



CSCI 1102 Computer Science 2

Meeting 9: Thursday 2/25/2021
Memory Organization

Mostly Random Access
Memory (RAM)

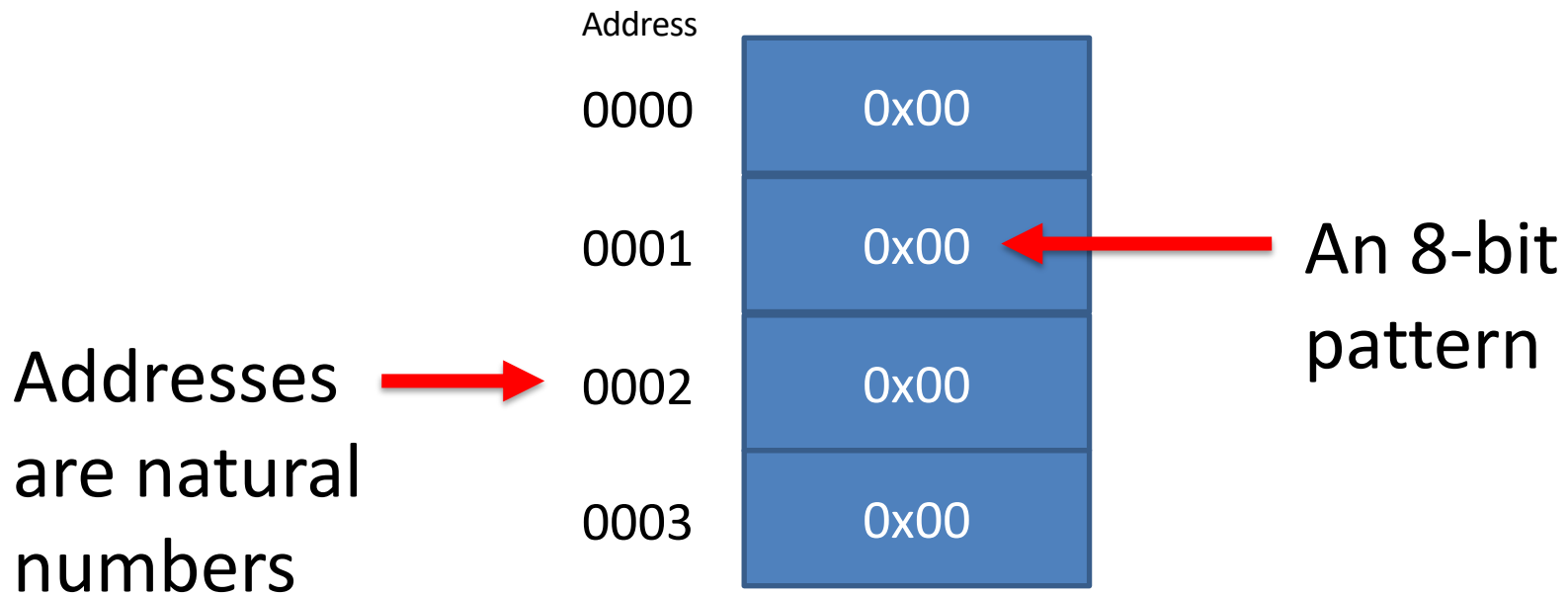
Flash Drive



Ephemeral
Memory

Persistent
Storage

RAM: Contiguously allocated bytes – each byte has a numerical address



32-bit **words** – 4 consecutive bytes

Address

0x0000

0x 03 02 01 00

0x0004

0x 07 06 05 04

0x0008

0x 0B 0A 09 08

0x000C

0x 0F 0E 0D 0C

64-bit **words** – 8 consecutive bytes

Address

0x0000

0x 07 06 05 04 03 02 01 00

0x0008

0x 07 06 05 04 03 02 01 00

0x0010

0x 07 06 05 04 03 02 01 00

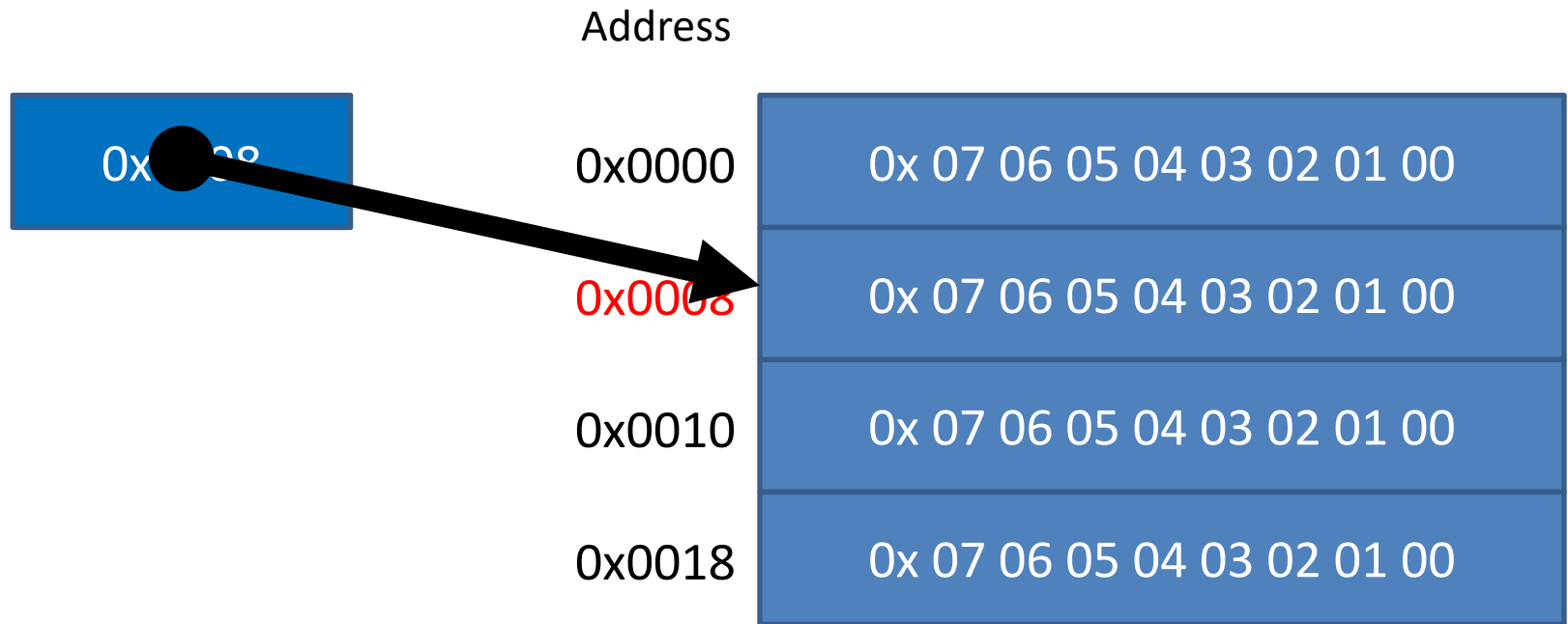
0x0018

0x 07 06 05 04 03 02 01 00

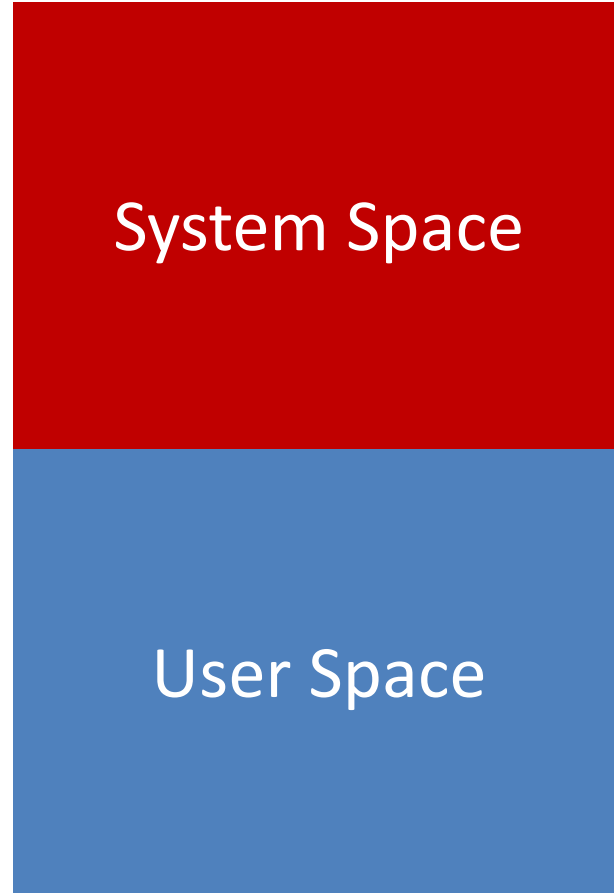
Variables often hold Addresses

0x0008	Address	
	0x0000	0x 07 06 05 04 03 02 01 00
	0x0008	0x 07 06 05 04 03 02 01 00
	0x0010	0x 07 06 05 04 03 02 01 00
	0x0018	0x 07 06 05 04 03 02 01 00

In diagrams, addresses are usually depicted abstractly as **arrows**, often called **pointers**



