

## CS340 Analysis of Algorithms Spring 2025

Worksheet: 7

Date:

Title: Recurrences

Professor: Dianna Xu

E-mail: dxu@cs.brynmawr.edu

URL: <http://cs.brynmawr.edu/cs340>

### 1 Recursive Binary Search

Consider the following pseudo code, which searches for a number  $x$  in an array  $A[1, \dots, m]$  by calling `BSearch(A, x, 1, m)`

```

Function BSearch(array A, integer x, integer left, integer right)
  if left > right
  then
    return -1
  end
  m = (left+right)/2
  if A[m] = x
  then
    return m
  end
  else
    if A[m] > x
    then
      return BSearch(A, x, left, m-1)
    end
  end
  else
    return BSearch(A, x, m+1, right)
  end

```

Note that the function operates on an input size of `right - left + 1`. Write the recurrence  $T(n)$  that denotes the maximum number of steps `BSearch` makes on an input of size  $n$ . In particular,

1. In the pseudocode above, highlight the non-recursive parts and estimate their running time in big-Oh notation.
2. In the pseudocode above, highlight the recursive parts and estimate their running time using  $T(n)$ .
3. Given the recurrence for  $T(n)$ . Don't worry about rounding. Do not forget the base case.

## 2 Master Theorem

Consider the following pseudo code segments, each of which has as input an array  $A$  of size  $n$ . State the recurrence run time for each, using  $T(n)$ . As before, (1) highlight the non-recursive parts and estimate their runtime, then (2) highlight the recursive parts and estimate their runtime using  $T(n)$ , then (3) finally evaluate  $T(n)$  using the Master Theorem.

```

Function F(array A[1...n])
  if n = 1
    then
      | return A[1]
    end
  perform  $O(n)$  steps to compute arrays B, C and D, of size  $\lfloor n/3 \rfloor$  each
  x = F(B)
  y = F(C)
  z = F(D)
  return x+y+z

```

```

Function G(array A[1...n])
  if n = 1
    then
      | return A[1]
    end
  perform  $O(n)$  steps to compute arrays B, C and D, of size  $\lfloor n/3 \rfloor$  each
  x = G(B)
  y = G(C)
  return x-y

```

```

Function H(array A[1...n])
  if n = 1
    then
      | return A[1]
    end
  perform  $O(n)$  steps to compute arrays B, C, D, E, of size  $\lfloor n/3 \rfloor$  each
  x = H(B)
  y = H(C)
  z = H(D)
  w = H(E)
  return x+y+z+w

```

```
Function I(array A[1...n])  
  if n = 1  
    then  
      | return A[1]  
    end  
  perform  $O(n^2)$  steps to compute arrays B, C, D, E, of size  $\lfloor n/2 \rfloor$  each  
  x = I(B)  
  y = I(C)  
  z = I(D)  
  w = I(E)  
  return x+y+z+w
```