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1 CHAIN COMPLEXES

1.1 CHAIN COMPLEXES

Want a more intuitive view of left/right exact functors, maybe in terms of lifts/ extensions.

Definition 1. A **chain complex** \mathcal{C} is a sequence of R -morphisms

$$\cdots \xrightarrow{d_{i+2}} C_{i+1} \xrightarrow{d_{i+1}} C_i \xrightarrow{d_i} C_{i-1} \xrightarrow{d_{i-1}} \cdots$$

such that $d^2 = 0$ for all i . **Cochain complexes** are the same, except the boundary maps take you up a level instead of down.

$$\cdots \xrightarrow{d_{i-1}} C^{i-1} \xrightarrow{d_i} C^i \xrightarrow{d_{i+1}} C^{i+1} \xrightarrow{d_{i+2}} \cdots$$

The map d_i is the **boundary operator**, as it is a generalization of the geometric concept of a boundary (note $d^2 = 0$). Thus an element of $\text{Im } d$ is a **boundary**. Since usual geometric cycles have no boundary, we call the elements of $\text{Ker } d$ **cycles**.

Example 1. Chain complexes generalize the concept of boundaries to objects that don't necessarily have clear cyclic geometric properties. Let $\Omega_n(M)$ denote the space of differential n -forms on a manifold M , then we have a cochain complex

$$\Omega_0(M) \xrightarrow{d} \Omega_1(M) \xrightarrow{d} \Omega_2(M) \xrightarrow{d} \cdots$$

where d is the exterior derivative. From this we see that the cycles of $\Omega_0(M)$ (the space of differentiable functions on M) are the constant functions.

A **morphism of complexes/chain morphism** $f : \mathcal{C} \rightarrow \mathcal{D}$ is a sequence of morphisms $f_i : C_i \rightarrow D_i$ respecting the boundary map, i.e. making the following diagram commute.

$$\begin{array}{ccc} C_i & \xrightarrow{d_{\mathcal{C}}} & C_{i-1} \\ f_i \downarrow & & \downarrow f_{i-1} \\ D_i & \xrightarrow{d_{\mathcal{D}}} & D_{i-1} \end{array}$$

1.2 CHAIN HOMOTOPIES

Definition 2. Given two chain complexes \mathcal{A}, \mathcal{B} , two chain morphisms $f, g : \mathcal{A} \rightarrow \mathcal{B}$ are **(chain) homotopic**, written $f \simeq g$, if there are morphisms $s_i : A_i \rightarrow B_{i-1}$ such that

$$d's + sd = f - g.$$

If \mathcal{A}, \mathcal{B} are cochain complexes instead, then $s_i : A_i \rightarrow B_{i+1}$.

$$\begin{array}{ccccc} A_{i-1} & \xrightarrow{d} & A_i & \xrightarrow{d} & A_{i+1} \\ & \nwarrow s_i & \downarrow f_i & \downarrow g_i & \nearrow s_{i+1} \\ B_{i-1} & \xrightarrow{d'} & B_i & \xrightarrow{d'} & B_{i+1} \end{array}$$

Motivation for this?

Definition 3. A chain morphism $f : \mathcal{A} \rightarrow \mathcal{B}$ is a **homotopy equivalence** if there's another chain morphism $g : \mathcal{B} \rightarrow \mathcal{A}$ such that $fg \simeq 1_B$ and $gf \simeq 1_A$.

$$\begin{array}{ccc} \mathcal{A} & \xrightarrow{f} & \mathcal{B} \\ & \xleftarrow{g} & \end{array}$$

Proposition 1. Additive functors preserve homotopy equivalence.

Proof. Let $f \simeq g$. If \mathcal{F} is additive and covariant, then $d's + sd = f - g \implies \mathcal{F}(d')\mathcal{F}(s) + \mathcal{F}(s)\mathcal{F}(d) = \mathcal{F}f - \mathcal{F}g$. Thus $\mathcal{F}f \simeq \mathcal{F}g$. If \mathcal{G} is additive and contravariant, then $\mathcal{G}(d)\mathcal{G}(s) = \mathcal{G}(s)\mathcal{G}(d') = \mathcal{G}f - \mathcal{G}g$. Since all the arrows are reversed, the LHS is the right form, so $\mathcal{G}f \simeq \mathcal{G}g$. \square

1.3 HOMOLOGY

Note 1. Big idea: given some module, we always have a way of getting a chain complex (take a resolution). Chain complexes by themselves aren't nice, though: they might be unwieldy or unnatural, similar objects might have dissimilar complexes, etc. Passing to homology makes these problems go away, though, giving us access to nice algebraic invariants.

