

OCaml: The Basics

Concepts of Programming Languages
Lecture 2

Announcements

- » You must attend **both** the in-class quizzes
- » We will follow the online book: [OCaml Programming: Correct + Efficient + Beautiful](#)
- » Hope you all have installed OCaml

Outline

- » Briefly outline the use of **Dune**
- » Cover the **basic expressions** we need to start programming in OCaml, look at some examples
- » Define more carefully the notion of an **inference rule**
- » Write basic OCaml programs on primitive types

Recall: Expressions

In OCaml and (functional languages in general),
everything is an expression

Functions, variables,
arguments, assignments, etc.

All expressions have a value

$$2 + (2 * 3)$$

if x = 3 then 4 else 5

$$H(f(f(f(x, y), 2), g(z)))$$

An OCaml Program is an Expression

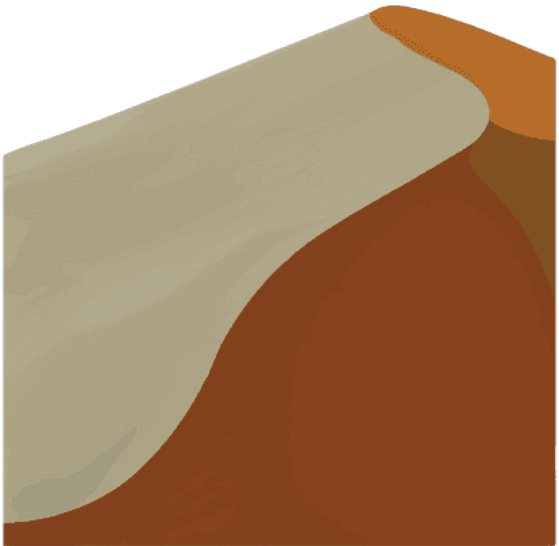
Therefore, it has a type

And a value

For every possible expression, we'll
define the syntax rules, the typing
rules, and the semantic rules

Working with OCaml

Dune



DUNE

Dune

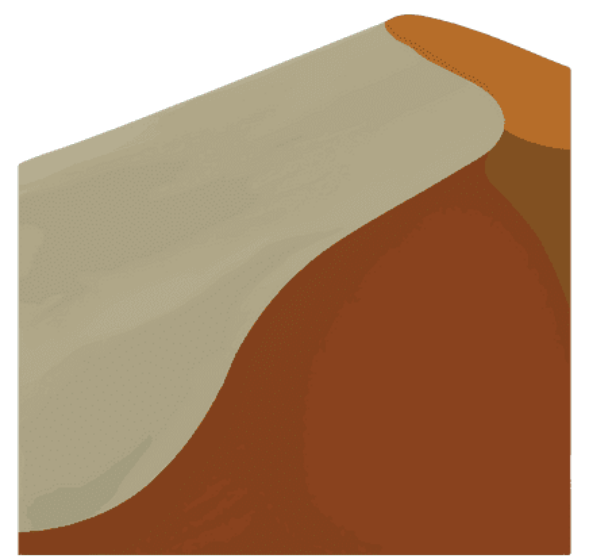
Dune is a build tool for OCaml



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations



DUNE

Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:

» **dune build:** type check your project



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:

» **dune build:** type check your project

» **dune utop:** open Utop in a project aware way



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:

- » **dune build:** type check your project
- » **dune utop:** open Utop in a project aware way
- » **dune test:** run a testing code associated with the project



Dune

Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:

- » **dune build:** type check your project
- » **dune utop:** open Utop in a project aware way
- » **dune test:** run a testing code associated with the project
- » **dune exec PROJ_NAME:** run the executable of your project



Dune



Dune is a build tool for OCaml

It allows us to specify project-level dependencies and configurations

We'll use it throughout the course for all assignments and projects

Cheatsheet:

- » **dune build:** type check your project
- » **dune utop:** open Utop in a project aware way
- » **dune test:** run a testing code associated with the project
- » **dune exec PROJ_NAME:** run the executable of your project
- » **dune clean:** removes files created by dune build (not so important but may come in handy)

UTop

```
Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directions:
#require "package";;      to load a package
#list;;                   to list the available packages
#camlp4o;;                to load camlp4 (standard syntax)
#camlp4r;;                to load camlp4 (revised syntax)
#predicates "p,q,...";;   to set these predicates
Topfind.reset();;         to force that packages will be reloaded
#thread;;                 to enable threads

Type #utop_help for help about using utop.

-( 23:00:06 )-< command 0 >
utop # 1 + 2;;
- : int = 3
-( 23:00:06 )-< command 1 >
utop #
```

Afl_instrument	Alias_analysis	Allocated_const	Annot	Arc
----------------	----------------	-----------------	-------	-----

