

# Introduction

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I. CONTEXT AND MOTIVATION

II. METHODOLOGY AND CONTRIBUTIONS

III. ORGANIZATION OF THE THESIS

REFERENCES

# State of the art

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## I. IoT APPLICATIONS [1]

- A. Transportation and logistics
- B. Healthcare
- C. Smart environnement
- D. personal and social
- E. Futuristic



Figure 1. Use cases.

Use cases			
Health Monitoring			
Water Distribution			
Electricity Distribution			
Smart Buildings			
Intelligent Transportation			
Surveillance			
Environmental Monitoring			

Table I. Use cases hancke\_role\_2012

voir [4]

Smart systems in smart cities [6]

- » Smart Mobility
- » Smart semaphores controle
- » Smart Red Swarm
- » Smart panels

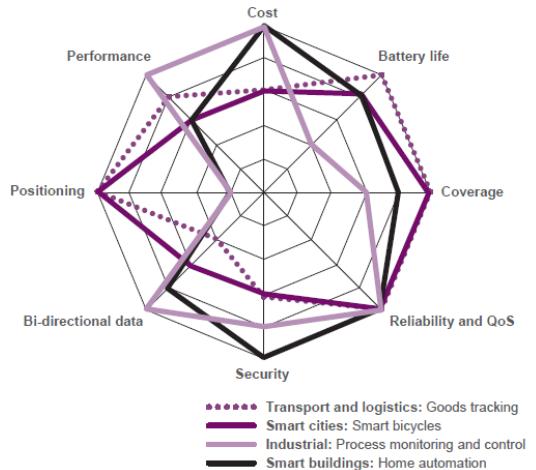


Figure 2. Use cases.

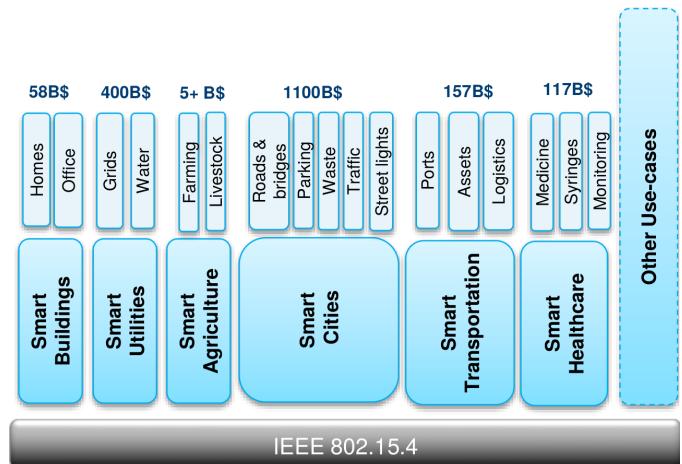


Figure 3. 802.15.4 use cases sarwar\_iot\_.

- » Smart bus scheduling
- » Smart EV management
- » Smart surface parking
- » Smart signs
- » Smart energy systems
- » Smart lighting
- » Smart water jet systems

Challenges-Applications	Gids	EHealth	Transportations	Cities	<b>Building</b>
Ressources constraints	+	+++	-	++	+
Mobility	+	++	+++	+++	-
<b>Heterogeneity</b>	++	++	++	+++	+
Scalability	+++	++	+++	+++	++
QoS constraints	++	++	+++	+++	+++
Data management	++	+	+++	+++	++
Lack of standardization	++	++	++	++	+++
Amount of attacks	+	+	+++	+++	+++
Safety	++	++	+++	++	+++

Table II. Main IoT challenges[2] + [3]

Use Case	Packet rate () [packet/day]	Minimum success rate (Ps,min )	Grouping
Wearables	10	90	Group A PL = 10/20B
Smoke Detectors	2	90	
Smart Grid	10	90	
White Goods	3	90	
Waste Management	24	90	
VIP/Pet Tracking	48	90	Group B PL = 50B
Smart Bicycle	192	90	
Animal Tracking	100	90	
Environmental Monitoring	5	90	
Asset Tracking	100	90	
Smart Parking	60	90	
Alarms/Actuators	5	90	Group C PL = 100/200B
Home Automation	5	90	
Machinery Control	100	90	
Water/Gas Metering	8	90	
Environmental Data Collection	24	90	
Medical Assisted Living	8	90	
Microgeneration	2	90	
Safety Monitoring	2	90	
Propane Tank Monitoring	2	90	
Stationary Monitoring	4	90	
Urban Lighting	5	90	Group D PL = 1KB
Vending Machines Payment	100	90	
Vending Machines General	1	90	

Table III. APPLICATION REQUIREMENTS FOR THE USE CASES OF INTEREST[5] [3].

- Smart residuals gathering
- Smart building construction
- Smart tourism
- Smart QRinfo
- Smart monitoring
- Smart hawkeye

## F. Summary and discussion

### II. IoT SENSORS

#### A. Software platform

The operating system is the foundation of the IoT technology as it provides the functions for the connectivity between the nodes. However, different types of nodes need different levels of OS complexity; a passive node generally only needs the communication stack and is not in need of any threading capabilities, as the program can handle all logic in one function. Active nodes and border routers need to have a much

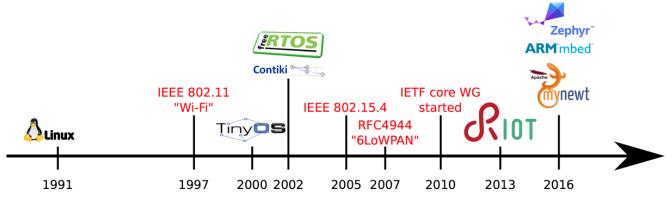


Figure 4. .

more complex OS, as they need to be able to handle several running threads or processes, e.g. routing, data collection and interrupts. To qualify as an OS suitable for the IoT, it needs to meet the basic requirements: Low Random-access memory (RAM) footprint Low Read-only memory (ROM) footprint Multi-tasking Power management (PM) Soft real-time These requirements are directly bound to the type of hardware designed for the IoT. As this type of hardware in general needs to have a small form factor and a long battery life, the on-board memory is usually limited to keep down size and energy consumption. Also, because of the limited amount of memory, the implementation of threads is usually a challenging task, as context switching needs to store thread or process variables to memory. The size of the memory also directly affects the energy consumption, as memory in general is very power hungry during accesses. To be able reduce the energy consumption, the OS needs some kind of power management. The power management does not only let the OS turn on and off peripherals such as flash memory, I/O, and sensors, but also puts the MCU itself in different power modes. As the nodes can be used to control and monitor consumer devices, either a hard or soft real-time OS is required. Otherwise, actions requiring a close to instantaneous reaction might be indefinitely delayed. Hard real-time means that the OS scheduler can guarantee latency and execution time, whereas Soft real-time means that latency and execution time is seen as real-time but can not be guaranteed by the scheduler. Operating systems that meet the above requirements are compared in table 2.1 and 2.2.

1) *Contiki*: Contiki is a embedded operating system developed for IoT written in C [12]. It supports a broad range of MCUs and has drivers for various transceivers. The OS does not only support TCP/IPv4 and IPv6 with the uIP stack [9], but also has support for the 6LoWPAN stack and its own stack called RIME. It supports threading with a thread system called Phototreads [13]. The threads are stack-less and thus use only two bytes of memory per thread; however, each thread is bound to one function and it only has permission to control its own execution. Included in Contiki, there is a range of applications such as a HTTP, Constrained Application Protocol (CoAP), FTP, and DHCP servers, as well as other useful programs and tools. These applications can be included in a project and can run simultaneously with the help of Phototreads. The limitations to what applications can be run is the amount of RAM and ROM the target MCU provides. A standard system with IPv6 networking needs about 10 kB

RAM and 30 kB ROM but as applications are added the requirements tend to grow.

Contiki is an open source operating system for the Internet of Things. Contiki connects tiny low-cost, low-power microcontrollers to the Internet.

2k RAM, 60k ROM; 10k RAM, 48K ROM Portable to tiny low-power micro-controllers I386 based, ARM, AVR, MSP430, ... Implements uIP stack IPv6 protocol for Wireless Sensor Networks (WSN) Uses the protothreads abstraction to run multiple process in an event based kernel. Emulates concurrency Contiki has an event based kernel (1 stack) Calls a process when an event happens

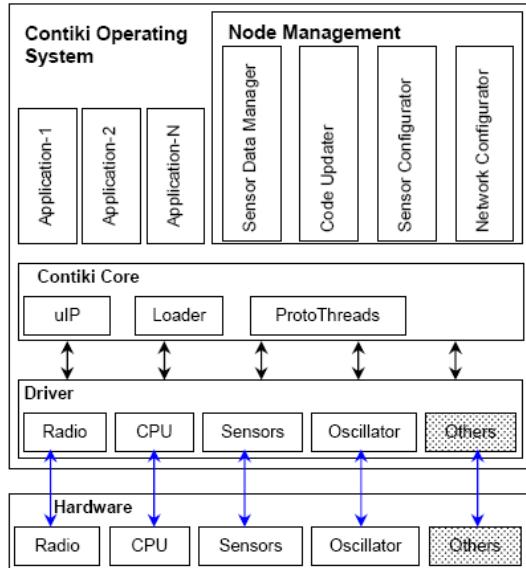


Figure 5. contiki .

*Contiki size:* One of the main aspect of the system, is the modularity of the code. Besides the system core, each program builds only the necessary modules to be able to run, not the entire system image. This way, the memory used from the system, can be reduced to the strictly necessary. This methodology makes more practical any change in any module, if it is needed. The code size of Contiki is larger than that of TinyOS, but smaller than that of the Mantis system. Contiki's event kernel is significantly larger than that of TinyOS because of the different services provided. While the TinyOS event kernel only provides a FIFO event queue scheduler , the Contiki kernel supports both FIFO events and poll handlers with priorities. Furthermore, the flexibility in Contiki requires more run-time code than for a system like TinyOS, where compile time optimization can be done to a larger extent.

The documentation in the doc folder can be compiled, in order to get the html wiki of all the code. It needs doxygen installed, and to run the command make html. This will create a new folder, doc/html, and in the index.html file, the wiki can be opened.

*Contiki Hardware:* Contiki can be run in a number of platforms, each one with a different CPU. Tab.7 shows the

hardware platforms currently defined in the Contiki code tree. All these platforms are in the platform folder of the code.

#### Kernel structure:

2) RIOT: RIOT is a open source embedded operating system supported by Freie Universität Berlin, INIRA, and Hamburg University of Applied Sciences [14]. The kernel is written in C but the upper layers support C++ as well. As the project originates from a project with real-time and reliability requirements, the kernel supports hard real-time multi-tasking scheduling. One of the goals of the project is to make the OS completely POSIX compliant. Overhead for multi-threading is minimal with less than 25 bytes per thread. Both IPv6 and 6LoWPAN is supported together with UDP, TCP, and IPv6 Routing Protocol for Low-Power and Lossy Networks (RPL); and CoAP and Concise Binary Object Representation (CBOR) are available as application level communication protocols.

3) TinyOS: TinyOS is written in Network Embedded Systems C (nesC) which is a variant of C [15]. nesC does not have any dynamic memory allocation and all program paths are available at compile-time. This is manageable thanks to the structure of the language; it uses modules and interfaces instead of functions [16]. The modules use and provide interfaces and are interconnected with configurations; this procedure makes up the structure of the program. Multitasking is implemented in two ways: through tasks and events. Tasks, which focus on computation, are non-preemptive, and run until completion. In contrast, events which focus on external events i.e. interrupts, are preemptive, and have separate start and stop functions. The OS has full support for both 6LoWPAN and RPL, and also have libraries for CoAP.

4) freeRTOS: One of the more popular and widely known operating systems is freeRTOS [17]. Written in C with only a few source files, it is a simple but powerful OS, easy to overview and extend. It features two modes of scheduling, pre-emptive and co-operative, which may be selected according to the requirements of the application. Two types of multitasking are featured: one is a lightweight Co-routine type, which has a shared stack for lower RAM usage and is thus aimed to be used on very small devices; the other is simply called Task, has its own stack and can therefore be fully pre-empted. Tasks also support priorities which are used together with the pre-emptive scheduler. The communication methods supported out-of-the-box are TCP and UDP.

Plan de contrôle	Plan de gestion	Plan de données
Contrôle d'admission	Contrôle et supervision de QoS	Contrôle du trafic
Réservation de ressources	Gestion de contrats	Façonnage du trafic
Routage	QoS mapping	Contrôle de congestion
Signalisation	Politique de QoS	Classification de paquets
		Marquage de paquets
		Ordonnancements des paquets
		Gestion de files d'attente

Table IV. An example table.

5) SDN platforms: Sensor OpenFlow [20,21] SDWN [60] Smart [14] SDN-WISE [78] SDCSN [88] TinySDN [69,118]

Virtual Overlay [59,87,90] Multi-task [122] SDWSN-RL [123]  
 Wireless power transfer [126] Function alternation [65] Statistical machine learning [24] Context-based [91,92] Soft-WSN [9]

- [7] Many studies have identified **SDN** as a potential solution to the WSN challenges, as well as a model for **heterogeneous** integration.
- [7] This **shortfall** can be resolved by using the **SDN approach**.
- [8] **SDN** also enhances better control of **heterogeneous** network infrastructures.
- [8] Anadiotis et al. define a **SDN operating system for IoT** that integrates SDN based WSN (**SDN-WISE**). This experiment shows how **heterogeneity** between different kinds of SDN networks can be achieved.
- [8] In cellular networks, OpenRoads presents an approach of introducing **SDN** based **heterogeneity** in wireless networks for operators.
- [9] There has been a plethora of (industrial) studies **synergising SDN in IoT**. The major characteristics of IoT are low latency, wireless access, mobility and **heterogeneity**.
- [9] Thus a bottom-up approach application of **SDN** to the realisation of **heterogeneous IoT** is suggested.
- [9] Perhaps a more complete IoT architecture is proposed, where the authors apply **SDN** principles in IoT **heterogeneous** networks.
- [10] it provides the **SDWSN** with a proper model of network management, especially considering the potential of **heterogeneity** in SDWSN.
- [10] We conjecture that the **SDN paradigm** is a good candidate to solve the **heterogeneity** in IoT.

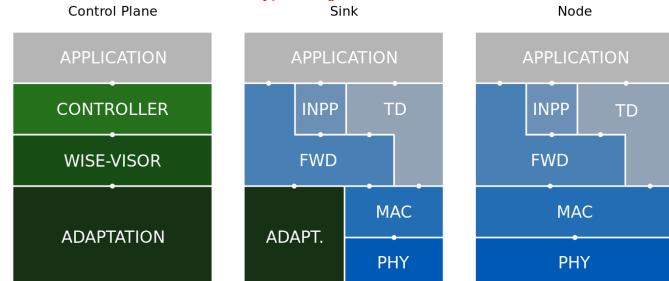


Figure 6. LPWAN connectivity.

## 6) Summary and conclusion:

### B. Hardware platform

1) *Processing Unit*: Even though the hardware is in one sense the tool that the OS uses to make IoT possible, it is still very important to select a platform that is future-proof and extensible. To be regarded as an extensible platform, the hardware needs to have I/O connections that can be used by external peripherals. Amongst the candidate interfaces are Serial Peripheral Interface (SPI), Inter-Integrated Circuit (I<sup>2</sup>C), and Controller Area Network (CAN). These interfaces allow developers to attach custom-made PCBs with sensors for monitoring or actuators for controlling the environment. The best practice is to implement an extension socket with a well-known form factor. A future-proof device is specified

Management architecture	Management feature	Controller configuration	Traffic Control	Configuration and monitoring	Scalability and localization	Communication management
[11] Sensor Open Flow	SDN support protocol	Distributed	in/out-band	✓	✓	✓
[12] SDWN	Duty cycling, aggregation, routing	Centralized	in-band	✓		
[13] SDN-WISE	Programming simplicity and aggregation	Distributed	in-band		✓	
deganeti smart 2014 Smart	in resource allocation	Distributed	in-band		✓	
SDCSN	Network reliability and QoS	Distributed	in-band		✓	
TinySDN	In-band-traffic control	Distributed	in-band		✓	
Virtual Overlay	Network flexibility	Distributed	in-band		✓	
Context based	Network scalability and performance	Distributed	in-band		✓	
CRLB	Node localization	Centralized	in-band			
Multi-hop	Traffic and energy control	Centralized	in-band			✓
Tiny-SDN	Network task measurement	-	in-band			

Table V. SDN-based network and topology management architectures. [9]

Architecture	LiteOS	Nano-RK	MANTIS	Contiki
Scheduling	Monolithic	Layered	Modular	Modular
Memory	Round Robin	Monotonic harmonized	Priority classes	Interrupts execute w.r.t.
Network	File	Socket abstraction	At Kernel COMM layer	uIP, Rime
Virtualization and Completion	Synchronization primitives	Serialized access semaphores	Semaphores	Serialized, Access
Multi-threading	✓	✓	✗	✓
Dynamic protection	✓	✗	✓	✓
Memory Stack	✓	✗	✗	✗

Table VI. Common operating systems used in IoT environment  
 al-fuqaha\_internet\_24

as a device that will be as attractive in the future as it is today. For hardware, this is very hard to achieve as there is constant development that follows Moores Law [4]; however, the most important aspects are: the age of the chip, its expected remaining lifetime, and number of current implementations i.e. its popularity. If a device is widely used by consumers, the lifetime of the product is likely to be extended. One last thing to take into consideration is the product family; if the chip

OS	Contiki	MANTIS	Nano-RK	LiteOS
<b>Architecture</b>	Modular	Modular	Layered	Monolithic
<b>Multi threading</b>	✓	✗	✓	✓
<b>Scheduling</b>	Interrupts	Priority	Monotonic	Round Robin
	execute w.r.t.	classes	harmonized	
<b>Dynamic Memory</b>	✓	✓	✗	✓
<b>Memory protection</b>	✗	✗	✗	✓
<b>Network Stack</b>	uIP/Rime	At Kernel-/COMM layer	Socket abstraction	file
<b>Virtualization and Completion</b>	Serialized/Access	Semaphores	Serialized/semaphores	Synchronization primitives

Table VII. Common operating systems used in IoT environment  
al-fuqaha\_internet\_24

belongs to a family with several members the transition to a newer chip is usually easier.

a) **OpenMote**: OpenMote is based on the TI CC2538 System on Chip (SoC), which combines an ARM Cortex-M3 with a IEEE 802.15.4 transceiver in one chip [18, 19]. The board follows the XBee form factor for easier extensibility, which is used to connect the core board to either the OpenBattery or OpenBase extension boards [20, 21]. It originates from the CC2538DK which was used by Thingsquare to demo their Mist IoT solution [22]. Hence, the board has full support for Contiki, which is the foundation of Thingsquare. It can run both as a battery-powered sensor board and as a border router, depending on what extension board it is attached to, e.g. OpenBattery or OpenBase. Furthermore, the board has limited support but ongoing development for RIOT and also full support for freeRTOS.

b) **MSB430-H**: The Modular Sensor Board 430-H from Freie Universität Berlin was designed for their ScatterWeb project [23]. As the university also hosts the RIOT project, the decision to support RIOT was natural. The main board has a TI MSP430F1612 MCU [24], a **TI CC1100 transceiver**, and a battery slot for dual AA batteries; it also includes a SHT11 temperature and humidity sensor and a MMA7260Q accelerometer to speed up early development. All GPIO pins and buses are connected to external pins for extensibility. Other modules with new peripherals can then be added by making a PCB that matches the external pin layout.

c) **Zolertia**: As many other Wireless Sensor Network (WSN) products, the Zolertia Z1 builds upon the MSP430 MCU [25, 26]. The communication is managed by the TI CC2420 which operates in the 2.4 GHz band. The platform includes two sensors: the SHT11 temperature and humidity sensor and the MMA7600Q accelerometer. Extensibility is ensured with: two connections designed especially for external sensors, an external connector with USB, Universal asynchronous **receiver/transmitter (UART)**, SPI, and I2C.

