



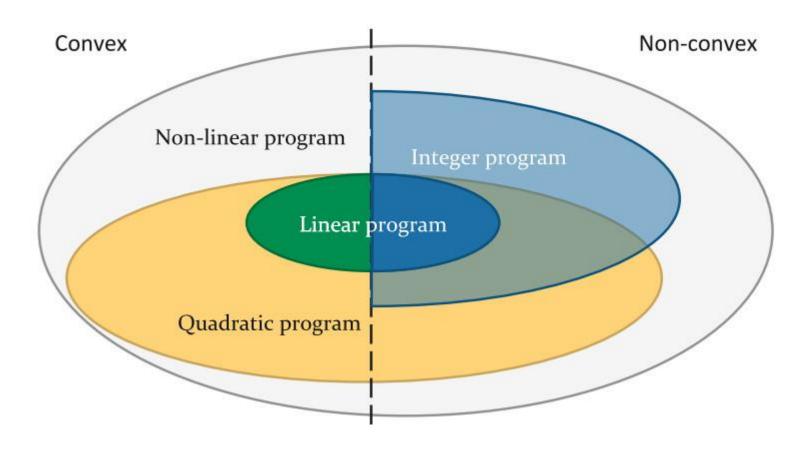
Energy systems modelling

Tutorial 11

Iegor Riepin



Introduction to complementarity problems



Source: [1]

Introduction to MCP formulations



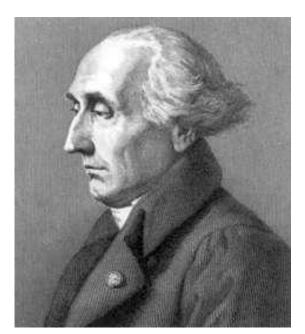
- MCP (mixed complementarity programming) is a common modelling approach to describe various energy markets around the world.
- Complementarity models generalize linear programs (LP), quadratic programs (QP) and (convex) nonlinear programs (NLPs)
- Complementarity problems are appropriate for modelling the regulated/deregulated, perfect/imperfect competition that characterizes today's energy markets



Method of Lagrange multipliers: problem definition

In mathematical optimization, the method of Lagrange multipliers <u>is a strategy for finding the local maxima and minima of a function subject to equality constraints:</u>

$$s.t.g(x,y) = c$$

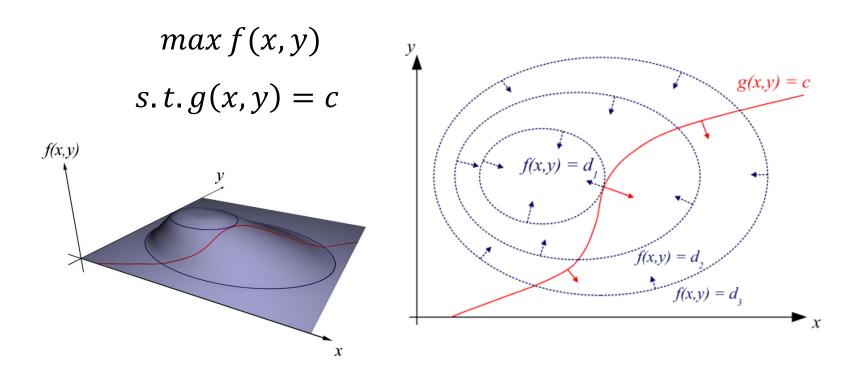


Joseph-Louis Lagrange



Method of Lagrange multipliers: problem definition

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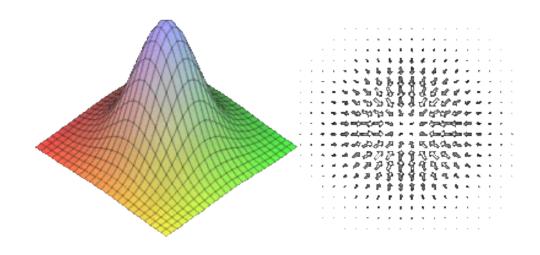
Method of Lagrange multipliers: gradient

The gradient is a generalization of the usual concept of derivative of a function in one dimension to a function in several dimensions.

