

CS151 Intro to Data Structures

Interfaces

Algorithm Analysis

Announcements

- HW02
 - Linked Lists
 - due Tuesday February 25th
 - start early!

Outline

- Objects
 - Comparing Objects
- Interfaces
- Algorithm Analysis

Java Memory Diagrams

Given a segment of code:

- Trace the program by going through line by line
 - Draw a memory diagram
 - As program variables are updated, update their values in the diagram

Example 2

Given a segment of code:

- Trace the program by going through line by line
 - Draw a memory diagram
 - As program variables are updated, update their values in the diagram

Interfaces

- An interface is a contract - A set of shared methods that users **must** implement
- create a program to calculate the area of different shapes, such as circles, rectangles, triangles etc.
- For each shape, you should be able to print the shape name and area
- Every time someone adds a new shape, they **must** include the methods for getName() and getArea()

Interfaces

- For any new shape that is created, we want to **enforce** that these methods are also implemented.

```
interface Shape {  
    public double getArea();  
    public String getName();  
}
```

```
class Circle implements Shape {
```

Interfaces

A contract - A set of shared methods that users **must** implement

A collection of method signatures with no bodies

A class can implement more than one interface

Interfaces

An interface is not a class!

A class is what an object **is**

An interface is what an object **does**

- can not be instantiated

- no constructors

- incomplete methods

Interface

No modifier - implicitly `public`

No instance variables except for constants (`static final`)

Object Comparison

Object Equality

A custom class must define (override) its own `equals`

Object Comparison

- What if we wanted to compare two students by GPA?

```
int compareTo(T o)
```

Parameters:

`o` - the object to be compared.

Returns:

a negative integer, zero, or a positive integer as this object is less than, equal to, or greater than the specified object.

compareTo

compareTo **returns an int, not a Boolean**

Why?

because it needs to convey three outcomes:

- `negative` if smaller compared to the parameter
- `0` if equal
- `positive` if larger compared to the parameter

Comparable interface

The Comparable interface is designed for objects that have an ordering

```
public interface Comparable<T> {  
    int compareTo(T o);  
}
```

Comparable interface

When would we want to use this? **Let's see in code :)**

Now, what if we wanted to sort from highest to lowest GPA

Custom Exceptions

Making Custom Exceptions

Often times we need to raise a custom exception

Extend `Exception` **or** `RuntimeException`

Custom Exceptions

What is the difference between extending from `Exception` rather than `RuntimeException`?

Subclass of `Exception` are checked exceptions – must be treated/caught

Subclass of `RuntimeException` are not checkable during compile time

Computational Complexity

Run Time Complexity

- Mathematical notation used to describe the performance or complexity of an algorithm.
- Hardware independent
- Represents the upper bound of the time complexity in the **worst-case scenario**.
- Helps us understand how the runtime of an algorithm grows ***as the input size increases***.

Big-O Example 1

```
int n = Integer.parseInt(args[0]);  
int power = 1;  
while (power < n) {  
    System.out.print(power + " ");  
    power *= 2;  
}
```

How does the runtime grow as a function of the input size?

$O(\log n)$

Big-O Example 2

```
int fetchFirstElement(int[] arr) {  
    return arr[0];  
}
```

How does the runtime grow as a function of the size of arr?

$O(1)$

Big-O Example 3

```
int n = Integer.parseInt(args[0]);
int tot = 0;
int i = 0;

while (i < n) {
    tot = tot * i;
    i++;

    for (int j=0; j<10000; j++) {
        System.out.println("hello");
    }
}
```

How does the runtime grow as a function of the input size?

Linearly!

$O(n)$

Big-O Example 4

```
int n = Integer.parseInt(args[0]);

for (int i = 0; i > (-1*n); i--) {
    for (int j = 0; j < n; j++) {
        System.out.println(i, j);
    }
}
```

How does the runtime grow as a function of the input size?

Quadratically!

$O(n^2)$

We do n operations n times

Big-O Example 5

```
String[] lst =  
    {"19", "12", "20", "15"};  
  
for (int i=0; i<100; i++) {  
    System.out.println(getNum(lst));  
}  
  
int getNum(int[] arr) {  
    return Integer.parseInt(arr[0]);  
}
```

How does the runtime grow as a function of the size of `lst`?

Constant! The runtime is not affected by the number of elements in `lst`

$O(1)$

Big-O Example 6

```
int[] lst = {1,2,3,4,5,6,7};

for (int i=0; i<lst.length; i++) {
    findMax(lst);
}

int findMax(int[] arr) {
    int max = Integer.MIN_VALUE;
    for (int i=0; i<arr.length; i++) {
        if (arr[i] > max) {
            max = arr[i];
        }
    }
    return max;
}
```

How does the runtime grow as a function of the size of lst?

