



*CS1699: Blockchain Technology and Cryptocurrency*

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## 3. Cryptographic Hashing Basics

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# What is a hash function?

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- ❖ A function which accepts some arbitrary input  $x$  and returns a fixed-length “summary”
- ❖ Let's look at “bad hash”

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# Bad Hash

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- ❖ Converts all values in a string to their ASCII values
- ❖ Sums up all of these values
- ❖ Returns value modulo 256 (results 0x0 - 0xFF)
- ❖ Note: this is hashing function, but not a good one for our purposes!
- ❖ See `./sample_code/hash/bad_hash.rb` to run an example, e.g. *`ruby bad_hash.rb "meow"`*



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# Hash vs Cryptographic Hash

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- ❖ Hash values can be used for a variety of purposes (e.g. hash maps, data distribution, nearest neighbor search etc.)
- ❖ These different kinds of hash functions will have different properties (e.g. continuous hash functions are great for nearest neighbor search, horrible for data distribution!)
- ❖ For our purposes, we are interested in using hash functions for *cryptographic hashes*

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# Cryptographic Hashes

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- ❖ Should have the following properties:
  - ❖ collision-free
  - ❖ hiding
  - ❖ puzzle-friendly

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# Collision-Free

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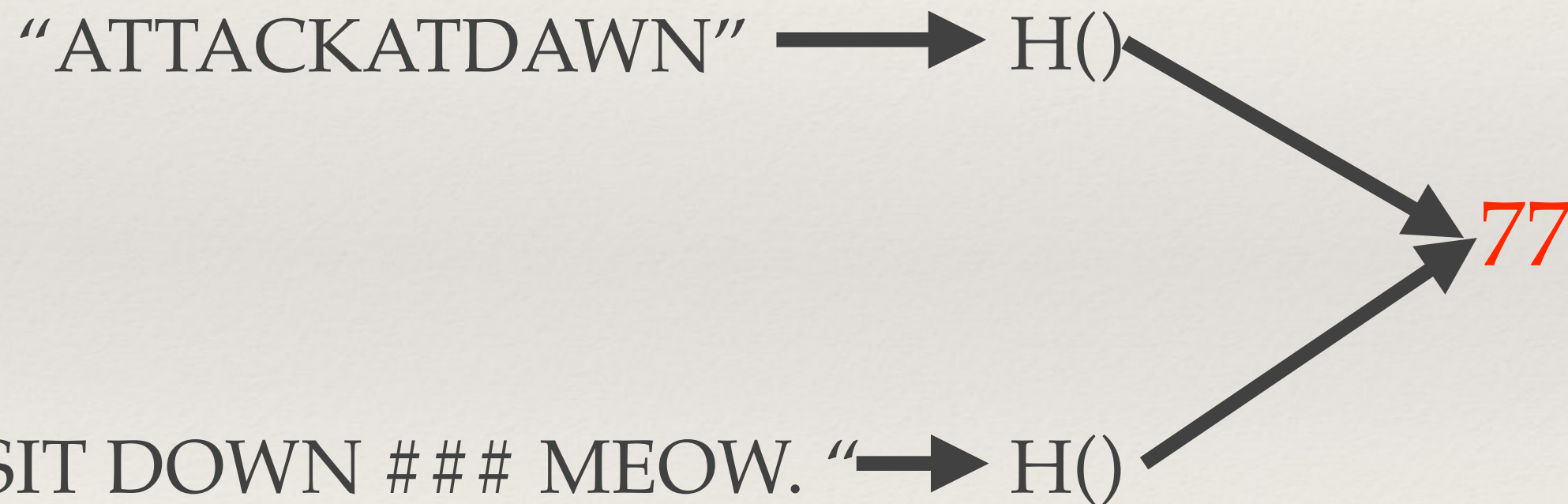
- ❖ For hash function  $H(x)$ , should be computationally infeasible to find a collision, that is, some  $y$  where :
  - ❖  $x \neq y$
  - ❖  $H(x) == H(y)$



# Collision Example with Bad Hash

```
(23477) $ ruby bad_hash.rb "ATTACKATDAWN"  
Hash of 'ATTACKATDAWN' is: 77
```

```
(23478) $ ruby bad_hash.rb "SIT DOWN ### MEOW. "  
Hash of 'SIT DOWN ### MEOW. ' is: 77
```



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# Collisions Always Exist!