Ephemeral Memory



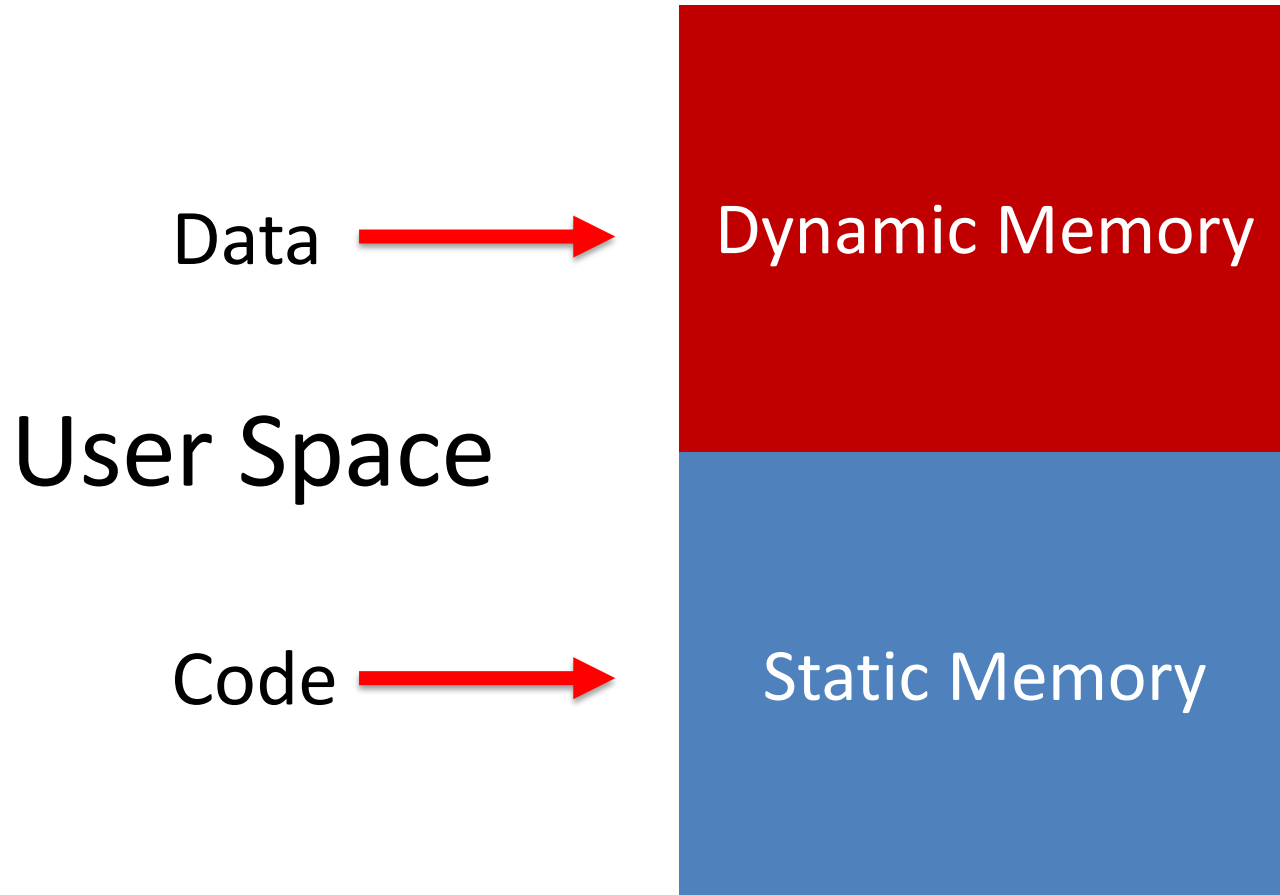
User Space

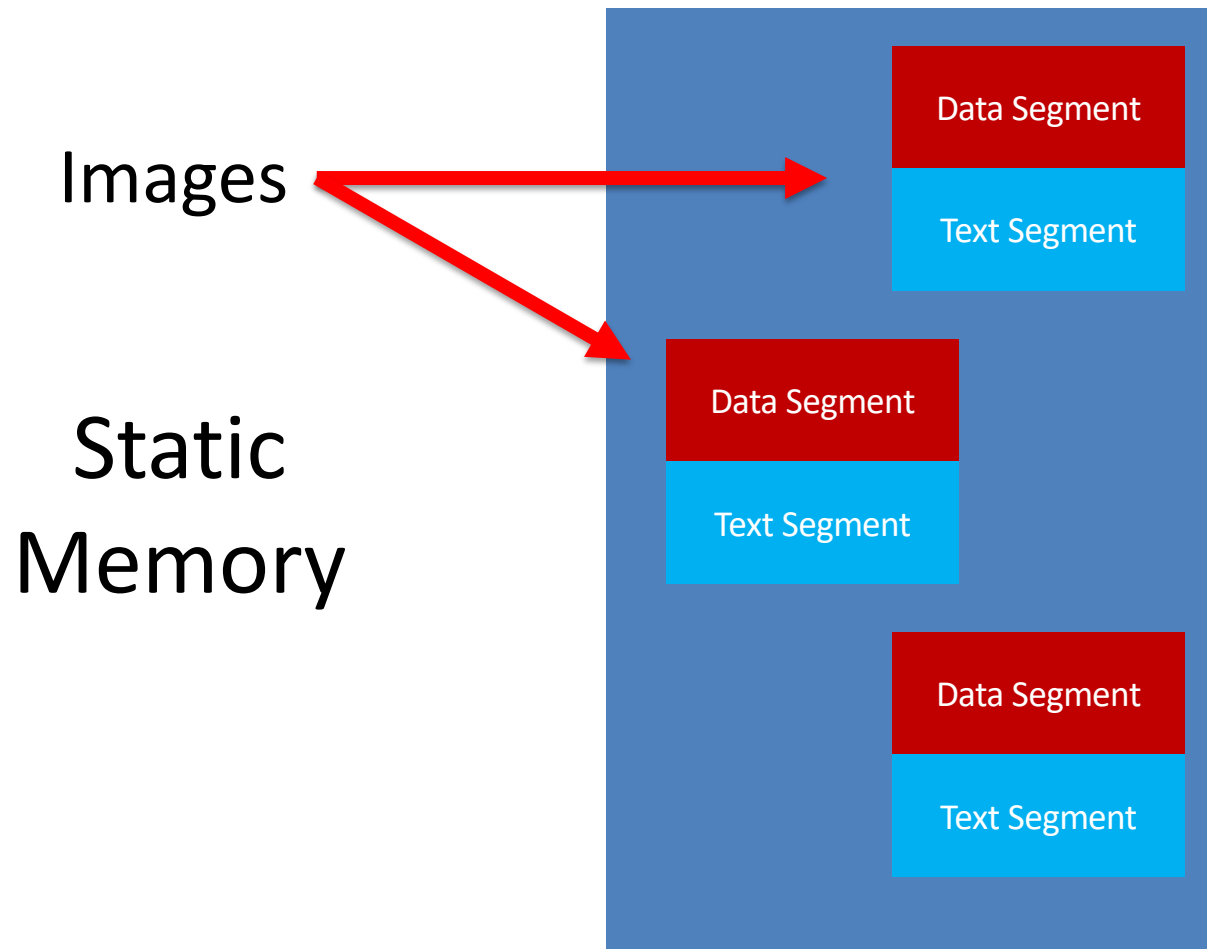


A vertical rectangle representing the User Space is divided into two horizontal sections. The top section is red and labeled 'Dynamic Memory'. The bottom section is blue and labeled 'Static Memory'.

Dynamic Memory

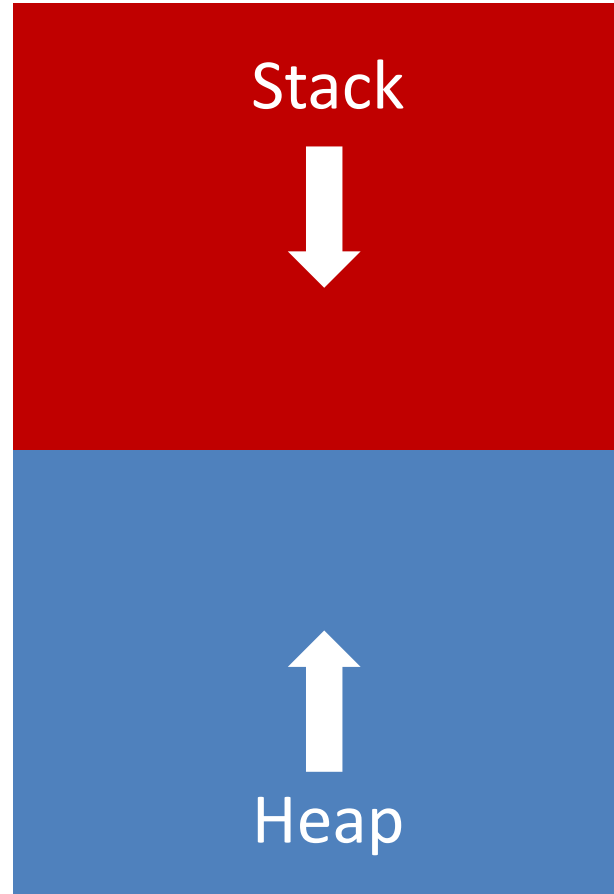
Static Memory



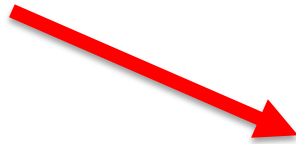


All function/method definitions
can be understood as **images**
residing in static memory

Dynamic Memory



Storage for
function
variables



Stack



Dynamic Memory

Storage for large
values (e.g., arrays)
& long-living values



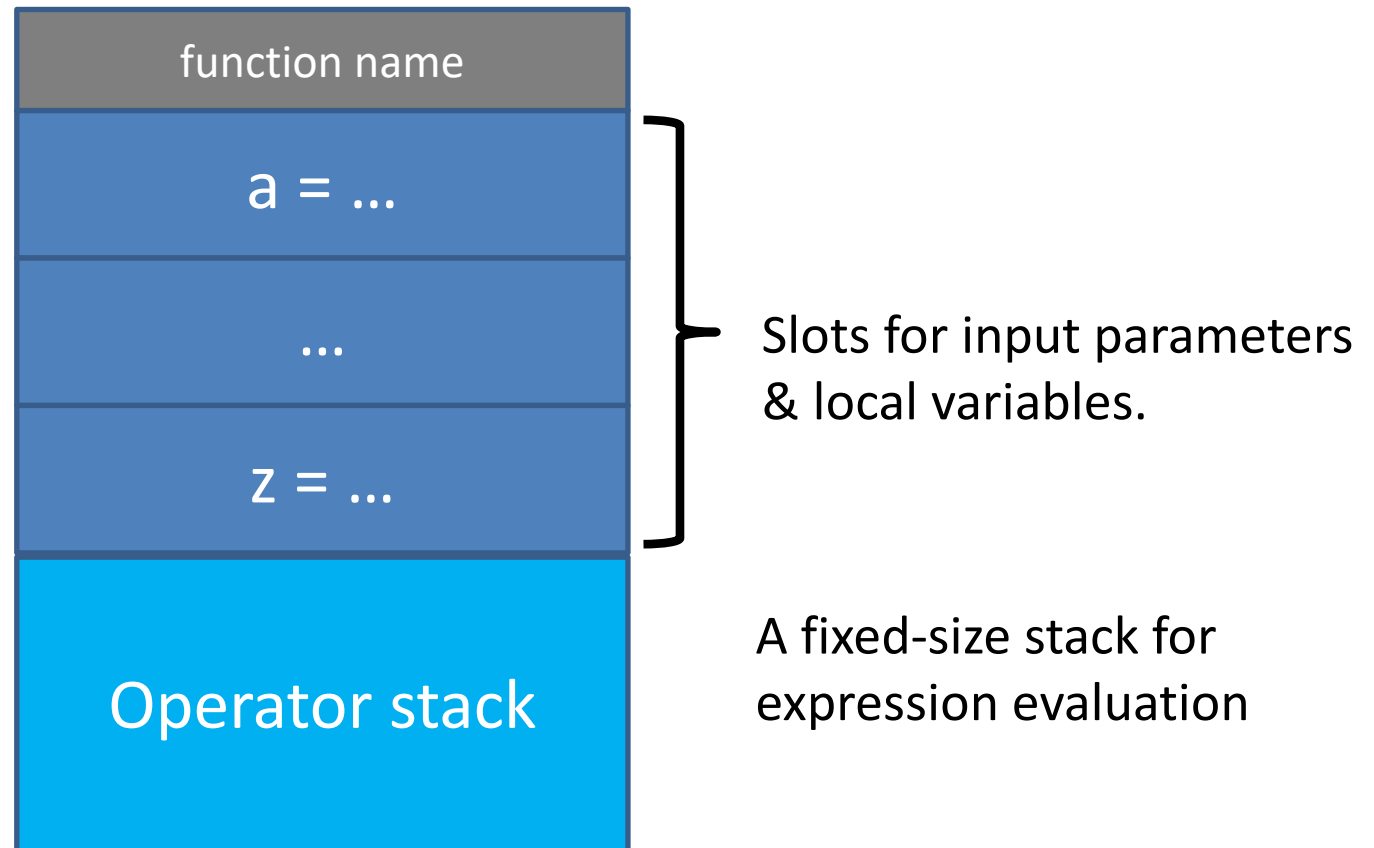
Heap



Managing Dynamic Memory

- The **stack**: ~~the compiler generates code that~~ JVM manages the allocation and deallocation of call frames on the call stack
- A call frame is placed on top of the stack on function call and removed on function return
- The **heap**: managed by a run-time support routine called a **garbage collector**.

The Call Stack: Call Frames/Activation Records



Body of function evaluated with respect to a call frame.

