2, 1, 4, 0$
- At $j=3$, $a[3]=4 > a[4]=0$, swap to get $a[]: 3, 2, 1, 0, 4$
- Ignore $a[4]=4$ from now on because it is the max no. correctly placed

- $i = 3, j = 0, 1, 2, 3$; result $a[]: 2, 1, 0, 3, 4$
- $i = 2, j = 0, 1, 2$; result $a[]: 1, 0, 2, 3, 4$
- $i = 1, j = 0$; result $a[]: 0, 1, 2, 3, 4$
- SORTING DONE!

Bubble.asm

- .orig x3000
- ...
- Size .FILL #10
- Mtable .FILL #9
- .FILL #8
- .FILL #7
- .FILL #6
- .FILL #5
- .FILL #4
- .FILL #3
- .FILL #2
- .FILL #1
- .FILL #0

Time Complexity of Bubble Sort

- For arrays of different sizes, count the accumulated no. of swaps and no. of comparisons
- Record in a table of array size and accumulated no. of swaps and total no. of comparisons
- Write down:
 - $s = f(n)$
 - $c = g(n)$

Array Size n	#Swaps s	#Comparisons c	Time t (sec)
10			
20			
30			

Task 2: Put Maximum Element at the End

- Improve bubble sort
- Do not swap two neighbor elements
- $i = \text{size}-1, \text{size}-2, \dots, 1$
 - Check all elements and find the maximum no. in the remaining elements
 - Then swap `nums[i]` with the maximum element
 - This puts the max. element in `nums[i]`
- Why would this reduce the no. of swaps?
- Repeat the experiment for size of 10, 20, 30 (or more if needed).
- Record #swaps and #comparisons in table.
- In this new algorithm, is #comparisons a better or worse measure of time complexity than #swaps? Why?
- Is this more efficient than bubble sort? Why?

Task 2: Example

- $i = 0, 1, 2, 3, 4, 5, 6, 7, 8, 9$
- $a[i] = 4, 3, 2, 1, 0, 9, 8, 7, 6, 5$
- $\text{Max} = 4$, i of $\text{max} = 0$
- At $i=1$, $\text{max} = 4 > a[1] = 3$, so no update
- At $i=2$, $\text{max} = 4 > a[2] = 2$, so no update
- At $i=3$, $\text{max} = 4 > a[3] = 1$, so no update
- At $i=4$, $\text{max} = 4 > a[4] = 0$, so no update
- At $i=5$, $\text{max} = 4 < a[5] = 9$, so update $\text{max} := 9$, $i_{\text{max}} := 5$
- At $i=6$, $\text{max} = 9 > a[6] = 8$, so no update
- ...
- At $i=9$, $\text{max} = 9 > a[9] = 5$, so no update

Example

- i: 0,1,2,3,4,5,6,7,8,9
- Data: 9,8,7,6,5,4,3,2,1,0
- Check i=0 to 9: max=a[0]=9, i_max=0
- Swap: 0,8,7,6,5,4,3,2,1, 9
- Check i=0 to 8: max=a[1]=8, i_max=1
- Swap: 0,1,7,6,5,4,3,2, 8,9
- Check i=0 to 7: max=a[2]=7, i_max=2
- Swap: 0,1,2,6,5,4,3, 7,8,9
- Check i=0 to 6: max=a[3]=6, i_max=3
- Swap: 0,1,2,3,5,4, 6,7,8,9
- Check i=0 to 5: max=a[4]=5, i_max=4
- Swap: 0,1,2,3,4, 5,6,7,8,9
- ...

swap_max.c

- void swap(int* array, int j, int k) {
- int temp;
- temp = array[j];
- array[j] = array[k];
- array[k] = temp;
- }
- int main(void) {
- int size=10, i, j, k, temp, swaps=0, max, i_max;
- int nums[size];
- for (i = 0; i<size; i = i+1)
- nums[i] = size-i-1;
- for (i = size-1; i>0; i = i-1) {
- max = nums[0];
- i_max = 0;
- for (j = 1; j <= i; j = j+1) {
- if (nums[j] > max) {
- max = nums[j];
- i_max = j;
- }
- }
- swap(nums, i, i_max);
- swaps = swaps + 1;
- }
- }

Record Experimental Data and Explain

- Record the experimental data in a table like the one on the right
- Compare the performance of this algorithm to the previous bubble sort algorithm
- Difference in #swaps?
- Difference in #comparisons?

Array Size n	#Swaps s	#Comparisons c	Time t (sec)
10			
20			
30			