

$\times \times \times$

Normal
 $\rightarrow w$

Plane

Apply when data
is linearly separable

- only for binary
classification

$0 0 0$
 $0 0 0 0$
 $0 0 0$
 0

$$w^* = \arg \max_w \sum_{i=1}^n y_i w^T x_i$$

optimization
problem

objective is to find optimal value for "w"

Squashing of Sigmoid function

$\pi: w_i: \text{Best / optimal } w \rightarrow w^*$

$$\arg \max_w \sum_{i=1}^n y_i w^T x_i$$

→ signed distance

$w^T x_i$: distance from "x_i" to " π " (w is unit vector)

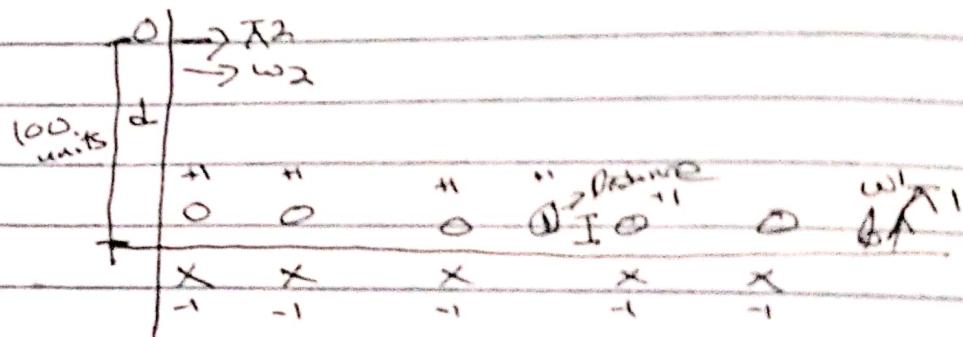
if

$y_i w^T x_i > 0 \Rightarrow \pi$ as defined by us correctly classifies " x_i "

$\Rightarrow -ve \Rightarrow$ incorrectly classifies.

Example:

assume distance of all the points from plane is ~~is~~ "1"



Case 1:- π_1 is class separating hyperplane

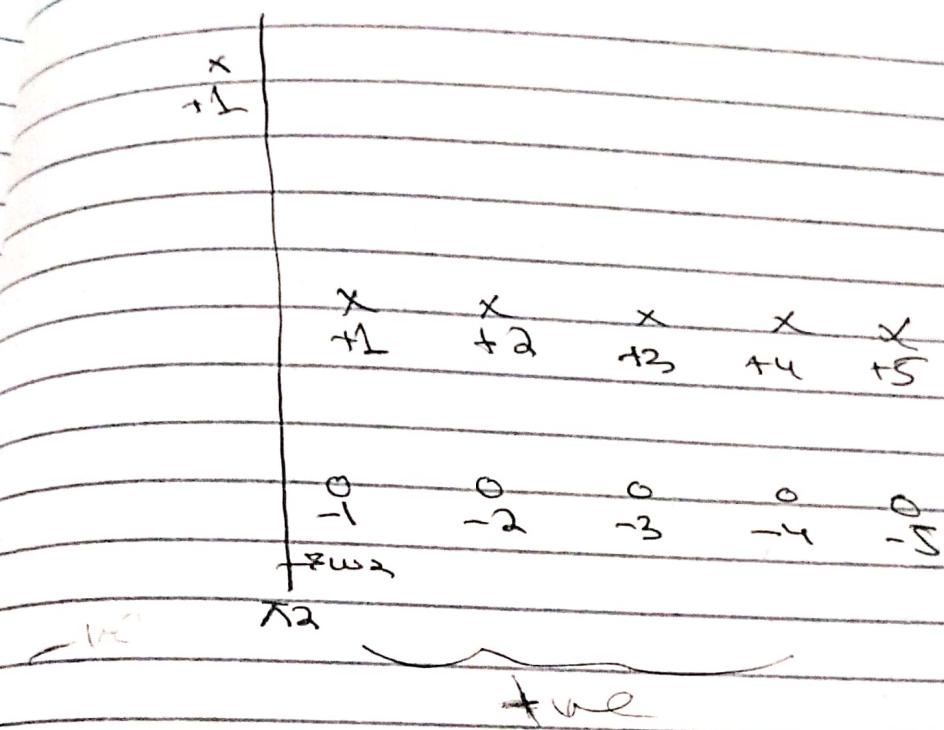
$$\sum_{i=1}^n y_i w^T x_i = 1+1+1+1+1 \quad (+ve) \\ 1+1+1+1+1 \quad (-ve) \\ -100 \\ = -90$$

