

Genetic Algorithm

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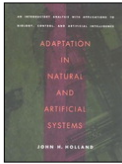
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Outline

- Genetic algorithm
- Binary-coded Genetic Algorithm
- Working of binary GA with an example
- Drawbacks of binary GA
- Real-coded Genetic Algorithm
- Working of real GA with an example
- Comparison of algorithms

Genetic Algorithm

Books > Adaptation in Natural and Art...



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1 Author(s) John H. Holland

Book Abstract

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Genetic Algorithms in Search, Optimization and Machine Learning

Author: [David E. Goldberg](#)

Publication:



• Book

Genetic Algorithms in Search, Optimization and Machine Learning

1st

Addison-Wesley Longman Publishing Co., Inc. Boston, MA, USA ©1989

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Multi-Objective Optimization Using Evolutionary Algorithms

Authors: [Kalyanmoy Deb](#)
[Deb Kalyanmoy](#)

Publication:



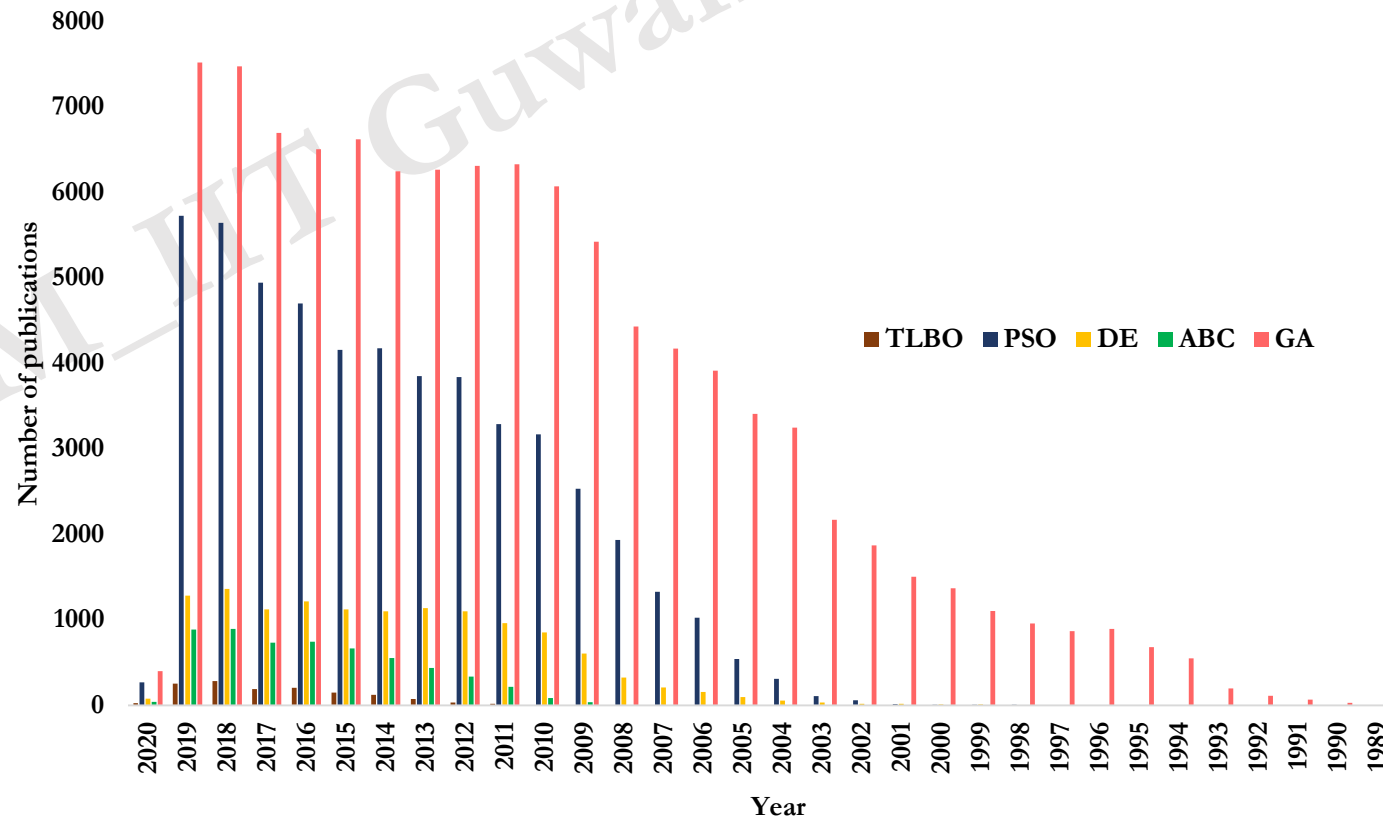
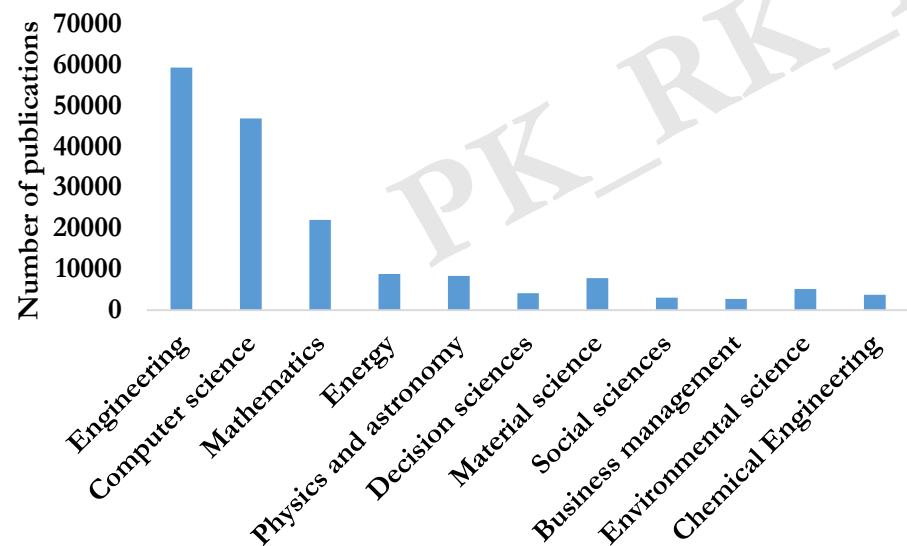
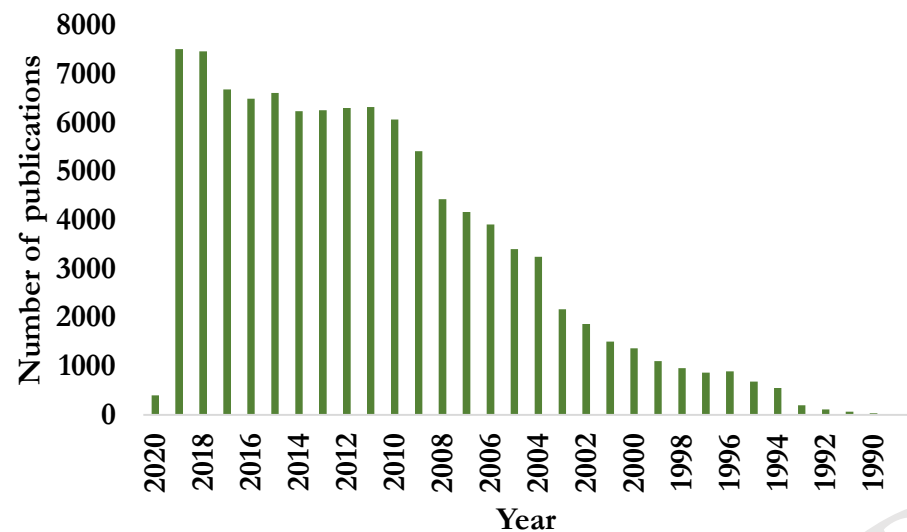
• Book

Multi-Objective Optimization Using Evolutionary Algorithms

John Wiley & Sons, Inc. New York, NY, USA ©2001

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Genetic Algorithm



Genetic Algorithm (GA)

- Inspired by the principles of natural genetics and selection (Darwin's principle of natural evolution).
- Solution vectors are termed as chromosome.
- **Parent**: Solution from which new solutions are generated.
- **Offspring**: Newly generated solutions.
- Offspring are generated through
 - reproduction (selection of good solutions for mating).
 - variation (crossover and mutation).
- Selection of good solutions after variation operator.
- Better candidate has more chance to survive in an environment of limited resources.
 - Good solutions are retained and bad solutions are eliminated.

Binary coded GA

- Real variables are encoded into binary variables (0 and 1).
- If the bit length (n) is specified for a variable, then exactly 2^n different solutions are possible between the bounds of the variable.

Let a binary string of length (n) 5 be used to represent a variable x with bounds $[x_{\min}, x_{\max}]$.