Definition 4. The n -th **homology group** $H_n(\mathcal{C})$ of a chain complex \mathcal{C} is the kernel of the map going out of C_n quotiented by the image of the map coming into C_n . Cochain complexes have **cohomology groups** $H^n(\mathcal{C})$ instead.

Proposition 2. A chain/cochain complex is exact \iff all its homology/cohomology groups are trivial.

Thus the (co)homology groups of a (co)chain complex measure how much it fails to be exact.

Proposition 3. A morphism of complexes $f : \mathcal{A} \rightarrow \mathcal{B}$ induces group morphisms between the complexes' homology/cohomology groups given by $[a] \mapsto [f_n(a)]$.

$$\begin{array}{ccccccc} A^{n-1} & \longrightarrow & A^n & \longrightarrow & A^{n+1} & & H^n(\mathcal{A}) \\ \downarrow & & \downarrow & & \downarrow & \rightsquigarrow & \downarrow \\ B^{n-1} & \longrightarrow & B^n & \longrightarrow & B^{n+1} & & H^n(\mathcal{B}) \end{array}$$

Proof. This follows from the morphism of complexes respecting the boundary map and thus mapping the kernels and images of the first complex to the kernels and images of the second. \square

Proposition 4. If f_* is the induced (co)homology map of f , then $(gf)_* = g_*f_*$.

Definition 5. $0 \rightarrow \mathcal{A} \rightarrow \mathcal{B} \rightarrow \mathcal{C} \rightarrow 0$ is a **short exact sequence of complexes** if each $0 \rightarrow A^n \rightarrow B^n \rightarrow C^n \rightarrow 0$ is short exact.

Lemma 1 (Snake Lemma). If the following diagram has exact rows,

$$\begin{array}{ccccccc} A & \longrightarrow & B & \twoheadrightarrow & C & \longrightarrow & 0 \\ & \downarrow \alpha & & \downarrow \beta & & \downarrow \gamma & \\ 0 & \longrightarrow & A' & \twoheadrightarrow & B' & \longrightarrow & C' \end{array}$$

then there is an induced exact sequence

$$\mathrm{Ker} \alpha \rightarrow \mathrm{Ker} \beta \rightarrow \mathrm{Ker} \gamma \rightarrow \mathrm{Coker} \alpha \rightarrow \mathrm{Coker} \beta \rightarrow \mathrm{Coker} \gamma.$$

Theorem 1 (Long Exact Sequence in Cohomology). If $0 \rightarrow \mathcal{A} \rightarrow \mathcal{B} \rightarrow \mathcal{C} \rightarrow 0$ is a short exact sequence of complexes, then there is a long exact sequence of cohomologies

$$\begin{aligned} 0 \rightarrow H^0(\mathcal{A}) &\rightarrow H^0(\mathcal{B}) \rightarrow H^0(\mathcal{C}) \\ &\rightarrow H^1(\mathcal{A}) \rightarrow H^1(\mathcal{B}) \rightarrow H^1(\mathcal{C}) \\ &\rightarrow H^2(\mathcal{A}) \rightarrow \cdots \end{aligned}$$

where the morphisms $H^n(\mathcal{C}) \rightarrow H^{n+1}(\mathcal{A})$ are the **connecting morphisms**.

Proof. **Intuition? Use snake lemma (have proof of this in spectral sequences paper).** \square

Corollary 1. If $0 \rightarrow \mathcal{A} \rightarrow \mathcal{B} \rightarrow \mathcal{C} \rightarrow 0$ is exact and any 2 of the complexes are exact themselves, then so is the third.

