

# Homework 11

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## 16

Initial Approach: Construct a TM  $U$  which takes an input  $x$ , runs  $T(x)$ , and then  $S(T(x))$  and returns the result.  $U$  correctly accepts  $x$  if and only if  $x \in A$ , however the space needed to store the temporary result  $T(x)$  as input to  $S$  could be more than  $\log(|x|)$  in space.

New Approach: Instead modify  $T$  to produce output for  $S$  on-demand. That is, have  $U$  use a modified  $T'$  which takes a string  $x$  and int  $i$  and only writes the  $i^{\text{th}}$  symbol of  $x$  to the output tape. Adding a counter to  $T$  to track this space only requires an additional  $\log(|x|)$  space. Next, have  $U$  use a modified  $S'$  which runs  $T'(x, i)$  when it needs the  $i^{\text{th}}$  symbol.

The new approach will have a TM  $U$  which runs  $T'$  and  $S'$  which are both  $\log(|x|)$  in space, so  $U$  is  $\log(|x|)$  in space while also accepting  $A$ .

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### a

Let  $C = \left\{ (M, I, 2^{|I|^k}) \mid \text{TM } M \text{ accepts } I \text{ in at most } 2^{|I|^k} \text{ space} \right\}$ .

First we will show that  $C \in \text{EXPSPACE}$ :

Construct a TM  $T$  that takes in the tuple  $(M, I, 2^{I^k})$ , reads  $I$ , and runs  $M$  on  $I$  with a tape of size  $2^{|I|^k}$ .

$T$  accepts if  $M$  accepts using less than  $2^{|I|^k}$  space and rejects otherwise.

Because the TM  $M$  uses exponential space,  $T$  must also have exponential space available and therefore  $C \in \text{EXPSPACE}$ .

Next, we will show that  $\forall C' \in \text{EXPSPACE}, C' \leq_{\text{poly}} C$ .

Let  $C' \in \text{EXPSPACE}$ .

Since  $C' \in \text{EXPSPACE}$ ,  $\exists$  TM  $M$  and  $k \in \mathbb{Z}$  such that  $M(x)$  accepts in  $2^{|x|^k}$  space iff  $x \in C'$

### b

The problem with the previous proof was that the space requirements for an arbitrary  $\langle M \rangle$  could grow faster than  $2^n$ . Consider a TM  $T$  such that  $L(T) = C$ . Since  $C$  passes in an arbitrary base  $c$ ,  $T$  needs to account for that difference in growth rate. But this  $c$  is negligible, because  $c^{n^k}$  can be expressed as  $2^{n^{k+\epsilon}}$  for some  $\epsilon$ . Since the number of steps can be written as  $2^{n^{x+\epsilon}}$ , and since  $T$  can handle any  $2^{n^k}$ , the base does not matter.