

Team Note of PS akgwi

Jeounghui Nah, Joowon Oh, Seongseo Lee

Compiled on October 19, 2023

Contents

1	DataStructure	1
1.1	Bipartite Union Find	1
1.2	Erasable Priority Queue	1
1.3	Convex Hull Trick	1
1.4	Persistent Segment Tree	2
1.5	Kinetic Segment Tree	2
1.6	Splay Tree, Link-Cut Tree	2
2	Geometry	3
2.1	Triangles	3
2.2	Rotating Calipers	3
2.3	Point in Convex Polygon	3
2.4	Segment Distance	4
2.5	Tangent Series	4
2.6	Intersect Series	4
2.7	Polygon Cut, Center, Union	5
2.8	Polygon Raycast	5
2.9	Shamos-Hoey	5
2.10	Half Plane Intersection	6
2.11	K-D Tree	6
2.12	Dual Graph	7
2.13	Bulldozer Trick (Rotating Sweep Line)	7
2.14	Smallest Enclosing Circle	7
2.15	Voronoi Diagram	7
3	Graph	8
3.1	Euler Tour	8
3.2	2-SAT	8
3.3	Horn SAT	9
3.4	BCC	9
3.5	Prufer Sequence	9
3.6	Maximum Clique	9
3.7	Tree Isomorphism	9
3.8	Complement Spanning Forest	10
3.9	Bipartite Matching, Konig, Dilworth	10
3.10	Push Relabel	10
3.11	LR Flow	11
3.12	Hungarian Method	11
3.13	Count/Find 3/4 Cycle	11
3.14	$O(V^3)$ Global Min Cut	11
3.15	Gomory-Hu Tree	12
3.16	Rectlinear MST	12
3.17	$O((V + E) \log V)$ Dominator Tree	12
3.18	$O(N^2)$ Stable Marriage Problem	12

3.19	$O(VE)$ Vizing Theorem	12
3.20	$O(E \log V)$ Directed MST	13
3.21	$O(E \log V + K \log K)$ K Shortest Path	13
3.22	Chordal Graph, Tree Decomposition	14
3.23	$O(V^3)$ General Matching	14
3.24	$O(V^3)$ Weighted General Matching	15
4	Math	
4.1	Extend GCD, CRT, Combination	
4.2	Diophantine	
4.3	Partition Number	
4.4	FloorSum	
4.5	XOR Basis(XOR Maximization)	
4.6	Stern Brocot Tree	
4.7	Gauss Jordan Elimination	
4.8	Berlekamp + Kitamasa	
4.9	Miller Rabin + Pollard Rho	
4.10	Linear Sieve	
4.11	Power Tower	
4.12	Discrete Log / Sqrt	
4.13	De Bruijn Sequence	
4.14	Simplex / LP Duality	
4.15	FFT, FWHT, MultipointEval, Interpolation, TaylorShift	
4.16	Matroid Intersection	
5	String	21
5.1	KMP, Hash, Manacher, Z	21
5.2	Aho-Corasick	21
5.3	$O(N \log N)$ SA + LCP	21
5.4	Suffix Automaton	22
5.5	Bitset LCS	22
5.6	Lyndon Factorization, Minimum Rotation	22
6	Misc	22
6.1	CMakeLists.txt	22
6.2	Ternary Search	23
6.3	Monotone Queue Optimization	23
6.4	Aliens Trick	23
6.5	Slope Trick	23
6.6	Hook Length Formula	23
6.7	Floating Point Add	23
6.8	Random, PBDS, Bit Trick, Bitset	23
6.9	Fast I/O, Fast Div/Mod, Hilbert Mo's	24
6.10	DP Opt, Tree Opt, Well-Known Ideas	24
6.11	Highly Composite Numbers, Large Prime	24
6.12	Catalan, Burnside, Grundy, Pick, Hall, Simpson, Kirchhoff, Area of Quadrangle, Fermat Point, Euler	24
6.13	inclusive and exclusive, Stirling Number, Bell Number	25
6.14	About Graph Matching(Graph with $ V \leq 500$)	25
6.15	Calculus, Newton's Method	25
6.16	Checklist	25

1 DataStructure

1.1 Bipartite Union Find

Usage: Union-Find with friend, enemy relations

```
int P[_Sz], E[_Sz]; // Parent, Enemy, iota(P, P+_Sz, 0);
memset(E, -1, sizeof E);
int find(int v){} bool merge(int u, int v){}
int set_friend(int u, int v){ return merge(u, v); }
int set_enemy(int u, int v){
    int ret = 0;
    if(E[u] == -1) E[u] = v; else ret += merge(E[u], v);
    if(E[v] == -1) E[v] = u; else ret += merge(u, E[v]);
    return ret;
}
```

1.2 Erasable Priority Queue

```
template<class T=int, class O=less<T>>
struct pq_set {
    priority_queue<T, vector<T>, O> q, del;
    const T& top() const { return q.top(); }
    int size() const { return int(q.size()-del.size()); }
    bool empty() const { return !size(); }
    void insert(const T x) { q.push(x); flush(); }
    void pop() { q.pop(); flush(); }
    void erase(const T x) { del.push(x); flush(); }
    void flush() { while(del.size() && q.top()==del.top())
        q.pop(), del.pop(); }
};
```

1.3 Convex Hull Trick

Usage: call init() before use

```
struct Line{
    ll a, b, c; // y = ax + b, c = line index
    Line(ll a, ll b, ll c) : a(a), b(b), c(c) {}
    ll f(ll x){ return a * x + b; }
};
vector<Line> v; int pv;
void init(){ v.clear(); pv = 0; }
int chk(const Line &a, const Line &b, const Line &c) const {
    return (__int128_t)(a.b - b.b) * (b.a - c.a) <=
        (__int128_t)(c.b - b.b) * (b.a - a.a);
}
void insert(Line l){
    if(v.size() > pv && v.back().a == l.a){
        if(l.b < v.back().b) l = v.back(); v.pop_back();
    }
    while(v.size() >= pv+2 && chk(v[v.size()-2], v.back(), l))
        v.pop_back();
    v.push_back(l);
}
p query(ll x){ // if min query, then v[pv].f(x) >= v[pv+1].f(x)
    while(pv+1 < v.size() && v[pv].f(x) <= v[pv+1].f(x)) pv++;
    return {v[pv].f(x), v[pv].c};
}
```

```
}
////// line container start (max query) ////
struct Line {
    mutable ll k, m, p;
    bool operator<(const Line& o) const { return k < o.k; }
    bool operator<(ll x) const { return p < x; }
}; // (for doubles, use inf = 1/.0, div(a,b) = a/b)
struct LineContainer : multiset<Line, less<>> {
    static const ll inf = LLONG_MAX;
    ll div(ll a, ll b) { return a / b - ((a ^ b) < 0 && a % b); }
    // floor
    bool isect(iterator x, iterator y) {
        if (y == end()) return x->p = inf, 0;
        if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
        else x->p = div(y->m - x->m, x->k - y->k);
        return x->p >= y->p;
    }
}
void add(ll k, ll m) {
    auto z = insert({k, m, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() && (--x->p >= y->p) isect(x,
        erase(y)));
}
ll query(ll x) { assert(!empty());
    auto l = *lower_bound(x);
    return l.k * x + l.m;
}
};
```

1.4 Persistent Segment Tree

Usage: call init(root[0], s, e) before use

```
struct PSTNode{
    PSTNode *l, *r; int v;
    PSTNode(){ l = r = nullptr; v = 0; }
};
PSTNode *root[101010];
PST({ memset(root, 0, sizeof root); }) // constructor
void init(PSTNode *node, int s, int e){
    if(s == e) return;
    int m = s + e >> 1;
    node->l = new PSTNode; node->r = new PSTNode;
    init(node->l, s, m); init(node->r, m+1, e);
}
void update(PSTNode *prv, PSTNode *now, int s, int e, int x){
    if(s == e){ now->v = prv ? prv->v + 1 : 1; return; }
    int m = s + e >> 1;
    if(x <= m){
        now->l = new PSTNode; now->r = prv->r;
        update(prv->l, now->l, s, m, x);
    }
    else{
        now->r = new PSTNode; now->l = prv->l;
        update(prv->r, now->r, m+1, e, x);
    }
    int t1 = now->l ? now->l->v : 0;
```

```
    int t2 = now->r ? now->r->v : 0;
    now->v = t1 + t2;
}
int kth(PSTNode *prv, PSTNode *now, int s, int e, int k){
    if(s == e) return s;
    int m = s + e >> 1, diff = now->l->v - prv->l->v;
    if(k <= diff) return kth(prv->l, now->l, s, m, k);
    else return kth(prv->r, now->r, m+1, e, k-diff);
}

1.5 Kinetic Segment Tree

struct line_t{
    ll a, b, v, idx;
    line_t() : line_t(0, nINF) {}
    line_t(ll a, ll b) : line_t(a, b, -1) {}
    line_t(ll a, ll b, ll idx) : a(a), b(b), v(b), idx(idx) {}
    void apply_heat(ll heat){ v += a * heat; }
    void apply_add(ll lz_add){ v += lz_add; }
    ll cross(const line_t &l) const {
        if(a == l.a) return pINF;
        ll p = v - l.v, q = l.a - a;
        if(q < 0) p = -p, q = -q;
        return p >= 0 ? (p + q - 1) / q : -p / q * -1;
    }
    ll cross_after(const line_t &l, ll temp) const {
        ll res = cross(l); return res > temp ? res : pINF;
    }
};
struct range_kinetic_segment_tree{
    struct node_t{
        line_t v;
        ll melt, heat, lz_add;
        node_t() : node_t(line_t()) {}
        node_t(ll a, ll b, ll idx) : node_t(line_t(a, b, idx)) {}
        node_t(const line_t &v) : v(v), melt(pINF), heat(0),
            lz_add(0) {}
        bool operator < (const node_t &o) const { return
            tie(v.v,v.a) < tie(o.v,v,o.v.a); }
        ll cross_after(const node_t &o, ll temp) const { return
            v.cross_after(o.v, temp); }
        void apply_lazy(){ v.apply_heat(heat); v.apply_add(lz_add);
            melt -= heat; }
        void clear_lazy(){ heat = lz_add = 0; }
        void prop_lazy(const node_t &p){ heat += p.heat; lz_add +=
            p.lz_add; }
        bool have_lazy() const { return heat != 0 || lz_add != 0; }
    };
    node_t T[SZ<<1]; range_kinetic_segment_tree(){ clear(); }
    void clear(){ fill(T, T+SZ*2, node_t()); }
    void pull(int node, int s, int e){
        if(s == e) return;
        const node_t &l = T[node<<1], &r = T[node<<1|1];
        assert(!l.have_lazy() && !r.have_lazy() &&
            !T[node].have_lazy());
        T[node] = max(l, r);
```

```
        T[node].melt = min({ l.melt, r.melt, l.cross_after(r, 0)
        });
    }
    void push(int node, int s, int e){
        if(!T[node].have_lazy()) return; T[node].apply_lazy();
        if(s != e) for(auto c : {node<<1, node<<1|1})
            T[c].prop_lazy(T[node]);
        T[node].clear_lazy();
    }
    void build(const vector<line_t> &lines, int node=1, int s=0,
        int e=SZ-1){
        if(s == e){ T[node] = s < lines.size() ? node_t(lines[s]) :
            node_t(); return; }
        int m = (s + e) / 2;
        build(lines,node*2,s,m); build(lines,node*2+1,m+1,e);
        pull(node, s, e);
    }
    void update(int x, const line_t &v, int node=1, int s=0, int
        e=SZ-1){
        push(node, s, e); int m = (s + e) / 2;
        if(s == e){ T[node] = v; return; }
        if(x <= m) update(x, v, node<<1, s, m), push(node<<1|1,
            m+1, e);
        else update(x, v, node<<1|1, m+1, e), push(node<<1, s, m);
        pull(node, s, e);
    }
    void add(int l, int r, ll v, int node=1, int s=0, int
        e=SZ-1){
        push(node, s, e); int m = (s + e) / 2;
        if(r < s || e < l) return;
        if(l <= s && e <= r){ T[node].lz_add += v; push(node, s,
            e); return; }
        add(l,r,v,node*2,s,m); add(l,r,v,node*2+1,m+1,e);
        pull(node, s, e);
    }
    void heaten(int l, int r, ll t, int node=1, int s=0, int
        e=SZ-1){
        push(node, s, e); int m = (s + e) / 2;
        if(r < s || e < l) return;
        if(l <= s && e <= r){ _heat(t, node, s, e); return; }
        heaten(l,r,t,node*2,s,m); heaten(l,r,t,node*2+1,m+1,e);
        pull(node, s, e);
    }
    void _heat(ll t, int node=1, int s=0, int e=SZ-1){
        push(node, s, e); int m = (s + e) / 2;
        if(T[node].melt > t){ T[node].heat += t; push(node, s, e);
            return; }
        _heat(t,node*2,s,m); _heat(t,node*2+1,m+1,e);
        pull(node, s, e);
    }
};
```

1.6 Splay Tree, Link-Cut Tree

```
struct Node{
    Node *l, *r, *p;
    bool flip; int sz;
```

```
T now, sum, lz;
Node(){ l = r = p = nullptr; sz = 1; flip = false; now = sum = lz = 0; }
bool IsLeft() const { return p && this == p->l; }
bool IsRoot() const { return !p || (this != p->l && this != p->r); }
friend int GetSize(const Node *x){ return x ? x->sz : 0; }
friend T GetSum(const Node *x){ return x ? x->sum : 0; }
void Rotate(){
    p->Push(); Push();
    if(IsLeft()) r && (r->p = p), p->l = r, r = p;
    else l && (l->p = p), p->r = l, l = p;
    if(!p->IsRoot()) (p->IsLeft() ? p->p->l : p->p->r) = this;
    auto t = p; p = t->p; t->p = this; t->Update(); Update();
}
void Update(){
    sz = 1 + GetSize(l) + GetSize(r); sum = now + GetSum(l) + GetSum(r);
}
void Update(const T &val){ now = val; Update(); }
void Push(){
    Update(now + lz); if(flip) swap(l, r);
    for(auto c : {l, r}) if(c) c->flip ^= flip, c->lz += lz;
    lz = 0; flip = false;
}
};
Node* rt;
Node* Splay(Node *x, Node *g=nullptr){
    for(g || (rt=x); x->p!=g; x->Rotate()){
        if(!x->p->IsRoot()) x->p->p->Push(); x->p->Push();
        x->Push();
        if(x->p->p != g) (x->IsLeft() ^ x->p->IsLeft() ? x : x->p)->Rotate();
    }
    x->Push(); return x;
}
Node* Kth(int k){
    for(auto x=rt; ; x=x->r){
        for(; x->Push(), x->l && x->l->sz > k; x=x->l);
        if(x->l) k -= x->l->sz;
        if(!k-- return Splay(x);
    }
}
Node* Gather(int s, int e){ auto t = Kth(e+1); return Splay(t, Kth(s-1))->l; }
Node* Flip(int s, int e){ auto x = Gather(s, e); x->flip ^= 1; return x; }
Node* Shift(int s, int e, int k){
    if(k >= 0){
        k %= e-s+1; if(k) Flip(s, e), Flip(s, s+k-1), Flip(s+k, e);
    }
    else{
        k = -k; k %= e-s+1; if(k) Flip(s, e), Flip(s, e-k), Flip(e-k+1, e);
    }
    return Gather(s, e);
}
```

```
int Idx(Node *x){ return x->l->sz; }
////////// Link Cut Tree Start //////////
Node* Splay(Node *x){
    for(; !x->IsRoot(); x->Rotate()){
        if(!x->p->IsRoot()) x->p->p->Push(); x->p->Push();
        x->Push();
        if(!x->p->IsRoot()) (x->IsLeft() ^ x->p->IsLeft() ? x : x->p)->Rotate();
    }
    x->Push(); return x;
}
void Access(Node *x){
    Splay(x); x->r = nullptr; x->Update();
    for(auto y=x; x->p; Splay(x)) y = x->p, Splay(y), y->r = x, y->Update();
}
int GetDepth(Node *x){ Access(x); x->Push(); return GetSize(x->l); }
Node* GetRoot(Node *x){
    Access(x); for(x->Push(); x->l; x->Push()) x = x->l; return Splay(x);
}
Node* GetPar(Node *x){
    Access(x); x->Push(); if(!x->l) return nullptr;
    x = x->l; for(x->Push(); x->r; x->Push()) x = x->r;
    return Splay(x);
}
void Link(Node *p, Node *c){ Access(c); Access(p); c->l = p; p->p = c; c->Update(); }
void Cut(Node *c){ Access(c); c->l->p = nullptr; c->l = nullptr; c->Update(); }
Node* GetLCA(Node *x, Node *y){
    Access(x); Access(y); Splay(x); return x->p ? x->p : x;
}
Node* Ancestor(Node *x, int k){
    k = GetDepth(x) - k; assert(k >= 0);
    for(; x->Push()){
        int s = GetSize(x->l); if(s == k) return Access(x), x;
        if(s < k) k -= s + 1, x = x->r; else x = x->l;
    }
}
void MakeRoot(Node *x){ Access(x); Splay(x); x->flip ^= 1; }
bool IsConnect(Node *x, Node *y){ return GetRoot(x) == GetRoot(y); }
void PathUpdate(Node *x, Node *y, T val){
    Node *root = GetRoot(x); // original root
    MakeRoot(x); Access(y); // make x to root, tie with y
    Splay(x); x->lz += val; x->Push();
    MakeRoot(root); // Revert
    Node *lca = GetLCA(x, y);
    Access(lca); Splay(lca); lca->Push();
    lca->Update(lca->now - val);
}
T VertexQuery(Node *x, Node *y){
    Node *l = GetLCA(x, y); T ret = l->now;
    Access(x); Splay(l); if(l->r) ret = ret + l->r->sum;
    Access(y); Splay(l); if(l->r) ret = ret + l->r->sum;
}
```

```
return ret;
}
Node* GetQueryResultNode(Node *u, Node *v){
    if(GetRoot(u) != GetRoot(v)) return 0;
    MakeRoot(u); Access(v); auto ret = v->l;
    while(ret->mx != ret->v){
        if (ret->l && ret->mx == ret->l->mx) ret = ret->l;
        else ret = ret->r;
    }
    Access(ret); return ret;
}
```

2 Geometry

2.1 Triangles

변 길이 $a, b, c; p = (a + b + c)/2$
넓이 $A = \sqrt{p(p-a)(p-b)(p-c)}$
외접원 반지름 $R = abc/4A$, 내접원 반지름 $r = A/p$
중선 길이 $m_a = 0.5\sqrt{2b^2 + 2c^2 - a^2}$
각 이등분선 길이 $s_a = \sqrt{bc(1 - \frac{a}{b+c}^2)}$
사인 법칙 $\frac{\sin A}{a} = 1/2R$, 코사인 법칙 $a^2 = b^2 + c^2 - 2bc \cos A$, 탄젠트 법칙 $\frac{a+b}{a-b} = \frac{\tan(A+B)/2}{\tan(A-B)/2}$
중심 좌표 $(\frac{\alpha x_a + \beta x_b + \gamma x_c}{\alpha + \beta + \gamma}, \frac{\alpha y_a + \beta y_b + \gamma y_c}{\alpha + \beta + \gamma})$

이름	α	β	γ	
외심	$a^2\mathcal{A}$	$b^2\mathcal{B}$	$c^2\mathcal{C}$	$\mathcal{A} = b^2 + c^2 - a^2$
내심	a	b	c	$\mathcal{B} = a^2 + c^2 - b^2$
무게중심	1	1	1	$\mathcal{C} = a^2 + b^2 - c^2$
수심	\mathcal{BC}	\mathcal{CA}	\mathcal{AB}	
방심(A)	$-a$	b	c	

2.2 Rotating Calipers

```
pair<Point, Point> RotatingCalipers(const vector<Point> &H){
    ll mx = 0; Point a, b;
    for(int i=0, j=0; i<H.size(); i++){
        while(j+1 < H.size() && CCW(0, H[i+1]-H[i], H[j+1]-H[j]) >= 0){
            if(ll now = D2(H[i], H[j]); mx < now) mx = now, a = H[i], b = H[j];
            j++;
        }
        if(ll now = D2(H[i], H[j]); mx < now) mx = now, a = H[i], b = H[j];
    }
    return {a, b};
}
```

2.3 Point in Convex Polygon

```
bool Check(const vector<Point> &v, const Point &pt){
    if(CCW(v[0], v[1], pt) < 0) return false; int l = 1, r = v.size() - 1;
    while(l < r){
```

```
int m = l + r + 1 >> 1;
if(CCW(v[0], v[m], pt) >= 0) l = m; else r = m - 1;
}
if(l == v.size() - 1) return CCW(v[0], v.back(), pt) == 0 &&
v[0] <= pt && pt <= v.back();
return CCW(v[0], v[l], pt) >= 0 && CCW(v[l], v[l+1], pt) >= 0
&& CCW(v[l+1], v[0], pt) >= 0;
}
```

2.4 Segment Distance

```
double Proj(Point a, Point b, Point c){
    ll t1 = (b - a) * (c - a), t2 = (a - b) * (c - b);
    if(t1 * t2 >= 0 && CCW(a, b, c) != 0)
        return abs(CCW(a, b, c)) / sqrt(Dist(a, b));
    else return 1e18;
}
double Dist(Point a[2], Point b[2]){
    double res = 1e18; // NOTE: need to check intersect
    for(int i=0; i<4; i++) res = min(res, sqrt(Dist(a[i/2],
        b[i%2])));
    for(int i=0; i<2; i++) res = min(res, Proj(a[0], a[1],
        b[i]));
    for(int i=0; i<2; i++) res = min(res, Proj(b[0], b[1],
        a[i]));
    return res;
}
```

2.5 Tangent Series

```
template<bool UPPER=true>
Point GetPoint(const vector<Point> &hull, real_t slope){
    auto chk = [slope](real_t dx, real_t dy){ return UPPER ? dy
    >= slope * dx : dy <= slope * dx; };
    int l = -1, r = hull.size() - 1;
    while(l + 1 < r){
        int m = (l + r) / 2;
        if(chk(hull[m+1].x - hull[m].x, hull[m+1].y -
            hull[m].y)) l = m; else r = m;
    }
    return hull[r];
}
int ConvexTangent(const vector<Point> &v, const Point &pt, int
up=1){ //given outer point
    auto sign = [&](ll c){ return c > 0 ? up : c == 0 ? 0 : -up;
    };
    auto local = [&](Point p, Point a, Point b, Point c){
        return sign(CCW(p, a, b)) <= 0 && sign(CCW(p, b, c)) >= 0;
    }; // assert(v.size() >= 2);
    int n = v.size() - 1, s = 0, e = n, m;
    if(local(pt, v[1], v[0], v[n-1])) return 0;
    while(s + 1 < e){
        m = (s + e) / 2;
        if(local(pt, v[m-1], v[m], v[m+1])) return m;
        if(sign(CCW(pt, v[s], v[s+1])) < 0){ // up
            if(sign(CCW(pt, v[m], v[m+1])) > 0) e = m;
            else if(sign(CCW(pt, v[m], v[s])) > 0) s = m; else e = m;
        }
    }
}
```

```
    }
    else{ // down
        if(sign(CCW(pt, v[m], v[m+1])) < 0) s = m;
        else if(sign(CCW(pt, v[m], v[s])) < 0) s = m; else e = m;
    }
}
if(s && local(pt, v[s-1], v[s], v[s+1])) return s;
if(e != n && local(pt, v[e-1], v[e], v[e+1])) return e;
return -1;
}
int Closest(const vector<Point> &v, const Point &out, int now){
    int prv = now > 0 ? now-1 : v.size()-1, nxt = now+1 <
    v.size() ? now+1 : 0, res = now;
    if(CCW(out, v[now], v[prv]) == 0 && Dist(out, v[res]) >
    Dist(out, v[prv])) res = prv;
    if(CCW(out, v[now], v[nxt]) == 0 && Dist(out, v[res]) >
    Dist(out, v[nxt])) res = nxt;
    return res; // if parallel, return closest point to out
} // int point_idx = Closest(convex_hull, pt,
ConvexTangent(hull + hull[0], pt, +-1) % N);
int tangent(circle &A, circle &B, pdd des[4]){ // return angle
    int top = 0; // outer
    double d = size(A.O - B.O), a = polar(B.O - A.O), b = PI + a;
    double t = sq(d) - sq(A.r - B.r);
    if (t >= 0){
        t = sqrt(t);
        double p = atan2(B.r - A.r, t);
        des[top++] = pdd(a + p + PI / 2, b + p - PI / 2);
        des[top++] = pdd(a - p - PI / 2, b - p + PI / 2);
    }
    t = sq(d) - sq(A.r + B.r); // inner
    if (t >= 0){ t = sqrt(t);
        double p = atan2(B.r + A.r, t);
        des[top++] = pdd(a + p - PI / 2, b + p - PI / 2);
        des[top++] = pdd(a - p + PI / 2, b - p + PI / 2);
    }
    return top;
}
```

