

# Homework 14

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### a

Consider that every Boolean circuit can be represented by AND, OR, and NOT gates.

Consider a Boolean circuit  $C$  containing  $n$  gates. There exists a corresponding Boolean formula  $F$  of size  $n$ , because AND, OR, and NOT gates can be represented in a Boolean formula as logical operators  $\wedge$ 's,  $\vee$ 's, and  $\neg$ 's. Since there is a one-to-one mapping of gates to logical operators, there exists a Boolean formula  $F$  with an output equivalent to  $C$  of size  $n$ .

Any Boolean function  $F$  for an input size  $n$  can be encoded into a string  $s$  of size  $2^n$ , where each bit represents an output for a given input. For example, if  $n = 4$ , the first bit of  $s$ , either a 0 or 1, represents the output for the input 0000. The second bit of  $s$  represents the output for input 0001... and the last bit of  $s$  represents the output for input 1111.

The formula for  $F$  can be naively implemented through the following algorithm:

List all binary numbers of length  $n$  from 0 to  $2^n$ . Let  $B$  be this list. For each number  $b$  in  $B$ :

for each bit  $b_i$  in  $b$ :

if  $b_i = 0$ , write out  $\neg x_i$

if  $b_i = 1$ , write out  $x_i$

Join together all of these boolean expressions with  $\wedge$ 's

Join together all of these boolean expressions with  $\vee$ 's

Here is an example Boolean function  $f$  encoded as 1011. This means  $f$  has the following outputs:

$$f(00) = 1$$

$$f(01) = 0$$

$$f(10) = 1$$

$$f(11) = 1$$

The corresponding boolean formula using the naive implementation is:  $f(x_1x_2) = (\neg x_1 \wedge \neg x_2) \vee \neg(\neg x_1 \wedge x_2) \vee (x_1 \wedge \neg x_2) \vee (x_1 \wedge x_2)$

So every boolean function on  $n$  bits has a formula of size at most  $2^n$ . Since  $\wedge$ ,  $\vee$ , and  $\neg$  exist as gates in circuits, this formula can be converted into a circuit with size at most  $2^n$  gates.

### b

First, show that  $f$  computed by circuit  $|C| = S \rightarrow f$  computed by  $S$ -line program:

For each Boolean gate (NOT, AND, OR) in  $C$ , there exists an equivalent straight-line program statement with a left-side assignment variable corresponding to the output of the gate and a right side consisting of either a boolean operation ( $\wedge$ ,  $\vee$ ) performing the gate's operation on input variables corresponding to the input of the gate, or the negation ( $\neg$ ) of an input variable corresponding to the input of the gate.

Therefore, given a circuit  $|C| = S$ , we can construct an equivalent straight-line program with  $S$  statements by iterating over each gate in  $C$  and converting it to an equivalent straight-line program statement.

Next, show that  $f$  computed by  $S$ -line program  $\rightarrow f$  computed by circuit  $|C| = S$ :

For every straight-line program statement there exists an equivalent Boolean gate (NOT, AND, OR) with inputs corresponding to the right-side variable(s) of the statement that performs the corresponding Boolean operation ( $\neg$ ,  $\wedge$ ,  $\vee$ ) of the right-side of the statement and an output corresponding to the left-side assignment variable of the statement.

Therefore, given a straight-line program with  $S$  statements, we can construct an equivalent circuit  $|C| = S$  by iterating over each statement and converting it to an equivalent Boolean gate.

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