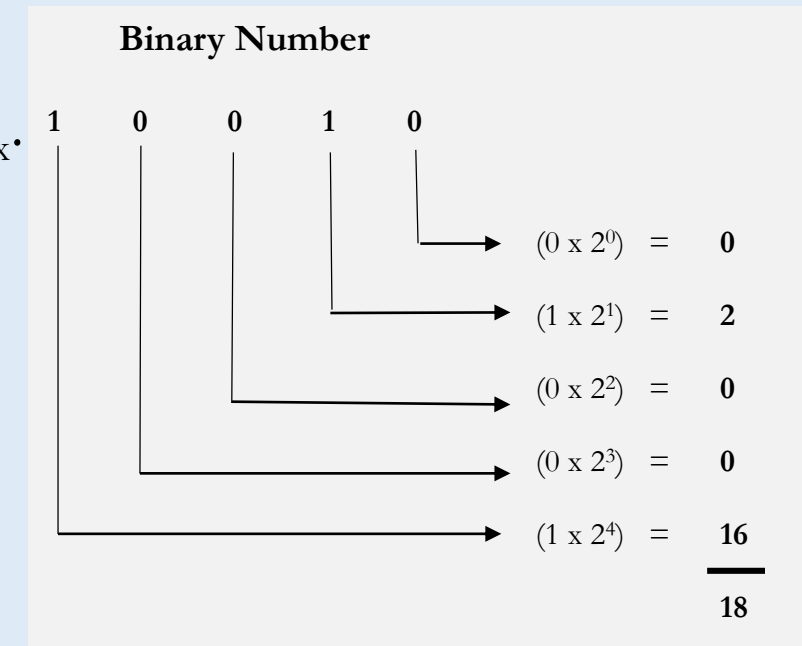
String $[0\ 0\ 0\ 0\ 0]$ will map to x_{\min} and $[1\ 1\ 1\ 1\ 1]$ will map to x_{\max} .

Decoded value of a string can be determined as

$$DV = b_n 2^{n-1} + b_{n-1} 2^{n-2} + \dots + b_2 2^1 + b_1 2^0$$

where binary string $s = [b_n \ b_{n-1} \ \dots \ b_2 \ b_1]$

Maximum decoded value of binary string is 31 ($2^n - 1$) and the minimum is 0.



Binary-coded GA

➤ For a fixed bit length (n), the value of a real variable x within bounds $[x_{\min}, x_{\max}]$ is

$$x = x_{\min} + \left(\frac{x_{\max} - x_{\min}}{2^n - 1} \right) DV$$

$$Precision = \frac{x_{\max} - x_{\min}}{2^n - 1}$$

➤ For evaluating fitness, x is used.

Example: Consider a variable x with bounds $[5, 30]$ to be represented by a 4 bit binary string.

We will be having 16 ($= 2^4$) different solutions between 5 and 30.

Let the binary string be, $s = [0 \ 1 \ 1 \ 0]$, then $DV = 6$

$$x = 5 + \left(\frac{(30 - 5)}{2^4 - 1} \right) 6 = 15$$

$$(0 \times 2^3) + (1 \times 2^2) + (1 \times 2^1) + (0 \times 2^0) = 6$$

Precision = $25/15 = 1.67$

For higher precision, n should be increased.

Binary to real values

$$x_{min} = 5, x_{max} = 30$$

$n = 3, 2^n = 8$ possible solutions

$$\text{precision} = \left(\frac{30 - 5}{2^3 - 1} \right) = 3.57$$

Binary string	Decoded value	Actual value
000	0	5
001	1	8.57
010	2	12.14
011	3	15.71
100	4	19.29
101	5	22.86
110	6	26.43
111	7	30

$$5 + 0 \times 3.57$$

$$5 + 1 \times 3.57$$

$n = 4, 2^n = 16$ possible solutions

$$\text{precision} = \left(\frac{30 - 5}{2^4 - 1} \right) = 1.67$$

Binary string	Decoded value	Actual value	Binary string	Decoded value	Actual value
0000	0	5	1000	8	18.33
0001	1	6.67	1001	9	20
0010	2	8.33	1010	10	21.67
0011	3	10	1011	11	23.33
0100	4	11.67	1100	12	25
0101	5	13.33	1101	13	26.67
0110	6	15	1110	14	28.33
0111	7	16.67	1111	15	30

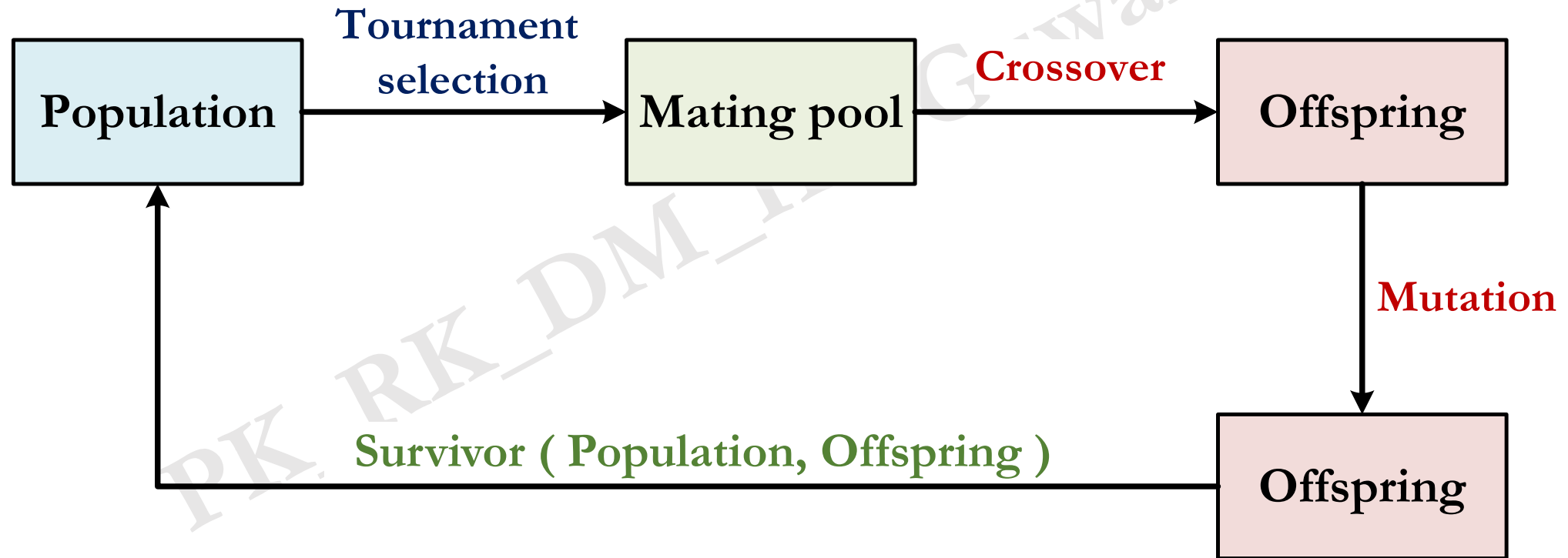
$$5 + 0 \times 1.67$$

$$5 + 1 \times 1.67$$

$$x = x_{min} + \underbrace{\left(\frac{x_{max} - x_{min}}{2^n - 1} \right)}_{\text{precision}} DV$$

By increasing the binary string length by 1 bit, the precision is increased from **3.57** to **1.67**

GA basic flowchart



Population initialization

- Let population size, $N_p = 6$ and the dimension of the problem, $D = 2$.
- Let the number of bits (n) for representing each variable be 5.
- Size of the population matrix will be $N_p \times nD$, (i.e., $= 6 \times 10$).
- Initial population is generated randomly with 0 or 1.

Equal n

$$P = \begin{bmatrix} \begin{matrix} x_1 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 \end{matrix} & \begin{matrix} x_2 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \end{matrix} \end{bmatrix}$$

Unequal n

$$P = \begin{bmatrix} \begin{matrix} x_1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{matrix} & \begin{matrix} x_2 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \end{matrix} \end{bmatrix}$$

Reproduction operator

- Identify good (usually above average) solutions in the population.
- Eliminate bad solutions from the population so that multiple copies of good solutions are considered for variation operator.
- Group of solutions selected through reproduction operator comprise the mating pool.
- **Reproduction:** Selection of good solutions before variation operator.
- Reproduction/selection operators reduces the diversity of the population.

Tournament selection

- **Pool size:** number of members required to be in the mating pool (usually N_p).
- **Tournament size (k):** number of members that participate in a tournament (commonly $k = 2$).
- Tournaments are played between 'k' members and the member with best fitness is selected for mating.
- N_p number of tournaments has to be played.

Let $k = 3$, the solutions selected randomly from the population ($N_p = 6$) be $\{4, 2, 6\}$ and their fitness $f_4 = 27$, $f_2 = 89$ and $f_6 = 12$.

For maximization problem: Solution 2 will win.

For minimization problem: Solution 6 will win.