➤ Gradient points in the <u>direction of the greatest rate of increase</u> of the function and its <u>magnitude is the slope of the graph</u> in that direction

$$\nabla f = \frac{\partial f}{\partial x_1} \mathbf{e}_1 + \dots + \frac{\partial f}{\partial x_n} \mathbf{e}_n$$

where the ei are the orthogonal unit vectors pointing in the coordinate directions.



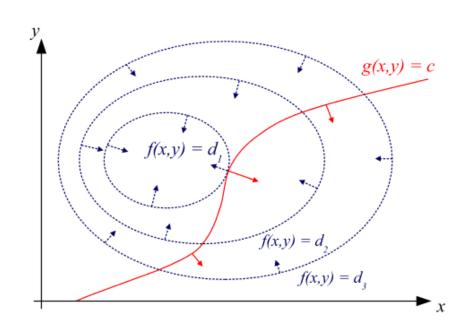
Method of Lagrange multipliers



The method of Lagrange multipliers <u>is a strategy for finding the local maxima</u> and minima of a function subject to equality constraints:

$$s.t.g(x,y) = c$$

Important observation: if 2 curves are tangent at the same point -> they have the same slope



Method of Lagrange multipliers



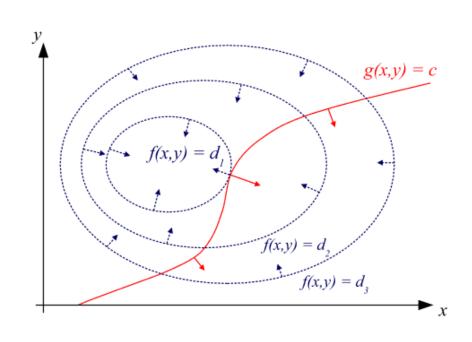
The method of Lagrange multipliers <u>is a strategy for finding the local maxima</u> and minima of a function subject to equality constraints:

$$s.t.g(x,y) = c$$

$$\nabla f(x,y) = \begin{bmatrix} \frac{\partial f}{\partial x} \\ \frac{\partial f}{\partial y} \end{bmatrix}$$

$$\nabla g(x, y) = \begin{vmatrix} \frac{\partial g}{\partial x} \\ \frac{\partial g}{\partial y} \end{vmatrix}$$

Important observation: if 2 curves are tangent at the same point -> they have the same slope







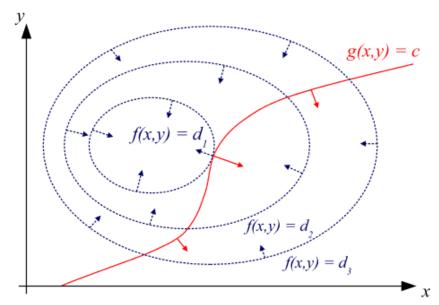
In mathematical optimization, the method of Lagrange multipliers <u>is a strategy for finding the local maxima and minima of a function subject to equality constraints:</u>

$$s.t.g(x,y) = c$$

$$\nabla f(x,y) = \lambda \cdot \nabla g(x,y)$$

This condition ensures that isolines (contour curves) are tangent

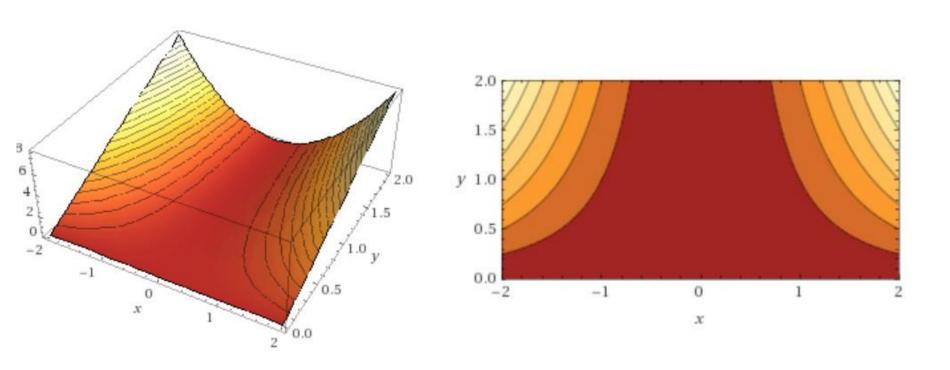
Important observation: if 2 curves are tangent at the same point -> they have the same slope -> their gradient vectors are parallel





