CP-Algorithms



Dynamic Programming on Broken Profile. Problem "Parquet"

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Common problems solved using DP on broken profile include:

- finding number of ways to fully fill an area (e.g. chessboard/grid) with some figures (e.g. dominoes)
- finding a way to fill an area with minimum number of figures
- finding a partial fill with minimum number of unfilled space (or cells, in case of grid)
- finding a partial fill with the minimum number of figures, such that no more figures can be added

Problem "Parquet"

Problem description. Given a grid of size $N \times M$. Find number of ways to fill the grid with figures of size

 2×1 (no cell should be left unfilled, and figures should not overlap each other).

Let the DP state be: D[i][mask], where i = 1 ... N, and mask = 0 ... 2^M-1. i respresents number of rows in the current grid, and mask is the state of last row of current grid. If j-th bit of mask is 0 then the corresponding cell is filled, otherwise it is unfilled. Clearly, the answer to the problem will be D[N][0].

We will be building the DP state by iterating over each i = 1 ... N and each mask = 0 ... 2^M-1, and for each mask we will be only transitioning forward, that is, we will be adding figures to the current grid.

Implementation

```
else
         {
             calc (x, y+1, mask, next_mask | my
             if (y+1 < m \&\& ! (mask \& my_mask))
                  calc (x, y+2, mask, next_mask)
         }
    }
}
int main()
{
    cin >> n >> m;
    d.resize (n+1, vector<long long> (1<<m));</pre>
    d[0][0] = 1;
    for (int x=0; x< n; ++x)
         for (int mask=0; mask<(1<<m); ++mask)</pre>
             calc (x, 0, mask, 0);
    cout << d[n][0];</pre>
}
```

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