Binary tournament selection

- If tournaments are played systematically, each solution will have two chances to play.
 - best solution will have two copies.
 - worst solution will never be selected.
 - other solutions can have 0, 1 or 2 copies.
- As the tournament size is increased, the selection pressure of each solution increases.

k	Remarks
2	Worst solution will not be in the mating pool
3	Worst and second worst solution will not be in the mating pool
n	The worst $n-1$ solutions will not be in the mating pool

Variation operator: Single point crossover

- Responsible for generating offspring.
- Two binary strings are chosen randomly from the population to perform crossover.
- Some portion of parent strings are exchanged to generate two offspring.
- Crossover site is a random integer chosen between 1 and nD .
- Crossover of two parents occurs with a probability (p_c).

Let $n = 4$, $D = 2$, crossover site = 3 and the randomly chosen parents be

$$parent_1 = [1 \ 0 \ 0 \mid 0 \ 0 \ 1 \ 0 \ 1]$$

$$parent_2 = [0 \ 1 \ 0 \mid 1 \ 0 \ 1 \ 1 \ 1]$$

$$offspring_1 = [1 \ 0 \ 0 \mid 1 \ 0 \ 1 \ 1 \ 1]$$

$$offspring_2 = [0 \ 1 \ 0 \mid 0 \ 0 \ 1 \ 0 \ 1]$$

Variation operator: Bit-wise mutation

- Change the bit 1 to 0 and vice versa with a mutation probability (p_m).
- Single solution is involved in mutation.
- Every member in the population can potentially undergo mutation.
- For each bit in a string, probability for mutation is checked.

Let $p_m = 0.3$ and the parent string be

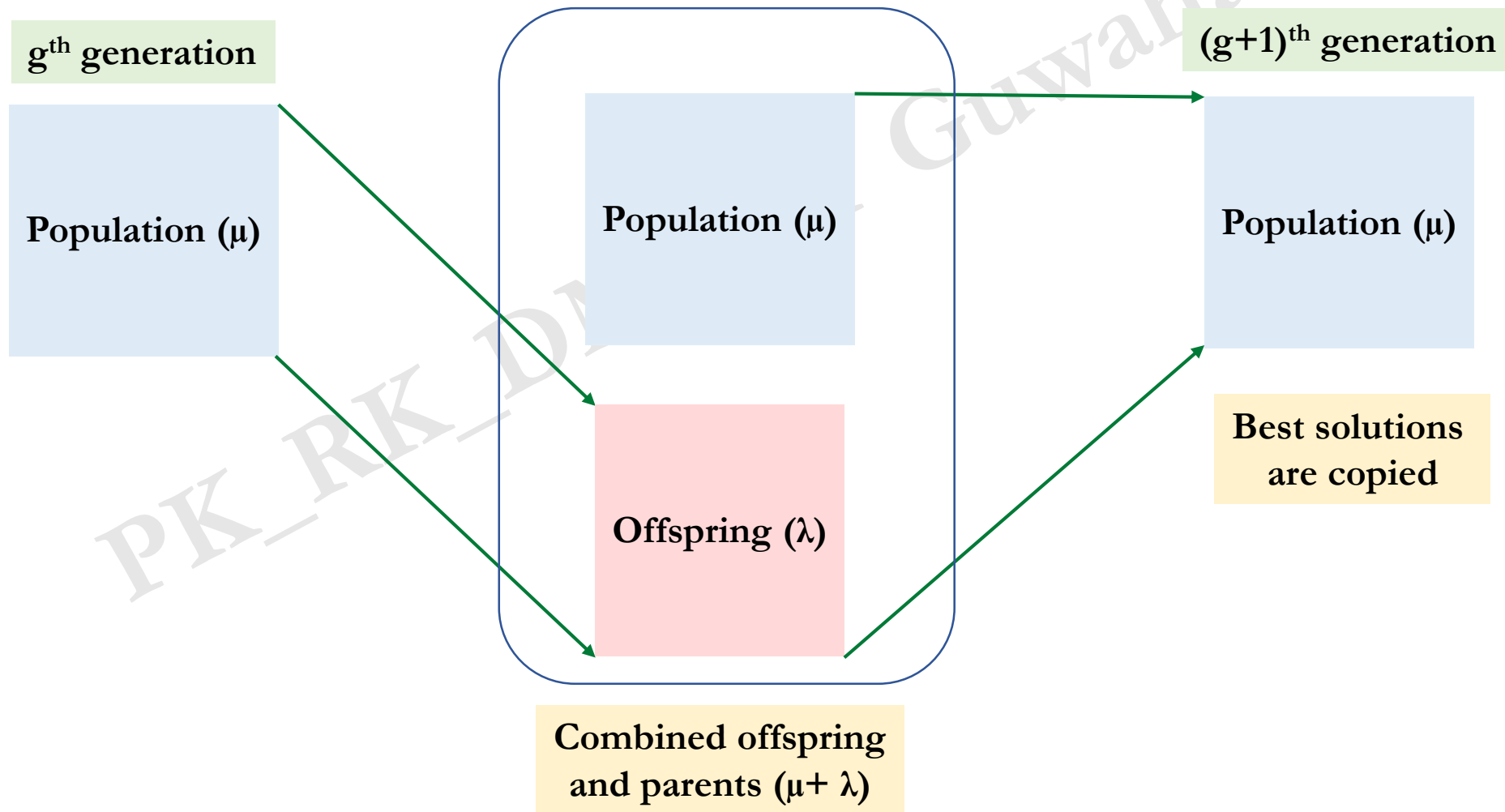
$$offspring = [1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 0 \ 0]$$

Let the random number for each bit be $r = [0.2 \ 0.5 \ 0.6 \ 0.8 \ 0.7 \ 0.1 \ 0.4 \ 0.9]$.

Mutation occurs for 1st ($r_1 < p_m$) and 6th ($r_6 < p_m$) variable and the offspring is

$$offspring = [0 \ 0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 0]$$

Survival of the fittest: $(\mu + \lambda)$ selection



Working of genetic algorithm: Sphere function

Consider $\min f(x) = \sum_{i=1}^4 x_i^2; \quad 0 \leq x_i \leq 30, \quad i = 1, 2$

$$f(x) = x_1^2 + x_2^2$$

- Decision variables: x_1 and x_2
- Step 1: Fix the population size, maximum iterations, bit length, crossover probability, mutation probability

$$N_p = 4, \quad T = 10, \quad n = 4, \quad p_c = 0.8, \quad p_m = 0.3$$

$$x = x_{\min} + \frac{x_{\max} - x_{\min}}{2^n - 1} (DV)$$

$$x = 0 + \frac{30 - 0}{2^4 - 1} (DV) = 2 \times DV$$

- Step 2: Generate random binary solutions

$$P = \begin{bmatrix} \begin{matrix} x_1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \end{matrix} & \begin{matrix} x_2 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{matrix} \end{bmatrix}$$

$$P_D = \begin{bmatrix} 9 & 12 \\ 3 & 7 \\ 4 & 5 \\ 6 & 10 \end{bmatrix}$$

$$P_A = \begin{bmatrix} 18 & 24 \\ 6 & 14 \\ 8 & 10 \\ 12 & 20 \end{bmatrix}$$

$$f = \begin{bmatrix} 900 \\ 232 \\ 164 \\ 544 \end{bmatrix}$$

Selection: Tournament selection

- Step 3: Select two random candidates for tournament.

Let the two candidates be

$$P_2 = [6 \quad 14] \quad f_2 = 232$$

$$P_4 = [12 \quad 20] \quad f_4 = 544$$

- Step 4: Compare fitness to select winner.

$$f_2 < f_4 \rightarrow f_2$$

- Step 5: After four tournaments

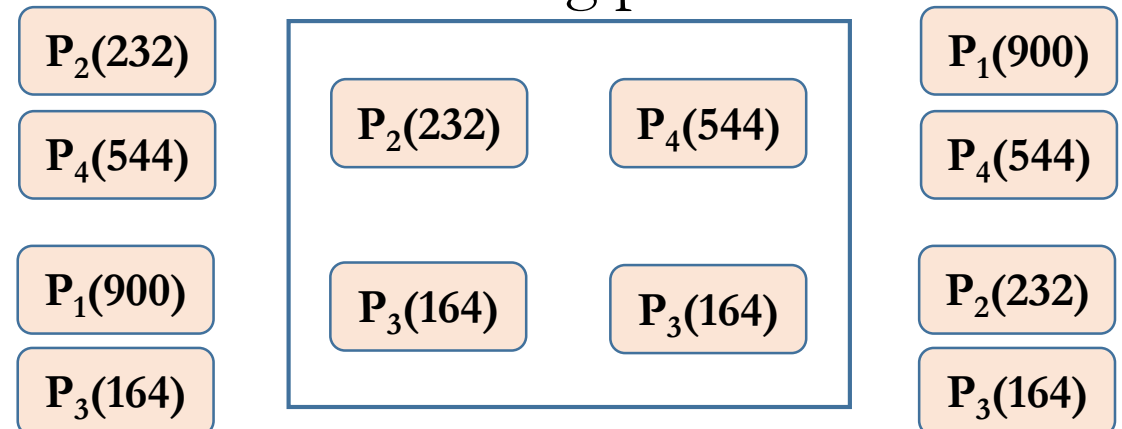
Best solution has two copies.

