Experiment No. 4

Views are frequently used in computer programming of a typical example is the creation of a job queue by an operating system does not use operating system does not use preorities then gobs are processed in order they enter the system. White Ctt program for stimulating gob queue. White functions to add gob and delete job from queue.

Pre-requiste:
Barics of Queue

Different operations that can be performed on queue

Objectives?
To perform addition & deletion operation on queue.

Toput: 5:20 of queue Element in queue.

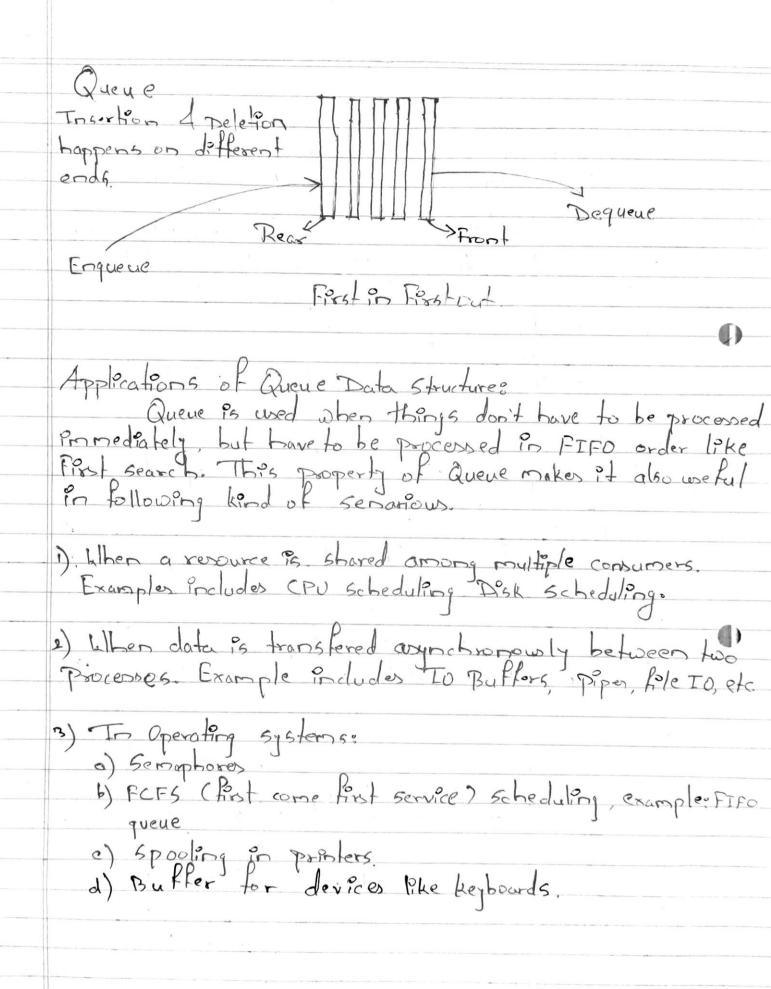
Dutcome:
Result of addition of job operation on queue.
Result of deletion of job operation on queue.

Theory.

A queue is a linear structure which follows a

Porticular order in which operations are performed. The

order in First In First out (FTFO). A good example of a queue is any queue of consumer for a resource where the consumer that came first is served first. The difference between stacks and queue, is in removing. In stack we remove the item the most recently added. In a queue, we remove the Hem the least recently added.



4) In Networks.

a) Queues in souters (switcher)

b) Mail Queues.

5) Voxiations:

Pequeue, Priority Queue Doubly ended priority queue. Boose Operations:

Queue operations may involve initializing or defining
the queue utilizing it of their completely examing it from
memory. Here we shall try to understand the basic
operations associated with queues.

Enqueue() - add (store) an item to queue Dequeuel) - remove (occess) an item from queux. mentioned queue operation efficient. peck() - gets the element at front of queue without removing it.

95 full() - check if queue is full. ?s empty () - check it queue is empty. Enqueue Operation:
Queues maintain two data pointers front 4 rear.
Therefore, its operation are comparitavely difficult to implement

that than of stacks.

(Propert) dater into a queue.
(insert) dater into a queue.
Step II: Check if queue is hell. Step II: If queue is full, produce overflow error if exit. Step III: If queue is not full, increment rear pointer to point the next empty space
step II: Il queue 95 tull, Produce overtow error 4 ent.
Step III: If queue is not full, increment near pointer to point
the next empty space.
Step TV: Add data element to queue location, where year is pointing
The next empty space. Step TV: Add data element to queue location, where year is pointing Step V: Return success.
Real front
Rear Front To B A before
Defore.
Rear - French
4
DIC BIA ofter
Queue Erqueue.
,

, . .

