

Binary Search—Complexity Class: $O(\log N)$

* Only works if the list is sorted

1. Compare the element at the middle position in the list to the target value.
2. If the target value is equal to the element at the middle position, then you are done.
3. If the target value is less than the element at the middle position, then repeat the procedure starting from step 1 for the left half of the list.
4. If the target value is greater than the element at the middle position, then repeat the procedure starting from step 1 for the right half of the list.

Note: If either the left or right sides of the list are empty for steps 3 or 4, then the target value is not contained in the list.

Target Value: 9

Index	0	1	2	3	4	5	6	7	8
	-3	6	9	12	15	18	21	24	27

↑
LOW↑
MID↑
HIGH

Since 9 is less than 15 and the list is sorted, we know 9 can't possibly be in the second half of the list. So we only continue searching in the first half.

Index	0	1	2	3	4	5	6	7	8
	-3	6	9	12	15	18	21	24	27

↑
LOW↑
MID↑
HIGH

9 is greater than 6, so we continue searching in the second half of the list.

Index	0	1	2	3	4	5	6	7	8
	-3	6	9	12	15	18	21	24	27

↑
LOW↑
MID↑
HIGH

We found 9!

Sorting Algorithms:

	Complexity	The Steps	Visual Representation																																																																																																				
Selection	$O(N^2)$	<ol style="list-style-type: none">1. Look through the entire list for the smallest value.2. Swap the smallest value with the value at the current index (Unless current index contains the smallest value).3. Increase current index.4. Look through the rest of the list for the smallest value.5. Swap this value with the value at current index.6. Repeat for the rest of the list.	<p>(Shaded boxes indicate swapped values)</p> <table><tr><th>Index</th><td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td><td>7</td><td>8</td></tr><tr><td></td><td>7</td><td>4</td><td>2</td><td>16</td><td>22</td><td>13</td><td>15</td><td>31</td><td>0</td></tr></table> <p>↑ Current Index ↑ Smallest Value</p> <table><tr><th>Index</th><td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td><td>7</td><td>8</td></tr><tr><td></td><td>0</td><td>4</td><td>2</td><td>16</td><td>22</td><td>13</td><td>15</td><td>31</td><td>7</td></tr></table> <p>↑ Current Index ↑ Smallest Value</p> <table><tr><th>Index</th><td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td><td>7</td><td>8</td></tr><tr><td></td><td>0</td><td>2</td><td>4</td><td>16</td><td>22</td><td>13</td><td>15</td><td>31</td><td>7</td></tr></table> <p>↑ Current Index In this case, 4 is the smallest value so we don't need to swap anything.</p> <table><tr><th>Index</th><td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td><td>7</td><td>8</td></tr><tr><td></td><td>0</td><td>2</td><td>4</td><td>16</td><td>22</td><td>13</td><td>15</td><td>31</td><td>7</td></tr></table> <p>↑ Current Index ↑ Smallest Value</p> <table><tr><th>Index</th><td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td><td>7</td><td>8</td></tr><tr><td></td><td>0</td><td>2</td><td>4</td><td>7</td><td>22</td><td>13</td><td>15</td><td>31</td><td>16</td></tr></table> <p>This process continues until you've reached the end of the list.</p>	Index	0	1	2	3	4	5	6	7	8		7	4	2	16	22	13	15	31	0	Index	0	1	2	3	4	5	6	7	8		0	4	2	16	22	13	15	31	7	Index	0	1	2	3	4	5	6	7	8		0	2	4	16	22	13	15	31	7	Index	0	1	2	3	4	5	6	7	8		0	2	4	16	22	13	15	31	7	Index	0	1	2	3	4	5	6	7	8		0	2	4	7	22	13	15	31	16
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Mergesort	$O(N \log N)$	<ol style="list-style-type: none"> 1. Repeatedly divide the list into two equal parts until each part is a single element of the list 2. Combine the parts in sorted order, until the list is completely reconstructed. 	<p>The diagram illustrates the Mergesort algorithm. It starts with an array [8, 0, 7, 4]. The process involves repeatedly dividing the array into two halves until each element is in its own box. Then, the elements are combined and sorted in pairs, and finally, the entire array is reconstructed in sorted order.</p>
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Insertion	$O(N^2)$	<ol style="list-style-type: none"> 1. Divide list into two imaginary lists: sorted (initially empty) and unsorted (the rest of the elements). 2. Take the first element from the unsorted list and place into sorted. 3. Take the next element from the unsorted list and insert the value into the correct location. 4. Repeat until the unsorted part is empty. 	<p>The diagram illustrates the Insertion Sort algorithm. It shows an 'Unsorted List' and a 'Sorted List'. Elements are moved from the unsorted list to the sorted list one by one, inserting them into their correct position. The final sorted list is [-1, 0, 4, 5, 8, 13].</p>
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