# ACM 模板 for 北湖咸鱼

### 三条咸鱼——zhber、strawberry、miamiao

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# 一、头文件

```
1 #include <bits/stdc++.h>
   using namespace std;
3 #define eps 1e-6
   #define e exp(1.0)
4
5
   #define pi acos(-1.0)
6
   #define fi first
7
   #define se second
8 #define pb push back
   #define mp make_pair
9
10 #define SZ(x) ((int)(x).size())
#define All(x) (x).begin(),(x).end()
12 #define rep(i,a,n) for (int i=a;i<n;i++)</pre>
13 #define per(i,a,n) for (int i=n-1;i>=a;i--)
14 #define Close() ios::sync_with_stdio(0),cin.tie(0)
15 #define INF 0x3f3f3f3f
16 typedef vector<int> VI;
17 typedef pair<int, int> PII;
18 typedef long long ll;
19 typedef long double ld;
20 typedef unsigned long long ull;
21 const int maxn = 1e5 + 10;
22 const int maxm = 3e3 + 5;
```

```
23 const int mod = 1e9 + 7;
```

注:①部分模板可能会重新定义新类型;②不同模板的 maxn、maxm 大小取决于题目要求

### 二、算法

#### 1.快速读入

```
①普通快速读
1
   inline ll read() {
          11 x=0,f=1;char ch=getchar();
2
3
          while(ch<'0'||ch>'9'){if(ch=='-')f=-1;ch=getchar();}
          while(ch>='0'&&ch<='9'){x=x*10+ch-'0';ch=getchar();}</pre>
4
5
          return x*f;
6
   }
  ②文件读入流
   namespace fastIO {
1
2
   #define BUF_SIZE 100000
3
          //fread -> read
4
          bool IOerror = 0;
           inline char nc() {
5
                  static char buf[BUF_SIZE], *p1 = buf + BUF_SIZE, *pend = buf
6
7
   + BUF SIZE;
8
                  if(p1 == pend) {
9
                         p1 = buf;
                         pend = buf + fread(buf, 1, BUF_SIZE, stdin);
10
11
                         if(pend == p1) {
12
                                IOerror = 1;
13
                                return -1;
14
15
                  } return *p1++;
16
           inline bool blank(char ch) {
17
                  return ch == ' ' || ch == '\n' || ch == '\r' || ch == '\t';
18
19
20
           inline void read(l1 &x) {
21
                  char ch;
22
                  11 f = 1;
                  while(blank(ch = nc()));
23
                  while (ch == '-')f = -f, ch = nc();
24
25
                  if(IOerror)return;
                  for(x = ch - 0; (ch = nc()) >= 0 && ch <= 9; x = x * 10
26
   + ch - '0');
27
                  x *= f;
28
29
           inline void read(int &x) {
30
31
                  char ch;
                  int f = 1;
32
```

```
33
                 while(blank(ch = nc()));
34
                 while (ch == '-')f = -f, ch = nc();
                 if(IOerror)return;
35
                 for(x = ch - '0'; (ch = nc()) >= '0' && ch <= '9'; x = x * 10
36
   + ch - '0');
37
                 x *= f;
38
39
          inline void read(char *s) {
40
41
                 char ch;
42
                 while(blank(ch = nc()));
                 if(IOerror)return;
43
                 int i = 0;
44
45
                 for (; !blank(ch); s[i++] = ch, ch = nc());
                 s[i] = '\0';
46
47
48
   #undef BUF_SIZE
49
   };
50 using namespace fastIO;
2.快速幂
1
   11 quickmod(11 a, 11 b, 11 m) {
2
          ll ans = 1;
3
          while(b) { //用一个循环从右到左便利 b 的所有二进制位
4
                 if(b & 1) { //判断此时 b[i]的二进制位是否为 1
                        ans = (ans * a) % m; //乘到结果上,这里 a 是 a^(2^i)%m
5
                        b--;//把该为变 0
6
7
                 }
8
                 b /= 2;
9
                 a = a * a % m;
10
          }
11
          return ans;
12
3.在线求 LCA
   void RMQ() {
1
2
          for (int j = 1; (1 << j) <= n; j++)
                 for (int i = 1; i <= n; i++)
3
4
                        if (~pa[i][j - 1])
5
                                pa[i][j] = pa[pa[i][j - 1]][j - 1];
6
   }
7
   int LCA(int x, int y) {
8
          if (deep[x] < deep[y]) swap(x, y);</pre>
9
          int i, j;
10
          for (i = 0; (1 << i) <= deep[x]; i++); i--;
          for (j = i; j >= 0; j--) // 把深度统一
11
12
                 if (deep[x] - (1 << j) >= deep[y])
13
                         x = pa[x][j];
14
          if (x == y) return x;
15
          for (j = i; j >= 0; j--) //去找 LCA
```

```
if (pa[x][j] != -1 && pa[x][j] != pa[y][j]) {
16
17
                         x = pa[x][j];
18
                         y = pa[y][j];
19
20
           return pa[x][0];
21 }
4. 二维 RMQ
   int num[maxn][maxn];
2
   int st[maxn][maxn][9][9];
3
   void build st(int n, int m) {