UTop

UTop is an interface for the OCaml toplevel (use `utop` in terminal)

```
Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directives:
#require "package";;      to load a package
#list;;                   to list the available packages
#camlp4o;;                to load camlp4 (standard syntax)
#camlp4r;;                to load camlp4 (revised syntax)
#predicates "p,q,...";;  to set these predicates
Topfind.reset();;         to force that packages will be reloaded
#thread;;                 to enable threads

Type #utop_help for help about using utop.

-( 23:00:06 )-< command 0 >
utop # 1 + 2;;
- : int = 3

-( 23:00:06 )-< command 1 >
utop #
```

Afl_instrument	Alias_analysis	Allocated_const	Annot	Arc
----------------	----------------	-----------------	-------	-----

UTop

UTop is an interface for the OCaml toplevel (use `utop` in terminal)

It's basically an **OCaml calculator**.
It can *evaluate* OCaml expressions
(useful for debugging)

```
Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directi
#require "package";;      to load a package
#list;;                  to list the available packag
#camlp4o;;                to load camlp4 (standard syn
#camlp4r;;                to load camlp4 (revised synt
#predicates "p,q,...";;  to set these predicates
Topfind.reset();;         to force that packages will
#thread;;                 to enable threads

Type #utop_help for help about using utop.

-( 23:00:06 )-< command 0 >
utop # 1 + 2;;
- : int = 3
-( 23:00:06 )-< command 1 >
utop # █

Afl_instrument | Alias_analysis | Allocated_const | Annot | Arc
```

UTop

UTop is an interface for the OCaml toplevel (use `utop` in terminal)

It's basically an **OCaml calculator**.
It can *evaluate* OCaml expressions
(useful for debugging)

Cheatsheet:

Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directi

#require "package";;

to load a package

#list;;

to list the available packag

#camlp4o;;

to load camlp4 (standard syn

#camlp4r;;

to load camlp4 (revised synt

#predicates "p,q,...";;

to set these predicates

Topfind.reset();;

to force that packages will

#thread;;

to enable threads

Type #utop_help for help about using utop.

-(23:00:06)-< command 0 >

utop # 1 + 2;;

- : int = 3

-(23:00:06)-< command 1 >

utop #

Afl_instrument	Alias_analysis	Allocated_const	Annot	Arc
----------------	----------------	-----------------	-------	-----

UTop

UTop is an interface for the OCaml toplevel (use `utop` in terminal)

It's basically an **OCaml calculator**.
It can *evaluate* OCaml expressions
(useful for debugging)

Cheatsheet:

» expressions must be followed
with two semicolons

```
Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directi
#require "package";;      to load a package
#list;;                  to list the available packag
#camlp4o;;               to load camlp4 (standard syn
#camlp4r;;               to load camlp4 (revised synt
#predicates "p,q,...";;  to set these predicates
Topfind.reset();;        to force that packages will
#thread;;                to enable threads

Type #utop_help for help about using utop.

-( 23:00:06 )-< command 0 >
utop # 1 + 2;;
- : int = 3
-( 23:00:06 )-< command 1 >
utop # █

Afl_instrument | Alias_analysis | Allocated_const | Annot | Arc
```

UTop

UTop is an interface for the OCaml toplevel (use `utop` in terminal)

It's basically an **OCaml calculator**.
It can *evaluate* OCaml expressions
(useful for debugging)

Cheatsheet:

» expressions must be followed
with two semicolons

» `#quit;;` or (Ctl-D) leaves UTop

```
Welcome to utop version %%VERSION%%

Findlib has been successfully loaded. Additional directi
#require "package";;      to load a package
#list;;                  to list the available packag
#camlp4o;;               to load camlp4 (standard syn
#camlp4r;;               to load camlp4 (revised synt
#predicates "p,q,...";;  to set these predicates
Topfind.reset();;        to force that packages will
#thread;;                to enable threads

Type #utop_help for help about using utop.

-( 23:00:06 )-< command 0 >
utop # 1 + 2;;
- : int = 3
-( 23:00:06 )-< command 1 >
utop # █

Afl_instrument | Alias_analysis | Allocated_const | Annot | Arc
```

Expressions (Informally)

High Level View

Each expression has a **form**
(e.g. `2 + 7`, `true`, `"true"`)
[defined by **syntax rules**]

Each expression has a **unique type**
(e.g. `2 + 7 : int`, `true : bool`, `"true" : string`)
[defined by typing rules]

Each expression has a **unique value**
(e.g. `2 + 7 ↓ 9`, `true ↓ true`, `"true" ↓ "true"`)
[defined by semantics rules]

Primitive Types and Literals

As with any PL, OCaml has a collection of standard types and literals

Type	Literals	Operators
int	0, -2, 13, -023	+, -, *, /, mod
float	3., -1.01	+. , -. , *. , /.
bool	true, false	&&, , not
char	'b', 'c'	
string	"word", "@*&#"	^

demo

(simple use of primitive types)