## 2) Radio Unit:

a) **Lora Tranceiver**: To limit the complexity of the radio unit:

- ➡ limiting message size: maximum application payload size between 51 and 222 bytes, depending on the spreading

factor

- ➡ using simple channel codes: Hamming code
- ➡ supporting only half-duplex operation
- ➡ using one transmit-and-receive antenna
- ➡ on-chip integrating power amplifier (since transmit power is limited)
- ➡ message size: maximum application payload size between 51 and 222 bytes, depending on the spreading factor
- ➡ using simple channel codes: Hamming code supporting only half-duplex operation using one transmit-and-receive antenna
- ➡ on-chip integrating power amplifier (since transmit power is limited)

Ref	Module	Frequency MHz	Tx power	Rx power	Sensitivity	Channels	Distance
[14]	Semtech SX1272	863-870 (EU) 902-928 (US)	14 dBm	dBm	-134 dBm	8 13	22+ km
[14]	rn2483						

Table VIII

## 3) Sensing Unit:

- a) **GPS**:
- b) **Humidity**:
- c) **Temperature**:

## C. Summary and discussion

### III. IoT COMMUNICATION PROTOCOLS

Application protocol	DDS	CoAP	AMQP	MQTT	MQTT-SN	XMP	DNS-SD
Service discovery		mDNS					
Transport							
Network		IPv6 RPL					IPv4/IPv6
		6LowPan					RFC 2464
MAC		IEEE 802.15.4	IEEE 802.11 (Wi-Fi)	IEEE 802.3 (Ethernet)			
		2.4GHz, 915, 868MHz	2.4, 5GHz				
		DSS, FSK, OFDM	CSMA/CA	CUTP, FO			

Table IX. Standardization efforts that support the IoT

## A. Application

- 1) **LwM2M**:
- 2) **CBOR**:
- 3) **DTLS**:
- 4) **OSCOAP**:

5) **COAP (COnstrained Application Protocol)**: The Constrained Application Protocol (CoAP) is a specialized web transfer protocol for use with constrained nodes and constrained networks in the Internet of Things. More detailed information about the protocol is given in the Contiki OS CoAP section.

a) **Overview**: Like HTTP, CoAP is a document transfer protocol. Unlike HTTP, CoAP is designed for the needs of constrained devices. The packets are much smaller than HTTP TCP flows. Packets are simple to generate and can be parsed in place without consuming extra RAM in constrained devices. CoAP runs over UDP, not TCP. Clients and servers communicate through connectionless datagrams. Retries and reordering are implemented in the application stack. It follows

a client/server model. Clients make requests to servers, servers send back responses. Clients may GET, PUT, POST and DELETE resources. CoAP implements the REST model from HTTP, with the primitives GET, POST, PUT and DELETE.

*b) Coap Methods:* CoAP extends the HTTP request model with the ability to observe a resource. When the observe flag is set on a CoAP GET request, the server may continue to reply after the initial document has been transferred. This allows servers to stream state changes to clients as they occur. Either end may cancel the observation. CoAP defines a standard mechanism for resource discovery. Servers provide a list of their resources (along with metadata about them) at /.well-known/core. These links are in the application/link-format media type and allow a client to discover what resources are provided and what media types they are.

*c) Coap Transactions:*

*d) Coap Messages:* The CoAP message structure is designed to be simpler than HTTP, for reduced transmission data. Each field responds to a specific purpose.

- » Constrained Application Protocol
- » The IETF Constrained RESTful Environments
- » CoAP is bound to UDP
- » CoAP can be divided into two sub-layers
  - » messaging sub-layer
  - » request/response sub-layer
    - a) Confirmable.
    - b) Non-confirmable.
    - c) Piggybacked responses.
    - d) Separate response
- » CoAP, as in HTTP, uses methods such as:
  - » GET, PUT, POST and DELETE to
  - » Achieve, Create, Retrieve, Update and Delete
  - » Ex: the GET method can be used by a server to inquire the clients temperature

*6) MQTT:*

- » Message Queue Telemetry Transport
- » Andy Stanford-Clark of IBM and Arlen Nipper of Arcom
  - » Standardized in 2013 at OASIS
- » MQTT uses the publish/subscribe pattern to provide transition flexibility and simplicity of implementation
- » MQTT is built on top of the TCP protocol
- » MQTT delivers messages through three levels of QoS
- » Specifications
  - » MQTT v3.1 and MQTT-SN (MQTT-S or V1.2)
  - » MQTT v3.1 adds broker support for indexing topic names
- » The publisher acts as a generator of interesting data.

*7) XMPP:*

- » Extensible Messaging and Presence Protocol
- » Developed by the Jabber open source community
- » An IETF instant messaging standard used for:
  - » multi-party chatting, voice and telepresence
- » Connects a client to a server using a XML stanzas
- » An XML stanza is divided into 3 components:
  - » message: fills the subject and body fields
  - » presence: notifies customers of status updates

» iq (info/query): pairs message senders and receivers

» Message stanzas identify:

- » the source (from) and destination (to) addresses
- » types, and IDs of XMPP entities

*8) AMQP:*

- » Advanced Message Queuing Protocol
- » Communications are handled by two main components
  - » exchanges: route the messages to appropriate queues.
  - » message queues: Messages can be stored in message queues and then be sent to receivers
- » It also supports the publish/subscribe communications.
- » It defines a layer of messaging on top of its transport layer.
- » AMQP defines two types of messages
  - » bare messages: supplied by the sender
  - » annotated messages: seen at the receiver
- » The header in this format conveys the delivery parameters:
  - » durability, priority, time to live, first acquirer & delivery count.
- » AMQP frame format
  - Size the frame size.

DOFF the position of the body inside the frame.

Type the format and purpose of the frame.

- \* Ex: 0x00 show that the frame is an AMQP frame
- \* Ex: 0x01 represents a SASL frame.

*9) DDS:*

- » Data Distribution Service
- » Developed by Object Management Group (OMG)
- » Supports 23 QoS policies:
  - » like security, urgency, priority, durability, reliability, etc
- » Relies on a broker-less architecture
  - » uses multicasting to bring excellent Quality of Service
  - » real-time constraints
- » DDS architecture defines two layers:

DLRL Data-Local Reconstruction Layer

- \* serves as the interface to the DCPS functionalities

DCPS Data-Centric Publish/Subscribe

- \* delivering the information to the subscribers

- » 5 entities are involved with the data flow in the DCPS layer:
  - » Publisher: disseminates data
  - » DataWriter: used by app to interact with the publisher
  - » Subscriber: receives published data and delivers them to app
  - » DataReader: employed by Subscriber to access received data
  - » Topic: relate DataWriters to DataReaders
- » No need for manual reconfiguration or extra administration
- » It is able to run without infrastructure
- » It is able to continue working if failure happens.
- » It inquires names by sending an IP multicast message to all the nodes in the local domain
  - » Clients asks devices that have the given name to reply back

Application protocol	Rest-Full	Transport	Publish-/Subscribe	Request/Response	Security	QoS	Header size (Byte)
COAP	✓	UDP	✓	✓	DTLS	✓	4
MQTT	✗	TCP	✓	✗	SSL	✓	2
MQTT-SN	✗	TCP	✓	✗	SSL	✓	2
XMPP	✗	TCP	✓	✓	SSL	✗	-
AMQP	✗	TCP	✓	✗	SSL	✓	8
DDS	✗	UDP TCP	✓	✗	SSL DTLS	✓	-
HTTP	✓	TCP	✗	✓	SSL	✗	-

Table X. Application protocols comparison

- the target machine receives its name and multicasts its IP @
- Devices update their cache with the given name and IP @

#### 10) mDNS:

- Requires zero configuration aids to connect machine
- It uses mDNS to send DNS packets to specific multicast addresses through UDP
- There are two main steps to process Service Discovery:
  - finding host names of required services such as printers
  - pairing IP addresses with their host names using mDNS
- Advantages
  - IoT needs an architecture without dependency on a configuration mechanism
  - smart devices can join the platform or leave it without affecting the behavior of the whole system
- Drawbacks
  - Need for caching DNS entries

### B. Network

#### 1) 6TiSCH:

#### 2) OLSRv2:

#### 3) AODVv2:

#### 4) LoRaWAN:

#### 5) ROHC:

#### 6) IPHC:

#### 7) SCHC:

#### 8) NHC:

#### 9) ROLL:

10) RPL: RPL is a Distance Vector IPv6 routing protocol for LLNs that specifies how to build a Destination Oriented Directed Acyclic Graph (DODAG) using an objective function and a set of metrics/constraints. The objective function operates on a combination of metrics and constraints to compute the best path.

An RPL Instance consists of multiple Destination Oriented Directed Acyclic Graphs (DODAGs). Traffic moves either up towards the DODAG root or down towards the DODAG leafs. The graph building process starts at the root or LBR (LowPAN Border Router). There could be multiple roots configured in the system. The RPL routing protocol specifies a set of ICMPv6 control messages to exchange graph related information. These messages are called DIS (DODAG Information Solicitation), DIO (DODAG Information Object) and DAO (DODAG Destination Advertisement Object). The root starts

advertising the information about the graph using the DIO message. The nodes in the listening vicinity (neighbouring nodes) of the root will receive and process DIO messages potentially from multiple nodes and makes a decision based on certain rules (according to the objective function, DAG characteristics, advertised path cost and potentially local policy) whether to join the graph or not. Once the node has joined a graph it has a route toward the graph (DODAG) root. The graph root is termed as the parent of the node. The node computes the rank of itself within the graph, which indicates the coordinates of the node in the graph hierarchy. If configured to act as a router, it starts advertising the graph information with the new information to its neighbouring peers. If the node is a leaf node, it simply joins the graph and does not send any DIO message. The neighbouring peers will repeat this process and do parent selection, route addition and graph information advertisement using DIO messages. This rippling effect builds the graph edges out from the root to the leaf nodes where the process terminates. In this formation each node of the graph has a routing entry towards its parent (or multiple parents depending on the objective function) in a hop-by-hop fashion and the leaf nodes can send a data packet all the way to root of the graph by just forwarding the packet to its immediate parent. This model represents a MP2P (Multipoint-to-point) forwarding model where each node of the graph has reachability toward the graph root. This is also referred to as UPWARD routing. Each node in the graph has a rank that is relative and represents an increasing coordinate of the relative position of the node with respect to the root in graph topology. The notion of rank is used by RPL for various purposes including loop avoidance. The MP2P flow of traffic is called the up direction in the DODAG.

The DIS message is used by the nodes to proactively solicit graph information (via DIO) from the neighbouring nodes should it become active in a stable graph environment using the poll or pull model of retrieving graph information or in other conditions. Similar to MP2P or up direction of traffic, which flows from the leaf towards the root there is a need for traffic to flow in the opposite or down direction. This traffic may originate from outside the LLN network, at the root or at any intermediate nodes and destined to a (leaf) node. This requires a routing state to be built at every node and a mechanism to populate these routes. This is accomplished by the DAO (Destination Advertisement Object) message. DAO messages are used to advertise prefix reachability towards the leaf nodes in support of the down traffic. These messages carry prefix information, valid lifetime and other information about the distance of the prefix. As each node joins the graph it will send DAO message to its parent set. Alternately, a node or root can poll the sub-dag for DAO message through an indication in the DIO message. As each node receives the DAO message, it processes the prefix information and adds a routing entry in the routing table. It optionally aggregates the prefix information received from various nodes in the subdag and sends a DAO message to its parent set. This process continues until the prefix information reaches the root and a complete path to the

prefix is setup. Note that this mode is called the storing mode of operation where intermediate nodes have available memory to store routing tables. RPL also supports another mode called non-storing mode where intermediate node do not store any routes.

11) *6LowPAN*: 6LoWPAN is a networking technology or adaptation layer that allows IPv6 packets to be carried efficiently within a small link layer frame, over IEEE 802.15.4 based networks. As the full name implies, IPv6 over Low-Power Wireless Personal Area Networks, it is a protocol for connecting wireless low power networks using IPv6.

As the full name implies, IPv6 over Low-Power Wireless Personal Area Networks, it is a protocol for connecting wireless low power networks using IPv6.

a) *Characteristics*:

- ➡ Compression of IPv6 and UDP/ICMP headers
- ➡ Fragmentation / reassembly of IPv6 packets
- ➡ Mesh addressing
- ➡ Stateless auto configuration
- ➡

b) *Encapsulation Header format*: All LowPAN encapsulated datagrams are prefixed by an encapsulation header stack. Each header in the stack starts with a header type field followed by zero or more header fields.

c) *Fragment Header*: The fragment header is used when the payload is too large to fit in a single IEEE 802.15.4 frame. The Fragment header is analogous to the IEEE 1394 Fragment header and includes three fields: Datagram Size, Datagram Tag, and Datagram Offset. Datagram Size identifies the total size of the unfragmented payload and is included with every fragment to simplify buffer allocation at the receiver when fragments arrive out-of-order. Datagram Tag identifies the set of fragments that correspond to a given payload and is used to match up fragments of the same payload. Datagram Offset identifies the fragments offset within the unfragmented payload and is in units of 8-byte chunks.

d) *Mesh addressing header*: The Mesh Addressing header is used to forward 6LoWPAN payloads over multiple radio hops and support layer-two forwarding. The mesh addressing header includes three fields: Hop Limit, Source Address, and Destination Address. The Hop Limit field is analogous to the IPv6 Hop Limit and limits the number of hops for forwarding. The Hop Limit field is decremented by each forwarding node, and if decremented to zero the frame is dropped. The source and destination addresses indicate the end-points of an IP hop. Both addresses are IEEE 802.15.4 link addresses and may carry either a short or extended address.

e) *Header compression (RFC4944)*: RFC 4944 defines HC1, a stateless compression scheme optimized for link-local IPv6 communication. HC1 is identified by an encoding byte following the Compressed IPv6 dispatch header, and it operates on fields in the upper-layer headers. 6LoWPAN elides some fields by assuming commonly used values. For example, it compresses the 64-bit network prefix for both source and destination addresses to a single bit each when they carry the well-known link-local prefix. 6LoWPAN compresses the Next

Header field to two bits whenever the packet uses UDP, TCP, or ICMPv6. Furthermore, 6LoWPAN compresses Traffic Class and Flow Label to a single bit when their values are both zero. Each compressed form has reserved values that indicate that the fields are carried inline for use when they dont match the elided case. 6LoWPAN elides other fields by exploiting cross-layer redundancy. It can derive Payload Length which is always elided from the 802.15.4 frame or 6LoWPAN fragmentation header. The 64-bit interface identifier (IID) for both source and destination addresses are elided if the destination can derive them from the corresponding link-layer address in the 802.15.4 or mesh addressing header. Finally, 6LoWPAN always elides Version by communicating via IPv6.

The HC1 encoding is shown in Figure 11. The first byte is the dispatch byte and indicates the use of HC1. Following the dispatch byte are 8 bits that identify how the IPv6 fields are compressed. For each address, one bit is used to indicate if the IPv6 prefix is linklocal and elided and one bit is used to indicate if the IID can be derived from the IEEE 802.15.4 link address. The TF bit indicates whether Traffic Class and Flow Label are both zero and elided. The two Next Header bits indicate if the IPv6 Next Header value is 7UDP, TCP, or ICMP and compressed or carried inline. The HC2 bit indicates if the next header is compressed using HC2. Fully compressed, the HC1 encoding reduces the IPv6 header to three bytes, including the dispatch header. Hops Left is the only field always carried inline.

RFC 4944 uses stateless compression techniques to reduce the overhead of UDP headers. When the HC2 bit is set in the HC1 encoding, an additional 8-bits is included immediately following the HC1 encoding bits that specify how the UDP header is compressed. To effectively compress UDP ports, 6LoWPAN introduces a range of wellknown ports (61616 - 61631). When ports fall in the well-known range, the upper 12 bits may be elided. If both ports fall within range, both Source and Destination ports are compressed down to a single byte. HC2 also allows elision of the UDP Length, as it can be derived from the IPv6 Payload Length field.

The best-case compression efficiency occurs with link-local unicast communication, where HC1 and HC2 can compress a UDP/IPv6 header down to 7 bytes. The Version, Traffic Class, Flow Label, Payload Length, Next Header, and linklocal prefixes for the IPv6 Source and Destination addresses are all elided. The suffix for both IPv6 source and destination addresses are derived from the IEEE 802.15.4 header.

However, RFC 4944 does not efficiently compress headers when communicating outside of link-local scope or when using multicast. Any prefix other than the linklocal prefix must be carried inline. Any suffix must be at least 64 bits when carried inline even if derived from a short 802.15.4 address. As shown in Figure 8, HC1/HC2 can compress a link-local multicast UDP/IPv6 header down to 23 bytes in the best case. When communicating with nodes outside the LoWPAN, the IPv6 Source Address prefix and full IPv6 Destination Address must be carried inline.

*f) Header compression Improved (draft-hui-6lowpan-hc-01):* To provide better compression over a broader range of scenarios, the 6LoWPAN working group is standardizing an improved header compression encoding format, called HC. The format defines a new encoding for compressing IPv6 header, called IPHC. The new format allows Traffic Class and Flow Label to be individually compressed, Hop Limit compression when common values (E.g., 1 or 255) are used, makes use of shared-context to elide the prefix from IPv6 addresses, and supports multicast addresses most often used for IPv6 ND and SLAAC. Contexts act as shared state for all nodes within the LoWPAN. A single context holds a single prefix. IPHC identifies the context using a 4-bit index, allowing IPHC to support up to 16 contexts simultaneously within the LoWPAN. When an IPv6 address matches a contexts stored prefix, IPHC compresses the prefix to the contexts 4-bit identifier. Note that contexts are not limited to prefixes assigned to the LoWPAN but can contain any arbitrary prefix. As a result, share contexts can be configured such that LoWPAN nodes can compress the prefix in both Source and Destination addresses even when communicating with nodes outside the LoWPAN.