* *

	Dequeue Operation:
	Accessing data from queue is a process of two twoks - access the date where front is pointing and remove the data after access. The following steps are taken to perform dequeue operation.
	remove the data after access. The following stops are
No.	Step II: The queue is empty, produce underflow error & exit. Step III: If queue is empty, produce underflow error & exit. Step III: If queue is not empty, accordate where front
)——	Step III: If queue is not empty, accordate where front
	Step TV: Increment front pointer to point the next available
	Step V: Return success.
	Rear Front
	DCBA
	D
	Rear Front dequeue.
	TD 10 1.B after
	*
	Queue dequeue.
	queue dequeue.

_	
Alposithm of peck () function:	, , , , , , , , , , , , , , , , , , ,
begin procedure Deck	
return queue [front] end procedure	
end procedure	at a series of the
1	
	,
Algorithm of isfull() function	6
begin Procedure 95 hill	and of the second
return true	3
return true	A-5, 18%
else redurn false	to make the time?
return folse	
end procedure.	
\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	
1 ou P	
Algorithm of eisempty () hund	TODO
100 - 1 0 1	
peges procedure isompty	4
of facility design of	las Pia
To be the second	I'OR front is greater than sear
return false	
else return false endif	
end procedure	
\	

•

	Algorithm of enqueue () operations	
	procedure enqueue (dota).	
	return overflow	
	rear = rear +1	
<u></u>	queue (rear) + date	
	end procedure	
	end procedure.	
	Algorithm of dequeuell operation:	
	Johnson of State of Charles.	
	Procedure dequeue	
	if queue is eronty	
	return unterflow	
	endit	
	data = quare (front) front - front + I	
	return true	
	\	
	end procedure	

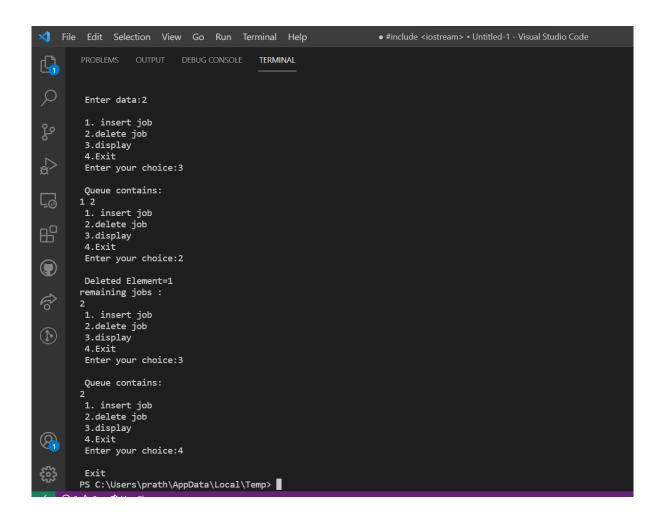
•

Program Code:

```
#include <iostream>
#define MAX 10
using namespace std;
struct queue
     int data[MAX];
       int front, rear;
};
class Queue
   struct queue q;
 public:
   Queue(){q.front=q.rear=-1;}
   int isempty();
   int isfull();
   void enqueue(int);
   int delqueue();
   void display();
};
int Queue::isempty()
       return(q.front==q.rear)?1:0;
int Queue::isfull()
{ return(q.rear==MAX-1)?1:0;}
void Queue::enqueue(int x)
{q.data[++q.rear]=x;}
int Queue::delqueue()
{return q.data[++q.front];}
void Queue::display()
{ int i;
  cout << "\n";
  for(i=q.front+1;i<=q.rear;i++)
          cout<<q.data[i]<<" ";
}
int main()
   Queue obj;
       int ch,x;
             cout<<"\n 1. insert job\n 2.delete job\n 3.display\n 4.Exit\n Enter your
choice:";
            cin>>ch;
       switch(ch)
       { case 1: if (!obj.isfull())
                 { cout<<"\n Enter data:";
                      cin>>x;
                      obj.enqueue(x);
                 }
```

```
else
                   cout << "Queue is overflow";
              break;
         case 2: if(!obj.isempty())
                         cout<<"\n Deleted Element="<<obj.delqueue();</pre>
                 else
                       { cout<<"\n Queue is underflow"; }
                 cout<<"\nremaining jobs :";</pre>
                 obj.display();
              break;
        case 3: if (!obj.isempty())
             { cout<<"\n Queue contains:";
                   obj.display();
            else
                     cout<<"\n Queue is empty";
            break;
        case 4: cout<<"\n Exit";
    }while(ch!=4);
return 0;
```

Program Output:



Conclusion:

By this way we can perform different operations on queue.

8 11 5 - 8 5

The sent face re-

3.7 ° 9.8 ° × × · · •

\$ s