CP-Algorithms

Search

Placing Bishops on a Chessboard

Table of Contents

- Algorithm
- Implementation

Find the number of ways to place K bishops on an N imes N chessboard so that no two bishops attack each other.

Algorithm

This problem can be solved using dynamic programming.

Let's enumerate the diagonals of the chessboard as follows: black diagonals have odd indices, white diagonals have even indices, and the diagonals are numbered in non-decreasing order of the number of squares in them. Here is an example for a 5×5 chessboard.

```
2
       5
           6
               9
1
2
   5
       6
           9
               8
5
   6
       9
           8
               7
           7
6
   9
       8
               4
       7
               3
9
   8
           4
```

Let **D**[i][j] denote the number of ways to place j bishops on diagonals with indices up to i which have the same color as diagonal i. Then i = 1...2N-1 and j = 0...K.

We can calculate D[i][j] using only values of D[i-2] (we subtract 2 because we only consider diagonals of the same color as i). There are two ways to get D[i] [j]. Either we place all j bishops on previous diagonals: then there are D[i-2][j] ways to achieve this. Or we place one bishop on diagonal i and j-1 bishops on previous diagonals. The number of ways to do this equals the number of squares in diagonal i minus j-1, because each of j-1 bishops placed on previous diagonals will block one square on the current diagonal. The number of squares in diagonal i can be calculated as follows:

```
int squares (int i) {
   if (i & 1)
      return i / 4 * 2 + 1;
   else
      return (i - 1) / 4 * 2 + 2;
}
```

```
The base case is simple: D[i][0] = 1, D[1][1] = 1.
```

Once we have calculated all values of <code>D[i][j]</code>, the answer can be obtained as follows: consider all possible numbers of bishops placed on black diagonals <code>i=0...K</code>, with corresponding numbers of bishops on white diagonals <code>K-i</code>. The bishops placed on black and white diagonals never attack each other, so the placements can be done independently. The index of the last black diagonal is <code>2N-1</code>, the last white one is <code>2N-2</code>. For each <code>i</code> we add <code>D[2N-1][i] * D[2N-2][K-i]</code> to the answer.

Implementation

```
int bishop_placements(int N, int K)
{
    if (K > 2 * N - 1)
        return 0;

    vector<vector<int>>> D(N * 2, vector<int>)(K
    for (int i = 0; i < N * 2; ++i)
        D[i][0] = 1;

    D[1][1] = 1;

    for (int i = 2; i < N * 2; ++i)
        for (int j = 1; j <= K; ++j)
            D[i][j] = D[i-2][j] + D[i-2][j-1]

    int ans = 0;
    for (int i = 0; i <= K; ++i)
        ans += D[N*2-1][i] * D[N*2-2][K-i];</pre>
```

return ans;
}

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