2.6 Intersect Series

```
// 0: not intersect, -1: infinity, 1: cross
// flag, xp, xq, yp, yq : (xp / xq, yp / yq)
using T = __int128_t; // T <= 0(COORD^3)
tuple<int,T,T,T,T> SegmentIntersect(Point s1, Point e1, Point
s2, Point e2){
    if(!Intersect(s1, e1, s2, e2)) return {0, 0, 0, 0, 0};
    auto det = (e1 - s1) / (e2 - s2);
    if(!det){
        if(s1 > e1) swap(s1, e1);
        if(s2 > e2) swap(s2, e2);
        if(e1 == s2) return {1, e1.x, 1, e1.y, 1};
        if(e2 == s1) return {1, e2.x, 1, e2.y, 1};
        return {-1, 0, 0, 0, 0};
    }
    T p = (s2 - s1) / (e2 - s2), q = det;
    T xp = s1.x * q + (e1.x - s1.x) * p, xq = q;
```

```
T yp = s1.y * q + (e1.y - s1.y) * p, yq = q;
if(xq < 0) xp = -xp, xq = -xq;
if(yq < 0) yp = -yp, yq = -yq;
T xg = __gcd(abs(xp), xq), yg = __gcd(abs(yp), yq);
return {1, xp/xg, xq/xg, yp/yg, yq/yg};
}
bool circleIntersect(P a,P b,double r1,double r2,pair<P, P>*
out){
    if (a == b) { assert(r1 != r2); return false; }
    P vec = b-a; double d2 = vec.dist2(), sum = r1+r2, dif =
    r1-r2;
    double p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
    if (sum*sum < d2 || dif*dif > d2) return false; // use EPS
    plz...
    P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2) / d2);
    *out = {mid + per, mid - per}; return true;
}
vector<P> circleLine(P c, double r, P a, P b){
    P ab = b - a, p = a + ab * (c-a).dot(ab) / ab.dist2();
    double s = a.cross(b, c), h2 = r*r - s*s / ab.dist2();
    if (h2 < 0) return {}; if (h2 == 0) return {p};
    P h = ab.unit() * sqrt(h2); return {p - h, p + h};
}
double circlePoly(P c, double r, vector<P> ps){ // return area
    auto tri = [&](P p, P q) { // ps must be ccw polygon
        auto r2 = r * r / 2; P d = q - p;
        auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2();
        auto det = a * a - b;
        if (det <= 0) return arg(p, q) * r2;
        auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
        if (t < 0 || 1 <= s) return arg(p, q) * r2;
        P u = p + d * s, v = p + d * t;
        return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
    };
    auto sum = 0.0;
    rep(i,0,sz(ps)) sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] -
    c);
    return sum;
}
// extrVertex: point of hull, max projection onto line
#define cmp(i,j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
int extrVertex(vector<P>& poly, P dir) {
    int n = sz(poly), lo = 0, hi = n;
    if (extr(0)) return 0;
    while (lo + 1 < hi) {
        int m = (lo + hi) / 2; if (extr(m)) return m;
        int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
        (ls < ms || (ls == ms && ls == cmp(lo, m)) ? hi : lo) = m;
    }
    return lo;
}
//(-1,-1): no collision
//(i,-1): touch corner
//(i,i): along side (i,i+1)
//(i,j): cross (i,i+1)and(j,j+1)
//(i,i+1): cross corner i
```



```
// 0(log n), ccw no colinear point convex polygon
// P perp() const { return P(-y, x); }
#define cmpL(i) sgn(a.cross(poly[i], b))
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
    int endA = extrVertex(poly, (a - b).perp());
    int endB = extrVertex(poly, (b - a).perp());
    if (cmpL(endA) < 0 || cmpL(endB) > 0) return {-1, -1};
    array<int, 2> res;
    rep(i,0,2) {
        int lo = endB, hi = endA, n = sz(poly);
        while ((lo + 1) % n != hi) {
            int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
            (cmpL(m) == cmpL(endB) ? lo : hi) = m;
        }
        res[i] = (lo + !cmpL(hi)) % n;
        swap(endA, endB);
    }
    if (res[0] == res[1]) return {res[0], -1};
    if (!cmpL(res[0]) && !cmpL(res[1]))
        switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
            case 0: return {res[0], res[0]};
            case 2: return {res[1], res[1]};
        }
    return res;
}
```

2.7 Polygon Cut, Center, Union

```
// Returns the polygon on the left of line l
// *: dot product, ^: cross product
// l = p + d*t, l.q() = l + d
// doubled_signed_area(p,q,r) = (q-p) ^ (r-p)
template<class T> vector<point<T>> polygon_cut(const
vector<point<T>> &a, const line<T> &l){
    vector<point<T>> res;
    for(auto i = 0; i < (int)a.size(); ++ i){
        auto cur = a[i], prev = i ? a[i - 1] : a.back();
        bool side = doubled_signed_area(l.p, l.q(), cur) > 0;
        if(side != (doubled_signed_area(l.p, l.q(), prev) > 0))
            res.push_back(l.p + (cur - l.p ^ prev - cur) / (l.d ^
            prev - cur) * l.d);
        if(side) res.push_back(cur);
    }
    return res;
}

P polygonCenter(const vector<P>& v){ // center of mass
P res(0, 0); double A = 0;
for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
    res = res + (v[i] + v[j]) * v[j].cross(v[i]);
    A += v[j].cross(v[i]);
} return res / A / 3;
}

// 0(points^2), area of union of n polygon, ccw polygon
double rat(P a, P b) { return sgn(b.x) ? a.x/b.x : a.y/b.y; }
double polyUnion(vector<vector<P>>& poly) {
    double ret = 0;
    rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
```

```
P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
rep(j,0,sz(poly)) if (i != j) {
    rep(u,0,sz(poly[j])) {
        P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])];
        int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
        if (sc != sd) {
            double sa = C.cross(D, A), sb = C.cross(D, B);
            if (min(sc, sd) < 0)
                segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
        } else if (!sc && !sd && j < i && sgn((B-A).dot(D-C))>0){
            segs.emplace_back(rat(C - A, B - A), 1);
            segs.emplace_back(rat(D - A, B - A), -1);
        }
    }
}
sort(all(segs));
for (auto& s : segs) s.first = min(max(s.first, 0.0), 1.0);
double sum = 0;
int cnt = segs[0].second;
rep(j,1,sz(segs)) {
    if (!cnt) sum += segs[j].first - segs[j - 1].first;
    cnt += segs[j].second;
}
ret += A.cross(B) * sum;
}
return ret / 2;
}
```

2.8 Polygon Raycast

```
// ray A + kd and CCW polygon C, return events {k, event_id}
// 0: out->line / 1: in->line / 2: line->out / 3: line->in
// 4: pass corner outside / 5: pass corner inside / 6: out ->
in / 7: in -> out
// WARNING: C.push_back(C[0]) before working
struct frac{
    ll first, second; frac(){}
    frac(ll a, ll b) : first(a), second(b) {
        if( b < 0 ) first = -a, second = -b; // operator cast
        int128
    } double v(){ return 1.*first/second; } // operator <,<=,==
};
frac raypoints(vector<pii> &C, pii A, pii d, vector<pair<frac,
int>> &R){
    assert(d != pii(0, 0));
    int g = gcd(abs(d.first), abs(d.second));
    d.first /= g, d.second /= g;
    vector<pair<frac, int>> L;
    for(int i = 0; i+1 < C.size(); i++){
        pii v = C[i+1] - C[i];
        int a = sign(d/(C[i]-A)), b = sign(d/(C[i+1]-A));
        if( a == 0 ) L.emplace_back(frac(d*(C[i]-A)/size2(d), 1),
        b);
        if( b == 0 ) L.emplace_back(frac(d*(C[i+1]-A)/size2(d), 1),
        a);
        if( a*b == -1 ) L.emplace_back(frac((A-C[i])/v, v/d), 6);
```

```

    }
    sort(L.begin(), L.end());
    int sz = 0;
    for(int i = 0; i < L.size(); i++){
        // assert(i+2 >= L.size() || !(L[i].first ==
        L[i+2].first));
        if( i+1 < L.size() && L[i].first == L[i+1].first &&
        L[i].second != 6){
            int a = L[i].second, b = L[i+1].second;
            R.emplace_back(L[i+1].first, a*b ? a*b > 0? 4: 6:
            (1-a-b)/2);
        }
        else R.push_back(L[i]);
    }
    int state = 0; // 0: out, 1: in, 2: line+ccw, 3: line+cw
    for(auto &e : R){
        int &n = e.second;
        if( n == 6 ) n ^= state, state ^= 1;
        else if( n == 4 ) n ^= state;
        else if( n == 0 ) n = state, state ^= 2;
        else if( n == 1 ) n = state^(state>>1), state ^= 3;
    } return frac(g, 1);
}

bool visible(vector<pii> &C, pii A, pii B){
    if( A == B ) return true;
    char I[4] = "356", O[4] = "157";
    vector<pair<frac, int>> R; vector<frac> E;
    frac s = frac(0, 1), e = raypoints(C, A, B-A, R);
    for(auto e : R){
        int &n = e.second, m;
        if(*find(O, O+3, n+'0')) E.emplace_back(e.first);
        if(*find(I, I+3, n+'0')) E.emplace_back(e.first);
    }
    for(int j = 0; j < E.size(); j += 2) if( !(e <= E[j] ||
    E[j+1] <= s) ) return false;
    return true;
}
```

2.9 Shamos-Hoeey

```
struct Line{
    static ll CUR_X; ll x1, y1, x2, y2, id;
    Line(Point p1, Point p2, int id) : id(id) {
        if(p1 > p2) swap(p1, p2);
        tie(x1,y1) = p1; tie(x2,y2) = p2;
    } Line() = default;
    int get_k() const { return y1 != y2 ? (x2-x1)/(y1-y2) : -1; }
    void convert_k(int k){ // x1,y1,x2,y2 = 0(COORD^2), use i128
    in ccw
        Line res;
        res.x1 = x1 + y1 * k; res.y1 = -x1 * k + y1;
        res.x2 = x2 + y2 * k; res.y2 = -x2 * k + y2;
        x1 = res.x1; y1 = res.y1; x2 = res.x2; y2 = res.y2;
        if(x1 > x2) swap(x1, x2), swap(y1, y2);
    }
    ld get_y(ll offset=0) const { // OVERFLOW
        ld t = ld(CUR_X-x1+offset) / (x2-x1);
```

```
        return t * (y2 - y1) + y1;
    }
    bool operator < (const Line &l) const {
        return get_y() < l.get_y();
    }
    // strict
    /* bool operator < (const Line &l) const {
        auto le = get_y(), ri = l.get_y();
        if(abs(le-ri) > 1e-7) return le < ri;
        if(CUR_X == x1 || CUR_X == l.x1) return get_y(1) <
l.get_y(1);
        else return get_y(-1) < l.get_y(-1);
    } */
}; ll Line::CUR_X = 0;
struct Event{ // f=0 st, f=1 ed
    ll x, y, i, f; Event() = default;
    Event(Line l, ll i, ll f) : i(i), f(f) {
        if(f==0) tie(x,y) = tie(l.x1,l.y1);
        else tie(x,y) = tie(l.x2,l.y2);
    }
    bool operator < (const Event &e) const {
        return tie(x,f,y) < tie(e.x,e.f,e.y);
        // strict
        // return make_tuple(x,-f,y) < make_tuple(e.x,-e.f,e.y);
    }
};
tuple<bool,int,int> ShamosHoeey(vector<array<Point,2>> v){
    int n = v.size(); vector<int> use(n+1);
    vector<Line> lines; vector<Event> E; multiset<Line> T;
    for(int i=0; i<n; i++){
        lines.emplace_back(v[i][0], v[i][1], i);
        if(int t=lines[i].get_k(); 0<=t && t<=n) use[t] = 1;
    }
    int k = find(use.begin(), use.end(), 0) - use.begin();
    for(int i=0; i<n; i++){
        lines[i].convert_k(k);
        E.emplace_back(lines[i], i, 0);
        E.emplace_back(lines[i], i, 1);
    } sort(E.begin(), E.end());
    for(auto &e : E){
        Line::CUR_X = e.x;
        if(e.f == 0){
            auto it = T.insert(lines[e.i]);
            if(next(it) != T.end() && Intersect(lines[e.i],
*next(it))) return {true, e.i, next(it)->id};
            if(it != T.begin() && Intersect(lines[e.i], *prev(it)))
return {true, e.i, prev(it)->id};
        }
        else{
            auto it = T.lower_bound(lines[e.i]);
            if(it != T.begin() && next(it) != T.end() &&
Intersect(*prev(it), *next(it))) return {true,
prev(it)->id, next(it)->id};
            T.erase(it);
        }
    }
    return {false, -1, -1};
}
```

```
}

2.10 Half Plane Intersection

Usage: Line :  $ax + by + c = 0$ 

double CCW(p1, p2, p3); bool same(double a, double b); const
Point o = Point(0, 0);
struct Line{
    double a, b, c; Line() : Line(0, 0, 0) {}
    Line(double a, double b, double c) : a(a), b(b), c(c) {}
    bool operator < (const Line &l) const {
        bool f1 = Point(a, b) > o, f2 = Point(l.a, l.b) > o;
        if(f1 != f2) return f1 > f2;
        double cw = CCW(o, Point(a, b), Point(l.a, l.b));
        return same(cw, 0) ? c * hypot(l.a, l.b) < l.c * hypot(a,
b) : cw > 0;
    }
    Point slope() const { return Point(a, b); }
};
Point LineIntersect(Line a, Line b){
    double det = a.a*b.b - b.a*a.b, x = (a.c*b.b - a.b*b.c) /
det, y = (a.a*b.c - a.c*b.a) / det;
    return Point(x, y);
}
bool CheckHPI(Line a, Line b, Line c){
    if(CCW(o, a.slope(), b.slope()) <= 0) return 0;
    Point v = LineIntersect(a, b); return v.x*c.a + v.y*c.b >=
c.c;
}
vector<Point> HPI(vector<Line> v){
    sort(v.begin(), v.end());
    deque<Line> dq; vector<Point> ret;
    for(auto &i : v){
        if(dq.size() && same(CCW(o, dq.back().slope(), i.slope()),
0)) continue;
        while(dq.size() >= 2 && CheckHPI(dq[dq.size()-2],
dq.back(), i)) dq.pop_back();
        while(dq.size() >= 2 && CheckHPI(i, dq[0], dq[1]))
dq.pop_front();
        dq.push_back(i);
    }
    while(dq.size() > 2 && CheckHPI(dq[dq.size()-2], dq.back(),
dq[0])) dq.pop_back();
    while(dq.size() > 2 && CheckHPI(dq.back(), dq[0], dq[1]))
dq.pop_front();
    for(int i=0; i<dq.size(); i++){
        Line now = dq[i], nxt = dq[(i+1)%dq.size()];
        if(CCW(o, now.slope(), nxt.slope()) <= eps) return
vector<Point>();
        ret.push_back(LineIntersect(now, nxt));
    }
    for(auto &[x,y] : ret) x = -x, y = -y;
    return ret;
}
```

```
2.11 K-D Tree

T GetDist(const P &a, const P &b){ return (a.x-b.x) * (a.x-b.x)
+ (a.y-b.y) * (a.y-b.y); }
struct Node{
    P p; int idx;
    T x1, y1, x2, y2;
    Node(const P &p, const int idx) : p(p), idx(idx), x1(1e9),
y1(1e9), x2(-1e9), y2(-1e9) {}
    bool contain(const P &pt) const{ return x1 <= pt.x && pt.x <=
x2 && y1 <= pt.y && pt.y <= y2; }
    T dist(const P &pt) const { return idx == -1 ? INF :
GetDist(p, pt); }
    T dist_to_border(const P &pt) const {
        const auto [x,y] = pt;
        if(x1 <= x && x <= x2) return min((y-y1)*(y-y1),
(y2-y)*(y2-y));
        if(y1 <= y && y <= y2) return min((x-x1)*(x-x1),
(x2-x)*(x2-x));
        T t11 = GetDist(pt, {x1,y1}), t12 = GetDist(pt, {x1,y2});
        T t21 = GetDist(pt, {x2,y1}), t22 = GetDist(pt, {x2,y2});
        return min({t11, t12, t21, t22});
    }
};
template<bool IsFirst = 1> struct Cmp {
    bool operator() (const Node &a, const Node &b) const {
        return IsFirst ? a.p.x < b.p.x : a.p.y < b.p.y;
    }
};
struct KDTree { // Warning : no duplicate
    constexpr static size_t NAIVE_THRESHOLD = 16;
    vector<Node> tree;
    KDTree() = default;
    explicit KDTree(const vector<P> &v) {
        for(int i=0; i<v.size(); i++) tree.emplace_back(v[i], i);
        Build(0, v.size());
    }
    template<bool IsFirst = 1>
    void Build(int l, int r) {
        if(r - l <= NAIVE_THRESHOLD) return;
        const int m = (l + r) >> 1;
        nth_element(tree.begin()+l, tree.begin()+m, tree.begin()+r,
Cmp<IsFirst>{});
        for(int i=l; i<r; i++){
            tree[m].x1 = min(tree[m].x1, tree[i].p.x); tree[m].y1 =
min(tree[m].y1, tree[i].p.y);
            tree[m].x2 = max(tree[m].x2, tree[i].p.x); tree[m].y2 =
max(tree[m].y2, tree[i].p.y);
        }
        Build<!IsFirst>(l, m); Build<!IsFirst>(m + 1, r);
    }
    template<bool IsFirst = 1>
    void Query(const P &p, int l, int r, Node &res) const {
        if(r - l <= NAIVE_THRESHOLD){
            for(int i=l; i<r; i++) if(p != tree[i].p && res.dist(p) >
tree[i].dist(p)) res = tree[i];
        }
        else{
```

```
const int m = (1 + r) >> 1;
const T t = IsFirst ? p.x - tree[m].p.x : p.y -
tree[m].p.y;
if(p != tree[m].p && res.dist(p) > tree[m].dist(p)) res =
tree[m];
if(!tree[m].contain(p) && tree[m].dist_to_border(p) >=
res.dist(p)) return;
if(t < 0){
    Query<!IsFirst>(p, l, m, res);
    if(t*t < res.dist(p)) Query<!IsFirst>(p, m+1, r, res);
}
else{
    Query<!IsFirst>(p, m+1, r, res);
    if(t*t < res.dist(p)) Query<!IsFirst>(p, l, m, res);
}
}
}
int Query(const P& p) const {
    Node ret(make_pair<T>(1e9, 1e9), -1); Query(p, 0,
tree.size(), ret); return ret.idx;
}
};
```

2.12 Dual Graph

```
constexpr int quadrant_id(const Point p){
    constexpr int arr[9] = { 5, 4, 3, 6, -1, 2, 7, 0, 1 };
    return arr[sign(p.x)*3+sign(p.y)+4];
}
pair<vector<int>, int> dual_graph(const vector<Point> &points,
const vector<pair<int,int>> &edges){
    int n = points.size(), m = edges.size();
    vector<int> uf(2*m); iota(uf.begin(), uf.end(), 0);
    function<int(int)> find = [&](int v){ return v == uf[v] ? v :
uf[v] = find(uf[v]); };
    function<bool(int,int)> merge = [&](int u, int v){ return
find(u) != find(v) && (uf[uf[u]]=uf[v], true); };
    vector<vector<pair<int,int>>> g(n);
    for(int i=0; i<m; i++){
        g[edges[i].first].emplace_back(edges[i].second, i);
        g[edges[i].second].emplace_back(edges[i].first, i);
    }
    for(int i=0; i<n; i++){
        const auto base = points[i];
        sort(g[i].begin(), g[i].end(), [&](auto a, auto b){
            auto p1 = points[a.first] - base, p2 = points[b.first] -
base;
            return quadrant_id(p1) != quadrant_id(p2) ?
quadrant_id(p1) < quadrant_id(p2) : p1.cross(p2) > 0;
        });
        for(int j=0; j<g[i].size(); j++){
            int k = j ? j - 1 : g[i].size() - 1;
            int u = g[i][k].second << 1, v = g[i][j].second << 1 | 1;
            auto p1 = points[g[i][k].first], p2 =
points[g[i][j].first];
            if(p1 < base) u ^= 1; if(p2 < base) v ^= 1;
            merge(u, v);
        }
    }
}
```

```
    }
}
vector<int> res(2*m);
for(int i=0; i<2*m; i++) res[i] = find(i);
auto comp = res; compress(comp);
for(auto &i : res) i = IDX(comp, i);
int mx_idx = max_element(points.begin(), points.end()) -
points.begin();
return {res, res[g[mx_idx].back().second << 1 | 1]};
}
```

2.13 Bulldozer Trick (Rotating Sweep Line)

```
struct Line{
    ll i, j, dx, dy; // dx >= 0
    Line(int i, int j, const Point &pi, const Point &pj)
        : i(i), j(j), dx(pj.x-pi.x), dy(pj.y-pi.y) {}
    bool operator < (const Line &l) const {
        return make_tuple(dy*l.dx, i, j) < make_tuple(l.dy*dx, l.i,
l.j);
    }
    bool operator == (const Line &l) const {
        return dy * l.dx == l.dy * dx;
    }
};
void Solve(){
    sort(A+1, A+N+1); iota(P+1, P+N+1, 1);
    vector<Line> V; V.reserve(N*(N-1)/2);
    for(int i=1; i<=N; i++) for(int j=i+1; j<=N; j++)
V.emplace_back(i, j, A[i], A[j]);
    sort(V.begin(), V.end());
    for(int i=0, j=0; i<V.size(); i=j){
        while(j < V.size() && V[i] == V[j]) j++;
        for(int k=i; k<j; k++){
            int u = V[k].i, v = V[k].j; // point id, index -> Pos[id]
            swap(Pos[u], Pos[v]); swap(A[Pos[u]], A[Pos[v]]);
            if(Pos[u] > Pos[v]) swap(u, v);
            // @TODO
        }
    }
}
```