Worst solution has zero copies.

$$P = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

$$P_A = \begin{bmatrix} 18 & 24 \\ 6 & 14 \\ 8 & 10 \\ 12 & 20 \end{bmatrix} \quad f = \begin{bmatrix} 900 \\ 232 \\ 164 \\ 544 \end{bmatrix}$$

Mating pool



Crossover: Single point crossover

- Step 6: Randomly select a pair of parents for crossover.

Let the two parents be

$$\text{Parent}_1 = [0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 1 \ 1]$$

$$\text{Parent}_2 = [0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 1]$$

$$p_c = 0.8$$

$$\text{Parent} = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \end{bmatrix} \begin{matrix} P_2 \\ P_3 \\ P_4 \\ P_3 \end{matrix}$$

- Step 7: Generate a random number to check if crossover is to be performed.

$$\text{Let } r = 0.2$$

$$r < p_c \rightarrow \text{perform crossover}$$

- Step 8: Select a random crossover site.

$$\text{Let } r = 3$$

$$\text{Parent}_1 = [0 \ 0 \ 1 \ | \ 1 \ 0 \ 1 \ 1 \ 1] \quad \text{offspring}_1 = [0 \ 0 \ 1 \ | \ 0 \ 0 \ 1 \ 0 \ 1]$$

$$\text{Parent}_2 = [0 \ 1 \ 0 \ | \ 0 \ 0 \ 1 \ 0 \ 1] \quad \text{offspring}_2 = [0 \ 1 \ 0 \ | \ 1 \ 0 \ 1 \ 1 \ 1]$$

Crossover: Single point crossover

- Step 9: Randomly select a pair of parents for crossover

Let the two parents be

$$\text{Parent}_3 = [0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0]$$

$$\text{Parent}_4 = [0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 1]$$

$$p_c = 0.8$$

$$\text{Parent} = \begin{bmatrix} 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}$$

- Step 10: Select a random number to check if crossover has to be performed

$$\text{Let } r = 0.6$$

$$r < p_c \rightarrow \text{perform crossover}$$

- Step 11: Select a random crossover site

$$\text{Let } r = 5$$

$$\text{Parent}_3 = [0 \ 1 \ 1 \ 0 \ 1 \ | \ 0 \ 1 \ 0] \quad \text{offspring}_3 = [0 \ 1 \ 1 \ 0 \ 1 \ | \ 1 \ 0 \ 1]$$

$$\text{Parent}_4 = [0 \ 1 \ 0 \ 0 \ 0 \ | \ 1 \ 0 \ 1] \quad \text{offspring}_4 = [0 \ 1 \ 0 \ 0 \ 0 \ | \ 0 \ 1 \ 0]$$

Mutation: bit-wise mutation

- Step 12: Select first offspring for mutation

$$\text{offspring}_1 = [0 \quad 0 \quad 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1]$$

$$p_m = 0.3$$

$$\text{offspring} = O = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

- Step 13: Select random numbers to check if mutation is to be performed

$$\text{Let } r = [0.6 \quad 0.1 \quad 0.5 \quad 0.4 \quad 0.9 \quad 0.7 \quad 0.3 \quad 0.8]$$

$$r < p_m \rightarrow \text{perform mutation}$$

$$\text{offspring}_1 = [0 \quad 0 \quad 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1] \rightarrow \text{offspring}_1 = [0 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1]$$

Mutation: bit-point mutation

- Step 14: Perform mutation for all offspring

$$p_m = 0.3$$

	Offspring	Random number for mutation	New offspring
O_2	[0 1 0 1 0 1 1 1]	[0.5 0.1 0.6 0.3 0.5 0.4 0.6 0.3]	[0 0 0 1 0 1 1 1]
O_3	[0 1 1 0 1 1 0 1]	[0.4 0.5 0.1 0.8 0.5 0.7 0.4 0.2]	[0 1 0 0 1 1 0 0]
O_4	[0 1 0 0 0 0 1 0]	[0.8 0.6 0.3 0.9 0.7 0.5 0.4 0.6]	[0 1 0 0 0 0 1 0]

$$f(x) = x_1^2 + x_2^2$$

$$O = \begin{matrix} & x_1 & & x_2 \\ \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

$$x = x_{\min} + \frac{x_{\max} - x_{\min}}{2^n - 1} (DV)$$

$$\Rightarrow x = 2 \times DV$$

$$O_A = \begin{bmatrix} 12 & 10 \\ 2 & 14 \\ 8 & 24 \\ 8 & 4 \end{bmatrix}$$

$$f = \begin{bmatrix} 244 \\ 200 \\ 640 \\ 80 \end{bmatrix}$$

Survival

- Step 15: Combine all the solutions and select best N_p ($N_p = 4$) solutions

$$P = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \end{bmatrix} \quad f = \begin{bmatrix} 900 \\ 232 \\ 164 \\ 544 \end{bmatrix}$$

$$O = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \quad f_o = \begin{bmatrix} 244 \\ 200 \\ 640 \\ 80 \end{bmatrix}$$

$$P_{combined} = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \quad f_{combined} = \begin{bmatrix} 900 \\ 232 \\ 164 \\ 544 \\ 244 \\ 200 \\ 640 \\ 80 \end{bmatrix}$$

$\leftarrow 4$
 $\leftarrow 2$
 $\leftarrow 3$
 $\leftarrow 1$

$$P = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \end{bmatrix} \quad f = \begin{bmatrix} 80 \\ 164 \\ 200 \\ 232 \end{bmatrix}$$

Population for the next iteration

Pseudocode

Input: Fitness function, lb, ub, N_p , T , n , p_c , p_m , k

1. Initialize a random population (P) of binary string (size: $N_p \times nD$)

2. Evaluate fitness (f) of P \leftarrow FE = N_p

for $t = 1$ to T

Perform **tournament selection** of tournament size, k

for $i = 1$ to $N_p/2$

Randomly choose two parents

if $r < p_c$

Select the crossover site

Generate two offspring using **single-point crossover**

else

Copy the selected parents and their fitness to offspring population

end

end

for $i = 1$ to N_p

Generate D random numbers between 0 and 1

Perform **bit-wise mutation** of i^{th} solution in offspring population

Evaluate the fitness of offspring \leftarrow Max FE = N_p

end

Combine population and offspring population to perform **($\mu + \lambda$)**

end

In one iteration, max #FE = N_p

For T iterations, max #FE = $N_p + N_p T$

crossover site = 3

$parent_1 = [1 \ 0 \ 0 \ | \ 0 \ 0 \ 1 \ 0 \ 1]$

$parent_2 = [0 \ 1 \ 0 \ | \ 1 \ 0 \ 1 \ 1 \ 1]$

$Offspring_1 = [1 \ 0 \ 0 \ | \ 1 \ 0 \ 1 \ 1 \ 1]$

$Offspring_2 = [0 \ 1 \ 0 \ | \ 0 \ 0 \ 1 \ 0 \ 1]$

Generation

$r_6 < p_m$

$parent = [1 \ 0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 0]$

$offspring = [1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 0 \ 0]$

Survival of fittest

Drawbacks of binary GA

➤ Binary GA makes the search space discrete.

➤ Unable to achieve any arbitrary precision.

$$x = \frac{x^{\max} - x^{\min}}{2^n - 1}$$

- If n bits are used to represent a decision variable, then 2^n different values are possible between the lower and upper bound of the variable.
- To increase the precision, n has to be increased.
- Increase in n results in larger dimension and population size.

➤ **Hamming cliffs:** transition to neighboring solution (in real space) needs change in multiple bits (Example: 01111 (= 15) to 10000 (= 16)).

Real coded Genetic Algorithm

- Encoding of real variables to binary is not required.
- Decision variables can be directly used to compute the objective function value.
- Selection operator used in binary GA can also be employed in real GA.
- Naive crossover operators such as single point crossover might fail to perform well.
 - Search within the current values of the decision variables.
 - Depends on mutation operator for a new value of decision variable.
- Modification in the variation operator is required to explore the search space.