The improved header compression encoding is shown in Figure 8. The first three bits (011) form the header type and indicate the use of IPHC. The TF bits indicate whether the Traffic Class and/or Flow Label fields are compressed. The HLIM bits indicate whether the Hop Limit takes the value 1 or 255 and compressed, or carried inline.

Bits 8-15 of the IPHC encoding indicate the compression methods used for the IPv6 Source and Destination Addresses. When the Context Identifier (CID) bit is zero, the default context may be used to compress Source and/or Destination Addresses. This mode is typically when both Source and Destination Addresses are assigned to nodes in the same LoWPAN. When the CID bit is one, two additional 4-bit fields follow the IPHC encoding to indicate which one of 16 contexts is in use for the source and destination addresses. The Source Address Compression (SAC) indicates whether stateless compression is used (typically for link-local communication) or stateful context-based compression is used (typically for global communication). The Source Address Mode (SAM) indicates whether the full Source Address is carried inline, upper 16 or 64-bits are elided, or the full Source Address is elided. When SAC is set and the Source Addresses prefix is elided, the identified context is used to restore those bits. The Multicast (M) field indicates whether the Destination Address is a unicast or multicast address. When the Destination Address is a unicast address, the DAC and DAM bits are analogous to the SAC and SAM bits. When the Destination Address is a multicast address, the DAM bits indicate different forms of multicast compression. HC also defines a new framework for compressing arbitrary next headers, called NHC. HC2 in RFC 4944 is only capable of compressing UDP, TCP, and ICMPv6 headers, the latter two are not yet defined. Instead, the NHC header defines a new variable length Next Header identifier, allowing for future definition of arbitrary next header compression encodings. HC initially defines a compression

encoding for UDP headers, similar to that defined in RFC 4944. Like RFC 4944, HC utilizes the same well-known port range (61616-61631) to effectively compress UDP ports down to 4-bits each in the best case. However, HC no longer provides an option to carry the Payload Length in line, as it can always be derived from the IPv6 header. Finally, HC allows elision of the UDP Checksum whenever an 10upper layer message integrity check covers the same information and has at least the same strength. Such a scenario is typical when transportor application-layer security is used. As a result, the UDP header can be compressed down to two bytes in the best case.

Routing protocol	Control Cost	Link Cost	Node Cost
OSPF/IS-IS	✗	✓	✗
OLSRv2	?	✓	✓
RIP	✓	?	✗
DSR	✓	✗	✗
RPL	✓	✓	✓

Table XI. Routing protocols comparison \_rpl2\_

- ➡ Routing over low-power and lossy links (ROLL)
  - ➡ Support minimal routing requirements.
  - ➡ like multipoint-to-point, point-to-multipoint and point-to-point.
  - ➡ A Destination Oriented Directed Acyclic Graph (DODAG)
    - ➡ Directed acyclic graph with a single root.
    - ➡ Each node is aware of its parents
    - ➡ but not about related children
  - ➡ RPL uses four types of control messages
    - ➡ DODAG Information Object (DIO)
    - ➡ Destination Advertisement Object (DAO)
    - ➡ DODAG Information Solicitation (DIS)
    - ➡ DAO Acknowledgment (DAO-ACK)
  - ➡ Standard topologies to form IEEE 802.15.4e networks are Star contains at least one FFD and some RFDs
- Mesh contains a PAN coordinator and other nodes communicate with each other
- Cluster consists of a PAN coordinator, a cluster head and normal nodes.
- ➡ The IEEE 802.15.4e standard supports 2 types of network nodes
- FFD Full function device: serve as a coordinator
- \* It is responsible for creation, control and maintenance of the net
  - \* It stores a routing table in their memory and implement a full MAC
- RFD Reduced function devices: simple nodes with restricted resources
- \* They can only communicate with a coordinator
  - \* They are limited to a star topology

Routing protocol	Control Cost	Link Cost	Node Cost
OSPF/IS-IS	✗	✓	✗
OLSRv2	?	✓	✓
RIP	✓	?	✗
DSR	✓	✗	✗
RPL	✓	✓	✓

Table XII. Routing protocols comparison \_rpl2\_

### C. MAC

Channel based	FDMA	OFDMA WDMA SC-FDMA		
	TDMA	MF-TDMA STDMA		
	CDMA	W-CDMA TD-CDMA TD-SCDMA DS-CDMA FH-CDMA MC-CDMA		
	SDMA	HC-SDMA		
Packet-based	Collision recovery	ALOHA Slotted ALOHA R-ALOHA AX.25 CSMA/CD		
	Collision avoidance	MACA MACAW CSMA CSMA/CA DCF PCF HCF CSMA/CARP		
	Collision-free	Token ring Token bus MS-ALOHA		
Duplexing methods	Delay and disruption tolerant	MANET VANET DTN Dynamic Source Routing		

Table XIII

1) *Sharing the channel:*

a) TDMA, FDMA, CDMA, TSMA:

2) *Transmitting information:*

a) TFDM, TDSSS, TFHSS:

### D. Radio

1) *Digital modulation:*

a) ASK, APSK, CPM, FSK, MFSK, MSK, OOK, PPM, PSK, QAM, SC-FDE, TCM WDM:

2) *Hierarchical modulation:*

a) QAM, WDM:

3) *Spread spectrum:*

a) SS, DSSS, FHSS, THSS:

4) *Radio performance:*

a) *Power Level (dB):* The dB measures the power of a signal as a function of its ratio to another standardized value. The abbreviation dB is often combined with other abbreviations to represent the values that are compared. Here are two examples:

- ⇒ dBm The dB value is compared to 1 mW.
- ⇒ dBw The dB value is compared to 1 W.

$$Power(indB) = 10 * \log(10)(Signal/Reference) \quad (1)$$

Where:

- ⇒ log(10) is logarithm base 10.
- ⇒ Signal is the power of the signal.
- ⇒ Reference is the reference power.

For example, if you want to calculate the power in dB of 50 mW:

$$\text{Power (in dB)} = 10 * \log(10) (50/1) = 10 * 1.7 = 17 \text{ dBm}$$

b) *Receive Signal Strength Indicator RSSI:* Receiver sensitivity is defined as the minimum signal power level with an acceptable Bit Error Rate (in dBm or mW) that is necessary for the receiver to accurately decode a given signal. This is usually expressed in a negative number depending on the data rate. For example an Access Point may require an RSSI of at least negative -91 dBm at 1 MB and an even higher strength RF power -79 dBm to decode 54 MB.

c) *Signal to Noise Ratio SNR:* Noise is any signal that interferes with your signal. Noise can be due to other wireless devices such as cordless phones, microwave devices, etc. This value is measured in decibels from 0 (zero) to -120 (minus 120). Noise level is the amount of interference in your wireless signal, so lower is generally good for WLAN deployments. Typical environments range between -90dBm and -98dBm with little ambient noise. This value may be even higher if there is a lot of RF interference coming in from other non-802.11 devices on the same spectrum. Signal to Noise Ratio or SNR is defined as the ratio of the transmitted power from the AP to the ambient (noise floor) energy present. To calculate the SNR value, we add the Signal Value to the Noise Value to get the SNR ratio. A positive value of the SNR ratio is always better. For example, say your Signal value is -55dBm and your Noise value is -95dBm. The difference of Signal (-55dBm) + Noise (-95dBm) = 40db. This means you have an SNR of 40. Note that in the above equation you are not merely adding two numbers, but are interested in the difference between the Signal and Noise values, which is usually a positive number. The lower the number, the lower the difference between noise and transmitted power, which in turn means lower quality of signal. The higher the difference between Signal and Noise means that the transmitted signal is of much higher power than the noise floor, thereby making it easier for a WLAN client to decode the signal.

d) *Signal Attenuation:* Signal attenuation or signal loss occurs even as the signal passes through air. The loss of signal strength is more pronounced as the signal passes through different objects. A transmit power of 20 mW is equivalent to 13 dBm. Therefore if the transmitted power at the entry point of a plasterboard wall is at 13 dBm, the signal strength will be reduced to 10 dBm when exiting that wall. Some common examples are shown in Table 10-5.

### E. Summary and discussion

## IV. IOT NORMS & STANDARDS

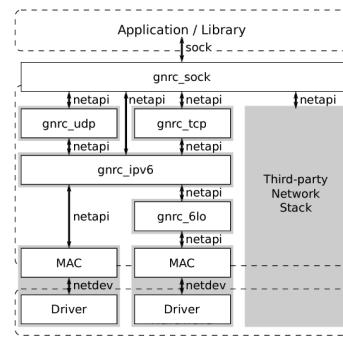


Figure 7. LPWAN.

Several different wireless communication protocols, such as Wireless LAN (WLAN), BLE, 6LoWPAN, and ZigBee may be suitable for IoT applications. They all operate in the 2.4GHz frequency band and this, together with the limited output power in this band, means that they all have a similar

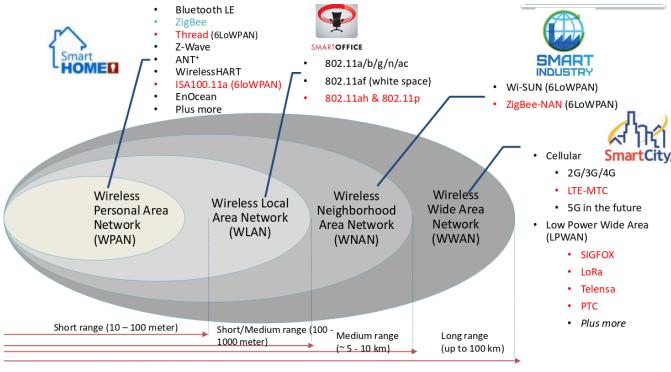


Figure 8. LPWAN.

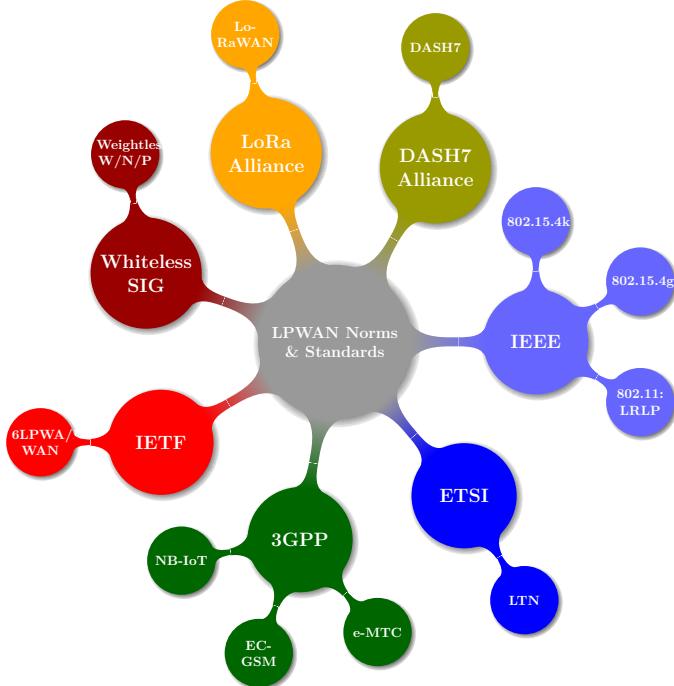


Figure 9. LPWAN.

range. The main differences are located in the MAC, PHY, and network layer. WLAN is much too power hungry as seen in table 2.6 and is only listed as a reference for the comparisons.

#### A. SigFox

#### B. IETF

1) **6LoWPAN:** is a relatively new protocol that is maintained by the Internet Engineering Task Force (IETF) [7, 6]. The purpose of the protocol is to enable IPv6 traffic over a IEEE 802.15.4 network with as low overhead as possible; this is achieved by compressing the IPv6 and UDP header. A full size IPv6 + UDP header is 40+8 bytes which is tild 38% of a IEEE 802.15.4 frame, but with the header compression this overhead can be reduced to 7 bytes, thus reducing the overhead to tild 5%, as seen in figures 2.3 and 2.4.

#### C. 3GPP

##### 1) NB-IoT:

#### 2) EC-GSM:

#### 3) e-MTC:

#### D. IEEE

##### 1) IEEE 802.11:

2) **IEEE 802.15.4:** At present days, there are several technology standards. Each one is designed for a specific need in the market. For the Wireless Sensor Networks, the aim is to transmit little information, in a small range, with a small power consumption and low cost. The IEEE 802.15.4 standard offers physical and media access control layers for low-cost, low-speed, low-power Wireless Personal Networks (WPANs)

a) **Physical Layer:** The standard operates in 3 different frequency bands: - 16 channels in the 2.4GHz ISM band - 10 channels in the 915MHz ISM band - 1 channel in the European 868MHz band

b) **Definitions:** Coordinator: A device that provides synchronization services through the transmission of beacons. PAN Coordinator: The central coordinator of the PAN. This device identifies its own network as well as its configurations. There is only one PAN Coordinator for each network. Full Function Device (FFD): A device that implements the complete protocol set, PAN coordinator capable , talks to any other device. This type of device is suitable for any topology. Reduced Function Device (RFD): A device with a reduced implementation of the protocol, cannot become a PAN Coordinator. This device is limited to leafs in some topologies.

c) **Topologies:** Star topology: All nodes communicate via the central PAN coordinator , the leafs may be any combination of FFD and RFD devices. The PAN coordinator usually uses main power.

Peer to peer topology: Nodes can communicate via the central PAN coordinator and via additional point-to-point links . All devices are FFD to be able to communicate with each other.

Combined Topology: Star topology combined with peer-to-peer topology. Leafs connect to a network via coordinators (FFDs) . One of the coordinators serves as the PAN coordinator .

**IEEE 802.15.4:** The IEEE 802.15.4 standard defines the PHY and MAC layers for wireless communication [6]. It is designed to use as little transmission time as possible but still have a decent payload, while consuming as little power as possible. Each frame starts with a preamble and a start frame delimiter; it then continues with the MAC frame length and the MAC frame itself as seen in figure 2.2. The overhead for each PHY packet is only 4+1+1 133 tild 4.5%; when using the maximum transmission speed of 250kbit/s, each frame can be sent 133byte in 250kbit/s = 4.265ms. Furthermore, it can also operate in the 868MHz and 915MHz bands, maintaining the 250kbit/s transmission rate by using Offset quadrature phase-shift keying (O-QPSK).

Several network layer protocols are implemented on top of IEEE 802.15.4. The two that will be examined are 6LoWPAN and ZigBEE.

3) *ZigBee*: is a communication standard initially developed for home automation networks; it has several different protocols designed for specific areas such as lighting, remote control, or health care [27, 6]. Each of these protocols uses their own addressing with different overhead; however, there is also the possibility of direct IPv6 addressing. Then, the overhead is the same as for uncompressed 6LoWPAN, as seen in figure 2.5.

A new standard called ZigBee 3.0 aims to bring all these standards together under one roof to simplify the integration into IoT. The release date of this standard is set to Q4 2015.

#### E. LoRaWAN

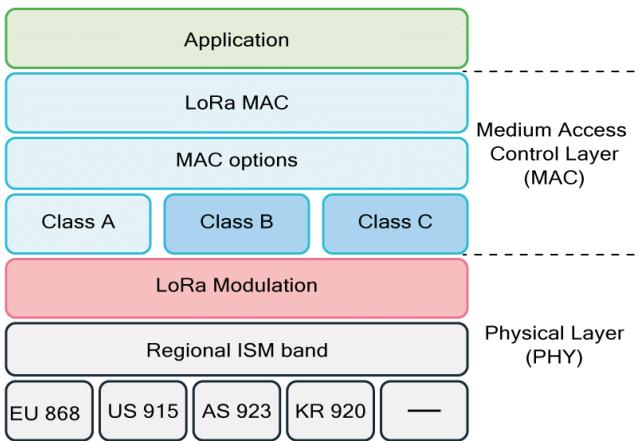


Figure 10. uhuhuh.

LoRa (Long Range) is a proprietary spread spectrum modulation technique by Semtech. It is a derivative of Chirp Spread Spectrum (CSS). The LoRa physical layer may be used with any MAC layer; however, LoRaWAN is the currently proposed MAC which operates a network in a simple star topology.

As LoRa is capable to transmit over very long distances it was decided that LoRaWAN only needs to support a star topology. Nodes transmit directly to a gateway which is powered and connected to a backbone infrastructure. Gateways are powerful devices with powerful radios capable to receive and decode multiple concurrent transmissions (up to 50). Three classes of node devices are defined: (1) Class A enddevices: The node transmits to the gateway when needed. After transmission the node opens a receive window to obtain queued messages from the gateway. (2) Class B enddevices with scheduled receive slots: The node behaves like a Class A node with additional receive windows at scheduled times. Gateway beacons are used for time synchronisation of end-devices. (3) Class C end-devices with maximal receive slots: these nodes are continuous listening which makes them unsuitable for battery powered operations.

##### 1) ALIANCE:

###### a) Class-A:

**Uplink:** LoRa Server supports Class-A devices. In Class-A a device is always in sleep mode, unless it has something to transmit. Only after an uplink transmission by the device, LoRa Server is able to schedule a downlink transmission. Received frames are de-duplicated (in case it has been received by multiple gateways), after which the mac-layer is handled by LoRa Server and the encrypted application-playload is forwarded to the application server.

**Downlink:** LoRa Server persists a downlink device-queue for to which the application-server can enqueue downlink payloads. Once a receive window occurs, LoRa Server will transmit the first downlink payload to the device.

**Confirmed data:** LoRa Server sends an acknowledgement to the application-server as soon one is received from the device. When the next uplink transmission does not contain an acknowledgement, a nACK is sent to the application-server.

**Note:** After a device (re)activation the device-queue is flushed.

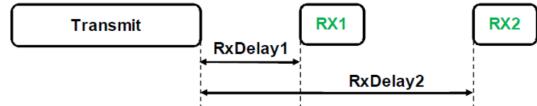


Figure 11. Class A.

b) *Class-B*: LoRa Server supports Class-B devices. A Class-B device synchronizes its internal clock using Class-B beacons emitted by the gateway, this process is also called a beacon lock. Once in the state of a beacon lock, the device negotiates its ping-interval. LoRa Server is then able to schedule downlink transmissions on each occurring ping-interval.

**Downlink:** LoRa Server persists all downlink payloads in its device-queue. When the device has acquired a beacon lock, it will schedule the payload for the next free ping-slot in the queue. When adding payloads to the queue when a beacon lock has not yet been acquired, LoRa Server will update all device-queue to be scheduled on the next free ping-slot once the device has acquired the beacon lock.

**Confirmed data:** LoRa Server sends an acknowledgement to the application-server as soon one is received from the device. Until the frame has timed out, LoRa Server will wait with the transmission of the next downlink Class-B payload.