The Size of a Call Frame's Operator Stack is Fixed

$a + (b + (c + (d + e)))$

a b c d e + + + +

e
d
c
b
a

Static Function Example

```
1 public class StatExample {
2
3     public static int g(int k) {
4         return k * 2;
5     }
6
7     public static int f(int j, int k) {
8         return j * g(k);
9     }
10
11     public static void main(String[] args) {
12         int result = f(2, 3);
13         System.out.format("The answer is %d.\n", result);
14     }
15 }
16 }
```

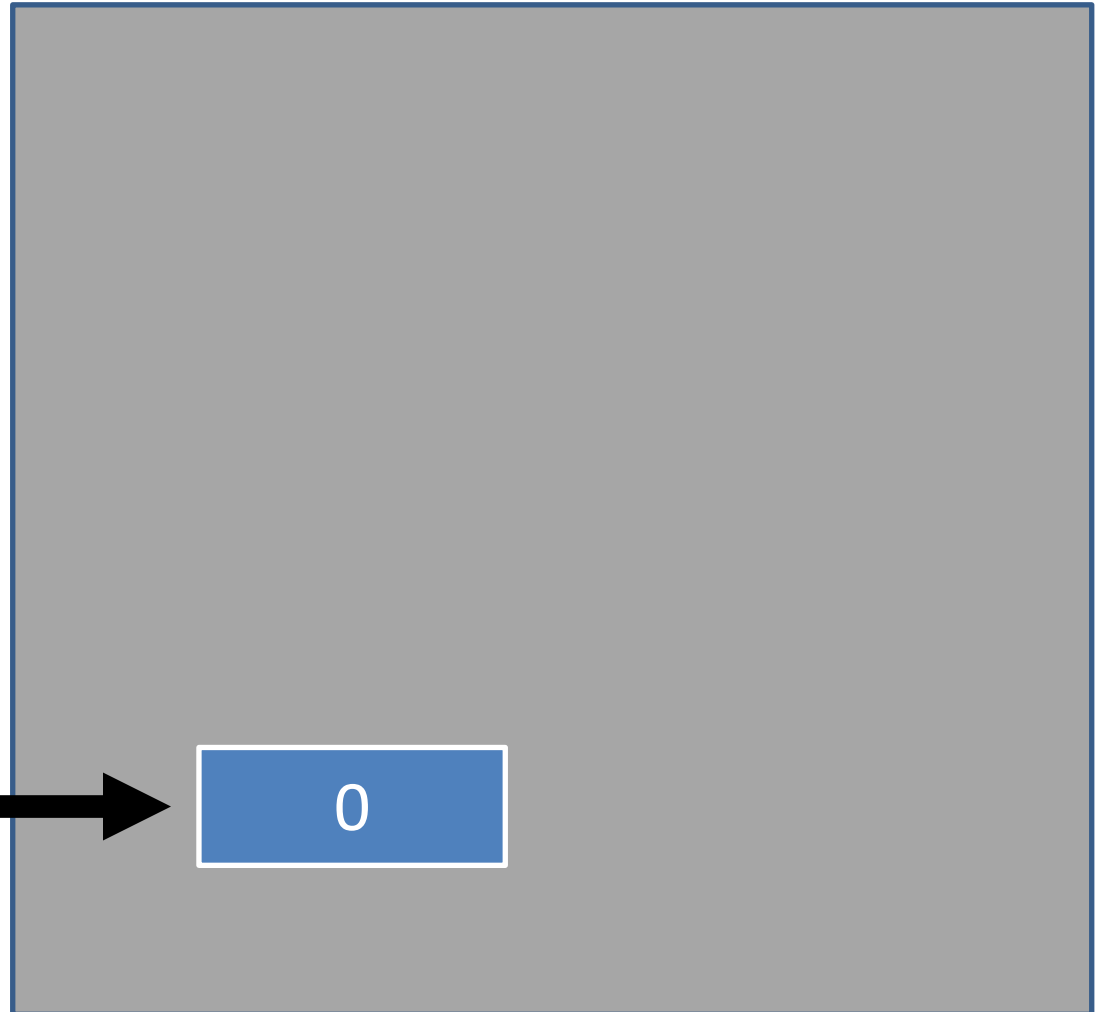
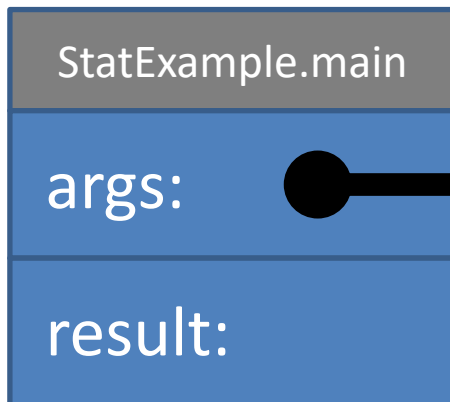
Stack

