

# CS151 Intro to Data Structures

Quick Sort  
Balanced Binary Search Trees

# Announcements

HW7 and Lab9 due Friday 4/18

Lab9 Manual checkoff

# Outline

Warmup

Sorting:

1. QuickSort

Balanced BSTs!

# Which data structure would you use?

You are implementing a system to track and manage a library's collection of books. Each book has a unique ISBN number, and the system needs to efficiently support the following operations:

- **Add a book:** Insert a new book using its ISBN number as the key.
- **Remove a book:** Delete a book from the system by its ISBN number.
- **Find a book:** Retrieve details about a book by its ISBN number.
- **Get all books in sorted order:** Return a list of all books, sorted by their ISBN numbers.
- **Find the book with the closest higher ISBN:** Given an ISBN, find the next highest ISBN in the collection.

Design a data structure to efficiently support these operations. Justify your choice and explain the time complexity for each operation.

# Quicksort

# Quicksort

- Divide and conquer
- **Divide:** select a *pivot* and create three sequences:
  - a. L: stores elements less than the pivot
  - b. E: stores elements equal to the pivot
  - c. G: stores elements greater than the pivot
- **Conquer:** recursively sort L and G
- **Combine:** L + E + G is a sorted list

# Quick Sort

Sort [2, 6, 5, 3, 8, 7, 1, 0]

1. choose a pivot
2. swap pivot to the end of the array
3. Find two items:
  - a. left which is larger than our pivot
  - b. right which is smaller than our pivot
1. swap left and right
2. repeat 3 and 4 until right < left
3. swap left and pivot
4. Sort L E and R recursively

# Quick Sort - Choosing a pivot

What if we chose our pivot to be the smallest element?

We want a pivot that divides our list as evenly as possible.

Median-of-three: look at the first, middle, and last elems in the array, and pick the middle element.



# Quicksort runtime complexity

Bad pivot:

$$O(n^2)$$

Good pivot:

$$O(n \log n)$$

# Summary of Sorting Algorithms

Algorithm	Time
selection-sort	
heap-sort	
merge-sort	
quick-sort	

# Balanced Binary Trees

# What can go wrong?

Complexity?

**Search**

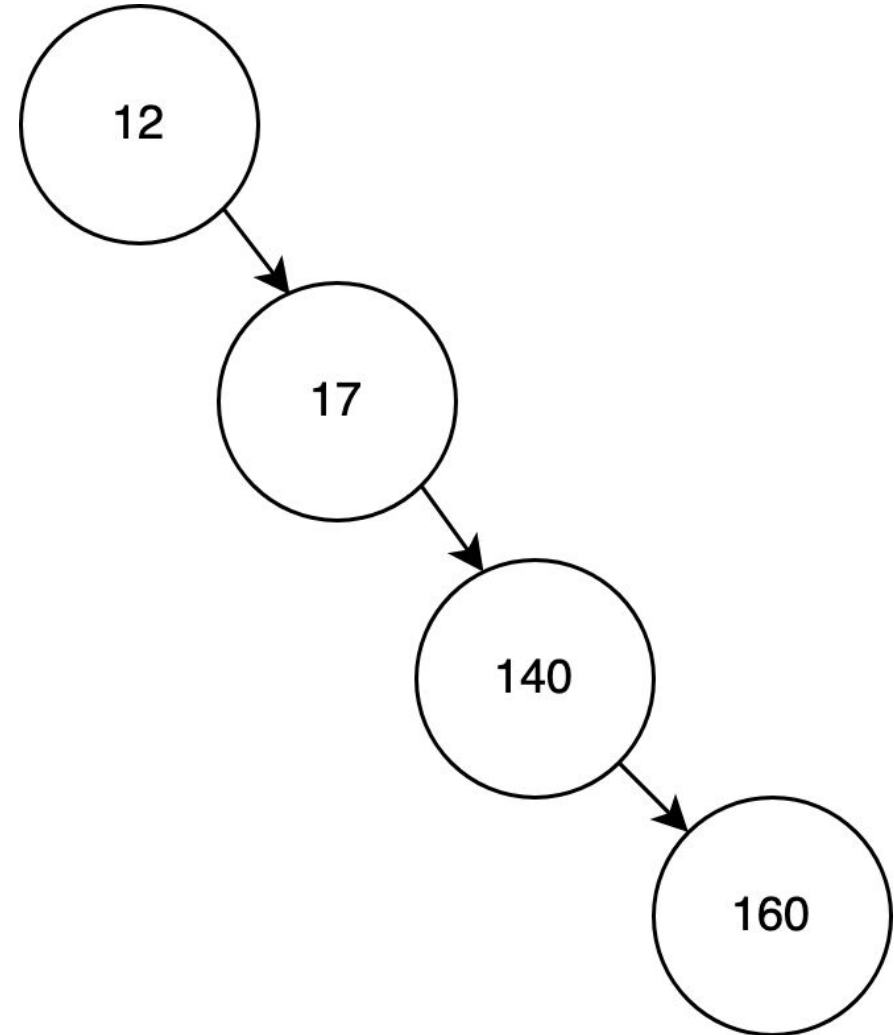
$O(n)$

**Insertion:**

$O(n)$

**Deletion:**

$O(n)$



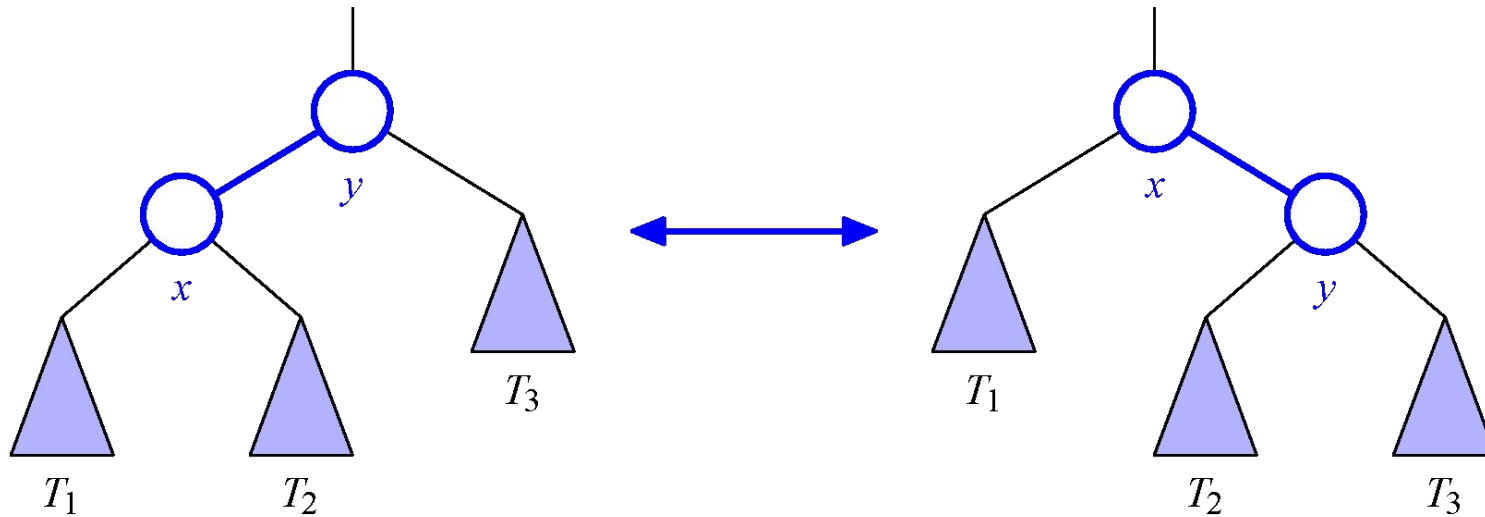
# Balanced Binary Trees

- Difference of heights of left and right subtrees at any node is at most 1
- Add an operation to BSTs to maintain balance:
  - **Rotation**