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- ❖ A mapping from an infinite domain ( $x$ ) to a finite codomain ( $H(x)$ ) will have, in fact, infinite collisions
- ❖ So while collision-free is our goal, usually we say “collision-resistant”
- ❖ But can they be found in a computationally feasible way (i.e., a reasonable amount of time)?



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# Finding Collisions

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- ❖ Brute force is always possible, but timescales for large hash functions are massive (millions of years using today's computers or even networks of computers)
- ❖ BUT - holes are always possible in any hash function  $H(x)$  - and have been exploited / broken in the past!
  - ❖ Broken = collisions can be found in a way significantly faster than brute force
  - ❖ SHA-1, GOST, MD2, MD4, MD5, Panama
- ❖ No hash function has been PROVEN secure

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# Can We Find A Collision?

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- ❖ With bad hash, sure
- ❖ Easy to calculate result - can move up or down knowing ASCII values of characters in original message  $m$
- ❖ Only 256 possible hash values, thus vulnerable to brute force

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# Using A Hash As a Message Digest

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- ❖ Generally speaking, if  $H(x) == H(y)$ , we can be reasonably certain that  $x == y$
- ❖ So if we want to check, say, if we have seen something before, we can store its hash instead of the actual data
- ❖ Example: git (Side note: git uses the broken SHA-1 algorithm! See <https://shattered.io/> and <https://blog.github.com/2017-03-20-sha-1-collision-detection-on-github-com/>)



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# Hiding

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- ❖ Given the result of a hash function  $H(x)$ , it is infeasible to “go backwards” to determine  $x$
- ❖ That is,  $H(x)$  should be a one-way function

$$41 \longrightarrow H(x) = x + 1 \longrightarrow 42$$

*No hiding: given  $H(x)$ , very easy to find  $x$ !*

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# Hiding: A Formal Definition

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- ❖ “A hash function  $H$  is said to be hiding when a secret value  $r$  is chosen from a probability distribution that has *high min-entropy*, then, given  $H(r \parallel x)$ , it is infeasible to find  $x$ ”
- ❖ min-entropy: a measure of how predictable an outcome is
- ❖ In other words, for any given input, if  $r$  is chosen uniformly, the result is likely to be any one of the possible output values

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# Bad Hash - Is it hiding?

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- ❖ No, assuming low-ASCII (first 128 possibilities)
- ❖ Adding an additional character will strongly favor one half of the possibilities in our output!



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# Hiding

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```
(23480) $ ruby bad_hash.rb "AAAA"  
Hash of 'AAAA' is: 04
```