           for (int i = 1; i <= n; ++i)
4
5
                  for (int j = 1; j <= m; ++j)</pre>
6
                         st[i][j][0][0] = num[i][j];
7
           for (int g = 0; g < 10; ++g) {
8
                  for (int t = 0; t < 10; ++t) {
9
                         if(g == 0 \&\& t == 0)
10
                                 continue;
                         for (int i = 1; i <= n; ++i) {</pre>
11
12
                                 for (int j = 1; j <= m; ++j) {
13
                                        if(i + (1 << g) - 1 > n || j + (1 << t)
14
   -1 > m)
15
                                               break;
16
                                        if(g == 0)
17
                                               st[i][j][g][t] =
   \max(st[i][j][g][t - 1], st[i][j + (1 << (t - 1))][g][t - 1]);
18
19
                                        else
20
                                               st[i][j][g][t] = max(st[i][j][g -
   1][t], st[i + (1 << (g - 1))][j][g - 1][t]);
21
22
                                 }
23
                         }
24
                  }
25
           }
26
27
   int query(int a, int b, int c, int d) {
28
           int k = (c - a + 1);
29
           int t = (d - b + 1);
           k = log2(1.0 * k) + eps;
30
          t = log2(1.0 * t) + eps;
31
           return \max(\max(st[a][b][k][t], st[c - (1 << k) + 1][b][k][t]),
32
   \max(st[a][d - (1 << t) + 1][k][t], st[c - (1 << k) + 1][d - (1 << t) +
34 1][k][t]));
35 }
5.Kmp
  ①Char 版本
1
   char n[N], m[N];
2
   int f[N];
   void compute(char *P) {
```

```
4
           int len = strlen(P);
5
           memset(f, 0, sizeof(f));
           for (int i = 1; i < len; i++) {</pre>
6
7
                  int j = f[i];
8
                  while (j && P[i] != P[j])
9
                          j = f[j];
                  f[i + 1] = (P[i] == P[j]) ? j + 1 : 0;
10
11
           }
12
   }
   bool kmp(char *T, char *P) {
13
14
           int n = strlen(T);
15
           int m = strlen(P);
16
           compute(P);
           int q = 0;
17
18
           for (int i = 0; i < n; i++) {
19
                  while (q && P[q] != T[i])
20
                          q = f[q];
21
                  if (P[q] == T[i])
22
                          q++;
23
                  if (q == m)
24
                                         //修改这里可以实现很多功能
                          return 1;
25
           }
26
           return 0;
27 }
  ②String 版本
1
   string n, m;
2
   int f[maxn];
3
   void compute(string &P) {
           int len = P.length();
4
5
           memset(f, 0, sizeof(f));
           for (int i = 1; i < len; i++) {</pre>
6
7
                  int j = f[i];
8
                  while (j && P[i] != P[j])
9
                          j = f[j];
10
                  f[i + 1] = (P[i] == P[j]) ? j + 1 : 0;
11
           }
12
   bool kmp(string &T, string &P) {
13
14
           int n = T.length();
15
           int m = P.length();
           compute(P);
16
17
           int q = 0;
           for (int i = 0; i < n; i++) {
18
19
                  while (q && P[q] != T[i])
20
                          q = f[q];
21
                  if (P[q] == T[i])
22
                          q++;
23
                  if (q == m)
```

```
24
                         return 1;
25
26
          return 0;
27 }
6.背包
   int dp[maxn];
2
   //0-1 背包,代价为 cost,获得的价值为 weight
   void ZeroOnePack(int cost, int weight, int nValue) {
4
          for (int i = nValue; i >= cost; i--)
5
                 dp[i] = max(dp[i], dp[i - cost] + weight);
6
   }
7
   //完全背包,代价为 cost,获得的价值为 weight
   void CompletePack(int cost, int weight, int nValue) {
          for (int i = cost; i <= nValue; i++)</pre>
9
10
                 dp[i] = max(dp[i], dp[i - cost] + weight);
11
   }
12
   //多重背包
   void MultiplePack(int cost, int weight, int amount, int nValue) {
14
          if (cost * amount >= nValue)
15
                 CompletePack(cost, weight, nValue);
16
          else {
                 int k = 1;
17
18
                 while (k < amount) {</pre>
                        ZeroOnePack(k * cost, k * weight, nValue);
19
20
                         amount -= k;
21
                         k <<= 1;
22
23
                 ZeroOnePack(amount * cost, amount * weight, nValue);
24
          }
25 }
```