Crossover site = 3

Parent 1:	5.9	2.6	7.3	3.5	2.7
Parent 2:	6.5	4.3	3.2	1.1	1.9

Simulated Binary Crossover (SBX)

- Simulates the single-point crossover on binary strings.
- Requires two parents to generate two offspring.
- For a fixed η_c , the offspring have a spread which is proportional to that of parents.

$$O_1 - O_2 = \beta (P'_a - P'_b) \quad \eta_c \text{ is distribution index}$$

- Compute β

$$\beta = \begin{cases} (2u)^{1/(\eta_c+1)} & \text{if } u \leq 0.5 \\ \left(\frac{1}{2(1-u)} \right)^{1/(\eta_c+1)} & \text{otherwise} \end{cases}$$

- Generate offspring

$$O_a = 0.5 \left[(1 + \beta) P'_a + (1 - \beta) P'_b \right]$$

$$O_b = 0.5 \left[(1 - \beta) P'_a + (1 + \beta) P'_b \right]$$

P'_a	Parent 1	O_a	Offspring 1
P'_b	Parent 2	O_b	Offspring 2

Simulated Binary Crossover (SBX)

- Crossover is performed with high probability.
- Two offspring are symmetric about the parents.
- Avoids the bias towards any particular parent in a single crossover operation.
- For a constant β :
 - **Distant parents** result in largely spread offspring and
 - **Near parents** result in closer offspring

Case 1: Let the parents be $P'_a = 2$ and $P'_b = 8$

$$O_1 = 0.5[(1+0.8) \times 2 + (1-0.8) \times 8] = 2.6$$

$$O_2 = 0.5[(1-0.8) \times 2 + (1+0.8) \times 8] = 7.4$$

Case 2: Let the parents be $P'_a = 4$ and $P'_b = 5$

$$O_1 = 0.5[(1+0.8) \times 4 + (1-0.8) \times 5] = 4.1$$

$$O_2 = 0.5[(1-0.8) \times 4 + (1+0.8) \times 5] = 4.9$$

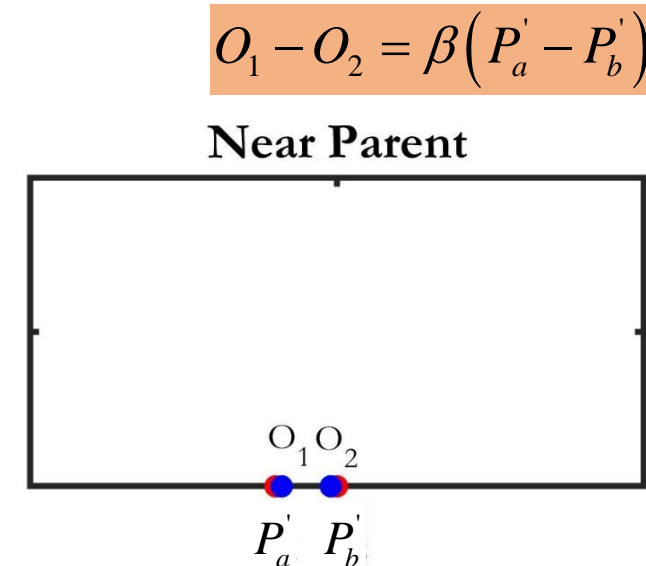
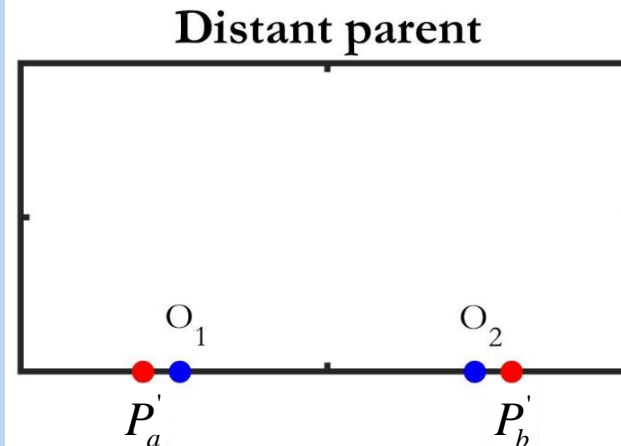


Illustration: Impact of varying β

➤ Consider two solutions $P'_a = 2$ and $P'_b = 5$

➤ Case 1: Contracting crossover ($\beta < 1$)

Take $\beta = 0.6$

$$O_1 = 0.5[(1+0.6) \times 2 + (1-0.6) \times 5] = 2.6$$

$$O_2 = 0.5[(1-0.6) \times 2 + (1+0.6) \times 5] = 4.4$$

Offspring are closer

➤ Case 2: Stationary crossover ($\beta = 1$)

Take $\beta = 1$

$$O_1 = 0.5[(1+1) \times 2 + (1-1) \times 5] = 2$$

$$O_2 = 0.5[(1-1) \times 2 + (1+1) \times 5] = 5$$

Offspring and parents are identical

➤ Case 3: Expanding crossover ($\beta > 1$)

Take $\beta = 1.4$

$$O_1 = 0.5[(1+1.4) \times 2 + (1-1.4) \times 5] = 1.4$$

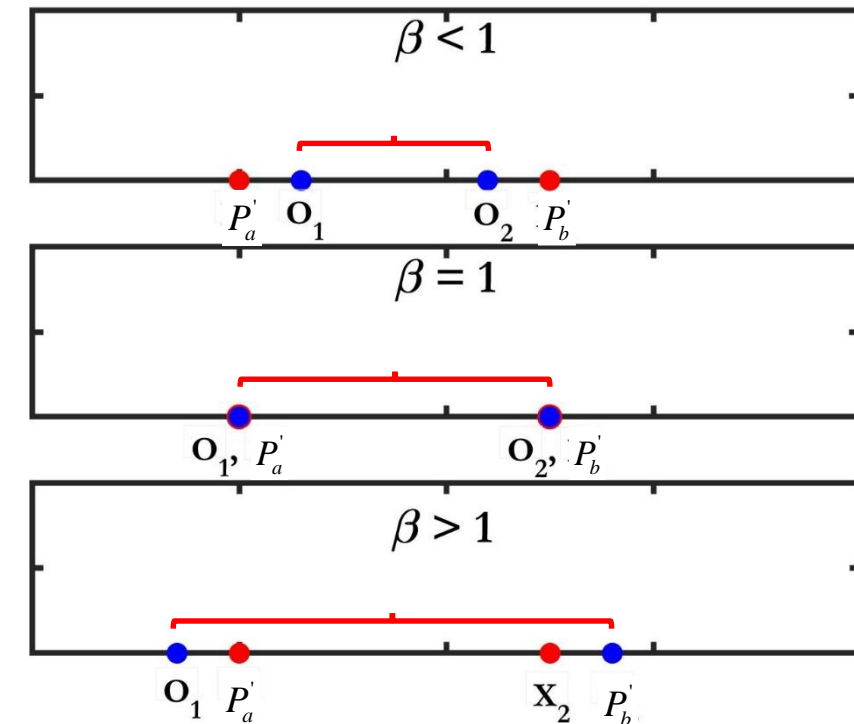
$$O_2 = 0.5[(1-1.4) \times 2 + (1+1.4) \times 5] = 5.6$$

Offspring are far apart

$$O_a = 0.5[(1+\beta)P'_a + (1-\beta)P'_b]$$

$$O_b = 0.5[(1-\beta)P'_a + (1+\beta)P'_b]$$

$$O_a - O_b = \beta(P'_a - P'_b)$$



Polynomial Mutation

➤ Mutation is performed with low probability.

➤ Compute δ as

$$\delta = \begin{cases} (2r)^{1/(\eta_m+1)} - 1 & \text{if } r < 0.5 \\ 1 - [2(1-r)]^{1/(\eta_m+1)} & \text{if } r \geq 0.5 \end{cases}$$

η_m is distribution index

➤ Generate offspring

$$O = O + (ub - lb)\delta$$

O Offspring solution

ub upper bound

lb lower bound

➤ One offspring is generated from an offspring.

Working of genetic algorithm: Sphere function

Consider $\min f(x) = \sum_{i=1}^4 x_i^2; \quad 0 \leq x_i \leq 10, \quad i = 1, 2, 3, 4$

Decision variables: x_1, x_2, x_3 and x_4

$$f(x) = x_1^2 + x_2^2 + x_3^2 + x_4^2$$

- Step 1: Fix the population size, crossover probability, mutation probability, maximum iterations, distribution index for crossover and mutation

$$N_p = 6, p_c = 0.8, p_m = 0.2, T = 10, \eta_c = 20, \eta_m = 20$$

- Step 2: Generate random solutions within the domain of the decision variables

$$P = \begin{bmatrix} 4 & 0 & 0 & 8 \\ 3 & 1 & 9 & 7 \\ 0 & 3 & 1 & 5 \\ 2 & 1 & 4 & 9 \\ 6 & 2 & 8 & 3 \\ 5 & 8 & 1 & 3 \end{bmatrix}$$

$$f = \begin{bmatrix} 80 \\ 140 \\ 35 \\ 102 \\ 113 \\ 99 \end{bmatrix}$$

Selection: Tournament selection

- Step 3: Select two random candidates for tournament

Let the two candidates be

$$P_3 = [0 \ 3 \ 1 \ 5] \quad f_3 = 35$$

$$P_2 = [3 \ 1 \ 9 \ 7] \quad f_2 = 140$$

$$P = \begin{bmatrix} 4 & 0 & 0 & 8 \\ 3 & 1 & 9 & 7 \\ 0 & 3 & 1 & 5 \\ 2 & 1 & 4 & 9 \\ 6 & 2 & 8 & 3 \\ 5 & 8 & 1 & 3 \end{bmatrix} \quad f = \begin{bmatrix} 80 \\ 140 \\ 35 \\ 102 \\ 113 \\ 99 \end{bmatrix}$$

