## DBMJ

| 1) | Views &   |
|----|---|
|    | - Views in sql or Kind of vistual tables. A view also has         |
|    | yours and columns as they are in a seal table in the database.    |
|    | We can execute a view by selecting fields from one or more        |
|    | tables present in the dotabase. A view can extres have all the    |
|    | solus of a table or specific solus based on certain condition     |
|    | 2 6   |
|    | i) (regting views:  |
|    | - We can create view using CREATE VIEW STATEMENT. A               |
|    | you an be mated from a single table or multiple tables            |
|    | 1.  |
|    | ii) Deloting views:   |
|    | - Sal alows be to delete an existing view we can delete of        |
|    | drup a view using the DROP Statement.                             |
|    |   |
|    | ii) Vedating views:   |
|    | - We an update the view by using replace Statement                |
|    |   |
| 2  |   |
| 2  | - Kigger :  |
|    | - A trigger is a stored procedure in database which automotically |
|    | invokes whenever a special event in the dotabase occur. Fox       |
|    | ex, a trigger can be invoked when a row is invoked into a         |
|    | specified table or when textain table tolumns are being           |
|    | updated   |
|    | Syntax: (seak trigger (trigger name)                              |
|    | [betore after]  |
|    | 2 injert 1 update 1 delett 3                                      |
|    | on [table_name]   |
|    | [ fox each you]   |
|    | [trigger-body]  |

| 3  | Stored procedures:                                       |
|----|--|
|    | - Stored procedures are executed to pertorm one or more  |
|    | DML operations on Database It il nothing but the group   |
|    | of SQL Statements that accepts some input in The form of |
|    | parameters and performs some task and may or may         |
|    | hot Jeturn a Value                                       |
|    | - The most important part is parameters parameters are   |
|    | wild to pass values to the procedure. There are 3        |
|    | different type of parameters                             |
|    | 1) IN a Default parameter, societed value                |
|    | 2) OUT: Sendy value to the calling program               |
|    | 3) IN OUT: It receives and sends yalves                  |
| 7  |  |
| 4  | Properties of the transactions [ACID]:                   |
|    | 21 1   |
|    | 1) Atomicity   |
|    | ii) Consistency  |
|    | iii) Isolation   |
|    | in Duxability  |
| -  |  |
| 5) | Transaction States:                                      |
|    | 1 1  |
|    | i) Active  |
|    | iif Pantially Committed                                  |
|    | nij Failed   |
|    | iv] Aboxted  |
|    | v) (opmitted   |
|    |  |
|    |  |

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| 6 | Seria lisability:  |
|   | - Saxializability is a conjust that helps us to check          |
|   | Which scheduly are consulizable of Senarzaba executive         |
|   | the one that always leaves the database in Consistent          |
|   | State  |
|   | a) Conflict Sexializability:                                   |
|   | - Since we are dealing with only sead and winter               |
|   | instructions there are four cases that we need to consider.    |
|   | i] T= sead (Q), J = sead (Q)                                   |
|   | ii) I = read (Q), J = work (Q)                                 |
|   | ii) I= work (Q), J= read (Q)                                   |
|   | 11) I= work (Q), I= work (Q)                                   |
|   | 3  |
| 1 | Caradeles Schedules:   |
|   | - The phenomenon, in which a single transaction failure        |
|   | leads to a series of transaction sollback, is Called           |
|   | Caxading vollback.   |
|   | - (availing sollback is undestrable, since it leads to         |
|   | The undoing of a significant comount of work It is             |
|   | desirable to restrict the schedules to those where cascading   |
|   | Sollbacks cannot occur such schedules ask Called               |
|   | Cascadoles Schedules.  |
|   |  |
| 8 | Transaction Isolation levels:                                  |
|   | - The uslation lovels specified by sal standard are of follows |
|   |  |
|   | 1) Senalizable   |
|   | ii) Repeatable xead  |
|   | III) Read Compositted  |
|   | is) head uncommitted   |
|   | ii) Repeatable read  |

a lock-Based Potocols. - There are various modes in which a data term may be forked In this section, there are two modes: i) Shared : Ti can read but cannot write, a in Exclusive: Ti can both knad and write Q 10] Two-phase Locking protocol: - This protocol requires that each transaction issue lock and unlock requests in Just Phases: i) Granulag phase: May objain locks but cannot reliage it ii Shonking phase may release locks, but cannot obtain't 1) Deadlock Handling: - A system is in a deadlock state it there exists a set of transaction such that every transaction in the set it waiting for another transaction the Set. - The only simedy to this undestrable Situation is for the System to govoke some drashiz action, such as xolling back some of the transactions govolved in deadlax - There are two principal methods for dealing with the deadlock problem are i) deadlock prevention ii) deadlock detection & xerovery Scheme

| 12) | Recovery Concept:  - The Recovery from deadlock, the three actions head  to be taken axe: |
|-----|---|
|     | i) Selection of a viction   |
|     | ii] Rollback  |
|     | hij Starration  |
| 13) | (gregories of Storage:  |
|     | - The three categories of Storage are   |
|     |   |
|     | 1) Volatile Storage   |
|     | ii) Non-volatile Storage  |
|     | III Stable Storage  |
| 147 | Recovery Algorithm:   |
|     | - i) Transaction Rollback:  |
|     | -ii] Recovery After a system (xash:   |
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