

Homework 29

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The general outline for this algorithm is as follows:

At the first time step, use $n - 1$ processors to write the numbers $1, 2, 3, 4, \dots, n - 1$ in memory location M . These represent values for k such that there exists a matching prefix and suffix of size k .

At the second time step, use n^2 processors to compare the prefixes and suffixes of all lengths for k from 1 to $n - 1$. If the prefix and suffix aren't equal to each other, then "zero out" the location in memory. For example, if the first 3 characters do not match the last 3 characters, then the numbers in M become $1, 2, 0, 4, \dots, n - 1$.

Finally, at the third time step, use n^2 processors to find the maximum number left in M . This will return the maximum valid k , and it can be done in constant time as discussed in class.

Below is an algorithm for one of the $(n - 1)^2$ processors, i, j . (Without loss of generality, the processors are labelled with two numbers, $i \in [1 \dots n - 1]$ and $j \in [1 \dots n - 1]$ for easier usage).

Algorithm 1 CRCW Common $O(1)$ algorithm

Require: A string C of size n , a processor $p_{i,j}$, a memory location M of size $n - 1$, a memory location T of size $(n - 1)^2$, and a memory location And of size $n - 1$.

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 $M[i] \leftarrow i$  ▷ First, copy the numbers  $1, 2, 3, \dots, n - 1$  into  $M$ .
if  $C[j] \neq C[n - i + j]$  then
     $M[i] \leftarrow 0$  ▷ If any of the pairs of characters don't match, then that  $k$  isn't viable.
end if
if  $M[i] \geq M[j]$  then ▷ Perform all possible pairwise comparisons of  $M$ .
     $T[i, j] \leftarrow 1$  ▷ Record which indices are greater.
else
     $T[i, j] \leftarrow 0$ 
end if
 $And[i] \leftarrow 1$  ▷ Perform an EREW AND operation to find a row in  $T$  of all 1s.
if  $T[i, j] = 0$  then
     $And[i] \leftarrow 0$ 
end if
if  $And[i] = 1$  then
    Output  $i$  ▷ That row is the maximum  $k$ . A maximum always exists, so this will always output.
end if
```

This algorithm runs in $O(1)$ time, since each processor only performs a constant number of operations, as described above.

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