A Couple Notes on Operators

A Couple Notes on Operators

Operators for int and float are *different*,
e.g., `+` (integer addition) and `+.` (float addition)

A Couple Notes on Operators

Operators for int and float are *different*,
e.g., `+` (integer addition) and `+.` (float addition)

OCaml has **no operator overloading** (bad feature)

A Couple Notes on Operators

Operators for int and float are *different*,
e.g., `+` (integer addition) and `+.` (float addition)

OCaml has **no operator overloading** (bad feature)

Comparison operators are standard, e.g., `<`, `<=`, `>`, `>=`, and
can be used to compare any expressions of the same type

A Couple Notes on Operators

Operators for int and float are *different*,
e.g., `+` (integer addition) and `+.` (float addition)

OCaml has **no operator overloading** (bad feature)

Comparison operators are standard, e.g., `<`, `<=`, `>`, `>=`, and
can be used to compare any expressions of the same type

Note that equality check is just `=` (not `==`) and inequality
is `<>` (not `!=`)

A Note on Type Annotations

```
let x : int = 2 + 7  
let y : bool = true  
let z : string = "true"
```

A Note on Type Annotations

```
let x : int = 2 + 7  
let y : bool = true  
let z : string = "true"
```

OCaml has type inference which means we rarely have to *specify* the types of expression in our program

A Note on Type Annotations

```
let x : int = 2 + 7
let y : bool = true
let z : string = "true"
```

OCaml has type inference which means we rarely have to *specify* the types of expression in our program

That said, you **should** include type annotations, especially at the beginning, because they're useful for *documentation* and for *code clarity*