Method of Lagrange multipliers: an example

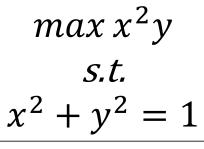
$max x^2 y$

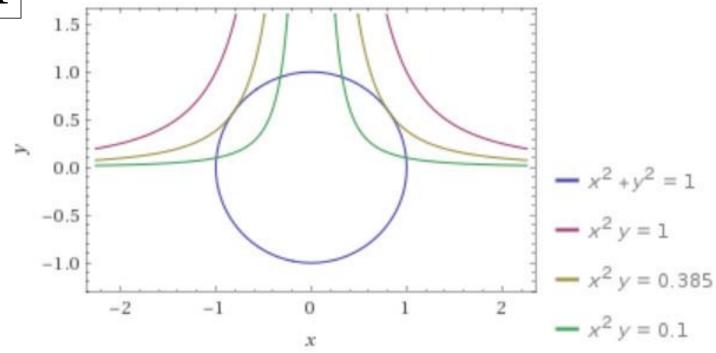


Own visualisations



Method of Lagrange multipliers: an example





Own visualisations



Method of Lagrange multipliers: economics

✓ In economics the optimal profit to a player is calculated subject to a constrained space of actions, where a Lagrange multiplier is the change in the <u>optimal value of the objective function (profit)</u> due to the <u>relaxation</u> of a given constraint

$$L(x, y, \lambda) = f(x, y) - \lambda \cdot g(x, y)$$



in such a context λ is the <u>marginal cost</u> of the constraint, and is referred as the <u>shadow price</u>





- ✓ The Karush–Kuhn–Tucker (KKT) conditions <u>are first order necessary</u> <u>conditions for a solution in nonlinear programming to be optimal</u>, provided that some regularity conditions are satisfied.
- ✓ Allowing inequality constraints, the KKT approach applied to nonlinear programming generalizes the method of Lagrange multipliers, which allows only equality constraints.

Karush-Kuhn-Tucker conditions



Let us consider the problem:

$$\min_{x} F(x)
s.t. \quad g_{i}(x) \leq 0 \quad (\lambda_{i} \geq 0) \quad \forall i = 1, ... n
h_{j}(x) = 0 \quad (\mu_{j} free) \quad \forall j = 1, ... m$$
[1.1]

> For this problem, the KKT conditions are:

$$\nabla f(x) + \sum_{i=1}^{n} \lambda_i \nabla g_i(x) + \sum_{j=1}^{n} \mu_i \nabla h_j(x) \le 0 \perp x \ge 0$$

$$0 \ge g_i(x) \perp \lambda_i \ge 0 \qquad \forall i = 1, \dots n$$

$$0 = h_i(x) \quad \mu_i \text{ free} \qquad \forall j = 1, \dots m$$
[1.4]
$$[1.4]$$

The solution stationarity is ensured by the equation [1.4]. Equations [1.5] and [1.6] ensure complementarity and feasibility of a solution



Example: perfectly competitive market two producers & single demand

Producer's objective function is to maximize profits P by selling quantity of product q at price p bearing costs of production C(q):

$$\max_{q_i \ge 0} P_i = q_i p - C_i(q_i) \ \forall i$$
 [1.7*]

$$s.t.q_i \le Q_i \ \forall i$$

where:

q - quantity of product sold by producer i at price p $C_i(q_i)$ - production costs

^{*}if price is an exogenous variable from producers perspective;