# Rotation Operation

Move a child above its parent and relink subtrees

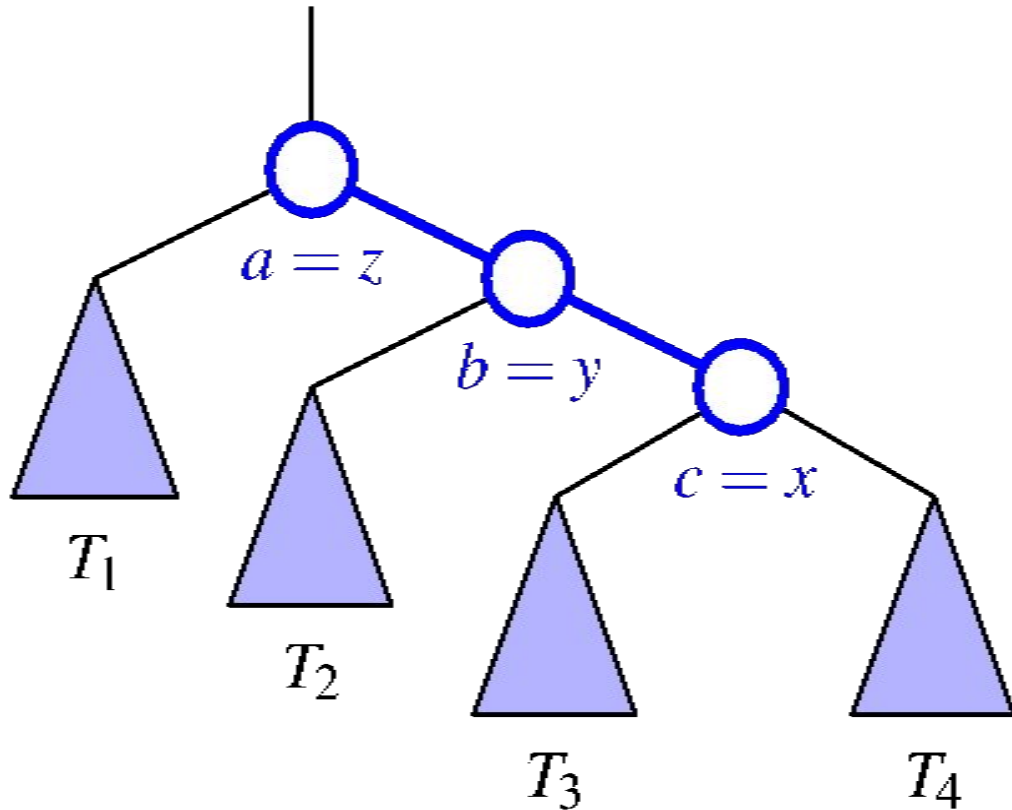
Maintains BST order



# Rotation Operation

- Used to maintain balance
- When should **rotate** be invoked?
  - Difference of heights of left and right subtrees at any node is  $> 1$

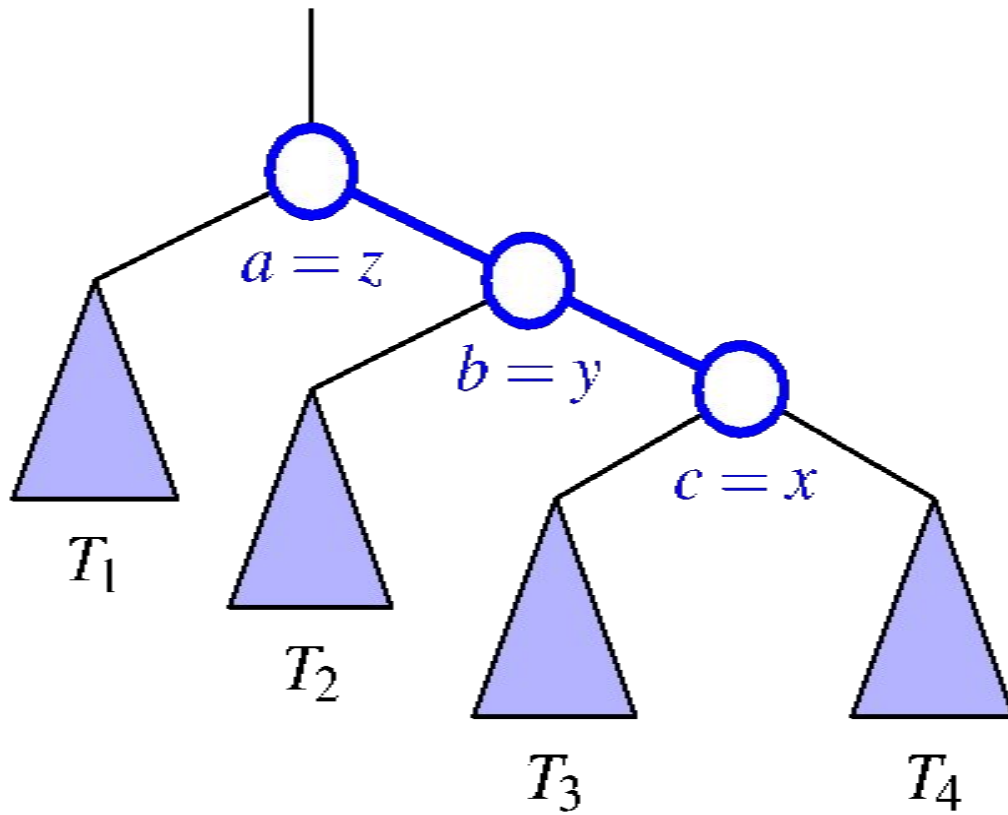
# Rotation Operation



- Assume heights of subtrees are equal
  - $h(T_1) = h(T_2) = h(T_3) = h(T_4)$
- What is the height of the entire tree?
  - $h(T_3) + 2$
- What is the height of the left subtree of  $a$ ?
  - $h(T_1)$
- What is the height of the right subtree of  $a$ ?
  - $h(T_4) + 2$
- Is this tree balanced?