# 三、数据结构

#### 1.单调栈

```
int n, a[N], 1[N];
2
   int st[N], t;
3
   void monotone() {
           t = 0;
4
5
           for (int i = 1; i <= n; i++) {
6
                  while (t > 0 && a[st[t - 1]] ? = a[i])
7
                  l[i] = t == 0 ? 0 : st[t - 1];
8
9
                  st[t++] = i;
10
           }
11 }
```

#### 2.线段树

```
注意:这里每个节点存了左和右边界,在极限情况可以省略。
  ①结构体版本
   //这是一个区间加和 RMO
1
   class SegmentTree {
2
   public:
3
4
   #define lson (root << 1)</pre>
5
   #define rson (root << 1 | 1)</pre>
   #define lent (t[root].r - t[root].l + 1)
6
7
   #define lenl (t[lson].r - t[lson].l + 1)
   #define lenr (t[rson].r - t[rson].l + 1)
8
9
           struct Tree {
                  int 1, r, val, lazy;
10
11
           } t[maxn << 4];</pre>
12
13
           void pushup(int root) {
                  t[root].val = t[lson].val + t[rson].val;
14
15
           }
16
17
           void pushdown(int root) {
18
                  if (t[root].lazy) {
                         t[lson].lazy += t[root].lazy;
19
20
                         t[rson].lazy += t[root].lazy;
21
                         t[lson].val += lenl * t[root].lazy;
22
                         t[rson].val += lenr * t[root].lazy;
                         t[root].lazy = 0;
23
24
                  }
25
           }
26
27
           void build(int 1, int r, int root) {
                  t[root].1 = 1;
28
29
                  t[root].r = r;
30
                  t[root].lazy = 0;
31
                  if (1 == r) {
32
                         t[root].val = 0;
33
                         return;
34
                  int mid = 1 + r \gg 1;
35
                  build(1, mid, lson);
36
37
                  build(mid + 1, r, rson);
                  pushup(root);
38
39
           }
40
           void update(int 1, int r, int val, int root) {
41
42
                  if (1 <= t[root].1 && t[root].r <= r) {</pre>
43
                         t[root].val += lent * val;
44
                         t[root].lazy += val;
45
                         return;
```

```
}
46
47
                  pushdown(root);
                  int mid = t[root].l + t[root].r >> 1;
48
                  if (1 <= mid) update(1, r, val, lson);</pre>
49
                  if (r > mid) update(l, r, val, rson);
50
                  pushup(root);
51
52
           }
53
           int query(int 1, int r, int root) {
54
55
                  if (1 <= t[root].1 && t[root].r <= r)</pre>
56
                         return t[root].val;
                  pushdown(root);
57
58
                  int mid = t[root].l + t[root].r >> 1;
                  int ans = 0;
59
60
                  if (1 <= mid) ans += query(1, r, lson);</pre>
                  if (r > mid) ans += query(l, r, rson);
61
62
                  return ans;
63
64 #undef lenr
65 #undef lenl
66 #undef lent
67 #undef rson
68 #undef lson
69 };
  ②数组版本(占坑)
3.树状数组
  ①一维树状数组
   int c[maxn];
1
   inline int lowbit(int x) { return x & (-x); }
   inline void modify(int x, int y, int data) {
3
          while (x \le y) {
4
5
                  c[x] += data;
6
                  x += lowbit(x);
7
           }
8
   }
   inline int query(int x) {
9
           int sum = 0;
10
           while (x > 0) {
11
                  sum += c[x];
12
13
                  x -= lowbit(x);
14
           }
15
           return sum;
16
   }
  ②二维树状数组——记得更新四个点
1
   class BTree {
2
           ll a[N][N];
3
           int n;
4
           void init() {
```

```
5
                  n = 0;
6
                  memset(a, 0, sizeof a);
7
           }
           void lowbit(int x) {
8
9
                  return x & (-x);
10
           }
           void modify(int x, int y, int d) {
11
12
                  for (int i = x; i <= n; i += lowbit(i))</pre>
13
                         for (int j = y; j <= n; j += lowbit(j))</pre>
14
                                 a[i][j] += d;
15
           }
           11 query(int x, int y) {
16
17
                  11 \text{ res} = 0;
18
                  for (int i = x; i \ge 0; i = lowbit(i))
19
                         for (int j = y; j >= 0; j -= lowbit(j))
20
                                 res += a[i][j];
21
                  return res;
22
           }
23 };
4.树链剖分
  ①点权式
1
   // 点权式 - 配合线段树的单点修改+区间查询
2
   struct Edge {
3
           int to, next;
4
           Edge() {}
5
           Edge(int a, int b) { to = a; next = b; }
6
   } E[maxn << 1];
7
   int head[maxn], cnt, tot;
   int top[maxn], son[maxn], size[maxn], deep[maxn], pa[maxn], id[maxn];
8
9
   int a[maxn];
   void init() {
10
11
           memset(head, -1, sizeof head);
12
          tot = cnt = 0;
13
   }
   void addedge(int u, int v) {
14
15
           E[cnt].to = v;
16
           E[cnt].next = head[u];
17
          head[u] = cnt++;
18
   }
19
   void dfs1(int u, int fa, int d) {
           size[u] = 1; deep[u] = d; son[u] = 0;
20
21
           for (int i = head[u]; ~i; i = E[i].next) {
                  int v = E[i].to;
22
23
                  if (v != fa) {
24
                         dfs1(v, u, d + 1);
25
                         pa[v] = u;
26
                         size[u] += size[v];
27
                         if (size[v] > size[son[u]]) son[u] = v;
```

```
28
                  }
29
           }
30
   }
   void dfs2(int u, int first) {
31
32
           top[u] = first;
33
           id[u] = ++tot;
           if (son[u]) dfs2(son[u], first);
34
35
           for (int i = head[u]; ~i; i = E[i].next) {
36
                  int v = E[i].to;
37
                  if (v != pa[u] && v != son[u]) dfs2(v, v);
38
           }
39
   }
   void solve(int u, int v) { //注意改的是这里
40
           int x = top[u], y = top[v], res = 0;
41
           while (x != y) {
42
43
                  if (deep[x] < deep[y]) {</pre>
44
                         swap(u, v);
45
                         swap(x, y);
                  }
46
47
                  res += T.query(id[x], id[u], 1);
48
                  u = pa[x];
49
                  x = top[u];
50
51
           if (deep[u] > deep[v]) swap(u, v);
52
           res += T.query(id[u], id[v], 1);
53
           printf("%d\n", res);
54 }
  ②边权式
   // 边权式: 把边权偏移到深度较大的点, 化成点权式
1
2
   struct Point {
3
           int from, to, val;
4
   } p[maxn];
5
   struct Edge {
6
           int to, next;
7
   } E[maxn << 1];
   int head[maxn], cnt, tot;
8
   int top[maxn], son[maxn], size[maxn], deep[maxn], pa[maxn], id[maxn];
9
10
   int a[maxn];
11
   void init() {
12
           memset(head, -1, sizeof head);
13
          tot = cnt = 0;
14
15
   void addedge(int u, int v) {
           E[cnt].to = v;
16
17
           E[cnt].next = head[u];
18
          head[u] = cnt++;
19
   void dfs1(int u, int fa, int d) {
20
21
           size[u] = 1; deep[u] = d; son[u] = 0;
```

```
22
           for (int i = head[u]; ~i; i = E[i].next) {
23
                  int v = E[i].to;
24
                  if (v != fa) {
25
                          dfs1(v, u, d + 1);
26
                          pa[v] = u;
27
                          size[u] += size[v];
28