$O(n^2)$

Space (Memory) Complexity

How much memory a program needs

The space requirements time typically grows with input size. Expressed as a size of the input. (Big O notation)

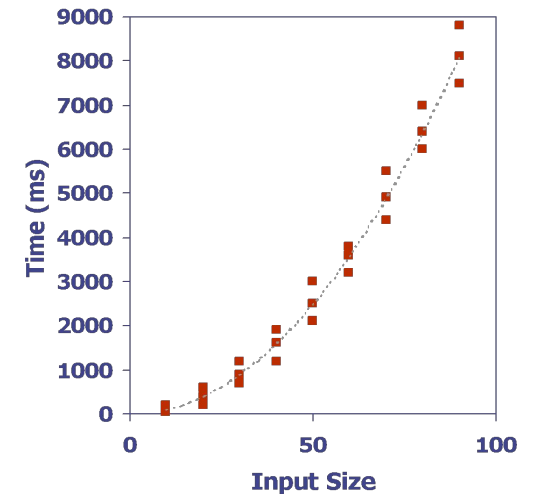
We focus on *worst case* analysis

- how much space will it take in the worst case?

Big O Notation and Theoretical Analysis

- Why do we express runtime notation with Big O notation? Why not just say the run time in number of seconds?

```
1 long startTime = System.currentTimeMillis();           // record the starting time
2 /* (run the algorithm) */
3 long endTime = System.currentTimeMillis();           // record the ending time
4 long elapsed = endTime - startTime;                  // compute the elapsed time
```



- Answer: comparing two algorithms requires exact same hardware and software environments

Constant Time Operations

- Constant time operations require the same amount of time, regardless of the size of the input
- Examples:
 - Basic computations: Assigning variables, adding, multiplying, boolean operators
 - What were some constant time operations in `ExapandableArray`?
 - `LinkedList`?

Linear Time Algorithms: $O(n)$

- The runtime grows linearly as the size of the input grows
- Processes the input in a single pass spending constant time on each item
- Examples:
 - A single loop over an array
 - ExpandableArray?
 - LinkedList?

Example: Find Max

Worst case: $4n + 1 \implies O(n)$

Best case: $3n + 2 \implies O(n)$

Quadratic Time: $O(n^2)$

Nested loops...

Example:

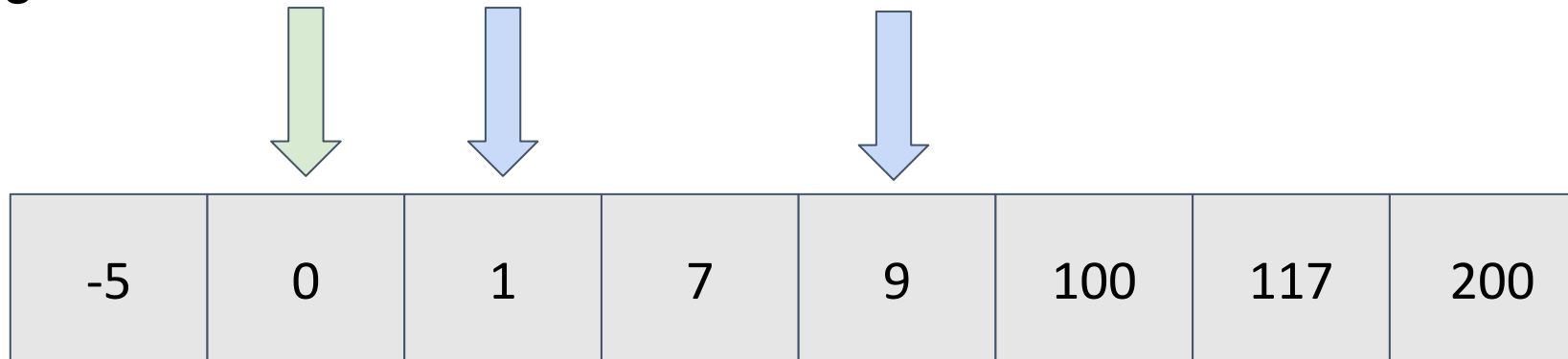
worst case: $4 + 3n^2$

best case: 7

$O(n \log n)$ time

Example: Binary Search!

find 0



How many elements did we touch?

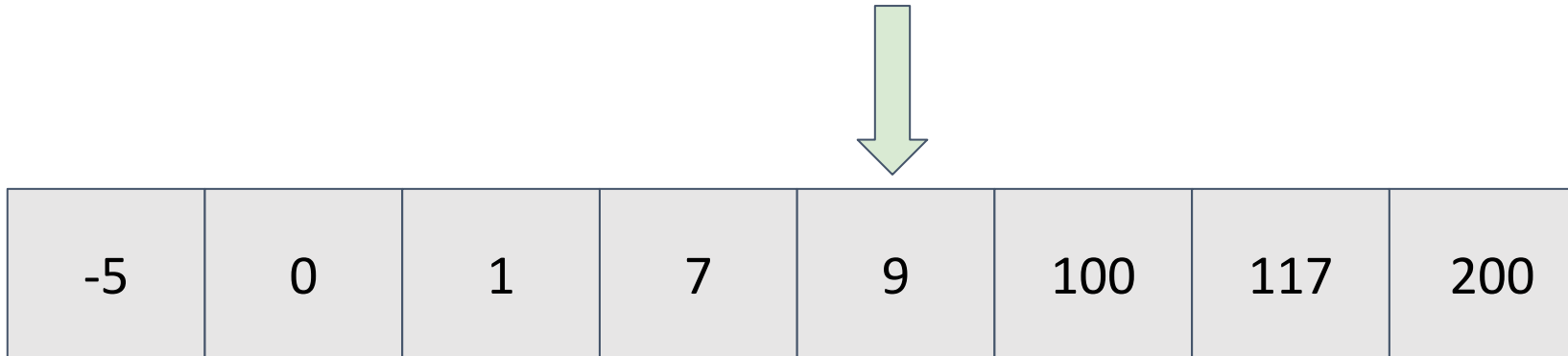
$3 = \log(8)$

Where did the n come from?