# Rotation Operation

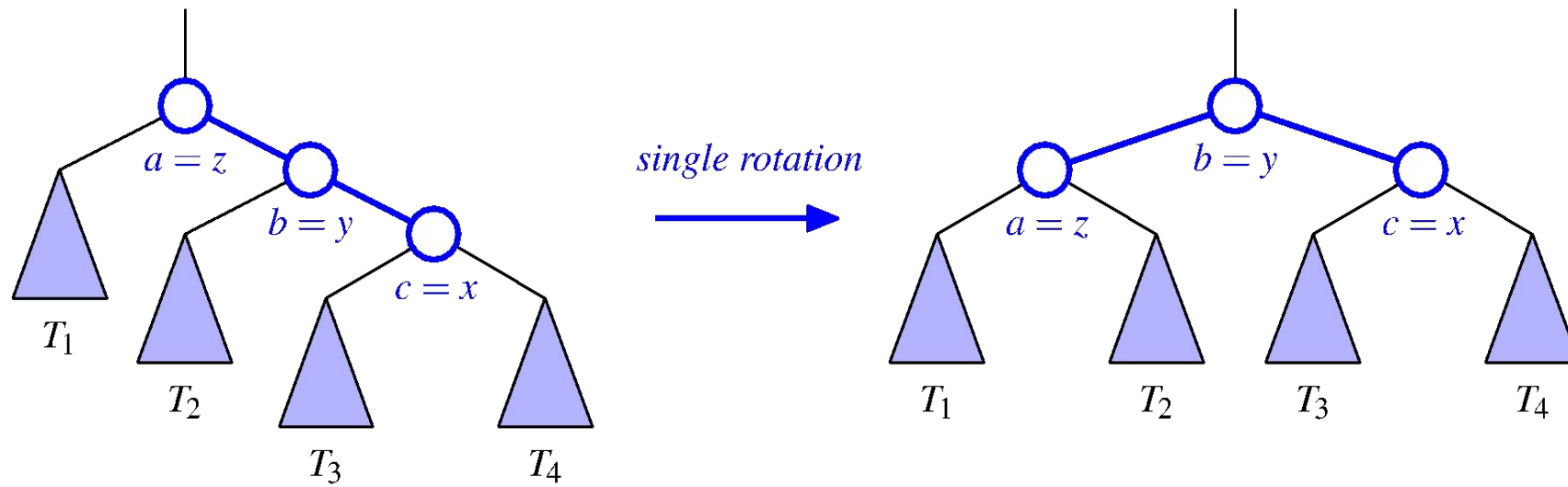


Right subtree is too large!

How can we rotate to fix this?

What should we make the root?

# Single Rotation (around $z$ )



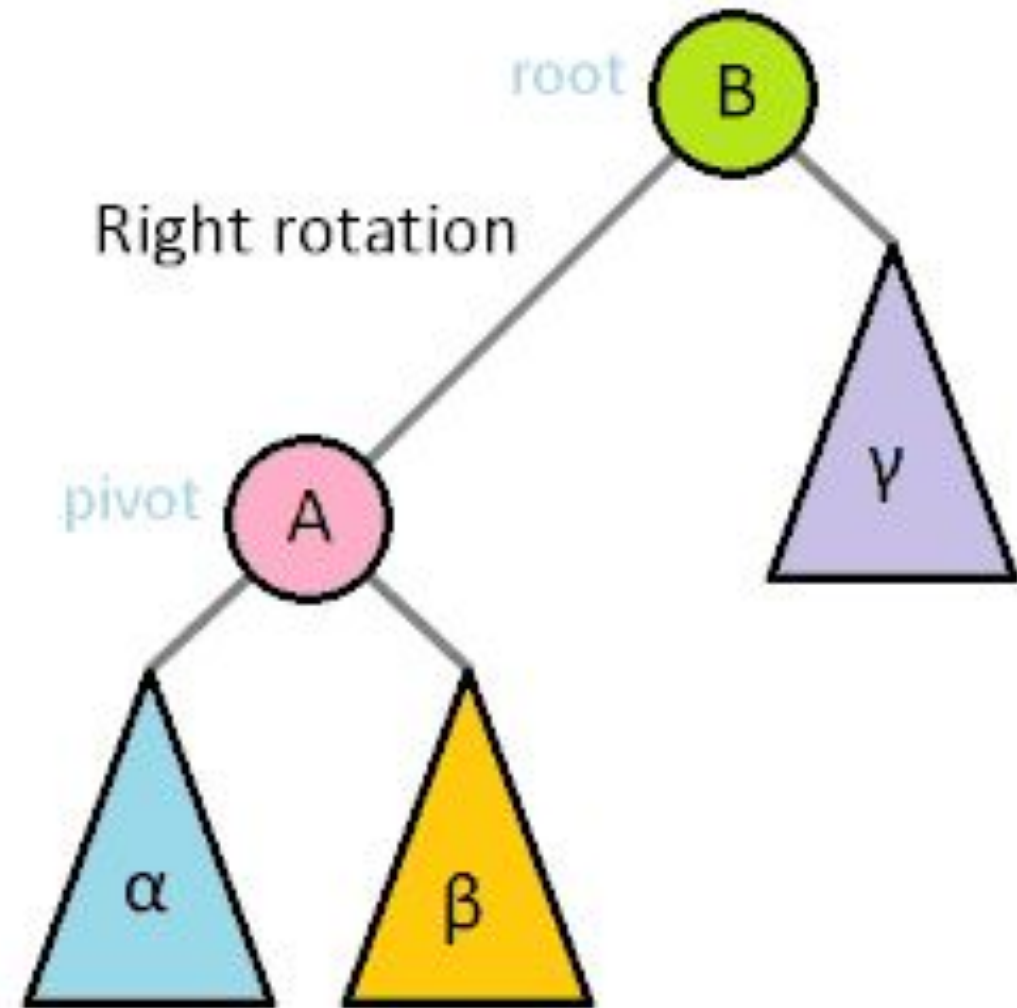
# Rotations

Right rotation:

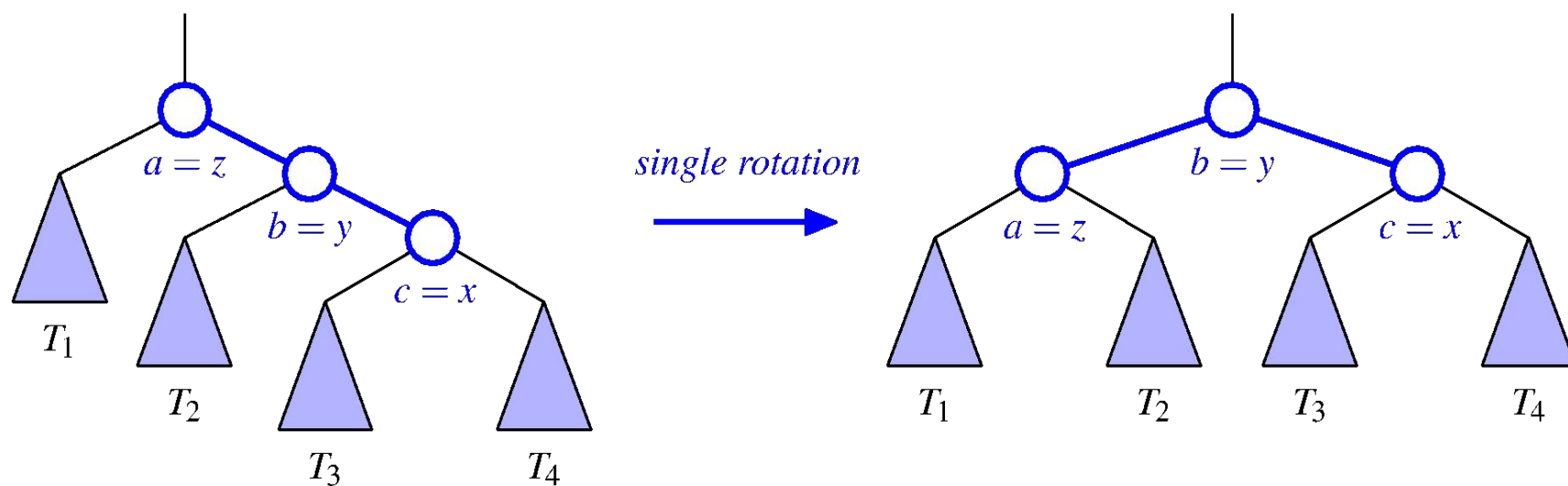
- Performed when left side is heavier
- left child becomes root