Proof. The LES of cohomologies becomes all 0, except for each $H^n(\mathcal{X})$, where \mathcal{X} is the third complex. Now $0 \rightarrow H^n(\mathcal{X}) \rightarrow 0$ exact $\implies H^n(\mathcal{X}) \cong 0$, so \mathcal{X} is exact. \square

Definition 6. A morphism of complexes is a **quasi-isomorphism** if the (co)homology maps it induces are all iso.

Lemma 2. If $f \simeq g$, then they induce the same (co)homology maps, i.e. $f_* = g_*$.

Proof. Suppressing subscripts, suppose $f = d's + sd$, then the induced map is

$$[a] \mapsto [f(a)] = [(d's)(a) + (sd)(a)] = [d'(s(a)) + s(0)] = [0].$$

Then if $f \simeq g$, we have $[f(a)] = [(g + d's + sd)(a)] = [g(a)]$. \square

Proposition 5. A homotopy equivalence is a quasi-iso.

Proof. Suppose f and g are inverse chain homotopies, then by the lemma, $f_*g_* = (fg)_* = (1_B)_* = 1_{H(\mathcal{B})}$ and, similarly, $g_*f_* = 1_{H(\mathcal{A})}$. Thus $H^n(\mathcal{A}) \cong H^n(\mathcal{B})$ for all n . \square

2 DERIVED FUNCTORS

2.1 RESOLUTIONS

Definition 7. Suppose A is an R -module, then a **projective resolution** over A is an exact sequence of projective R -modules

$$\cdots \longrightarrow P_n \xrightarrow{d_n} P_{n-1} \xrightarrow{d_{n-1}} \cdots \xrightarrow{d_1} P_0 \xrightarrow{\varepsilon} A \longrightarrow 0$$

and a **injective resolution** over A is an exact sequence of injective R -modules

$$0 \longrightarrow A \xrightarrow{\varepsilon} I_0 \xrightarrow{d_1} \cdots \xrightarrow{d_{n-1}} I_{n-1} \xrightarrow{d_n} I_n \longrightarrow \cdots$$

Theorem 2 (Existence). Every R -module has a projective **and injective** resolution.

Proof. Let A be an R -module, and let P_0 be free on A . Then there's a unique R -morphism $\varepsilon : P_0 \rightarrow A$ extending 1_A . It's clearly epic, so $P_0 \xrightarrow{\varepsilon} A \rightarrow 0$ is exact. Now let P_1 be free on $\text{Ker } \varepsilon$, then there's a unique R -morphism $d_1 : P_1 \rightarrow \text{Ker } \varepsilon$, so $P_1 \xrightarrow{d_1} P_0 \xrightarrow{\varepsilon} A$ is exact. Continue inductively, with P_{n+1} free on $\text{Ker } d_n$. Since each P_i is free, each is projective. \square

Proposition 6. R -morphisms lift to morphisms of projective/injective resolutions that are unique up to chain homotopy.

$$\begin{array}{ccccccc} \cdots & \longrightarrow & P_1 & \longrightarrow & P_0 & \longrightarrow & A \longrightarrow 0 \\ & & \downarrow f_1 & & \downarrow f_0 & & \downarrow f \\ \cdots & \longrightarrow & P'_1 & \longrightarrow & P'_0 & \longrightarrow & A' \longrightarrow 0 \\ \\ 0 & \longrightarrow & A & \longrightarrow & I_0 & \longrightarrow & I_1 \longrightarrow \cdots \\ & & \downarrow g & & \downarrow g_0 & & \downarrow g_1 \\ 0 & \longrightarrow & A' & \longrightarrow & I'_0 & \longrightarrow & I'_1 \longrightarrow \cdots \end{array}$$

Proof. Inductively use the fact that all the P_i, P'_i are projective and I_i, I'_i are injective to get each f_n and g_n . **Chain homotopy bit.** \square

Corollary 2. Any two projective/injective resolutions of the same module are homotopy equivalent.

