Simple recursive substitution parser for CFGs

Main theory:

- 1. Start with the target string and the start symbol of G as the first sentential form
- 2. Find all substitution you can make
- 3. Recursively feed the new sentential forms into the parser
- 4. If the target string appears, return True

Now, that's not always (ever?) going to work, because the parser can get lost in an infinite derivation.

Of course, you can play against that with a clever sorting of the rules (recursive one last), but as soon as you have more than one recursive rule, you're obviously lost.

So we're introducing a depth limiter for the parse that just cuts of longer branches and returns False. So, False means "not with this depth" - nothing more.

Simple Rule class: LHS and RHS, nothing more

```
In [3]: class Rule:
    def __init__(self,lhs,rhs):
        self.lhs = lhs
        self.rhs = rhs

def __str__(self):
        return self.lhs + ' --> ' + self.rhs

def __repr__(self):
        return self.__str__()
```

Little wrapper for Sentential Forms

It just wraps the string but also allows for computation of all derivations given a rule set

```
In [4]: class SententialForm:
```

```
# Typical functional style: filter out all rules from "ruleSet" with "symbol" on the LHS
@classmethod
def select(cls,ruleSet,symbol):
    return filter(lambda rule: rule.lhs == symbol, ruleSet)
def init (self, string):
    self.string = string
def str (self):
    return self.string
# Just a combinatorial problem:
# 1. Check the string for non-terminals (always uppercase)
# 2. Find all rules with this non-terminal on the LHS
# 3. Produce the next sentential form bu substituting the LHS for the symbol
# 4. Collect everything and return it
def possibleDerivs(self, rules):
    result = []
    for i in range(len(self.string)): # Check every symbol
        c = self.string[i]
        if c.isupper():
                                      # Uppercase --> non-terminal
            # Check all rules with c on the LHS
            for rule in SententialForm.select(rules,c):
                # The magic substitution step!
                newForm = self.string[:i] + rule.rhs + self.string[i+1:]
                result.append(SententialForm(newForm))
    return result
```

CFG

Simple CFG class

Contains just the start symbol and a set of rules

Well, it also contains a recursive, brute force, depth first, left to right (DFS) parser

AND NOW even a BFS parser

```
In [5]:
    class CFG:
        def __init__(self, start, rules):
            self.start = start
            self.rules = rules
            self.queue = []

    def addRule(self, rule):
        self.rules.append(rule)
```

```
# The mighty parser Part 1: Depth First Search version
# Step 1: Wrap the start symbol into a Sentential Form and jump in
def parseDFS(self,string,maxDepth = 10):
    return self.parseDFSAux(SententialForm(self.start),string,maxDepth)
# That's the recursive workhorse
def parseDFSAux(self,current,string,maxDepth):
    # Recursion base cases - trivial
   if current.string == string: return True
                                                # Got it!
    if maxDepth == 0: return False
                                                # Give up!
    # Here we go down. Compute all possible next steps
    nextForms = current.possibleDerivs(self.rules)
    # For each check if there's a successful parse to the target string
    for form in nextForms:
        print(current.string, '==>', form)
        # Ok. Did this work? If yes, break out with True and Leave the recursion
        if self.parseDFSAux(form, string, maxDepth-1): return True
    # Didn't find anything, so that's a blind alley
    return False
# The mighty parser Part 2: Breadth First Search version
# Step 1: Wrap the start symbol into a Sentential Form, initialize queue and jump in
def parseBFS(self,string):
    self.queue = [SententialForm(self.start)]
    # Anything Left to check?
    while self.queue:
        # Extract first (head) element from queue (dequeue)
        current = self.queue.pop(0)
        print(current)
        # Check for success
        if current.string == string: return True # Got it! Get outa here...
        # Compute all possible next steps and place them at end of queue (enqueue)
        self.queue.extend(current.possibleDerivs(self.rules))
    # Didn't find anything, queue empty - think about this for a moment!
    return False
```

CFG

```
In [6]: cfg = CFG('S',[])
#cfg.addRule(Rule('S', 'aSa'))
```

```
#cfg.addRule(Rule('S', 'bSb'))
         cfg.addRule(Rule('S', 'aSb'))
         cfg.addRule(Rule('S','c'))
         print(cfg.rules)
         print(cfg.parseDFS('aaaaaaacbbbbbbbb',8))
         cfg.parseBFS('aaaaaaacbbbbbbbb')
         #cfq.parseBFS('aabcbaa')
        [S --> aSb, S --> c]
        S ==> aSb
        aSb ==> aaSbb
        aaSbb ==> aaaSbbb
         aaaSbbb ==> aaaaSbbbb
         aaaaSbbbb ==> aaaaaSbbbbb
         aaaaaSbbbbbb ==> aaaaaaSbbbbbb
        aaaaaaSbbbbbbb ==> aaaaaaaSbbbbbbbb
         aaaaaaaSbbbbbbbb ==> aaaaaaaaSbbbbbbbbb
        aaaaaaaSbbbbbbb ==> aaaaaaacbbbbbbb
        True
        S
         aSb
        C
         aaSbb
        acb
         aaaSbbb
        aacbb
         aaaaSbbbb
        aaacbbb
         aaaaaSbbbbb
        aaaacbbbb
         aaaaaaSbbbbbb
         aaaaacbbbbb
         aaaaaaaSbbbbbbbb
         aaaaaacbbbbbb
         aaaaaaaaSbbbbbbbbb
        aaaaaaacbbbbbbb
Out[6]:
In [ ]:
```