```
(23481) $ ruby bad_hash.rb "AAAA "  
Hash of 'AAAA ' is: 24
```

```
(23482) $ ruby bad_hash.rb "AAAAZ"  
Hash of 'AAAAZ' is: 5E
```

```
(23483) $ ruby bad_hash.rb "AAAAz"  
Hash of 'AAAAz' is: 7E
```

```
(23484) $ ruby bad_hash.rb "AAAA~"  
Hash of 'AAAA~' is: 82
```

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# Application of Hiding: Commitment

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- ❖ Assume you want to wager on something happening, but you don't want others to know your guess
- ❖ e.g., "The price of one bitcoin will be over \$10,000 in one year"

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# Commitment Scheme

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$\text{com} := \text{commit}(m, \text{nonce})$   
 $\text{verify}(\text{com}, \text{msg}, \text{nonce})$

A nonce is random value that can be used only once (i.e., the  $r$  in our hiding definition). Nonces will come up very often in our exploration of Bitcoin!

Assumes the commit function is hiding and binding

*Binding* = collision-resistant

Given two pairs  $m / \text{nonce}$  and  $m' / \text{nonce}'$ , it is infeasible to find  $m \neq m'$  and  $\text{commit}(m, \text{nonce}) == \text{commit}(m', \text{nonce}')$



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# Commitment Scheme

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- ❖ To commit to something:
  - ❖  $(com, key) = commit(m)$
  - ❖ Publish  $com$  publicly
- ❖ To “open the envelope”
  - ❖ Publish  $key$  and  $m$
  - ❖ Anyone can use  $verify$  function to verify  $m$

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# Puzzle Friendliness

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- ❖ “For every possible  $n$ -bit output value  $y$ , if  $k$  is chosen from a distribution with high min-entropy, then it is infeasible to find  $x$  such that  $H(K \parallel x) == y$  in time significantly less than  $2^n$ ”

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# Puzzle Friendliness

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- ❖ In other words, if someone wants the hash value result  $H(x)$  to be a particular value  $y$ , if  $H(x)$  produces a result of size  $n$  bits, then finding some value  $x$  such that  $H(x) == y$  should require at least  $2^n$  attempts (  $O(2^n)$  )
- ❖ In other other words, the best way to find the result of a hash function is to just do the hash with different inputs, and brute force it



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# Puzzle Friendliness Application

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- ❖ Proof of work
- ❖ I want to prove that you have done some work before you get a prize
- ❖ So I may ask you to find a value  $x$  where  $H(x)$  gives me 0
- ❖ Best way to do it SHOULD be to just try every possible input

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# Is Bad Hash Puzzle Friendly?

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- ❖ No
- ❖ Very easy to determine output and modify input to modify output

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# A Shortcut for Bad Hash

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```
(23485) $ ruby bad_hash.rb "AAAA"  
Hash of 'AAAA' is: 04
```

```
(23486) $ ruby bad_hash.rb "AAA<"  
Hash of 'AAA<' is: FF
```

```
(23487) $ ruby bad_hash.rb "AAA?"  
Hash of 'AAA?' is: 02
```

```
(23488) $ ruby bad_hash.rb "AAA="  
Hash of 'AAA=' is: 00
```



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# Hash Functions Used in Bitcoin

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- ❖ SHA-256: Created by NSA, 256-bit output
- ❖ RIPEMD-160: Developed by researchers at Katholieke Universiteit Leuven, 160-bit output

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# SHA 256

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- ❖ SHA 256 is a “good” hash for our purposes
- ❖ As far as we know (it has not been proven!), it is:
  - ❖ Collision-resistant
  - ❖ Hiding
  - ❖ Puzzle-friendly

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# SHA-256 Shortcut?

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```
(23489) $ ruby sha256_hash.rb "AAAA"
```

```
Hash of 'AAAA' is:
```

```
63c1dd951ffedf6f7fd968ad4efa39b8ed584f162f46e715114ee184f8de9201
```

```
(23490) $ ruby sha256_hash.rb "AAA<"
```

```
Hash of 'AAA<' is:
```

```
ad587d6af700f2625287ce9b0c66120a120a0b56fd3c8049cc9d0596832d60a6
```

```
(23491) $ ruby sha256_hash.rb "AAAB"
```

```
Hash of 'AAAB' is:
```

```
5e01d7a18db038714fcbb6de5708f3f3c36ef61c206c72ce5dd91c149ecc910c
```

```
(23492) $ ruby sha256_hash.rb "AAA="
```

```
Hash of 'AAA=' is:
```

```
e2d4768b3472b90ca749600da34e622122d16f61517c459a5d7915693ec4ccf3
```



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# SHA-256 in Bitcoin

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- ❖ Mining: SHA-256 hashing
- ❖ Internal integrity checks
- ❖ Hash pointers for Merkle tree
- ❖ SHA-256 hashing of ECDSA as part of key generation