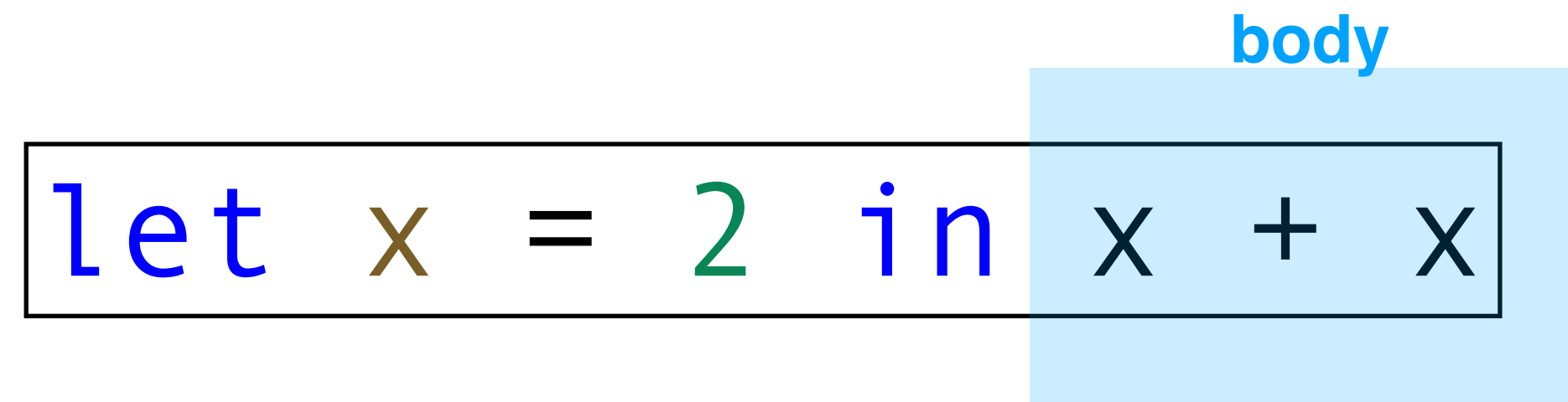
First Abstraction of the Day

Let Expressions and Variables

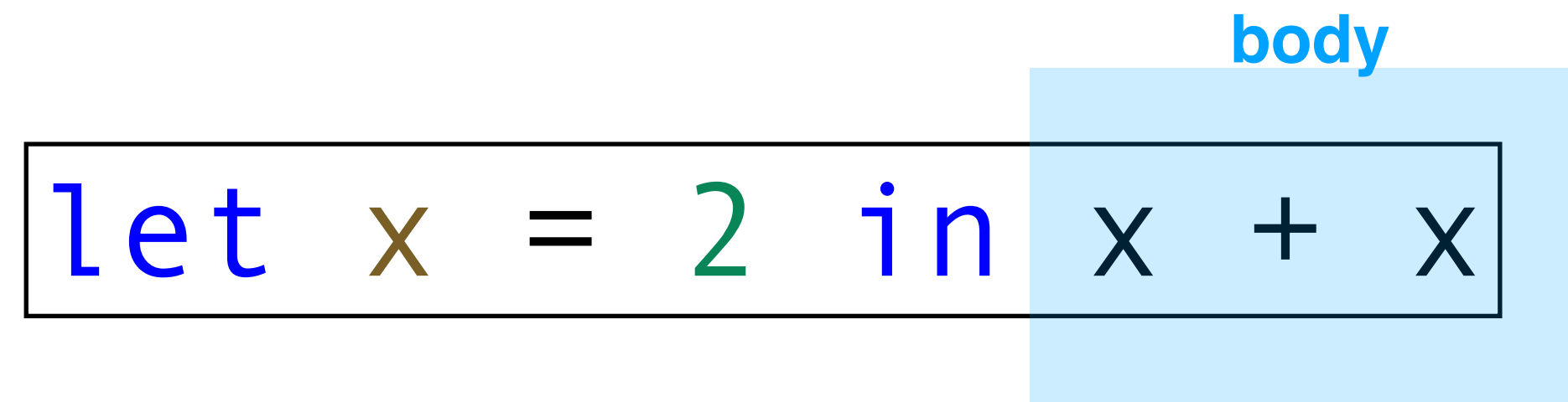
Let Expressions and Variables

body

```
let x = 2 in x + x
```

The diagram shows the code 'let x = 2 in x + x'. The text 'body' is written in blue above a light blue rectangular box. This box highlights the expression 'x + x' at the end of the line, which is the body of the let expression. The entire line of code is enclosed in a thin black rectangular border.

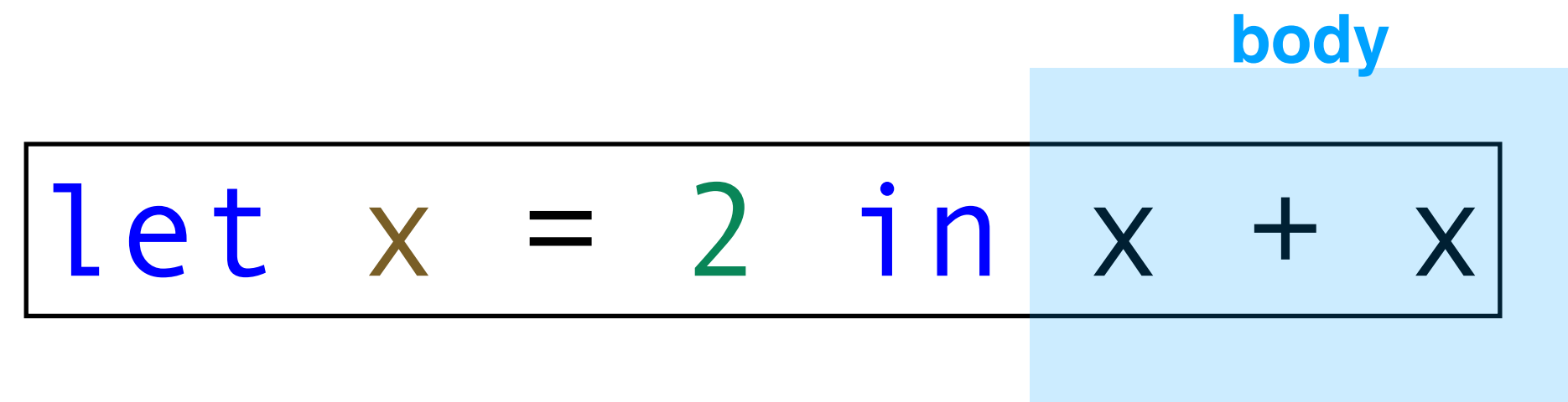
Let Expressions and Variables



The diagram shows the code `let x = 2 in x + x` enclosed in a thin black rectangular border. The word `let` is blue, `x` is brown, `=` is black, `2` is green, `in` is blue, `x` is black, `+` is black, and `x` is black. A light blue rectangular box highlights the expression `x + x`. The word `body` in blue is positioned above the right side of this light blue box.