Left rotation:

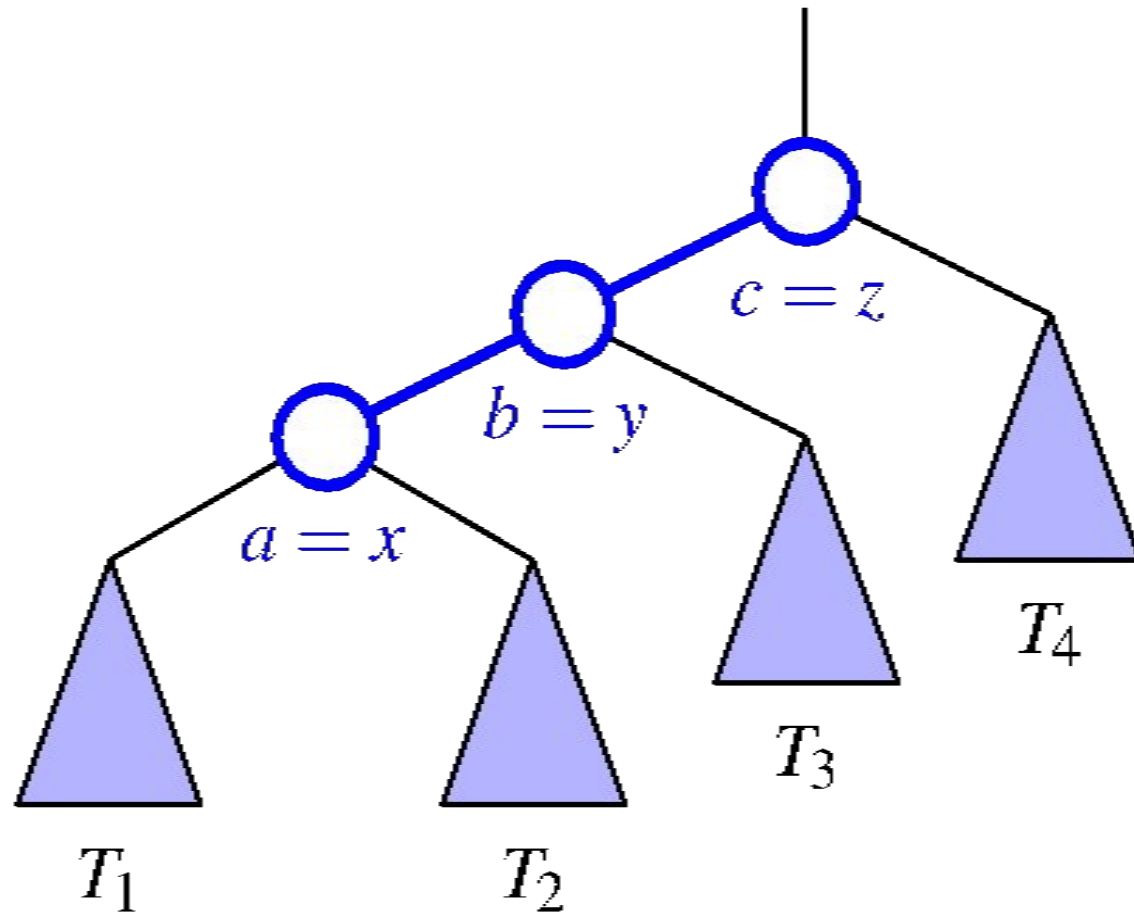
- Performed when right side is heavier
- right child becomes root



# Left or Right rotation?



## Example 2:

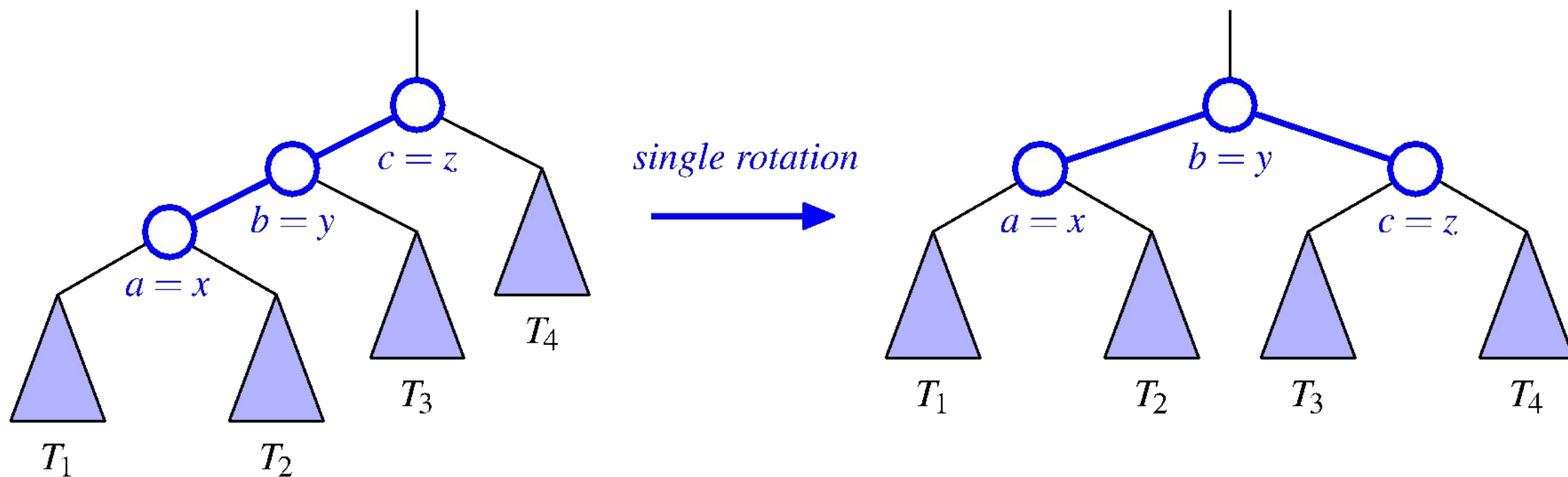


Should we do a left or right rotation?