                          if (size[v] > size[son[u]]) son[u] = v;
29
                  }
30
           }
31
32
   void dfs2(int u, int first) {
33
           top[u] = first;
34
           id[u] = ++tot;
           if (son[u]) dfs2(son[u], first);
35
           for (int i = head[u]; ~i; i = E[i].next) {
36
37
                  int v = E[i].to;
38
                  if (v != pa[u] && v != son[u]) dfs2(v, v);
39
           }
40
   }
41
   int solve(int u, int v) {
42
           int x = top[u], y = top[v];
43
           int ans = 0;
           while (x != y) {
44
45
                  if (deep[x] < deep[y]) {</pre>
46
                          swap(u, v);
47
                          swap(x, y);
48
                  }
49
                  ans += T.query(id[x], id[u], 1);
50
                  u = pa[x];
                  x = top[u];
51
52
           }
53
           //
                  printf("%d : %d - %d : ", ans, u, v);
54
           if (deep[u] > deep[v])
55
                  swap(u, v);
56
           if (u != v) ans += T.query(id[son[u]], id[v], 1);
57
           return ans;
58 }
5.Trie
   struct Trie {
1
2
   #define sigma size 26
3
           int ch[maxn][sigma size];
           int val[maxn];
4
5
           int sz;
6
7
           void init() {
8
                  sz = 1;
                  memset(ch[0], 0, sizeof ch[0]);
9
10
           }
```

```
11
12
                             int idx(char c) {
13
                                                return c - 'a';
14
                             }
15
                            void insert(char *s, int v) {
16
                                                int u = 0, len = strlen(s);
17
18
                                                rep(i, 0, len) {
                                                                    int c = idx(s[i]);
19
20
                                                                    if (!ch[u][c]) {
21
                                                                                       memset(ch[sz], 0, sizeof ch[sz]);
22
                                                                                        val[sz] = 0;
23
                                                                                        ch[u][c] = sz++;
24
                                                                    }
25
                                                                    u = ch[u][c];
26
27
                                                val[u] = v;
28
29 #undef sigma size
30 };
 6.Treap
           class Treap{
1
2
           public:
3
                      struct treap point{
4
                                int 1,r,x,rep,son,rnd;
5
                      }t[N];
6
                      int treesz,root;
7
                      inline void update(int k){t[k].son=t[t[k].l].son+t[t[k].r].son+t[k].rep;}
                      inline void right_rotate(int &k){int
8
9
           tt=t[k].1;t[k].1=t[tt].r;t[tt].r=k;update(k);update(tt);k=tt;}
                      inline void left rotate(int &k){int
10
11
           tt=t[k].r;t[k].r=t[tt].l;t[tt].l=k;update(k);update(tt);k=tt;}
12
                      inline void insert to tree(int &k,int x)
13
                      {
14
                                if
            (!k)\{k=++\text{treesz};t[k].x=x;t[k].rnd=rand();t[k].son=t[k].rep=1;t[k].l=t[k].r=0;ret[k].rep=1;t[k].l=t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep=1;t[k].rep
15
16
                                t[k].son++;
17
                                if (x==t[k].x){t[k].rep++;return;}
18
                                if (x<t[k].x){insert to tree(t[k].1,x);if</pre>
19
            (t[t[k].1].rnd<t[k].rnd)right rotate(k);}</pre>
20
                                if (x>t[k].x){insert to tree(t[k].r,x);if
21
            (t[t[k].r].rnd<t[k].rnd)left_rotate(k);}</pre>
22
23
                      inline bool delete from tree(int &k,int x)
24
25
                                if (!k)return 0;
26
                                if (x==t[k].x)
27
                                {
```

```
28
                if (t[k].rep>1){t[k].rep--;t[k].son--;return 1;}
29
                if (!t[k].1||!t[k].r){k=t[k].1+t[k].r;return 1;}
30
                if (t[t[k].1].rnd<t[t[k].r].rnd){right_rotate(k);return</pre>
31
    delete_from_tree(k,x);}
32
                left rotate(k);return delete from tree(k,x);
33
            }
34
            int res;
35
            if (x<t[k].x){res=delete_from_tree(t[k].1,x);if (res)t[k].son--;return re</pre>
            res=delete_from_tree(t[k].r,x);if (res)t[k].son--;return res;
36
37
        inline int get succ in tree(int k,int x)
38
39
40
            if (!k)return -1;
            int sv;
41
42
            if (x<t[k].x){sv=get_succ_in_tree(t[k].l,x);return sv==-1?t[k].x:sv;}</pre>
43
            return get_succ_in_tree(t[k].r,x);
44
45
        inline int get_pred_in_tree(int k,int x)
46
47
            if(!k)return -1;
48
            int sv;
49
            if (x>t[k].x){sv=get_pred_in_tree(t[k].r,x);return sv==-1?t[k].x:sv;}
            return get_pred_in_tree(t[k].1,x);
50
51
        }
52
        inline int ask_kth_in_tree(int k,int x)
53
        {
54
            if (!k)return -1;
55
            if (x<=t[t[k].1].son)return ask_kth_in_tree(t[k].1,x);</pre>
56
            if (x>t[t[k].1].son+t[k].rep)return ask_kth_in_tree(t[k].r,x-t[t[k].1].so
57
    t[k].rep);
58
            return t[k].x;
59
        }
        inline int ask rank in tree(int k,int x)
60
61
        {
62
            if (!k)return -1;
63
            if (x==t[k].x)return t[t[k].1].son+1;
            if (x<t[k].x)return ask rank in tree(t[k].1,x);</pre>
64
65
            return t[t[k].1].son+t[k].rep+ask_rank_in_tree(t[k].r,x);
        }
66
67
        /*
68
        inline void bianli(int k)//for debug
69
70
            if (!k)return;
71
            bianli(t[k].1);
72
            printf("i=%d l=%d r=%d x=%d rep=%d
73
    son=%d\n'',k,t[k].1,t[k].r,t[k].x,t[k].rep,t[k].son);
74
            bianli(t[k].r);
75
        }
76
```

```
77
       inline void put(){puts("\n------
    ");bianli(root);puts("-----\n");}//for debug
78
79
       */
80
       inline void init(){treesz=root=0;}
81
       inline void insert(int x){insert to tree(root,x);}
       inline int del(int x){return delete_from_tree(root,x);}
82
       inline int succ(int x){return get succ in tree(root,x);}
83
       inline int pred(int x){return get pred in tree(root,x);}
84
85
       inline int ask_kth(int x){return ask_kth_in_tree(root,x);}
86
       inline int ask rank(int x){return ask rank in tree(root,x);}
87
    }T;
    int main()
88
89
    {
       int n,x,op;read(n);
90
       T.init();
91
92
       for (int i=1;i<=n;i++)</pre>
93
       {
94
           read(op);read(x);
           if (op==1)T.insert(x);
95
96
           if (op==2)T.del(x);
97
           if (op==3)printf("%d\n",T.ask rank(x));
98
           if (op==4)printf("%d\n",T.ask_kth(x));
           if (op==5)printf("%d\n",T.pred(x));
99
           if (op==6)printf("%d\n",T.succ(x));
100
101
       }
    };
```