Case 2:- π_2 is class separating plane

Obj: is to find best π which maximize the sum of signed distance

One Single Outlier is changing
the old model which is
very terrible Situation.

distance from
all points
is least



$$w_1^T x_i = 1 + 2 + 3 + 4 + 5 = 15$$

$$w_2^T x_i = -1 - 2 - 3 - 4 - 5 = -15$$

+1 outlier

Obj: find
w such that
maximize
Signed
distance

we get value of w_1 using L1

$$|+1| > |-1|$$

w_1 is my classifier

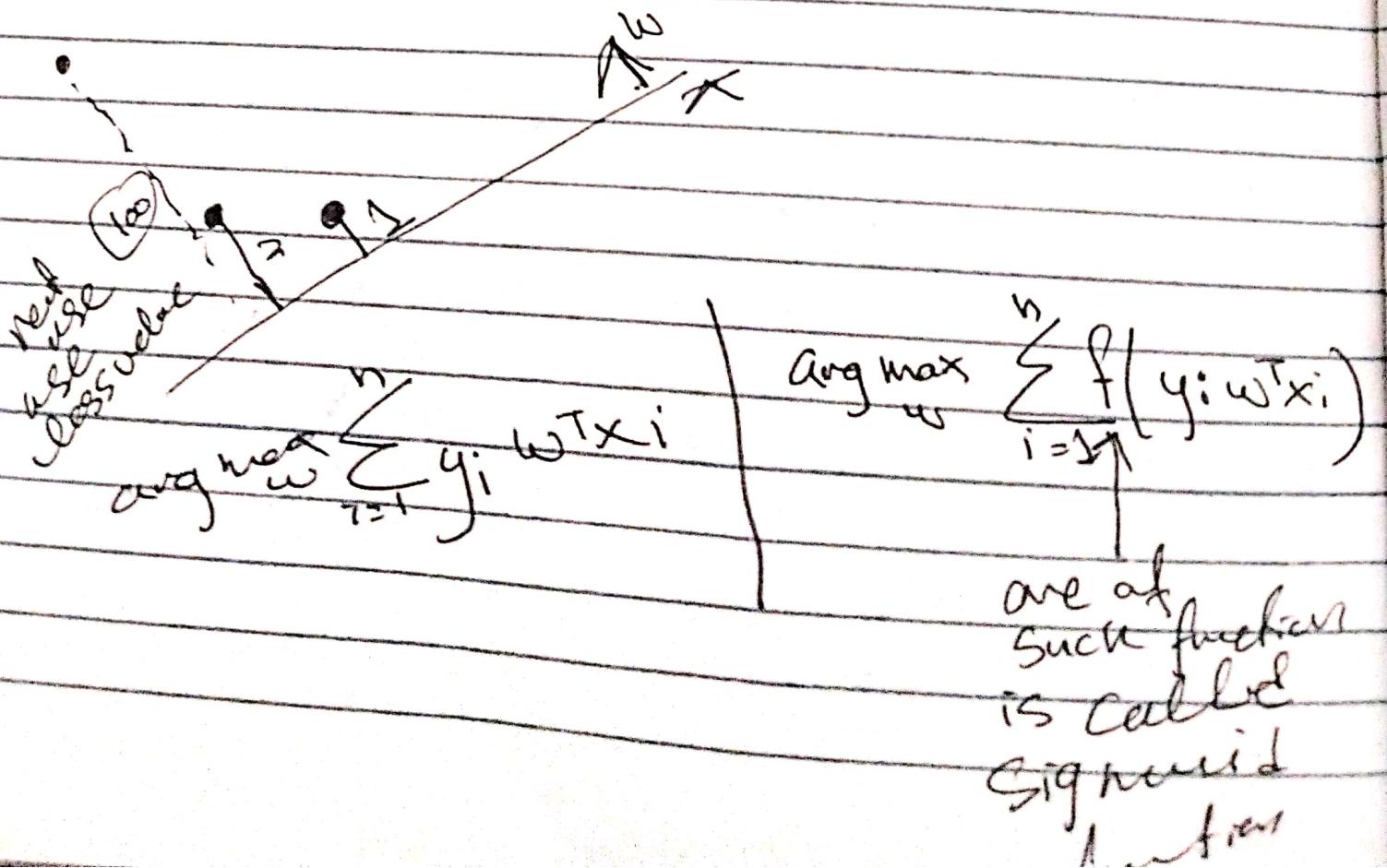
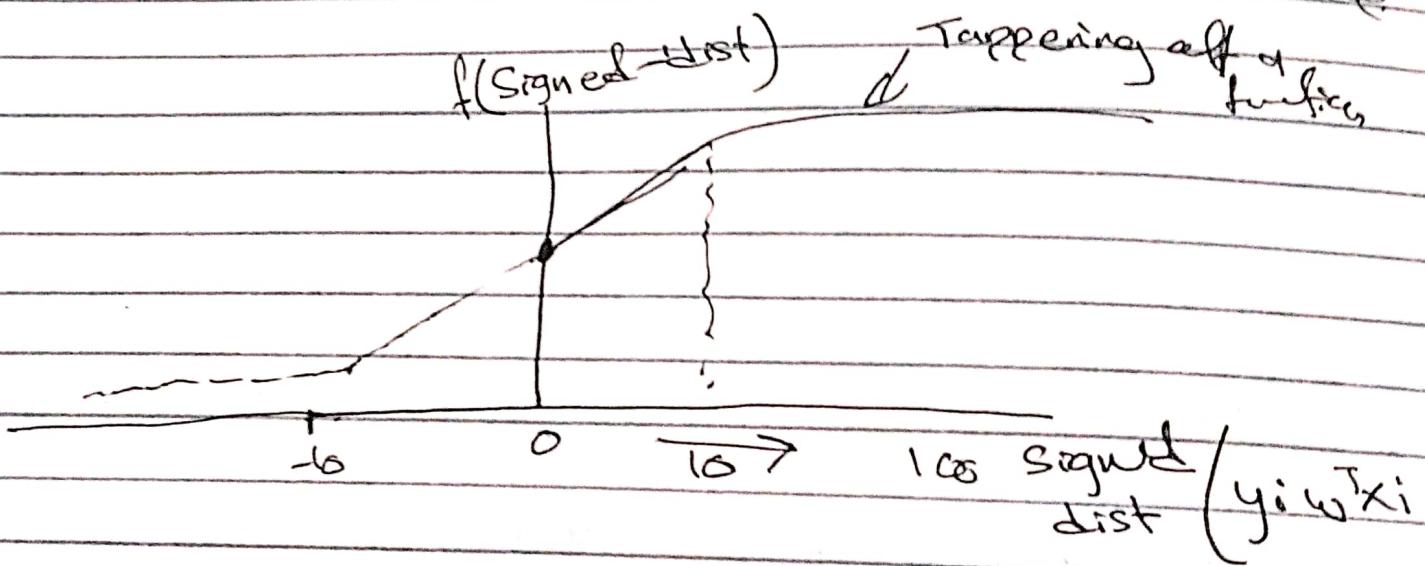
Max Sum of Signed dist