$O(n \log n)$ time

Example: Binary Search!

Best case?



Exponential Time: $O(2^n)$

- Generate all possible subsets

$\{a, b, c\} = \dots$

How many subsets are there?

$\{\emptyset\}, \{a\}, \{b\}, \{c\}, \{a,b\}, \{b,c\}, \{a,c\}, \{a,b,c\}$

8

$$2^3 = 8$$

Growth Rate

n	$\log n$	n	$n \log n$	n^2	n^3	2^n
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Growth Rate

n	$\log n$	n	$n \log n$	n^2	n^3	2^n
8	3	8	24	64	512	256

Growth Rate

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8	3	8	24	64	512	256
16	4	16	64	256	4,096	65,536

Growth Rate

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8	3	8	24	64	512	256
16	4	16	64	256	4,096	65,536
32	5	32	160	1,024	32,768	4,294,967,296

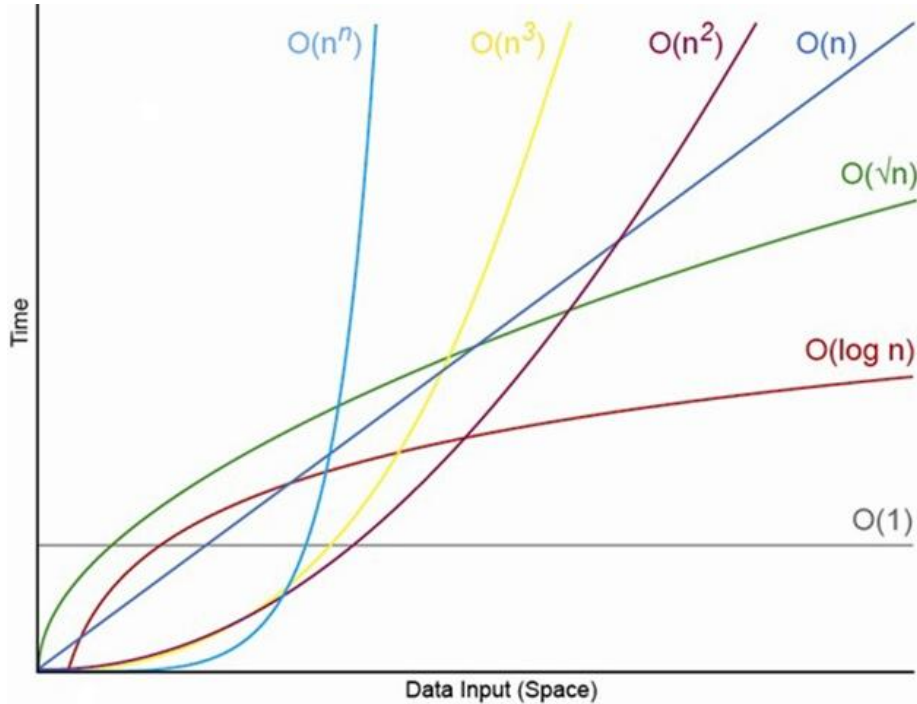
Growth Rate

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16	4	16	64	256	4,096	65,536
32	5	32	160	1,024	32,768	4,294,967,296
64	6	64	384	4,096	262,144	1.84×10^{19}

Growth Rate

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8	3	8	24	64	512	256
16	4	16	64	256	4,096	65,536
32	5	32	160	1,024	32,768	4,294,967,296
64	6	64	384	4,096	262,144	1.84×10^{19}
128	7	128	896	16,384	2,097,152	3.40×10^{38}
256	8	256	2,048	65,536	16,777,216	1.15×10^{77}
512	9	512	4,608	262,144	134,217,728	1.34×10^{154}

Asymptotic Notation



As the number of elements approaches infinity, only the dominant term matters

That is why we simplify $O(n+1)$ to $O(n)$ etc.

Big- O Analysis

1. Write a polynomial in terms of input size n

- Only loops contribute
- Each nested factor is multiplied
- Each sequential factor is summed

2. Simplify the polynomial

- Identify dominant term – highest degree polynomial
- Polynomials beat polylogs
- Exponentials beat polynomials
- Discard constants

Summary

- Every non-primitive is an OBJECT in Java
- Every Object is a REFERENCE
- Interfaces are a *contract* of which methods you will implement
- You can compare objects with compareTo