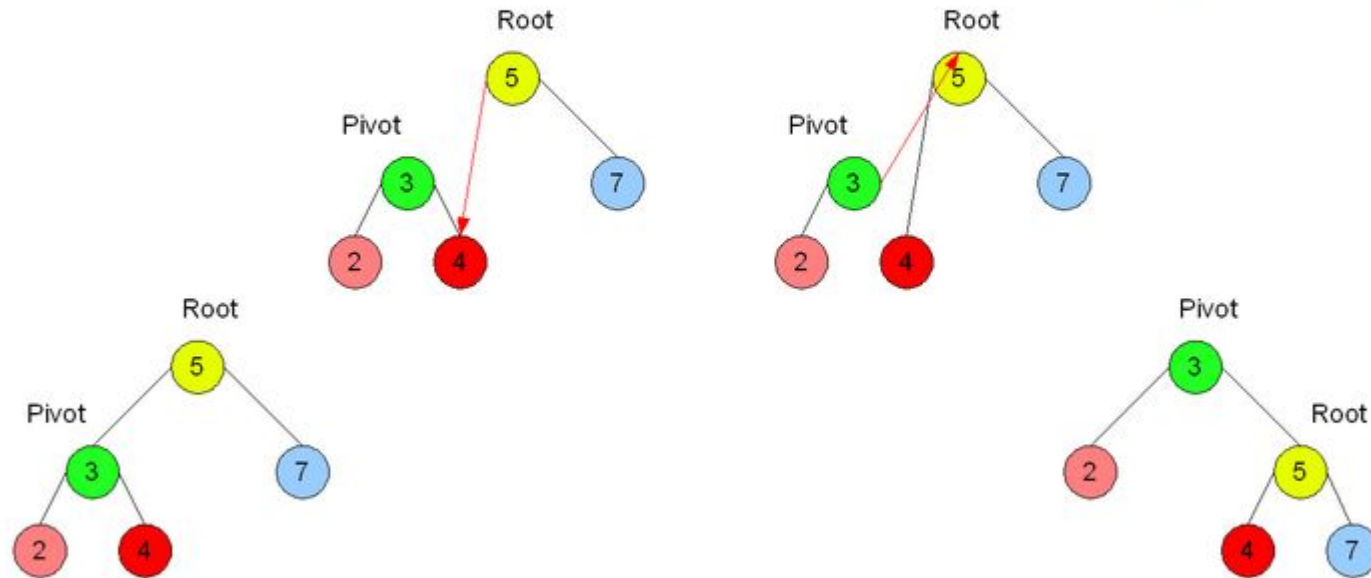
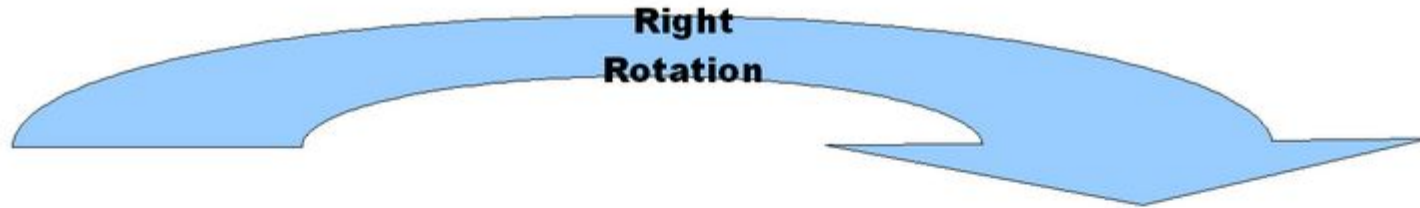
What will become the root?

Let's draw what it will look like after rotation

## Example 2: Rotate Right



# RotateRight Algorithm



Initial state

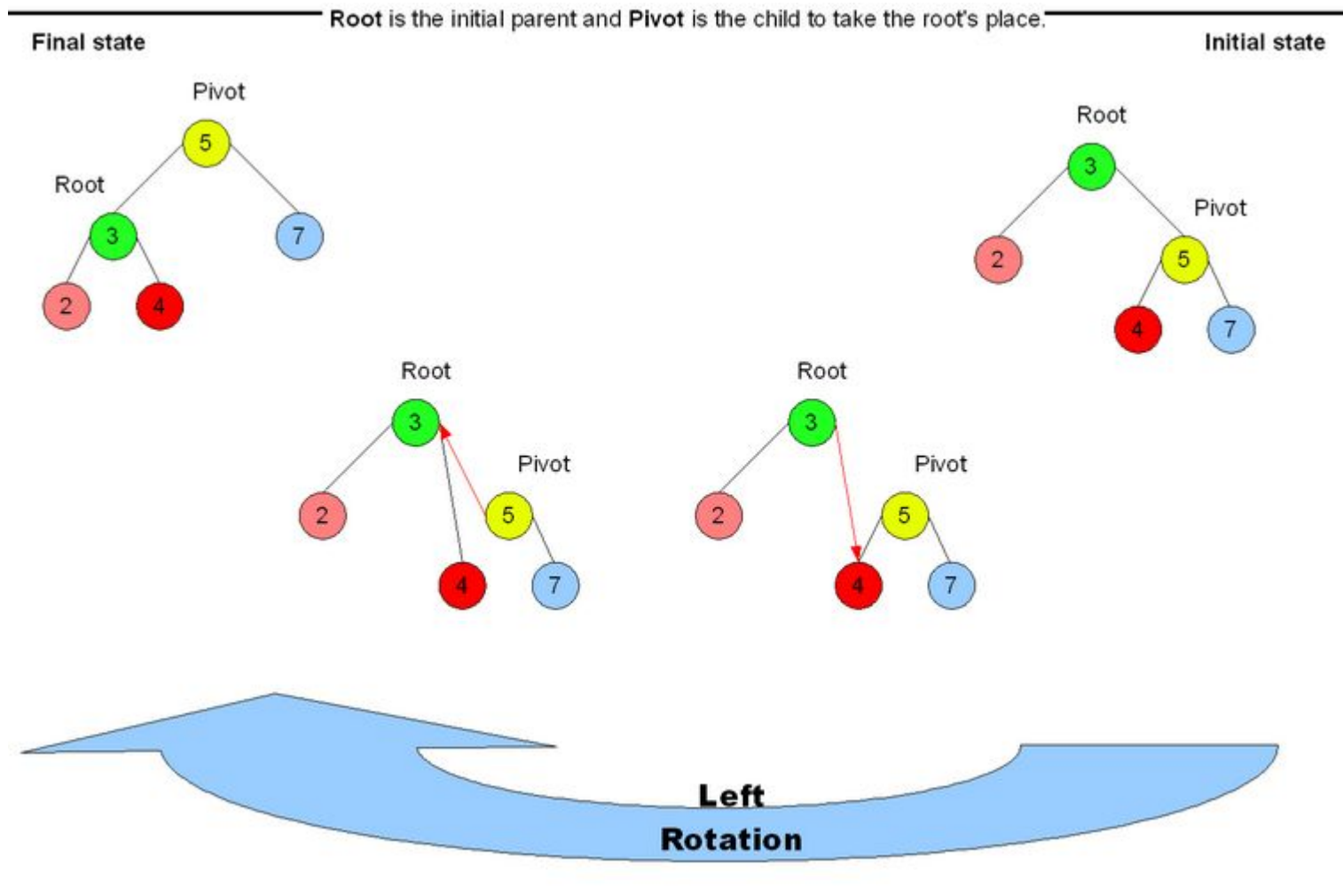
Final state

Root is the initial parent and Pivot is the child to take the root's place.

1. `Root.left = Pivot.right`

1. `Pivot.right = root`

# RotateLeft Algorithm



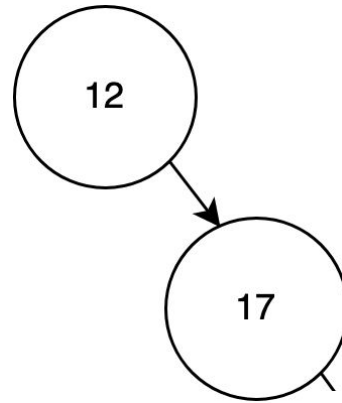
1. `Root.right = Pivot.left`

1. `Pivot.left = root`



# Example:

1. What is the height of the right and left subtrees?
1. Is this tree balanced?
1. Insert 140. Now, revisit questions (1) and (2)
1. Rotate? Which one?