#### 7.后缀数组

①Da

②Dc3

## 四、图论

#### 1.最短路

```
①Floyd
   #define inf 0x3fffffff
                             //注意,太大会溢出
1
                                    //最大点数
2
   #define M
3
   int n, dist[M][M];
                             //n: 实际点数
4
                             //有时候需要初始化
5
   void init ()
6
7
      int i, j;
      for (i = 1; i <= n; i++)
8
9
          for (j = i + 1; j \le n; j++)
10
             dist[i][j] = dist[j][i] = inf;
```

```
11 }
12
13 void floyd ()
14 {
15
       int i, j, k;
       for (k = 1; k <= n; k++)
16
17
           for (i = 1; i <= n; i++)
18
              for (j = 1; j <= n; j++)
                                         //有的题目会溢出就要自己变通了
19
                  if (dist[i][k] + dist[k][j] < dist[i][j])</pre>
20
                             dist[i][j] = dist[i][k] + dist[k][j];
21 }
  ②Dijkstra
    普通:
   #define inf 0x3fffffff
1
2
   #define M 105
   int dist[M], map[M][M], n;
3
4
   bool mark[M];
   void init () {
5
          int i, j;
6
7
          for (i = 1; i \le n; i++) //i==j 的时候也可以初始化为 0,只是有时候不合适
8
                  for (j = 1; j <= n; j++)
9
                         map[i][j] = inf;
10
   }
   void dijkstra (int u) {
11
12
          int i, j, mins, v;
          for (i = 1; i <= n; i++) {
13
                 dist[i] = map[u][i];
14
15
                 mark[i] = false;
16
          }
17
          mark[u] = true;
                          //既然上面的 map 当 i==j 时不是 0,就要这句
          dist[u] = 0;
18
19
          while (1) {
20
                 mins = inf;
21
                  for (j = 1; j <= n; j++)
22
                         if (!mark[j] && dist[j] < mins)</pre>
23
                                mins = dist[j], v = j;
24
                  if (mins == inf)
25
                         break;
                 mark[v] = true;
26
27
                  for (j = 1; j <= n; j++)
28
                         if (!mark[j] && dist[v] + map[v][j] < dist[j])</pre>
                                dist[j] = dist[v] + map[v][j];
29
30
          }
31 }
     堆优化:
   struct node {
2
          int to, val;
```

```
3
          friend bool operator< (const node &a, const node &b) {
4
                  return a.val > b.val;
5
          }
6
   };
7
   vector<node> g[maxn];
   int dis[maxn];
8
   int dijkstra(int s, int t, int n) {
9
          fill(dis, dis + n + 1, INF);
10
          priority_queue<node> q;
11
12
          dis[s] = 0;
13
          q.push({s, 0});
          while (!q.empty()) {
14
                  int u = q.top().to; q.pop();
15
                  int sz = SZ(g[u]);
16
17
                  rep(i, 0, sz) {
18
                         int v = g[u][i].to, w = g[u][i].val;
19
                         if (dis[v] > dis[u] + w)
20
                                q.push(\{v, dis[v] = dis[u] + w\});
21
                  }
22
23
          return dis[t];
24 }
  3Spfa
     普通队列:
     双端队列:
2.二分图匹配
  ①匈牙利算法
  ● 链表:
   vector<int> g[maxn];
1
   int match[maxn];
2
   bool vis[maxn];
3
4
   int n;
5
   bool dfs(int u) {
          vis[u] = 1;
6
7
          int sz = g[u].size();
8
          rep(i, 0, sz) {
9
                  int v = g[u][i];
                  if (match[v] == -1 || !vis[v] && dfs(v)) {
10
11
                         match[v] = u;
12
                         match[u] = v;
13
                         return 1;
14
                  }
15
          }
16
          return 0;
17
   int bipartite_matching() {
18
19
          int ans = 0;
```

```
memset(match, -1, sizeof match);
20
21
          rep(u, 1, n + 1) {
22
                 if (match[u] == -1) {
23
                        memset(vis, 0, sizeof vis);
24
                         if (dfs(u))
25
                                ans++;
26
                 }
27
          }
28
          return ans;
29 }
     邻接矩阵:
   int g[N][N], match[N];
1
   bool vis[N];
2
3
   int n, m; // n 个点 m 个边
4
   bool dfs(int u) {
5
          rep(v, 1, n + 1) {
6
                 if (g[u][v] && !vis[v]) {
7
                        vis[v] = 1;
                         if (match[v] == -1 || dfs(match[v])) {
8
9
                                match[v] = u;
10
                                return 1;
11
                         }
12
                 }
13
          }
14
          return 0;
15
   }
16
   int hungary() {
17
          int ans = 0;
          memset(match, -1, sizeof match);
18
19
          rep(u, 1, m + 1) {
                 memset(vis, 0, sizeof vis);
20
                 if (dfs(u))
21
22
                        ans++;
23
          }
24
          return ans;
25 }
3.KM——计算最大权匹配
   int nx, ny; //两边的点数
1
2
   int g[N][N];//二分图描述
3
   int match[N], lx[N], ly[N]; //y 中各点匹配状态, x,y 中的点标号
   int slack[N];
4
5
   bool visx[N], visy[N];
6
7
   bool DFS(int x) {
8
          visx[x] = true;
9
          for(int y = 0; y < ny; y++) {
10
                 if(visy[y])continue;
```

```
11
                   int tmp = 1x[x] + 1y[y] - g[x][y];
12
                   if(tmp == 0) {
13
                           visy[y] = true;
14
                           if(match[y] == -1 \mid\mid DFS(match[y])) {
15
                                   match[y] = x;
16
                                   return true;
17
                           }
                   } else if(slack[y] > tmp)
18
19
                           slack[y] = tmp;
20
21
           return false;
22
   }
23
   int KM() {
           memset(match, -1, sizeof(match));
24
25
           memset(ly, 0, sizeof(ly));
26
           for(int i = 0; i < nx; i++) {</pre>
27
                   lx[i] = -INF;
                   for(int j = 0; j < ny; j++)</pre>
28
29
                           if(g[i][j] > lx[i])
30
                                   lx[i] = g[i][j];
31
           }
           for(int x = 0; x < nx; x++) {
32
33
                   for(int i = 0; i < ny; i++)</pre>
34
                           slack[i] = INF;
35
                   while(true) {
                           memset(visx, false, sizeof(visx));
36
37
                           memset(visy, false, sizeof(visy));
38
                           if(DFS(x))break;
39
                           int d = INF;
                           for(int i = 0; i < ny; i++)</pre>
40
41
                                   if(!visy[i] && d > slack[i])
42
                                          d = slack[i];
43
                           for(int i = 0; i < nx; i++)</pre>
44
                                   if(visx[i])
45
                                          lx[i] -= d;
                           for(int i = 0; i < ny; i++) {</pre>
46
47
                                   if(visy[i])ly[i] += d;
48
                                   else slack[i] -= d;
49
                           }
50
                   }
51
52
           int res = 0;
53
           for(int i = 0; i < ny; i++)</pre>
54
                   if(match[i] != -1)
55
                           res += g[match[i]][i];
56
           return res;
57 }
```