**Note:** The timeout of a confirmed Class-B downlink can be configured through the device-profile. This should be set to a value less than the interval between two ping-slots.

##### Requirements:

Device The device must be able to operate in Class-B mode. This feature has been tested against the develop branch of the Semtech LoRaMac-node source.

**Gateway** The gateway must have a GNSS based time-source and must use at least the Semtech packet-forwarder version 4.0.1 or higher. It also requires LoRa Gateway Bridge 2.2.0 or higher.

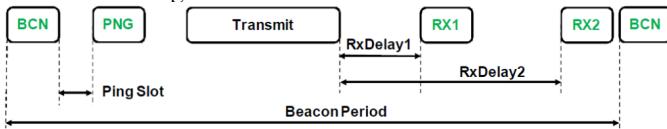


Figure 12. Class B.

c) *Class-C:*

**Downlink:** LoRa Server supports Class-C devices and uses the same Class-A downlink device-queue for Class-C downlink transmissions. The application-server can enqueue one or multiple downlink payloads and LoRa Server will transmit these (semi) immediately to the device, making sure no overlap exists in case of multiple Class-C transmissions.

**Confirmed data:** LoRa Server sends an acknowledgement to the application-server as soon one is received from the device. Until the frame has timed out, LoRa Server will wait with the transmission of the next downlink Class-C payload.

**Note:** The timeout of a confirmed Class-C downlink can be configured through the device-profile.

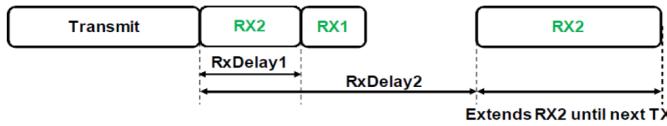


Figure 13. Class C.

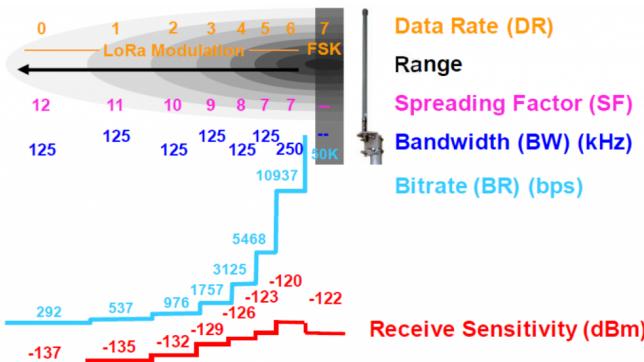


Figure 14. LoraWan Parameters.

2) *SEMTECH:* LoRa has four configurable parameters:

**BW Bandwidth:** Bandwidth (BW) is the range of frequencies in the transmission band. Higher BW gives a higher data rate (thus shorter time on air), but a lower sensitivity (due to integration of additional noise). A lower BW gives a higher sensitivity, but a lower data rate. Lower BW also requires more accurate crystals (less ppm). Data is send out at a chip rate equal to the bandwidth. So, a bandwidth

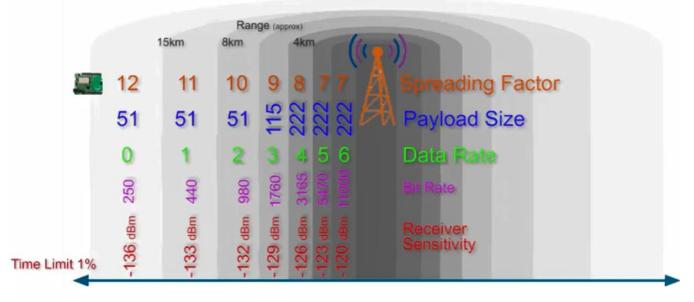


Figure 15. .

of 125 kHz corresponds to a chip rate of 125 kcps. The SX1272 has three programmable bandwidth settings: 500 kHz, 250 kHz and 125 kHz. The Semtech SX1272 can be programmed in the range of 7.8 kHz to 500 kHz, though bandwidths lower than 62.5 kHz requires a temperature compensated crystal oscillator (TCXO).

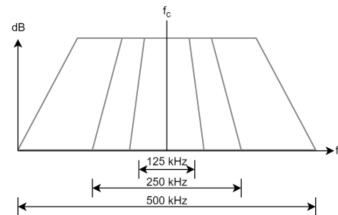


Figure 16. .

**CF Carrier Frequency:** Carrier Frequency (CF) is the centre frequency used for the transmission band. For the SX1272 it is in the range of 860 MHz to 1020 MHz, programmable in steps of 61 Hz. The alternative radio chip Semtech SX1276 allows adjustment from 137 MHz to 1020 MHz.

**CR Coding Rate:** Coding Rate (CR) is the FEC rate used by the LoRa modem and offers protection against bursts of interference. A higher CR offers more protection, but increases time on air. Radios with different CR (and same CF/SF/BW), can still communicate with each other. CR of the payload is stored in the header of the packet, which is always encoded at 4/8.

**SF Spreading Factor:** SF is the ratio between the symbol rate and chip rate. A higher spreading factor increases the Signal to Noise Ratio (SNR), and thus sensitivity and range, but also increases the air time of the packet. The number of chips per symbol is calculated as  $2^{sf}$ . For example, with an SF of 12 (SF12) 4096 chips/symbol are used. Each increase in SF halves the transmission rate and, hence, doubles transmission duration and ultimately energy consumption. Spreading factor can be selected from 6 to 12. SF6, with the highest rate transmission, is a special case and requires special operations. For example, implicit headers are required. Radio communications with different SF are orthogonal to each other and network

separation using different SF is possible.

#### Tx Transmition power:

Payload Payload leght:

$$\text{LoRa} = \frac{2^{SF}}{BW} \left( (NP + 4.25) + \left( SW + \max \left( \left[ \frac{8PL - 4SF + 28 + 16CRC + 20IH}{4(SF - DE)} \right] * \frac{4}{CR} \right) \right) \right)$$

$$\text{Lora} = n_s = 8 + \max \left( \left[ \frac{8PL - 4SF + 8 + CRC + H}{4 * (SF - DE)} \right] * \frac{4}{CR} \right)$$

$$\text{Lora} = \frac{1}{R_s} \left( n_{preamble} + \left( SW + \max \left( \left[ \frac{8PL - 4SF + 28 + 16CRC + 20IH}{4(SF - DE)} \right] * \frac{4}{CR} \right) \right) \right)$$

$$\text{GFSK} = \frac{8}{R_{GFSK}} \left( L_{preamble} + SW + PL + 2CRC \right)$$

$$\text{GFSK} = \frac{8}{DR} (NP + SW + PL + 2CRC)$$

2) *BACnet*:

3) *Z-WAze*:

4) *Bluetooth LE*: BLE is developed to be backwards compatible with Bluetooth, but with lower data rate and power consumption [28]. Featuring a data rate of 1Mbit/s with a peak current consumption (CR less than 4), 15mA, it is a very efficient protocol for small amounts of data. Each frame can

(2) be transmitted 47bytes in 1Mbit/s = 376Mus; thanks to the

short transmission time, the transceivers consumes less power as the transceiver can be in receive mode or completely off

(3) most of the time. BLE uses its own addressing methods and as the MAC frame size (figure 2.6) is only 39bytes, thus IPv6

addressing is not possible (CR + 4), 0

(4) Starting from Bluetooth version 4.2, there is support for

IPv6 addressing with the Internet Protocol Support Profile; the new version allows the BLE frame to be variable between 2 257 bytes. The network set-up is controlled by the standard

Bluetooth methods, whereas IPv6 addressing is handled by 6LoWPAN as specified in IPv6 over Bluetooth Low Energy [29].

#### G. Summary and discussion

#### V. IOT SDN

SDN is the equipement virtualisation, NFV is the function virtualisation

A. Application

B. Controle

C. Data

D. Summary and discussion

#### VI. IOT SECURITY

A. Application

1) *MQTT*:

2) *Blockchain*: Blockchain Layers

- ➡ Transaction & contract layer
- ➡ Validation layer (forward validation request)
- ➡ Block Generation Layer (PoW,PoC, PoA PoS, PBFT)
- ➡ Distribution Layer
- Consensus algorithms
- ➡ Proof of Work (PoW)
- ➡ Proof of Capacity (PoC)
- ➡ Proof of Authority (PoA)
- ➡ Proof of Stake (PoS)
- ➡ Proof of Byzantine Fault Tolerant (PBFT)

B. Network

C. MAC/Radio

1) *LoraWan*:

2) *Zigbee*:

3) *802.15.4*:

SF	07	08	09	10	11	12	07	08	09	10	11	12	07	08	09	10	11	12
BW	125						250						500					
125	x							x							x			x
	x							x										
		x							x									
			x							x								
				x							x							
					x							x						
250					x				x						x			x
					x				x							x		
				x					x								x	
					x					x								x
						x				x								
							x				x							
500									x				x		x			
						x			x					x				
					x				x						x			
						x			x							x		
							x			x							x	
								x			x							x

Table XIV. uyuyuy

Module	SX1261/62/68	SX1272/73	SX1276/77/78/79
Modem	LoRa & FSK	LoRa	LoRa
Link budget	170dB	157dB	168dB
Power amplifier	/61: +15dBm 62/68:+22dBm	+14 dBm	+14dBm
Rx current	4.6 mA	10 mA	10 mA
Bit rate	62.5kbps-LoRa 300kbps-FSK	300 kbps	300 kbps
Sensitivity	-148 dBm	-137 dBm	-148 dBm
Blocking immunity	88 dB	89 dB	Excellent
Frequency	150-960 MHz	860-1000 MHz	137-1020 MHz
RSSI		127 dB	127 dB

Table XV. [15]

#### F. Divers

1) *IPLC*:

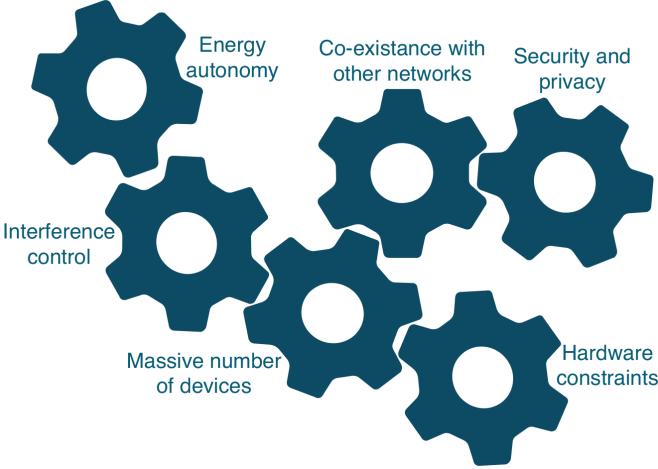


Figure 17. LPWAN Challenges [16].

#### D. Summary and discussion

#### VII. LPWAN CHALLENGES

- A. Energy autonomy
- B. Interference control
- C. Massive number of devices (Scalability)
- D. Co-existance with other networks
- E. Security and privacy
- F. Hardware constraints

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# x-Testbed

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## Abstract—

## I. INTRODUCTION [1]

### A. Problem Statement

### B. Background

Internet of Things (IoT) is a concept aiming at connecting all things to the Internet [1]. The different kinds of devices range from simple sensor devices to complex machines such as industry robots. Home automation has been available for a few years in the forms of timers and remotely controlled devices, such as lights, garage door, and climate control equipment. Also in the industry and workplace, there are current systems that have some of the functionality of IoT, e.g., sensors in robots and machines which keep track of the system status so that maintenance can be scheduled at the right time. However, these systems or sensors rarely communicate with each other or make decisions based on other sensor values; instead they depend on input from a user. In the same way cellphones connected people and made them constantly connected to the Internet, IoT will connect devices and make them constantly connected to the Internet [2]. In theory, this could lead to a future with autonomous technology all around us. The benefits could be huge as it would save time and energy for both the individual at home and for the industry [3]. IoT could be used in industry to automate power-heavy tasks to run when the electricity price is low. This principle can also be applied for the home user with laundry machines and charging of e.g. electric cars. This practice would lead to reduced energy consumption and thus a reduced environmental footprint. i3tex AB wants to investigate potential fields of applicability of this upcoming technology. i3tex AB has customers in the automotive, communication, and pulp industries; those customers have made inquiries on how to integrate IoT and sensor networks into production. As technology evolves, size and energy consumption of the IoT devices will decrease and computation power will increase [4]. This reduction in size and energy consumption, together with the increased computing power, will open up new fields for IoT. Thus, i3tex AB want to have an IoT platform to present to their current and potential customers. The interest in IoT is rapidly increasing, and thus, in the near future, the number of devices connected to the Internet is expected to increase rapidly. To support this huge

increase in both number of connected devices and the sheer amount of data that will be sent over both wired and wireless networks, the communication technology must be ready [5].

### C. Purpose (Goal)

The purpose of this project is to find and examine a communication method for devices that are made to be a part of IoT. This will be done by examining the available technologies and then developing a prototype based on the findings, which will be used for examining the communication method. This project will examine the physical, link, and network layers [6, 7] of the Open Systems Interconnection model (OSI model) [8], in order to find suitable technologies on the market. As IoT is still only defined as a concept, there are several technologies to take into consideration and examine in further detail. The prototype will be delivered to i3tex AB together with appropriate documentation, e.g. technical specification, hardware manual, software manual, and API specification.

### D. Limitations

To be able to achieve the project goal within the available time, limitations need to be defined in the three main areas of: Operating System (OS), hardware, and communication method. The OS will not be custom-made, but rather selected amongst those already on the market. Thus, to simplify the hardware selection, only those OSs which already have hardware support that meets the requirements will be taken into consideration. Furthermore, support for either 6LoWPAN, ZigBee, or Bluetooth Low Energy (BLE) as communication method is required, since development to make those standards available is outside the scope of the project. On the hardware side, the limitations will be to only use existing devices and parts as there will be no time for developing hardware or Printed Circuit Boards (PCBs). However, the hardware does not need to have an integrated radio transceiver, but needs to support at least one transceiver supporting IEEE802.15.4 [6]. Thus, communication methods will be primarily selected from specifications building on the IEEE 802.15.4.

### E. Method

To ensure that the right technologies were selected and investigated, the first phase of the project was a literature

study. The study served as a foundation when developing and performing the evaluation of the communication methods. At the end of the phase, a requirements specification was formulated to serve as a platform for the next phase. After the literature study, a selection process was performed, where the most promising technologies that met the requirements were examined in further detail and brought into the development phase. This process included the selection of development tools and other decisions bound to the product development. In the development phase, the chosen set-up was configured and assembled to prepare for testing; it was then tested according to throughput, range, latency, and energy consumption. Throughput was measured in kilobyte per second (KB/s) and tested by transferring data of different sizes in both congested and uncongested network set-ups to simulate real world and lab environments. The same set-up was used to measure the latency of a transmission, which was measured in microseconds (ms). Range was calculated instead of measured, with meters (m) as the unit. The power consumption was measured in watts (W). Each week a meeting with the company supervisors was performed, to keep the work on the right track. Here, feedback was given and other issues and questions handled.

## II. RELATED WORK

## III. BACKGROUND

The goal of the implementation phase is to have a working prototype for future assessment. To make the process of implementing the prototype possible, the first part of the implementation process will be to create a set of requirements. When these are set, the process will continue by comparing the data from chapter 2 to find the candidates that fulfil the requirements. After the technologies are selected, the process will continue with setting up the workspace, which includes the platform for development and the required tools to build, debug and test the prototype. Finally, when these three steps have been performed, the next step will be to start with the actual prototype development.

### *Requirements*

As the time dedicated for development is limited, the requirements have to make sure that the development process does not run into any major obstacles. All parts of the total prototype need to fit together seamlessly. However, the hardware platform and the operating system are tied most closely together. Therefore, they need requirements that complement each other, so that they can act as a platform for software development. Naturally, both the hardware and the operating system requirements might have to be altered slightly to enable the best match.

### *A. Hardware*

Each platform examined in chapter 2 has different strengths and weaknesses. When looking at the MCU, OpenMote has a ARM Cortex-M3 which is more powerful compared to the other two alternatives: it features a 32bit 32MHz core with 32KB RAM and 512KB flash memory compared to

the MSP430 16bit 25MHz core with 10KB/100KB memory configuration. In terms of peripherals all three platforms are comparable, with similar amount of DAC, GPIO pins, and external busses. All of the platforms have a temperature sensor and an accelerometer, but OpenMote also features an light/uv-light sensor and a voltage sensor built into the MCUs ADC. The MSB430 platform has a somewhat lower power consumption in active RF mode thanks to the less power-hungry sub-GHz transceiver. On the other hand, the less powerful MSP430 MCU has a better deep sleep power consumption, but as the radio is not integrated in the chip as it is in the CC2538 SoC, that advantage is offset by the external transceiver. Comparing the transceivers, there are two 2.4GHz models and one subGHz model; the sub-GHz CC1100 has a higher transmit power of 10dBm compared to 0dBm for CC2420 and 7dBm for CC2538. Also the sensitivity is similar for all the alternatives but gives a slight advantage to CC2538 with -97dBm compared to -95dBm and -93dBm for CC2420 and CC1100. Using these numbers and Friis range equation (equation 4.1) the range of each transceiver with a Fade margin (FM) of 20dB can be seen in figure 3.1. The benefits of working with lower frequencies can clearly be seen as the theoretical range of CC1100 is almost 3 times longer than the 2.4GHz transceivers.

Adding all this information together, the choice of platform will land on OpenMote with the CC2538SoC. It both has a MCU with more memory and better performance, and a transceiver with really good characteristics both in terms of energy consumption and range. Also OpenMote is the only option that can act as a border router using OpenBase; it lets the SoC interface with USB, UART, JTAG, and Ethernet, which enables the standalone border router mode without the need to be connected to a computer or other hardware. The OpenBattery extension lets the SoC operate as a node in a mesh network and provides a dual AAA battery slot connected to the PCB.

*Resume:* Transceiver with IEEE 802.15.4 or IEEE 802.15.1 Integrated sensor/sensors MCU with low power mode under 5MicroA Wakeup from low power mode with timer Border router ability Joint Test Action Group (JTAG) support

### *B. Operating system*

As a modern operating system can be compiled to match almost any hardware, the most important thing to have in mind is the out-of-the-box hardware support. Only RIOT and Contiki have full support for the ARM Cortex-M3 of the considered operating systems and thus both TinyOS and freeRTOS are directly eliminated as developing the support would take too much time. Compared to RIOT, Contiki also has full driver support for the sensors and transceiver, which should decrease the implementation time significantly. When compiled into binary form, RIOT uses less RAM and ROM and thus probably is a bit faster compared to Contiki, which could be important if the application consumes much resources. The lower memory usage might also give RIOT an advantage in being future-proof. Contiki also has support for soft real-time scheduling

compared to the hard real-time scheduling of RIOT; this is however not crucial, as the software that will be running on the OS does not have any hard real-time constraints. Both RIOT and Contiki have support for 6LoWPAN but no support for either ZigBee or BLE; this is due to the fact that these are proprietary stacks. Support could be added but would take some time to customize for the given OS. What gives Contiki the largest advantage is that it also have border router software ready for deployment, which in the RIOT case would have to be developed. All in all, as the project has such a limited time frame, Contiki will be selected as the OS; this, mainly because Contiki comes with most advantages time-wise; this choice means that the focus of the software development will be the creation of a test and evaluation system.