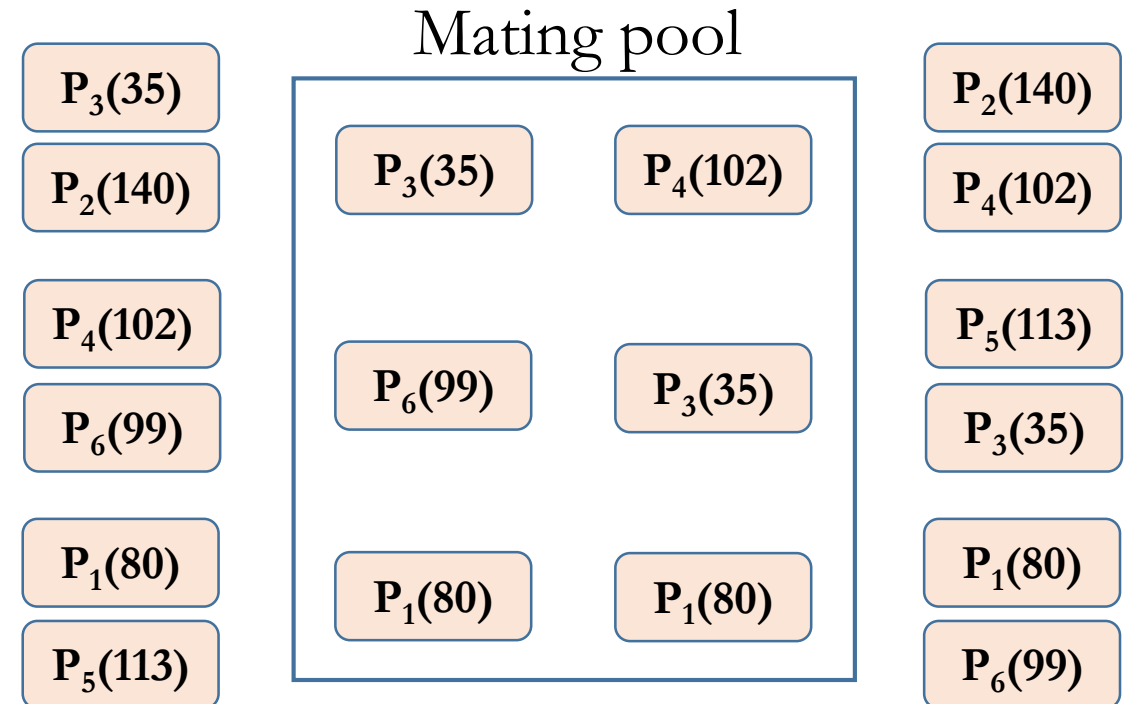
- Step 4: Compare fitness to select winner

$$f_3 < f_2 \rightarrow f_3$$

- Step 5: After six tournaments

Best solution has two copies

Worst solution has zero copies



Procedure for SBX crossover

Input: P , D , p_c , η_c

1. Randomly select a pair of parents (say P'_a and P'_b) from mating pool.
2. Generate a random number (r) between 0 and 1.
3. If $r \geq p_c$, then copy the parent solutions as offspring.
4. If $r < p_c$, generate D random numbers (u) for each variable
5. Determine β of each variables.
6. Generate two offspring (O_a and O_b) using

$$O_a = 0.5 \left[(1 + \beta) P'_a + (1 - \beta) P'_b \right]$$
$$O_b = 0.5 \left[(1 - \beta) P'_a + (1 + \beta) P'_b \right]$$

$$\beta = \begin{cases} (2u)^{1/(\eta_c+1)} & \text{if } u \leq 0.5 \\ \left(\frac{1}{2(1-u)} \right)^{1/(\eta_c+1)} & \text{otherwise} \end{cases}$$

Crossover

- Step 6: Select first pair of parents randomly for crossover

$$p_c = 0.8, \quad \eta_c = 20$$

Let the two parents be

$$\text{Parent}_1 = [0 \quad 3 \quad 1 \quad 5] \quad \text{and} \quad \text{Parent}_3 = [4 \quad 0 \quad 0 \quad 8]$$

Parent =

0	3	1	5
5	8	1	3
4	0	0	8
2	1	4	9
0	3	1	5
4	0	0	8

- Step 7: Select a random number to check if crossover is to be performed

$$\text{Let } r = 0.2$$

$$r < p_c \rightarrow \text{Perform crossover}$$

- Step 8: Select a random number for each variable to perform crossover

$$\text{Let } u = [0.2 \quad 0.6 \quad 0.1 \quad 0.8]$$

$$(u \leq 0.5)$$

$$(u > 0.5)$$

$$\beta = (2 \times 0.2)^{1/(20+1)} = 0.96 \quad \beta = \left(\frac{1}{2 \times (1 - 0.8)} \right)^{1/(20+1)} = 1.04$$

$$\beta = \begin{cases} (2u)^{1/(\eta_c+1)} & \text{if } u \leq 0.5 \\ \left(\frac{1}{2(1-u)} \right)^{1/(\eta_c+1)} & \text{otherwise} \end{cases}$$

Crossover

$$\beta = [0.96 \quad 1.01 \quad 0.93 \quad 1.04]$$

Step 9: Create two offspring

$$\begin{aligned} \text{offspring}_1 &= 0.5 \times \begin{pmatrix} (1 + [0.96 \quad 1.01 \quad 0.93 \quad 1.04]) \times [0 \quad 3 \quad 1 \quad 5] \\ (1 - [0.96 \quad 1.01 \quad 0.93 \quad 1.04]) \times [4 \quad 0 \quad 0 \quad 8] \end{pmatrix} \\ &= [0.08 \quad 3.01 \quad 0.97 \quad 4.94] \end{aligned}$$

$$\begin{aligned} \text{offspring}_2 &= 0.5 \times \begin{pmatrix} (1 - [0.96 \quad 1.01 \quad 0.93 \quad 1.04]) \times [0 \quad 3 \quad 1 \quad 5] \\ (1 + [0.96 \quad 1.01 \quad 0.93 \quad 1.04]) \times [4 \quad 0 \quad 0 \quad 8] \end{pmatrix} \\ &= [3.92 \quad -0.02 \quad 0.03 \quad 8.06] \end{aligned}$$

$$O_1 = 0.5[(1 + \beta)P'_a + (1 - \beta)P'_b]$$

$$\text{Parent} = P' = \begin{bmatrix} 0 & 3 & 1 & 5 \\ 5 & 8 & 1 & 3 \\ 4 & 0 & 0 & 8 \\ 2 & 1 & 4 & 9 \\ 0 & 3 & 1 & 5 \\ 4 & 0 & 0 & 8 \end{bmatrix}$$

$$O_2 = 0.5[(1 - \beta)P'_a + (1 + \beta)P'_b]$$

Step 10: Check for bound violation

$$\text{offspring}_2 = [3.92 \quad -0.02 \quad 0.03 \quad 8.06] \rightarrow [3.92 \quad 0 \quad 0.03 \quad 8.06]$$

No bound violation in offspring₁

$$0 \leq x_i \leq 10$$

Crossover

- Step 11: Select second pair of parents randomly for crossover

Let the two parents be

$$\text{Parent}_2 = [5 \ 8 \ 1 \ 3] \quad \text{and} \quad \text{Parent}_6 = [4 \ 0 \ 0 \ 8]$$

$$p_c = 0.8, \quad \eta_c = 20$$

$$\text{Parent} = \begin{bmatrix} 0 & 3 & 1 & 5 \\ 5 & 8 & 1 & 3 \\ 4 & 0 & 0 & 8 \\ 2 & 1 & 4 & 9 \\ 0 & 3 & 1 & 5 \\ 4 & 0 & 0 & 8 \end{bmatrix}$$

- Step 12: Select a random number to check if crossover is to be performed

Let $r = 0.9$

$$r > p_c \rightarrow \text{no crossover}$$

- Step 13: Copy the parents as offspring solutions

$$\text{offspring}_3 = \text{Parent}_2 = [5 \ 8 \ 1 \ 3]$$

$$\text{offspring}_4 = \text{Parent}_6 = [4 \ 0 \ 0 \ 8]$$

$$\text{offspring} = O = \begin{bmatrix} 0.08 & 3.01 & 0.97 & 4.94 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 8 & 1 & 3 \\ 4 & 0 & 0 & 8 \end{bmatrix}$$

Crossover

- Step 14: Select third pair of parents randomly for crossover

$$p_c = 0.8, \eta_c = 20$$

Let the two parents be $\text{Parent}_4 = [2 \ 1 \ 4 \ 9]$ $\text{Parent}_5 = [0 \ 3 \ 1 \ 5]$