2.14 Smallest Enclosing Circle

```
pt getCenter(pt a, pt b){ return pt((a.x+b.x)/2, (a.y+b.y)/2);
}
pt getCenter(pt a, pt b, pt c){
    pt aa = b - a, bb = c - a;
    auto c1 = aa*aa * 0.5, c2 = bb*bb * 0.5, d = aa / bb;
    auto x = a.x + (c1 * bb.y - c2 * aa.y) / d;
    auto y = a.y + (c2 * aa.x - c1 * bb.x) / d;
    return pt(x, y);
}
Circle solve(vector<pt> v){
    pt p = {0, 0};
    double r = 0; int n = v.size();
    for(int i=0; i<n; i++) if(dst(p, v[i]) > r + EPS){
```

```
    p = v[i]; r = 0;
    for(int j=0; j<i; j++) if(dst(p, v[j]) > r + EPS){
        p = getCenter(v[i], v[j]); r = dst(p, v[i]);
        for(int k=0; k<j; k++) if(dst(p, v[k]) > r + EPS){
            p = getCenter(v[i], v[j], v[k]); r = dst(v[k], p);
        }
    }
}
return {p, r};
}
```

2.15 Voronoi Diagram

```
/*
input: order will be changed, sorted by (y,x) order
vertex: voronoi intersection points, degree 3, may duplicated
edge: may contain inf line (-1)
area
    - (a,b) = i-th element of area
    - (u,v) = i-th element of edge
    - input[a] is located CCW of u->v line
    - input[b] is located CW of u->v line
    - u->v line is a subset of perpendicular bisector of input[a]
to input[b] segment
    - Straight line {a, b}, {-1, -1} through midpoint of input[a]
and input[b]
*/
const double EPS = 1e-9;
int dcmp(double x){ return x < -EPS? -1 : x > EPS ? 1 : 0; }
// sq(x) = x*x, size(p) = hypot(p.x, p.y)
// sz2(p) = sq(p.x)+sq(p.y), r90(p) = (-p.y, p.x)
double sq(double x){ return x*x; }
double size(pdd p){ return hypot(p.x, p.y); }
double sz2(pdd p){ return sq(p.x) + sq(p.y); }
pdd r90(pdd p){ return pdd(-p.y, p.x); }
pdd line_intersect(pdd a, pdd b, pdd u, pdd v){ return u +
(((a-u)/b) / ((v/b))*v); }
pdd get_circumcenter(pdd p0, pdd p1, pdd p2){
    return line_intersect(0.5 * (p0+p1), r90(p0-p1), 0.5 *
(p1+p2), r90(p1-p2)); }
double pb_int(pdd left, pdd right, double sweepline){
    if(dcmp(left.y - right.y) == 0) return (left.x + right.x) /
2.0;
    int sign = left.y < right.y ? -1 : 1;
    pdd v = line_intersect(left, right-left, pdd(0, sweepline),
pdd(1, 0));
    double d1 = sz2(0.5 * (left+right) - v), d2 = sz2(0.5 *
(left-right));
    return v.x + sign * sqrt(std::max(0.0, d1 - d2)); }
struct Beachline{
    struct node{ node(){}
        node(pdd point, int idx):point(point), idx(idx), end(0),
link{0, 0}, par(0), prv(0), nxt(0) {}
        pdd point; int idx; int end;
        node *link[2], *par, *prv, *nxt; };
    node *root;
    double sweepline;
```

```
Beachline() : sweepline(-1e20), root(NULL){ }
inline int dir(node *x){ return x->par->link[0] != x; }
void rotate(node *n){
    node *p = n->par; int d = dir(n);
    p->link[d] = n->link[!d];
    if(n->link[!d]) n->link[!d]->par = p;
    n->par = p->par; if(p->par) p->par->link[dir(p)] = n;
    n->link[!d] = p; p->par = n;
} void splay(node *x, node *f = NULL){
    while(x->par != f){
        if(x->par->par == f){
            if(x->par->par == f);
        } else if(dir(x) == dir(x->par)) rotate(x->par);
        else rotate(x);
        rotate(x); }
    if(f == NULL) root = x;
} void insert(node *n, node *p, int d){
    splay(p); node* c = p->link[d];
    n->link[d] = c; if(c) c->par = n;
    p->link[d] = n; n->par = p;
    node *prv = !d?p->prv:p, *nxt = !d?p->nxt:
    n->prv = prv; if(prv) prv->nxt = n;
    n->nxt = nxt; if(nxt) nxt->prv = n;
} void erase(node* n){
    node *prv = n->prv, *nxt = n->nxt;
    if(!prv && !nxt){ if(n == root) root = NULL; return; }
    n->prv = NULL; if(prv) prv->nxt = nxt;
    n->nxt = NULL; if(nxt) nxt->prv = prv;
    splay(n);
    if(!nxt){
        root->par = NULL; n->link[0] = NULL;
        root = prv; }
    else{
        splay(nxt, n); node* c = n->link[0];
        nxt->link[0] = c; c->par = nxt; n->link[0] = NULL;
        n->link[1] = NULL; nxt->par = NULL;
        root = nxt; }
} bool get_event(node* cur, double &next_sweep){
    if(!cur->prv || !cur->nxt) return false;
    pdd u = r90(cur->point - cur->prv->point);
    pdd v = r90(cur->nxt->point - cur->point);
    if(dcmp(u/v) != 1) return false;
    pdd p = get_circumcenter(cur->point, cur->prv->point,
    cur->nxt->point);
    next_sweep = p.y + size(p - cur->point); return true;
} node* find_bl(double x){
    node* cur = root;
    while(cur){
        double left = cur->prv ? pb_int(cur->prv->point,
        cur->point, sweepline) : -1e30;
        double right = cur->nxt ? pb_int(cur->point,
        cur->nxt->point, sweepline) : 1e30;
        if(left <= x && x <= right){ splay(cur); return cur; }
        cur = cur->link[x > right]; }
}
}; using BNode = Beachline::node;
static BNode* arr;
static int sz;
```

```
static BNode* new_node(pdd point, int idx){
    arr[sz] = BNode(point, idx); return arr + (sz++); }
struct event{
    event(double sweep, int idx):type(0), sweep(sweep),
    idx(idx){}
    event(double sweep, BNode* cur):type(1), sweep(sweep),
    prv(cur->prv->idx), cur(cur), nxt(cur->nxt->idx){}
    int type, idx, prv, nxt; BNode* cur; double sweep;
    bool operator>(const event &l)const{ return sweep > l.sweep;
    }
};
void VoronoiDiagram(vector<pdd> &input, vector<pdd> &vertex,
vector<pii> &edge, vector<pii> &area){
    Beachline bl = Beachline();
    priority_queue<event, vector<event>, greater<event>> events;
    auto add_edge = [&](int u, int v, int a, int b, BNode* c1,
    BNode* c2){
        if(c1) c1->end = edge.size()*2;
        if(c2) c2->end = edge.size()*2 + 1;
        edge.emplace_back(u, v);
        area.emplace_back(a, b);
    };
    auto write_edge = [&](int idx, int v){ idx%2 == 0 ?
    edge[idx/2].x = v : edge[idx/2].y = v; };
    auto add_event = [&](BNode* cur){ double nxt;
    if(bl.get_event(cur, nxt)) events.emplace(nxt, cur); };
    int n = input.size(), cnt = 0;
    arr = new BNode[n*4]; sz = 0;
    sort(input.begin(), input.end(), [](const pdd &l, const pdd
    &r){
        return l.y != r.y ? l.y < r.y : l.x < r.x; });
    BNode* tmp = bl.root = new_node(input[0], 0), *t2;
    for(int i = 1; i < n; i++){
        if(dcmp(input[i].y - input[0].y) == 0){
            add_edge(-1, -1, i-1, i, 0, tmp);
            bl.insert(t2 = new_node(input[i], i), tmp, 1);
            tmp = t2;
        }
        else events.emplace(input[i].y, i);
    }
    while(events.size()){
        event q = events.top(); events.pop();
        BNode *prv, *cur, *nxt, *site;
        int v = vertex.size(), idx = q.idx;
        bl.sweepline = q.sweep;
        if(q.type == 0){
            pdd point = input[idx];
            cur = bl.find_bl(point.x);
            bl.insert(site = new_node(point, idx), cur, 0);
            bl.insert(prv = new_node(cur->point, cur->idx), site, 0);
            add_edge(-1, -1, cur->idx, idx, site, prv);
            add_event(prv); add_event(cur);
        }
        else{
            cur = q.cur, prv = cur->prv, nxt = cur->nxt;
            if(!prv || !nxt || prv->idx != q.prv || nxt->idx !=
            q.nxt) continue;
```

```
vertex.push_back(get_circumcenter(prv->point, nxt->point,
cur->point));
write_edge(prv->end, v); write_edge(cur->end, v);
add_edge(v, -1, prv->idx, nxt->idx, 0, prv);
bl.erase(cur);
add_event(prv); add_event(nxt);
    }
}
delete arr;
}
```

3 Graph

3.1 Euler Tour

```
// Not Directed / Cycle
constexpr int SZ = 1010;
int N, G[SZ][SZ], Deg[SZ], Work[SZ];
void DFS(int v){
    for(int &i=Work[v]; i<=N; i++) while(G[v][i]) G[v][i]--,
    G[i][v]--, DFS(i);
    cout << v << " ";
}
// Directed / Path
void DFS(int v){
    for(int i=1; i<=pv; i++) while(G[v][i]) G[v][i]--, DFS(i);
    Path.push_back(v);
}
void Get(){
    for(int i=1; i<=pv; i++) if(In[i] < Out[i]){ DFS(i); return;
    }
    for(int i=1; i<=pv; i++) if(Out[i]){ DFS(i); return; }
}
```

3.2 2-SAT

```
int SZ; vector<vector<int>> G1, G2;
void Init(int n){ SZ = n; G1 = G2 = vector<vector<int>>(SZ*2);
}
int New(){
    for(int i=0; i<2; i++) G1.emplace_back(), G2.emplace_back();
    return SZ++;
}
inline void AddEdge(int s, int e){ G1[s].push_back(e);
G2[e].push_back(s); }
// T(x) = x << 1, F(x) = x << 1 | 1, I(x) = x ^ 1
inline void AddCNF(int a, int b){ AddEdge(I(a), b);
AddEdge(I(b), a); }
void MostOne(vector<int> vec){
    compress(vec);
    for(int i=0; i<vec.size(); i++){
        int now = New();
        AddEdge(vec[i], T(now)); AddEdge(F(now), I(vec[i]));
        if(i == 0) continue;
        AddEdge(T(now-1), T(now)); AddEdge(F(now), F(now-1));
        AddEdge(T(now-1), I(vec[i])); AddEdge(vec[i], F(now-1));
```



```
}
}
```

3.3 Horn SAT

```
/* n : number of variance
{ }, 0 : x1
{0, 1}, 2 : (x1 and x2) => x3, (-x1 or -x2 or x3)
fail -> empty vector */
vector<int> HornSAT(int n, const vector<vector<int>>> &cond,
const vector<int>> &val){
    int m = cond.size();
    vector<int> res(n, margin(m), stk;
    vector<vector<int>>> gph(n);
    for(int i=0; i<m; i++){
        margin[i] = cond[i].size();
        if(cond[i].empty()) stk.push_back(i);
        for(auto j : cond[i]) gph[j].push_back(i);
    }
    while(!stk.empty()){
        int v = stk.back(); stk.pop_back();
        int h = val[v];
        if(h < 0) return vector<int>();
        if(res[h]) continue; res[h] = 1;
        for(auto i : gph[h]) if(!--margin[i]) stk.push_back(i);
    }
    return res;
}
```

3.4 BCC

Usage: call tarjan() before use

```
vector<int> G[MAX_V]; int In[MAX_V], Low[MAX_V], P[MAX_V];
void addEdge(int s, int e){ G[s].push_back(e);
G[e].push_back(s); }
void tarjan(int n){ /// Pre-Process
    int pv = 0;
    function<void(int,int)> dfs = [&pv,&dfs](int v, int b){
        In[v] = Low[v] = ++pv; P[v] = b;
        for(auto i : G[v]){
            if(i == b) continue;
            if(!In[i]) dfs(i, v), Low[v] = min(Low[v], Low[i]); else
                Low[v] = min(Low[v], In[i]);
        }
    };
    for(int i=1; i<=n; i++) if(!In[i]) dfs(i, -1);
}
vector<int> cutVertex(int n){
    vector<int> res; array<char,MAX_V> isCut; isCut.fill(0);
    function<void(int)> dfs = [&dfs,&isCut](int v){
        int ch = 0;
        for(auto i : G[v]){
            if(P[i] != v) continue; dfs(i); ch++;
            if(P[v] == -1 && ch > 1) isCut[v] = 1; else if(P[v] != -1
                && Low[i] >= In[v]) isCut[v]=1;
        }
    };
};
```

```
for(int i=1; i<=n; i++) if(P[i] == -1) dfs(i);
for(int i=1; i<=n; i++) if(isCut[i]) res.push_back(i);
return move(res);
}
vector<PII> cutEdge(int n){
    vector<PII> res;
    function<void(int)> dfs = [&dfs,&res](int v){
        for(int t=0; t<G[v].size(); t++){
            int i = G[v][t]; if(t != 0 && G[v][t-1] == G[v][t])
                continue;
            if(P[i] != v) continue; dfs(i);
            if((t+1 == G[v].size() || i != G[v][t+1]) && Low[i] >
                In[v]) res.emplace_back(min(v,i), max(v,i));
        }
    };
    for(int i=1; i<=n; i++) sort(G[i].begin(), G[i].end()); //
    multi edge -> sort
    for(int i=1; i<=n; i++) if(P[i] == -1) dfs(i);
    return move(res); // sort(all(res));
}
vector<int> BCC[MAX_V]; // BCC[v] = components which contains v
void vertexDisjointBCC(int n){ // allow multi edge, not allow
self loop
    int cnt = 0; array<char,MAX_V> vis; vis.fill(0);
    function<void(int,int)> dfs = [&dfs,&vis,&cnt](int v, int c){
        vis[v] = 1; if(c > 0) BCC[v].push_back(c);
        for(auto i : G[v]){
            if(vis[i]) continue;
            if(In[v] <= Low[i]) BCC[v].push_back(++cnt), dfs(i, cnt);
            else dfs(i, c);
        }
    };
    for(int i=1; i<=n; i++) if(!vis[i]) dfs(i, 0);
    for(int i=1; i<=n; i++) if(BCC[i].empty())
        BCC[i].push_back(++cnt);
}
void edgeDisjointBCC(int n){ // remove cut edge, do flood fill
```

3.5 Prufer Sequence

```
vector<pair<int,int>> PruferSequence(int n, vector<int> a){ //
a : [1,n]^(n-2)
    if(n == 1) return {}; if(n == 2) return { make_pair(1, 2)
    };
    vector<int> deg(n+1); for(auto i : a) deg[i]++;
    vector<pair<int,int>> res; priority_queue<int> pq;
    for(int i=n; i; i--) if(!deg[i]) pq.emplace(i);
    for(auto i : a){
        res.emplace_back(i, pq.top()); pq.pop();
        if(!--deg[i]) pq.push(i);
    }
    int u = pq.top(); pq.pop(); int v = pq.top(); pq.pop();
    res.emplace_back(u, v); return res;
}
```

3.6 Maximum Clique

```
int N, M; ull G[40], MX, Clique; // 0-index, adj list with
bitset, 0(3~N/3)
void get_clique(int R = 0, ull P = (1ULL<<N)-1, ull X = 0, ull
V=0){
    if((P/X) == 0){ if(R > MX) MX = R, Clique = V; return; }
    int u = __builtin_ctzll(P/X); ll c = P&G[u];
    while(c){
        int v = __builtin_ctzll(c);
        get_clique(R + 1, P&G[v], X&G[v], V | 1ULL << v);
        P ^= 1ULL << v; X |= 1ULL << v; c ^= 1ULL << v;
    }
}
```

3.7 Tree Isomorphism

```
struct Tree{ // (M1,M2)=(1e9+7, 1e9+9), P1,P2 = random int
array(sz >= N+2)
    int N; vector<vector<int>>> G; vector<pair<int,int>>> H;
    vector<int> S, C; // size,centroid
    Tree(int N) : N(N), G(N+2), H(N+2), S(N+2) {}
    void addEdge(int s, int e){ G[s].push_back(e);
G[e].push_back(s); }
    int getCentroid(int v, int b=-1){
        S[v] = 1; // do not merge if-statements
        for(auto i : G[v]) if(i!=b) if(int now=getCentroid(i,v);
            now<=N/2) S[v]+=now; else break;
        if(N - S[v] <= N/2) C.push_back(v); return S[v] = S[v];
    }
    int init(){
        getCentroid(1); if(C.size() == 1) return C[0];
        int u = C[0], v = C[1], add = ++N;
        G[u].erase(find(G[u].begin(), G[u].end(), v));
        G[v].erase(find(G[v].begin(), G[v].end(), u));
        G[add].push_back(u); G[u].push_back(add);
        G[add].push_back(v); G[v].push_back(add);
        return add;
    }
    pair<int,int> build(const vector<ll> &P1, const vector<ll>
&P2, int v, int b=-1){
        vector<pair<int,int>>> ch; for(auto i : G[v]) if(i != b)
            ch.push_back(build(P1, P2, i, v));
        ll h1 = 0, h2 = 0; stable_sort(ch.begin(), ch.end());
        if(ch.empty()){ return {1, 1}; }
        for(int i=0; i<ch.size(); i++)
            h1=(h1+(ch[i].first^P1[P1.size()-1-i])*P1[i])%M1,
            h2=(h2+(ch[i].second^P2[P2.size()-1-i])*P2[i])%M2;
        return H[v] = {h1, h2};
    }
    int build(const vector<ll> &P1, const vector<ll> &P2){
        int rt = init(); build(P1, P2, rt); return rt;
    }
};
```

3.8 Complement Spanning Forest

```
vector<pair<int,int>> ComplementSpanningForest(int n, const
vector<pair<int,int>> &edges){ // V+ElgV
    vector<vector<int>> g(n);
    for(const auto &[u,v] : edges) g[u].push_back(v),
    g[v].push_back(u);
    for(int i=0; i<n; i++) sort(g[i].begin(), g[i].end());
    set<int> alive;
    for(int i=0; i<n; i++) alive.insert(i);
    vector<pair<int,int>> res;
    while(!alive.empty()){
        int u = *alive.begin(); alive.erase(alive.begin());
        queue<int> que; que.push(u);
        while(!que.empty()){
            int v = que.front(); que.pop();
            for(auto it=alive.begin(); it!=alive.end(); ){
                if(auto t=lower_bound(g[v].begin(), g[v].end(), *it); t
                != g[v].end() && *it == *t) ++it;
                else que.push(*it), res.emplace_back(u, *it), it =
                alive.erase(it);
            }
        }
    }
    return res;
}
```

3.9 Bipartite Matching, Konig, Dilworth

```
struct HopcroftKarp{
    int n, m;
    vector<vector<int>> g;
    vector<int> dst, le, ri;
    vector<char> visit, track;
    HopcroftKarp(int n, int m) : n(n), m(m), g(n), dst(n), le(n,
    -1), ri(m, -1), visit(n), track(n+m) {}
    void add_edge(int s, int e){ g[s].push_back(e); }
    bool bfs(){
        bool res = false; queue<int> que;
        fill(dst.begin(), dst.end(), 0);
        for(int i=0; i<n; i++)if(le[i] == -1)que.push(i),dst[i]=1;
        while(!que.empty()){
            int v = que.front(); que.pop();
            for(auto i : g[v]){
                if(ri[i] == -1) res = true;
                else
                    if(!dst[ri[i]])dst[ri[i]]=dst[v]+1,que.push(ri[i]);
            }
        }
        return res;
    }
    bool dfs(int v){
        if(visit[v]) return false; visit[v] = 1;
        for(auto i : g[v]){
            if(ri[i] == -1 || !visit[ri[i]] && dst[ri[i]] == dst[v] +
            1 && dfs(ri[i])){
                le[v] = i; ri[i] = v; return true;
            }
        }
    }
}
```

```
    }
    return false;
}
int maximum_matching(){
    int res = 0; fill(all(le), -1); fill(all(ri), -1);
    while(bfs()){
        fill(visit.begin(), visit.end(), 0);
        for(int i=0; i<n; i++) if(le[i] == -1) res += dfs(i);
    }
    return res;
}
vector<pair<int,int>> maximum_matching_edges(){
    int matching = maximum_matching();
    vector<pair<int,int>> edges; edges.reserve(matching);
    for(int i=0; i<n; i++) if(le[i] != -1)
        edges.emplace_back(i, le[i]);
    return edges;
}
void dfs_track(int v){
    if(track[v]) return; track[v] = 1;
    for(auto i : g[v]) track[n+i] = 1, dfs_track(ri[i]);
}
tuple<vector<int>, vector<int>, int> minimum_vertex_cover(){
    int matching = maximum_matching(); vector<int> lv, rv;
    fill(track.begin(), track.end(), 0);
    for(int i=0; i<n; i++) if(le[i] == -1) dfs_track(i);
    for(int i=0; i<n; i++) if(!track[i]) lv.push_back(i);
    for(int i=0; i<m; i++) if(track[n+i]) rv.push_back(i);
    return {lv, rv, lv.size() + rv.size()}; // s(lv)+s(rv)=mat
}
tuple<vector<int>, vector<int>, int>
maximum_independent_set(){
    auto [a,b,matching] = minimum_vertex_cover();
    vector<int> lv, rv; lv.reserve(n-a.size());
    rv.reserve(m-b.size());
    for(int i=0, j=0; i<n; i++){
        while(j < a.size() && a[j] < i) j++;
        if(j == a.size() || a[j] != i) lv.push_back(i);
    }
    for(int i=0, j=0; i<m; i++){
        while(j < b.size() && b[j] < i) j++;
        if(j == b.size() || b[j] != i) rv.push_back(i);
    } // s(lv)+s(rv)=n+m-mat
    return {lv, rv, lv.size() + rv.size()};
}
vector<vector<int>> minimum_path_cover(){ // n == m
    int matching = maximum_matching();
    vector<vector<int>> res; res.reserve(n - matching);
    fill(track.begin(), track.end(), 0);
    auto get_path = [&](int v) -> vector<int> {
        vector<int> path{v}; // ri[v] == -1
        while(le[v] != -1) path.push_back(v=le[v]);
        return path;
    };
    for(int i=0; i<n; i++) if(!track[n+i] && ri[i] == -1)
        res.push_back(get_path(i));
    return res; // sz(res) = n-mat
}
```

```
    }
    vector<int> maximum_anti_chain(){ // n = m
        auto [a,b,matching] = minimum_vertex_cover();
        vector<int> res; res.reserve(n - a.size() - b.size());
        for(int i=0, j=0, k=0; i<n; i++){
            while(j < a.size() && a[j] < i) j++;
            while(k < b.size() && b[k] < i) k++;
            if((j == a.size() || a[j] != i) && (k == b.size() || b[k]
            != i)) res.push_back(i);
        }
        return res; // sz(res) = n-mat
    }
};
```

