

Homework 10

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14

Let J_i be the i^{th} binary-encoded Java program, and MJ_i be the i^{th} binary-encoded mini-Java program.

Define TM D : For an input x , run $MJ_x(x)$ and return the opposite. (Note that all MJ programs halt on all inputs).

The language accepted by D cannot be accepted by any MJ_i by its construction. Any MJ_i will return the opposite of D on input i , for any i .

Need to show that language accepted by D can be accepted by some J_i .

15

a

A TM M can be defined as follows so that $L(M) = A$:

On input x :

 Instantiate $c = 0$ on the working tape.

 for each character $x_i \in x$:

 if $x_i = '('$, then increment c .

 else, if $x_i = ')'$, then decrement c .

 if $c < 0$ reject. (There is a right paren before a left one).

 Accept if and only if the final value of $c = 0$.

$L(M) = A$ because $L(M)$ only accepts when the number of left parentheses matches the number of right parentheses. M runs in logspace because in the worst case, the input x to M will be n number of left parentheses. So the working tape has to count up to n . But by using the standard base-2 binary encoding of n , the working tape will only use a maximum of $\log(n)$ cells.

b

A TM N can be defined as follows so that $L(N) = B$:

Assume N has two work tapes.

On input x :

 First pass: For each character x_i, x_{i+1} in x :

 If $x_i, x_{i+1} = '(,)'$, reject,

 or if $x_i, x_{i+1} = '[, ']'$, reject.

 Set $c_1 = 0$ on the first working tape.

 Set $c_2 = 0$ on the second working tape.

 Second pass: For each character x_i in x :

 If $x = '('$, increment c_1

 or if $x = ')'$, decrement c_1

 or if $x = '['$, increment c_2

 or if $x = ']'$, decrement c_2

If $c_1 < 0$ or $c_2 < 0$, then immediately reject.

If $c_1 = 0$ and $c_2 = 0$, then accept. Otherwise, reject.

N runs in logspace, even though it uses two work tapes, since $2 \log(n) = O(\log(n))$.