As with any reasonable PL, we can define local variables in OCaml

Let Expressions and Variables

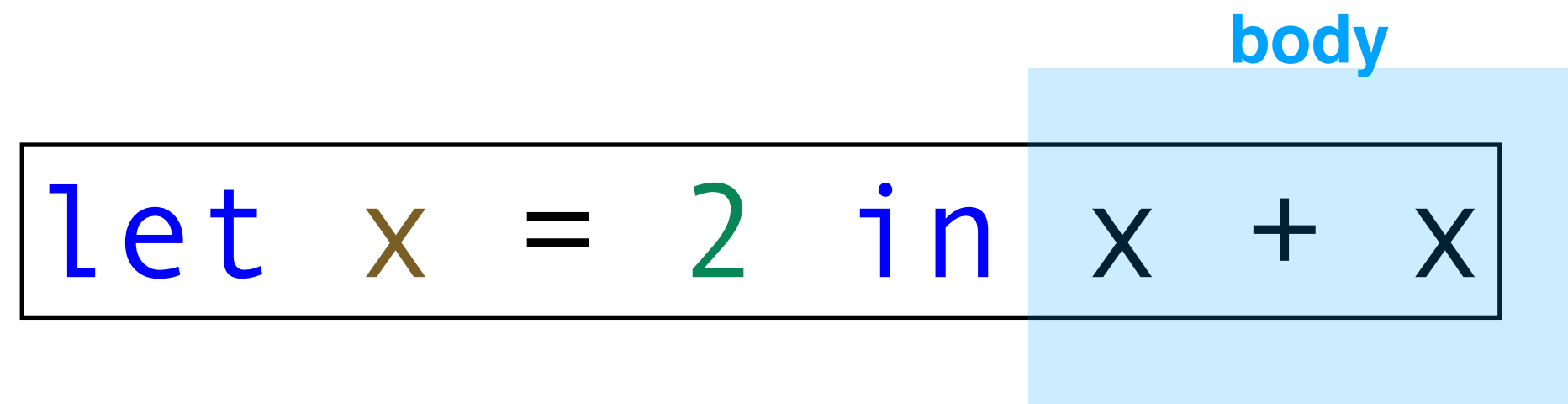


The diagram shows a code snippet `let x = 2 in x + x` enclosed in a black rectangular border. The word `let` is blue, `x` is brown, `=` is black, `2` is green, `in` is blue, `x` is black, `+` is black, and `x` is black. A light blue rectangular box highlights the expression `x + x`. The word `body` is written in blue text above the right side of this box.

As with any reasonable PL, we can define local variables in OCaml

This is useful for writing better abstractions

Let Expressions and Variables



The diagram shows the code `let x = 2 in x + x` enclosed in a black rectangular box. A light blue rectangular box highlights the expression `x + x` on the right side of the `in` keyword. The word `body` is written in blue text above the light blue box.

As with any reasonable PL, we can define local variables in OCaml

This is useful for writing better abstractions

Note that it reads like a sentence: *let x stand for 2 in the expression x + x*

Multiple Local Variables

```
let x_squared = x * x in  
let y_squared = y * y in  
x_squared + y_squared
```

It's very easy to use multiple local variables, we just *nest* local variables

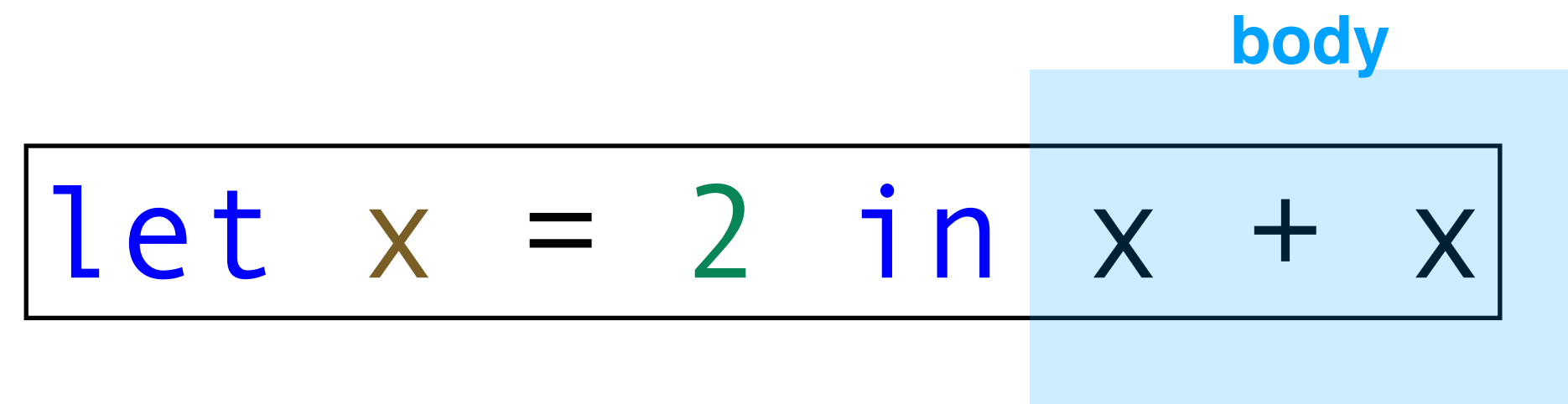
(If it helps, think of in as a semicolon ;)

IMPORTANT: `let x = e1 in e2` is an *expression* so it can be the body of a let expression.

