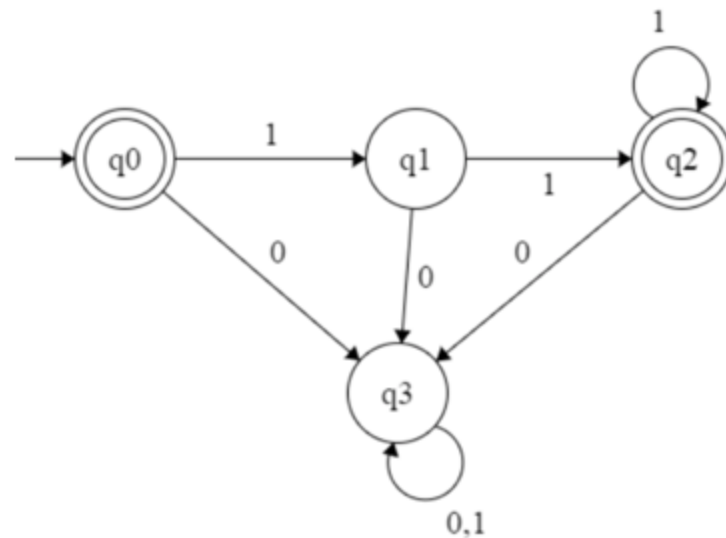


The finite automaton pictured in the figure below represents what language?



- A-) Changes the sign bit ☐
- B-) Increments a given bit pattern by 1 ☐
- C-) Complements a given bit pattern ☐
- D-) None of them ☒
- E-) Finds 2's complement of a given bit pattern ☐

We know that we can find a Turing machine M that decides the language L .

Which one of the following is FALSE?

- A-) None of them. ☐
- B-) M is a recognizer for L . ☐
- C-) M is a decider for L . ☐
- D-) M may loop forever for some strings not in L . ☐
- E-) M does not loop forever for all strings in L . ☐

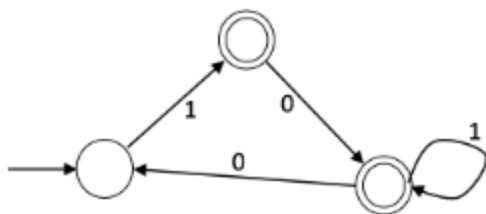
Fill in the blank with the correct option.

_____ Thus, there are languages that are not Turing recognizable.

- A-) Turing recognizable languages does not include Turing enumerable languages. ☐
- B-) Set of all languages is uncountable, but the set of Turing recognizable languages is countable. ☐
- C-) Turing recognizable languages does not include regular languages. ☐
- D-) Set of all languages is not decidable, but the set of Turing recognizable language is decidable. ☐
- E-) All the options can be used to fill in the blank. ☒

Let ϵ be empty string.

Which one of the following context-free grammars belongs to the automaton given below?



$S \rightarrow 1A$

A-)

$A \rightarrow 0B$

$B \rightarrow 1B \mid 0S$

$S \rightarrow 1A$

B-)

$A \rightarrow 0B \mid \epsilon$

$B \rightarrow 1B \mid 0S \mid \epsilon$

$S \rightarrow 1A$

C-)

$A \rightarrow 0B \mid \epsilon$

$B \rightarrow 1B \mid 0S$

D-)

None of them.

$S \rightarrow 1A$

E-)

$A \rightarrow 0B$

$B \rightarrow 1B \mid 0S \mid \epsilon$



Definition: A Turing machine is a 7-tuple given by

$$(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$$

where Q, Σ, Γ are all finite sets and

1. Q denotes the set of states,
2. Σ denotes the input alphabet not containing the blank symbol $\$,$
3. Γ denotes the tape alphabet, where $\$ \in \Gamma$ and $\Sigma \subseteq \Gamma,$
4. $\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$ is the transition function, where L and R denote the moves to the left and right, respectively.
5. $q_0 \in Q$ is the start state,
6. $q_{\text{accept}} \in Q$ is the accept state, and
7. $q_{\text{reject}} \in Q$ is the reject state, where $q_{\text{reject}} \neq q_{\text{accept}}.$

Given the above definition of a standard Turing machine, which one of the following is TRUE?