#### 4.网络流

```
(1)Dinic
   struct Edge {
1
2
           int to, cap, next;
3
   } E[maxn];
   int head[N], pa[N], vis[N], cnt;
4
5
   void init() {
6
           memset(head, -1, sizeof head);
7
           cnt = 0;
8
   }
9
   void addedge(int u, int v, int w) {
10
           E[cnt].to = v; E[cnt].cap = w; E[cnt].next = head[u]; head[u] =
   cnt++;
11
12
           E[cnt].to = u; E[cnt].cap = 0; E[cnt].next = head[v]; head[v] =
13
   cnt++; // 0 or w
14
   }
15
   bool bfs(int s, int t) {
           memset(vis, -1, sizeof vis);
16
17
           queue<int> q;
           vis[s] = 0;
18
19
           q.push(s);
           while (!q.empty()) {
20
21
                  int u = q.front();
22
                  q.pop();
23
                  for (int i = head[u]; i != -1; i = E[i].next) {
24
                          int v = E[i].to;
25
                          if (E[i].cap && vis[v] == -1) {
26
                                 vis[v] = vis[u] + 1;
27
                                 q.push(v);
28
                          }
29
                  }
30
31
           return vis[t] != -1;
32
33
   int dfs(int u, int t, int flow) {
34
           if (u == t) return flow;
35
           for (int &i = pa[u]; i != -1; i = E[i].next) {
36
                  int v = E[i].to;
37
                  if (E[i].cap && vis[v] == vis[u] + 1) {
38
                          int res = dfs(v, t, min(flow, E[i].cap));
39
                          if (res) {
40
                                 E[i].cap -= res;
41
                                 E[i ^ 1].cap += res;
42
                                 return res;
43
                          }
44
                  }
45
           }
46
           return 0;
```

```
47
   }
48
   int Dinic(int s, int t) {
49
           int max_flow = 0;
           while (bfs(s, t)) {
50
                  memcpy(pa, head, sizeof head);
51
52
                  int res;
53
                  do {
54
                          res = dfs(s, t, INF);
55
                          max_flow += res;
56
                  } while (res);
57
           }
58
           return max_flow;
59
   }
60
  ②ISAP
1
    struct Edge {
2
            int from, to, cap, flow;
3
    };
4
    class ISAP {
5
    public:
    #define max V 202
6
7
            int n, m, s, t;
8
            vector<Edge> edges;
9
            vector<int> g[max_V];
            bool vis[max_V];
10
            int d[max_V], cur[max_V], p[max_V], num[max_V];
11
12
13
            void Addedge(int u, int v, int w) {
14
                   edges.pb((Edge) {
                           u, v, w, 0
15
16
                   });
17
                   edges.pb((Edge) {
18
                           v, u, 0, 0
19
                   });
20
                   m = SZ(edges);
21
                   g[u].pb(m - 2);
22
                   g[v].pb(m - 1);
23
            }
24
25
            bool Bfs() {
26
                   memset(vis, 0, sizeof vis);
27
                   queue<int> q;
28
                   q.push(t);
29
                   vis[t] = 1;
30
                   d[t] = 0;
                   while (!q.empty()) {
31
32
                           int u = q.front(); q.pop();
33
                           int sz = SZ(g[u]);
34
                           rep(i, 0, sz) {
```

```
Edge &E = edges[g[u][i] ^1];
35
36
                                   if (!vis[E.from] && E.cap > E.flow) {
37
                                          vis[E.from] = 1;
38
                                          d[E.from] = d[u] + 1;
39
                                          q.push(E.from);
                                  }
40
                           }
41
42
                   }
43
                   return vis[s];
44
            }
45
            void ClearAll(int n) {
46
47
                   this->n = n;
                   rep(i, 0, n) g[i].clear();
48
49
                   edges.clear();
50
            }
51
52
            void ClearFlow() {
53
                   int sz = SZ(edges);
54
                   rep(i, 0, sz) edges[i].flow = 0;
55
            }
56
            int Augment() {
57
                   int x = t, a = INF;
58
59
                   while (x != s) {
60
                           Edge &E = edges[p[x]];
61
                           a = min(a, E.cap - E.flow);
62
                           x = edges[p[x]].from;
63
                   }
64
                   x = t;
65
                   while (x != s) {
                           edges[p[x]].flow += a;
66
67
                           edges[p[x] ^ 1].flow -= a;
68
                           x = edges[p[x]].from;
69
                   }
70
                   return a;
71
            }
72
73
            int MaxFlow(int s, int t) {
74
                   this->s = s; this->t = t;
75
                   int flow = 0;
76
                   Bfs();
77
                   memset(num, 0, sizeof num);
78
                   rep(i, 0, n) num[d[i]]++;
79
                   int x = s;
                   memset(cur, 0, sizeof cur);
80
81
                   while (d[s] < n) {
82
                           if (x == t) {
83
                                  flow += Augment();
```

```
84
                                   x = s;
85
