

数据库原理CH16作业

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16.5

Consider the relations $r_1(A, B, C)$, $r_2(C, D, E)$, and $r_3(E, F)$, with primary keys A , C , and E , respectively. Assume that r_1 has 1000 tuples, r_2 has 1500 tuples, and r_3 has 750 tuples. Estimate the size of $r_1 \bowtie r_2 \bowtie r_3$, and give an efficient strategy for computing the join.

Answer:

因为C是 r_2 的主码，所以应该先将 r_1 和 r_2 连接形成 r_{12} ，

因为E是 r_3 的主码，所以将 r_{12} 与 r_3 连接，

最后的顺序应该为 $(r_1 \bowtie r_2) \bowtie r_3$ 。

16.8

16.8 Consider the query:

```
select *  
from  $r, s$   
where upper( $r.A$ ) = upper( $s.A$ );
```

where “upper” is a function that returns its input argument with all lowercase letters replaced by the corresponding uppercase letters.

- Find out what plan is generated for this query on the database system you use.
- Some database systems would use a (block) nested-loop join for this query, which can be very inefficient. Briefly explain how hash-join or merge-join can be used for this query.

Answer:

a.

经过上网搜索，MYSQL 5.5版本以前都使用的是**Nested-Loop Join**。

5.5版本以后，MYSQL使用**Block Nested-Loop Join**。

b.

先执行 $\text{upper}(r.A)$ 和 $\text{upper}(s.A)$ ，再执行 hash-join 或者 merge-join。