Parent =

0	3	1	5
5	8	1	3
4	0	0	8
2	1	4	9
0	3	1	5
4	0	0	8

- Step 15: Select a random number to check if crossover is to be performed

Let $r = 0.4$

$r < p_c \rightarrow$ Perform crossover

- Step 16: Select a random number for each variable to perform crossover

Let $u = [0.3 \ 0.1 \ 0.8 \ 0.6]$

$$\beta = \begin{bmatrix} \overset{(u \leq 0.5)}{0.98} & \overset{(u \leq 0.5)}{0.93} & \overset{(u > 0.5)}{1.04} & \overset{(u > 0.5)}{1.01} \end{bmatrix}$$

$$\beta = \begin{cases} (2u)^{1/(\eta_c+1)} & \text{if } u \leq 0.5 \\ \left(\frac{1}{2(1-u)} \right)^{1/(\eta_c+1)} & \text{otherwise} \end{cases}$$

- Step 17: Create two offspring

offspring₅ = [1.98 1.07 4.06 9.02]

offspring₆ = [0.02 2.93 0.94 4.98]

$$O_1 = 0.5 \left[(1 + \beta) P'_a + (1 - \beta) P'_b \right]$$

$$O_2 = 0.5 \left[(1 - \beta) P'_a + (1 + \beta) P'_b \right]$$

Procedure for polynomial mutation

Input: P , D , p_m , η_m

1. Generate a random number (u) between 0 and 1
2. If $u \geq p_m$, then no change in the offspring
3. If $u < p_m$, generate D random numbers (r) corresponding to each variable
4. Determine δ of each variables

$$\delta = \begin{cases} (2r)^{1/(\eta_m+1)} - 1 & \text{if } r < 0.5 \\ 1 - [2(1-r)]^{1/(\eta_m+1)} & \text{if } r \geq 0.5 \end{cases}$$

5. Modify offspring using $O = O + (x^u - x^l) \times \delta$

Mutation

- Step 18: Select first offspring for mutation

$$\text{offspring}_1 = [0.08 \quad 3.01 \quad 0.97 \quad 4.94]$$

- Step 19: Select a random number to check if mutation happens

$$\text{Let } r = 0.1$$

$$r < p_m \rightarrow \text{Perform mutation}$$

- Step 20: Select a random number for each variable to perform mutation

$$\text{Let } r = [0.6 \quad 0.1 \quad 0.2 \quad 0.8]$$

$$(r \geq 0.5)$$

$$\begin{aligned} \delta &= 1 - (2 \times (1 - 0.6))^{1/(20+1)} \\ &= 0.01 \end{aligned}$$

$$(r < 0.5)$$

$$\begin{aligned} \delta &= (2 \times 0.2)^{1/(20+1)} - 1 \\ &= -0.04 \end{aligned}$$

$$\delta = [0.01 \quad -0.07 \quad -0.04 \quad 0.04]$$

$$O = \begin{bmatrix} 0.08 & 3.01 & 0.97 & 4.94 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 8 & 1 & 3 \\ 4 & 0 & 0 & 8 \\ 1.98 & 1.07 & 4.06 & 9.02 \\ 0.02 & 2.93 & 0.94 & 4.98 \end{bmatrix}$$

$$p_m = 0.2, \eta_m = 20$$

$$\delta = \begin{cases} (2r)^{1/(\eta_m+1)} - 1 & \text{if } r < 0.5 \\ 1 - [2(1-r)]^{1/(\eta_m+1)} & \text{if } r \geq 0.5 \end{cases}$$

Mutation

$$\delta = [0.01 \quad -0.07 \quad -0.04 \quad 0.04]$$

$$0 \leq x_i \leq 10, \quad i = 1, 2, 3, 4$$

- Step 21: Generate new offspring

$$\text{offspring}_1 = [0.08 \quad 3.01 \quad 0.97 \quad 4.94] +$$

$$([10 \quad 10 \quad 10 \quad 10] - [0 \quad 0 \quad 0 \quad 0]) \times [0.01 \quad -0.07 \quad -0.04 \quad 0.04]$$

$$= [0.18 \quad 2.31 \quad 0.57 \quad 5.34]$$

$$O_1 = O_1 + (x^u - x^l) \times \delta$$

- Step 22: Select second offspring for crossover

$$\text{offspring}_2 = [3.92 \quad 0 \quad 0.03 \quad 8.06]$$

$$O = \begin{bmatrix} 0.18 & 2.31 & 0.57 & 5.34 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 8 & 1 & 3 \\ 4 & 0 & 0 & 8 \\ 1.98 & 1.07 & 4.06 & 9.02 \\ 0.02 & 2.93 & 0.94 & 4.98 \end{bmatrix}$$

- Step 23: Select a random number to check if mutation happens

$$\text{Let } r = 0.2$$

$$r = p_m \rightarrow \text{No mutation}$$

- Step 24: No change in the offspring

Mutation

- Step 25: Perform mutation for rest of the offspring

$$p_m = 0.2, \eta_m = 20$$

Offspring	Random number for mutation probability	Random number for mutation	New offspring
[5 8 1 3]	0.1	[0.5 0.1 0.6 0.3]	[5 7.3 1.1 2.8]
[4 0 0 8]	0.6	No mutation ($r > p_m$)	[4 0 0 8]
[1.98 1.07 4.06 9.02]	0.3	No mutation ($r > p_m$)	[1.98 1.07 4.06 9.02]
[0.02 2.93 0.94 4.98]	0.8	No mutation ($r > p_m$)	[0.02 2.93 0.94 4.98]

- Step 26: Evaluate fitness of all the offspring solutions

$$O = \begin{bmatrix} 0.18 & 2.31 & 0.57 & 5.34 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 7.3 & 1.1 & 2.8 \\ 4 & 0 & 0 & 8 \\ 1.98 & 1.07 & 4.06 & 9.02 \\ 0.02 & 2.93 & 0.94 & 4.98 \end{bmatrix}$$

$$f_o = \begin{bmatrix} 34.21 \\ 80.33 \\ 87.34 \\ 80 \\ 102.91 \\ 34.27 \end{bmatrix}$$

Survival

- Step 27: Combine all the solutions and select best N_p ($N_p = 6$) solutions

$$P = \begin{bmatrix} 4 & 0 & 0 & 8 \\ 3 & 1 & 9 & 7 \\ 0 & 3 & 1 & 5 \\ 2 & 1 & 4 & 9 \\ 6 & 2 & 8 & 3 \\ 5 & 8 & 1 & 3 \end{bmatrix}$$

$$f = \begin{bmatrix} 80 \\ 140 \\ 35 \\ 102 \\ 113 \\ 99 \end{bmatrix}$$

$$P_{combined} = \begin{bmatrix} 0.18 & 2.31 & 0.57 & 5.34 \\ 0.02 & 2.93 & 0.94 & 4.98 \\ 0 & 3 & 1 & 5 \\ 4 & 0 & 0 & 8 \\ 4 & 0 & 0 & 8 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 7.3 & 1.1 & 2.8 \\ 5 & 8 & 1 & 3 \\ 2 & 1 & 4 & 9 \\ 1.98 & 1.07 & 4.06 & 9.02 \\ 6 & 2 & 8 & 3 \\ 3 & 1 & 9 & 7 \end{bmatrix}$$

$$f_{combined} = \begin{bmatrix} 34.21 \\ 34.27 \\ 35 \\ 80 \\ 80 \\ 80.33 \\ 87.34 \\ 99 \\ 102 \\ 102.91 \\ 113 \\ 140 \end{bmatrix}$$

$$O = \begin{bmatrix} 0.18 & 2.31 & 0.57 & 5.34 \\ 3.92 & 0 & 0.03 & 8.06 \\ 5 & 7.3 & 1.1 & 2.8 \\ 4 & 0 & 0 & 8 \\ 1.98 & 1.07 & 4.06 & 9.02 \\ 0.02 & 2.93 & 0.94 & 4.98 \end{bmatrix}$$

$$f_o = \begin{bmatrix} 34.21 \\ 80.33 \\ 87.34 \\ 80 \\ 102.91 \\ 34.27 \end{bmatrix}$$