↓
Prove about it

Squashing :-

Idea: instead of using simple signed distance we have to modify a little bit.

if Signed distance is small : use it as it is
if " " " larger - make it a smaller value

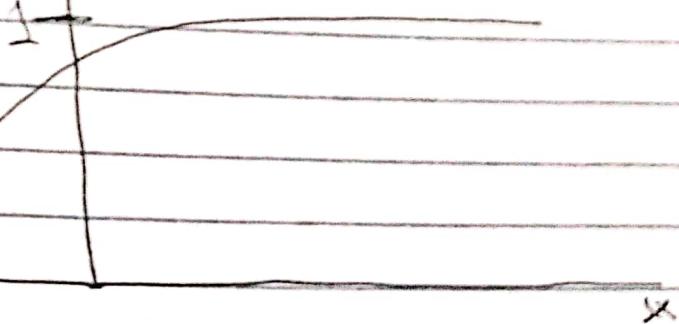


Sigmoid function

$$\delta(u) = \frac{1}{1+e^{-x}}$$

max value = 1
min value = 0

x = Signed distance



$$\delta(0) = 0.5$$

The reason we choose Sigmoid function because it has nice probabilistic interpretation.

$w^T x_i$ is very large

$$\delta(w^T x_i) = 1$$

+

Pointedly
or very
close

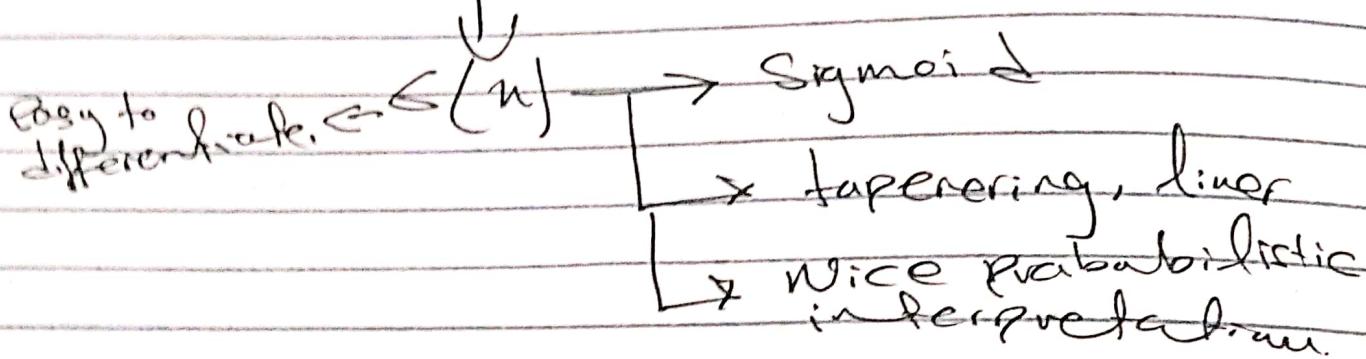
$$w^T x_i > 0$$

$$\delta(w^T x_i) = 0$$

$$\delta(w^T x_i) = 0$$

Deriving Logistic regression using geometry.

max. Sum. of Signed dist \rightarrow outlier problem



Max. Sum of Transformed signed dist.

$$w^* = \arg \max_w \zeta(w)$$

$$w^* = \arg \max_w \sum_{i=1}^n \zeta(y_i w^T x_i)$$



$$w^* = \arg \max_w \sum_{i=1}^n \frac{1}{1 + \exp(y_i w^T x_i)}$$

This function is called Squashing because it squeezes the value which is $(-\infty, +\infty)$ to $(0, 1)$.

Optimization Problem:

$$w^* = \arg \max_w \sum_{i=1}^n \frac{1}{1 + \exp(-y_i w^T x_i)}$$

optimization problem

→ Monotonic function.

