



# From Boot to Root

Richard Weinberger - sigma star gmbh 2024-11-11



### Richard Weinberger

- Co-founder of sigma star gmbh
- Linux kernel developer and maintainer
- Strong focus on Linux kernel, lowlevel components, virtualization, security, code audits

#### sigma star gmbh

- Software Development & Security
   Consulting
- Main areas: Embedded Systems, Linux Kernel & Security
- Contributions to Linux Kernel and other OSS projects

## Bootloader 101

- > Started by boot ROM (BIOS/UEFI on x86)
- On embedded systems, boot ROM is part of the SoC
- › Bootloader loads operating system (the kernel)
- > ±fancy UI and configuration

# Get root On a Typical Linux System

- > Edit bootloader config
- > Add init=/bin/sh to kernel command line
- > Solution: Lockdown bootloader (plus config!)

### **Bootloader Lockdown**

- > No way to change boot config
- No shell
- > Input only from trusted sources or fully authenticated input
- > Sounds easier than it is

#### Chain of Trust

- › Boot ROM authenticates bootloader signature
- › Bootloader authenticates OS kernel signagure
- > OS kernel authenticates userspace
- Common on security focused systems (UEFI Secure Boot, etc.)
- > Hello CRA (Cyber Resilience Act), hello NIS-2 (Network and Information Security)

## The Weakest Link: The Bootloader

- > Break the bootloader and control the rest of the system:
  - > Start our own code
  - > Extract secrets (key material)
  - > Impersonate the device
  - Basically become root

## **U-Boot and Barebox**

- > Extremely common bootloaders for embdded Linux
- > Load and authenticate files from a filesystem
- > I started auditing their critical code paths

## **Critical Code Paths**

- > Config file parsing (AKA boot environment)
- > Parsing other state, e.g. boot counter on EEPROM
- > Loading boot files, kernel, device tree, ...
- Most inputs are authenticated
- > The elephant in the room: filesystems

# Filesystems at Boot Stage

- > Data on the filesystems is authenticated
- > The filesystem itself is *not*
- > Filesystems drivers in bootloades:
  - > Good enough to read a file
  - Not more
- > Filesystems can get manipulated by an attacker

- > Integer overflow in ext4 symlink code
- Results in attacker driven out of bounds write
- > Unauthenticated attacker can trigger it
- > Both U-Boot and Barebox affected

- > Integer overflow in squashfs symlink code, like vulnerability #1.
- > Results in attacker driven out of bounds write
- > Unauthenticated attacker can trigger it
- > Both U-Boot and Barebox affected
- > Although they have different squashfs implementations

- > Stack overflow in squashfs symlink code
- Code follows symlinks recursively
- Results in attacker driven stack smashing
- > Unauthenticated attacker can trigger it
- Only U-Boot affected

```
int sqfs_size(const char *filename, loff_t
   *size)
  switch (get_unaligned_le16(&base-> _/
      inode_type)) {
  case SQFS_LSYMLINK_TYPE:
  symlink = (struct squashfs_symlink_inode
       *)ipos;
  filename):
  ret = sqfs_size(resolved, size):
  free(resolved);
  break;
```

- Multiple integer overflows in memory allocator
- You ask for N bytes but get much less
- Can get triggered by most filesystem drivers
- > Unauthenticated attacker can trigger it
- > Both U-Boot and Barebox affected
- They use Doug Lea's Malloc, but broke it 25 years ago
- Bonus: Another integer overflow in their sbrk()
- Bonus #2: ptrdiff\_t too small on x86\_64, more overflows

```
#define request2size(req) \
 (((long)((req) + (SIZE_SZ +
     MALLOC ALIGN MASK)) < \
  (long)(MINSIZE + MALLOC_ALIGN_MASK)) ?
      MINSIZE : \
   (((req) + (SIZE_SZ + MALLOC_ALIGN_MASK))
//

         & ~(MALLOC_ALIGN_MASK)))
Void_t* mALLOc_impl(size_t bytes)
  if ((long)bytes < 0) return NULL:</pre>
  nb = request2size(bytes); /* padded
      request size: */
```

## Outcome



- At least four beefy vulnerabilities that allow full compromise
- Sent bug reports and patches for all vulnerabilities, all merged
- Improved (fixed) U-Boot's ASAN integration

# Just Update the Damn Bootloader?!

- Think of downgrade attacks!
- > Attacker can always install the old vulnerable bootloader
- Mitigations:
  - Have a revoke list (hard!)
  - > Have an authenticated version counter in hardware



> CVF-2024-2312



- CVE-2024-2312
- > CVE-2024-1048



- > CVF-2024-2312
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- > CVE-2023-4693



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# Discussion: How to Improve the Situation?

- > Reusing Linux implementations is hard
  - Needs a Linux VFS
  - Offer more than needed
  - Code size matters
- > Toy implementations are always broken
- > kexec is problematic on embedded systems
- > Idea: Provide sane libfs{ext4, squashfs, ...} for bare metal
  - How to sync with Linux?
  - > Funding?
  - > What license?

## FIN



# Thank you!

Questions, Comments?

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