# Runtime Complexity

Runtime Complexity of rotation?

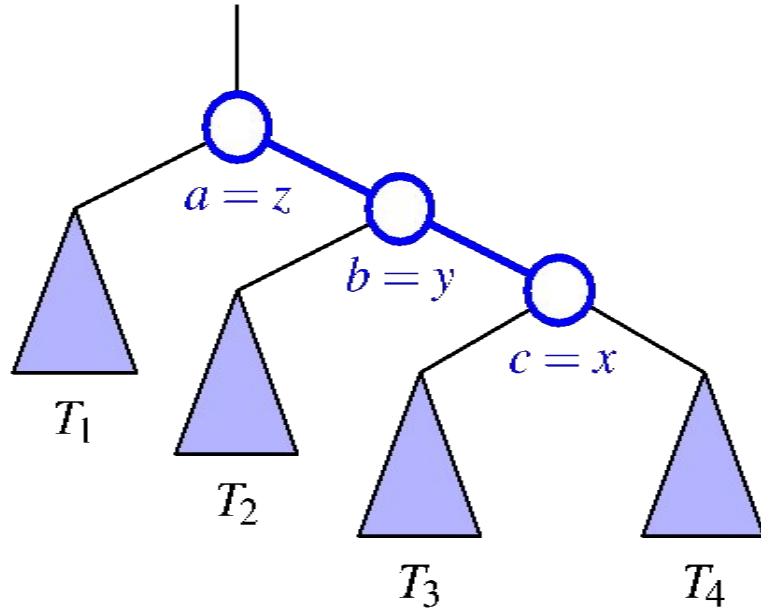
- $O(1)$

Constant time... we're just updating links

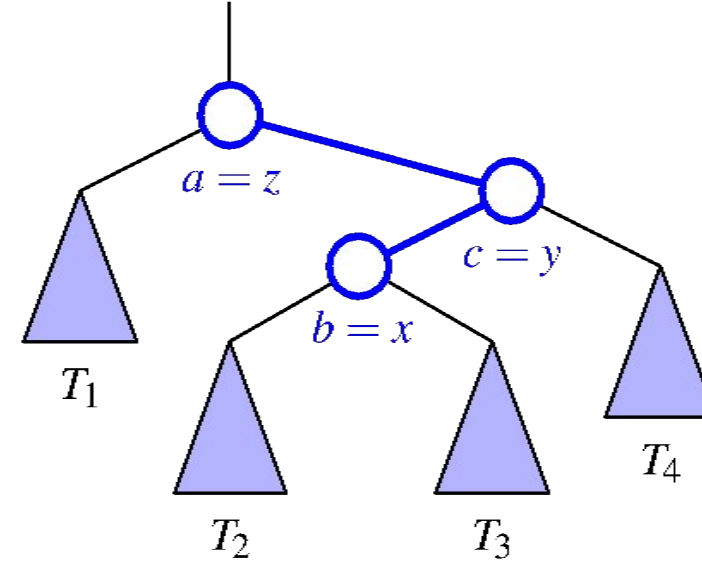
# Double Rotation

Sometimes a single rotation is not enough to restore balance

# Double Rotation

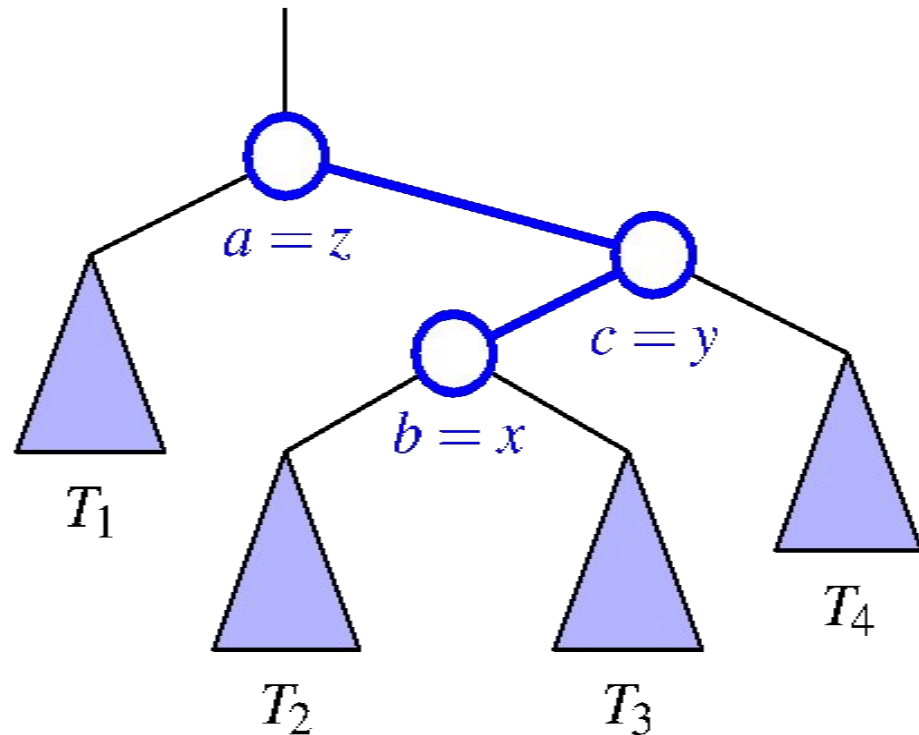


**Right** child of  $a$  is too heavy.. because  
**Right subtree** of  $b$  is too heavy..  
Single Left rotation on the root needed



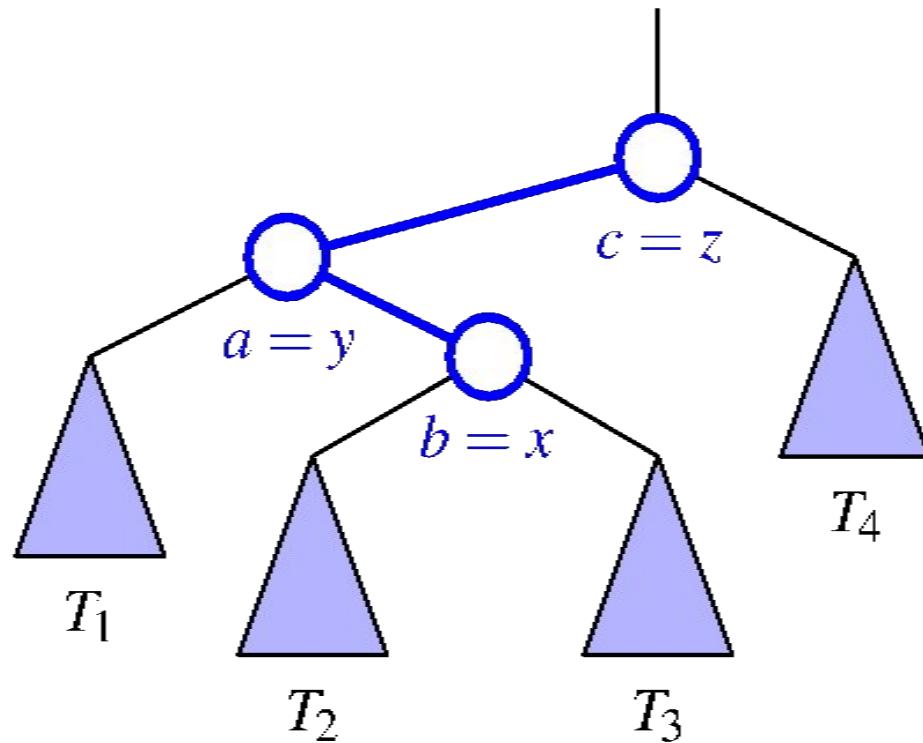
**Right** child of  $a$  is too heavy... because  
**Left subtree** of  $c$  is too heavy  
**Is a single rotation enough?**