Satisfaction of termination criterion

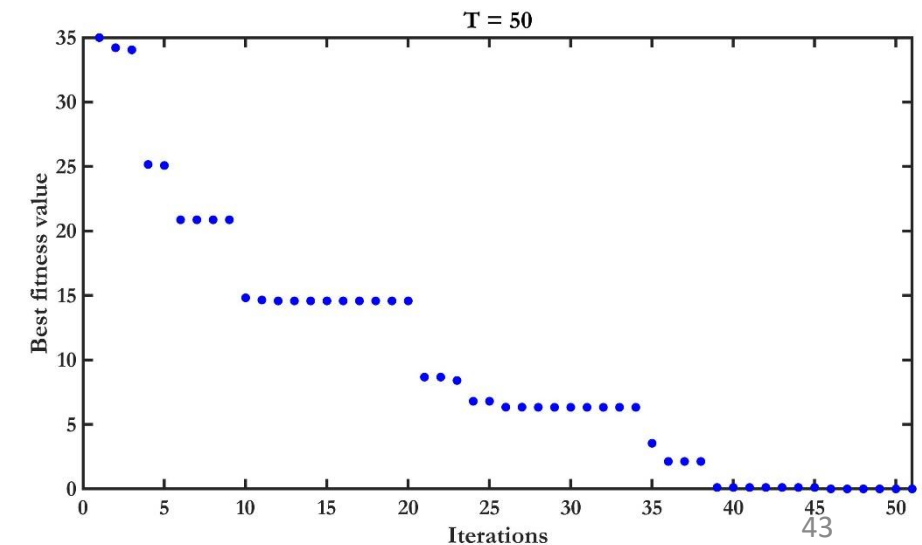
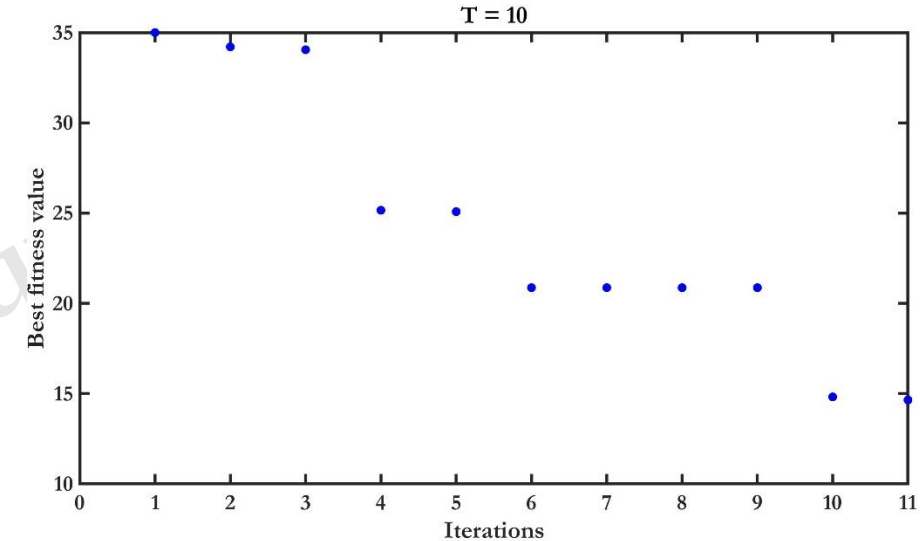
$$\min f(x) = \sum_{i=1}^4 x_i^2; \quad 0 \leq x_i \leq 10, \quad i = 1, 2, 3, 4$$

After completion of 10 iterations

$$P = \begin{bmatrix} 0.3 & 1.23 & 0.61 & 3.56 \\ 0.3 & 1.24 & 0.61 & 3.58 \\ 0.3 & 1.24 & 0.61 & 3.59 \\ 0 & 1.74 & 0.61 & 4.18 \\ 0 & 1.74 & 0.61 & 4.18 \\ 0 & 1.74 & 0.61 & 4.18 \end{bmatrix} \quad f = \begin{bmatrix} 14.65 \\ 14.82 \\ 14.89 \\ 20.87 \\ 20.87 \\ 20.87 \end{bmatrix}$$

The minimum value of the function is 0

The optima is achieved by increasing T



Pseudocode

Input: Fitness function, lb, ub, N_p , T , p_c , p_m , η_c , η_m , k

```

1. Initialize a random population ( P )
2. Evaluate fitness ( f ) of P ← FE =  $N_p$ 
   for t = 1 to T
     Perform tournament selection of tournament size, k
     for i = 1 to  $N_p/2$ 
       Randomly choose two parents
       if  $r < p_c$ 
         Generate two offspring using SBX-crossover
         Bound the offspring and store them in offspring population
       else
         Copy the selected parents and their fitness to offspring population
       end
     end
   end
   for i = 1 to  $N_p$ 
     if  $r < p_m$ 
       Perform polynomial mutation of  $i^{\text{th}}$  solution in offspring population
       Bound the mutated offspring
     else
       No change in  $i^{\text{th}}$  solution of offspring
     end
   end
   end
   Evaluate the fitness ← Max #FE =  $N_p$ 
   Combine population and offspring population to perform ( $\mu + \lambda$ )
end
  
```

In one iteration, **max** #FE = N_p
 For T iterations, **max** #FE = $N_p + N_p T$

Generation

$$\beta = \begin{cases} (2u)^{1/(\eta_c+1)} & \text{if } u \leq 0.5 \\ \left(\frac{1}{2(1-u)} \right)^{1/(\eta_c+1)} & \text{otherwise} \end{cases}$$

$$O_1 = 0.5[(1+\beta)X_1 + (1-\beta)X_2]$$

$$O_2 = 0.5[(1-\beta)X_1 + (1+\beta)X_2]$$

$$\delta = \begin{cases} (2r)^{1/(\eta_m+1)} & \text{if } r < 0.5 \\ 1 - [2(1-r)]^{1/(\eta_m+1)} & \text{if } r \geq 0.5 \end{cases}$$

$$y = O + (ub - lb)\delta$$

Survival of fittest

Pseudocode

Input: Fitness function, lb, ub, N_p , T , n , p_c , p_m , k

1. Initialize a random population (P) of **binary string (size: $N_p \times nD$)**

2. Evaluate fitness (f) of P

for $t = 1$ to T

 Perform **tournament selection** of tournament size, k

for $i = 1$ to $N_p/2$

 Randomly choose two parents

if $r < p_c$

 Select the crossover site

 Generate two offspring using **single-point crossover**

else

 Copy the selected parents as offspring

end

end

for $i = 1$ to N_p

 Generate **D random numbers between 0 and 1**

 Perform **bit-wise mutation** of i^{th} offspring

end

 Evaluate the fitness of offspring

 Combine population and offspring population to perform **($\mu + \lambda$)**

end

BGA

Input: Fitness function, lb, ub, N_p , T , p_c , p_m , η_c , η_m , k

1. Initialize a random population (P)

2. Evaluate fitness (f) of P

for $t = 1$ to T

 Perform **tournament selection** of tournament size, k

for $i = 1$ to $N_p/2$

 Randomly choose two parents

if $r < p_c$

 Generate two offspring using **SBX-crossover**

 Bound the offspring and store them in offspring population

else

 Copy the selected parents and their fitness to offspring population

end

end

for $i = 1$ to N_p

if $r < p_m$

 Perform **polynomial mutation** of i^{th} offspring

 Bound the mutated offspring

else

 No change in i^{th} offspring

end

end

 Evaluate the fitness of offspring

 Combine population and offspring population to perform **($\mu + \lambda$)**

end

Real variables are encoded to binary in BGA
Encoding is not required in RGA

In RGA mutation probability is checked only once
In BGA, it is checked for each variable

RGA

Comparison of techniques

	TLBO	PSO	DE	ABC	GA
Phases	Teacher, Learner	No phases (Position and velocity update)	No phases (Mutation and crossover)	Employee, Onlooker and Scout	No phases (Mutation and crossover)
Convergence	Monotonic	Monotonic (with g_{best} & p_{best})	Monotonic	Monotonic (with globalized memory)	Monotonic
Parameters	Population size, termination criteria	Population size, termination criteria, w , c_1 and c_2	Population size, termination criteria, P_c , F	Population size, termination criteria, limit	Population size, termination criteria, p_c , p_m , and other parameters of variation operators (η_c and η_m)
Generation of solution	Using other solutions, mean and best solution	Using velocity vector, p_{best} and g_{best}	Using other solutions	Using other solutions	Using other solutions
Best solution	Part of population	Need not be part of population	Part of population	Need not be part of population	Part of population
Fitness function	Objective function	Objective function	Objective function	Inversely related to objective function	Objective function
Population update	Twice	Once	Once	Twice or thrice (Scout phase)	Once
Selection	Greedy	Always accept new solution into the population (μ, λ)	Greedy	Greedy and (μ, λ) (in scout phase)	Survival of the fittest ($\mu+\lambda$)
#FE	$N_p + 2N_pT$	$N_p + N_pT$	$N_p + N_pT$	Max #FE = $N_p + 2N_pT + T$	Max #FE = $N_p + N_pT$

Further reading

- Genetic Algorithms in search, optimization and machine learning, Addison-Wesley publishing company, 1989
- An introduction to genetic algorithms, **Sadhana**, 24, Parts 4 & 5, 293-315, 1999
- Multi-Objective Optimization Using Evolutionary Algorithms, John Wiley & Sons Inc., 2001
- Simulated Binary Crossover for Continuous Search Space, **Complex systems**, 9(2), 115-148, 1995
- A fast and elitist multiobjective genetic algorithm: NSGA-II, **IEEE Transactions on Evolutionary Computation**, 6(2), 182-197, 2002

Thank You !!!