*Resume:* Support for the SoC/MCU and transceiver 6LoWPAN, ZigBee or BLE stack Soft real-time RAM and ROM footprint matching the hardware

### C. Communication protocol

The OpenMote platform has a IEEE 802.15.4 transceiver and thus supports both ZigBee and 6LoWPAN; this means that BLE is not an option. As ZigBee does not have full IPv6 support yet and is not integrated into Contiki, the natural choice is 6LoWPAN. This choice will not only save some development time but also enables evaluation of the header compression. As seen in figure 3.2, the 6LoWPAN stack in Contiki will replace the IP stack while maintaining the same functionality. As the functionality is the same, TCP and HTTP will work with 6LoWPAN, but including them in the source increases the OS build size considerably. On top of the UDP layer, Contiki also has a working implementation of CoAP that can be used for retrieving data from the nodes in a power efficient manner. CoAP is a stateless protocol that uses the HTTP response headers to achieve a very low overhead in transmissions while using application level reliability methods to ensure packet delivery.

*Resume:* IPv6 addressing support Existing OS support Network type is mesh UDP

### D. Workspace and tools

The ARM Cortex-M3 chip that OpenMote and CC2538 is built upon requires the GCC ARM Embedded compiler. This tool-chain is free and runs on both Linux, OSX, and Windows; however, there is no bundled development application so a secondary application for programming is needed. In Windows, there are several Integrated Development Environments (IDEs) such as IAR Workbench ARM [30], Code Composer Studio [31] and the Eclipse plug-in ARM DS-5 [32, 33]; these IDEs use various proprietary toolchains and have a price tag ranging from free to several thousand SEK. Most of the IDEs also have a code size restriction for the free versions. To minimize the costs, the development machine development machine used in this project will run Ubuntu 14.04 LTS, the used tool-chain is GCC ARM Embedded, and Geany is used as the code development application. To analyse the network traffic in real-time, the open source tool Wireshark is used together

with a IEEE 802.15.4 packet sniffer. Together with a laptop, the packet sniffer will grant the ability to traverse the mesh network and analyse the network in real-time as it is seen by the nodes.

## IV. PROPOSED ...

The goal of the development process is to have a functional border router and at least two nodes to be able to test how response time and throughput differs with each hop in a mesh network. To be able to measure response time and throughput, each node needs to have a CoAP server which can respond to ping and also receive an arbitrary amount of data for throughput measurement. It is desirable for each node to be able to send information about each sensor so the project can be used as a tech-demo. The first part of the development was to set-up of the workspace and tools mentioned in section 3.1.5. Ubuntu OS was installed in a VirtualBox Virtual Machine to make it easier to duplicate and backup; this procedure gave a noticeable decrease in performance and it is recommended to have a dedicated native Ubuntu machine for this type of development. Even though Ubuntu uses an easy-to-use package system, there were some problems in finding a version of GCC ARM Embedded tool-chain that was compatible with Contikis built-in simulator Cooja [34, 35]; eventually, version 4.82 was used to successfully build Contiki. Cooja is a useful tool for testing and debugging network configurations but does not have support for the CC2538 MCU; instead, nodes called Cooja Motes are simulated with generic hardware. As Cooja is written in Java and runs in a JVM, Oracle Java 1.8 was also installed.

### A. Drivers and firmware

Figure 3.3 shows an overview of the full system. The foundation is the SoC with the MCU, transceiver, and sensors. The Contiki operating system implements the soft real-time kernel together with the firmware for the SoC/MCU and the drivers for the peripherals and sensors. The last part is the communication stack, which provides TCP and UDP connectivity over 6LoWPAN. On top of the TCP and UDP protocols, HTTP and/or CoAP can be implemented. The firmware required for the OS to work properly on the hardware platform was already implemented. However, the drivers for the I<sup>2</sup>C bus and the sensors were not implemented. The I<sup>2</sup>C driver is required for the sensor drivers which in turn enables the MCU to communicate with the sensors on the OpenBattery platform.

### B. CoAP server

In order to make each nodes sensor data accessible, a CoAP sever was implemented as an application running on top of Contiki. A CoAP server in general can handle any number of resources; in this implementation, one resource was made for each sensor value i.e. temperature, light, humidity, and core voltage. The temperature, light, and humidity sensors all work in a similar fashion. When their value is requested, the I<sup>2</sup>C bus is initialized and then a request is sent over the bus. When the response with the value arrives, that data is put

into either a plain-text or JavaScript Object Notation (JSON) formatted message depending on the request and then sent back to the requester. As the core voltage sensor is part of the MCUs ADC, that value is retrieved by simply getting data from a register (a somewhat faster operation). As a buffer for testing throughput speed was also needed, a resource with a circular buffer was implemented. This resource is configured with CoAPs block-wise transfer functionality for arbitrary data size; however, the buffer in itself is only 1024Kb to allow the program variables to fit into the ultra low leakage SRAM. For testing purposes, the data could have been discarded instead of actually saved into the buffer, but then the transfer can not be verified. Resources are defined by paths as CoAP works in a very similar way as HTTP.

Each resource is registered in the server with its path, media type, and content type. When a package arrives on the CoAP port, the server starts to break down the package to be able to direct it to the right resource. It starts with verifying that the package is actually a CoAP package, and then it checks the path and sends it to the correct resource. The resource then inspects the method field in the package header to direct the incoming data to the right function. CoAP package method can be either GET, PUT, POST or DELETE. This function then inspects the request media type and answer content type so that the function can parse the request and send a correctly formatted answer. If the resource does not implement the received method, the server responds with 405 Method not allowed and if the content/media type is not supported the answer is 415 Unsupported Media Type. The content/media types are text/plain, application/json, application/exi, and application/xml.

*1) Testing:* Contiki is shipped with a simulation tool called Cooja which is written in Java; it can simulate an arbitrary number of nodes with different roles and configurations. All simulation data, such as radio packages and node serial output may be viewed through different windows and exported to various formats. Unfortunately Cooja did not have support for ARM Cortex-M3, but the general set-up was still tested by using Cooja Motes, which are nodes without specified hardware, and MSP430 nodes such as Wismote or Skymotes. With this simulator the basic understanding of the communication between nodes was gained; also, before the hardware arrived, early testing was performed to test the OS and application software.

*2) Final prototype:* The final prototype consists of four OpenBattery nodes and one OpenBase border router. Both the nodes and the border router are deployed with Contiki. Each node runs a CoAP server, described in section 3.2.2, on top of the OS in its own thread. The border router runs a router software called 6lbr that acts as translator between Ethernet and IEEE 802.15.4 [36]. Both types of hardware are configured with a 8Hz Radio Duty Cycle (RDC) driver to keep the power consumption to a minimum. RDC is a OS driver that cycles the listening mode of the transceiver to reduce power consumption. As Contiki puts the MCU into Low Power Mode (LPM) when no function is running and

the transceiver is off, the RDC driver indirectly controls when the MCU is in LPM. When using the RDC protocol, the nodes repeatedly send messages until the target node wakes up and sends an Acknowledge packet (ACK); this makes communication seamless, even though most of the time the nodes transceivers are not active. Also, an always-on RDC driver, where the transceiver is constantly listening, will be used to be able to look at the performance impact of the 8Hz RDC.

## V. EXPERIMENTATION

In this chapter the results from each type of assessment are presented. The first assessment is range, followed by response time, after that connection speed, and finally the power consumption. The only assessment that is not performed on the prototype is the range assessment.

### A. Range

Range is very hard to measure without advanced equipment and isolated rooms but can be roughly estimated with equation 4.1 called Friis range equation [37].  $P_t$  is the sender transmit power,  $P_r$  the receiver sensitivity,  $d$  is the distance between the antennas in meters,  $f$  is the signal frequency in hertz, and  $\lambda$  is the wavelength.  $G_t$  and  $G_r$  is the antenna gain for the transmitter and the receiver. The last term in equation 4.1, when inverted, is the Free-space path loss (FSPL) and can be expanded as shown in equation 4.3.  $P_r(\text{dB}) = P_t + G_t + G_r + 20 \log_{10}(\lambda) - 4d \cdot c \cdot f / \text{FSPL}(\text{dB}) = 20 \log_{10}(d) + 20 \log_{10}(\lambda) + 147.56$  = (4.1) (4.2) (4.3) Unlike Friis range equation, the Link budget equation 4.4 also takes external loss like FM into account [38]. This is needed to make a correct estimation of the actual range as there are several things in the environment that obstructs and distorts the signal.  $P_r = P_t + G_t + G_r - \text{FM} - \text{FSPL}$  (4.4) Combining equation 4.1, 4.3 and 4.4 gives us the equation for the estimated distance as seen in equation 4.5.  $d = 10 \times P_t + G_t + G_r - P_r - \text{FM} - 147.56 - 20 \log_{10}(\lambda)$

With this equation an estimation of the transceiver range can be made for different FMs and transmit powers. When deployed, the transceiver is configured to only accept packages with a signal strength of -70dBm and above to minimize packet loss and corruption. The antenna gain for OpenMote is 0dBi and can thus be omitted. Figure 4.1 shows a comparison between three different levels of FM: 0dB, 10dB, and 20dB. A FM of 0dB means that there is no signal loss except the FSPL and this is very hard to achieve outside of a lab environment. When increasing the FM to 10dB, which corresponds to a normal home environment, the maximum range drops to 22m. However, in these kind of environments the desired range is usually around 10m which would let the device reduce the transmit power to around 0dBm. Finally, the FM is increased to 20dB which is roughly what it would be in a office or industrial environment. The maximum range in this environment is now reduced to only 7m when transmitting at maximum power.

**Response time** Before measuring the response time, some theoretical estimations are needed to be able to evaluate the real values. The theoretical values are based upon the radio duty cycle (RDC) and the average response time to reach a node can thus be derived from equation 4.6, 4.8, 4.9 and 4.10. As each node only

### B. Response time

Before measuring the response time, some theoretical estimations are needed to be able to evaluate the real values. The theoretical values are based upon the radio duty cycle (RDC) and the average response time to reach a node can thus be derived from equation 4.6, 4.8, 4.9 and 4.10. As each node only checks the radio every 125ms, this duration combined with the data packet send time of 4ms (equation 4.7.) and ACK send time corresponds to the worst case delivery, as the node needs to wait a whole cycle before being able to send the package to the desired node. When the target node is already listening, the best case delivery time is 5ms. Thus, the average theoretical delivery time to reach any adjacent node is 67.5ms. Radio duty cycle: Transfer time:  $1s = 0.125s = 125ms$  8 (4.6)  $133B + 4B = 4ms$   $31.25KB/s$  (4.7) Worst case delivery:  $125ms + 4ms + 1ms$  (4.8) Best case delivery:  $4ms + 1ms$  (4.9)  $130ms + 5ms = 67.5ms$  (4.10) 2 The delivery time is only calculating the time to send a packet over a link, but when calculating the response time, the acknowledge (ACK) response has to be included in the calculation. Each ACK also needs to wait for the target node to be awake, adding one more instance of average delivery time, resulting in 125ms in average response time. This time will multiply with each hop, resulting in equation 4.11, 4.12 and 4.13. Avg. delivery: Avg. response time:  $(2 \times 67.5ms) \times \text{hops}$  (4.11) Best case response time:  $(2 \times 5ms) \times \text{hops}$  (4.12) Worst case response time:  $(2 \times 130ms) \times \text{hops}$  (4.13) After doing a test with real nodes set-up with a 8Hz RDC with three hops, as seen in figure 4.2, the values in table 4.1 were obtained. Each node was pinged 200 times at a one minute interval to simulate some traffic on the network. What can clearly be seen in the average field of the table is that the average of 765ms is much higher than the expected average of 135ms; the difference is mainly due to the worst-case pings that in some cases had response times up to 30 seconds. However, when looking at the geometrical mean which is better at smoothing out big spikes seen in figure 4.3, the observed response time is still 265ms which is a bit longer than the expected worst case response time for one hop. Also, for two and three hops the observed average is high, but the geometrical mean shows that this is due to the spikes. The estimated response time for two hops is 270ms which as seen in the geometrical mean table 4.1 is way off by 500ms. The same observation goes for three hops where the observed geometrical mean response time is 1181ms which compared to the estimated response time of 405ms is significantly higher.

With these observations in mind, the estimation could be described much better with equation 4.14. which would result in an average response time of 266, 532 and 1064 ms for one, two and tree hops. However, this would mean that the

response time is doubled for one hop and then doubled for each consequent hop making the response time exponential which should not be the case.

With the RDC disabled, i.e. the transceiver is always listening and the MCU does not go into sleep mode, the response time is completely different. As seen in table 4.2 the average response time is around 12ms per hop and the spikes seen in the response time for 8Hz RDC is gone. Furthermore, the estimated best case response time of 10ms is very close to the observed average response time. The response time also scales to the number of hops as expected and is roughly 12ms per hop. It appears there might be a problem in

### C. Connection speed

Connection speeds can be measured in several ways each with their own different pros and cons. One of the most popular ways is throughput, i.e. the amount of data over the link is divided by the time it took to reach the target. However, this gives a false picture of how fast the connection actually is from the developers point of view, as the measured data does not only contain application data but also headers and checksums. IEEE 802.15.4 has a theoretical data rate of 250kb/s as seen in equation 4.15. but this is only a measure of how many bits per second the transceiver is able to output. The application data part, when using no header compression, is only 41% of the total transfer. Thus, resulting in a theoretical application data rate, also called goodput, of only 12.81KB/s. Data rate:  $250kb/s = 31.25KB/s$   $133B - 54B = 0.59 \times 133B$  Theoretical goodput:  $31.25KB/s \times (1 - 0.59) = 12.81KB/s$  Overhead: (4.15) (4.16) (4.17) When using CoAP as the application level protocol, each package can carry either 32 or 64 bytes of application data. In practice, the 64B mode is only applicable when sending packages between nodes on the same mesh network, as the addressing fields then can be fully compressed. When using applications outside the mesh network, each package can only carry 32B of data, resulting in a packet size of 111B as shown in equation 4.18; this does not affect the theoretical data rate but has a noticeable impact on the goodput due to the large overhead of 71as shown in equation 4.19. To be able to use the full data rate, the application needs to use a protocol without handshakes, i.e. UDP, as the transceiver then can send the packets as fast as physically possible. CoAP is implemented on top of UDP and thus has a low transport layer overhead, but uses its own mechanism for handshaking, delivery and ordering. The theoretical CoAP application throughput can be estimated by looking at the average response time of the node and then add that time to the the data delivery time. Each package needs to be acknowledged before the next package is sent, and thus the node response time needs to be taken into consideration. When doing so, the throughput as calculated in equation 4.20. is only 1.64KB/s and thus the theoretical goodput is reduced from 12.81KB/s down to 0.48KB/s as shown in equation 4.21. Packet size:  $133B - 54B + 32B = 111B$  Actual overhead:  $79B = 0.71 \times 111B$  (4.18) (4.19)

Using the packet size of 111B together with the theoretical response time from equation 4.11 would give the results shown in figure 4.4. To verify these calculations, the same test set-up as shown in figure 4.2, which also was used in section 4.2, was used to test throughput and goodput at different number of hops. Each node was sent 1KB data each minute for 200 minutes; the time from the first package sent to the final acknowledge packet received was measured for each 1KB transmission. The first test was performed with an RDC of 8Hz and resulted in the values shown in figure 4.5. As the chart shows, the theoretical throughput and goodput is much higher than the observed values, but this is due to the fact that the actual average response time is higher than the theoretical one. With some calculations made the observed throughput and goodput are within range of what is expected, given the observed response times in table 4.1. Equation 4.22 uses the observed values to calculate the average response time, given the values in figure 4.5.

#### D. Power consumption

To measure power on devices that use very low power and also changes the power consumption very rapidly and frequently is not an easy task. According to the currency specification from the CC2538, the different power modes have the consumption seen in table 4.3 using the built in voltage regulator TSP6750 that switches the input voltage down from the 3V to 2.2V. The components on the OpenBattery supplied directly by the 3V batteries have the current and power consumption specifications as seen in table 4.4.

Given these power profiles, combined with the time it takes to receive and transmit packages, and retrieve a measurement, the theoretical power for one RDC cycle results in the chart seen in figure 4.7. The node starts in sleep mode using 112.81W and after 109ms wakes up and goes into RX mode where a request for a sensor value is received. The node then switches off the radio and fetches the sensor value. After the value is retrieved from the sensor, the radio is once again put in to RX mode for a Channel Clear Assessment (CCA) before entering TX mode and sending the payload. The transmission is successful and the node goes into RX mode to listen for the ACK, when it is received the node enters sleep mode again. For this cycle the average power consumption is 4.8mW which would drain the 2250mWh batteries in 19 days.

However, as the nodes have a RDC running at 8Mz, most of the time there will be no package for the node to receive and thus no measuring and transmitting, as seen in figure 4.8. This cycling reduces the average power consumption to 0.47mW, which would make the batteries last for ca 200 days. The goal is to have a node that can run for one year without having to change the batteries and to be able to do this on 2xAAA batteries with 750mAh the average consumption has to be under 257W as calculated in

To verify these assumptions, we used a Keithley 2280S power supply [39] to measure the total current draw of the prototype. The node was connected to the power supply, which was set up to make 277 measures each second with a supply

voltage of 3V. Several measurements were performed. One of the most interesting ones can be seen in figure 4.9. In this picture, we can clearly see the different operating modes, as the node performs 3 transmissions during the interval. In the first transmission at the 1.6s mark, the strobing feature of the RDC protocol is seen as the package is sent 5 times before the receiving node is awake and can receive the package. In the two following transmissions, the package is delivered on the first try. As our measurement is limited to 277Hz, the current peaks when only waking up to listen for traffic are sometimes missed, and the peak value is hard to extract; but the 8Hz RDC cycle is still visible. The average power consumption for these cycles is 8mW, which would make the batteries only last for 11 days. However, when taking the average of a measuring series without any transmissions, the average goes down to 4mW, which increases the battery time to 23 days. The theoretical sleep power of 0.11mW compared to the measured of 3mW is what makes the average power consumption that high.