3.10 Push Relabel

```
template<typename flow_t> struct Edge {
    int u, v, r; flow_t c, f;
    Edge() = default;
    Edge(int u, int v, flow_t c, int r) : u(u), v(v), r(r), c(c),
    f(0) {}
};
template<typename flow_t, size_t _Sz> struct PushRelabel {
    using edge_t = Edge<flow_t>;
    int n, b, dist[_Sz], count[_Sz+1];
    flow_t excess[_Sz]; bool active[_Sz];
    vector<edge_t> g[_Sz]; vector<int> bucket[_Sz];
    void clear(){ for(int i=0; i<_Sz; i++) g[i].clear(); }
    void addEdge(int s, int e, flow_t x){
        g[s].emplace_back(s, e, x, (int)g[s].size());
        if(s == e) g[s].back().r++;
        g[e].emplace_back(e, s, 0, (int)g[s].size()-1);
    }
    void enqueue(int v){
        if(!active[v] && excess[v] > 0 && dist[v] < n){
            active[v] = true; bucket[dist[v]].push_back(v); b =
            max(b, dist[v]);
        }
    }
    void push(edge_t &e){
        flow_t fl = min(excess[e.u], e.c - e.f);
        if(dist[e.u] == dist[e.v] + 1 && fl > flow_t(0)){
            e.f += fl; g[e.v][e.r].f -= fl; excess[e.u] -= fl;
            excess[e.v] += fl; enqueue(e.v);
        }
    }
    void gap(int k){
        for(int i=0; i<n; i++){
            if(dist[i] >= k) count[dist[i]]--, dist[i] = max(dist[i],
            n), count[dist[i]]++; enqueue(i);
        }
    }
    void relabel(int v){
        count[dist[v]]--; dist[v] = n;
        for(const auto &e : g[v]) if(e.c - e.f > 0) dist[v] =
        min(dist[v], dist[e.v] + 1);
    }
};
```

```
count[dist[v]]++; enqueue(v);
}
void discharge(int v){
    for(auto &e : g[v]) if(excess[v] > 0) push(e); else break;
    if(excess[v] > 0) if(count[dist[v]] == 1) gap(dist[v]);
    else relabel(v);
}
flow_t maximumFlow(int _n, int s, int t){
    memset(dist, 0, sizeof dist); memset(excess, 0, sizeof excess);
    memset(count, 0, sizeof count); memset(active, 0, sizeof active);
    n = _n; b = 0;
    for(auto &e : g[s]) excess[s] += e.c;
    count[s] = n; enqueue(s); active[t] = true;
    while(b >= 0){
        if(bucket[b].empty()) b--;
        else{
            int v = bucket[b].back(); bucket[b].pop_back();
            active[v] = false; discharge(v);
        }
    }
    return excess[t];
}
};
```

3.11 LR Flow

```
addEdge(t, s, inf) // 기존 싱크 -> 기존 소스 inf
addEdge(s, nt, 1) // s -> 새로운 싱크 1
addEdge(ns, e, 1) // 새로운 소스 -> e 1
addEdge(a, b, r-1) // s -> e (r-1)
// ns -> nt의 max flow == 1들의 합 확인
// maxflow : s -> t 플로우 찾을 수 있을 때까지 반복
```

3.12 Hungarian Method

```
// 1-based, only for minimum matching, maximum matching may get TLE
template<typename cost_t=int, cost_t _INF=0x3f3f3f3f>
struct Hungarian{
    int n; vector<vector<cost_t>> mat;
    Hungarian(int n) : n(n), mat(n+1, vector<cost_t>(n+1, _INF)) {}
    void addEdge(int s, int e, cost_t x){ mat[s][e] = min(mat[s][e], x); }
    pair<cost_t, vector<int>> run(){
        vector<cost_t> u(n+1), v(n+1), m(n+1);
        vector<int> p(n+1), w(n+1), c(n+1);
        for(int i=1,a,b; i<=n; i++){
            p[0] = i; b = 0; fill(m.begin(), m.end(), _INF);
            fill(c.begin(), c.end(), 0);
            do{
                int nxt; cost_t delta = _INF; c[b] = 1; a = p[b];
                for(int j=1; j<=n; j++){
                    if(c[j]) continue;
                    cost_t t = mat[a][j] - u[a] - v[j];
```

```
                    if(t < m[j]) m[j] = t, w[j] = b;
                    if(m[j] < delta) delta = m[j], nxt = j;
                }
            }while(p[b] != 0);
            do{ int nxt = w[b]; p[b] = p[nxt]; b = nxt; }while(b != 0);
        }
        vector<int> assign(n+1); for(int i=1; i<=n; i++)
            assign[p[i]] = i;
        return {-v[0], assign};
    }
};
```

3.13 Count/Find 3/4 Cycle

```
vector<tuple<int,int,int>> Find3Cycle(int n, const vector<pair<int,int>> &edges){ // N+MsqrtN
    int m = edges.size();
    vector<int> deg(n), pos(n), ord; ord.reserve(n);
    vector<vector<int>> gph(n), que(m+1), vec(n);
    vector<vector<tuple<int,int,int>>> tri(n);
    vector<tuple<int,int,int>> res;
    for(auto [u,v] : edges) deg[u]++, deg[v]++;
    for(int i=0; i<n; i++) que[deg[i]].push_back(i);
    for(int i=m; i>=0; i--) ord.insert(ord.end(), que[i].begin(), que[i].end());
    for(int i=0; i<n; i++) pos[ord[i]] = i;
    for(auto [u,v] : edges) gph[pos[u]].push_back(pos[v]),
        gph[pos[v]].push_back(pos[u]);
    for(int i=0; i<n; i++){
        for(auto j : gph[i]){
            if(i > j) continue;
            for(int x=0, y=0; x<vec[i].size() && y<vec[j].size(); ){
                if(vec[i][x] == vec[j][y]) res.emplace_back(ord[i], ord[j], ord[vec[i][x]]), x++, y++;
                else if(vec[i][x] < vec[j][y]) x++; else y++;
            }
            vec[j].push_back(i);
        }
    }
    for(auto &[u,v,w] : res){
        if(pos[u] < pos[v]) swap(u, v);
        if(pos[u] < pos[w]) swap(u, w);
        if(pos[v] < pos[w]) swap(v, w);
        tri[u].emplace_back(u, v, w);
    }
    res.clear();
    for(int i=n-1; i>=0; i--) res.insert(res.end(), tri[ord[i]].begin(), tri[ord[i]].end());
    return res;
}
bitset<500> B[500]; // N3/w
```

```
long long Count3Cycle(int n, const vector<pair<int,int>> &edges){
    long long res = 0;
    for(int i=0; i<n; i++) B[i].reset();
    for(auto [u,v] : edges) B[u].set(v), B[v].set(u);
    for(int i=0; i<n; i++) for(int j=i+1; j<n; j++)
        if(B[i].test(j)) res += (B[i] & B[j]).count();
    return res / 3;
}
// 0(n + m * sqrt(m) + th) for graphs without loops or multiedges
void Find4Cycle(int n, const vector<array<int, 2>> &edge, auto process, int th = 1){
    int m = (int)edge.size();
    vector<int> deg(n), order, pos(n);
    vector<vector<int>> appear(m+1), adj(n), found(n);
    for(auto [u, v] : edge) ++deg[u], ++deg[v];
    for(auto u=0; u<n; u++) appear[deg[u]].push_back(u);
    for(auto d=m; d>=0; d--) order.insert(order.end(), appear[d].begin(), appear[d].end());
    for(auto i=0; i<n; i++) pos[order[i]] = i;
    for(auto i=0; i<m; i++){
        int u = pos[edge[i][0]], v = pos[edge[i][1]];
        adj[u].push_back(v), adj[v].push_back(u);
    }
    T res = 0; vector<int> cnt(n);
    for(auto u=0; u<n; u++){
        for(auto v : adj[u]) if(u < v) for(auto w : adj[v]) if(u < w) cnt[w] = 0;
        for(auto v : adj[u]) if(u < v) for(auto w : adj[v]) if(u < w) res += cnt[w] ++;
    }
    for(auto u=0; u<n; u++){
        for(auto v : adj[u]) if(u < v) for(auto w : adj[v]) if(u < w) found[w].clear();
        for(auto v : adj[u]) if(u < v) for(auto w : adj[v]) if(u < w) {
            for(auto x : found[w]){
                if(!th--) return;
                process(order[u], order[v], order[w], order[x]);
            }
            found[w].push_back(v);
        }
    }
}
```

3.14 $O(V^3)$ Global Min Cut

```
int vertex, g[S][S], dst[S], chk[S], del[S];
void init(){
    memset(g, 0, sizeof g); memset(del, 0, sizeof del);
}
void addEdge(int s, int e, int x){ g[s][e] = g[e][s] = x; }
int minCutPhase(int &s, int &t){
    memset(dst, 0, sizeof dst);
    memset(chk, 0, sizeof chk);
    int mincut = 0;
```

```
for(int i=1; i<=vertex; i++){
    int k = -1, mx = -1;
    for(int j=1; j<=vertex; j++) if(!del[j] && !chk[j])
        if(dst[j] > mx) k = j, mx = dst[j];
    if(k == -1) return mincut;
    s = t, t = k;
    mincut = mx, chk[k] = 1;
    for(int j=1; j<=vertex; j++){
        if(!del[j] && !chk[j]) dst[j] += g[k][j];
    }
}
return mincut;
}
int getMinCut(int n){
    vertex = n; int mincut = 1e9+7;
    for(int i=1; i<vertex; i++){
        int s, t;
        int now = minCutPhase(s, t);
        mincut = min(mincut, now); del[t] = 1;
        if(mincut == 0) return 0;
        for(int j=1; j<=vertex; j++){
            if(!del[j]) g[s][j] = (g[j][s] += g[j][t]);
        }
    }
    return mincut;
}
```

3.15 Gomory-Hu Tree

```
// 0-based, S-T cut in graph == S-T cut in gomory-hu tree (path minimum)
vector<Edge> GomoryHuTree(int n, const vector<Edge> &e){
    Dinic<int,100> Flow;
    vector<Edge> res(n-1); vector<int> pr(n);
    for(int i=1; i<n; i++, Flow.clear()){
        for(const auto &[s,e,x] : e) Flow.AddEdge(s, e, x); // bi-directed
        int fl = Flow.MaxFlow(pr[i], i);
        for(int j=i+1; j<n; j++){
            if(!Flow.Level[i] == !Flow.Level[j] && pr[i] == pr[j]) pr[j] = i;
        }
        res[i-1] = Edge(pr[i], i, fl);
    }
    return res;
}
```

3.16 Rectlinear MST

```
template<class T> vector<tuple<T, int, int>>
rectilinear_minimum_spanning_tree(vector<point<T>> a){
    int n = a.size();
    vector<int> ind(n);
    iota(ind.begin(), ind.end(), 0);
    vector<tuple<T, int, int>> edge;
    for(int k=0; k<4; k++){
        sort(ind.begin(), ind.end(), [&](int i,int j){return
            a[i].x-a[j].x < a[j].y-a[i].y;});
```

```
map<T, int> mp;
for(auto i: ind){
    for(auto it=mp.lower_bound(-a[i].y); it!=mp.end();
        it=mp.erase(it)){
        int j = it->second; point<T> d = a[i] - a[j];
        if(d.y > d.x) break;
        edge.push_back({d.x + d.y, i, j});
    }
    mp.insert({-a[i].y, i});
}
for(auto &p: a) if(k & 1) p.x = -p.x; else swap(p.x, p.y);
}
sort(edge.begin(), edge.end());
disjoint_set dsu(n);
vector<tuple<T, int, int>> res;
for(auto [x, i, j]: edge) if(dsu.merge(i, j))
    res.push_back({x, i, j});
return res;
}
```

3.17 $O((V + E) \log V)$ Dominator Tree

```
vector<int> DominatorTree(const vector<vector<int>> &g, int
src){ // // 0-based
    int n = g.size();
    vector<vector<int>> rg(n), buf(n);
    vector<int> r(n), val(n), idom(n, -1), sdom(n, -1), o, p(n),
    u(n);
    iota(all(r), 0); iota(all(val), 0);
    for(int i=0; i<n; i++) for(auto j : g[i]) rg[j].push_back(i);
    function<int(int)> find = [&](int v){
        if(v == r[v]) return v;
        int ret = find(r[v]);
        if(sdom[val[v]] > sdom[val[r[v]]]) val[v] = val[r[v]];
        return r[v] = ret;
    };
    function<void(int)> dfs = [&](int v){
        sdom[v] = o.size(); o.push_back(v);
        for(auto i : g[v]) if(sdom[i] == -1) p[i] = v, dfs(i);
    };
    dfs(src); reverse(all(o));
    for(auto &i : o){
        if(sdom[i] == -1) continue;
        for(auto j : rg[i]){
            if(sdom[j] == -1) continue;
            int x = val[find(j), j];
            if(sdom[i] > sdom[x]) sdom[i] = sdom[x];
        }
        buf[o.size() - sdom[i] - 1].push_back(i);
        for(auto j : buf[p[i]]) u[j] = val[find(j), j];
        buf[p[i]].clear();
        r[i] = p[i];
    }
    reverse(all(o)); idom[src] = src;
    for(auto i : o){ // WARNING : if different, takes idom
        if(i != src) idom[i] = sdom[i] == sdom[u[i]] ? sdom[i] :
            idom[u[i]];
```

```

    }
    for(auto i : o) if(i != src) idom[i] = o[idom[i]];
    return idom; // unreachable -> ret[i] = -1
}
```

3.18 $O(N^2)$ Stable Marriage Problem

```
// man : 1~n, woman : n+1~2n
struct StableMarriage{
    int n; vector<vector<int>> g;
    StableMarriage(int n) : n(n), g(2*n+1) { for(int i=1; i<=n+n;
        i++) g[i].reserve(n); }
    void addEdge(int u, int v){ g[u].push_back(v); } // insert
    in decreasing order of preference.
    vector<int> run(){
        queue<int> q; vector<int> match(2*n+1), ptr(2*n+1);
        for(int i=1; i<=n; i++) q.push(i);
        while(q.size()){
            int i = q.front(); q.pop();
            for(int &p=ptr[i]; p<g[i].size(); p++){
                int j = g[i][p];
                if(!match[j]){ match[i] = j; match[j] = i; break; }
                int m = match[j], u = -1, v = -1;
                for(int k=0; k<g[j].size(); k++){
                    if(g[j][k] == i) u = k; if(g[j][k] == m) v = k;
                }
                if(u < v){
                    match[m] = 0; q.push(m); match[i] = j; match[j] = i;
                    break;
                }
            }
            ptr[i]++;
        }
        return match;
    }
};
```

3.19 $O(VE)$ Vizing Theorem

```
// Graph coloring with (max-degree)+1 colors,  $O(N^2)$ 
int C[MX][MX] = {}, G[MX][MX] = {}; // MX ~= 2500
void solve(vector<pii> &E, int N, int M){
    int X[MX] = {}, a, b;
    auto update = [&](int u){ for(X[u] = 1; C[u][X[u]]; X[u]++);
    };
    auto color = [&](int u, int v, int c){
        int p = G[u][v]; G[u][v] = G[v][u] = c;
        C[u][c] = v; C[v][c] = u; C[u][p] = C[v][p] = 0;
        if( p ) X[u] = X[v] = p; else update(u), update(v);
        return p; }; // end of function : color
    auto flip = [&](int u, int c1, int c2){
        int p = C[u][c1], q = C[u][c2];
        swap(C[u][c1], C[u][c2]);
        if( p ) G[u][p] = G[p][u] = c2;
        if( !C[u][c1] ) X[u] = c1; if( !C[u][c2] ) X[u] = c2;
        return p; }; // end of function : flip
    for(int i = 1; i <= N; i++) X[i] = 1;
```



```

for(int t = 0; t < E.size(); t++){
    int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c
    = c0, d;
    vector<pii> L; int vst[MX] = {};
    while(!G[u][v0]){
        L.emplace_back(v, d = X[v]);
        if(!C[v][c]) for(a = (int)L.size()-1; a >= 0; a--) c =
        color(u, L[a].first, c);
        else if(!C[u][d])for(a=(int)L.size()-1;a>=0;a--)
        color(u,L[a].first,L[a].second);
        else if( vst[d] ) break;
        else vst[d] = 1, v = C[u][d];
    }
    if( !G[u][v0] ){
        for(;v; v = flip(v, c, d), swap(c, d));
        if(C[u][c0]){
            for(a = (int)L.size()-2; a >= 0 && L[a].second != c;
            a--);
            for(; a >= 0; a--) color(u, L[a].first, L[a].second);
        } else t--;
    }
}
}
}

```

3.20 $O(E \log V)$ Directed MST

```

struct Edge{
    int s, e; cost_t x;
    Edge() = default;
    Edge(int s, int e, cost_t x) : s(s), e(e), x(x) {}
    bool operator < (const Edge &t) const { return x < t.x; }
};

struct UnionFind{
    vector<int> P, S;
    vector<pair<int, int>> stk;
    UnionFind(int n) : P(n), S(n, 1) { iota(P.begin(), P.end(),
    0); }
    int find(int v) const { return v == P[v] ? v : find(P[v]); }
    int time() const { return stk.size(); }
    void rollback(int t){
        while(stk.size() > t){
            auto [u,v] = stk.back(); stk.pop_back();
            P[u] = u; S[v] -= S[u];
        }
    }
    bool merge(int u, int v){
        u = find(u); v = find(v);
        if(u == v) return false;
        if(S[u] > S[v]) swap(u, v);
        stk.emplace_back(u, v);
        S[v] += S[u]; P[u] = v;
        return true;
    }
};

struct Node{
    Edge key;
    Node *l, *r;

```

```

    cost_t lz;
    Node() : Node(Edge()) {}
    Node(const Edge &edge) : key(edge), l(nullptr), r(nullptr),
    lz(0) {}
    void push(){
        key.x += lz;
        if(l) l->lz += lz;
        if(r) r->lz += lz;
        lz = 0;
    }
    Edge top(){ push(); return key; }
};

Node* merge(Node *a, Node *b){
    if(!a || !b) return a ? a : b;
    a->push(); b->push();
    if(b->key < a->key) swap(a, b);
    swap(a->l, (a->r = merge(b, a->r)));
    return a;
}

void pop(Node* &a){ a->push(); a = merge(a->l, a->r); }

// 0-based
pair<cost_t, vector<int>> DirectMST(int n, int rt, vector<Edge>
&edges){
    vector<Node*> heap(n);
    UnionFind uf(n);
    for(const auto &i : edges) heap[i.e] = merge(heap[i.e], new
    Node(i));
    cost_t res = 0;
    vector<int> seen(n, -1), path(n), par(n);
    seen[rt] = rt;
    vector<Edge> Q(n), in(n, {-1,-1, 0}), comp;
    deque<tuple<int, int, vector<Edge>>> cyc;
    for(int s=0; s<n; s++){
        int u = s, qi = 0, w;
        while(seen[u] < 0){
            if(!heap[u]) return {-1, {}};
            Edge e = heap[u]->top();
            heap[u]->lz -= e.x; pop(heap[u]);
            Q[qi] = e; path[qi++] = u; seen[u] = s;
            res += e.x; u = uf.find(e.s);
            if(seen[u] == s){ // found cycle, contract
                Node* nd = 0;
                int end = qi, time = uf.time();
                do nd = merge(nd, heap[w = path[--qi]]);
                while(uf.merge(u, w));
                u = uf.find(u); heap[u] = nd; seen[u] = -1;
                cyc.emplace_front(u, time, vector<Edge>{&Q[qi],
                &Q[end]});
            }
        }
        for(int i=0; i<qi; i++) in[uf.find(Q[i].e)] = Q[i];
    }
    for(auto& [u,t,comp] : cyc){
        uf.rollback(t);
        Edge inEdge = in[u];
        for (auto& e : comp) in[uf.find(e.e)] = e;
    }
}

```

```

        in[uf.find(inEdge.e)] = inEdge;
    }
    for(int i=0; i<n; i++) par[i] = in[i].s;
    return {res, par};
}

```

3.21 $O(E \log V + K \log K)$ K Shortest Path

```

int rnd(int l, int r){ /* return random int [l,r] */ }
struct node{
    array<node*, 2> son; pair<ll, ll> val;
    node() : node(make_pair(-1e18, -1e18)) {}
    node(pair<ll, ll> val) : node(nullptr, nullptr, val) {}
    node(node *l, node *r, pair<ll, ll> val) : son({l,r}),
    val(val) {}
};
node* copy(node *x){ return x ? new node(x->son[0], x->son[1],
x->val) : nullptr; }
node* merge(node *x, node *y){ // precondition: x, y both
points to new entity
    if(!x || !y) return x ? x : y;
    if(x->val > y->val) swap(x, y);
    int rd = rnd(0, 1);
    if(x->son[rd]) x->son[rd] = copy(x->son[rd]);
    x->son[rd] = merge(x->son[rd], y); return x;
}

struct edge{
    ll v, c, i; edge() = default;
    edge(ll v, ll c, ll i) : v(v), c(c), i(i) {}
};
vector<vector<edge>> gph, rev;
int idx;
void init(int n){ gph = rev = vector<vector<edge>>(n); idx = 0;
}
void add_edge(int s, int e, ll x){
    gph[s].emplace_back(e, x, idx);
    rev[e].emplace_back(s, x, idx);
    assert(x >= 0); idx++;
}

vector<int> par, pae; vector<ll> dist; vector<node*> heap;
void dijkstra(int snk){ // replace this to SPFA if edge weight
is negative
    int n = gph.size();
    par = pae = vector<int>(n, -1);
    dist = vector<ll>(n, 0x3f3f3f3f3f3f3f3f);
    heap = vector<node*>(n, nullptr);
    priority_queue<pair<ll,ll>, vector<pair<ll,ll>>, greater<>>
    pq;
    auto enqueue = [&](int v, ll c, int pa, int pe){
        if(dist[v] > c) dist[v] = c, par[v] = pa, pae[v] = pe,
        pq.emplace(c, v);
    }; enqueue(snk, 0, -1, -1); vector<int> ord;
    while(!pq.empty()){
        auto [c,v] = pq.top(); pq.pop(); if(dist[v] != c) continue;
        ord.push_back(v); for(auto e : rev[v]) enqueue(e.v, c+e.c,
        v, e.i);
    }
}

```

```
for(auto &v : ord){
    if(par[v] != -1) heap[v] = copy(heap[par[v]]);
    for(auto &e : gph[v]){
        if(e.i == pae[v]) continue;
        ll delay = dist[e.v] + e.c - dist[v];
        if(delay < 1e18) heap[v] = merge(heap[v], new
            node(make_pair(delay, e.v)));
    }
}
vector<ll> run(int s, int e, int k){
    using state = pair<ll, node*>; dijkstra(e); vector<ll> ans;
    priority_queue<state, vector<state>, greater<state>> pq;
    if(dist[s] > 1e18) return vector<ll>(k, -1);
    ans.push_back(dist[s]);
    if(heap[s]) pq.emplace(dist[s] + heap[s]->val.first,
        heap[s]);
    while(!pq.empty() && ans.size() < k){
        auto [cst, ptr] = pq.top(); pq.pop(); ans.push_back(cst);
        for(int j=0; j<2; j++) if(ptr->son[j])
            pq.emplace(cst-ptr->val.first + ptr->son[j]->val.first,
                ptr->son[j]);
        int v = ptr->val.second;
        if(heap[v]) pq.emplace(cst + heap[v]->val.first, heap[v]);
    }
    while(ans.size() < k) ans.push_back(-1);
    return ans;
}
```