Proof. Consider the following lifts to projective resolutions (the case of injective resolutions is similar).

$$\begin{array}{ccccc}
 \mathcal{P} & \longrightarrow & A & \longrightarrow & 0 \\
 \downarrow f_\bullet & & \downarrow 1_A & & \\
 \mathcal{P}' & \longrightarrow & A & \longrightarrow & 0 \\
 \downarrow g_\bullet & & \downarrow 1_A & & \\
 \mathcal{P} & \longrightarrow & A & \longrightarrow & 0
 \end{array}$$

The identity map is a valid candidate for each $g_n f_n$, so $g_n f_n \simeq 1$. Now flip the diagram upside down and use the same f_n and g_n maps, yielding $f_n g_n \simeq 1$. \square

Corollary 3. Any two projective/injective resolutions have isomorphic (co)homology groups.

Proof. They're homotopy equivalent, so they're quasi-isomorphic by Proposition 5. \square

Maybe start this section more generally, with left and right resolutions...

2.2 DERIVED FUNCTORS

Note 2. Big idea: a left exact covariant functor \mathcal{F} can turn a SES into a short left exact sequence, but there is only one canonical way to extend this to a LES, and that's with right derived functors. Left derived functors do the same for right exact functors. **Why do we want a LES, though?**

Suppose we apply an additive functor \mathcal{F} to some projective/injective resolution of a module A . The (co)homology groups of the resulting complex are unique up to isomorphism:

- Let \mathcal{X}, \mathcal{Y} be projective/injective resolutions of A , then they're homotopy equivalent by Corollary 2.
- By Proposition 1, additive functors preserve homotopy equivalence, so $\mathcal{F}\mathcal{X}$ and $\mathcal{F}\mathcal{Y}$ are also homotopy equivalent.
- By Proposition 5, homotopy equivalent complexes have isomorphic (co)homology groups.

Thus we can define the (co)homology groups of the sequence gotten by applying an additive functor \mathcal{F} to a projective/injective resolution of A using *any* resolution.

Definition 8. Let \mathcal{F} be a functor and A an R -module, then choose a resolution of A from the following chart based on \mathcal{F} .

Left exact, covariant	injective
Left exact, contravariant	projective
Right exact, covariant	projective
Right exact, contravariant	injective

Apply \mathcal{F} to the resolution, remove the $\mathcal{F}A$ term from it, then take (co)homologies. If \mathcal{F} is left exact, the cohomologies $R^i\mathcal{F}$ are the **right derived functors** of \mathcal{F} . If \mathcal{F} is right exact, the homologies $L_i\mathcal{F}$ are the **left derived functors** of \mathcal{F} .

I don't think the derived functors depend on A at all, they can just be applied to A , etc. to get new objects...

With left exact functors, we end up with induced sequences of the form

$$0 \rightarrow \mathcal{F}X_0 \rightarrow \mathcal{F}X_1 \rightarrow \mathcal{F}X_2 \rightarrow \cdots,$$

thus why the derived functors are “right”. Similarly, for right exact functors, we end up with induced sequences of the form

$$\mathcal{F}X_2 \rightarrow \mathcal{F}X_1 \rightarrow \mathcal{F}X_0 \rightarrow 0,$$

thus why the derived functors are “left”. As examples, see the next two propositions.

Can you get both sets of sequences at once if \mathcal{F} is exact?

They're actually functors.... b/c (co)homology is a functor.

More detail about *why* we choose inj or proj resolution.

Proposition 7. If \mathcal{F} is left exact, then $R^0\mathcal{F} = \mathcal{F}$.

Proof. **Covariant:** If $0 \rightarrow A \xrightarrow{f} I_0 \xrightarrow{g} I_1$ is exact, then so is $0 \rightarrow \mathcal{F}A \xrightarrow{\mathcal{F}f} \mathcal{F}I_0 \xrightarrow{\mathcal{F}g} \mathcal{F}I_1$. Then $R^0\mathcal{F}(A) = \text{Ker}(\mathcal{F}g) = \text{Im}(\mathcal{F}f) \cong \mathcal{F}(A)$ (by the 1st iso theorem since $\mathcal{F}f$ monic).

Contravariant: Use a projective resolution instead. The process is the same. \square

Proposition 8. If \mathcal{F} is right exact, then $L_0\mathcal{F} = \mathcal{F}$.