Heap

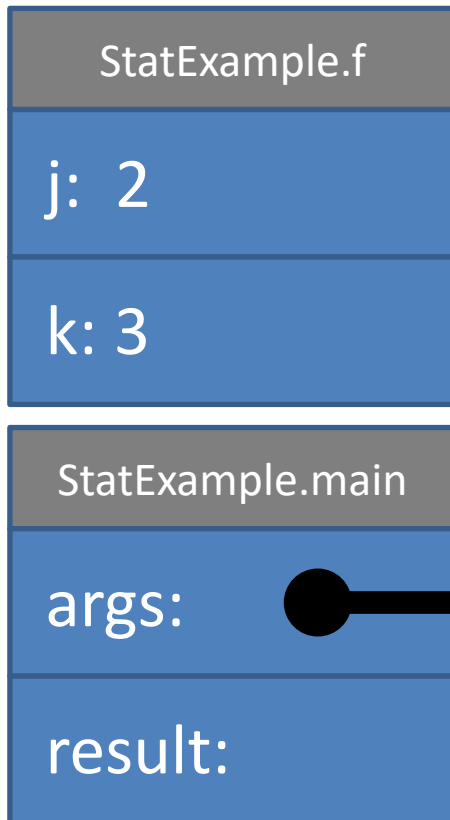


Stack

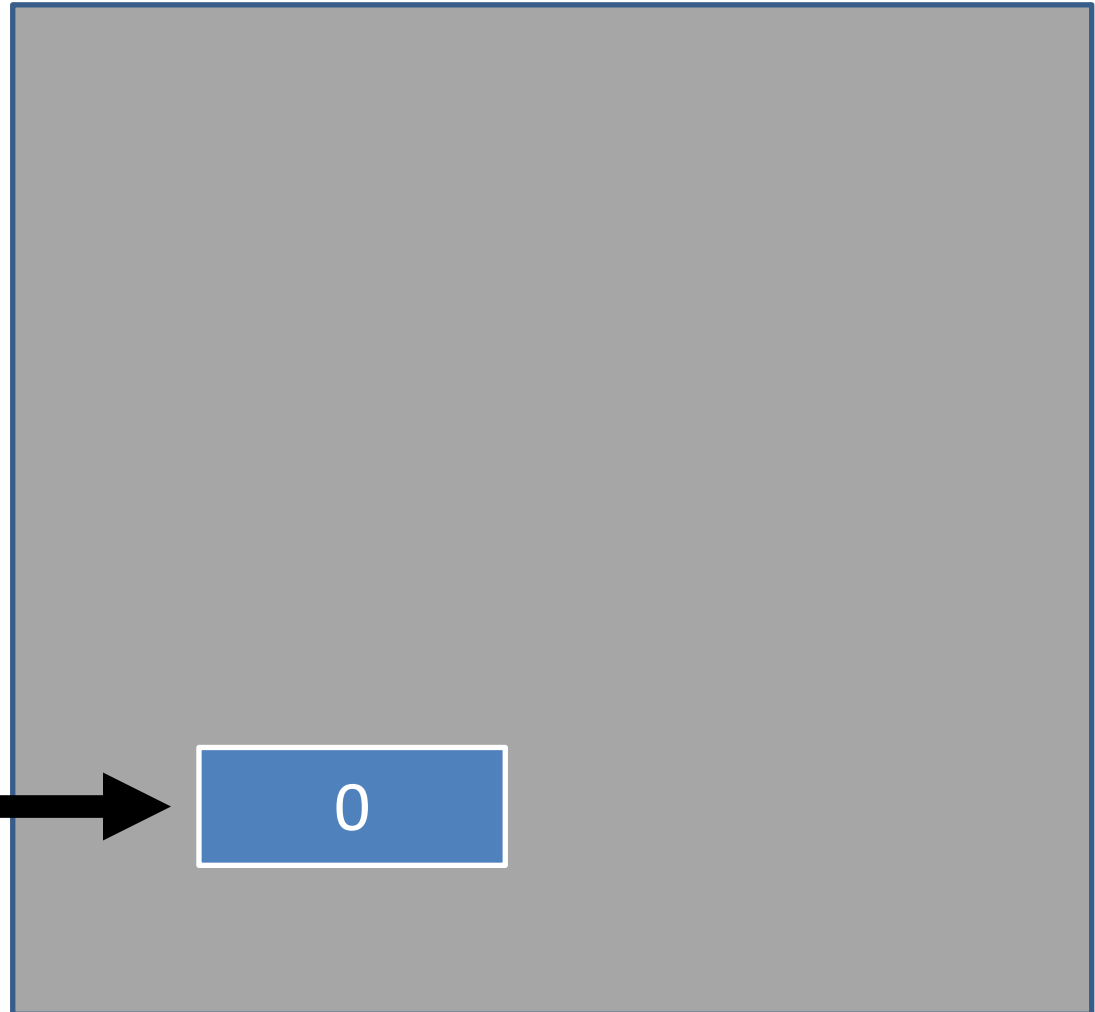
Heap



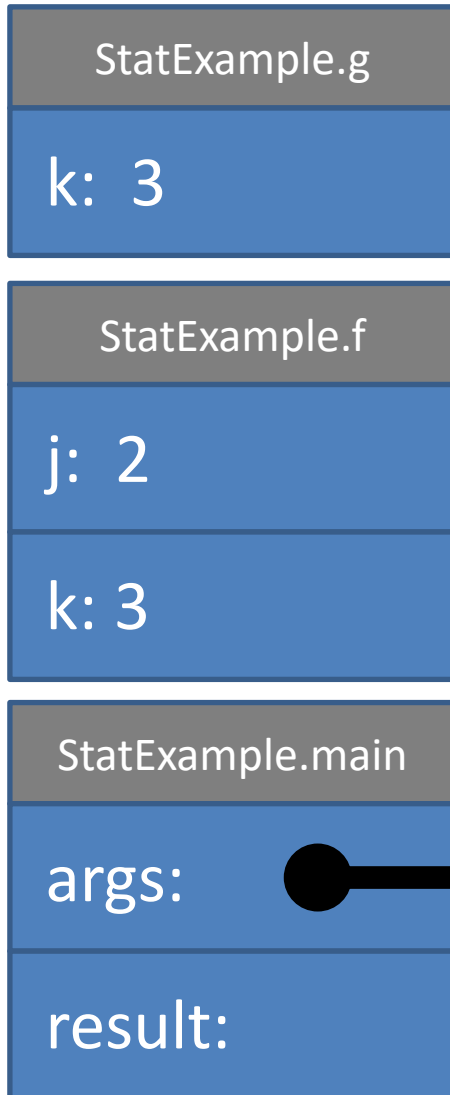
Stack



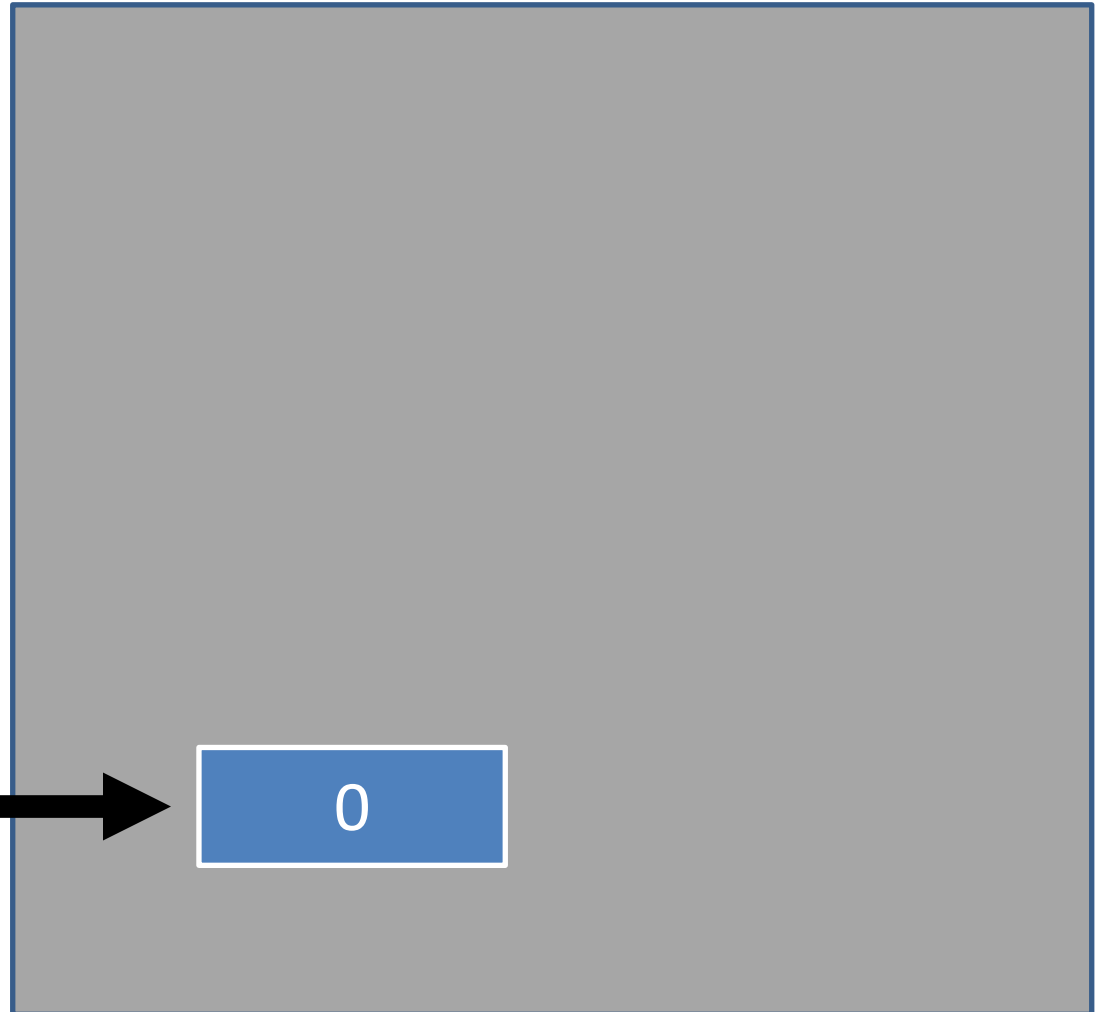
Heap



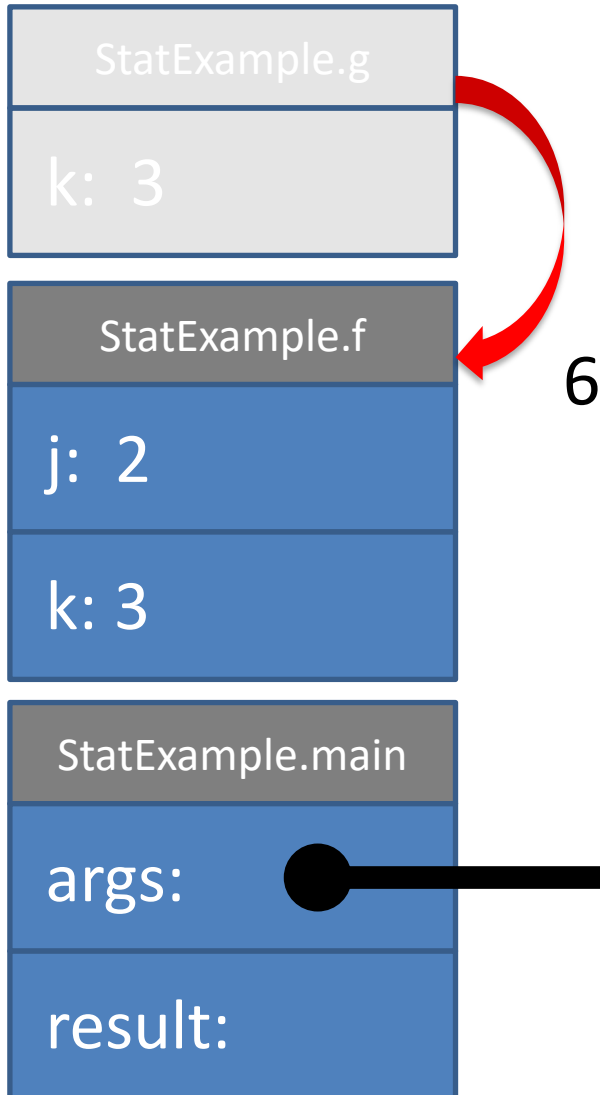
Stack