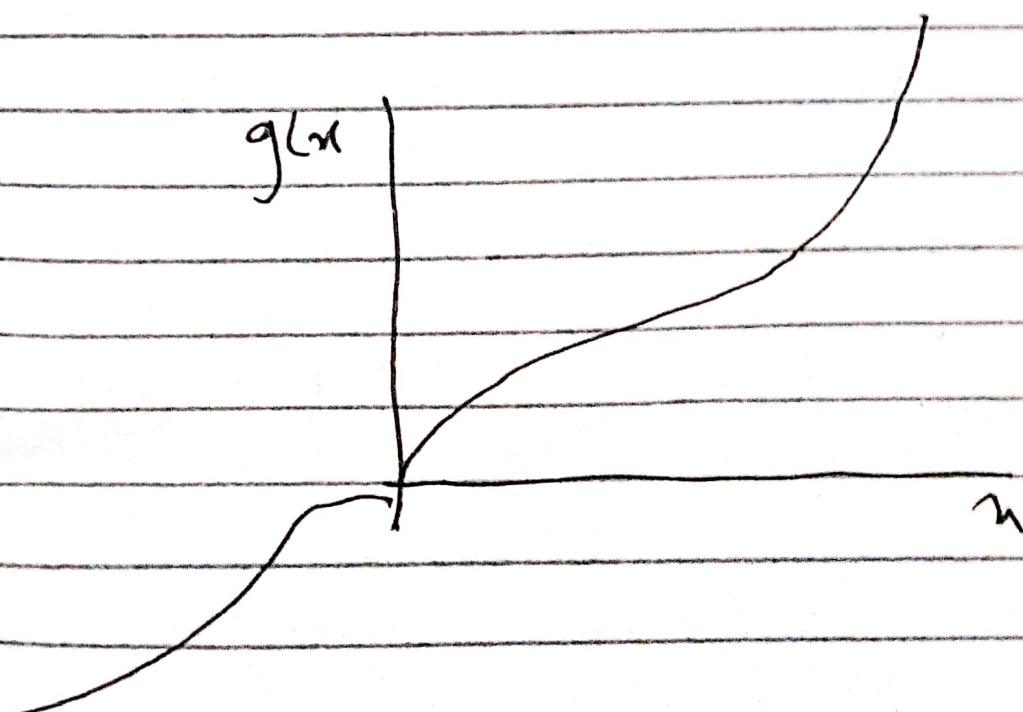
we can solve this problem easily

a function $f(n)$ is said to be increasing

if n increase $f(n)$ increase

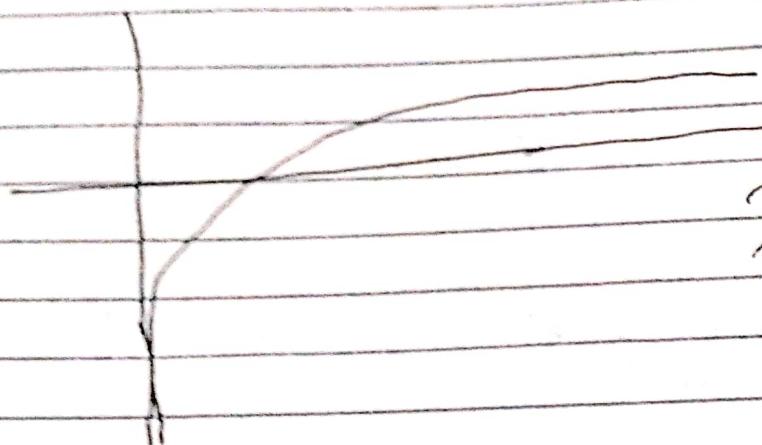
if $n_1 > n_2$ then $g(n_1) > g(n_2)$

then $g(n)$ is called monotonically increasing function.



$\log(n) \quad n > 0$

$\log(n); n > 0$



$n \uparrow \log(n) \uparrow$
 $n \downarrow \log(n) \downarrow$

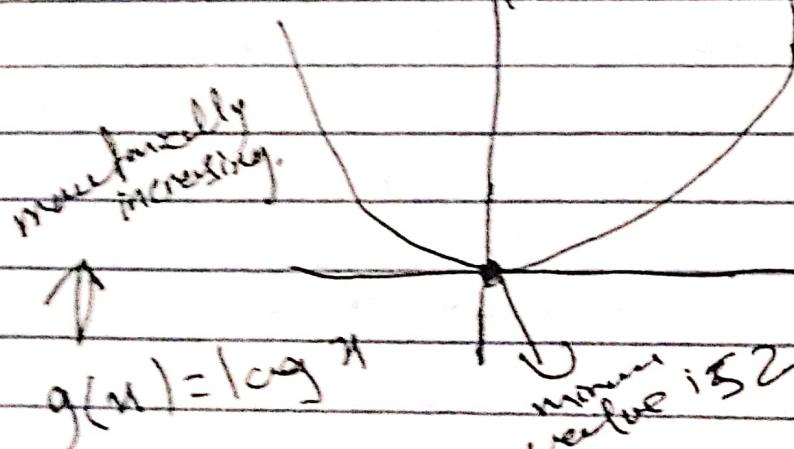
opt - Prob:

$$x^* = \underset{x}{\operatorname{argmin}} \quad x^2$$

"we want to find
n which minimizes
 x^2 "

① $\rightarrow x^*$

$$f(n) = n^2$$



x^2 is monotonically increasing when $n > 0$
and monotonically decreasing when $n < 0$

$$x^* = \underset{n}{\operatorname{argmin}} \quad f(x) \quad ; \quad f(n) = n^2 \text{ problem 1}$$

$$x^* = \underset{n}{\operatorname{argmax}} \quad g(f(n)) \quad ; \quad f(x) = n^2 \text{ problem 2}$$

$$g(n) = \log(n)$$

$$f(n) = n^2$$

$$n^* = \arg \min_n f(n)$$

$$n' = \arg \min_n g(f(n)) = \log(n^2)$$

n^* and n' is same because $g(n)$ is monotonically increasing $f(n)$.

