# $\begin{array}{c} {\rm Second~Exam} \\ {\rm CS~1102~Computer~Science~2} \end{array}$

## Fall 2018

Tuesday November 20, 2018 Instructor Muller

## KEY

Before reading further, please arrange to have an empty seat on either side of you. Now that you are seated, please write your name **on the back** of this exam.

This is a closed-notes and closed-book exam. Computers, calculators, and books are prohibited.

- Partial credit will be given so be sure to show your work.
- Feel free to write helper functions if you need them.
- Please write neatly.

Problem	Points	Out Of
1		5
2		2
3		3
4		3
5		3
6		4
Total		20

# Part 1: (5 Points) Binary Trees

The problems in this section relate to binary trees of the form

```
class Node {
  int key;
  Node left, right;
  public Node(int key, Node left, Node right) {
    this.key = key;
    this.left = left;
    this.right = right;
  }
}
```

A binary tree is either empty (in Java, this is normally represented by null) or it is a Node with a key field and left and right fields.

1. (1 Point) Write a function boolean isLeaf(Node t) such that a call isLeaf(t) returns true if t is a leaf. Otherwise isLeaf should return false.

## Answer:

```
public boolean isLeaf(Node t) {
  return (t != null) && (t.left == null) && (t.right == null);
}
```

2. (2 Points) Write a function int height(Node t) such that the call height(t) returns the integer height of t.

```
public int height(Node t) {
  if (t == null || isLeaf(t)) return 0;
  else return 1 + Math.max(height(t.left), height(t.right));
}
```

3. (2 Points) Write a function boolean isBST(Node t) such that a call isBST(t) returns true if t is a binary search tree. Otherwise isBST should return false.

```
public boolean isBST(Node t) {
   if (t == null) return true;
   boolean tIsOK = allLess(t.left, t.key) && allGreater(t.right, t.key);
   return tIsOK && isBST(t.left) && isBST(t.right);
}

private boolean allLess(Node t, int key) {
   if (t == null) return true;
   return (t.key < key) && allLess(t.left, key) && allLess(t.right, key);
}

// Efficient solution due to Bob Dondero. (From SW BST.java)

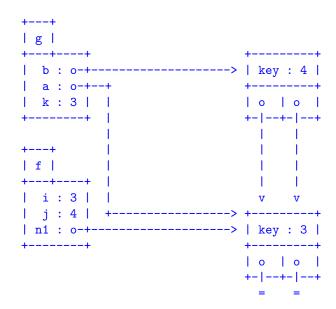
//
boolean isBST(Node t) { return isBST(t, null, null); }

boolean isBST(Node t, Key min, Key max) {
   if (t == null) return true;
   if (min != null && t.key.compareTo(min) <= 0) return false;
   if (max != null && t.key.compareTo(max) >= 0) return false;
   return isBST(t.left, min, t.key) && isBST(t.right, t.key, max);
}
```

## Part 2: (2 Points) Storage Diagrams

Using the class Node from the previous problem, consider the following code and show the state of the stack and heap after (1) has executed but before (2) has executed (i.e., before g returns). Execution is initiated with the call f(3, 4).

stack heap



# Part 3: (3 Points) Sets and Relations

For the questions below let  $A = \{a, b, c\}, B = \{\spadesuit, \clubsuit, \heartsuit\}$  and let

- $R_0 = \{\}$
- $R_1 = \{(a, a), (b, b), (c, c)\}$
- $R_2 = R_1 \cup \{(a,b)\}$
- $R_3 = \{(a, \clubsuit), (b, \clubsuit)\}$
- $R_4 = A \times A = \{(a, a), (b, b), (c, c), (a, b), (a, c), (b, a), (b, c), (c, a), (c, b)\}$
- $R_5 = \{(a, \clubsuit), (b, \heartsuit), (c, \clubsuit)\}$
- $\bullet \ R_6 = A \times B = \{(a, \clubsuit), (a, \clubsuit), (a, \heartsuit), (b, \spadesuit), (b, \clubsuit), (b, \heartsuit), (c, \clubsuit), (c, \clubsuit), (c, \heartsuit)\}$

The following are half-point problems.

1. List the reflexive relations on A. (E.g.,  $R_8$ ,  $R_{12}$  etc).

Answer:  $R_1$ ,  $R_2$  and  $R_4$ .

2. List the transitive relations on A.

Answer:  $R_0$ ,  $R_1$ ,  $R_2$  and  $R_4$ .

3. List the total maps from A to B.

Answer: Only  $R_5$ .

4. List the partial maps from A to B.

Answer:  $R_0$ ,  $R_3$  and  $R_5$ .

5. List the equivalence relations on A.

Answer:  $R_1$  and  $R_4$ .

6. List the partial orders on A.

Answer:  $R_1$  and  $R_2$ .

# Part 4: (3 Points) Hash Tables

- 1. (2 Points) In Java, every value of a reference type (i.e., created with new) comes equipped with an equals function boolean equals (Object other) and a hashCode function int hashCode().
  - (a) True or false: if x.hashCode() == y.hashCode() then it must be the case that x.equals(y).

Answer: False. It is common for unequal keys to hash to the same table location.

(b) True or false: if x.equals(y) then it must be the case that x.hashCode() == y.hashCode().

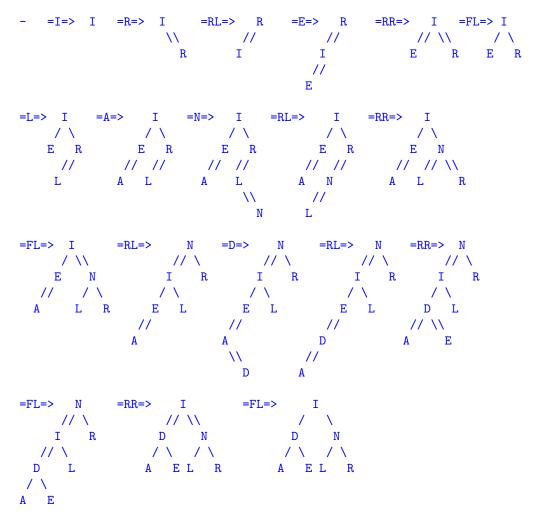
Answer: True. If two keys are equal as judged by the equals equivalence relation, then they should hash to the same table location.

2. (1 Point) In a sentence or two, what are the advantages and disadvantages in resolving hash collisions using chaining?

Answer: Chaining is simple to implement and has very fast insert. If the load factor of the table is large, the chains can get long and the find operation can take time proportional to the chain length. Chains also have poor locality.

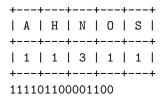
# Part 5: (3 Points) Red/Black Trees

Show all of the successive trees that result from the left-to-right insertion of the letters IRELAND into an empty left-leaning Red/Black tree.



## Part 6: (4 Points) Huffman Coding

A zip file contains the following frequency table and bit sequence.



The file was constructed with the same assumptions as in problem sets 7 and 8:

- 1. Letters are initially entered into the PQ in alphabetical order;
- 2. Ties are broken by placing the newly inserted entry behind all entries with the same priority;
- 3. In a Huffman Tree traversal, left means 0 and right means 1.

What is the uncompressed text? Please show all of your work.

## Answer:

## PQ:

Η

A N N O N