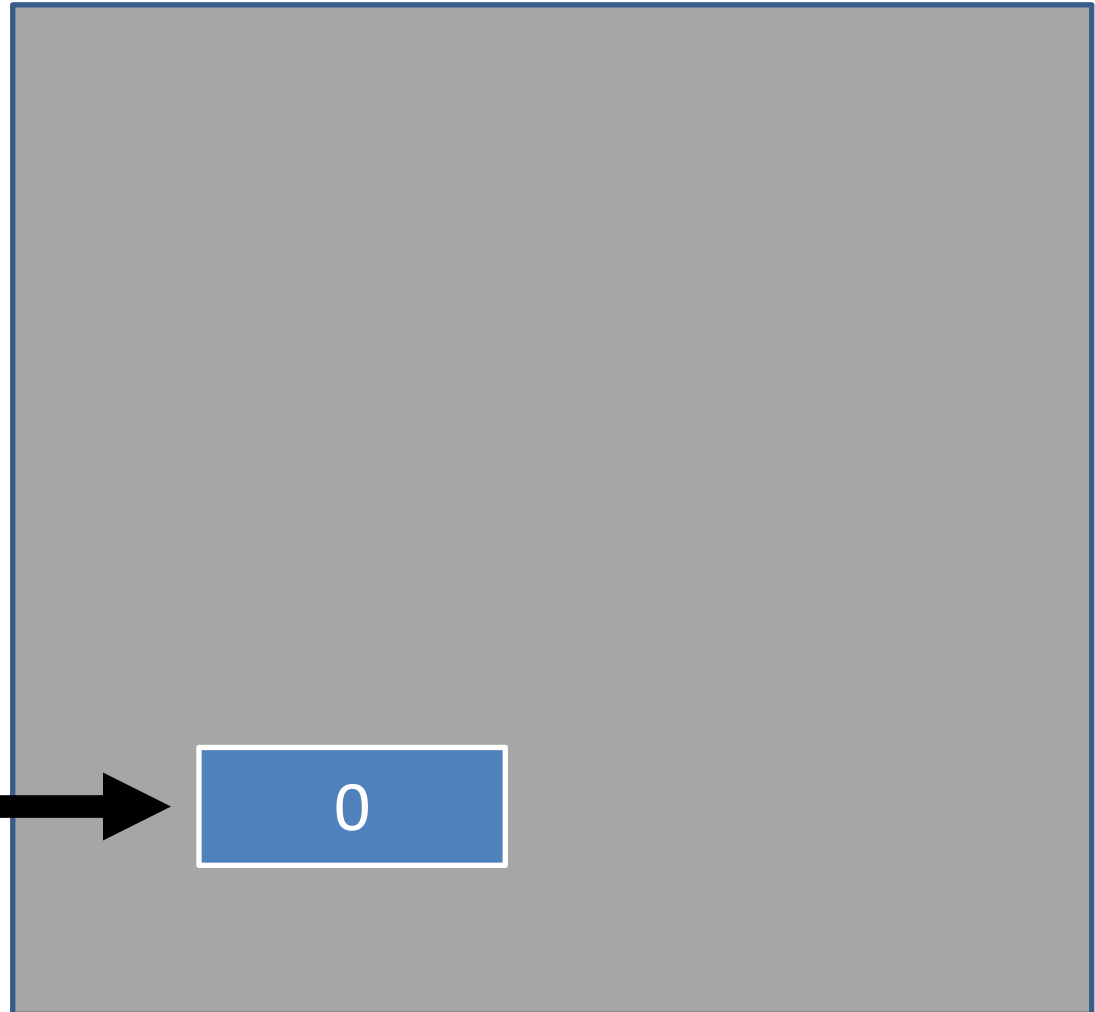
Heap



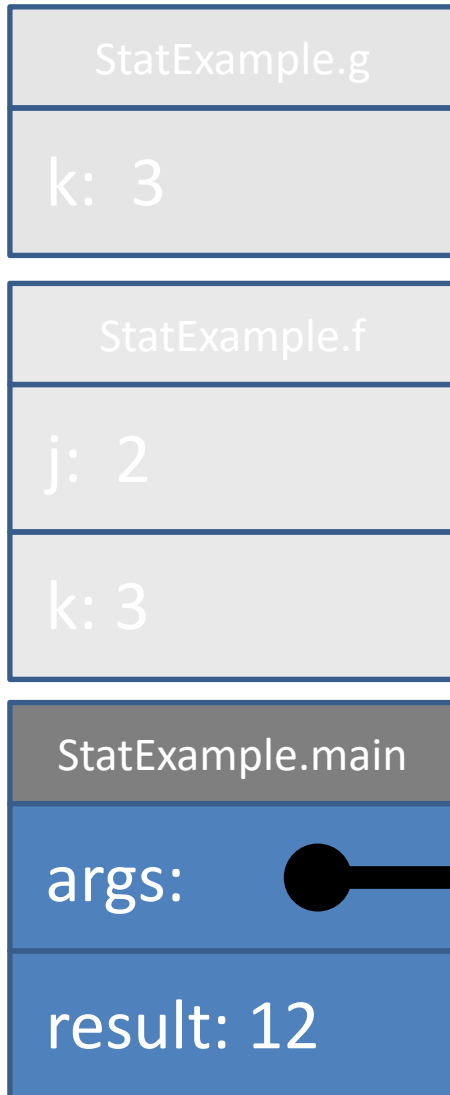
Stack



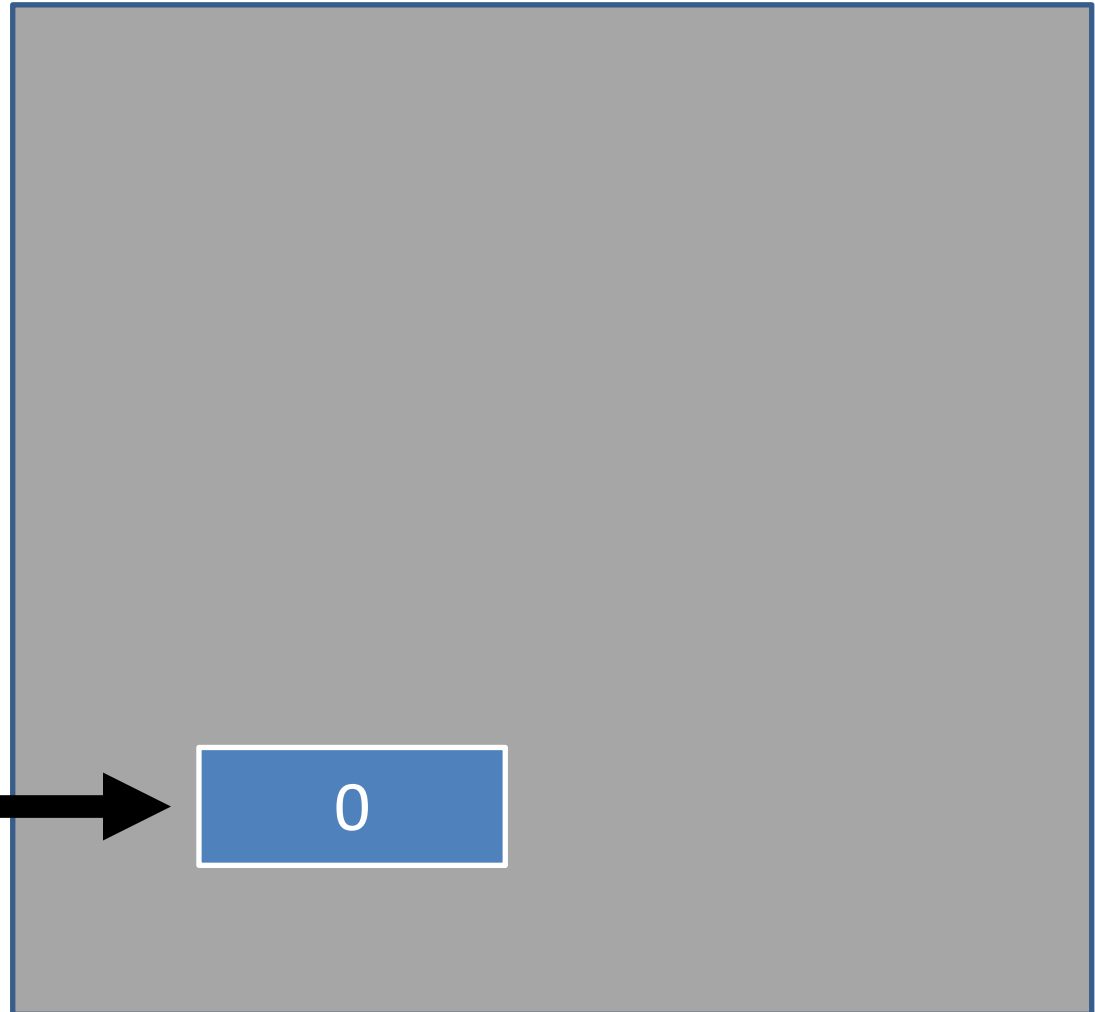
Heap



Stack

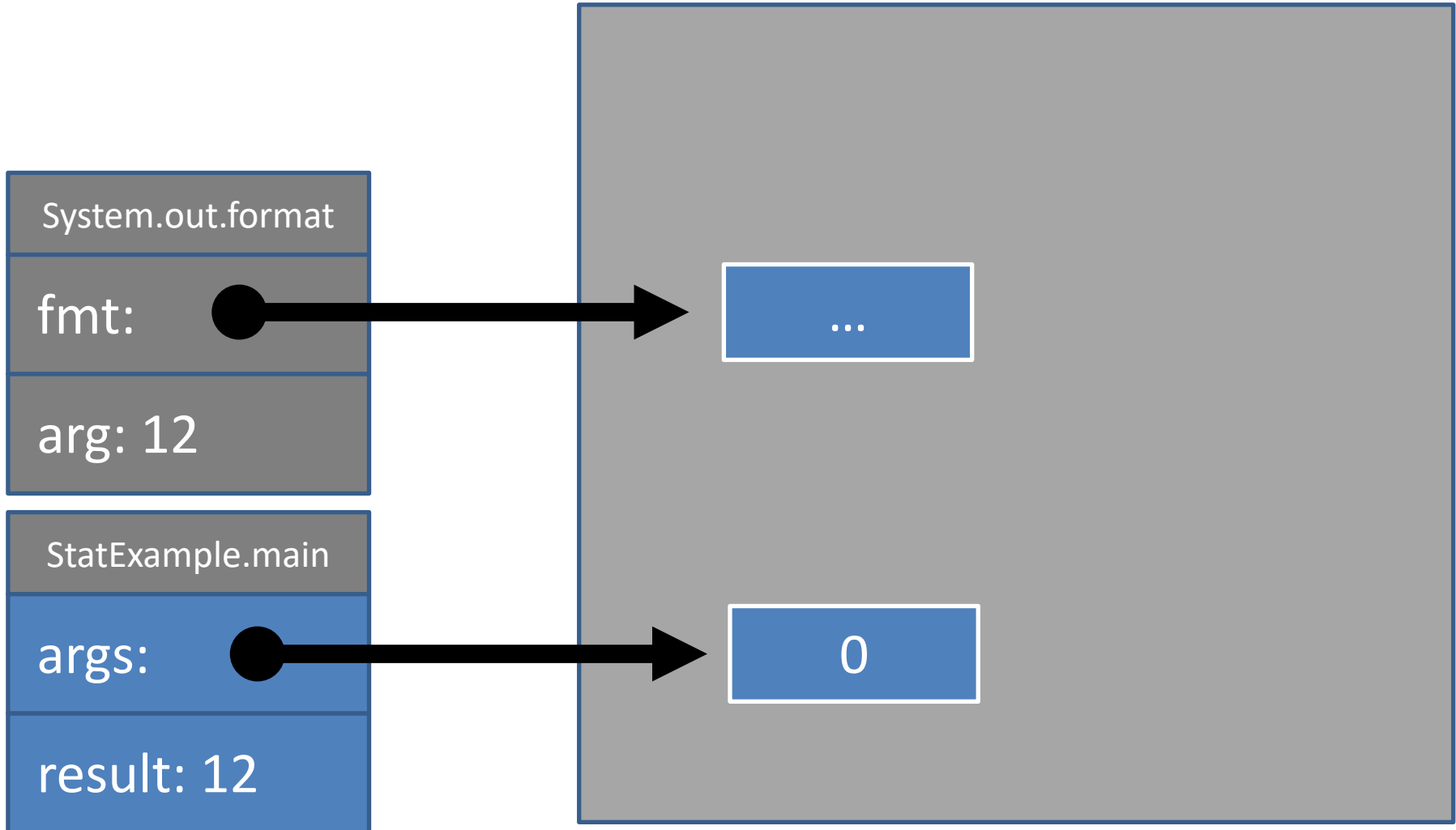


Heap



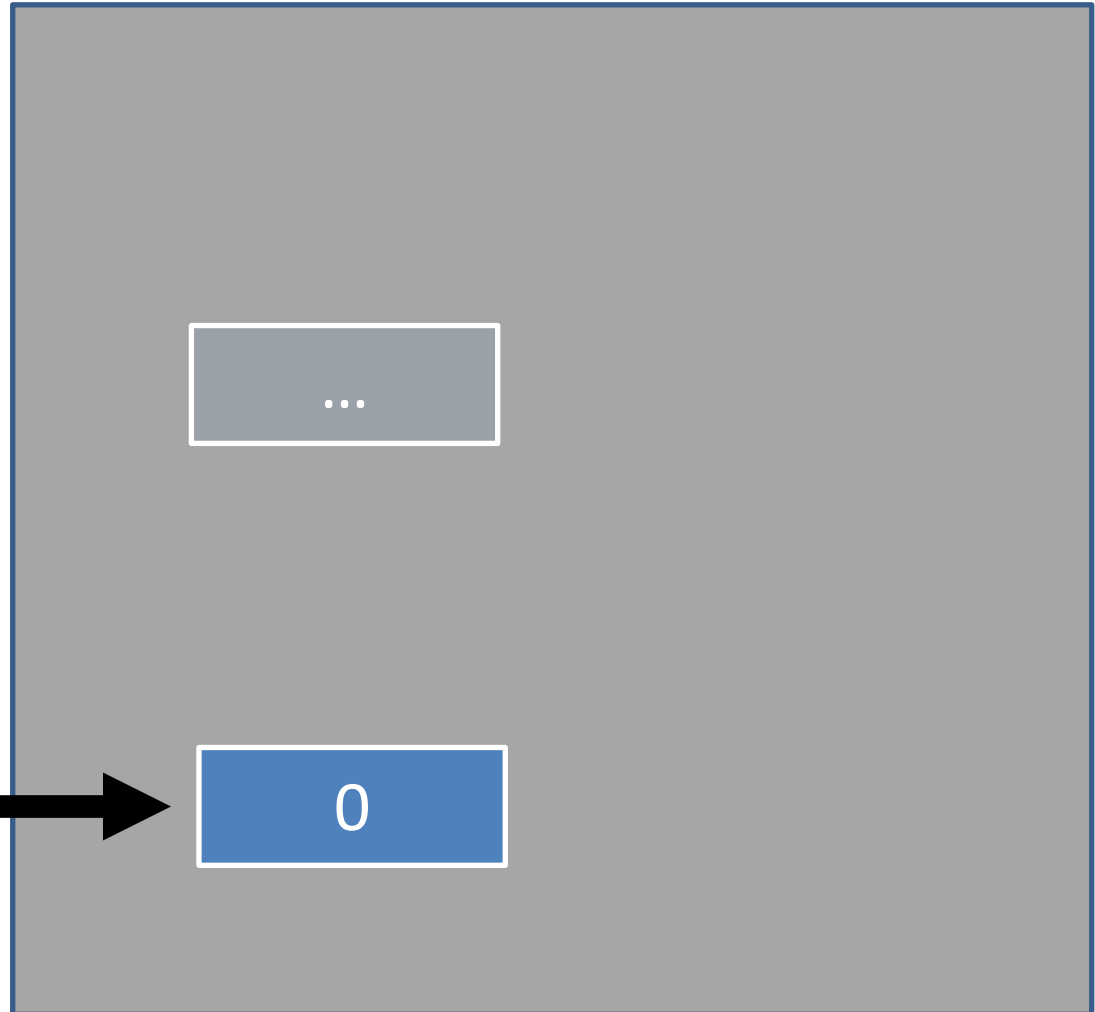
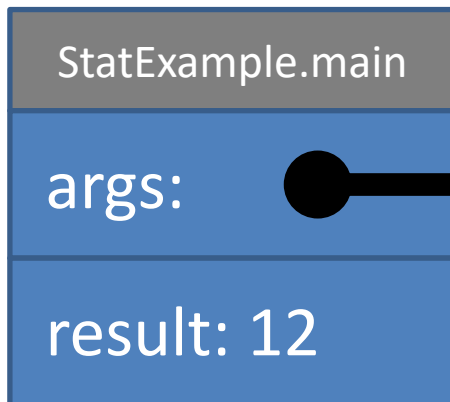
Stack

