

8-6

a. ~~C_n^2~~
 $C_n^n = \frac{2n \cdot (2n-1) \cdot (2n-2) \dots n}{n \cdot (n-1) \dots 1}$

b. ~~C_n^n~~
 since $C_n^n \leq 2^n$ (h is the number of comparisons)

~~C_n^n~~
 So $h \geq \lg \frac{2n(2n-1)(2n-2) \dots n}{n(n-1) \dots 1}$
 ~~$\Rightarrow \lg \frac{2n(2n-1)(2n-2) \dots n}{n(n-1) \dots 1}$~~
 ~~$\Rightarrow \lg \frac{2n(2n-1)(2n-2) \dots n}{n(n-1) \dots 1}$~~

$= \lg \frac{(2n)!}{(n!)^2}$

~~$= \theta(n \lg n)$~~
 $= \lg(n \cdot (n-1) \cdot (1)) - \lg(n \cdot (n-1) \cdot (n-2) \dots 1)$

$= \theta(n \lg n) - 2\theta(n \lg n) = \theta(n)$

8-7

a. Since $A[p]$ is the smallest value in A , $A[q]$ is definitely greater than $A[p]$.
 $B[p] = 0$. Because $A[p] = A[p]$
 so the satisfy the first condition.

$B[q] = 1$ Because $A[q] > A[p]$, B

b. Algorithm fails to sort. Since $A[p]$ is in the array location. And the array location that $A[p]$ should go is taken by $A[q]$. For an $i < p$. Since $A[p]$ is the smallest value in A . then $A[i] \geq A[p]$.
 so $B[i] = 1$ for an value $j > i$, if $A[j] \neq A$
 $A[p]$ is the smallest value in A . then B would be an array contains only
 then 1 would contain only

