

Homework 17

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$$\rho = \frac{1}{\pi}$$

The complete binary representation of $\frac{1}{\pi}$ is infinitely long. Any finite representation is not completely accurate, since $\frac{1}{\pi}$ is irrational.

Since $\frac{1}{\pi}$ is infinitely long, the probability that it contains an encoding of a computation history as a substring is 1.

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Show that $(NP \cup \text{co-NP}) \subseteq PP$.

Proof:

We will show that Boolean Satisfiability $\in PP$. Since Boolean Satisfiability is NP-complete, then any $L \in NP \rightarrow L \in PP$.

Let M be a probabilistic Turing machine which takes a Boolean formula F as input. Then does the following:

1. Pick a random assignment of variables for F
2. Verify if F is satisfied
3. If F is satisfied, then accept
4. If F is not satisfied, accept with $P(1/2)$ or reject with $P(1/2)$

First note that each step can clearly be done in polynomial time. $P(M \text{ accepts}) = 1/2 + P(\text{random assignment is satisfied})$ and since $P(\text{random assignment is satisfied}) > 0$, then $P(M \text{ accepts}) > 1/2$ and $P(M \text{ rejects}) < 1/2$. Thus Boolean Satisfiability $\in PP$ since M is a Turing machine which meets the requirements for PP. So NP is contained in PP.

Since PP is closed under complement and contains NP, it must also contain co-NP. Finally, $(NP \cup \text{co-NP}) \subseteq PP$.