Reducing this power consumption by a tenfold would result in an average consumption of 0.39mW, which is closer to the theoretical average power consumption. A discovery made when measuring the power was that the nodes consumed less power when supplied with a lower input voltage. Simply by reducing the voltage from 3V to 2.6V reduced the power consumption in LPM by 15%. However, this reduction could affect the range of the nodes.

Internet of Things can be realised in several ways as there are still many viable options on the market, mainly in terms of hardware, operating systems, and communication standards. Given the recent development in the field, Thingsquare recently released a technology demo using the same practices as used in this thesis; the choices taken are on track with the latest development [40]. Also, both Google and Microsoft have announced that they are developing IoT OSs. When these products are released, it would be very interesting to compare them with Contiki. It would be exciting to see if an open-source project can surpass the commercial offerings in terms of speed, RAM and ROM footprint, and device support. Furthermore, an in-depth comparison between RIOT and Contiki would give much insight into the kind of OS practices that benefit IoT development the most. Google have also started to develop a substitute for 6LoWPAN and UDP that they have named Thread [41]. As 6LoWPAN and ZigBee, it runs on top of IEEE 802.15.4 and thus might be able to out-compete the existing implementations. Google promises lower latencies and power consumption compared to the existing technologies.

#### The prototype

The prototype development took more time than initially planned; mostly because of the complexity of the OS, but also due to bugs in the untested drivers. The prototype combines the technology from each field, i.e. hardware, OS, and communication protocol, and fulfils the requirements set in section 3.1.1. Even though the OS is relatively simple, compared to Linux, Windows, and OS X, understanding the mechanics

of the RDC driver and the LPM driver was difficult, but necessary to be able to interpret the test results. The prototype worked very well during most of the testing, with only a few unforeseen deviations. One occurred during the power measurement, where the power consumption in low power mode tripled in one of the test series; this behaviour could not be reproduced and is therefore not included in the results. Also, in the early stages when working with the 8Hz RDC driver, packet losses over 50% were recorded for packets with more than one hop; this problem was solved, when a new version of radio driver was released by the OS development team. Selecting OpenMote to be the hardware platform together with Contiki as the OS, was a very good choice as companies are starting to build their IoT solutions around Contiki and similar hardware platforms [42, 43]. Already in the beginning of the development, several benefits were noticed; new drivers and bug-fixes were released increasing the stability and functionality of the OS. The active community around the combination of OpenMote and Contiki was really helpful when developing the drivers for the I<sup>2</sup>C and sensor drivers. Example projects for other platforms could be used as references, giving much insight to how the programming for this type of OS worked. It would have been interesting to examine the differences between two operating systems; not only to test which one has the better performance, but also to compare which one that has the more favourable code structure and development procedure.

## VI. RESULTS

Collecting the data went well and were reasonably straight forward; it was easy to transition between the two different test set-ups and thus making several test scenarios. Assessments were made in the areas of range, response time, connection speed, and power consumption. In each area, the theoretical values were first calculated and then compared to the retrieved measurements; except in the range case, as the required equipment for measuring was not economically justifiable to purchase.

### A. Range

The theoretical range for OpenMote when transmitting at full power in an office environment is only 7m. As measuring the range was not a viable option due to the cost of measuring equipment, only distance estimations from the placement of the nodes when maintaining a stable connection can be used as a reference. Using a map of the office and the position of the nodes the range seems to be around 10m, which would mean that the effective FM of the office is around 16dB using the always-on RDC. The FM changed a bit when using the 8Hz RDC as more packages congested the air and the range dropped to somewhere around 5m; resulting in an effective FM of 23dB. To increase the range of the transceiver, a switch to the 860MHz frequency band would be the most effective solution; with a FM of 23dB, the theoretical range would increase to 14m with the same transceiver properties, and with a FM of 16dB the range would be 31m. Usually,

transceivers with a lower frequency output also have a lower power consumption while transmitting. Working in sub-GHz also gives the benefit of less interference as fewer other devices uses those frequencies. Changing to a sub-GHz band would thus decrease the power consumption and increase the range, without changing the functionality of the nodes.

### B. Response time

Initially when measuring the response time the always-on RDC was used and the measured response time was very close to the theoretical value. However, when using the 8Hz RDC protocol the values started to drastically differ from the theory. This behaviour is likely to originate from the way the RDC driver predicts the next time when the target node should be awake. The procedure is called phase optimization; when enabled, the node saves the time when the node was last seen, it then uses this value to predict the next time the node should be awake based on the RDC cycle. However, this prediction is based on the nodes internal clock. As the clock can differ from those of the other nodes, misalignments seem to occur, resulting in misses when trying to reach the target node. Each misalignment increases the time it takes to reach the target node as the node then needs to strobe the package until the target nodes wakes up again. In theory, when sending strobos the target node should wake up and receive the package within one cycle (125ms); however, this is not guaranteed as other transmissions might occupy the air, further increasing the response time. If the phase optimization could be improved to guarantee the alignment between the nodes, the response time should get much closer to the theoretical value; as the time to reach the node would be maximum one cycle and the air would not be as congested by nodes sending strobos.

### C. Connection speed

The connection speed, when using CoAP or any other protocol with perpacket ACK, is directly bound to the response time. IEEE 802.15.4 has a relatively low data-rate, only 250kbps, compared to other solutions, e.g. BLE (1Mbps) and WLAN (>54Mbps). As throughput is based on datarate over a longer period of time, both the overhead and the response time is needed to make a good estimation. CoAP has a very low header size compared to many other communication protocols, but due to the very small frame size, the overhead is still relatively high. As of now, the results clearly show that when a reliable transfer is desired the connection speed of IEEE 802.15.4 and CoAP is only sufficient for data exchanges around 32 bytes. When the nodes use the always-on RDC, the goodput is less than 3KB/s for one hop and is halved for every hop; however, when the 8Hz RDC is enabled, the goodput is reduced to under 0.1KB/s. Using messages without per-packet ACK, thus removing the response time from the equation, would let the nodes transfer real-time audio and maybe even highly compressed video. However, using messages without the per-packet ACK disables the reliable transmission guarantee, and thus it can only be used with data streams where packet loss is acceptable.

#### D. Power consumption

Making a rough estimation of the power consumption of the platform was straight forward task and so was measuring the actual consumption. When comparing, the two the values differed by a factor of 30, which was not expected. The reason probably originates from the clock interrupt which is triggered every 8ms. Initially, this interrupt was assumed to be disabled when the system entered the lower power modes, but this was not the case. As the interrupt fires at 125Hz and the time to wake up and go back to LPM is only 272s, the power spikes from these interrupts were not seen on the measuring instruments. As seen in figure 4.9, even the peaks from the listening cycles were hard to record and those lasted for at least 4ms; instead, the power consumption from the clock timer looks like an increased LPM power consumption. At the time this was discovered there was no time to fix it, but doing so should decrease the average power consumption to within the limits, granting the nodes the ability to run on battery power for a year. As no delays from calculation could be observed, the clock speed on MCU could, in all probability, have been reduced to save power on the nodes. However, this reduction would only have affected the consumption when the node was in active mode, which is only a few percent of the total cycle time. The OpenMote chip has a step-down DC-DC converter for this purpose which is switched off in LPM mode to reduce quiescent currents; however, as most of the time is spent in LPM, reducing the input voltage to 2.1V by changing battery type and removing the step-down converter would be preferable as it would reduce the power consumption. These changes could affect the range of the device, but this has to be assessed.

#### E. Project execution

Looking at the time plan and the milestones, as seen in Appendix A and B, each milestone matches a task or transition in the time plan. The planning report was not submitted to the examiner until the 6/2-15, which is two weeks behind schedule, exceeding the time planned for milestone M1. The first draft was submitted before deadline, but several revisions were necessary. In retrospect, the litterature study should probably have been planned in parallel with the planning report, as the information from the study helped with the report. Milestone M2 marks the switch from the literature study and selection of technology to the development phase. This milestone was met and development could begin in the following week. As seen in Appendix C, the development phase have several risks to consider. The only risk encountered in this phase was R4, as one of the hardware platforms was delivered with a broken sensor. However, this malfunction did not affect the time plan as the development could continue regardless of the malfunction. The end of the development phase was defined by milestone M3, approval of prototype, which was completed ahead of schedule granting an early transition into the assessment phase. In the assessment phase, it could be argued that risk R9 was encountered when measuring the power consumption, as the results from those measurements

did not properly show the wake-ups from the clock timer. This phase contained milestone M4 and M5, of which of only M5 was done in time. The Half-time presentation, milestone M4, was performed on the 8/4-15 in the form of a meeting, where the progress, results and continuation plan were discussed. Also, a half-time version of the report was sent the 17/4-15 and approved by the examiner. Milestone M6, deliver the final prototype, was completed a few days before the set deadline which eliminated risk R11 and gave more time to work on the writing and the presentation. Both of the oral presentations were attended on the 1/6-15 to grant some experience in how the presentation and opposition are carried out, thus now following the time plan. However, there were not many presentations to watch during the planned weeks, as the presentation schedule follow the academic semesters. The presentation for this thesis was not performed until the 3/6-15, thus being two weeks behind schedule. However, it was scheduled on the first available date suggested by the institution. The final version of the report will be submitted to the examiner before the 19/6-15, thus successfully completing milestone M7.

## VII. DISCUSSION

The purpose of the project was to find and examine a communication protocol that could be suitable for IoT applications, by investigating the current hardware, OS, and communication protocols and building a prototype from the selected choices. What can be said about the investigation is that it is difficult to examine all candidates in detail; this means that a rough selection has to be made based on initial knowledge potentially discarding good options. The general feeling is, however, that all of the examined candidates in this project were relevant and added valuable insights to the current technology status. The assessment gave relevant and interesting results that improved the understanding in what IoT can be used for, and what further areas of investigation could be. One of the most interesting areas of further investigation would be the RDC driver, as it directly affects the response time and thus also the connection speed. Even though the power consumption was not in line with the expectations, the reason has been found and can be resolved. Another conclusion is that IoT is not ready for real-time applications as the latency is much higher than expected, for the technologies assessed in this thesis, and also has a high spread. As the latency increases for each subsequent network hop and the minimum observed latency per hop is 11ms, when using the always-on RDC, this type of communication will probably only be used for applications where response time can vary greatly, without affecting the functionality. CoAP as a communication protocol shows a lot of promise when combined with 6LoWPAN and IEEE 802.15.4. It performs well given its simplicity but has one disadvantage: the large overhead which comes from the MAC addressing fields in the IEEE 802.15.4 frame. If this overhead could be reduced from the current 71% to only 30% the goodput would double. A solution would be to use a similar mechanism as BLE where the packet size varies depending

on application. Each node also has computing time left as the MCU is more powerful than needed for the given application; an improvement would be to use a less powerful MCU, like the ARM Cortex-M0+, to reduce the clock speed as suggested in the discussion. When looking at the future-proof aspect the later suggestion is probably the better, as the clock then could be increased if more computing power is needed. In the future, batteries will hopefully be able to store more energy, thus increasing the time between battery changes or reducing the battery size.

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# Genetic Algorithm For LoRa

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**Abstract**—The exponential growth of Internet of things (*IoT*) applications in both industry and academic research raises many questions in wireless sensor networks. Heterogeneous networks of IoT devices strongly depend on the ability of IoT devices to adapt their data transmission parameters to each application requirement. One of the most important problem of the emerging IoT networks is the limitation in terms of energy consumption and computation capability. These limitations could be addressed by using the edge computing to unload IoT devices from additional computation tasks. Our work is motivated by the idea of matching each transmission configuration with a reward and cost values to satisfy applications constraints. Our goal is to make IoT devices able to select the optimal configuration and send their data to the gateway with the QoS required by IoT applications. In this work, we use LoRa network to evaluate the efficiency of our algorithm. Determining the best configuration among 6720 LaRa transmission settings is challenging. The difficulty is mainly due to the lack of tools that could take all applications requirements into account to select the best settings. To address this problem, we use a genetic algorithm in an edge computing to select the transmission parameters needed by the application. Each LoRa configuration represents a feature that needs to be selected to match better the QoS criteria. Particularly, we analyze the impact of selecting one configuration in 3 kinds of applications: text, voice and image transmission by modeling a new adaptive data rate selection process.

**Keywords**—Genetic algorithm; Fuzzy logic; LoRaWAN; Adaptive Data Rate (ADR).

## I. INTRODUCTION

The need of Low Power Wide Area Networks (*LPWAN*) increased significantly these five last years. The main factor is that IoT devices require low power consumption to transmit data in a wide area. LoRa, Sigfox and Narrowband IoT (*NB – IoT*) are the most known technologies that satisfy these requirements. Applications like smart building and smart environment are one of hundreds use cases that need to be deployed with such technologies. Unlike Sigfox and NB-IoT, LoRa is more open for academic research because the specification that governs it is relatively open. The transmission could be configured with 4 parameters: Spreading Factor (*SF*), Transmission Power (*Tx*), Coding Rate (*CR*) and Bandwidth (*BW*), to achieve better performance.

The main LPWAN research directions are about link optimization, adaptability and large scale networks to support massive number of devices. The selection of an appropriate transmission parameter for IoT networks typically depends on

the nature of the application. Thus heterogeneous transmission configuration and *SF* allocation strategies need to be studied. In this paper, we investigate the performance of heterogeneous networks (i.e., when each IoT device selects its LoRa transmission parameters according to its link budget and the application requirements). For that purpose, we have developed a LoRa transmission adaptation mechanism. Both ns-3 simulator and the Low cost LoRa Gateway [1] are used to validate our approach. The computation tasks of the selection process will run on the Gateway device (Raspberry-pi) and the required settings will be sent to nodes for the next transmission.

This paper is organized as follows. Section II elucidates summary of related works. In Section III, we propose our approach to solve LoRa parameter selection problem. Our experiments is presented in Section IV. Section V concludes this paper.

## II. RELATED WORK

Transmission parameter configuration mechanisms, such as Adaptive Data Rate (*ADR*) scheme [2] need to be developed to fit each application requirement in terms of power consumption, delay and packet delivery ratio. Solutions running on LoRa node should be less complex to match computation limitation of *IoT* devices as required in LoRaWAN specification. However, LoRa network server could run complex management mechanism, which can be developed to improve network performance. In this paper, we focus on the server-side mechanisms.

The basic *ADR* scheme [2] provided by LoRaWAN predicts channel conditions using the maximum received Signal Noise Rate (*SNR*) in the last 20 packets. The basic *ADR* scheme is sufficient when the variance of the channel is low, it reduces the interference compared with the static data rate [3][4]. However, their simplicity causes many potential drawbacks. First, the diversity of LoRa Gateway models that measures *SNR* make the measurement inaccurate as a result of hardware calibration and interfering transmissions. Second, selecting the maximum *SNR* each 20 packets received could be a very long period in many IoT applications that require less uplink transmission. Third, transmission parameters adjustment considers only the link of a single node. If many LoRa nodes are connected to the near gateway, all nodes connected to this one will use the fastest data rate. In this case, the number of

LoRa nodes using the same data rate will increase and the probability of collisions also increases dramatically.

For example, the authors in [4] slightly modify the basic *ADR* scheme by replacing the maximum *SNR* with the average function. In this paper, we focus on building a framework that help IoT devices to adapt their transmission parameters to the application requirements in a server side.

### III. PROPOSED FRAMEWORK

The selection process scheme illustrated in Figure 1 can be described following these five steps:

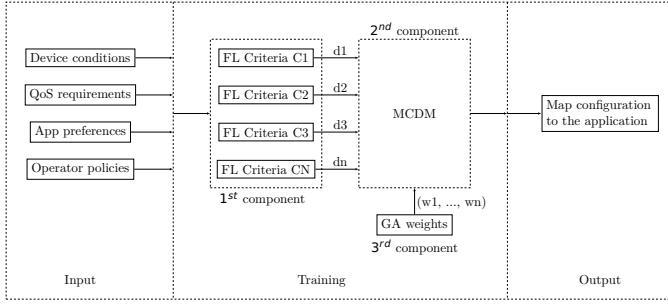
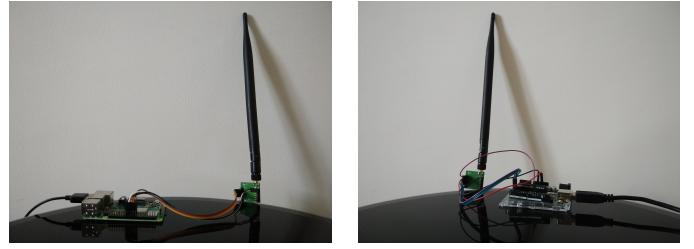


Figure 1. The proposed scheme for LoRa transmission parameters selection based on *GA*, *FL* and Multi-Criteria Decision Making *MCDM*.

- 1) According to the Semtech SX1276 LoRa transceiver [5], there are 6720 possible settings ( $s_1, \dots, s_{6720}$ ) and the framework has to select the most optimal one or to rank them according to their relevance.
- 2) The first step of the selection process depends on multiple criteria up to  $i$  ( $c_1, \dots, c_i$ ). Different type of criteria can be measured from different sources to cover the maximum point of views, as an example, the network server requirements, the applications requirements and the devices conditions.
- 3) The Fuzzy Logic (FL) based subsystem gives an initial score for each configuration that reflects its relevance. The different sets of scores ( $d_1, \dots, d_i$ ) are sent to the *MCDM* in the 5<sup>th</sup> step.
- 4) At the same time, the *GA* [6] assigns a suitable weight ( $w_1, \dots, w_n$ ) for each initial selection decision, this selection is made according to the objective function that is required by the application.
- 5) Using the initial scores coming from the 3<sup>rd</sup> step and the weights using the 4<sup>th</sup> step, the multi criteria decision making *MCDM* will select the most relevant settings and rank them according to their reward.

### IV. EXPERIMENTS

For our experiments we use both real environment (Figure 2) and ns-3 simulator with SX1276 LoRa module. However, to test the scalability of genetic algorithm with numerous IoT devices in a real environment, we use FIT IoT-LAB platform among other platforms presented in [7]. This choice is motivated by the number of devices supported by this platform (up to 2000 nodes).



(a) Gateway (Raspberry-pi). (b) Sensor node (Arduino).

Figure 2. Gateway & Sensor node.

Figure 2a presents the LoRa gateway that we build using a low cost LoRa gateway [1] on a Raspberry-pi. Figure 2b presents one of the two Arduino boards equipped with an antenna that cover both 868 and 433 MHz band with a SX1276 LoRa Transceiver.

Due to the energy constraints of LoRa nodes of class A that we use, our framework will send commands through the FCtrl fields to ask nodes to adapt their transmission behavior to the new application or the new environment conditions.

### V. DISCUSSION

Our main contribution was to build 3 applications that requires 3 different levels of QoS, such as text, sound and image transmission. We used a low cost LoRa gateway on a Raspberry-pi with 2 Arduino boards equipped with 2 LoRa Transceivers based on the Semtech SX1276 specification. The main challenge addressed in this work was to explore the application of genetic algorithm in LoRa transmission parameter selection. The efficiency of such algorithms is measured by the ability to satisfy each application requirement. To measure the accuracy of applying genetic algorithm in an edge computing we expect to compare our approach with other adaptive data rate solutions.