- A-) A standard Turing machine cannot write the blank symbol on its tape. ☐
- B-) In a standard Turing machine, the tape alphabet can be same as the input alphabet. ☐
- C-) A standard Turing machine can not contain just a single state. ☐
- D-) A standard Turing machine's head cannot be in the same location in successive steps. ☐

Let the alphabet be $\Sigma=\{0,1\}$, and let ε be empty string.

Let A be the language of strings $w \in \Sigma^*$ containing exactly one double symbol. We say that a string has a double symbol if it contains 00 or 11 as a substring. The string 10010 has exactly one double symbol, but 100010 has two double symbols.

Which one of the following is the regular expression for the language A.

A-)

$$(\varepsilon+1)(01)^*00(10)^*(\varepsilon+1) + (\varepsilon+0)(10)^*11(01)^*(\varepsilon+0)$$

☐

B-)

$$(\varepsilon+1)(11)^*00(10)^*(\varepsilon+1) + (\varepsilon+0)(10)^*11(00)^*(\varepsilon+0)$$

☐

C-)

$$(\varepsilon+1)(00)^*00(10)^*(\varepsilon+1) + (\varepsilon+0)(10)^*11(11)^*(\varepsilon+0)$$

☐

D-)

$$(\varepsilon+1)(00)^*01(10)^*(\varepsilon+1) + (\varepsilon+0)(10)^*10(11)^*(\varepsilon+0)$$

☒

Which one of the following statements are TRUE?

I. Every multi-tape Turing machine has an equivalent single-tape Turing machine.

II. Every context-free language is Turing-recognizable.

III. Every language is Turing-recognizable.

- A-)** I, II and III ☐
- B-)** II and III ☐
- C-)** I and III ☐
- D-)** I and II ☒
- E-)** None of them. ☐

Which one of the following are correct?

- A-) The language $\{a^n b^n \mid 0 \leq n \leq 3\}$ is non-regular. ☐
- B-) The class of context-free languages is closed under intersection. ☒
- C-) If A and B are regular languages, then $(A \cup B)^*$ is regular. ☐
- D-) If A is a regular language, then A is finite. ☐
- E-) The language $\{a^n b^n a^n \mid 0 \leq n\}$ is regular. ☐

Let w denote a string.

Which one of the following is not decidable?

A-)

$L = \{ \langle M, w \rangle \mid M \text{ is a non-deterministic finite automaton that accepts } w \}$

☐

B-)

$L = \{ \langle M, w \rangle \mid M \text{ is a deterministic finite automaton that accepts } w \}$

☒

C-)

$L = \{ \langle M, w \rangle \mid M \text{ is a push down automaton that accepts } w \}$

☐

D-)

None of them.

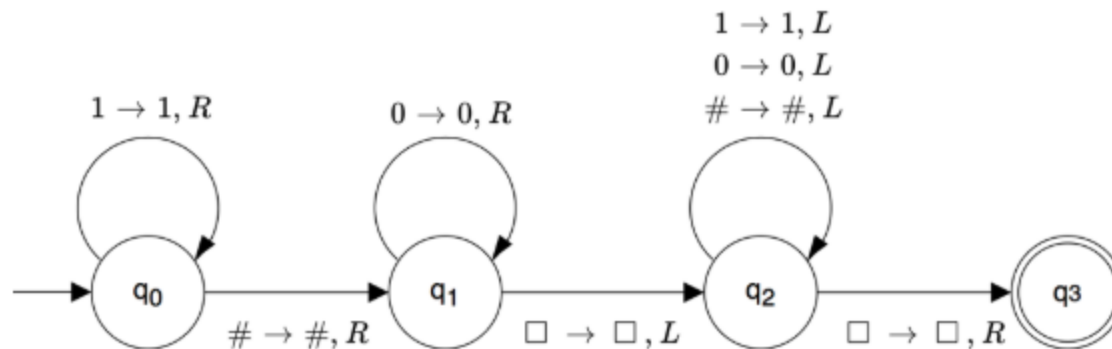
☐

E-)