Let Expressions (Informal)

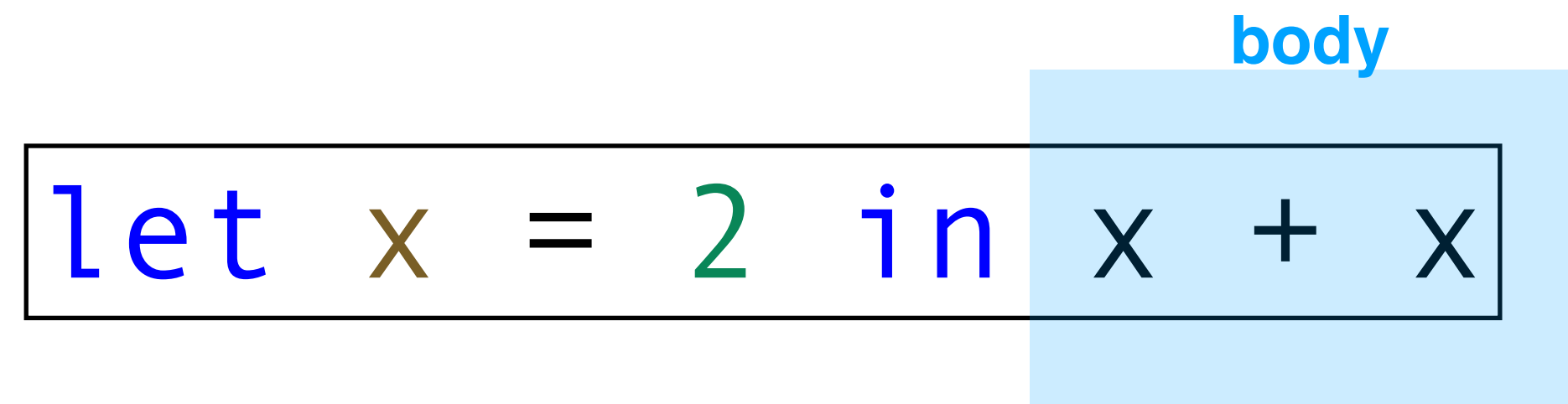
body

```
let x = 2 in x + x
```



The diagram illustrates the structure of a let expression. The text "let x = 2 in x + x" is shown within a black rectangular border. The portion "x + x" is highlighted by a light blue rectangular background. Above this blue area, the word "body" is written in blue text, indicating that the expression following the binding is the body of the let expression.

Let Expressions (Informal)

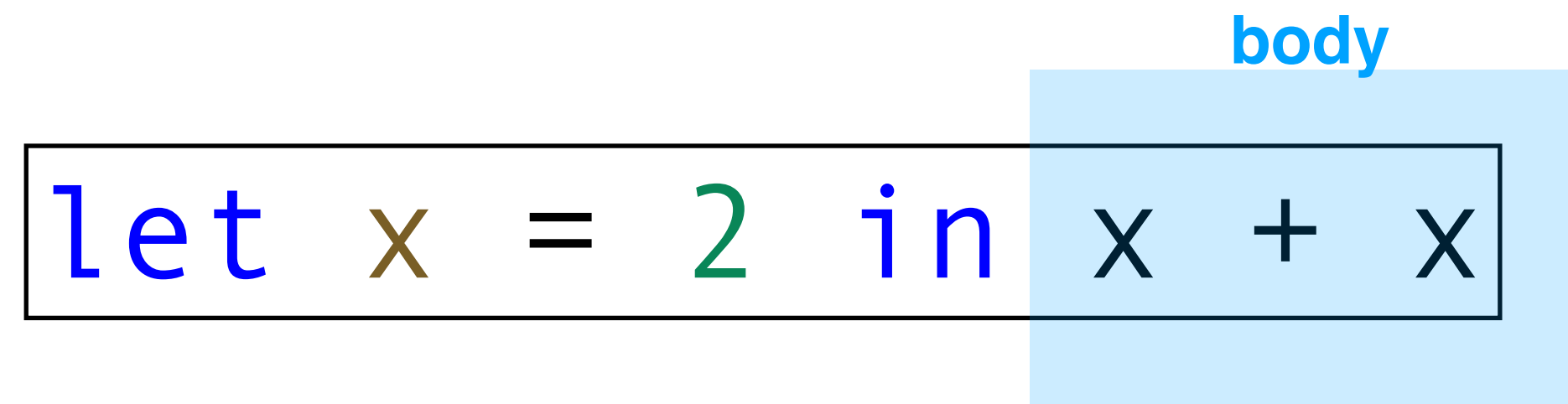


The diagram shows the code `let x = 2 in x + x` enclosed in a thin black rectangular border. A light blue rectangular background highlights the entire expression. The word `body` is written in blue text above the right side of the expression, specifically above the `x + x` part.

```
let x = 2 in x + x
```