3.22 Chordal Graph, Tree Decomposition

```
struct Set {
    list<int> L; int last;
    Set() { last = 0; }
};
struct PEO {
    int N;
    vector<vector<int>> > g;
    vector<int> vis, res;
    list<Set> L;
    vector<list<Set>::iterator> ptr;
    vector<list<int>::iterator> ptr2;
    PEO(int n, vector<vector<int>> > _g) {
        N = n; g = _g;
        for (int i = 1; i <= N; i++) sort(g[i].begin(),
            g[i].end());
        vis.resize(N + 1); ptr.resize(N + 1); ptr2.resize(N + 1);
        L.push_back(Set());
        for (int i = 1; i <= N; i++) {
            L.back().L.push_back(i);
            ptr[i] = L.begin(); ptr2[i] = prev(L.back().L.end());
        }
    }
    pair<bool, vector<int>> Run() {
        // lexicographic BFS
        int time = 0;
        while (!L.empty()) {
```

```
if (L.front().L.empty()) { L.pop_front(); continue; }
        auto it = L.begin();
        int n = it->L.front(); it->L.pop_front();
        vis[n] = ++time;
        res.push_back(n);
        for (int next : g[n]) {
            if (vis[next]) continue;
            if (ptr[next]->last != time) {
                L.insert(ptr[next], Set()); ptr[next]->last = time;
            }
            ptr[next]->L.erase(ptr2[next]); ptr[next]--;
            ptr[next]->L.push_back(next);
            ptr2[next] = prev(ptr[next]->L.end());
        }
    }
    // PEO existence check
    for (int n = 1; n <= N; n++) {
        int mx = 0;
        for (int next : g[n]) if (vis[n] > vis[next]) mx =
            max(mx, vis[next]);
        if (mx == 0) continue;
        int w = res[mx - 1];
        for (int next : g[n]) {
            if (vis[w] > vis[next] && !binary_search(g[w].begin(),
                g[w].end(), next)){
                vector<int> chk(N+1), par(N+1, -1); // w와 next가
                이어져 있지 않다면 not chordal
                deque<int> dq{next}; chk[next] = 1;
                while (!dq.empty()) {
                    int x = dq.front(); dq.pop_front();
                    for (auto y : g[x]) {
                        if (chk[y] || y == n || y != w &&
                            binary_search(g[n].begin(), g[n].end(), y))
                            continue;
                        dq.push_back(y); chk[y] = 1; par[y] = x;
                    }
                }
                vector<int> cycle{next, n};
                for (int x=w; x!=next; x=par[x]) cycle.push_back(x);
                return {false, cycle};
            }
        }
    }
    reverse(res.begin(), res.end());
    return {true, res};
}
};
bool vis[200201]; // 배열 크기 알아서 수정하자.
int p[200201], ord[200201], P = 0; // P=경점 개수
vector<int> V[200201], G[200201]; // V=bags, G=edges
void tree_decomposition(int N, vector<vector<int>> > g) {
    for(int i=1; i<=N; i++) sort(g[i].begin(), g[i].end());
    vector<int> peo = PEO(N, g).Run(), rpeo = peo;
    reverse(rpeo.begin(), rpeo.end());
    for(int i=0; i<peo.size(); i++) ord[peo[i]] = i;
    for(int n : rpeo) { // tree decomposition
        vis[n] = true;
```

```
if (n == rpeo[0]) { // 처음
        P++; V[P].push_back(n); p[n] = P; continue;
    }
    int mn = INF, idx = -1;
    for(int next : g[n]) if (vis[next] && mn > ord[next]) mn =
        ord[next], idx = next;
    assert(idx != -1); idx = p[idx];
    // 두 set인 V[idx]와 g[n](visited ver)가 같나?
    // V[idx]의 모든 원소가 g[n]에서 나타나는지 판별로 충분하다.
    int die = 0;
    for(int x : V[idx]) {
        if (!binary_search(g[n].begin(), g[n].end(), x)) { die =
            1; break; }
    }
    if (!die) { V[idx].push_back(n), p[n] = idx; } // 기존
    집합에 추가
    else { // 새로운 집합을 자식으로 추가
        P++;
        G[idx].push_back(P); // 자식으로만 단방향으로 잇자.
        V[P].push_back(n);
        for(int next : g[n]) if (vis[next]) V[P].push_back(next);
        p[n] = P;
    }
}
for(int i=1; i<=P; i++) sort(V[i].begin(), V[i].end());
}
```

3.23 $O(V^3)$ General Matching

```
int N, M, R, Match[555], Par[555], Chk[555], Prv[555],
Vis[555];
vector<int> G[555];
int Find(int x){ return x == Par[x] ? x : Par[x] =
    Find(Par[x]); }
int LCA(int u, int v){ static int cnt = 0;
    for(cnt++; Vis[u]!=cnt; swap(u, v)) if(u) Vis[u] = cnt, u =
        Find(Prv[Match[u]]);
    return u;
}
void Blossom(int u, int v, int rt, queue<int> &q){
    for(; Find(u)!=rt; u=Prv[v]){
        Prv[u] = v; Par[u] = Par[v=Match[u]] = rt; if(Chk[v] & 1)
            q.push(v), Chk[v] = 2;
    }
}
bool Augment(int u){
    iota(Par, Par+555, 0); memset(Chk, 0, sizeof Chk); queue<int>
    Q; Q.push(u); Chk[u] = 2;
    while(!Q.empty()){
        u = Q.front(); Q.pop();
        for(auto v : G[u]){
            if(Chk[v] == 0){
                Prv[v] = u; Chk[v] = 1; Q.push(Match[v]); Chk[Match[v]]
                    = 2;
                if(!Match[v]){ for(; u; v=u) u = Match[Prv[v]],
                    Match[Match[v]=Prv[v]] = v; return true; }
            }
```

```
    }
    else if(Chk[v] == 2){ int l = LCA(u, v); Blossom(u, v, l,
    Q), Blossom(v, u, l, Q); }
}
return 0;
}
void Run(){ for(int i=1; i<=N; i++) if(!Match[i]) R +=
Augment(i); }
```

3.24 $O(V^3)$ Weighted General Matching

```
namespace weighted_blossom_tree{
#define d(x) (lab[x.u]+lab[x.v]-e[x.u][x.v].w*2)
const int N=403*2; using ll = long long; using T = int; //
sum of weight, single weight
const T inf=numeric_limits<T>::max()-1;
struct Q{ int u, v; T w; } e[N][N]; vector<int> p[N];
int n, m=0, id, h, t, lk[N], sl[N], st[N], f[N], b[N][N],
s[N], ed[N], q[N]; T lab[N];
void upd(int u, int v){ if (!sl[v] || d(e[u][v]) <
d(e[sl[v]][v])) sl[v] = u; }
void ss(int v){
    sl[v]=0; for(int u=1; u<=n; u++) if(e[u][v].w > 0 && st[u]
    != v && !s[st[u]]) upd(u, v);
}
void ins(int u){ if(u <= n) q[++t] = u; else for(int v :
p[u]) ins(v); }
void mdf(int u, int w){ st[u]=w; if(u > n) for(int v : p[u])
mdf(v, w); }
int gr(int u,int v){
    if ((v=find(p[u].begin(), p[u].end(), v) - p[u].begin()) &
    1){
        reverse(p[u].begin()+1, p[u].end()); return
        (int)p[u].size() - v;
    }
    return v;
}
void stm(int u, int v){
    lk[u] = e[u][v].v;
    if(u <= n) return; Q w = e[u][v];
    int x = b[u][w.u], y = gr(u,x);
    for(int i=0; i<y; i++) stm(p[u][i], p[u][i^1]);
    stm(x, v); rotate(p[u].begin(), p[u].begin()+y,
    p[u].end());
}
void aug(int u, int v){
    int w = st[lk[u]]; stm(u, v); if (!w) return;
    stm(w, st[f[w]]); aug(st[f[w]], w);
}
int lca(int u, int v){
    for(++id; u|v; swap(u, v)){
        if(!u) continue; if(ed[u] == id) return u;
        ed[u] = id; if(u = st[lk[u]]) u = st[f[u]]; // not ==
    }
    return 0;
}
}
```

```
void add(int u, int a, int v){
    int x = n+1; while(x <= m && st[x]) x++;
    if(x > m) m++;
    lab[x] = s[x] = st[x] = 0; lk[x] = lk[a];
    p[x].clear(); p[x].push_back(a);
    for(int i=u, j; i!=a; i=st[f[j]]) p[x].push_back(i),
    p[x].push_back(j=st[lk[i]]), ins(j);
    reverse(p[x].begin()+1, p[x].end());
    for(int i=v, j; i!=a; i=st[f[j]]) p[x].push_back(i),
    p[x].push_back(j=st[lk[i]]), ins(j);
    mdf(x, x); for(int i=1; i<=m; i++) e[x][i].w = e[i][x].w =
    0;
    memset(b[x+1, 0, n*sizeof b[0][0]);
    for (int u : p[x]){
        for(v=1; v<=m; v++) if(!e[x][v].w || d(e[u][v]) <
        d(e[x][v])) e[x][v] = e[u][v],e[v][x] = e[v][u];
        for(v=1; v<=n; v++) if(b[u][v]) b[x][v] = u;
    }
    ss(x);
}
void ex(int u){ // s[u] == 1
    for(int x : p[u]) mdf(x, x);
    int a = b[u][e[u][f[u]].u],r = gr(u, a);
    for(int i=0; i<r; i+=2){
        int x = p[u][i], y = p[u][i+1];
        f[x] = e[y][x].u; s[x] = 1; s[y] = 0; sl[x] = 0; ss(y);
        ins(y);
    }
    s[a] = 1; f[a] = f[u];
    for(int i=r+1; i<p[u].size(); i++) s[p[u][i]] = -1,
    ss(p[u][i]);
    st[u] = 0;
}
bool on(const Q &e){
    int u=st[e.u], v=st[e.v], a;
    if(s[v] == -1) f[v] = e.u, s[v] = 1, a = st[lk[v]], sl[v] =
    sl[a] = s[a] = 0, ins(a);
    else if(!s[v]){
        a = lca(u, v); if(!a) return aug(u,v), aug(v,u), true;
        else add(u,a,v);
    }
    return false;
}
bool bfs(){
    memset(s+1, -1, m*sizeof s[0]); memset(sl+1, 0, m*sizeof
    sl[0]);
    h = 1; t = 0; for(int i=1; i<=m; i++) if(st[i] == i &&
    !lk[i]) f[i] = s[i] = 0, ins(i);
    if(h > t) return 0;
    while (true){
        while (h <= t){
            int u = q[h++];
            if (s[st[u]] != 1) for (int v=1; v<=n; v++) if
            (e[u][v].w > 0 && st[u] != st[v])
                if(d(e[u][v])) upd(u, st[v]); else if(on(e[u][v]))
                return true;
        }
    }
}
```

```
T x = inf;
for(int i=n+1; i<=m; i++) if(st[i] == i && s[i] == 1) x =
min(x, lab[i]>>1);
for(int i=1; i<=m; i++) if(st[i] == i && sl[i] && s[i] !=
1) x = min(x, d(e[sl[i]][i])>>s[i]+1);
for(int i=1; i<=n; i++) if(~s[st[i]]) if((lab[i] +=
(s[st[i]]*2-1)*x) <= 0) return false;
for(int i=n+1; i<=m; i++) if(st[i] == i && ~s[st[i]])
lab[i] += (2-s[st[i]]*4)*x;
h = 1; t = 0;
for(int i=1; i<=m; i++) if(st[i] == i && sl[i] &&
st[sl[i]] != i && !d(e[sl[i]][i]) && on(e[sl[i]][i]))
return true;
for(int i=n+1; i<=m; i++) if(st[i] == i && s[i] == 1 &&
!lab[i]) ex(i);
}
return 0;
}
template<typename TT> pair<int,ll> run(int N, const
vector<tuple<int,int,TT>> &edges){ // 1-based
    memset(ed+1, 0, m*sizeof ed[0]); memset(lk+1, 0, m*sizeof
    lk[0]);
    n = m = N; id = 0; iota(st+1, st+n+1, 1); T wm = 0; ll r =
    0;
    for(int i=1; i<=n; i++) for(int j=1; j<=n; j++) e[i][j] =
    {i,j,0};
    for(auto [u,v,w] : edges) wm = max(wm,
    e[v][u].w=e[u][v].w=max(e[u][v].w,(T)w));
    for(int i=1; i<=n; i++) p[i].clear();
    for(int i=1; i<=n; i++) for (int j=1; j<=n; j++) b[i][j] =
    i*(i==j);
    fill_n(lab+1, n, wm); int match = 0; while(bfs()) match++;
    for(int i=1; i<=n; i++) if(lk[i]) r += e[i][lk[i]].w;
    return {match, r/2};
}
#undef d
} using weighted_blossom_tree::run, weighted_blossom_tree::lk;
```

4 Math

4.1 Extend GCD, CRT, Combination

```
// ll gcd(ll a, ll b), ll lcm(ll a, ll b), ll mod(ll a, ll b)
tuple<ll,ll,ll> ext_gcd(ll a, ll b){ // return [g,x,y] s.t.
ax+by=gcd(a,b)=g
    if(b == 0) return {a, 1, 0}; auto [g,x,y] = ext_gcd(b, a %
    b); return {g, y, x - a/b * y};
}
ll inv(ll a, ll m){ //return x when ax mod m = 1, fail -> -1
    auto [g,x,y] = ext_gcd(a, m); return g == 1 ? mod(x, m) : -1;
}
void DivList(ll n){ // {n/1, n/2, ... , n/n}, size <= 2 sqrt n
    for(ll i=1, j=1; i<=n; i=j+1) cout << i << " " << (j=n/(n/i))
    << " " << n/i << "\n";
}
pair<ll,ll> crt(ll a1, ll m1, ll a2, ll m2){
```

```
ll g = gcd(m1, m2), m = m1 / g * m2;
if((a2 - a1) % g) return {-1, -1};
ll md = m2/g, s = mod((a2-a1)/g, m2/g);
ll t = mod(get<1>(ext_gcd(m1/g%md, m2/g)), md);
return { a1 + s * t % md * m1, m };
}

pair<ll,ll> crt(const vector<ll> &a, const vector<ll> &m){
    ll ra = a[0], rm = m[0];
    for(int i=1; i<m.size(); i++){
        auto [aa,mm] = crt(ra, rm, a[i], m[i]);
        if(mm == -1) return {-1, -1}; else tie(ra,rm) = tie(aa,mm);
    }
    return {ra, rm};
}

struct Lucas{ // init : O(P), query : O(log P)
    const size_t P;
    vector<ll> fac, inv;
    ll Pow(ll a, ll b){ /* return a^b mod P */ }
    Lucas(size_t P) : P(P), fac(P), inv(P) {
        fac[0] = 1; for(int i=1; i<P; i++) fac[i] = fac[i-1] * i % P;
        inv[P-1] = Pow(fac[P-1], P-2); for(int i=P-2; ~i; i--)
            inv[i] = inv[i+1] * (i+1) % P;
    }
    ll small(ll n, ll r) const { return r <= n ? fac[n] * inv[r] % P * inv[n-r] % P : 0LL; }
    ll calc(ll n, ll r) const {
        if(n < r || n < 0 || r < 0) return 0;
        if(!n || !r || n == r) return 1; else return small(n%P, r%P) * calc(n/P, r/P) % P;
    }
};

template<ll p, ll e> struct CombinationPrimePower{ // init : O(p^e), query : O(log p)
    vector<ll> val; ll m;
    CombinationPrimePower(){
        m = 1; for(int i=0; i<e; i++) m *= p; val.resize(m); val[0] = 1;
        for(int i=1; i<m; i++) val[i] = val[i-1] * (i % p ? i : 1) % m;
    }
    pair<ll,ll> factorial(int n){
        if(n < p) return {0, val[n]};
        int k = n / p; auto v = factorial(k);
        int cnt = v.first + k, kp = n / m, rp = n % m;
        ll ret = v.second * Pow(val[m-1], kp % 2, m) % m * val[rp] % m;
        return {cnt, ret};
    }
    ll calc(int n, int r){
        if(n < 0 || r < 0 || n < r) return 0;
        auto v1 = factorial(n), v2 = factorial(r), v3 = factorial(n-r);
        ll cnt = v1.first - v2.first - v3.first;
        ll ret = v1.second * inv(v2.second, m) % m * inv(v3.second, m) % m;
        if(cnt >= e) return 0;
    }
};

ll calc(int n, int r){
    if(n < 0 || r < 0 || n < r) return 0;
    auto v1 = factorial(n), v2 = factorial(r), v3 = factorial(n-r);
    ll cnt = v1.first - v2.first - v3.first;
    ll ret = v1.second * inv(v2.second, m) % m * inv(v3.second, m) % m;
    if(cnt >= e) return 0;
}
```

```
for(int i=1; i<=cnt; i++) ret = ret * p % m;
return ret;
}
};
```

4.2 Diophantine

```
// solutions to ax + by = c where x in [xlow, xhigh] and y in [ylow, yhigh]
// cnt, leftsol, rightsol, gcd of a and b
template<class T> array<T, 6> solve_linear_diophantine(T a, T b, T c, T xlow, T xhigh, T ylow, T yhigh){
    T g, x, y = euclid(a >= 0 ? a : -a, b >= 0 ? b : -b, x, y);
    array<T, 6> no_sol{0, 0, 0, 0, 0, g};
    if(c % g) return no_sol; x *= c / g, y *= c / g;
    if(a < 0) x = -x; if(b < 0) y = -y;
    a /= g, b /= g, c /= g;
    auto shift = [&](T &x, T &y, T a, T b, T cnt){ x += cnt * b, y -= cnt * a; };
    int sign_a = a > 0 ? 1 : -1, sign_b = b > 0 ? 1 : -1;
    shift(x, y, a, b, (xlow - x) / b);
    if(x < xlow) shift(x, y, a, b, sign_b);
    if(x > xhigh) return no_sol;
    T lx1 = x; shift(x, y, a, b, (xhigh - x) / b);
    if(x > xhigh) shift(x, y, a, b, -sign_b);
    T rx1 = x; shift(x, y, a, b, -(ylow - y) / a);
    if(y < ylow) shift(x, y, a, b, -sign_a);
    if(y > yhigh) return no_sol;
    T lx2 = x; shift(x, y, a, b, -(yhigh - y) / a);
    if(y > yhigh) shift(x, y, a, b, sign_a);
    T rx2 = x; if(lx2 > rx2) swap(lx2, rx2);
    T lx = max(lx1, lx2), rx = min(rx1, rx2);
    if(lx > rx) return no_sol;
    return {(rx - lx) / (b >= 0 ? b : -b) + 1, lx, (c - lx * a) / b, rx, (c - rx * a) / b, g};
}
```

4.3 Partition Number

```
for(int j=1; j*(3*j-1)/2<=i; j++) P[i] += (j%2?-1)*P[i-j*(3*j-1)/2], P[i] %= MOD;
for(int j=1; j*(3*j+1)/2<=i; j++) P[i] += (j%2?-1)*P[i-j*(3*j+1)/2], P[i] %= MOD;
```

4.4 FloorSum

```
// sum of floor((A*i+B)/M) over 0 <= i < N in O(log(N+M+A+B))
ll FloorSum(ll N, ll M, ll A, ll B){ // 1 <= N,M <= 1e9, 0 <= A,B < M
    ll R = 0;
    if(A >= M) R += N * (N - 1) / 2 * (A / M), A %= M;
    if(B >= M) R += B / M * N, B %= M;
    ll Y = (A * N + B) / M, X = Y * M - B;
    if(Y == 0) return R;
    R += (N - (X + A - 1) / A) * Y;
    R += FloorSum(Y, A, M, (A - X % A) % A);
    return R;
}
```

4.5 XOR Basis(XOR Maximization)

```
vector<ll> basis; // ascending
for(int i=0; i<n; i++){
    ll x; cin >> x;
    for(int j=(int)basis.size()-1; j>=0; j--) x = min(x, basis[j]^x);
    if(x) basis.insert(lower_bound(basis.begin(), basis.end(), x), x);
} // if xor maximization, reverse -> for(auto i:basis) r = max(r,r^i);
```

4.6 Stern Brocot Tree

```
pair<ll,ll> Solve(ld l, ld r){ // find 1 < p/q < r -> min q -> min p
    auto g = [](ll v, pair<ll,ll> a, pair<ll,ll> b) -> pair<ll,ll> {
        return { v * a.first + b.first, v * a.second + b.second };
    };
    auto f = [g](ll v, pair<ll,ll> a, pair<ll,ll> b) -> ld {
        auto [p,q] = g(v, a, b); return ld(p) / q;
    };
    pair<ll,ll> s(0, 1), e(1, 0);
    while(true){
        pair<ll,ll> m(s.first+e.first, s.second+e.second);
        ld v = 1.L * m.first / m.second;
        if(v >= r){
            ll ks = 1, ke = 1; while(f(ke, s, e) >= r) ke *= 2;
            while(ks <= ke){
                ll km = (ks + ke) / 2;
                if(f(km, s, e) >= r) ks = km + 1; else ke = km - 1;
            } e = g(ke, s, e);
        }
        else if(v <= l){
            ll ks = 1, ke = 1; while(f(ke, e, s) <= l) ke *= 2;
            while(ks <= ke){
                ll km = (ks + ke) / 2;
                if(f(km, e, s) <= l) ks = km + 1; else ke = km - 1;
            } s = g(ke, e, s);
        }
        else return m;
    }
}
```

4.7 Gauss Jordan Elimination

```
template<typename T> // return {rref, rank, det, inv}
tuple<vector<vector<T>>, T, T, vector<vector<T>>>>
Gauss(vector<vector<T>> a, bool square=true){
    int n = a.size(), m = a[0].size(), rank = 0;
    vector<vector<T>> out(n, vector<T>(m, 0)); T det = T(1);
    for(int i=0; i<n; i++) if(square) out[i][i] = T(1);
    for(int i=0; i<m; i++){
        if(rank == n) break;
```



```
if(IsZero(a[rank][i])){
    T mx = T(0); int idx = -1; // fucking precision error
    for(int j=rank+1; j<n; j++) if(mx < abs(a[j][i])) mx =
        abs(a[j][i]), idx = j;
    if(idx == -1 || IsZero(a[idx][i])){ det = 0; continue; }
    for(int k=0; k<m; k++){
        a[rank][k] = Add(a[rank][k], a[idx][k]);
        if(square) out[rank][k] = Add(out[rank][k],
            out[idx][k]);
    }
    det = Mul(det, a[rank][i]);
    T coeff = Div(T(1), a[rank][i]);
    for(int j=0; j<m; j++) a[rank][j] = Mul(a[rank][j], coeff);
    for(int j=0; j<m; j++) if(square) out[rank][j] =
        Mul(out[rank][j], coeff);
    for(int j=0; j<n; j++){
        if(rank == j) continue;
        T t = a[j][i]; // Warning: [j][k], [rank][k]
        for(int k=0; k<m; k++) a[j][k] = Sub(a[j][k],
            Mul(a[rank][k], t));
        for(int k=0; k<m; k++) if(square) out[j][k] =
            Sub(out[j][k], Mul(out[rank][k], t));
    }
    rank++;
}
return {a, rank, det, out};
}
```

