

# Homework 15

Brian Knotten, Brett Schreiber, Brian Falkenstein

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## 20

A decision tree is as follows: The decision tree can be thought of as follows for each node  $v$  at depth  $i$ :

- $v$  designates each of  $p_1, p_2, \dots, p_i$  into lists  $A$  and  $B$ , which represent the routes that either taxi  $A$  or taxi  $B$  travels.
- $v$ 's left child copies the assignments of  $v$  to  $A$  and  $B$ , but appends point  $p_{i+1}$  to  $A$ .
- $v$ 's right child copies the assignments of  $v$  to  $A$  and  $B$ , but appends point  $p_{i+1}$  to  $B$ .
- At depth  $i$ ,  $i$  points have been visited.

Because each point needs to be visited by at least one taxi, and since there are no benefits for having both taxis visit a point, each point is either visited by taxi  $A$  or by taxi  $B$ . So there are  $2^n$  possible point designations at the leaves of the decision tree.

The decision tree can be pruned using the following rules:

1. If two nodes  $u$  and  $v$  have the same total distance,  $\text{dist}(u_A) + \text{dist}(u_B) = \text{dist}(v_A) + \text{dist}(v_B)$ , prune  $v$ .

```
taxi(i, a, b):
    if i > n:
        return 0

    return min(
        dist(p[a], p[i]) + taxi(i + 1, i, b)
        dist(p[b], p[i]) + taxi(i + 1, a, i)
    )
```

Let  $\text{taxi}[i, a, b]$  = the minimum total distance of designating  $A$  and  $B$  over points  $p_1 \dots p_i$ , and  $A$ 's last stop is  $p_a$ , and  $B$ 's last stop is  $p_b$ . The dynamic solution algorithm is:

```
for a = 1 to n do:
    for b = 1 to n do:
        for i = n to 1 do:

            taxi[a, b, i] = min(
                taxi[i, b, i + 1] + dist(p[a], p[i]),
                taxi[a, i, i + 1] + dist(p[b], p[i])
            )
```

## 21

The decision tree can be defined as:

- Every node  $v$  contains the cumulative response time up until that point in the path, and the path taken to get to  $v$ .
- Every node having at most  $n$  children, where each child represents the next point to travel to

- The leaves being paths covering all  $n$  points
- The solution being the minimum response time found at the leaves

Without pruning, the size of the tree will be  $n^n$  (height of  $n$  since every path must cover all points, and a branching factor of  $n$ ). The tree can then be pruned using the following rules:

1. If a node visits a point in its path that it has already visited, prune it.
2. If 2 nodes at the same depth have visited the same points in its path, prune the one with the longer cumulative response time.