Example: perfectly competitive market two producers & single demand

Optimization problem:

$$\max_{q_i \ge 0} P_i = q_i p - C_i(q_i) \ \forall i$$
$$s. t. q_i \le Q_i \ \forall i$$

The KKTs for producer i are as follows:

$$0 \le q_i \perp p - C'_i(q_i) + \lambda_i \le 0$$

$$0 \le \lambda_i \perp (q_i - Q_i) \le 0$$
[1.9]

Symbol \perp states orthogonality

$$0 < a \perp b > 0$$

Implies:

$$a \ge 0, b \ge 0$$
 and $ab = 0$

Eq. [1.9] is a short way to express the following:

$$0 \le q_i$$

$$p - C_i'(q_i) + \lambda_i \le 0$$

$$q_i(p - C_i'(q_i) + \lambda_i) = 0$$



Hometask - simple LP / MCP problem in GAMS /I know it might be challenging.../

Solve the following problem using LP and MCP formulations:

$$\max_{x,y} 4x + 5y$$

s. t.

$$x + y \leq 24$$

$$\boldsymbol{x} \in R_{+}, \boldsymbol{y} \in R_{+}$$

Compare the marginal value of the constraint to the values of Lagrange multipliers in MCP model

$$\nabla f(x,y) = \lambda \cdot \nabla g(x,y)$$

 $\nabla f(x,y) = \lambda \cdot \nabla g(x,y)$ in Such context A is the <u>integrate cost</u> of the constraint, and is referred as the <u>shadow</u> price in such context λ is the <u>marginal cost</u> of the



See you next class!

Appendix A: Duality concept



Some (classes of) optimization problems have a "twin" problem,

⇒ This is called the "dual problem"

Illustration using a simple linear program:

Mathematical formulation

$$\min_{x \in \mathbb{R}} a^T x$$

s.t.
$$Ax \le b$$
 (y)

$$\max_{y \in \mathbb{R}_+} b^T y$$

s.t.
$$A^T y = a$$
 (x)

Example of interpretation

Minimize cost of supplying electricity subject to engineering and power flow constraints

Maximize pay-off such that the dual constraints are satisfied

⇒ If the optimal objective values are identical, we call it strong duality





The "classical" formulation of Karush-Kuhn-Tucker conditions:

$$0 = \nabla_{u} f(x^{*}) + \sum_{i} \lambda_{i}^{*} \nabla_{u} g_{i}(x^{*}) + \sum_{j} \mu_{j}^{*} \nabla_{u} h_{j}(x^{*}) , \quad x_{u}^{*} \text{ (free)}$$

$$0 \ge g_{i}(x^{*}) \perp \lambda_{i}^{*} \ge 0$$

$$0 = h_{j}(x^{*}) , \quad \mu_{j}^{*} \text{ (free)}$$

One alternative formulation (of many):

$$0 \leq \nabla_{u} f(x^{*}) + \sum_{i} \lambda_{i}^{*} \nabla_{u} g_{i}(x^{*}) + \sum_{j} \mu_{j}^{*} \nabla_{u} h_{j}(x^{*}) \perp x_{u}^{*} \geq 0$$

$$0 \leq -g_{i}(x^{*}) \perp \lambda_{i}^{*} \geq 0$$

$$0 = h_{j}(x^{*}) , \mu_{j}^{*} \text{ (free)}$$



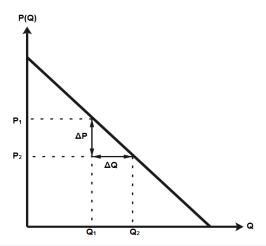
Appendix C: inverse demand function

The affine inverse demand function is commonly expressed in the following way:

$$P(Q) = a + b \cdot Q$$

where P(Q) represents the price of a good as a function of quantity demanded (Q). The constant b represents a slope of the function and the constant a is an intersection point with the vertical axis.

Inverse demand function is plotted on a coordinate system with the price on the vertical axis and quantity on the horizontal axis:







- 1. Huppmann, D., 2014. One and Two-level Energy Market Equilibrium Modelling
- 2. Bertsekas, D.P., 1999. Nonlinear programming.