syntax: `let VARIABLE = EXPRESSION in BODY`

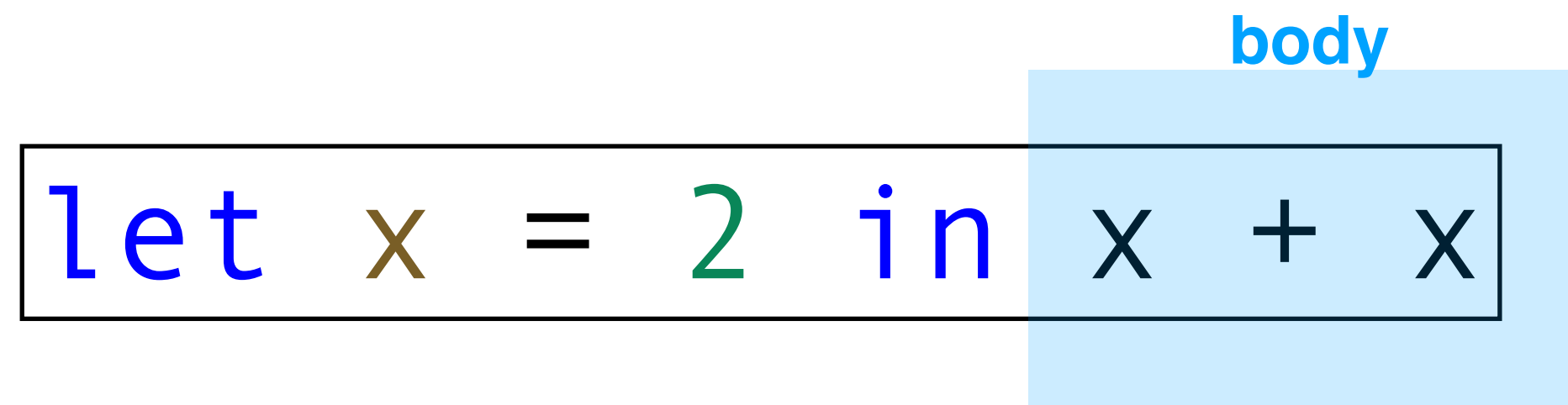
Let Expressions (Informal)

`let x = 2 in x + x`

syntax: `let VARIABLE = EXPRESSION in BODY`

typing: *compute type of* EXPRESSION; *assume*
VARIABLE *has that type; compute type of* BODY

Let Expressions (Informal)

`let x = 2 in x + x`

syntax: `let VARIABLE = EXPRESSION in BODY`

typing: *compute type of* EXPRESSION; *assume*
VARIABLE *has that type; compute type of* BODY

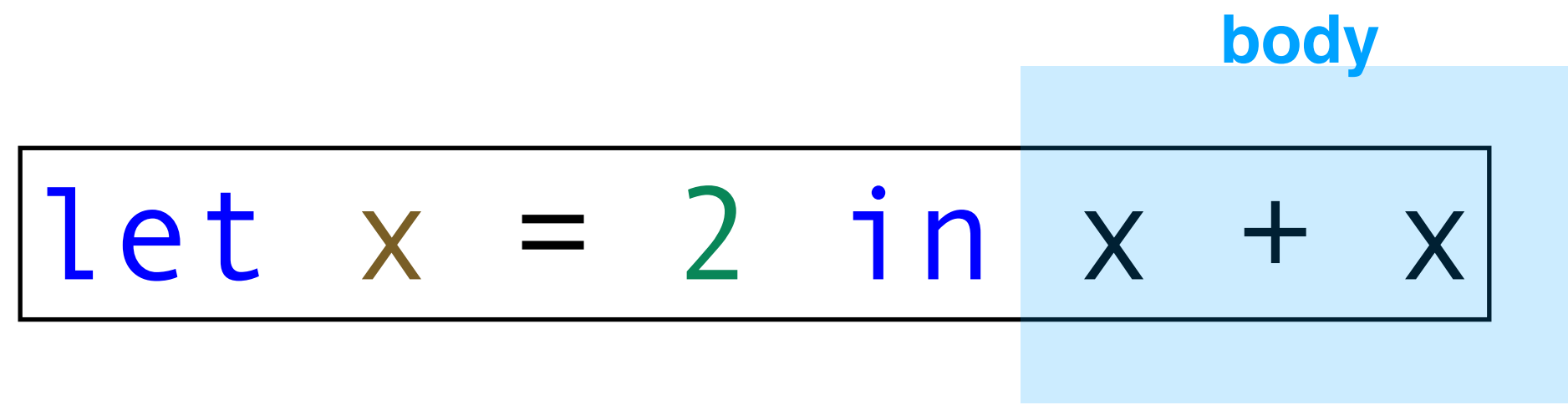
semantics: *compute value of* EXPRESSION; *substitute*
that value for VARIABLE *in* BODY

Syntax Examples

Syntax Examples

body

```
let x = 2 in x + x
```

A diagram illustrating the syntax of a `let` expression. The code `let x = 2 in x + x` is shown. The `let` and `in` keywords are blue, `x` is brown, `=` is black, and `2` is green. A light blue rectangular box highlights the expression `x + x`, which is the body of the `let` binding. The word `body` is written in blue above the box.