$L = \{ \langle M, w \rangle \mid M \text{ is a Turing machine that accepts } w \}$

☐

Consider following Turing machine, wherein \square denotes the blank symbol, and input tape is infinite in both ends. Which of the following is FALSE?



- A-) One purpose of q_2 is to return to the starting position. ☐
- B-) One purpose of q_0 is to consume starting 1's. ☐
- C-) This Turing machine accepts the string 1#000 ☐
- D-) This Turing machine accepts the string 10#00 ☒
- E-) After the Turing machine accepts a string, the head goes to its starting position. ☐

Let the alphabet be $\Sigma=\{a,b,c\}$.

Which one of the following grammars is in Chomsky normal form?

A-)

$S \rightarrow TT \mid a$

$T \rightarrow c$



B-)

$S \rightarrow TT \mid ac$

$T \rightarrow c$



C-)

$S \rightarrow TTa \mid a$

$T \rightarrow c$



D-)

$S \rightarrow TTTa \mid a$

$T \rightarrow c$



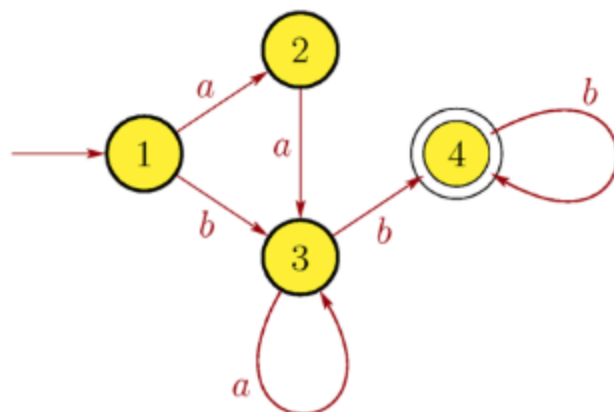
E-)

$S \rightarrow Ta \mid a$

$T \rightarrow c$



Let ε be empty string. Give a regular expression for the language recognized by the non-deterministic finite automata (NFA) given below.



- A-) $(aaa^* + ba^*)a^*bb^*$ ☐
- B-) $(aa + b)a^*bb^*$ ☒
- C-) $(ba^* + \varepsilon)a^*bb^*$ ☐
- D-) $(aa + \varepsilon)a^*bb^*$ ☐
- E-) $(ba^* + a^*b)a^*bb^*$ ☐

Let $L = \{ w \in (0+1)^* \mid w \text{ has even number of 1's} \}$, i.e. L is the set of all bit strings with even number of 1's.

Which one of the regular expressions below represents L ?

A-)

$0^*1(10^*1)10^*$

☐

B-)

$(0^*10^*1)^*$

☐

C-)

$0^*(10^*1^*)^*0^*$

☐

D-)

$0^*(0^*10^*10^*)^*0^*$

☒

E-)

None of them

☐

Let A be a language.

Which one of the following is **FALSE**.

A-)

If A is finite, then A has to be Turing-recognizable.

☐

B-)

If A is regular, then A has to be Turing-recognizable.

☐

C-)

If A is context-free, then A has to be Turing-decidable.

☐

D-)

If A is infinite, then A has to be Turing-acceptable.

☒

E-)

If A is context-free, then A has to be Turing-recognizable.

