

17.15

17.15 Consider the following two transactions:

```
T13: read(A);
      read(B);
      if A = 0 then B := B + 1;
      write(B).
T14: read(B);
      read(A);
      if B = 0 then A := A + 1;
      write(A).
```

Let the consistency requirement be  $A = 0 \vee B = 0$ , with  $A = B = 0$  as the initial values.

- a. Show that every serial execution involving these two transactions preserves the consistency of the database.
- b. Show a concurrent execution of  $T_{13}$  and  $T_{14}$  that produces a nonserializable schedule.
- c. Is there a concurrent execution of  $T_{13}$  and  $T_{14}$  that produces a serializable schedule?

Answer:

a.

假如顺序为 $T_{13} \rightarrow T_{14}$ , 那么最终结果为 $A = 0, B = 1$ , 满足 $A = 0 \vee B = 0$

假如顺序为 $T_{14} \rightarrow T_{13}$ , 那么最终结果为 $A = 1, B = 0$ , 满足 $A = 0 \vee B = 0$

b.

T13	T14
read(A)	
	read(B);
	read(A);
	if B=0 then A:=A+1;
	write(A);
read(B);	
if A=0 then B:=B+1;	
write(B);	

这种并行会破坏一致性

C.

不存在任何并行执行顺序可以不破坏一致性，原因如下：

假设 $T_{14}$ 在 $T_{13}$ 执行完成前并行执行，无论在什么时候执行 $read(B)$ ，得到的 $B$ 值均为0，结果导致 $T_{14}$ 将 $A$ 修改为1，最后结果为 $A = 1, B = 1$ 。

$T_{13}$ 在 $T_{14}$ 执行完成前并行也会导致一样的结果，所以任何并行执行都会破坏一致性条件 $A = 0 \vee B = 0$

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