Heap



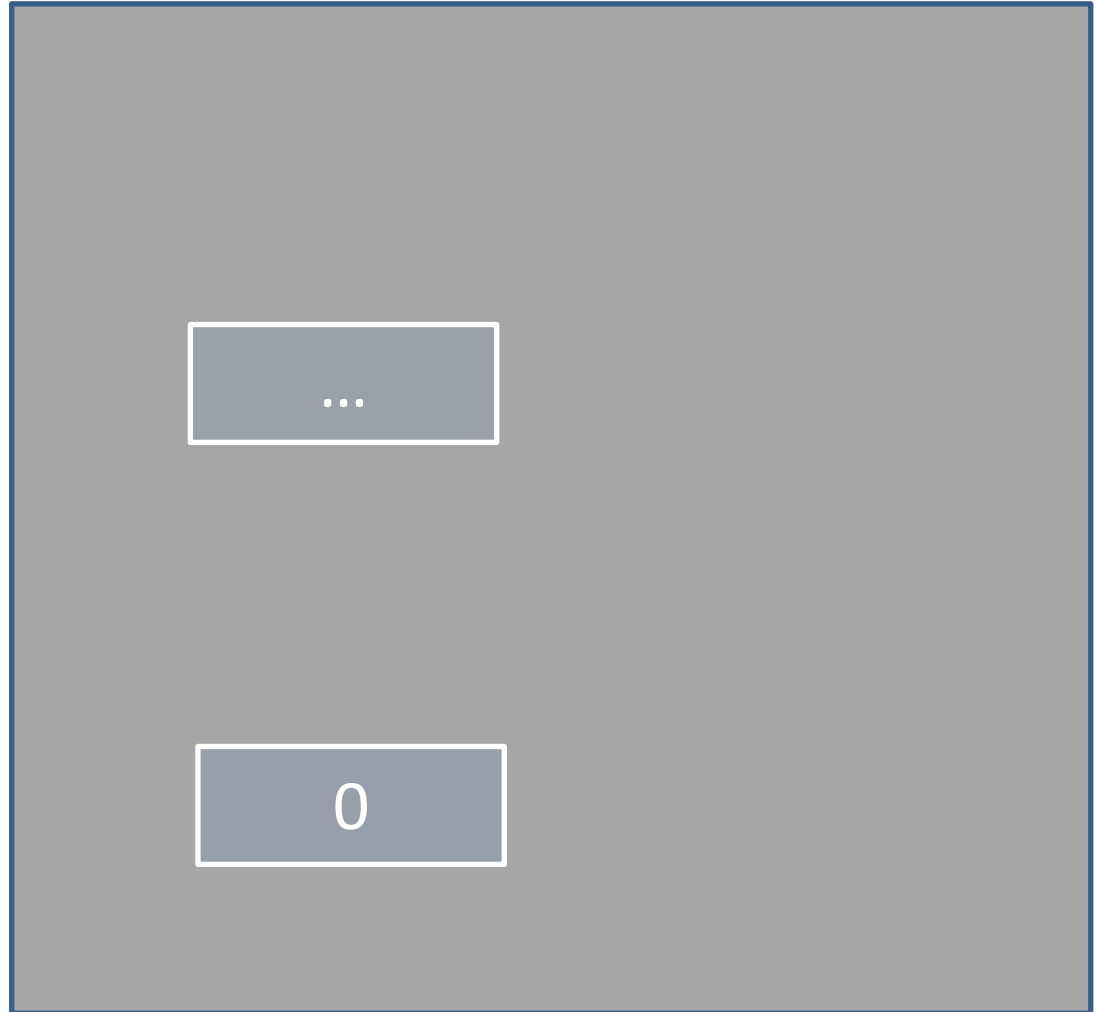
Stack

Heap



Stack

Heap

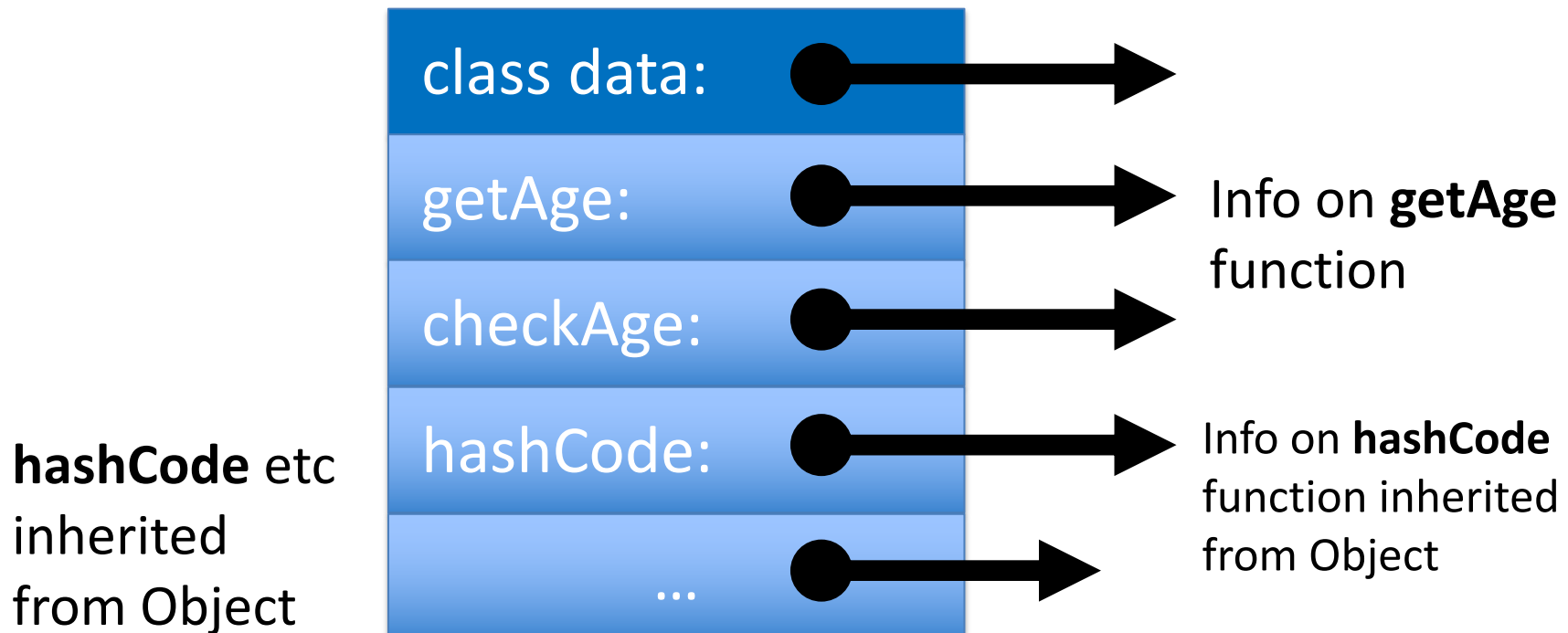


```
1 public class DynExample {
2
3     private int age;
4
5     public DynExample(int age) { this.age = age; }
6
7     private int getAge() { return this.age; }
8
9     private void checkAge() {
10         if (this.age < 0 || this.age > 120)
11             throw new RuntimeException("Bad age");
12     }
13
14     public static void main(String[] args) {
15         DynExample de = new DynExample(21);
16         de.checkAge();
17         System.out.format("Age is %d.\n", de.getAge());
18     }
19 }
20 }
```

Dynamic Function Example

Every Class has a *Dispatch Table*

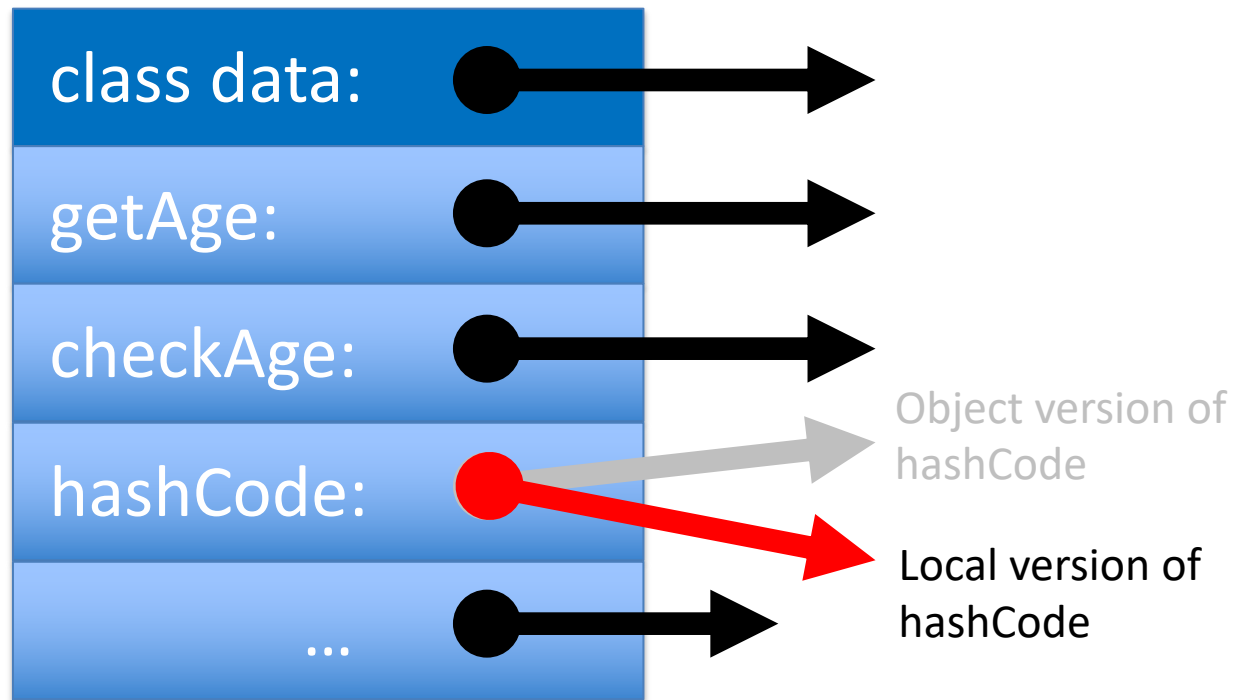
```
18 public static void main(String[] args) {  
19  
20     DynExample de = new DynExample(21);  
    }
```



Dispatch Tables Support Inheritance and Override

```
18 public static void main(String[] args) {  
19  
20     DynExample de = new DynExample(21);  
    }
```

Override e.g.,
hashCode?
Replace the
arrow in the
table

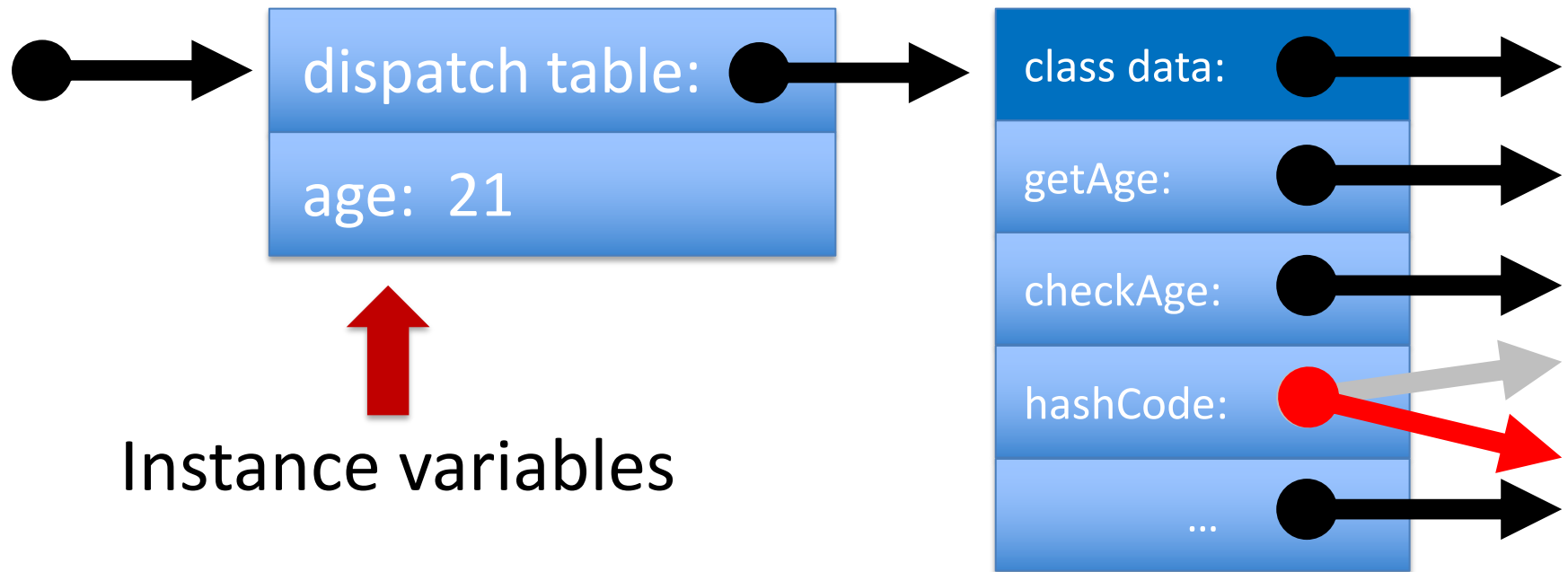


Dispatch Tables

- All objects created with **new** can share the same dispatch table
- Dispatch tables are *fixed*, they can be stored in static memory or in the heap.