4.8 Berlekamp + Kitamasa

Time Complexity: $O(NK + N \log \text{mod})$, $O(N^2 \log X)$

```
const int mod = 1e9+7; ll pw(ll a, ll b){ /* return a^b mod m
*/ }
vector<int> berlekamp_massey(vector<int> x){
    vector<int> ls, cur; int lf, ld;
    for(int i=0; i<x.size(); i++){
        ll t = 0;
        for(int j=0; j<cur.size(); j++) t = (t + 111 * x[i-j-1] *
            cur[j]) % mod;
        if((t - x[i]) % mod == 0) continue;
        if(cur.empty()){ cur.resize(i+1); lf = i; ld = (t - x[i]) %
            mod; continue; }
        ll k = -(x[i] - t) * pw(ld, mod - 2) % mod;
        vector<int> c(i-lf-1); c.push_back(k);
        for(auto &j : ls) c.push_back(-j * k % mod);
        if(c.size() < cur.size()) c.resize(cur.size());
        for(int j=0; j<cur.size(); j++) c[j] = (c[j] + cur[j]) %
            mod;
        if(i-lf+(int)ls.size())>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) % mod);
        }
        cur = c;
    }
    for(auto &i : cur) i = (i % mod + mod) % mod; return cur;
}
int get_nth(vector<int> rec, vector<int> dp, ll n){
```

```
int m = rec.size(); vector<int> s(m), t(m);
s[0] = 1; if(m != 1) t[1] = 1; else t[0] = rec[0];
auto mul = [&rec](vector<int> v, vector<int> w){
    int m = v.size();
    vector<int> t(2 * m);
    for(int j=0; j<m; j++) for(int k=0; k<m; k++){
        t[j+k] += 111 * v[j] * w[k] % mod;
        if(t[j+k] >= mod) t[j+k] -= mod;
    }
    for(int j=2*m-1; j>=m; j--) for(int k=1; k<=m; k++){
        t[j-k] += 111 * t[j] * rec[k-1] % mod;
        if(t[j-k] >= mod) t[j-k] -= mod;
    }
    t.resize(m); return t;
};
while(n){
    if(n & 1) s = mul(s, t);
    t = mul(t, t); n >>= 1;
}
ll ret = 0;
for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;
return ret % mod;
}
int guess_nth_term(vector<int> x, ll n){
    if(n < x.size()) return x[n];
    vector<int> v = berlekamp_massey(x);
    if(v.empty()) return 0;
    return get_nth(v, x, n);
}
struct elem{int x, y, v;}; // A_-(x, y) <- v, 0-based. no
duplicate please..
vector<int> get_min_poly(int n, vector<elem> M){
    // smallest poly P such that A^i = sum_{j < i} {A^j \times
    P_j}
    vector<int> rnd1, rnd2, gobs; mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){ return
    uniform_int_distribution<int>(lb, ub)(rng); };
    for(int i=0; i<n; i++) rnd1.push_back(randint(1, mod-1)),
        rnd2.push_back(randint(1, mod-1));
    for(int i=0; i<2*n+2; i++){
        int tmp = 0;
        for(int j=0; j<n; j++) tmp = (tmp + 111 * rnd2[j] *
            rnd1[j]) % mod;
        gobs.push_back(tmp); vector<int> nxt(n);
        for(auto &j : M) nxt[j.x] = (nxt[j.x] + 111 * j.v *
            rnd1[j.y]) % mod;
        rnd1 = nxt;
    }
    auto sol = berlekamp_massey(gobs); reverse(sol.begin(),
        sol.end()); return sol;
}
int det(int n, vector<elem> M){
    vector<int> rnd; mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){ return
    uniform_int_distribution<int>(lb, ub)(rng); };
    for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));
    for(auto &i : M) i.v = 111 * i.v * rnd[i.y] % mod;
```

```
auto sol = get_min_poly(n, M)[0]; if(n % 2 == 0) sol = mod
- sol;
for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) %
mod;
return sol;
}
```

4.9 Miller Rabin + Pollard Rho

```
constexpr int SZ = 10'000'000; bool PrimeCheck[SZ+1];
vector<int> Primes;
void Sieve(){ memset(PrimeCheck, true, sizeof PrimeCheck); /*
Sieve */ }
ull MulMod(ull a, ull b, ull c){ return ((__uint128_t)a * b % c;
}
// 32bit : 2, 7, 61
// 64bit : 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool MillerRabin(ull n, ull a){
    if(a % n == 0) return true;
    int cnt = __builtin_ctzll(n - 1);
    ull p = PowMod(a, n >> cnt, n);
    if(p == 1 || p == n - 1) return true;
    while(cnt--){ if((p=MulMod(p,p,n)) == n - 1) return true;
        return false;
    }
}
bool IsPrime(ll n){
    if(n <= SZ) return PrimeCheck[n];
    if(n <= 2) return n == 2;
    if(n % 2 == 0 || n % 3 == 0 || n % 5 == 0 || n % 7 == 0 || n
        % 11 == 0) return false;
    for(int p : {2, 325, 9375, 28178, 450775, 9780504,
        1795265022}) if(!MillerRabin(n, p)) return false;
    return true;
}
ll Rho(ll n){
    while(true){
        ll x = rand() % (n - 2) + 2, y = x, c = rand() % (n - 1) +
            1;
        while(true){
            x = (MulMod(x,x,n)+c) % n; y = (MulMod(y,y,n)+c) % n; y =
                (MulMod(y,y,n)+c) % n;
            ll d = __gcd(abs(x - y), n); if(d == 1) continue;
            if(IsPrime(d)) return d; else{ n = d; break; }
        }
    }
}
vector<pair<ll,ll>> Factorize(ll n){
    vector<pair<ll,ll>> v;
    int two = __builtin_ctzll(n);
    if(two > 0) v.emplace_back(2, two), n >>= two;
    if(n == 1) return v;
    while(!IsPrime(n)){
        ll d = Rho(n), cnt = 0; while(n % d == 0) cnt++, n /= d;
        v.emplace_back(d, cnt); if(n == 1) break;
    }
}
```

```
    if(n != 1) v.emplace_back(n, 1); return v;
}
```

4.10 Linear Sieve

```
// sp : 최소 소인수, 소수라면 0
// tau : 약수 개수, sigma : 약수 합
// phi : n 이하 자연수 중 n과 서로소인 개수
// mu : non square free이면 0, 그렇지 않다면 (-1)^(소인수 종류)
// e[i] : 소인수분해에서 i의 지수
vector<int> prime;
int sp[sz], e[sz], phi[sz], mu[sz], tau[sz], sigma[sz];
phi[1] = mu[1] = tau[1] = sigma[1] = 1;
for(int i=2; i<=n; i++){
    if(!sp[i]){
        prime.push_back(i);
        e[i] = 1; phi[i] = i-1; mu[i] = -1; tau[i] = 2; sigma[i] = i+1;
    }
    for(auto j : prime){
        if(i*j >= sz) break;
        sp[i*j] = j;
        if(i % j == 0){
            e[i*j] = e[i]+1; phi[i*j] = phi[i]*j; mu[i*j] = 0;
            tau[i*j] = tau[i]/e[i*j]*(e[i*j]+1);
            sigma[i*j] = sigma[i]*(j-1)/(pw(j, e[i*j])-1)*(pw(j, e[i*j]+1)-1)/(j-1);//overflow
            break;
        }
        e[i*j] = 1; phi[i*j] = phi[i] * phi[j]; mu[i*j] = mu[i] * mu[j];
        tau[i*j] = tau[i] * tau[j]; sigma[i*j] = sigma[i] * sigma[j];
    }
}
```

4.11 Power Tower

```
bool PowOverflow(ll a, ll b, ll c){
    __int128_t res = 1;
    bool flag = false;
    for(; b; b >>= 1, a = a * a){
        if(a >= c) flag = true, a %= c;
        if(b & 1){
            res *= a;
            if(flag || res >= c) return true;
        }
    }
    return false;
}

ll Recursion(int idx, ll mod, const vector<ll> &vec){
    if(mod == 1) return 1;
    if(idx + 1 == vec.size()) return vec[idx];
    ll nxt = Recursion(idx+1, phi[mod], vec);
    if(PowOverflow(vec[idx], nxt, mod)) return Pow(vec[idx], nxt, mod) + mod;
    else return Pow(vec[idx], nxt, mod);
}
```

```
}
ll PowerTower(const vector<ll> &vec, ll mod){ //
vec[0]^(vec[1]^(vec[2]^(...)))
    if(vec.size() == 1) return vec[0] % mod;
    else return Pow(vec[0], Recursion(1, phi[mod], vec), mod);
}
```

4.12 Discrete Log / Sqrt

Time Complexity: Log : $O(\sqrt{P} \log P)$, $O(\sqrt{P})$ with hash set
Sqrt : $O(\log^2 P)$, $O(\log P)$ in random data

```
// Given A, B, P, solve A^x === B mod P
ll DiscreteLog(ll A, ll B, ll P){
    __gnu_pbds::gp_hash_table<ll, __gnu_pbds::null_type> st;
    ll t = ceil(sqrt(P)), k = 1; // use binary search?
    for(int i=0; i<t; i++) st.insert(k), k = k * A % P;
    ll inv = Pow(k, P-2, P);
    for(int i=0, k=1; i<t; i++, k=k*inv%P){
        ll x = B * k % P;
        if(st.find(x) == st.end()) continue;
        for(int j=0, k=1; j<t; j++, k=k*A%P){
            if(k == x) return i * t + j;
        }
    }
    return -1;
}

// Given A, P, solve X^2 === A mod P
ll DiscreteSqrt(ll A, ll P){
    if(A == 0) return 0;
    if(Pow(A, (P-1)/2, P) != 1) return -1;
    if(P % 4 == 3) return Pow(A, (P+1)/4, P);
    ll s = P - 1, n = 2, r = 0, m;
    while(~s & 1) r++, s >>= 1;
    while(Pow(n, (P-1)/2, P) != P-1) n++;
    ll x = Pow(A, (s+1)/2, P), b = Pow(A, s, P), g = Pow(n, s, P);
    for(;; r=m){
        ll t = b;
        for(m=0; m<r && t!=1; m++) t = t * t % P;
        if(!m) return x;
        ll gs = Pow(g, 1LL << (r-m-1), P);
        g = gs * gs % P;
        x = x * gs % P;
        b = b * g % P;
    }
}
```

4.13 De Bruijn Sequence

```
// Create cyclic string of length k^n that contains every
length n string as substring. alphabet = [0, k - 1]
int res[10000000], aux[10000000]; // >= k^n
int de_bruijn(int k, int n) { // Returns size (k^n)
    if(k == 1) { res[0] = 0; return 1; }
    for(int i = 0; i < k * n; i++) aux[i] = 0;
    int sz = 0;
    function<void(int, int)> db = [&](int t, int p) {
```

```
    if(t > n) {
        if(n % p == 0) for(int i = 1; i <= p; i++) res[sz++] = aux[i];
    }
    else {
        aux[t] = aux[t - p]; db(t + 1, p);
        for(int i = aux[t - p] + 1; i < k; i++) aux[t] = i, db(t + 1, t);
    }
};
db(1, 1);
return sz;
}
```

4.14 Simplex / LP Duality

// Solves the canonical form: maximize c^T x, subject to ax <= b and x >= 0.

```
template<class T> // T must be of floating type
struct linear_programming_solver_simplex{
    int m, n; vector<int> nn, bb; vector<vector<T>> mat;
    static constexpr T eps = 1e-8, inf = 1/.0;
    linear_programming_solver_simplex(const vector<vector<T>> &a,
    const vector<T> &b, const vector<T> &c) : m(b.size()),
    n(c.size()), nn(n+1), bb(m), mat(m+2, vector<T>(n+2)){
        for(int i=0; i<m; i++) for(int j=0; j<n; j++) mat[i][j] = a[i][j];
        for(int i=0; i<m; i++) bb[i] = b[i], mat[i][n] = -1,
        mat[i][n + 1] = b[i];
        for(int j=0; j<n; j++) nn[j] = j, mat[m][j] = -c[j];
        nn[n] = -1; mat[m + 1][n] = 1;
    }
    void pivot(int r, int s){
        T *a = mat[r].data(), inv = 1 / a[s];
        for(int i=0; i<m+2; i++) if(i != r && abs(mat[i][s]) > eps)
        {
            T *b = mat[i].data(), inv2 = b[s] * inv;
            for(int j=0; j<n+2; j++) b[j] -= a[j] * inv2;
            b[s] = a[s] * inv2;
        }
        for(int j=0; j<n+2; j++) if(j != s) mat[r][j] *= inv;
        for(int i=0; i<m+2; i++) if(i != r) mat[i][s] *= -inv;
        mat[r][s] = inv; swap(bb[r], nn[s]);
    }
    bool simplex(int phase){
        for(auto x=m+phase-1; ; ){
            int s = -1, r = -1;
            for(auto j=0; j<n+1; j++) if(nn[j] != -phase) if(s == -1
            || pair(mat[x][j], nn[j]) < pair(mat[x][s], nn[s])) s = j;
            if(mat[x][s] >= -eps) return true;
            for(auto i=0; i<m; i++){
                if(mat[i][s] <= eps) continue;
                if(r == -1 || pair(mat[i][n + 1] / mat[i][s], bb[i]) <
                pair(mat[r][n + 1] / mat[r][s], bb[r])) r = i;
            }
        }
    }
}
```

```
        if(r == -1) return false;
        pivot(r, s);
    }
}
// Returns -inf if no solution, {inf, a vector satisfying the constraints}
// if there are abritrarily good solutions, or {maximum c^T x, x} otherwise.
// O(n m (# of pivots)), O(2 ^ n) in general.
pair<T, vector<T>> solve(){
    int r = 0;
    for(int i=1; i<m; i++) if(mat[i][n+1] < mat[r][n+1]) r = i;
    if(mat[r][n+1] < -eps){
        pivot(r, n);
        if(!simplex(2) || mat[m+1][n+1] < -eps) return {-inf, {}};
        for(int i=0; i<m; i++) if(bb[i] == -1){
            int s = 0;
            for(int j=1; j<n+1; j++) if(s == -1 || pair(mat[i][j], nn[j]) < pair(mat[i][s], nn[s])) s = j;
            pivot(i, s);
        }
    }
    bool ok = simplex(1);
    vector<T> x(n);
    for(int i=0; i<m; i++) if(bb[i] < n) x[bb[i]] = mat[i][n + 1];
    return {ok ? mat[m][n + 1] : inf, x};
}
};
```

Simplex Example

Maximize $p = 6x + 14y + 13z$

Constraints

- $0.5x + 2y + z \leq 24$

- $x + 2y + 4z \leq 60$

Coding

- $n = 2, m = 3, a = \begin{pmatrix} 0.5 & 2 & 1 \\ 1 & 2 & 4 \end{pmatrix}, b = \begin{pmatrix} 24 \\ 60 \end{pmatrix}, c = [6, 14, 13]$

LP Duality & Example

tableu를 대각선으로 뒤집고 음수 부호를 붙인 답 = -(원 문제의 답)

- Primal : $n = 2, m = 3, a = \begin{pmatrix} 0.5 & 2 & 1 \\ 1 & 2 & 4 \end{pmatrix}, b = \begin{pmatrix} 24 \\ 60 \end{pmatrix}, c = [6, 14, 13]$

- Dual : $n = 3, m = 2, a = \begin{pmatrix} -0.5 & -1 \\ -2 & -2 \\ -1 & -4 \end{pmatrix}, b = \begin{pmatrix} -6 \\ -14 \\ -13 \end{pmatrix}, c = [-24, -60]$

공식

- Primal : $\max_x c^T x, \text{Constraints } Ax \leq b, x \geq 0$

- Dual : $\min_y b^T y, \text{Constraints } A^T y \geq c, y \geq 0$

4.15 FFT, FWHT, MultipointEval, Interpolation, TaylorShift

// 104,857,601 = 25 * 2^22 + 1, w = 3 | 998,244,353 = 119 * 2^23 + 1, w = 3
// 2,281,701,377 = 17 * 2^27 + 1, w = 3 | 2,483,027,969 = 37 * 2^26 + 1, w = 3

```
// 2,113,929,217 = 63 * 2^25 + 1, w = 5 | 1,092,616,193 = 521 * 2^21 + 1, w = 3
using real_t = double; using cpx = complex<real_t>;
void FFT(vector<cpx> &a, bool inv_fft=false){
    int N = a.size(); vector<cpx> root(N/2);
    for(int i=1, j=0; i<N; i++){
        int bit = N / 2;
        while(j >= bit) j -= bit, bit >>= 1;
        if(i < (j += bit)) swap(a[i], a[j]);
    }
    long double ang = 2 * acosl(-1) / N * (inv_fft ? -1 : 1);
    for(int i=0; i<N/2; i++) root[i] = cpx(cosl(ang*i), sinl(ang*i));
    /*
    NTT : ang = pow(w, (mod-1)/n) % mod, inv_fft -> ang^{-1},
    root[i] = root[i-1] * ang
    XOR Convolution : set roots[*] = 1, a[j+k] = u+v, a[j+k+i/2]
    = u-v
    OR Convolution : set roots[*] = 1, a[j+k+i/2] += inv_fft ?
    -u : u;
    AND Convolution : set roots[*] = 1, a[j+k ] += inv_fft ? -v
    : v;
    */
    for(int i=2; i<=N; i<=1){
        int step = N / i;
        for(int j=0; j<N; j+=i) for(int k=0; k<i/2; k++){
            cpx u = a[j+k], v = a[j+k+i/2] * root[step * k];
            a[j+k] = u+v; a[j+k+i/2] = u-v;
        }
    }
    if(inv_fft) for(int i=0; i<N; i++) a[i] /= N; // skip for AND/OR convolution.
}
vector<ll> multiply(const vector<ll> &a, const vector<ll> &b){
    vector<cpx> a(all(_a)), b(all(_b));
    int N = 2; while(N < a.size() + b.size()) N <= 1;
    a.resize(N); b.resize(N); FFT(a); FFT(b);
    for(int i=0; i<N; i++) a[i] *= b[i];
    vector<ll> ret(N); FFT(a, 1); // NTT : just return a
    for(int i=0; i<N; i++) ret[i] = llround(a[i].real());
    while(ret.size() > 1 && ret.back() == 0) ret.pop_back();
    return ret;
}
vector<ll> multiply_mod(const vector<ll> &a, const vector<ll> &b, const ull mod){
    int N = 2; while(N < a.size() + b.size()) N <= 1;
    vector<cpx> v1(N), v2(N), r1(N), r2(N);
    for(int i=0; i<a.size(); i++) v1[i] = cpx(a[i] >> 15, a[i] & 32767);
    for(int i=0; i<b.size(); i++) v2[i] = cpx(b[i] >> 15, b[i] & 32767);
    FFT(v1); FFT(v2);
    for(int i=0; i<N; i++){
        int j = i ? N-i : i;
        cpx ans1 = (v1[i] + conj(v1[j])) * cpx(0.5, 0);
        cpx ans2 = (v1[i] - conj(v1[j])) * cpx(0, -0.5);
```

```
        cpx ans3 = (v2[i] + conj(v2[j])) * cpx(0.5, 0);
        cpx ans4 = (v2[i] - conj(v2[j])) * cpx(0, -0.5);
        r1[i] = (ans1 * ans3) + (ans1 * ans4) * cpx(0, 1);
        r2[i] = (ans2 * ans3) + (ans2 * ans4) * cpx(0, 1);
    }
    vector<ll> ret(N); FFT(r1, true); FFT(r2, true);
    for(int i=0; i<N; i++){
        ll av = llround(r1[i].real()) % mod;
        ll bv = ( llround(r1[i].imag()) + llround(r2[i].real()) ) % mod;
        ll cv = llround(r2[i].imag()) % mod;
        ret[i] = (av << 30) + (bv << 15) + cv;
        ret[i] %= mod; ret[i] += mod; ret[i] %= mod;
    }
    while(ret.size() > 1 && ret.back() == 0) ret.pop_back();
    return ret;
}
template<char op> vector<ll> FWHT_Conv(vector<ll> a, vector<ll> b){
    int n = max((int)a.size(), (int)b.size() - 1, 1);
    if(__builtin_popcount(n) != 1) n = 1 << (__lg(n) + 1);
    a.resize(n); b.resize(n); FWHT<op>(a); FWHT<op>(b);
    for(int i=0; i<n; i++) a[i] = a[i] * b[i] % M;
    FWHT<op>(a, true); return a;
}
vector<ll> SubsetConvolution(vector<ll> p, vector<ll> q){ // N log^2 N
    int n = max((int)p.size(), (int)q.size() - 1, 1), w = __lg(n);
    if(__builtin_popcount(n) != 1) n = 1 << (w + 1);
    p.resize(n); q.resize(n); vector<ll> res(n);
    vector<vector<ll>> a(w+1, vector<ll>(n)), b(a);
    for(int i=0; i<n; i++) a[__builtin_popcount(i)][i] = p[i];
    for(int i=0; i<n; i++) b[__builtin_popcount(i)][i] = q[i];
    for(int bit=0; bit<=w; bit++) FWHT<'|'|>(a[bit]), FWHT<'|'|>(b[bit]);
    for(int bit=0; bit<=w; bit++){
        vector<ll> c(n); // Warning : MOD
        for(int i=0; i<=bit; i++) for(int j=0; j<n; j++) c[j] += a[i][j] * b[bit-i][j] % M;
        for(auto &i : c) i %= M;
        FWHT<'|'|>(c, true);
        for(int i=0; i<n; i++) if(__builtin_popcount(i) == bit) res[i] = c[i];
    }
    return res;
}
vector<ll> Trim(vector<ll> a, size_t sz){
    a.resize(min(a.size(), sz)); return a; }
vector<ll> Inv(vector<ll> a, size_t sz){
    vector<ll> q(1, Pow(a[0], M-2, M)); // 1/a[0]
    for(int i=1; i<sz; i<=1){
        auto p = vector<ll>{2} - Multiply(q, Trim(a, i*2)); // polynomial minus
        q = Trim(Multiply(p, q), i*2);
```

```

    }
    return Trim(q, sz);
}
vector<ll> Division(vector<ll> a, vector<ll> b){
    if(a.size() < b.size()) return {};
    size_t sz = a.size() - b.size() + 1; auto ra = a, rb = b;
    reverse(ra.begin(), ra.end()); ra = Trim(ra, sz);
    reverse(rb.begin(), rb.end()); rb = Inv(Trim(rb, sz), sz);
    auto res = Trim(Multiply(ra, rb), sz);
    for(int i=sz-(int)a.size(); i>0; i--) res.push_back(0);
    reverse(res.begin(), res.end()); while(!res.empty() &&
    !res.back()) res.pop_back();
    return res;
}
vector<ll> Modular(vector<ll> a, vector<ll> b){ return a -
Multiply(b, Division(a, b)); }
ll Evaluate(const vector<ll> &a, ll x){
    ll res = 0;
    for(int i=(int)a.size()-1; i>=0; i--) res = (res * x + a[i])
    % M;
    return res >= 0 ? res : res + M;
}
vector<ll> Derivative(const vector<ll> &a){
    if(a.size() <= 1) return {};
    vector<ll> res(a.size() - 1);
    for(int i=0; i+1<a.size(); i++) res[i] = (i+1) * a[i+1] % M;
    return res;
}
vector<vector<ll>> PolynomialTree(const vector<ll> &x){
    int n = x.size(); vector<vector<ll>> tree(n*2-1);
    function<void(int,int,int)> build = [&](int node, int s, int
    e){
        if(e-s == 1){ tree[node] = vector<ll>{-x[s], 1}; return; }
        int m = s + (e-s)/2, v = node + (m-s)*2;
        build(node+1, s, m); build(v, m, e);
        tree[node] = Multiply(tree[node+1], tree[v]);
    }; build(0, 0, n); return tree;
}
vector<ll> MultipointEvaluation(const vector<ll> &a, const
vector<ll> &x){ // n log^2 n
    if(x.empty()) return {}; if(a.empty()) return
    vector<ll>(x.size(), 0);
    int n = x.size(); auto tree = PolynomialTree(x); vector<ll>
    res(n);
    function<void(int,int,int,vector<ll>>> eval = [&](int node,
    int s, int e, vector<ll> f){
        f = Modular(f, tree[node]);
        if(e-s == 1){ res[s] = f[0]; return; }
        if(f.size() < 150){ for(int i=s; i<e; i++) res[i] =
        Evaluate(f, x[i]); return; }
        int m = s + (e-s)/2, v = node + (m-s)*2;
        eval(node+1, s, m, f); eval(v, m, e, f);
    }; eval(0, 0, n, a);
    return res;
}
vector<ll> Interpolation(const vector<ll> &x, const vector<ll>
&y){ // n log^2 n