# Double Rotation

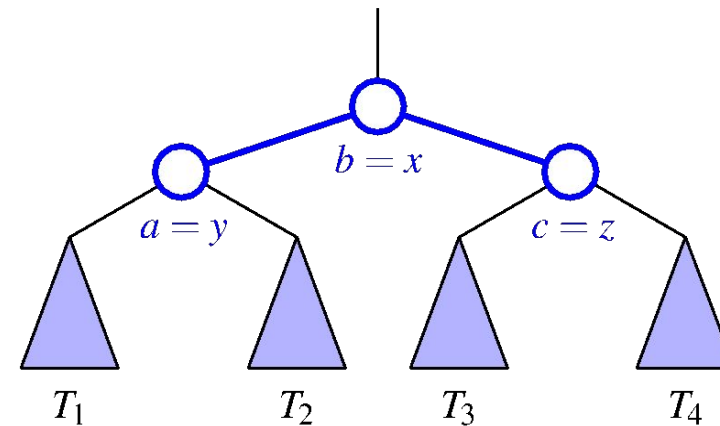


1. **Rotate Right** at  $c$  because right subtree of root is too heavy
2. **Rotate Left** at the root ( $a$ )

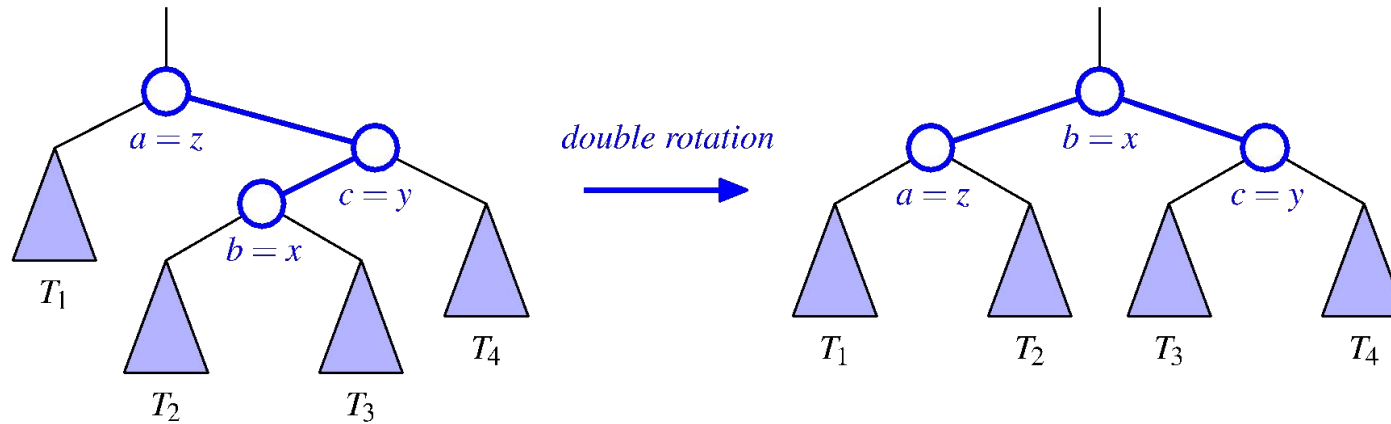
# Double Rotation Example 2:



1. **Rotate Left** at  $a$  because right subtree of root is too heavy
2. **Rotate right** at the root ( $c$ )



# Double Rotations

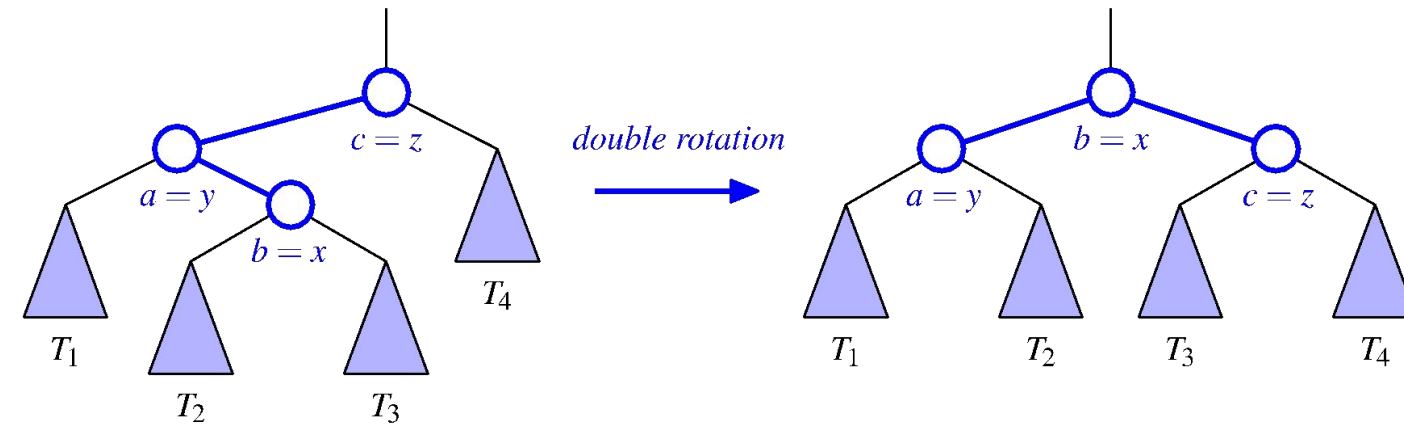


**Right** subtree is too heavy because of **left** subtree of  $c$

1. Rotate Right about  $c$
2. Rotate Left about  $a$

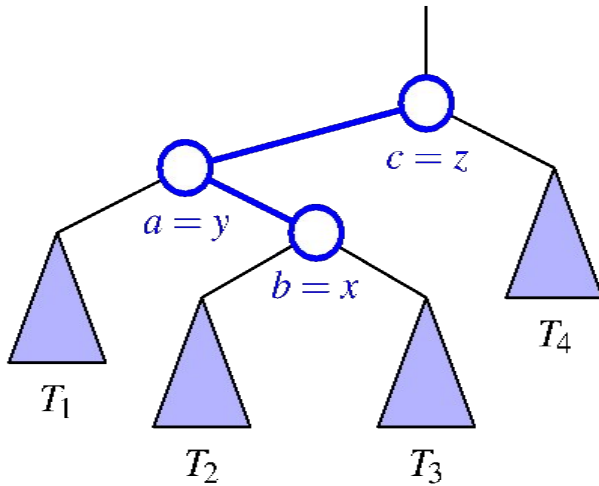
**Left** subtree is too heavy because of **right** subtree of  $a$