Object Representation in the Heap

```
18 public static void main(String[] args) {  
19  
20     DynExample de = new DynExample(21);  
    }
```



Message-Passing Style

- Dynamic functions use message-passing style;

`move(point, dx, dy)`

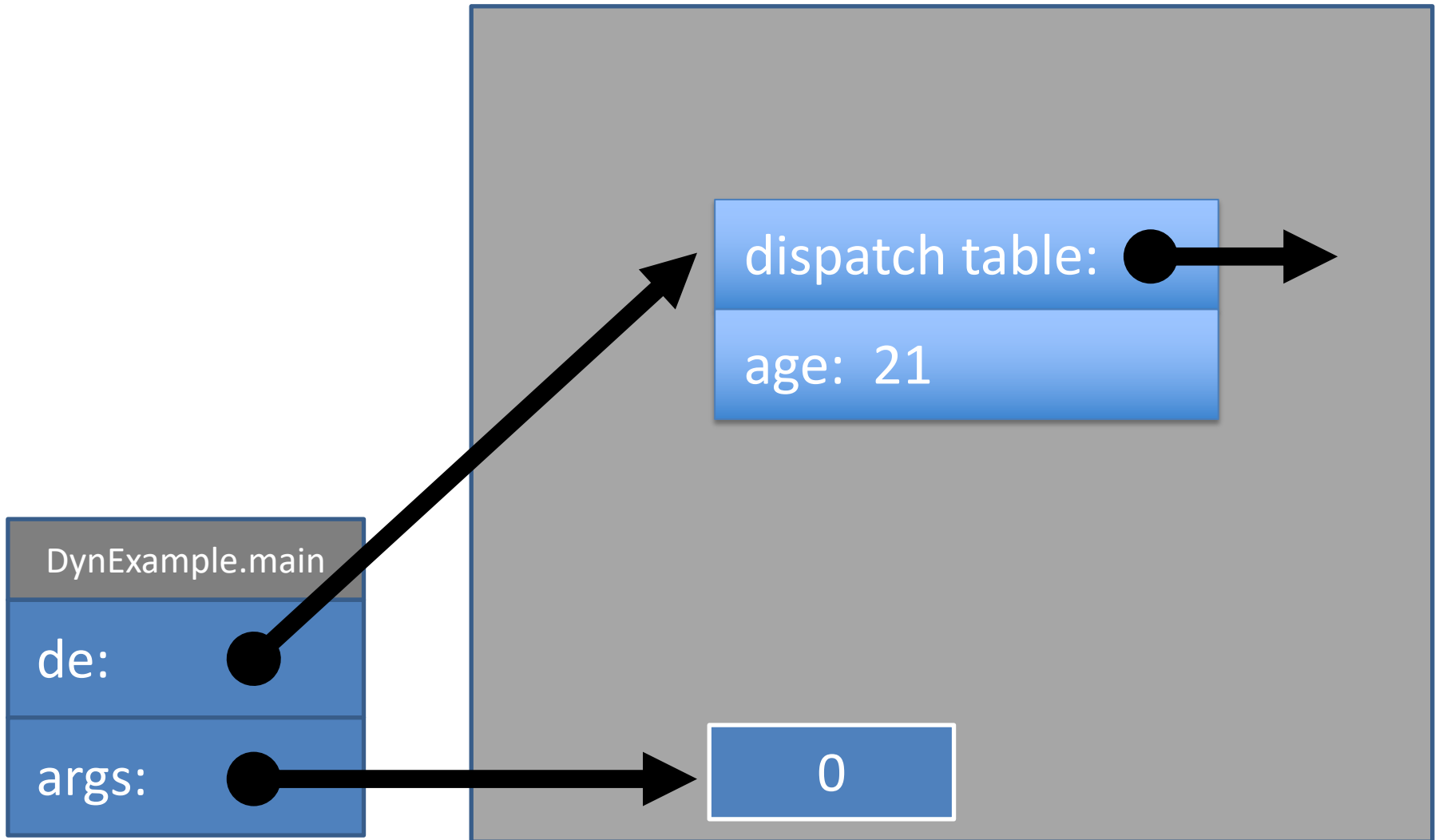
`point.move(dx, dy)`

- Message-passing style is orthogonal to inheritance

Message-Passing style Plumbing

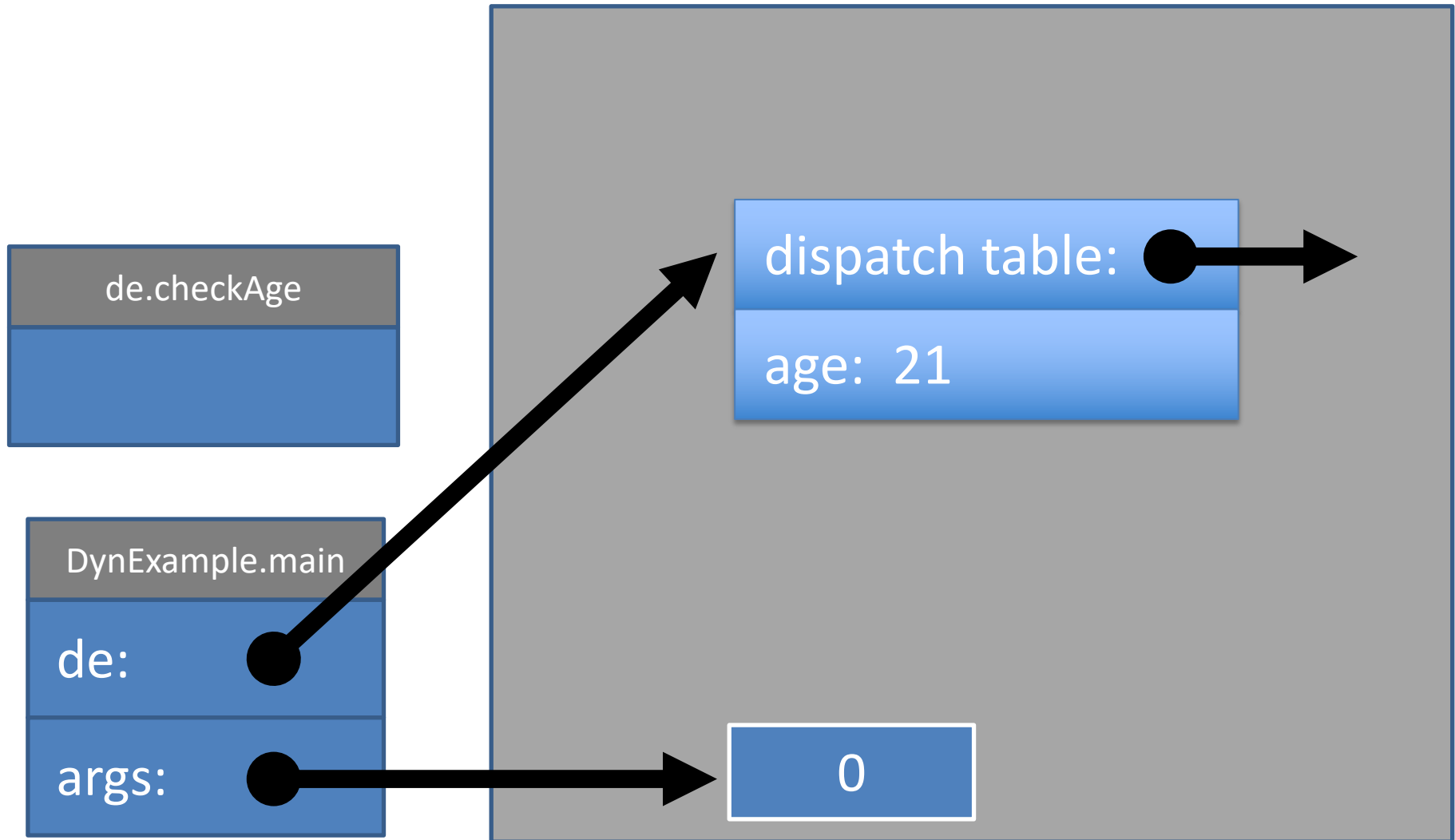
Stack

Heap



Stack

Heap



How can `de.checkAge` access the
instance variable `age`?

Replace all calls of dynamic functions

de.checkAge()

de.checkAge(de)

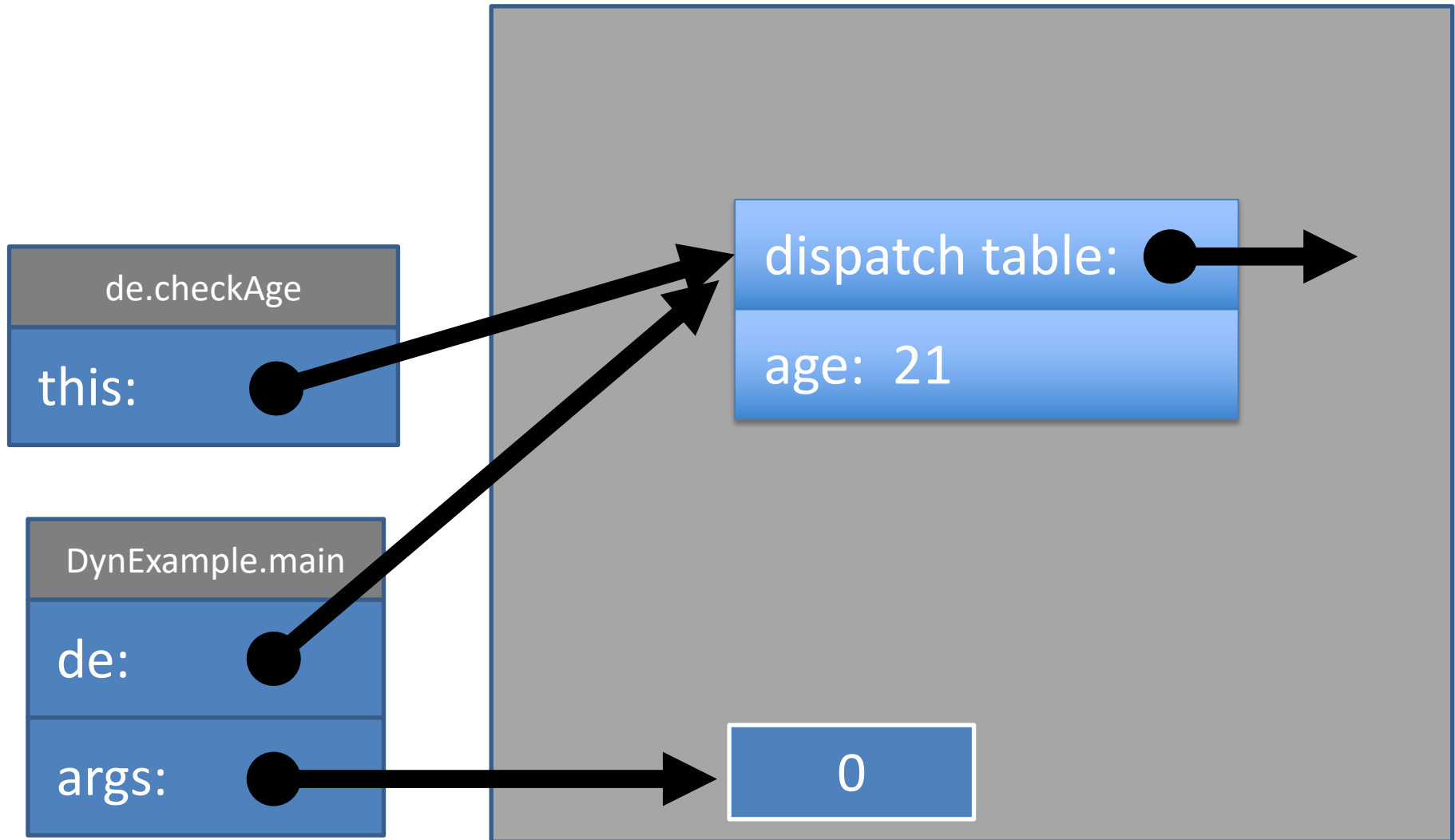
Replace all defs of dynamic functions

```
void checkAge() { ... }
```

```
void checkAge(DynExample this) { ... }
```

