CS151 Intro to Data Structures

Implementing Stacks

Queues

Announcements

HW2, Lab3, Lab4 due Friday

Lab5 today - due Oct 11th

Manual checkoff by me or TAs

Agenda

- Stack Review
- Linked List based Stack implementation
- Amortized Analysis
- Queues

Stacks - FILO

- First In Last Out

- *stack* of plates in the dining hall

- Operations:
    - push
    - pop
    - peek
    - isEmpty

Stack Review - what will this code print?

```
public static void main(String[] args) {
    Stack<Integer> stack = new Stack<Integer>();

    stack.push(10);
    stack.push(20);
    stack.push(30);

    int popped = stack.pop();
    System.out.println("Popped: " + popped);

    int top = stack.peek();
    System.out.println("Top: " + top);

    stack.push(40);
    System.out.println("New Top after push: " + stack.peek());

    while (!stack.isEmpty()) {
        System.out.println("Popped: " + stack.pop());
    }
}
```

Implementing a Stack with an Array

Goal: O(1) operations

Our class implementation:

- fixed size array (no expansions!)
- How did we implement push?
- How did we implement pop?
- How did we implement peek?

Now let’s implement stack with a linked list!

Goal: O(1) operations.

What to consider:

- When we PUSH where should we insert to?
    - Front, back, middle?
- When we POP where should we remove from?
    - Reminder: Stack should be FIFO

Linked List Stack Performance

Space complexity is

- O(n)

Runtime Complexity:

- push:
  - O(1)
- Pop:
  - O(1)
- Peek:
  - O(1)

Stack Summary

- FIFO wrapper around Array / Linked List

- Allows for limited data structures operations all with O(1) cost

- Real world applications: call stack, browser history, postfix calculator

Amortized Analysis

https://courses.cs.washington.edu/courses/cse373/17wi/summaries/amortized-runtime.pdf

Amortized Analysis

average run time complexity of an operation.

Compares the total cost of a series of operations with how many of those operations happened.

Amortized Runtime Analysis

$$\frac{\textit{total cost of operations}}{\textit{total number of operations}}$$

Where an “operation” is the operation a client is doing through your public interface, like insert(5) or pop() or add(3).

Exercise: Amortized Analysis of Array Insert

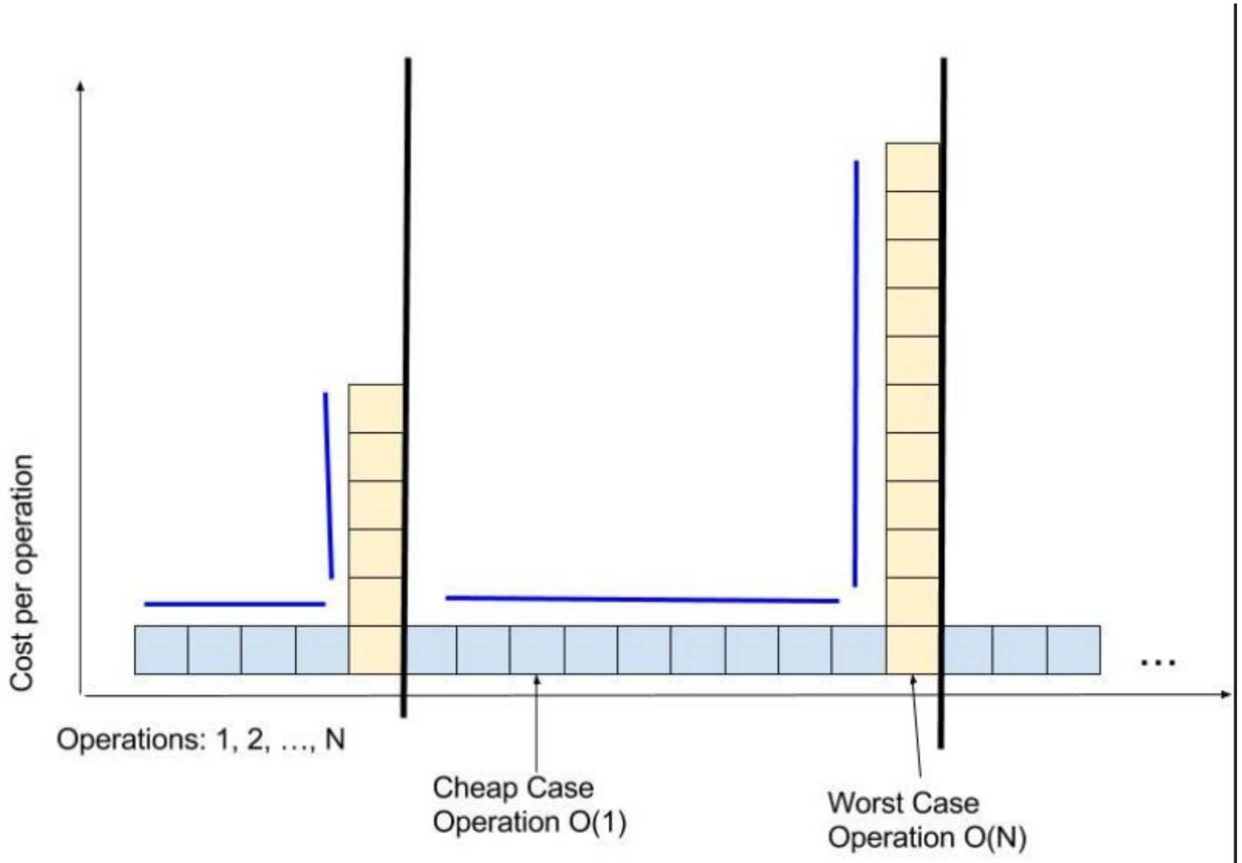
Let's say you have an expandable array with N elements: [1, 2, 3... N]