Proof. **Covariant:** If $P_1 \xrightarrow{f} P_0 \xrightarrow{g} A \rightarrow 0$ is exact, then so is $\mathcal{F}P_1 \xrightarrow{\mathcal{F}f} \mathcal{F}P_0 \xrightarrow{\mathcal{F}g} \mathcal{F}A \rightarrow 0$. Then $L_0\mathcal{F}(A) = \frac{\mathcal{F}P_0}{\text{Im } \mathcal{F}f} = \frac{\mathcal{F}P_0}{\text{Ker } \mathcal{F}g} \cong \mathcal{F}A$ (by the 1st iso theorem since $\mathcal{F}g$ epic).

Contravariant: Use an injective resolution instead. The process is the same. \square

Does “left derived functor of left exact functor” make sense? Are they all just trivial or something? To prove something like that, would you use a left exact variant of “exact functors preserves LES’s”?

Theorem 3 (LES of derived functors). **Do this.**

If derived functors measure the extent to which a functor fails to be exact, then an exact functor should have trivial derived functors. This turns out to be true.

Proposition 9. If \mathcal{F} is exact, then $R^i\mathcal{F} = L_i\mathcal{F} = 0$ for all $i > 0$.

Proof. **This uses the fact that exact functors preserves LES’s... prove this.** **Covariant:** Exact functors preserve exactness, so $0 \rightarrow A \rightarrow I_0 \rightarrow I_1 \rightarrow \cdots$ exact implies $0 \rightarrow \mathcal{F}A \rightarrow \mathcal{F}I_0 \rightarrow \mathcal{F}I_1 \rightarrow \cdots$ exact. Chopping off the $\mathcal{F}A$ term and taking cohomologies gives $L_i\mathcal{F} = 0$ when $i > 0$. Now repeat the argument with a projective resolution for the R^i .

Contravariant: Similar argument. \square

Proposition 10. Fix a functor \mathcal{F} . If A is projective/injective (depending on the type of \mathcal{F}), then $R^i\mathcal{F}(A)$ or $L_i\mathcal{F}(A)$ (whichever is correct for \mathcal{F}) is trivial when $i > 0$.

Proof. We consider the case when \mathcal{F} is left exact and covariant, but the other three cases are similar. Suppose A is injective, then $0 \rightarrow A \xrightarrow{\text{id}} A \rightarrow 0$ is an injective resolution of A . This induces the exact sequence $0 \rightarrow \mathcal{F}A \xrightarrow{\text{id}} \mathcal{F}A \rightarrow 0$. Chopping off the first $\mathcal{F}A$ term and taking cohomologies gives $R^i\mathcal{F}(A) = 0$. \square

2.3 THE EXT FUNCTOR

Note 3. Big idea: the hom functors are left exact, but we can use cohomology to measure how much they fail to be right exact.

Definition 9. The **Ext functors** are the (right) derived functors of the hom functors.

Given an R -module A , there are two equivalent ways to construct them:

1. **Using $\text{Hom}(-, M)$ (contravariant):** Take a projective resolution

$$\cdots \longrightarrow P_n \xrightarrow{d_n} P_{n-1} \xrightarrow{d_{n-1}} \cdots \xrightarrow{d_1} P_0 \xrightarrow{\varepsilon} A \longrightarrow 0$$

and apply $\text{Hom}(-, M)$ to it. Removing $\text{Hom}(A, M)$ from the sequence gives

$$0 \longrightarrow \text{Hom}(P_0, M) \xrightarrow{d_1^*} \text{Hom}(P_1, M) \xrightarrow{d_2^*} \cdots$$

This is a cochain complex since for any map f , applying d^* twice gives $d^{*2}(f) = f d^2 = 0$. Then $\text{Ext}_R^n(A, M)$ is the n -th cohomology group of this complex.

2. **Using $\text{Hom}(M, -)$ (covariant):** Take an injective resolution, apply $\text{Hom}(M, -)$, remove $\text{Hom}(M, A)$, then take homology. **Is this actually equivalent?**

Proposition 11. $\text{Ext}_R^0(A, M) \cong \text{Hom}_R(A, M)$.

Proof. The hom functors are left exact, so apply Proposition 7. □

Ext(A,B) is contra in A, cov in B.