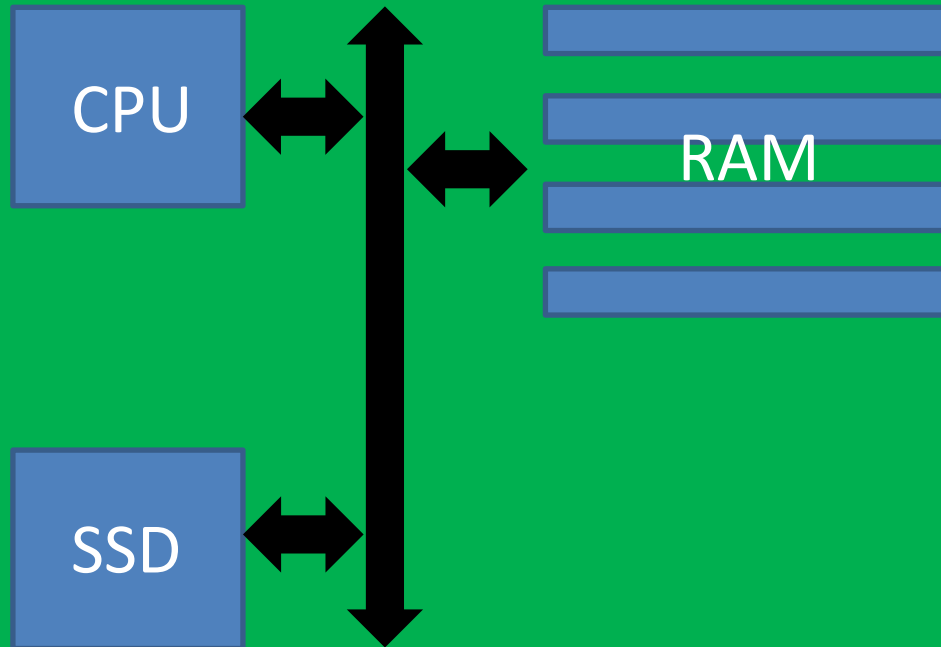
Stack

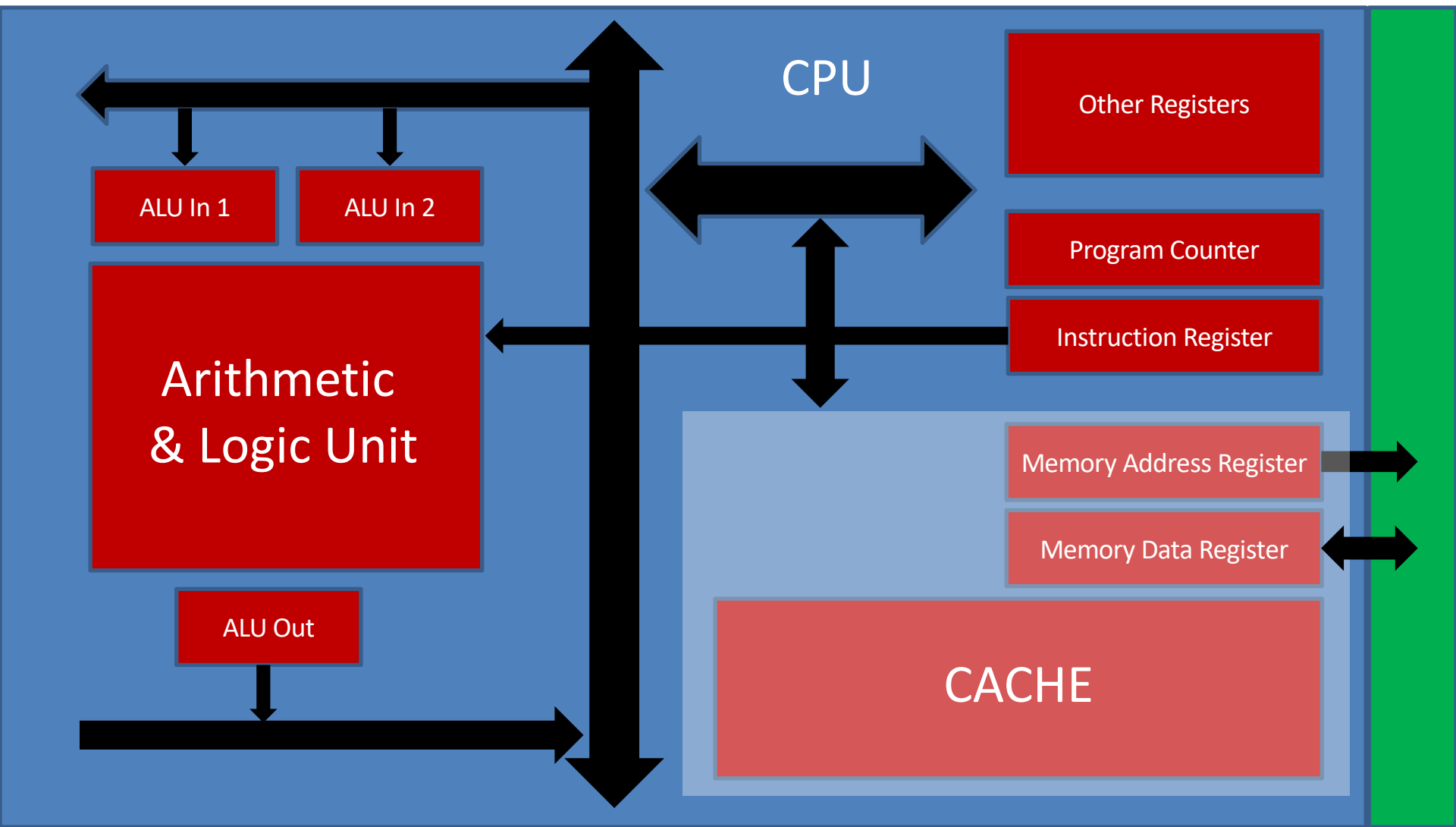
Heap



Locality

Motherboard





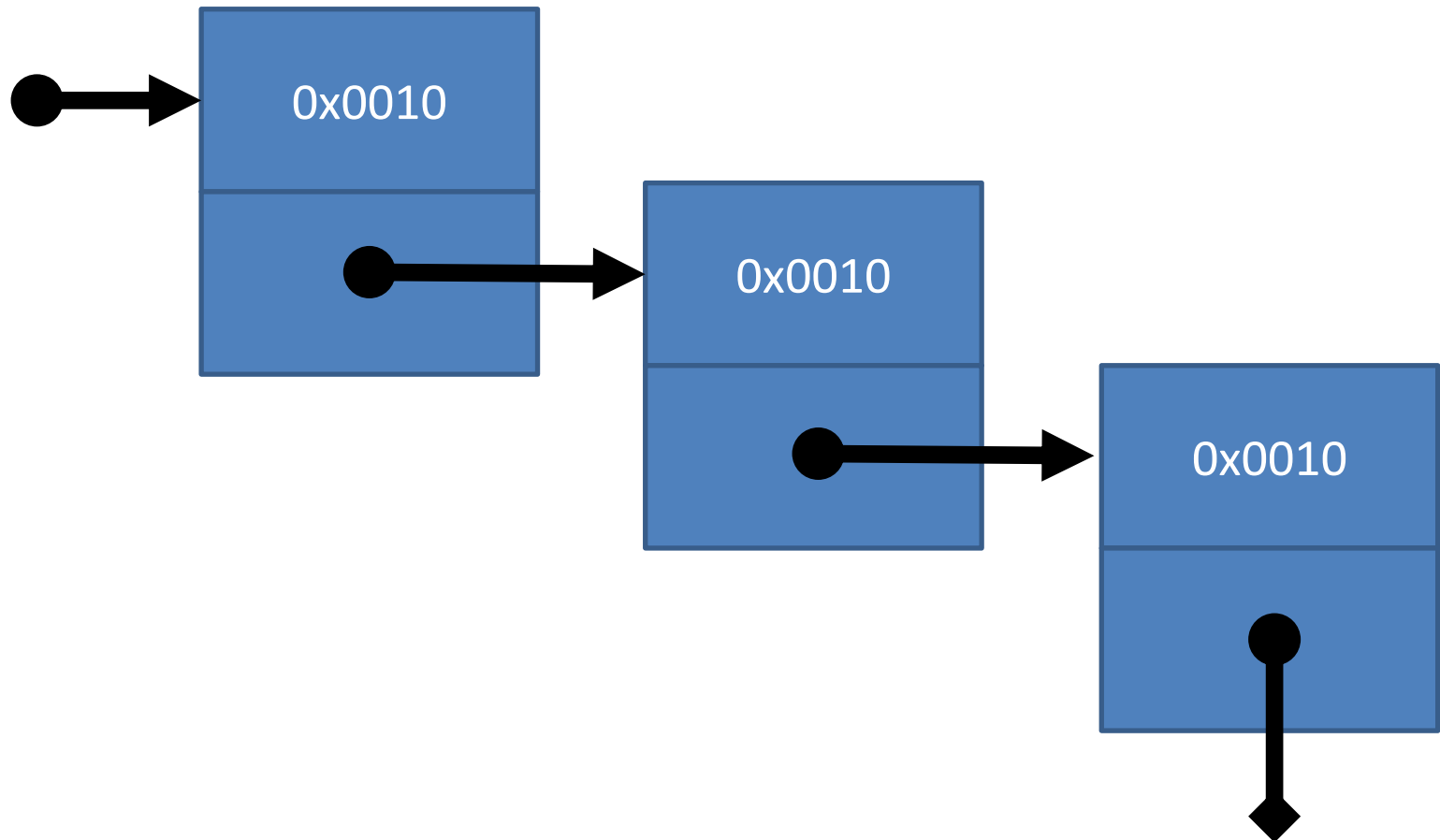
Latency numbers every programmer should know

L1 cache reference	0.5 ns	
Branch mispredict	5 ns	
L2 cache reference	7 ns	
Mutex lock/unlock	25 ns	
Main memory reference	100 ns	
Compress 1K bytes with Zippy	3,000 ns	= 3 μ s
Send 2K bytes over 1 Gbps network	20,000 ns	= 20 μ s
SSD random read	150,000 ns	= 150 μ s
Read 1 MB sequentially from memory	250,000 ns	= 250 μ s
Round trip within same datacenter	500,000 ns	= 0.5 ms
Read 1 MB sequentially from SSD*	1,000,000 ns	= 1 ms
Disk seek	10,000,000 ns	= 10 ms
Read 1 MB sequentially from disk	20,000,000 ns	= 20 ms
Send packet CA->Netherlands->CA	150,000,000 ns	= 150 ms

Linked Data Structures

- Pro: unlimited size with no linear resizing cost;
- Con: Poor locality

Block Storage & Linked Storage



Block Storage & Linked Storage