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# EES

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**Abstract**—The number of building sites increased considerably these few last years. Emerging countries invest more and more in building and construction sites to overcome the need of citizens like education, residential and administration buildings. Such sites should be equipped with evacuation systems to allow bricklayers and other persons to be evacuated quickly. Furthermore, evacuation systems should communicate with each other to coordinate evacuation emergency of people. One of the most important problem of such systems is the need of a powerful network technology with sufficient range that could make communication in sites of several kilometers squares. The lack of such mechanisms make the life of billion citizens in a danger. This problem could be addressed by building a novel evacuation emergency application in *IoT* devices with *LPWAN*. One of the known *LPWAN* networks is Long Range (*LoRa*) network. Our work is motivated by the deployment of an Emergency Evacuation Systems (*EES*) in building sites. Our goal is to measure the Quality of Service (*QoS*) of IEEE.802.15.4 and LoRa technologies and adapt their configurations to the emergency situation of the site.

## I. INTRODUCTION

Modern and even old cities have several issues regarding building management, urban mobility, disaster recovery, which can be supported by *IoT* applications. Actually, different kinds of services appear to help city management by automating tasks and predicting critical situations before they occur. *EES* for example could help to locate survivors and plans for rescue missions in case of disasters.

*EES* require high quality level of communication to ensure their proper operations in critical situations. This constraint is more complex when devices need to communicate in a large area like building sites. Thus, the need for a powerful network technology that could transmit data with a long range and sufficient Data Rate (*DR*) is very challenging.

First the architecture that we propose are based on three component, *IoT* devices, smart gateway and edge computing server. *IoT* applications, especially emergency ones in case of disasters, require data to be received and bounded in time intervals. So, service orchestrations is an important rigid constraint. Hence, the orchestration algorithms should adapt the network configuration to the emergency situation of the site. Furthermore, the delay to adapt network behavior to the emergency situation has to be optimized to react quickly when an alarm is triggered. For this reason, we opt for the use of

an edge computing and a smart gateway which could map transmission tuning to the emergency state.

This paper is organized as follows. Section II elucidates summary of related works. Section III highlights the required network technologies that could be used in *EES*. In Section IV, we propose our architecture scheme for *EES*. Section V concludes this paper.

## II. RELATED WORK

Research communities are paying more and more attention for *EES*, this particular attention is mainly due to the interest of building companies to deploy smart sensors in their construction sites. The main factors of a good deployment of such systems is about the quality of the network that makes all emergency agents communicate and collaborate to make good decisions. In another hand, the need of population of new tools to deal with the permanently dynamic changes of pedestrian and car paths in their city, becomes very challenging. For example, Qiu, et al [1] proposed an emergency response system with sensor nodes for real time application they proposed an iterative method for delay optimizing with the lossless transmission, authors tried to find a trade-off between the increased energy consumption and the reduced network lifetime. The framework proposed in by Al Turjman, et al [2] is a multi hop routing method for disaster management in real time. To achieve a higher efficiency, the system, they propose, is designed for using limited multiple hops with the help of left over energy may sometimes result with loss of data due the exhausted energy level.

## III. BACKGROUND

To select the wireless communication that best fit smart building application requirements, four main parameters are generally used: (i) cost; (ii) data rate; (iii) autonomy and (iv) communication range **lopes\_design\_2019**. As data rate and communication range are inversely proportional to autonomy, devices need to consume more energy to communicate at higher bit rates and greater distances. *IoT* devices in emerging evaluation system are particularly vulnerable to power consumption. This vulnerability is mainly due to the awareness aspect of their task and the need to be powered all the time without carrying about battery level for a long period of time.

Characteristics	$CF [Hz]$	6LoWPAN	LoRaWAN
Modulation	2.4G	O-QPSK	-
	915M	BPSK	LoRa
	868M	BPSK	LoRa/GFSK
Channels	2.4G	16	-
	915M	10	64+8↑, 8↓
	868M	1	10
$CF [MHz]$	2.4G	X	-
	915M	902-929	902-928
	868M	868-868.6	863-870 and 780
$BW [Hz]$	2.4G	5M	-
	915M	2M	125K-500K
	868M	600M	125K-250K
$DR [bps]$	2.4G	250M	-
	915M	40M	980-21.9K
	868M	20M	LoRa: 0.3K-37.5K FSK: 50K
$CR [dBm]$	2.4G	-85	-
	915M	-92	X
	868M	-92	-137

Table I. LoRa and 6LowPAN characteristics [3].

Table 1 compiles the two more promising standards that are being used in smart building applications, notably LoRa and 6LowPan.

Carrier Frequency ( $CF$ )  $DR$   $CR$   $BW$   
raza\_low\_22\_lopes\_design\_2019

#### IV. ARCHITECTURE

Research in *IoT* field usually pay attention to energy consumption by optimizing the data transmission rate or by aggregating the data. However, in an urgent situation like a disaster monitoring service, data accuracy becomes very important. Furthermore, these data are required to achieve the network server quickly over Constrained Application Protocol (*CoAP*) or Message Queuing Telemetry Transport (*MQTT*) as shown in Figure 1 to make the appropriate decisions to evacuate people. For this reason, we use an edge computing server that receives data from *IoT* devices by means of IEEE.802.15.4 and LoRa technologies. This requires to make the LoRa gateway able to forward received packets to the IEEE802.15.4 border router and vice versa.

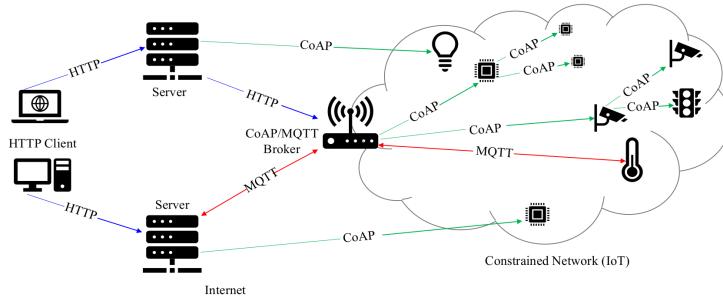


Figure 1. IoT architecture.

#### V. CONCLUSION

*LPWAN*, Wireless Sensor Network (*WSN*) and *IoT* architecture are the first candidates to enhanced disaster monitoring and management systems. Particularly, IEEE802.15.4

and LoRa networks gives new insight for effective *EES*. This work gives an overview of deployment of *IoT* architecture for *EES*. Such services and demand for edge computing in real-time poses new architectural and service orchestration challenges. As a future work, we plan to study the efficiency of using artificial intelligence/machine learning to adapt these two networks to the emergency situation of the building.

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# Template

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**Abstract**—One of the most important problem ... This limitation could be addressed by Our work is motivated by Our goal is to

The difficulty is mainly due to To address this problem,

Particularly,

We validate our approach by

Simulation results show that

Furthermore,

## I. INTRODUCTION

The need of The main factor is

The difficulty to build such system is

For that purpose, we ... In this work we ...

This paper is organized as follows. Section II elucidates summary of related works. Section ?? provide the required background. In Section IV, we ... our approach to solve ... Our experimentation is presented in Section V. Our findings are presented in section ?. Section VII concludes this paper.

## II. RELATED WORK

## III. BACKGROUND

*Selection of technology*

A. *Hardware*

B. *Operating system*

C. *Communication protocol*

D. *Workspace and tools*

## IV. APPROACH

## V. EXPERIMENTATION

## VI. RESULTS EXPLOITATION

A. *Range*

B. *Response time*

C. *Connection speed*

D. *Power consumption*

## VII. CONCLUSION

The main challenge of this work was to The efficiency of such algorithms

Our main contribution was to

To measure the accuracy of

As we find that ..., we plan to ..

## REFERENCES

# Conclusion

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I. CONCLUSION

II. PERSPECTIVES

REFERENCES

# Survey Platforms

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Heterogeneity of applications -> network adaptability to different application network slicing is a solution to -> fixed slicing strategies to maximize resource allocation efficiency -> (on-demand) dynamic slicing strategies maximize resource allocation efficiency

ultra-reliable low-latency communications (URLLC), enhanced mobile broadband (eMBB), and massive machine-type communications (mMTC))

## I. BY PAPERS ADR

### A. Learning

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table I. ADR solutions

### B. Multi-Armed Bandit

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes	Estimation	<b>Closeness</b> have a high degree of

Table II. ADR solutions

### C. Experiments

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table III. ADR solutions

### D. Fuzzy

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table IV. ADR solutions

### E. Game

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table V. ADR solutions

### F. MCDM

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table VI. ADR solutions

### G. Probability

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table VII. ADR solutions

### H. Utility

Year	Factors	Computation Model	Results interpretation
2018	EXPLoRa-SF [1] connected nodes time-on-air	Estimation	<b>Closeness</b> have a high degree of Correlation with <b>privacy score</b>

Table VIII. ADR solutions

## II. BY PAPERS ORCHESTRATION

Year	Factors	Computation Model	Results interpretation
2017	yuzehuang_timeaware_2017	Marcov	
2017	EMMA [2]		
2018	[3]		
2017	EMMA [2]	Service Orchestration	

Table IX. Social metrics

### III. BY TECHNOLOGIES

Chirp Spread Spectrum (Proprietary) (**CSS**) **CF** Forward error correction (**FEC**) Payload length (**PL**) Link Symmetry (**LS**)  
 Base Station (**BS**) **CSS** Direct Sequence Spread Spectrum (**DSSS**) Ultra narrow band (**UNB**) **DR ADR CR BW**

Characteristics	<b>CF</b> [ $\text{Hz}$ ]	6LoWPAN	LoRaWAN	SigFox	NB-IoT	INGENU	TELENS
Modulation	2.4G	O-QPSK	-	-	QSPSK $\downarrow$	RPMA $\uparrow$ , CDMA $\downarrow$	2-FSK
	915M	BPSK	LoRa	BPSK $\uparrow$ , GFSK $\downarrow$	QSPSK multi-tone $\uparrow$		2-FSK
	868M	BPSK	LoRa/GFSK	BPSK $\uparrow$ , GFSK $\downarrow$	/4-QPSK Single-tone		2-FSK
Chwidth [ $\text{kHz}$ ]	2.4G	16	500 - 125	-	180	40	$\times$
Channels	915M	10	64+8 $\uparrow$ , 8 $\downarrow$	$\times$	-		$\times$
	868M	1	10	360+40	$\times$		$\times$
<b>CF</b> [ $\text{MHz}$ ]	2.4G	$\times$	-	-	-	$\times$	ISM
	915M	902-929	902-928	902	$\times$		915M
	868M	868-868.6	863-870 and 780	868.18-868.22	$\times$		868M/430
<b>BW</b> [ $\text{Hz}$ ]	2.4G	5M	-	-	200K	1M	$\times$
	915M	2M	125K-500K	$\times$	$\times$	$\times$	$\times$
	868M	600M	125K-250K	0.1K-1.2K	$\times$	$\times$	$\times$
<b>DR</b> [ $\text{bps}$ ]	2.4G	250M	-	-	-	78K $\uparrow$ , 19,5K $\downarrow$	$\times$
	915M	40M	980-22K	$\times$	234.7 $\downarrow$ , 204.8 $\uparrow$	$\times$	$\times$
	868M	20M	LoRa: 0.3K-37.5K FSK: 50K	0.1K $\uparrow$ , 0.6K $\downarrow$	$\times$	$\times$	62.5 $\uparrow$ , 50
<b>CR</b> [ $\text{dBm}$ ]	2.4G	-85	-	-	-	$\times$	$\times$
	915M	-92	$\times$	$\times$	$\times$	$\times$	$\times$
	868M	-92	-137	-137	$\times$	$\times$	$\times$
<b>ChipR</b> [ $\text{chip/s}$ ]	2.4G						
	915M						
	868M						
Range	2.4G 915M 868M	10-100 m	5-15 Km	10-50 Km	1Km	15-? Km	1Km-?
Handover	2.4G	$\times$	-	-	-	$\times$	$\times$
	915M	$\times$	$\times$	$\times$	$\times$	$\times$	$\times$
	868M	$\times$	Multi <b>BS</b>	Multi <b>BS</b>	$\times$	$\times$	$\times$
msg/day	2.4G	$\times$	-	-	-	$\times$	$\times$
	915M	$\times$	$\times$	$\times$	$\times$	$\times$	$\times$
	868M	$\times$	Unlimited	140 $\uparrow$ , 4 $\downarrow$	Unlimited	$\times$	$\times$
<b>PL</b> B	2.4G	$\times$	-	-	-	$\times$	$\times$
	915M	$\times$	$\times$	$\times$	$\times$	$\times$	$\times$
	868M	$\times$	51 - 243	12 $\uparrow$ , 8 $\downarrow$	1600B	10KB	$\times$
Coding		DSSS	CSS	UNB	$\times$	DSSS	UNB
Proprietary		$\times$	$\times$	$\checkmark$	$\times$	$\times$	$\times$
Topology		$\times$	Star, Stars	Star	$\times$	Star, Tree	Star
<b>ADR</b>		$\times$	$\checkmark$	$\times$	$\times$	$\checkmark$	$\times$
Security		$\times$	AES 128b	$\times$	$\times$	AES 256B	$\times$
<b>LS</b>		$\times$	$\checkmark$	$\times$	$\times$	$\times$	$\times$
<b>FEC</b>		$\times$	AES 128b	$\times$	$\times$	$\checkmark$	$\times$
Battery		1-2 years	<10 years	<10 years	<10 years		
Cost		Free	35e	25e	1020e		
Standar		IETF	LoRa Alliance		3GPP		
Duplex			Half		Half		
Mob support			High and Simple		High and complex		
Mob latency			Low		High (1.6 - 10 s)		
$Tx$ [ $\text{dBm}$ ]			+14 - +27		20/23		

Real-Time			Class C	<b>X</b>	
Scalability			1M↑, 100K↓	55 k	
<i>Linkbudget</i> [dB]			157	154	
<i>Sensitivity</i> [dBm]			-124 - (-134)	-141	
Multi-hop supporter			<b>X</b>	<b>X</b>	
Addressing			Broadcast↑, Unicast↓	Unicast↑, Both↓	
Peak current			32 mA	120300 mA	
Sleep current			1 A	5 A	

Table X. LPWAN Characteristics [4], lopes\_design\_2019, raza\_low\_22, [5]

Characteristics	$CF_{[Hz]}$	ZigBee	LoRaWAN	SigFox	NB-IoT	INGENU	TELENSA
<b>Modulation</b>	2.4G 915M 868M	O-QPSK BPSK BPSK					
<b>Channels</b>	2.4G 915M 868M	16 10 1					
$CF_{[MHz]}$	2.4G 915M 868M	2.4835 902, 928 868, 868.6					
$BW_{[Hz]}$	2.4G 915M 868M						
$DR_{[b/s]}$	2.4G 915M 868M	250 kbps 40 kbps 20 kbps					
$CR_{[dBm]}$	2.4G 915M 868M						
$ChipR_{[chip/s]}$	2.4G 915M 868M	2M 600K 300K					
<b>Handover</b>	2.4G 915M 868M						
<b>msg/day</b>	2.4G 915M 868M						
<b>PL B</b>	2.4G 915M 868M						
<b>Coding</b>							
<b>Proprietary</b>							
<b>Topology</b>							
<b>ADR</b>							
<b>Security</b>							
<b>LS</b>							
<b>FEC</b>							
<b>Range</b>							
<b>Battery</b>							
<b>Cost</b>							
<b>Standar</b>	IEEE 802.15.4						

Table XI. LPWAN Characteristics [6]

Standard	802.15.4k	802.15.4g	Weightless-W	Weightless-N	Weightless-P	DASH 7 Alliance
Modulation	DSSS, FSK	MR-[FSK, OFDMA, OQPSK]	16-QAM, BPSK, QPSK, DBPSK	UNB DBPSK	GMSK, offset-QPSK	GFSK
BW	ISM S UB -GH Z, 2.4GHz	ISM S UB -GH Z, 2.4GHz	TV white spaces 470-790MHz	ISM S UB -GH Z EU (868MHz), US (915MHz)	S UB -GH Z ISM or licensed	UB -GH Z 433MHz, 868MHz, 915MHz
DR	1.5 bps-128 kbps	4.8 kbps-800 kbps	1 kbps-10 Mbps	30 kbps-100 kbps	200 bps-100kbps	9.6,55.6,166.7 kbps
Range	5 km ( URBAN )	up to several kms	5 km ( URBAN )	3 km ( URBAN )	2 km ( URBAN )	0-5 km ( URBAN )
MAC	CSMA/CA, CSMA/CA or A LOHA with PCA	CSMA/CA	TDMA/FDMA	slotted A LOHA	TDMA/FDMA	CSMA/CA
Topology	star	tar, mesh, peer-to-peer	star	star	star	tree, star
PL	2047B	2047B	>10B	20B	>10B	256B
Security	AES 128b	AES 128b	AES 128b	AES 128b	AES 128/256b	AES 128b
Forward error correction	✓	✓	✓	✗	✓	✓

Table XII. raza\_low\_22

Phy protocol	IEEE 802.15.4	BLE	EPCglobal	Z-Wave	LTE-M	ZigBee
Standard		IEEE 802.15.1				IEEE 802.15.4, ZigBee Alliance
BW(MHz)	868/915/2400	2400	860-960	868/908/2400	700-900	
MAC	TDMA, CSMA/CA	TDMA	ALOHA	CSMA/CA	OFDMA	
DR (bps)	20/40/250 K	1024K	varies 5-640K	40K	1G (up), 500M (down)	
Throughput				9.6, 40, 200kbps		
Scalability	65K nodes	5917 slaves	-	232 nodes	-	
Range	10-20m	10-100m				
Addressing	8bit	16bit				

Table XIII. IoT cloud platforms and their characteristics al-fuqaha\_internet\_24

	802.15.4	802.15.4e	802.15.4g	802.15.4f
CF	2.4Ghz (DSSS + oQPSK) 868Mhz (DSSS + BPSK) 915Mhz (DSSS + BPSK)	2.4Ghz (DSSS + oQPSK, CSS+DQPSK ) 868Mhz (DSSS + BPSK) 915Mhz (DSSS + BPSK)	2.4Ghz (DSSS + oQPSK, CSS+DQPSK ) 868Mhz (DSSS + BPSK) 915Mhz (DSSS + BPSK)	2.4Ghz (DSSS + oQPSK,CSS+DQPSK ) 868Mhz (DSSS + BPSK) 915Mhz (DSSS + BPSK) 3~10Ghz (BPM+BPSK )
DR Differences	Upto 250kbps -	Upto 800kbps Time sync and channel hopping	Up to 800kbps Phy Enhancements	Mac and Phy Enhancements
PL Range Goals	127 bytes 1 ~75+ m General Low-power Sensing/Actuating	N/A 1 ~75+ m Industrial segments	Up to 2047 bytes Upto 1km Smart utilities	N/A N/A Active RFID
Products	Many	Few	Connnode (6LoWPAN)	LeanTegra PowerMote

Table XIV. IEEE 802.15.4 standards [7]

Feature	Wi-Fi	802.11p	UMTS	LTE	LTE-A
Channel MHz	20	10	5	1,4, 3, 5, 10, 15, 20	<100
Frequency band(s) GHz	2.4 , 5.2	5.86-5.92	0.7-2.6	0.7-2.69	0.45-4.99
BR Mb/s	6-54	327	2	<300	<1000
Range km	<0.1	<1	<10	<30	<30
Capacity	Medium	Medium	✗	✓	✓
Coverage	Intermittent	Intermittent	Ubiquitous	Ubiquitous	Ubiquitous
Mobility support km/h	✗	Medium	✓	<350	<350
QoS support	EDCA Enhanced Distributed Channel Access	EDCA Enhanced Distributed Channel Access	QoS classes and bearer selection	QCI and bearer selection	QCI and bearer selection
Broadcast/multicast support	Native broadcast	Native broadcast	Through MBMS	Through eMBMS	Through eMBMS
V2I support	✓	✓	✓	✓	✓
V2V support	Native (ad hoc)	Native (ad hoc)	✗	✗	Through D2D
Market penetration	✓	✗	✓	✓	✓
DR	<640 kbps	250 kbps	106424 kbps	✓	✓

Table XV. An example table.