                           }
86
                           bool ok = 0;
87
                           int sz = SZ(g[x]);
                           rep(i, cur[x], sz) {
88
89
                                   Edge &E = edges[g[x][i]];
90
                                   if (E.cap > E.flow && d[x] == d[E.to] + 1) {
91
                                           ok = 1;
92
                                           p[E.to] = g[x][i];
93
                                           cur[x] = i;
94
                                           x = E.to;
95
                                           break;
96
                                   }
97
                           }
                           if (!ok) {
98
99
                                   int m = n - 1;
100
                                   rep(i, 0, sz) {
101
                                           Edge &E = edges[g[x][i]];
102
                                           if (E.cap > E.flow)
103
                                                  m = min(m, d[E.to]);
104
                                   }
105
                                   if (--num[d[x]] == 0) break; // gap
106
                                   num[d[x] = m + 1]++;
                                   cur[x] = 0;
107
108
                                   if (x != s) x = edges[p[x]].from;
109
                           }
110
                    }
111
                    return flow;
112
            }
113
114
            vector<int> Mincut() {
115
                    Bfs();
116
                    vector<int> ans;
117
                    int sz = SZ(edges);
118
                    rep(i, 0, sz) {
119
                           Edge &E = edges[i];
120
                           if (!vis[E.from] && vis[E.to] && E.cap > 0)
121
                                   ans.pb(i);
122
                    }
123
                    return ans;
124
            }
125
126
            void Reduce() {
127
                    int sz = SZ(edges);
128
                    rep(i, 0, sz) {
129
                           edges[i].cap -= edges[i].flow;
130
                    }
131
            }
132
```

```
void Print() {
133
134
                   puts("Graph:");
                   int sz = SZ(edges);
135
136
                   rep(i, 0, sz) {
137
                          printf("%d->%d, %d, %d\n", edges[i].from,
    edges[i].to, edges[i].cap, edges[i].flow);
138
139
                   }
140
            }
141
   #undef max V
142
    };
5.费用流
   struct Edge {
1
2
           int from, to, cap, flow, cost;
           Edge(int a, int b, int c, int d, int E) { from = a; to = b; cap = c;
3
4
   flow = d; cost = E; }
5
   };
6
   class MCMF {
7
   public:
   #define max V 202
8
9
           int n, m, s, t;
           vector<Edge> edges;
10
11
           vector<int> g[max_V];
12
           bool vis[max V];
           int d[max_V], p[max_V], a[max_V];
13
14
15
           void Init(int n) {
16
                  this->n = n;
17
                  rep(i, 0, n) g[i].clear();
                  edges.clear();
18
19
           }
20
21
           void Addedge(int u, int v, int c, int w) {
22
                  edges.pb(Edge(u, v, c, 0, w));
23
                  edges.pb(Edge(v, u, 0, 0, -w));
24
                  int m = SZ(edges);
25
                  g[u].pb(m - 2);
26
                  g[v].pb(m - 1);
27
           }
28
29
           bool Spfa(int s, int t, int &flow, int &cost) {
30
                  fill(d, d + n, INF);
31
                  memset(vis, 0, sizeof vis);
32
                  d[s] = 0; vis[s] = 1; p[s] = 0; a[s] = INF;
33
                  queue<int> q;
34
                  q.push(s);
35
                  while (!q.empty()) {
36
                         int u = q.front(); q.pop();
37
                         vis[u] = 0;
```

```
int sz = SZ(g[u]);
38
39
                          rep(i, 0, sz) {
40
                                 Edge& E = edges[g[u][i]];
                                 if (E.cap > E.flow && d[E.to] > d[u] + E.cost)
41
42
   {
                                         d[E.to] = d[u] + E.cost;
43
                                         p[E.to] = g[u][i];
44
45
                                         a[E.to] = min(a[u], E.cap - E.flow);
46
                                         if (!vis[E.to]) {
47
                                                vis[E.to] = 1;
48
                                                q.push(E.to);
49
                                         }
50
                                 }
                          }
51
52
53
                  if (d[t] == INF) return 0;
54
                  flow += a[t];
55
                  cost += d[t] * a[t];
                  int u = t;
56
57
                  while (u != s) {
58
                          edges[p[u]].flow += a[t];
59
                          edges[p[u] ^ 1].flow -= a[t];
60
                          u = edges[p[u]].from;
61
62
                  return 1;
63
           }
64
           int MincostMaxflow(int s, int t) {
65
                  int flow = 0, cost = 0;
66
67
                  while (Spfa(s, t, flow, cost));
68
                  return cost;
69
           }
70 #undef max V
71 };
```