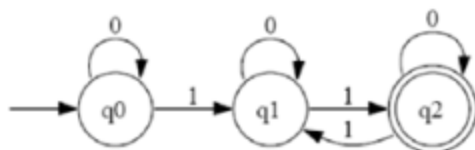
☐

Which one of the following deterministic finite automata (DFA) accepts the language

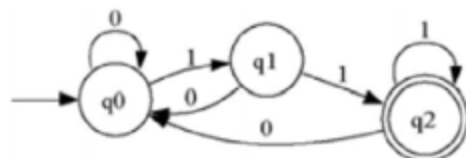
$$L = \{w \in \Sigma^* \mid \text{each } 1 \text{ in } w \text{ is immediately followed by a } 0\}$$

over alphabet $\Sigma = \{0, 1\}$.

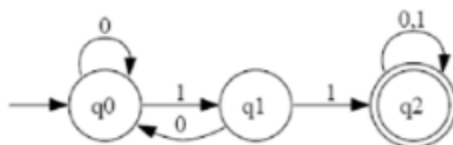
A-)


☐

B-)


☐

C-)

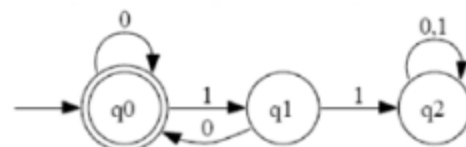

☐

D-)

None of them.

☐

E-)


☒

Which ones of the following statements are TRUE?

Statement I: If $A \subseteq B$ and A is a regular language, then B must be regular.

Statement II: If $A \subseteq B$ and B is a regular language, then A must be regular

- A-) Both statements are ambiguous. ☐
- B-) Both statements are FALSE. ☐
- C-) Only second statement is TRUE. ☒
- D-) Only first statement is TRUE. ☐
- E-) Both statements are TRUE. ☐

Which one of the following languages is context free?

A-)

$$L = \{a^m b^n (ac)^{m+n} \mid m, n \geq 1\}$$

☐

B-)

$$L = \{a^{m+n} (ab)^m c^{m+n} \mid m, n \geq 1\}$$

☐

C-)

$$L = \{(ac)^m c^m b^m \mid m, n \geq 1\}$$

☒

D-)

$$L = \{a^m (ac)^{m+n} b^m \mid m, n \geq 1\}$$

☐

E-)

$$L = \{a^m b^m (ac)^{m+n} \mid m, n \geq 1\}$$

☐

Let $A=\{a,b,c\}$ and $B=\{1,2,3,4\}$ be the domain and range of the function $f:A\rightarrow B$ such that

$$f(a)=1,$$

$$f(b)=2,$$

$$f(c)=1.$$

Then, which one of the following is FALSE about the function $f:A\rightarrow B$.

- A-) There exist a standard Turing machine that implements the function f . ☐
- B-) It is not an invertible function since $f(a)=f(c)=1$. ☐
- C-) It is not a correspondence because f is not one-to-one and onto. ☒
- D-) It is not an onto function because nothing in A maps to 3 or to 4, which are both in B . ☐
- E-) It is a one-to-one function since $f(a)=f(c)=1$. ☐

Consider the following languages.

$$L_1 = \{ 0^p 1^q 0^r \mid p, q, r \geq 0 \}$$

$$L_2 = \{ 0^p 1^q 0^r \mid p, q, r \geq 0 \text{ and } p \neq q \}$$

Which one of the following statements is FALSE?

- A-) L_2 is subset of L_1 . ☐
- B-) Complement of L_1 is context-free but not regular. ☐
- C-) L_1 intersection L_2 is context-free. ☐
- D-) L_2 is context-free. ☐
- E-) L_1 is context-free. ☒

Fill in the blank with the correct option.

$L = \{a^n b^n c^n \mid n \geq 0\}$ is not a context-free language because

_____.

A-) L is not a regular language.

☐

B-) $A = \{a^5 b^5 c^5\}$ is not a context-free language.

☐

C-) L is not Turing recognizable.

☐

D-) L is a context-free language. The question is wrong.

☐

E-) we can use Pumping Lemma to prove that L is not a context-free.

☒