When you try to add the N+1 th element, a resize occurs.

$$\frac{\textit{total cost of operations}}{\textit{total number of operations}}$$

Amortized Analysis

Intuition: you want to “build up enough credit” with a series of cheap operations, so that when you have one (or more) expensive operations, you can average out the cost of the expensive one



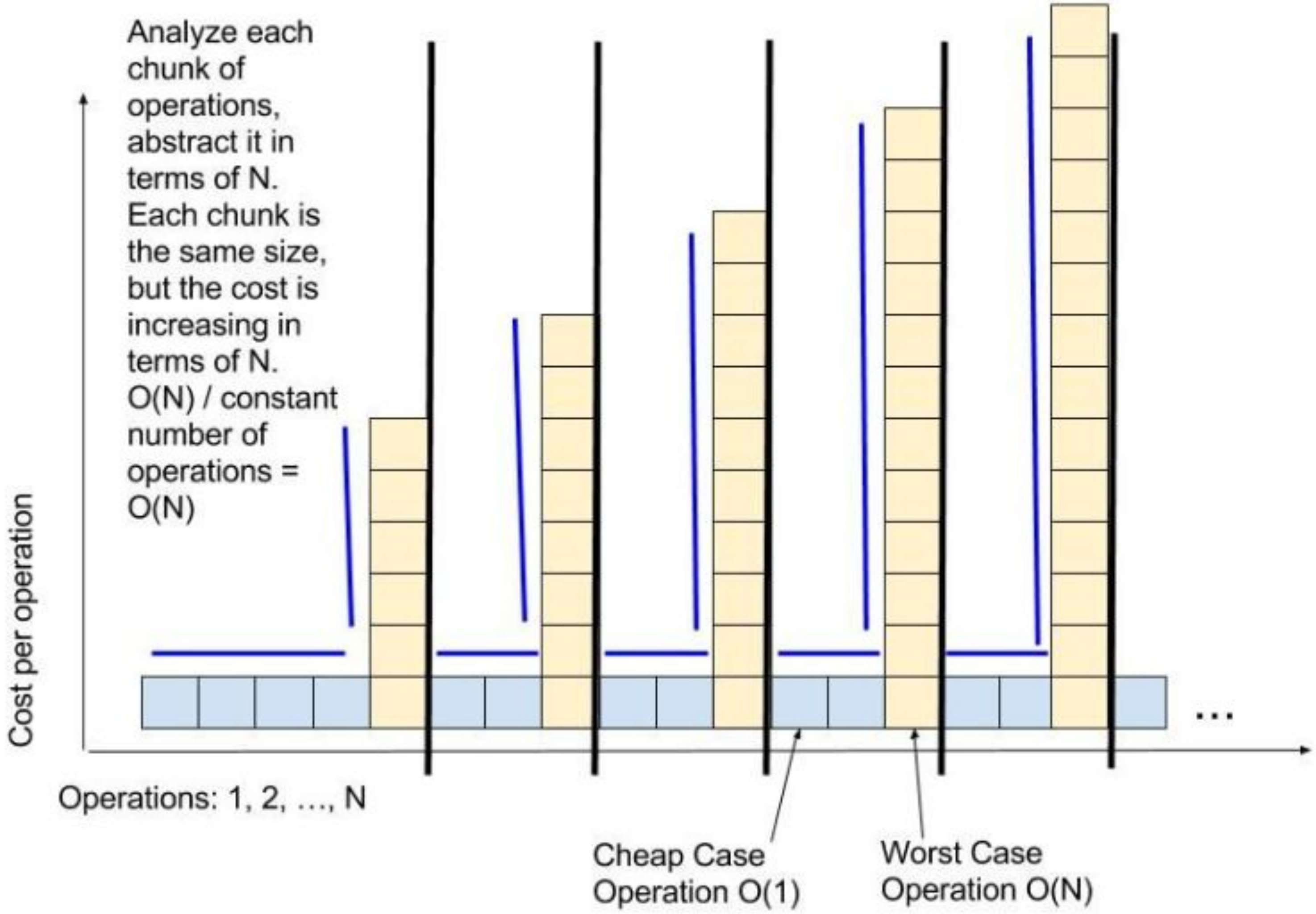
Exercise: Amortized Analysis of Array Insert