# 五、数论

#### 1.质因子分解

```
Miller Rabin and Pollard Rho
   //质因子分解,小数据用筛法直接判,大数据用 pollard rho
1
   //map[i]是含有多少个质因子i
2
3
   //map<LL,int>::iterator c,c->first 表示质因子, c->second 表示次方
4
5
   #define maxn for division 1000001
   bool is prime[maxn for division];
6
7
   vector<int>prime;
8
   map<LL, int>factor;
```

```
inline void get_prime() {
9
10
            for(int i = 0; i < maxn_for_division; i++)is_prime[i] = 1;</pre>
            is prime[0] = is prime[1] = 0;
11
12
            for(int i = 2; i < maxn_for_division; i++)</pre>
13
                   if (is prime[i]) {
14
                           prime.push_back(i);
15
                           for (int j = i; j < maxn for division; j +=</pre>
    i)is_prime[j] = 0;
16
17
                   }
18
    }
19
    inline LL gcd(LL a, LL b) {
            if (!b)return a;
20
21
            return gcd(b, a % b);
22
    }
23
    inline LL mod mul(LL a, LL b, LL p) {
24
            LL ans = 011;
25
            a %= p, b %= p;
            if (a > b)swap(a, b);
26
            while (b) {
27
28
                   if (b \& 1)ans = (ans + a) \% p;
29
                   a = (a + a) \% p;
30
                   b >>= 1;
31
            }
32
            return ans;
33
34
    LL mod pow(LL a, LL b, LL p) {
35
            LL ans = 111;
            a %= p;
36
37
            while (b) {
                   if (b & 1)ans = mod_mul(ans, a, p);
38
39
                   a = mod mul(a, a, p);
                   b >>= 1;
40
41
42
            return ans;
43
    bool witness(LL a, LL n) {
44
45
            LL m = n - 1;
46
            int j = 0;
47
            while(!(m & 1))j++, m >>= 1;
            LL x = mod pow(a, m, n);
48
49
            if (x == 1 || x == n - 1)return 0;
50
            while(j--) {
51
                   x = mod_mul(x, x, n);
52
                   if(x == n - 1)return 0;
53
            }
54
            return 1;
55
56
    #define rep_times 20
57
    bool Miller Rabin(LL n) { //判断 n 是否为素数
```

```
58
            srand(time(0));
59
            if(n < 2)return 0;</pre>
            if(n == 2)return 1;
60
            if (!(n & 1))return 0;
61
            for(int i = 0; i < rep_times; i++) {</pre>
62
                   LL a = rand() \% (n - 1) + 1;
63
                   if (witness(a, n))return 0;
64
65
            }
66
            return 1;
67
68
    #undef rep times
69
    LL Pollard Rho(LL n, int c) {
70
            LL x = 2, y = 2, d = 1;
71
            while (d == 1) {
72
                   x = mod_mul(x, x, n) + c;
73
                   y = mod_mul(y, y, n) + c;
74
                   y = mod_mul(y, y, n) + c;
75
                   d = gcd((x - y >= 0 ? x - y : y - x), n);
76
77
            if (d == n)return Pollard Rho(n, c + 1);
78
            return d;
79
    }
80
    bool Is_Prime(LL n) {
            return n < maxn for division && is prime[n] || n >=
81
82
    maxn for division && Miller Rabin(n);
83
84
    void Find_Factor(LL n) {
85
            if (Is Prime(n)) {
86
                   factor[n]++;
87
                   return;
88
            for (int i = 0; i < prime.size() && prime[i] <= n; i++)</pre>
89
90
                   if (n % prime[i] == 0) {
91
                           while (n % prime[i] == 0) {
                                  factor[prime[i]]++;
92
93
                                   n /= prime[i];
                           }
94
95
                   }
            if (n != 1) {
96
97
                   if (Is Prime(n))factor[n]++;
98
                   else {
99
                           LL p = Pollard Rho(n, 1);
100
                           Find_Factor(p);
101
                           Find_Factor(n / p);
102
                   }
103
            }
104
    }
105
    int main() {
106
            LL n;
```

```
get_prime();
107
         while(read(n), !IOerror) {
108
109
               factor.clear();
110
               Find_Factor(n);
111
               //----
                                 -----output-----
112
   -----
               for(map<LL, int>::iterator c = factor.begin(); c !=
113
114 factor.end();) {
                     printf("%lld^%d", c->first, c->second);
115
                     if((++c) != factor.end())printf("*");
116
117
               }
               //-----
118
119 -----
120
               puts("");
121
122 }
```

# 六、代数

# 七、组合数学