```

```

    assert(x.size() == y.size()); if(x.empty()) return {};
    int n = x.size(); auto tree = PolynomialTree(x);
    auto res = MultipointEvaluation(Derivative(tree[0]), x);
    for(int i=0; i<n; i++) res[i] = y[i] * Pow(res[i], M-2, M) %
    M; // y[i] / res[i]
    function<vector<ll>(int,int,int)> calc = [&](int node, int s,
    int e){
        if(e-s == 1) return vector<ll>{res[s]};
        int m = s + (e-s)/2, v = node + (m-s)*2;
        return Multiply(calc(node+1, s, m), tree[v]) +
        Multiply(calc(v, m, e), tree[node+1]);
    };
    return calc(0, 0, n);
}
vector<double> interpolate(vector<double> x, vector<double> y,
int n){ // n^2
    vector<double> res(n), temp(n);
    for(int k=0; k<n-1; k++) for(int i=k+1; i<n; i++) y[i] =
    (y[i] - y[k]) / (x[i] - x[k]);
    double last = 0; temp[0] = 1;
    for(int k=0; k<n; k++){
        for(int i=0; i<n; i++) res[i] += y[k] * temp[i], swap(last,
        temp[i]), temp[i] -= last * x[k];
    }
    return res;
}
vector<ll> Interpolation_0_to_n(vector<ll> y){ // n^2
    int n = y.size();
    vector<ll> res(n), tmp(n), x; // x[i] = i / (i+1)
    for(int i=0; i<n; i++) x.push_back(Pow(i+1, M-2));
    for(int k=0; k+1<n; k++) for(int i=k+1; i<n; i++)
        y[i] = (y[i] - y[k] + M) * x[i-k-1] % M;
    ll lst = 0; tmp[0] = 1;
    for(int k=0; k<n; k++) for(int i=0; i<n; i++) {
        res[i] = (res[i] + y[k] * tmp[i]) % M;
        swap(lst, tmp[i]);
        tmp[i] = (tmp[i] - lst * k) % M;
        if(tmp[i] < 0) tmp[i] += M;
    }
    return res;
}
vector<ll> Shift(const vector<ll> &f, ll c){ // O(n log n)
    if(f.size() <= 1 || c == 0) return f; // return f(x+c)
    ll n = f.size(), pw = 1; c = (c % M + M) % M;
    vector<ll> fac(n,1), inv(n,1), a(n), b(n);
    for(int i=2; i<n; i++) fac[i] = fac[i-1] * i % M;
    inv[n-1] = Pow(fac[n-1], M-2);
    for(int i=n-2; i>=2; i--) inv[i] = inv[i+1] * (i+1) % M;
    for(int i=0; i<n; i++, pw=pw*c%M)
        a[i] = f[i] * fac[i] % M, b[i] = pw * inv[i] % M;
    reverse(b.begin(), b.end()); a = Multiply(a, b);
    a = vector<ll>(a.begin()+n-1, a.begin()+n+n-1);
    for(int i=0; i<n; i++) a[i] = a[i] * inv[i] % M;
    return a;
}

```

4.16 Matroid Intersection

```

struct Matroid{
    virtual bool check(int i) = 0; // O(R^2N), O(R^2N)
    virtual void insert(int i) = 0; // O(R^3), O(R^2N)
    virtual void clear() = 0; // O(R^2), O(RN)
};
template<typename cost_t>
vector<cost_t> MI(const vector<cost_t> &cost, Matroid *m1,
Matroid *m2){
    int n = cost.size();
    vector<pair<cost_t, int>> dist(n+1);
    vector<vector<pair<int, cost_t>>> adj(n+1);
    vector<int> pv(n+1), inq(n+1), flag(n); deque<int> dq;
    auto augment = [&]() -> bool {
        fill(dist.begin(), dist.end(),
        pair(numeric_limits<cost_t>::max()/2, 0));
        fill(adj.begin(), adj.end(), vector<pair<int, cost_t>>());
        fill(pv.begin(), pv.end(), -1);
        fill(inq.begin(), inq.end(), 0);
        dq.clear(); m1->clear(); m2->clear();
        for(int i=0; i<n; i++) if(flag[i]) m1->insert(i),
        m2->insert(i);
        for(int i=0; i<n; i++){
            if(flag[i]) continue;
            if(m1->check(i)) dist[pv[i]=i] = {cost[i], 0},
            dq.push_back(i), inq[i] = 1;
            if(m2->check(i)) adj[i].emplace_back(n, 0);
        }
        for(int i=0; i<n; i++){
            if(!flag[i]) continue;
            m1->clear(); m2->clear();
            for(int j=0; j<n; j++) if(i != j && flag[j])
            m1->insert(j), m2->insert(j);
            for(int j=0; j<n; j++){
                if(flag[j]) continue;
                if(m1->check(j)) adj[i].emplace_back(j, cost[j]);
                if(m2->check(j)) adj[j].emplace_back(i, -cost[i]);
            }
        }
        while(dq.size()){
            int v = dq.front(); dq.pop_front(); inq[v] = 0;
            for(const auto &[i,w] : adj[v]){
                pair<cost_t, int> nxt{dist[v].first+w,
                dist[v].second+1};
                if(nxt < dist[i]){
                    dist[i] = nxt; pv[i] = v;
                    if(!inq[i]) dq.push_back(i), inq[i] = 1;
                }
            }
        }
        if(pv[n] == -1) return false;
        for(int i=pv[n]; ; i=pv[i]){
            flag[i] ^= 1; if(i == pv[i]) break;
        }
        return true;
    };
};

```



```
vector<int> res;
while(augment()){
    int now = 0;
    for(int i=0; i<n; i++) if(flag[i]) now += cost[i];
    res.push_back(now);
}
return res;
}
```

5 String

5.1 KMP, Hash, Manacher, Z

```
vector<int> getFail(const container &pat){
    vector<int> fail(pat.size());
    // match: pat[0..j] and pat[j-i..i] is equivalent
    // ins/del: manipulate corresponding range to pattern
    starts at 0
    // (insert/delete pat[i], manage pat[j-i..i])
    function<bool(int, int)> match = [&](int i, int j){ };
    function<void(int)> ins = [&](int i){ };
    function<void(int)> del = [&](int i){ };
    for(int i=1, j=0; i<pat.size(); i++){
        while(j && !match(i, j)){
            for(int s=i-j; s<i-fail[j-1]; s++) del(s);
            j = fail[j-1];
        }
        if(match(i, j)) ins(i), fail[i] = ++j;
    }
    return fail;
}

vector<int> doKMP(const container &str, const container &pat){
    vector<int> ret, fail = getFail(pat);
    // match: pat[0..j] and str[j-i..i] is equivalent
    // ins/del: manipulate corresponding range to pattern
    starts at 0
    // (insert/delete str[i], manage str[j-i..i])
    function<bool(int, int)> match = [&](int i, int j){ };
    function<void(int)> ins = [&](int i){ };
    function<void(int)> del = [&](int i){ };
    for(int i=0, j=0; i<str.size(); i++){
        while(j && !match(i, j)){
            for(int s=i-j; s<i-fail[j-1]; s++) del(s);
            j = fail[j-1];
        }
        if(match(i, j)){
            if(j+1 == pat.size()){
                ret.push_back(i-j);
                for(int s=i-j; s<i-fail[j]+1; s++) del(s);
                j = fail[j];
            }
            else ++j;
            ins(i);
        }
    }
    return ret;
}
```

```
// 1e5+3, 1e5+13, 131'071, 524'287, 1'299'709, 1'301'021
// 1e9-63, 1e9+7, 1e9+9, 1e9+103
template<ll P, ll M> struct Hashing {
    vector<ll> H, B;
    void Build(const string &S){
        H.resize(S.size()+1);
        B.resize(S.size()+1);
        B[0] = 1;
        for(int i=1; i<=S.size(); i++) H[i] = (H[i-1] * P +
            S[i-1]) % M;
        for(int i=1; i<=S.size(); i++) B[i] = B[i-1] * P % M;
    }
    ll sub(int s, int e){
        ll res = (H[e] - H[s-1] * B[e-s+1]) % M;
        return res < 0 ? res + M : res;
    }
};

// # a # b # a # a # b # a #
// 0 1 0 3 0 1 6 1 0 3 0 1 0
vector<int> Manacher(const string &inp){
    int n = inp.size() * 2 + 1;
    vector<int> ret(n);
    string s = "#";
    for(auto i : inp) s += i, s += "#";
    for(int i=0, p=-1, r=-1; i<n; i++){
        ret[i] = i <= r ? min(r-i, ret[2*p-i]) : 0;
        while(i-ret[i]-1 >= 0 && i+ret[i]+1 < n &&
            s[i-ret[i]-1] == s[i+ret[i]+1]) ret[i]++;
        if(i+ret[i] > r) r = i+ret[i], p = i;
    }
    return ret;
}

// input: manacher array, 1-based hashing structure
// output: set of pair(hash_val, length)
set<pair<hash_t, int>> UniquePalindrome(const vector<int> &dp,
    const Hashing &hashing){
    set<pair<hash_t, int>> st;
    for(int i=0, s, e; i<dp.size(); i++){
        if(!dp[i]) continue;
        if(i & 1) s = i/2 - dp[i]/2 + 1, e = i/2 + dp[i]/2 + 1;
        else s = (i-1)/2 - dp[i]/2 + 2, e = (i+1)/2 + dp[i]/2;

        for(int l=s, r=e; l<=r; l++, r--){
            auto now = hashing.get(l, r);
            auto [iter, flag] = st.emplace(now, r-l+1);
            if(!flag) break;
        }
    }
    return st;
}

//z[i]=match length of s[0,n-1] and s[i,n-1]
vector<int> Z(const string &s){
    int n = s.size();
    vector<int> z(n);
    z[0] = n;
    for(int i=1, l=0, r=0; i<n; i++){
        if(i < r) z[i] = min(r-i-1, z[i-l]);
```

```
        while(i+z[i] < n && s[i+z[i]] == s[z[i]]) z[i]++;
        if(i+z[i] > r) r = i+z[i], l = i;
    }
    return z;
}

5.2 Aho-Corasick

struct Node{
    int g[26], fail, out;
    Node() { memset(g, 0, sizeof g); fail = out = 0; }
};

vector<Node> T(2); int aut[100101][26];
void Insert(int n, int i, const string &s){
    if(i == s.size()){ T[n].out++; return; }
    int c = s[i] - 'a';
    if(T[n].g[c] == 0) T[n].g[c] = T.size(), T.emplace_back();
    Insert(T[n].g[c], i+1, s);
}

int go(int n, int i){ // DO NOT USE `aut` DIRECTLY
    int &res = aut[n][i]; if(res) return res;
    if(n != 1 && T[n].g[i] == 0) res = go(T[n].fail, i);
    else if(T[n].g[i] != 0) res = T[n].g[i];
    else res = 1;
    return res;
}

void Build(){
    queue<int> q; q.push(1); T[1].fail = 1;
    while(!q.empty()){
        int n = q.front(); q.pop();
        for(int i=0; i<26; i++){
            int next = T[n].g[i];
            if(next == 0) continue;
            if(n == 1) T[next].fail = 1;
            else T[next].fail = go(T[n].fail, i);
            q.push(next); T[next].out += T[T[next].fail].out;
        }
    }
}

bool Find(const string &s){
    int n = 1, ok = 0;
    for(int i=0; i<s.size(); i++){
        n = go(n, s[i] - 'a');
        if(T[n].out != 0) ok = 1;
    }
    return ok;
}
```

5.3 O(N log N) SA + LCP

```
pair<vector<int>, vector<int>> SuffixArray(const string &s){ //
O(N log N)
    int n = s.size(), m = max(n, 256);
    vector<int> sa(n), lcp(n), pos(n), tmp(n), cnt(m);
    auto counting_sort = [&]() {
        fill(cnt.begin(), cnt.end(), 0);
```

```
for(int i=0; i<n; i++) cnt[pos[i]]++;
partial_sum(cnt.begin(), cnt.end(), cnt.begin());
for(int i=n-1; i>=0; i--) sa[--cnt[pos[tmp[i]]]] = tmp[i];
};
for(int i=0; i<n; i++) sa[i] = i, pos[i] = s[i], tmp[i] = i;
counting_sort();
for(int k=1; ; k<=1){
    int p = 0;
    for(int i=n-k; i<n; i++) tmp[p++] = i;
    for(int i=0; i<n; i++) if(sa[i] >= k) tmp[p++] = sa[i] - k;
    counting_sort();
    tmp[sa[0]] = 0;
    for(int i=1; i<n; i++){
        tmp[sa[i]] = tmp[sa[i-1]];
        if(sa[i-1]+k < n && sa[i]+k < n && pos[sa[i-1]] ==
           pos[sa[i]] && pos[sa[i-1]+k] == pos[sa[i]+k]) continue;
        tmp[sa[i]] += 1;
    }
    swap(pos, tmp); if(pos[sa.back()] + 1 == n) break;
}
for(int i=0, j=0; i<n; i++, j=max(j-1,0)){
    if(pos[i] == 0) continue;
    while(sa[pos[i]-1]+j < n && sa[pos[i]]+j < n &&
          s[sa[pos[i]-1]+j] == s[sa[pos[i]]+j]) j++;
    lcp[pos[i]] = j;
}
return {sa, lcp};
}
auto [SA,LCP] = SuffixArray(S); RMQ<int> rmq(LCP);
vector<int> Pos(N); for(int i=0; i<N; i++) Pos[SA[i]] = i;
auto get_lcp = [&](int a, int b){
    if(Pos[a] > Pos[b]) swap(a, b);
    return a == b ? (int)S.size() - a : rmq.query(Pos[a]+1,
        Pos[b]);
};
vector<pair<int,int>> can; // common substring {start, lcp}
vector<tuple<int,int,int>> valid; // valid substring [start,
end_l~end_r]
for(int i=1; i<N; i++){
    if(SA[i] < X && SA[i-1] > X) can.emplace_back(SA[i], LCP[i]);
    if(i+1 < N && SA[i] < X && SA[i+1] > X)
        can.emplace_back(SA[i], LCP[i+1]);
}
for(int i=0; i<can.size(); i++){
    int skip = i > 0 ? min({can[i-1].second, can[i].second,
        get_lcp(can[i-1].first, can[i].first)}) : 0;
    valid.emplace_back(can[i].first, can[i].first + skip,
        can[i].first + can[i].second - 1);
}
}
```

5.4 Suffix Automaton

```
template<typename T, size_t S, T init_val>
struct initialized_array : public array<T, S> {
    initialized_array(){ this->fill(init_val); }
};
template<class Char_Type, class Adjacency_Type>
```

```
struct suffix_automaton{
    // Begin States
    // len: length of the longest substring in the class
    // link: suffix link
    // firstpos: minimum value in the set endpos
    vector<int> len{0}, link{-1}, firstpos{-1}, is_clone{false};
    vector<Adjacency_Type> next{{}};
    ll ans{0LL}; // 서로 다른 부분 문자열 개수
    // End States
    void set_link(int v, int lnk){
        if(link[v] != -1) ans -= len[v] - len[link[v]];
        link[v] = lnk;
        if(link[v] != -1) ans += len[v] - len[link[v]];
    }
    int new_state(int l, int sl, int fp, bool c, const
Adjacency_Type &adj){
        int now = len.size(); len.push_back(l); link.push_back(-1);
        set_link(now, sl); firstpos.push_back(fp);
        is_clone.push_back(c); next.push_back(adj); return now;
    }
    int last = 0;
    void extend(const vector<Char_Type> &s){
        last = 0; for(auto c: s) extend(c);
    }
    void extend(Char_Type c){
        int cur = new_state(len[last] + 1, -1, len[last], false,
            {}), p = last;
        while(~p && !next[p][c]) next[p][c] = cur, p = link[p];
        if(!~p) set_link(cur, 0);
        else{
            int q = next[p][c];
            if(len[p] + 1 == len[q]) set_link(cur, q);
            else{
                int clone = new_state(len[p] + 1, link[q], firstpos[q],
                    true, next[q]);
                while(~p && next[p][c] == q) next[p][c] = clone, p =
                    link[p];
                set_link(cur, clone);
                set_link(q, clone);
            }
        }
        last = cur;
    }
    int size() const { return (int)len.size(); } // # of states
}; suffix_automaton<int, initialized_array<int,26,0>> T;
// for(auto c : s) if((x=T.next[x][c]) == 0) return false;
```

5.5 Bitset LCS

```
#include <x86intrin.h>
template<size_t _Nw> void _M_do_sub(_Base_bitset<_Nw> &A, const
_Base_bitset<_Nw> &B){
    for(int i=0, c=0; i<_Nw; i++) c = _subborrow_u64(c,
        A._M_w[i], B._M_w[i], (ull*)&A._M_w[i]);
}
void _M_do_sub(_Base_bitset<1> &A, const _Base_bitset<1> &B){
    A._M_w -= B._M_w; }
```

```
template<size_t _Nb> bitset<_Nb>& operator==(bitset<_Nb> &A,
const bitset<_Nb> &B){
    _M_do_sub(A, B); return A;
}
template<size_t _Nb> inline bitset<_Nb> operator-(const
bitset<_Nb> &A, const bitset<_Nb> &B){
    bitset<_Nb> C(A); return C -= B;
}
char s[50050], t[50050];
int lcs(){ // O(NM/64)
    bitset<50050> dp, ch[26];
    int n = strlen(s), m = strlen(t);
    for(int i=0; i<m; i++) ch[t[i]-'A'].set(i);
    for(int i=0; i<n; i++){ auto x = dp | ch[s[i]-'A']; dp = dp ^
        (dp ^ x) & x; }
    return dp.count();
}
```

5.6 Lyndon Factorization, Minimum Rotation

```
// factorize string into w1 >= w2 >= ... >= wk, wi is smallest
cyclic shift of suffix.
vector<string> Lyndon(const string &s){ // O(N)
    int n = s.size(), i = 0, j, k;
    vector<string> res;
    while(i < n){
        for(j=i+1, k=i; i<n && s[k]<=s[j]; j++) k = s[k] < s[j] ? i
            : k + 1;
        for(; i<=k; i+=j-k) res.push_back(s.substr(i, j-k));
    }
    return res;
}
// rotate(v.begin(), v.begin()+min_rotation(v), v.end());
template<typename T> int min_rotation(T s){ // O(N)
    int a = 0, N = s.size();
    for(int i=0; i<N; i++) s.push_back(s[i]);
    for(int b=0; b<N; b++) for(int k=0; k<N; k++){
        if(a+k == b || s[a+k] < s[b+k]){ b += max(0, k-1); break; }
        if(s[a+k] > s[b+k]){ a = b; break; }
    }
    return a;
}
```

6 Misc

6.1 CMakeLists.txt

```
set(CMAKE_CXX_STANDARD 17)
set(CMAKE_CXX_FLAGS "-DLOCAL -lm -g -Wl,--stack,268435456")
add_compile_options(-Wall -Wextra -Winvalid-pch -Wfloat-equal
-Wno-sign-compare -Wno-misleading-indentation -Wno-parentheses)
# add_compile_options(-O3 -mavx -mavx2 -mfma)
```

6.2 Ternary Search

```
while(s + 3 <= e){ // get minimum / when multiple answer, find minimum `s`
    T l = (s + s + e) / 3, r = (s + e + e) / 3;
    if(Check(l) > Check(r)) s = l; else e = r;
}
T mn = INF, idx = s;
for(T i=s; i<=e; i++) if(T now = Check(i); now < mn) mn = now, idx = i;
```

6.3 Monotone Queue Optimization

```
template<class T, bool GET_MAX = false> // D[i] = func_{0 <= j < i} D[j] + cost(j, i)
pair<vector<T>, vector<int>> monotone_queue_dp(int n, const vector<T> &init, auto cost){
    assert((int)init.size() == n + 1); // cost function -> auto, do not use std::function
    vector<T> dp = init; vector<int> prv(n+1);
    auto compare = [](T a, T b){ return GET_MAX ? a < b : a > b; };
    auto cross = [&](int i, int j){
        int l = j, r = n + 1;
        while(l < r){
            int m = (l + r + 1) / 2;
            if(compare(dp[i] + cost(i, m), dp[j] + cost(j, m))) r = m - 1; else l = m;
        }
        return l;
    };
    deque<int> q{0};
    for(int i=1; i<=n; i++){
        while(q.size() > 1 && compare(dp[q[0]] + cost(q[0], i), dp[q[1]] + cost(q[1], i))) q.pop_front();
        dp[i] = dp[q[0]] + cost(q[0], i); prv[i] = q[0];
        while(q.size() > 1 && cross(q[q.size()-2], q.back()) >= cross(q.back(), i)) q.pop_back();
        q.push_back(i);
    }
    return {dp, prv};
}
```

6.4 Aliens Trick

```
// pair<T, vector<int>> f(T c): return opt_val, prv
// cost function must be multiplied by 2
template<class T, bool GET_MAX = false>
pair<T, vector<int>> AliensTrick(int n, int k, auto f, T lo, T hi){
    T l = lo, r = hi;
    while(l < r){
        T m = (l + r + (GET_MAX?1:0)) >> 1;
        vector<int> prv = f(m*2+(GET_MAX?-1:+1)).second;
        int cnt = 0; for(int i=n; i; i=prv[i]) cnt++;
        if(cnt <= k) (GET_MAX?l:r) = m;
        else (GET_MAX?r:l) = m + (GET_MAX?-1:+1);
    }
}
```

```
T opt_value = f(l*2).first / 2 - k*1;

vector<int> prv1 = f(l*2+(GET_MAX?1:-1)).second, p1{n};
vector<int> prv2 = f(l*2-(GET_MAX?1:-1)).second, p2{n};
for(int i=n; i; i=prv1[i]) p1.push_back(prv1[i]);
for(int i=n; i; i=prv2[i]) p2.push_back(prv2[i]);
reverse(p1.begin(), p1.end()); reverse(p2.begin(), p2.end());
assert(p2.size() <= k+1 && k+1 <=p1.size());
if(p1.size() == k+1) return {opt_value, p1};
if(p2.size() == k+1) return {opt_value, p2};

for(int i=1, j=1; i<p1.size(); i++){
    while(j < p2.size() && p2[j] < p1[i-1]) j++;
    if(p1[i] <= p2[j] && i - j == k+1 - (int)p2.size()){
        vector<int> res;
        res.insert(res.end(), p1.begin(), p1.begin()+i);
        res.insert(res.end(), p2.begin()+j, p2.end());
        return {opt_value, res};
    }
}
assert(false);
}
```