Theorem:

if $g(n)$ is monotonic function.

$$\arg \min_n f(n) = \arg \min_n g(f(n))$$

$$\arg \max_n f(n) = \arg \max_n g(f(n))$$

Original Problem.

$$w^* = \arg \max_w \sum_{i=1}^n \frac{1}{1 + e^{-y_i w^T x_i}}$$

$$g(n) = -\log(n)$$

$$w^* = \arg \max_w \sum_{i=1}^n \log(g(y_i w^T x_i))$$

$$w^* = \arg \max_w \sum_{i=1}^n \log \left[\frac{1}{1 + \exp(-y_i w^T x_i)} \right]$$

$$\log(1/n) = -\log(n)$$

$$w^* = \arg \max_w \sum_{i=1}^n -\log(1 + \exp(-y_i w^T x_i))$$

$$\arg \max_x f(x) = \arg \min_x -f(x)$$

$$\arg \max_x -f(x) = \arg \min_x f(x) \quad [-1 or 1]$$

$$w^* = \arg \min_w \sum_{i=1}^n \log \left[\frac{1}{1 + \exp(-y_i w^T x_i)} \right]$$

Solution of optimization problem

$$w^* = \operatorname{argmin}_w \sum_{i=1}^n \log(1 + \exp(-y_i w^T x_i))$$

$$\hat{w}^* = \operatorname{argmin}_w \sum_{i=1}^n -y_i w^T x_i$$

$$w^* = \operatorname{argmax}_w \sum_{i=1}^n y_i w^T x_i$$

↓
huge outliers problem

$$w^* = \operatorname{argmin}_w \sum_{i=1}^n \log(1 + \exp(-y_i w^T x_i))$$

+1,-1

$$\left\{ \begin{array}{l} w^* = \operatorname{argmin}_w \sum_i -y_i \log p_i - (1-y_i) \log(1-p_i) \\ p_i \in [w^T x_i] \end{array} \right.$$

↳ probabilistic method.

Weight vectors:

$$w^* = \underset{w}{\operatorname{argmin}} \sum_{i=1}^n \log(1 + \exp(-y_i w^T x_i))$$

w is called weight vector.

$w = \langle w_1, w_2, w_3, \dots, w_d \rangle$

\downarrow 2 features

$E[12]$

$$w = \langle w_1, w_2, w_3, \dots, w_d \rangle$$

$$f_1, f_2, f_3, \dots, f_d$$

for every feature we have 'weight' associated with it

decision: $y_q \rightarrow y_q$ if " $w^T x_q$ " > 0 then $y_q = +1$
 if " $w^T x_q$ " < 0 then $y_q = -1$

Probabilistic:

$$\delta(w^T x_q) = p(y_q = +1)$$

interpretation of w :

(1) if w_i is true, $x_{q,i} \uparrow \Rightarrow (w_i x_{q,i}) \uparrow$
 f_i
 $\Rightarrow \sum_{i=1}^d w_i x_{q,i} \uparrow$
 $\delta(w^T x_q) \uparrow$ then

interpretation of ...

Case 1:

if w_i is +ve : if $\sum w_i y_i \uparrow$

f_i

$$= (w_i y_i) \uparrow$$

$$= \left(\sum_{i=1}^d w_i y_i \right) \uparrow$$

$$= \delta(w^T x_q) \uparrow$$

$$= P(y_{q1}=1) \uparrow$$

Case 2:

If w_i is -ve : if $w_i y_i \uparrow$

$\Rightarrow w_i y_i \downarrow$

$$\Rightarrow \sum_{i=1}^d w_i y_i \downarrow$$

$$\Rightarrow \delta(w^T x_q) \downarrow$$

$$\Rightarrow P(y_{q1}=1) \downarrow$$

$$\Rightarrow P(y_{q1}=-1) \uparrow$$