1. Rotate Left about  $a$
2. Rotate Right about  $c$

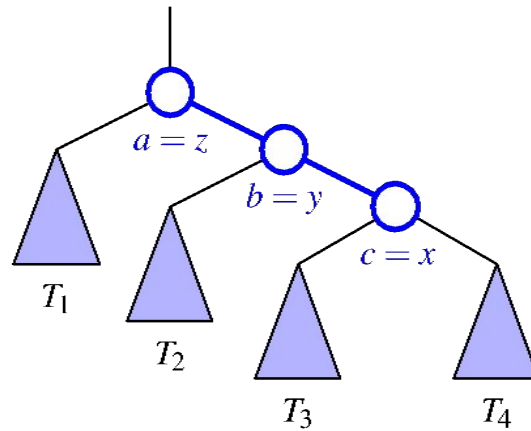


# Double Rotation

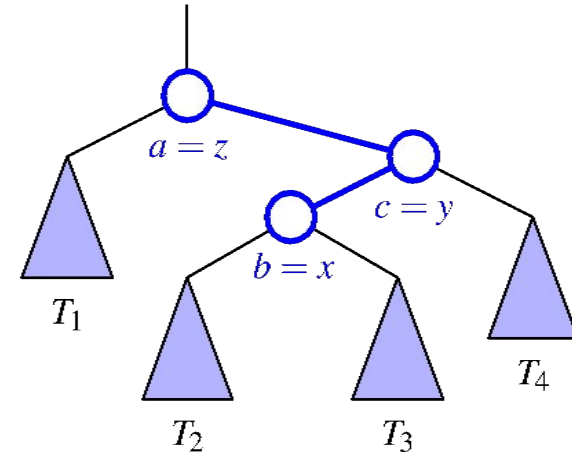
When do we need a double rotation vs a single rotation?



Double rotation



Single rotation



Double rotation

Look for zig-zag pattern!



# Double rotation

When do we need a double rotation?

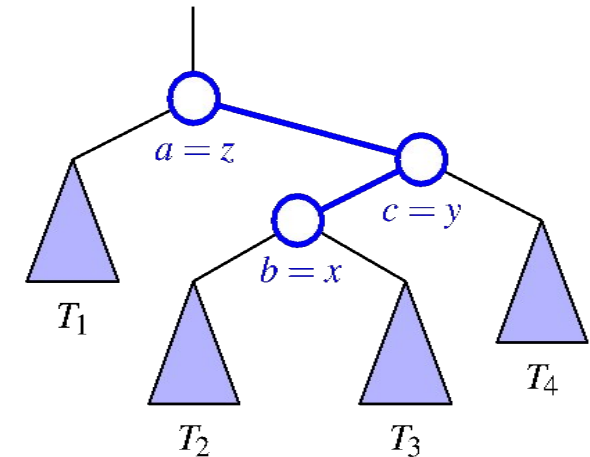
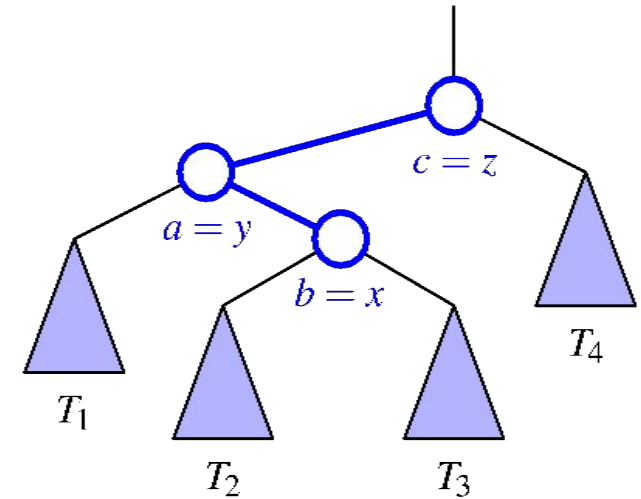
Left subtree is too heavy on the right side

`rotateLeftRight`

OR

Right subtree is too heavy on the left side

`rotateRightLeft`

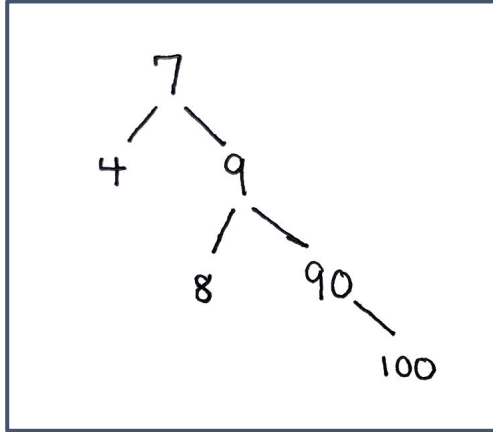


# Double Rotation Code

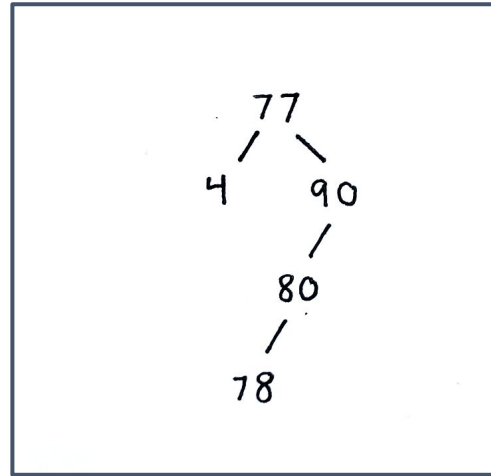
```
def rotateLeftRight(n)
    n.left = rotateLeft(n.left);
    n = rotateRight(n);
```

```
def rotateRightLeft(n)
    n.right = rotateRight(n.right);
    n = rotateLeft(n);
```

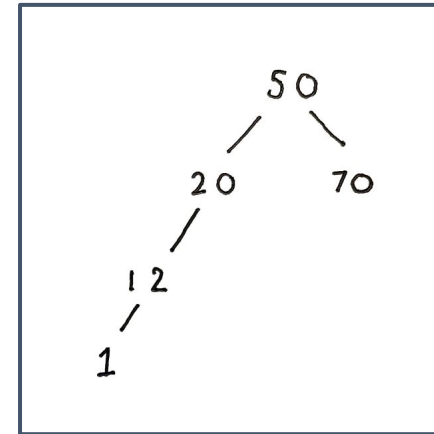
# Examples - which way should I rotate?



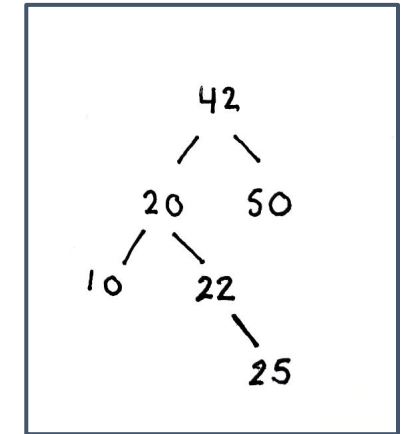
rotateLeft



rotateRightLeft



rotateRight



rotateLeftRight

# Summary: Tree rotation

- Can rotate to left or right
- Used to restore balance in height
- Rotation maintains BST order
- Runtime complexity of rotation?
  - $O(1)$

# Summary

- Quicksort and Mergesort are **recursive  $O(n \log n)$**  sorting algorithms
- In quicksort, good pivots are important in achieving  $O(n \log n)$  runtime complexity
- HashTable + Quicksort homework due Friday!

Rotations:

double rotation needed when

Left subtree is too heavy on the right side OR

Right subtree is too heavy on the left side (zig-zag pattern)

Rotations are constant time