6.5 Slope Trick

```
//NOTE: f(x)=min{f(x+i),i<a+|x-k|+m -> pf(k)sf(k)ab(-a,m)
//NOTE: sf_inc에 답구하는게 들어있어서, 반드시 한 연산에 대해 pf_dec->sf_inc순서로 호출
struct LeftHull{
    void pf_dec(int x){ pq.emplace(x-bias); } //x이하의 기울기들 -1
    int sf_inc(int x){ //x이상의 기울기들 +1, pop된 원소 반환(Right Hull관리에 사용됨)
        if(pq.empty() or argmin()<=x) return x; ans += argmin()-x; //이 경우 최솟값이 증가함
        pq.emplace(x-bias); /* 이하 -1*/int r=argmin(); pq.pop(); /*전체 +1*/
        return r;
    }
    void add_bias(int x,int y){ bias+=x; ans+=y; } int minval(){ return ans; } //x축 평행이동, 최소값
    int argmin(){return pq.empty()?-inf<int>():pq.top()+bias; } //최소값 x좌표
    void operator+=(LeftHull& a){ ans+=a.ans; while(sz(a.pq)) pf_dec(a.argmin()), a.pq.pop(); }
    int size()const{return sz(pq);} PQMax<int> pq; int ans=0, bias=0;
};
//NOTE: f(x)=min{f(x+i),a<i<b+|x-k|+m -> pf(k)sf(k)ab(-a,b,m)
struct SlopeTrick{
    void pf_dec(int x){ l.pf_dec(-r.sf_inc(-x)); }
    void sf_inc(int x){ r.pf_dec(-l.sf_inc(x)); }
    void add_bias(int lx,int rx,int y){ l.add_bias(lx,0),r.add_bias(-rx,0),ans+=y; }
    int minval(){return ans+l.minval()+r.minval(); }
    int argmin(){return {l.argmin(),-r.argmin()}; }
    void operator+=(SlopeTrick& a){
```

```
while(sz(a.l.pq)) pf_dec(a.l.argmin()),a.l.pq.pop();
l.ans+=a.l.ans;
while(sz(a.r.pq)) sf_inc(-a.r.argmin()),a.r.pq.pop();
r.ans+=a.r.ans; ans+=a.ans;
}
int size()const{return l.size()+r.size();} LeftHull l,r; int ans=0;
};
//LeftHull 역추적 방법: 스텝i의 argmin값을 am(i)라고 하자. 스텝n부터 스텝1까지 ans[i]=min(ans[i+1],am(i))하면 된다. 아래는 증명..은 아니고 간략한 이유
//am(i)<=ans[i+1]일때: ans[i]=am(i)
//x[i]>ans[i+1]일때: ans[i]=ans[i+1] 왜냐하면 f(i,a)는 a<x[i]에서 감소함수이므로 가능한 최대로 오른쪽으로 붙은 ans[i+1]이 최적.
//스텝i에서 add_bias(k,0)한다면 간격제한k가 있는것이므로 ans[i]=min(ans[i+1]-k,x[i])으로 수정.
//LR Hull 역추적은 케이스나뉘서 위 방법을 확장하면 될듯
```

6.6 Hook Length Formula

```
int HookLength(const vector<int> &young){
    if(young.empty()) return 1;
    vector<int> len(young[0]);
    ll num = 1, div = 1, cnt = 0;
    for(int i=(int)young.size()-1; i>=0; i--){
        for(int j=0; j<young[i]; j++){
            num = num * ++cnt % MOD;
            div = div * ((+len[j] + young[i] - j - 1) % MOD);
        }
        return num * Pow(div, MOD-2) % MOD;
    }
}
```

6.7 Floating Point Add

```
T Add(vector<T> v){
    vector<T> a, b; T r = 0;
    for(auto i : v) (i>0?a:b).push_back(i);
    sort(a.begin(), a.end());
    sort(b.begin(), b.end(), greater<>());
    for(int i=0, j=0; i<a.size() || j<b.size(); ){
        if(i < a.size() && (j == b.size() || r <= 0)) r += a[i++];
        else r += b[j++];
    }
    return r;
}
```

6.8 Random, PBDS, Bit Trick, Bitset

```
mt19937
rd((unsigned)chrono::steady_clock::now().time_since_epoch().count)
uniform_int_distribution<int> rnd_int(1, r); // rnd_int(rd)
uniform_real_distribution<double> rnd_real(0, 1); // rnd_real(rd)
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/rope>
```

```
using namespace __gnu_pbds; //ordered_set :
find_by_order(order), order_of_key(key)
using namespace __gnu_cxx; //crope : append(str), substr(s, e),
at(idx)
template <typename T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
tree_order_statistics_node_update>;
int __builtin_clz(int x); // number of leading zero
int __builtin_ctz(int x); // number of trailing zero
int __builtin_popcount(int x); // number of 1-bits in x
lsb(n): (n & -n); // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);
long long next_perm(long long v){
    long long t = v | (v-1);
    return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) +
1));
}
int frq(int n, int i) { // # of digit i in [1, n]
    int j, r = 0;
    for (j = 1; j <= n; j *= 10) if (n / j / 10 >= !i) r += (n /
10 / j - !i) * j + (n / j % 10 > i ? j : n / j % 10 == i ? n
% j + 1 : 0);
    return r;
}
bitset<17> bs; bs[1] = bs[7] = 1;
assert(bs._Find_first() == 1);
assert(bs._Find_next(0) == 1 && bs._Find_next(1) == 7);
assert(bs._Find_next(3) == 7 && bs._Find_next(7) == 17);
cout << bs._Find_next(7) << "\n";
template <int len = 1> // Arbitrary sized bitset
void solve(int n){
    if(len < n){ solve<std::min(len*2, MAXLEN)>(n); return; }
    // solution using bitset<len>
}
```

6.9 Fast I/O, Fast Div/Mod, Hilbert Mo’s

```
namespace io { // thanks to cgiosy
    const signed IS=1<<20;
    char I[IS+1],*J=I;
    inline void daer(){if(J>=I+IS-64){
        char*p=I;do*p++=*J++;
        while(J!=I+IS);p[read(0,p,I+IS-p)]=0;J=I;}}
    template<int N=10,typename T=int>inline T getu(){
        daer();T x=0;int k=0;do x=x*10+*J-'0';
        while(++J=='0'&&+k<N);++J;return x;}
    template<int N=10,typename T=int>inline T geti(){
        daer();bool e=*J=='-';J+=e;return(e?-1:1)*getu<N,T>();}
    struct f{f(I){I[read(0,I,IS)]=0;}}flu;
};
struct FastMod{ // typedef __uint128_t L;
    ull b, m;
    FastMod(ull b) : b(b), m(ull((L(1) << 64) / b)) {}
    ull reduce(ull a){ // can be proven that 0 <= r < 2*b
        ull q = (ull)((L(m) * a) >> 64), r = a - q * b;
        return r >= b ? r - b : r;
    }
};
```

```
}
};
ull mulmod(ull a, ull b, ull M){ // ~2x faster than int128
    ll ret = a * b - M * ull(1.L / M * a * b);
    return ret + M * (ret < 0) - M * (ret >= (ll)M);
} // safe for 64bit integer when long double is 80bit
inline int64_t hilbertOrder(int x, int y, int pow, int rotate)
{
    if(pow == 0) return 0;
    int hpow = 1 << (pow-1), seg = (x<hpow) ? ( (y<hpow) ? 0 : 3
) : ( (y<hpow) ? 1 : 2 );
    const int rotateDelta[4] = {3, 0, 0, 1}; seg = (seg + rotate)
& 3;
    int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
    int nrot = (rotate + rotateDelta[seg]) & 3;
    int64_t subSquareSize = int64_t(1) << (2*pow - 2);
    int64_t ans = seg * subSquareSize, add = hilbertOrder(nx, ny,
pow-1, nrot);
    ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add -
1); return ans;
}
struct Query{
    int s, e, x; ll order; void init(){ order = hilbertOrder(s,
e, 21, 0); }
    bool operator < (const Query &t) const { return order <
t.order; }
};
```

6.10 DP Opt, Tree Opt, Well-Known Ideas

```
// Quadrangle Inequality : C(a, c)+C(b, d) ≤ C(a, d)+C(b, c)
// Monotonicity : C(b, c) ≤ C(a, d)
// CHT, DnC Opt(Quadrangle), Knuth(Quadrangle and Monotonicity)
```

// 크기가 A, B인 두 서브트리의 결과를 합칠 때 O(AB)이면 O(N^3)이 아니라 O(N^2)
// 각 정점마다 sum(2 ~ C번째로 높이가 작은 정점의 높이)에 결과를 구할 수 있으면 O(N^2)이 아니라 O(N)

```
// IOI 16 Alien(Lagrange Multiplier), IOI 11 Elephant(sqrt
batch process)
// IOI 09 Region
// 서로소 합집합의 크기가 적당히 bound 되어 있을 때 사용
// 쿼리 메모이제이션 / 쿼리 하나에 O(A log B), 전체 O(N√Q log N)
```

6.11 Highly Composite Numbers, Large Prime

< 10^k	number	divisors	2	3	5	7	11	13	17	19	23	29	31	37
1	6	4	1	1										
2	60	12	2	1	1									
3	840	32	3	1	1	1								
4	7560	64	3	3	1	1								
5	83160	128	3	3	1	1	1							
6	720720	240	4	2	1	1	1	1						
7	8648640	448	6	3	1	1	1	1						
8	73513440	768	5	3	1	1	1	1	1					

9	735134400	1344	6	3	2	1	1	1	1					
10	6983776800	2304	5	3	2	1	1	1	1	1				
11	97772875200	4032	6	3	2	2	1	1	1	1				
12	963761198400	6720	6	4	2	1	1	1	1	1	1			
13	9316358251200	10752	6	3	2	1	1	1	1	1	1	1		
14	97821761637600	17280	5	4	2	2	1	1	1	1	1	1		
15	866421317361600	26880	6	4	2	1	1	1	1	1	1	1	1	
16	8086598962041600	41472	8	3	2	2	1	1	1	1	1	1	1	
17	74801040398884800	64512	6	3	2	2	1	1	1	1	1	1	1	1
18	897612484786617600	103680	8	4	2	2	1	1	1	1	1	1	1	1

< 10^k	prime	# of prime	< 10^k	prime
1	7	4	10	9999999967
2	97	25	11	99999999977
3	997	168	12	999999999989
4	9973	1229	13	9999999999971
5	99991	9592	14	9999999999973
6	999983	78498	15	99999999999989
7	9999991	664579	16	999999999999937
8	99999989	5761455	17	999999999999997
9	999999937	50847534	18	9999999999999989

6.12 Catalan, Burnside, Grundy, Pick, Hall, Simpson, Kirchhoff, Area of Quadrangle, Fermat Point, Euler

- 카탈란 수
1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900
 $C_n = \text{binomial}(n * 2, n) / (n + 1)$;
- 길이가 2n인 올바른 괄호 수식의 수
- n + 1개의 리프를 가진 풀 바이너리 트리의 수
- n + 2각형을 n개의 삼각형으로 나누는 방법의 수
- 여는 괄호 n개, 닫는 괄호 k(≤ n)개 경우의 수 $(n-k+1)/(n+1) \times \binom{n+k}{k}$
- Burnside’s Lemma
- 수식
G=(X,A): 집합X와 액션A로 정의되는 군G에 대해, $|A||X/A| = \text{sum}(|\text{Fixed points of a}|, \text{for all a in A})$
X/A 는 Action으로 서로 변형가능한 X의 원소들을 동치로 묶었을때 동치류(파티션) 집합
- 풀어쓰기
orbit: 그룹에 대해 두 원소 a,b와 액션f에 대해 f(a)=b인거에 간선연결한 컴포넌트(연결집합)
orbit개수 = sum(각 액션 g에 대해 f(x)=x인 x(고정점)개수)/액션개수
- 자유도 치트시트
회전 n개: 회전i의 고정점 자유도=gcd(n,i)
임의뒤집기 n=홀수: n개 원소중심축(자유도 (n+1)/2)
임의뒤집기 n=짝수: n/2개 원소중심축(자유도 n/2+1) + n/2개 원소안 지나는축(자유도 n/2)
- 알고리즘 게임
- Nim Game의 해법 : 각 더미의 돌의 개수를 모두 XOR했을 때 0 이 아니면 첫번째, 0 이면 두번째 플레이어가 승리.
- Grundy Number : 어떤 상황의 Grundy Number는, 가능한 다음 상황들의 Grundy Number를 모두 모은 다음, 그 집합에 포함 되지 않는 가장 작은 수가 현재 state의 Grundy Number가 된다. 만약 다음 state가 독립된 여러개의 state들로 나뉘는 경우, 각각의 state의 Grundy Number

의 XOR 합을 생각한다.

- Subtraction Game : 한 번에 k 개까지의 돌만 가져갈 수 있는 경우, 각 터미의 돌의 개수를 k + 1로 나눈 나머지를 XOR 합하여 판단한다.

- Index-k Nim : 한 번에 최대 k개의 터미를 골라 각각의 터미에서 아무렇게나 돌을 제거할 수 있을 때, 각 binary digit에 대하여 합을 k + 1로 나눈 나머지를 계산한다. 만약 이 나머지가 모든 digit에 대하여 0이라면 두번째, 하나라도 0이 아니라면 첫번째 플레이어가 승리.

- Misere Nim : 모든 돌 무더기가 1이면 N이 홀수일 때 후공 승, 그렇지 않은 경우 XOR 합 0이면 후공 승

• Pick’s Theorem

격자점으로 구성된 simple polygon이 주어짐. I 는 polygon 내부의 격자점 수, B 는 polygon 선분 위 격자점 수, A는 polygon의 넓이라고 할 때, 다음과 같은 식이 성립한다. $A = I + B/2 - 1$

```
// number of (x, y) : (0 <= x < n && 0 < y <= k/d x + b/d)
ll count_solve(ll n, ll k, ll b, ll d) { // argument
should be positive
    if (k == 0) {
        return (b / d) * n;
    }
    if (k >= d || b >= d) {
        return ((k / d) * (n - 1) + 2 * (b / d)) * n / 2 +
count_solve(n, k % d, b % d, d);
    }
    return count_solve((k * n + b) / d, d, (k * n + b) % d,
k);
}
```

• 홀의 결혼 정리 : 이분그래프(L-R)에서, 모든 L을 매칭하는 필요충분 조건 = L에서 임의의 부분집합 S를 골랐을 때, 반드시 (S의 크기) <= (S와 연결되어있는 모든 R의 크기)이다.

• Simpson 공식 (적분) : Simpson 공식, $S_n(f) = \frac{h}{3}[f(x_0) + f(x_n) + 4\sum f(x_{2i+1}) + 2\sum f(x_{2i})]$

- $M = \max |f^4(x)|$ 이라고 하면 오차 범위는 최대 $E_n \leq \frac{M(b-a)}{180}h^4$

• Kirchhoff’s Theorem : 그래프의 스패닝 트리 개수

- $m[i][j] := -(i-j \text{ 간선 개수})$ ($i \neq j$)

- $m[i][i] :=$ 정점 i의 degree

- $res = (m$ 의 첫 번째 행과 첫 번째 열을 없앤 $(n-1)$ by $(n-1)$ matrix의 행렬식)

• Tutte Matrix : 그래프의 최대 매칭

- $m[i][j] :=$ 간선 (i, j) 가 없으면 0, 있으면 $i < j ? r : -r$, $r \in [0, P)$ 구간의 임의의 정수

- $rank(m)/2$ 가 높은 확률로 최대 매칭

• 브라마굽타 : 원에 내접하는 사각형의 각 선분의 길이가 a, b, c, d 일 때 사각형의 넓이 $S = \sqrt{(s-a)(s-b)(s-c)(s-d)}$, $s = (a+b+c+d)/2$

• 브레치나이더 : 임의의 사각형의 각 변의 길이를 a, b, c, d 라고 하고, 마주보는 두 각의 합을 2로 나눈 값을 θ 라 하면, $S = \sqrt{(s-a)(s-b)(s-c)(s-d) - abcd \times \cos^2 \theta}$

• 페르마 포인트 : 삼각형의 세 꼭짓점으로부터 거리의 합이 최소가 되는 점 $2\pi/3$ 보다 큰 각이 있으면 그 점이 페르마 포인트, 그렇지 않으면 각 변마다 정삼각형 그린 다음, 정삼각형의 끝점에서 반대쪽 삼각형의 꼭짓점으로 연결한 선분의 교점

$2\pi/3$ 보다 큰 각이 없으면 거리의 합은 $\sqrt{(a^2 + b^2 + c^2 + 4\sqrt{3}S)}/2$, S 는 넓이

• 오일러 정리: 서로소인 두 정수 a, n 에 대해 $a^{\phi(n)} \equiv 1 \pmod n$

모든 정수에 대해 $a^n \equiv a^{n-\phi(n)} \pmod n$

$m \geq \log_2 n$ 이면 $a^m \equiv a^{m\% \phi(n) + \phi(n)} \pmod n$

• $g^0 + g^1 + g^2 + \dots + g^{p-2} \equiv -1 \pmod p$ iff $g = 1$, otherwise 0.

6.13 inclusive and exclusive, Stirling Number, Bell Number

• 공 구별 X, 상자 구별 O, 전사함수 : 포함배제 $\sum_{i=1}^k (-1)^{k-i} \times kCi \times i^n$

• 공 구별 O, 상자 구별 X, 전사함수 : 제 2종 스티어링 수 $S(n, k) = k \times S(n - 1, k) + S(n - 1, k - 1)$

포함배제하면 $O(K \log N)$, $S(n, k) = 1/k! \times \sum_{i=1}^k (-1)^{k-i} \times kCi \times i^n$

• 공 구별 O, 상자 구별 X, 제약없음 : 벨 수 $B(n, k) = \sum_{i=0}^k S(n, i)$ 몇 개의 상자를 버릴지 다 돌아보기

수식 정리하면 $O(\min(N, K) \log N)$ 에 됨. $B(n, n) = \sum_{i=0}^{n-1} (n-1)Ci \times B(i, i)$

$$B(n, k) = \sum_{j=0}^k S(n, j) = \sum_{j=0}^k 1/j! \sum_{i=0}^j (-1)^{j-i} jCi \times i^n = \sum_{j=0}^k \sum_{i=0}^j \frac{(-1)^{j-i}}{i!(j-i)!} i^n = \sum_{i=0}^k \sum_{j=i}^k \frac{(-1)^{j-i}}{i!(j-i)!} i^n = \sum_{i=0}^k \frac{i^n}{i!} \sum_{j=0}^{k-i} \frac{(-1)^j}{j!}$$

• Derangement: $D(n) = (n - 1)(D(n - 1) + D(n - 2))$

• Signed Stirling 1: $S_1(n, k) = (n - 1)S_1(n - 1, k) + S_1(n - 1, k - 1)$

• Unsigned Stirling 1: $C_1(n, k) = (n - 1)C_1(n - 1, k) + C_1(n - 1, k - 1)$

• Stirling 2: $S_2(n, k) = kS_2(n - 1, k) + S_2(n - 1, k - 1)$

• Stirling 2: $S_2(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$

• Partition: $p(n, k) = p(n - 1, k - 1) + p(n - k, k)$

• Partition: $p(n) = \sum (-1)^k p(n - k(3k - 1)/2)$

• Bell: $B(n) = \sum_{k=1}^n \binom{n-1}{k-1} B(n - k)$

• Catalan: $C_n = \frac{1}{n+1} \binom{2n}{n}$

• Catalan: $C_n = \binom{2n}{n} - \binom{2n}{n+1}$

• Catalan: $C_n = \frac{(2n)!}{n!(n+1)!}$

• Catalan: $C_n = \sum C_i C_{n-i}$

6.14 About Graph Matching(Graph with $|V| \leq 500$)

• **Game on a Graph** : s에 토큰이 있음. 플레이어는 각자의 턴마다 토큰을 인접한 정점으로 옮기고 못 옮기면 짐.

s를 포함하지 않는 최대 매칭이 존재함 \leftrightarrow 후공이 이김

• **Chinese Postman Problem** : 모든 간선을 방문하는 최소 가중치 Walk를 구하는 문제.

Floyd를 돌린 다음, 홀수 정점들을 모아서 최소 가중치 매칭 (홀수 정점은 짝수 개 존재)

• **Unweighted Edge Cover** : 모든 정점을 덮는 가장 작은(minimum cardinality/weight) 간선 집합을 구하는 문제

$|V| - |M|$, 길이 3짜리 경로 없음, star graph 여러 개로 구성

• **Weighted Edge Cover** : $sum_{v \in V} (w(v)) - sum_{(u,v) \in M} (w(u) + w(v) - d(u, v))$, $w(x)$ 는 x와 인접한 간선의 최소 가중치

• **NEERC’18 B** : 각 기계마다 2명의 노동자가 다뤄야 하는 문제.

기계마다 두 개의 정점을 만들고 간선으로 연결하면 정답은 $|M| - |기계|$ 임. 정답에 1/2씩 기여한다는 점을 생각해보면 좋음.

• **Min Disjoint Cycle Cover** : 정점이 중복되지 않으면서 모든 정점을 덮는 길이 3 이상의 사이클 집합을 찾는 문제.

모든 정점은 2개의 서로 다른 간선, 일부 간선은 양쪽 끝점과 매칭되어야 하므로 플로우를 생각할 수 있지만 용량 2짜리 간선에 유량을 1만큼 흘릴 수 있으므로 플로우는 불가능.

각 정점과 간선을 2개씩((v, v') , $(e_{i,u}, e_{i,v})$)로 복사하자. 모든 간선 $e = (u, v)$ 에 대해 e_u 와 e_v 를 잇는 가중치 w짜리 간선을 만들고(like NEERC18), $(u, e_{i,u}), (u', e_{i,u}), (v, e_{i,v}), (v', e_{i,v})$ 를 연결하는 가중치

0짜리 간선을 만들자. Perfect 매칭이 존재함 \leftrightarrow Disjoint Cycle Cover 존재. 최대 가중치 매칭 찾은 뒤 모든 간선 가중치 합에서 매칭 빼면 됨.

• **Two Matching** : 각 정점이 최대 2개의 간선과 인접할 수 있는 최대 가중치 매칭 문제.

각 컴포넌트는 정점 하나/경로/사이클이 되어야 함. 모든 서로 다른 정점 쌍에 대해 가중치 0짜리 간선 만들고, 가중치 0짜리 (v, v') 간선 만들면 Disjoing Cycle Cover 문제가 됨. 정점 하나만 있는 컴포넌트는 self-loop, 경로 형태의 컴포넌트는 양쪽 끝점을 연결한다고 생각하면 편함.

6.15 Calculus, Newton’s Method

- $(\arcsin x)' = 1/\sqrt{1 - x^2}$
- $(\tan x)' = 1 + \tan^2 x$
- $\int \tan ax = -\ln |\cos ax|/a$
- $(\arccos x)' = -1/\sqrt{1 - x^2}$
- $(\arctan x)' = 1/(1 + x^2)$
- $\int x \sin ax = (\sin ax - ax \cos ax)/a^2$
- Newton: $x_{n+1} = x_n - f(x_n)/f'(x_n)$
- $\oint_C (Ldx + Mdy) = \int_D (\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y}) dxdy$
- where C is positively oriented, piecewise smooth, simple, closed; D is the region inside C; L and M have continuous partial derivatives in D.

6.16 Checklist

- (예비소집) bits/stdc++.h, int128, long double 80bit, avx2 확인
- (예비소집) 스택 메모리 확인(지역 변수, 재귀 함수, 람다 재귀)
- 비슷한 문제를 풀어본 적이 있던가?
- 단순한 방법에서 시작할 수 있을까? (Brute Force)
- 내가 문제를 푸는 과정을 수식화할 수 있을까? (예제를 직접 해결해보면서)
- 문제를 단순화할 수 없을까?
- 그림으로 그려볼 수 있을까?
- 수식으로 표현할 수 있을까?
- 문제를 분해할 수 있을까?
- 뒤에서부터 생각해서 풀 수 있을까?
- 순서를 강제할 수 있을까?
- 특정 형태의 답만을 고려할 수 있을까? (정규화)
- 구간을 통째로 가져간다 : 플로우 + 적당한 자료구조 ($i, i + 1, k, 0$), $(s, e, 1, w)$, $(N, T, k, 0)$
- $a = b$: a만 움직이기, b만 움직이기, 두 개 동시에 움직이기, 반대로 움직이기
- 말도 안 되는 것들을 한 번은 생각해보기 / ”당연하다고 생각한 것” 다시 생각해보기
- Directed MST / Dominator Tree
- 일정 비율 충족 or 2 3개로 모두 커버 : 랜덤
- 확률 : DP, 이분 탐색(NYPC 2019 Finals C)
- 최대/최소 : 이분 탐색, 그리디(Prefix 고정, Exchange Argument), DP(순서 고정)