Let's say you have an expandable array **which adds two extra slots each time it is full**

Amortized cost of insert? Calculate for 9 inserts.

$$\frac{\textit{total cost of operations}}{\textit{total number of operations}}$$

Amortized Analysis



Amortized Analysis

Average case runtime analysis

Intuition: you want to “build up enough credit” with a series of cheap operations, so that when you have one (or more) expensive operations, you can average out the cost of the expensive one

Comparing the total cost of a series of operations with how many operations happened

Queues

FIFO Stacks

Stack Property

First-in Last-out (FILO)

Where might a FILO stack not make sense?

Line for the cash register

Printer Queue

FIFO: First-in First-out

The first item in, is the first item out

Add-to the back, remove from the front

This is a **Queue**

Inserting – "`enqueue`"

Removing - "`dequeue`"

`Queue` Interface

```
public interface Queue<E> {
   int size();
   boolean isEmpty();
   E first();
   void enqueue(E e);
   E dequeue();
}
```

- `null` is returned from `dequeue()` and `first()` when queue is empty

Queue Example

Cash register code

Example

Operation

Output *Q*

Example

Operation	*Output*	*Q*
enqueue(5)		

Example

Operation	*Output*	*Q*
enqueue(5)	–	(5)

Example

Operation	*Output*	*Q*
enqueue(5)	–	(5)
enqueue(3)	–	(5, 3)
dequeue()	*5*	(3)
enqueue(7)	–	(3, 7)
dequeue()	*3*	(7)
first()	*7*	(7)
dequeue()	*7*	()
dequeue()	*null*	()
isEmpty()	*true*	()
enqueue(9)	–	(9)
enqueue(7)	–	(9, 7)
size()	*2*	(9, 7)
enqueue(3)	–	(9, 7, 3)
enqueue(5)	–	(9, 7, 3, 5)
dequeue()	*9*	(7, 3, 5)

Example

Operation	*Output*	*Q*
enqueue(5)	–	(5)
enqueue(3)	–	(5, 3)
dequeue()		
enqueue(7)		
dequeue()		
first()		
dequeue()		
dequeue()		
isEmpty()		
enqueue(9)		
enqueue(7)		
size()		
enqueue(3)		
enqueue(5)		
dequeue()		

Example

Operation	*Output*	*Q*
enqueue(5)	–	(5)
enqueue(3)	–	(5, 3)
dequeue()	*5*	(3)
enqueue(7)	–	(3, 7)
dequeue()	*3*	(7)
first()	*7*	(7)
dequeue()	*7*	()
dequeue()	*null*	()
isEmpty()	*true*	()
enqueue(9)	–	(9)
enqueue(7)	–	(9, 7)
size()	*2*	(9, 7)
enqueue(3)	–	(9, 7, 3)
enqueue(5)	–	(9, 7, 3, 5)
dequeue()	*9*	(7, 3, 5)

Implementing a Queue with an Array

Goal: O(1) operations

How can we achieve this?

1. Fixed size underlying array (no expansions)

2. Where should we insert?
   a. front, back, middle?

3. Where should we remove from?
   a. we should preserve first in first out property!