Payload size (*PS*) Signal-to-interference & noise ratio (*SINR*)

<i>SF/BW</i>	125kHz					250kHz					500kHz				
	[8]	[9]	[10]	<i>PS</i>		Sensitivity	<i>BR</i>	Rx wind	<i>SINR</i>		Sensitivity	<i>BR</i>	Rx wind	<i>SINR</i>	
	[dBm]	[kb/s]	[ms]	[dB]	Byte	[dBm]	[kb/s]	[ms]	[dB]		[dBm]	[kb/s]	[ms]	[dB]	
-	-118	5.468	5.1	-7.5	242+13	-115					-111				
6	-123	3.125	10.2	-10	242+13	-120					-116				
7	-126	1.757	20.5	-12.5	115+13	-123					-119				
8	-129	0.976	41.0	-15	51+13	-125					-122				
9	-132	0.537	81.9	-17.5	51+13	-128					-125				
10	-133	0.293	163.8	-20	51+13	-130					-128				
11						-133					-130				
12															

Table XVI. Receiver sensitivity [dBm]

evaluation Nous avons vu en effet plus haut qu'il a été démontré que la méthode CSMA est plus efficace pour le traitement des faibles trafics, tandis que TDMA est nettement plus appropriée pour supporter les trafics intenses. *PS*

<i>DR</i>	Modulation			<i>PS</i>	<i>BR</i>
	<i>SF</i>	<i>BW</i> [kHz]	<i>CR</i>	Byte	x kbit/s
0	12	125	4/6	51+13	0.25
1	11	125	4/6	51+13	0.44
2	10	125	4/5	51+13	0.98
3	9	125	4/5	115+13	1.76
4	8	125	4/5	242+13	3.125
5	7	125	4/5	242+13	5.47
6	7	125	4/5	242+13	11
7		125	4/5	242+13	50

Table XVII. oioioi

#### IV. BY PAQUET

##### A. CoAP

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31																	
Ver	T	TKL	Code										Message ID																																			
Token																																																
Options																																																
11111111										Payload																																						

Figure 1. CoAP Header.

- **Ver:** is the version of CoAP
- **T:** is the type of Transaction
- **TKL:** Token length
- **Code:** represents the request method (1-10) or response code (40-255).
  - Ex: the code for GET, POST, PUT, and DELETE is 1, 2, 3, and 4, respectively.
- **Message ID:** is a unique identifier for matching the response.
- **Token:** Optional response matching token.

##### B. MQTT

0	1	2	3	4	5	6	7		
Message Type			UDP	QoS Level		Retain			
Remaining length									
Variable length header									
Variable length message payload									

Figure 2. CoAP Header.

- **Message type:** CONNECT (1), CONNACK (2), PUBLISH (3), SUBSCRIBE (8) and so on
- **DUP flag:** indicates that the message is duplicated
- **QoS Level:** identify the three levels of QoS for delivery assurance of Publish messages
- **Retain field:** retain the last received Publish message and submit it to new subscribers as a first message

##### C. XMPP

- Extensible Messaging and Presence Protocol
- Developed by the Jabber open source community
- An IETF instant messaging standard used for:
  - multi-party chatting, voice and telepresence
- Connects a client to a server using a XML stanzas
- An XML stanza is divided into 3 components:
  - message: fills the subject and body fields
  - presence: notifies customers of status updates
  - iq (info/query): pairs message senders and receivers
- Message stanzas identify:
  - the source (from) and destination (to) addresses
  - types, and IDs of XMPP entities

##### D. AMQP

- **Size:** the frame size.
- **DOFF:** the position of the body inside the frame.
- **Type:** the format and purpose of the frame.
  - Ex: 0x00 show that the frame is an AMQP frame
  - Ex: 0x01 represents a SASL frame.

#### *E. DDS*

- ➡ Data Distribution Service
- ➡ Developed by Object Management Group (OMG)
- ➡ Supports 23 QoS policies:
  - ➡ like security, urgency, priority, durability, reliability, etc
- ➡ Relies on a broker-less architecture
  - ➡ uses multicasting to bring excellent Quality of Service
  - ➡ real-time constraints
- ➡ DDS architecture defines two layers:
  - ➡ **DLRL:** Data-Local Reconstruction Layer
    - \* serves as the interface to the DCPS functionalities
  - ➡ **DCPS:** Data-Centric Publish/Subscribe
    - \* delivering the information to the subscribers
- ➡ 5 entities are involved with the data flow in the DCPS layer:
  - ➡ Publisher: disseminates data
  - ➡ DataWriter: used by app to interact with the publisher
  - ➡ Subscriber: receives published data and delivers them to app
  - ➡ DataReader: employed by Subscriber to access received data
  - ➡ Topic: relate DataWriters to DataReaders
- ➡ No need for manual reconfiguration or extra administration
- ➡ It is able to run without infrastructure
- ➡ It is able to continue working if failure happens.
- ➡ It inquires names by sending an IP multicast message to all the nodes in the local domain
  - ➡ Clients ask devices that have the given name to reply back
  - ➡ the target machine receives its name and multicasts its IP @
  - ➡ Devices update their cache with the given name and IP @

#### *F. mDNS*

- ➡ Requires zero configuration aids to connect machine
- ➡ It uses mDNS to send DNS packets to specific multicast addresses through UDP
- ➡ There are two main steps to process Service Discovery:
  - ➡ finding host names of required services such as printers
  - ➡ pairing IP addresses with their host names using mDNS
- ➡ Advantages
  - ➡ IoT needs an architecture without dependency on a configuration mechanism
  - ➡ smart devices can join the platform or leave it without affecting the behavior of the whole system
- ➡ Drawbacks
  - ➡ Need for caching DNS entries

#### *G. Lora*

Preamble		Sync msg		PHY Header		PHDR-CRC		PHY Payload								CRC					
Modulation	length	Sync msg	PHY Header	PHDR-CRC		MAC Header		MAC Payload						MIC	CRC Type	Polynomial					
Modulation	length	Sync msg	PHY Header	PHDR-CRC	MType	RFU	Major	FCnt	FCnt	FCnt	FCnt	FOpns	FCnt	Frame Payload	MIC	CRC Type	Polynomial				
Modulation	length	Sync msg	PHY Header	PHDR-CRC	MType	RFU	Major	Dev Address				FOpns	FCnt	Frame Payload	MIC	CRC Type	Polynomial				
Modulation	length	Sync msg	PHY Header	PHDR-CRC	MType	RFU	Major	NwkID	NwkAddr	ADR	ADRACKReq	ACK	FPendingRFU	FOpnsLen	Frame Payload	MIC	CRC Type	Polynomial			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21

- 0) **Modulation :**  
 ↗ LorA: 8 Symbols, 0x34 (Sync Word)  
 ↗ FSK: 5 Bytes, 0xC1194C1 (Sync Word)

**Length :**

**Sync msg :**

**PHY Header :** It contains:

↳ The Payload length (Bytes)

**The Code rate**

Optional 16bit CRC for payload

**Phy Header :** CRC If contains CRC of Physical Layer Header

**MType :** is the message type (uplink or a downlink)

whether or not it is a confirmed message (reqst ack)

0000 Join Request

0011 Join Accept

0100 Unconfirmed Data Up

0111 Unconfirmed Data Down

1000 Confirmed Data Up

1011 Confirmed Data Down

1100 RFU

1111 Proprietary

1) **RFU :** Reserved for Future Use

2) **Major :** is the LoRaWAN version; currently, only a value of zero is valid

3) **00** LoRaWAN R1

4) **01-11** RFU

5) **NwkID :** the short address of the device (Network ID): 31th to 25th

6) **ADR :** the short address of the device (Network Address): 24th to 0th

7) **ADR :** Network server will change the data rate through appropriate MAC commands

8) **1** To change the data rate

- ↳ 0 No change
- ↳ 1) **ADRAckReq :** (Adaptive Data Rate ACK Request): if network doesn't respond in 'ADR-ACK-DELAY' time, end-device switch to next lower data rate.
- ↳ 1 if (ADR-ACK-CNT)  $\geq$  (ADR-ACK-Limit)
- ↳ 0 otherwise
- ↳ 0) **ACK :** (Message Acknowledgement): If end-device is the sender then gateway will send the ACK in next receive window else if gateway is the sender then end-device will send the ACK in next transmission.
- ↳ 1 if confirmed data message
- ↳ 0 otherwise
- ↳ 1) **FPending↓/RFU↑ :** (Only in downlink), if gateway has more data pending to be send then it asks end-device to open another receive window ASAP
- ↳ 1 to ask for more receive windows
- ↳ 0 otherwise
- ↳ 2) **FOpnsLen :** is the length of the FOpns field in bytes ↴ 0000 to 1111
- ↳ 3) **FCnt :** 2 type of frame counters
  - ↳ FCntUp: counter for uplink data frame, MAX-FCNT-GAP
  - ↳ FCntDown: counter for downlink data frame, MAX-FCNT-GAP
- ↳ 4) **FOpns :** is used to piggyback MAC commands on a data message
- ↳ 5) **FPort :** a multiplexing port field
  - ↳ 0 the payload contains only MAC commands
  - ↳ 1 to 223 Application Specific
  - ↳ 224 & 225 RFU
- ↳ 6) **FRMPayload :** (Frame Payload) Encrypted (AES, 128 key length) Data
- ↳ 7) **MIC :** is a cryptographic message integrity code
- ↳ 8) **computed over the fields MHDR, FHDR, FPort and the encrypted FRMPayload.**
  - ↳ 9) **CRC :** (only in uplink,
    - ↳ 10) CCITT  $x^{16} + x^{12} + x^5 + 1$
    - ↳ 11) IBM  $x^{16} + x^{15} + x^5 + 1$

## V. BY CLOUD

Paper	Architecture	Availability	Reliability	Mobility	Performance	Management	Scalability	Interoperability	Security
IoT-A									
IoT@Work									
EBBITS									
BETaaS									
CALIPSO									
VITAL									
SENSAI									
RERUM									
RELEYonIT									
IoT6									
OpenIoT									
Apec IoV									
Smart Santander									
OMA Device									
OMA-DM									
LWM2M									
NETCONF Light									
Kura									
MASH									
IoT-iCore									
PROBE-JT									
OpenIoT									
LinkSmart									
IETF SOLACE									
BUTLER									
Codo									
SVELETE									

Table XVIII. An example table.

Platform	COAP	XMPP	MQTT
Arkessa			✓
Axeda			
Etherios			
LittleBits			
NanoService	✓		
Nimbits		✓	
Ninja blocks			
OnePlatformv	✓	✓	
RealTime.io			
SensorCloud			
SmartThings			
TempoDB			
ThingWorx			✓
Xively			✓
Ubidots			✓

Table XIX. IoT cloud platforms and their characteristics

## VI. DIVERS

Naïve modes	Instantaneous Hist. average Clustering
Parametric models	Rarely used Traffic Models Time Series Linear regression ARIMA Kalman filtering ATHENA SETAR Gaussian Maximum Likelihood
Non-Parametric models	k-Nearest Neighbor Locally Weighted Regression Fuzzy Logic Bayes Network Neural Network Include temporal/spatial patterns

Table XX. Taxonomy of prediction models short\_2007

$SINR$  Signal-to-Interference Ratio ( $SIR$ )  $BR$  Bit Error Rate ( $BER$ )  $DR$  Packet Error Rate ( $PER$ )  $BW$  Packet Reception Rate ( $PRR$ )  $PDR$  Packet delivery ratio ( $PLR$ )  $SNR$  Packet loss rate ( $RTT$ )  $Tx$   $PS$  Traffic congestion ( $TC$ )  $DC$  Duty cycle ( $SR$ ) Symbol Rate ( $SL$ ) Sleep time ( $Jit$ ) Jitter ( $CCI$ ) Co-channel Interference ( $ToA$ ) Time on Air ( $PL$ ) Mobility ( $Mob$ ) Throughput ( $Th$ ) Service Cost ( $SC$ ) Sensitivity ( $Sen$ ) Received Signal Strength Indication ( $RSSI$ ) Time Complexity ( $O_{time}$ ) Space Complexity ( $O_{space}$ ) Payload length ( $PktL$ ) Receiver Sensitivity ( $RS$ )

$$SF = \log_2 \frac{R_c}{SR} \quad (8)$$

$$R_c[\text{chips/s}] = BW \quad (9)$$

$$R_S[\text{symbols/s}] = \frac{R_c}{2^{SF}} = \frac{BW}{2^{SF}} \quad (10)$$

$$R_m = DR[\text{bps}] = SF \cdot RS = SF * \frac{BW[\text{Hz}]}{2^{SF}} * CR \quad (11)$$

$$(12)$$

Setting	Values	Rewards	Costs
$BW$	7.8 $\rightarrow$ 500 [ $kHz$ ]	$DR$	$RS$ , Range
$SF$	$2^6 \rightarrow 2^{12}$	$RS$ , Range	$DR$ , $SNR$ , $PktL$ , $Tx$
$CR$	4/5 $\rightarrow$ 4/8	Resilience	$PktL$ , $Tx$ , $ToA$
$Tx$	-4 $\rightarrow$ 20 [ $dBm$ ]	$SNR$	$Tx$

Table XXI. 1

Layer	Maximize (Reward)	Minimize (Cost)
Application Network	$\text{Sec}$ security $\text{PRR} = (1 - \text{BER})^L$ $\text{*PDR}$ $\text{*Range}$ $\text{*Th}$	$\text{SC}$ Service Cost $\text{*PS}$ $\text{*Jit}$ $\text{*TC}$ $\text{*RTT}$ $\text{PER}_{[\text{pps}]} = 1 - (1 - \text{BER})^{n_{\text{bits}}}$ $\text{*O}_{\text{time}}$ $\text{*O}_{\text{space}}$
Radio	$\text{SNR}_{[\text{dB}]} = 20 \cdot \log(\frac{S}{N})$ $\text{RSSI} = T x_{\text{power}} \cdot \frac{\text{Rayleigh power}}{PL}$ $\text{SR}_{[\text{sps}]} = \frac{BW}{2^{SF}}$ $\text{BR}_{[\text{bps}]} = SF * \frac{4}{\frac{4+CR}{BW}}$ $\text{PL} =  RSSI  + SNR + P_{TX} + G_{RX}$ <b>petajajarvi_coverage_2015</b> $\text{Sen}_{[\text{dBm}]} = -174 + 10 \log_{10} BW + NF + SNR$ [12] $\text{*DC}$ Duty cycle $\text{*Mob}$ $\text{*SINR}$ $\text{*SL}$ Sleep time $\text{*SIR}$	$\text{BER}_{[\text{bps}]} = 10^{\alpha \cdot e^{-\beta \text{SNR}}}$ $\text{BER}_{[\text{bps}]} = \frac{8}{15} \cdot \frac{1}{16} \cdot \sum k = 216 - 1^k (\frac{16}{k}) e^{20 \cdot \text{SINR}(\frac{1}{k} - 1)}$ $\text{*Tx}$ $\text{*CCI}$ $\text{ToA}_s = \frac{2^{SF}}{BW_{[\text{Hz}]}}$

Table XXII. Network selection inputs and classification of parameters [13] + QoS parameters [14] [15]

### fakhfakh\_deep\_2017

$$\text{MSE} = \frac{1}{n} \sum_{i=1}^n (p_i - r_i)^2 \quad (13)$$

$$\text{RMSE} = \sqrt{\frac{1}{n} \sum_{i=1}^n (p_i - r_i)^2} \quad (14)$$

$$\text{MAE} = \frac{1}{n} \sum_{i=1}^n |p_i - r_i| \quad (15)$$

$$\text{Recall} = \frac{TP}{TP + FN} \quad (16)$$

$$\text{Precision} = \frac{TP}{TP + FP} \quad (17)$$

$$F1\_Score = \frac{2 \times \text{Precision} \times \text{Recall}}{\text{precision} + \text{recall}} \quad (18)$$

$$\text{TPR} = \frac{TP}{TP + FN} \quad (19)$$

$$\text{FPR} = \frac{FP}{FP + TN} \quad (20)$$

$$\text{ROC} = (TPR, FPR) \quad (21)$$

$$\text{Novelty} = \sum_{i \in L} \frac{\log_2 P_i}{n} \text{ where } P_i = \frac{n - rank_i}{n - 1} \quad (22)$$

$$\text{Serendipity} = \frac{1}{n} \sum_{i \in n} \max(P_{\text{user}} - P_U, 0) \times rel_i \quad (23)$$

$$\text{diversity} = \frac{a}{c} \sum_{i=1}^c \frac{1}{n} \sum_{j=1}^n i_j \quad (24)$$

$$\text{Coverage} = 100 \times \frac{u}{U} \quad (25)$$

$$\text{Stability} = \frac{1}{P_2} \sum_{i \in P_2} |P_{2,i} - P_{1,i}| \quad (26)$$

$$\text{DCG} = rel_1 + \sum_{i=2}^{\text{pos}} \frac{rel_i}{\log_2 i} \quad (27)$$

$$\text{IDCG} = rel_1 + \sum_{i=2}^{|h|-1} \frac{rel_i}{\log_2 i} \quad (28)$$

$$\text{NDCG} = \frac{\text{DCG}}{\text{IDCG}} \quad (29)$$

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