Syntax Examples

body

```
let x = 2 in x + x
```

A diagram illustrating the syntax of a 'let' expression. The text 'let x = 2 in x + x' is shown. The words 'let', 'x', '=', and '2' are in blue, 'in' is in green, and 'x + x' is in black. A light blue rectangular box highlights the 'x + x' part, with the word 'body' in blue text positioned above it.

body

```
let x = true in if x then 3 else 4
```

A diagram illustrating the syntax of a 'let' expression. The text 'let x = true in if x then 3 else 4' is shown. The words 'let', 'x', '=', and 'true' are in blue, 'in' is in green, and 'if x then 3 else 4' is in black. A light blue rectangular box highlights the 'if x then 3 else 4' part, with the word 'body' in blue text positioned above it.

Syntax Examples

body

```
let x = 2 in x + x
```

body

```
let x = true in if x then 3 else 4
```

body

```
let x = 3.5 in 2. * . x
```

Syntax Examples

let *x* = 2 in *x* + *x*

body

let *x* = true in if *x* then 3 else 4

body

let *x* = 3.5 in 2. * . *x*

body

let *y* = (let *x* = 2 in *x* + *x*) in 4 * *y*

body

body

Example: Ill-Typed Let-Expression

```
let x = 2. in 3 + x
```

An ill-typed expression will throw a type error when you type it into `utop`

Note that the body of a let-expression may be ill-typed *depending on the value assigned to its variable*

Semantics Examples

Semantics Examples

`let x = 2 in x + x` \longrightarrow `2 + 2`

Semantics Examples

`let x = 2 in x + x` \longrightarrow `2 + 2`

`let x = true in
if x then 3 else 4` \longrightarrow `if true then 3 else 4`

Semantics Examples

`let x = 2 in x + x` \longrightarrow `2 + 2`

`let x = true in
if x then 3 else 4` \longrightarrow `if true then 3 else 4`

Formally, we write $[v/x]e$ to mean "substitute v for x in e ",
e.g., $[3/x](x + x)$ is the same as $3 + 3$

Semantics Examples

`let x = 2 in x + x` \longrightarrow `2 + 2`

`let x = true in
if x then 3 else 4` \longrightarrow `if true then 3 else 4`

Formally, we write $[v/x]e$ to mean "substitute v for x in e ",
e.g., $[3/x](x + x)$ is the same as $3 + 3$

Intuitively, substitution is simple: **replace the variable**

Semantics Examples

`let x = 2 in x + x` \longrightarrow `2 + 2`

`let x = true in
if x then 3 else 4` \longrightarrow `if true then 3 else 4`

Formally, we write $[v/x]e$ to mean "substitute v for x in e ",
e.g., $[3/x](x + x)$ is the same as $3 + 3$

Intuitively, substitution is simple: **replace the variable**

Turns out, this is **very hard** to do correctly, *it's subtle*
and a source of a lot of mistakes in PL implementations

demo

(simple use of lets)

Second Abstraction of the Day

If Expressions

If-Expressions

```
if x > 0 then x else -x
```

```
if x <> y then x+y else x-y
```

If-Expressions

```
if x > 0 then x else -x
```

```
if x <> y then x+y else x-y
```

As with any reasonable PL, OCaml has expressions for doing conditional reasoning

If-Expressions

```
if x > 0 then x else -x
```

```
if x <> y then x+y else x-y
```

As with any reasonable PL, OCaml has expressions for doing conditional reasoning

Note: The **else** case is *required* and the **then** and **else** cases must be the *same type* (why?)

If-Expressions

```
if x < 0 then
  "negative"
else if x = 0 then
  "zero"
else
  "positive"
```

Answer: Remember, all we have is expressions. So every if-expression must have a value and a type (and therefore, an **else** case of the same type)

We can do **else if** just by nesting if-expressions! (neat)

If-Expressions (Informal)

```
if x > 0 then x else -x
```


If-Expressions (Informal)

```
if x > 0 then x else -x
```

Syntax: if CONDITION then TRUE-CASE else FALSE-CASE

If-Expressions (Informal)

```
if x > 0 then x else -x
```

Syntax: if CONDITION then TRUE-CASE else FALSE-CASE

Typing: CONDITION *must be a Boolean; compute the types of*
TRUE-CASE *and* FALSE-CASE; *must be the same type; expression*
type is same as that of TRUE-CASE *and* FALSE-CASE

If-Expressions (Informal)

```
if x > 0 then x else -x
```

Syntax: if CONDITION then TRUE-CASE else FALSE-CASE

Typing: CONDITION *must be a Boolean; compute the types of* TRUE-CASE *and* FALSE-CASE; *must be the same type; expression type is same as that of* TRUE-CASE *and* FALSE-CASE

Semantics: *If* CONDITION *evaluates to true; evaluate* TRUE-CASE, *else evaluate* FALSE-CASE

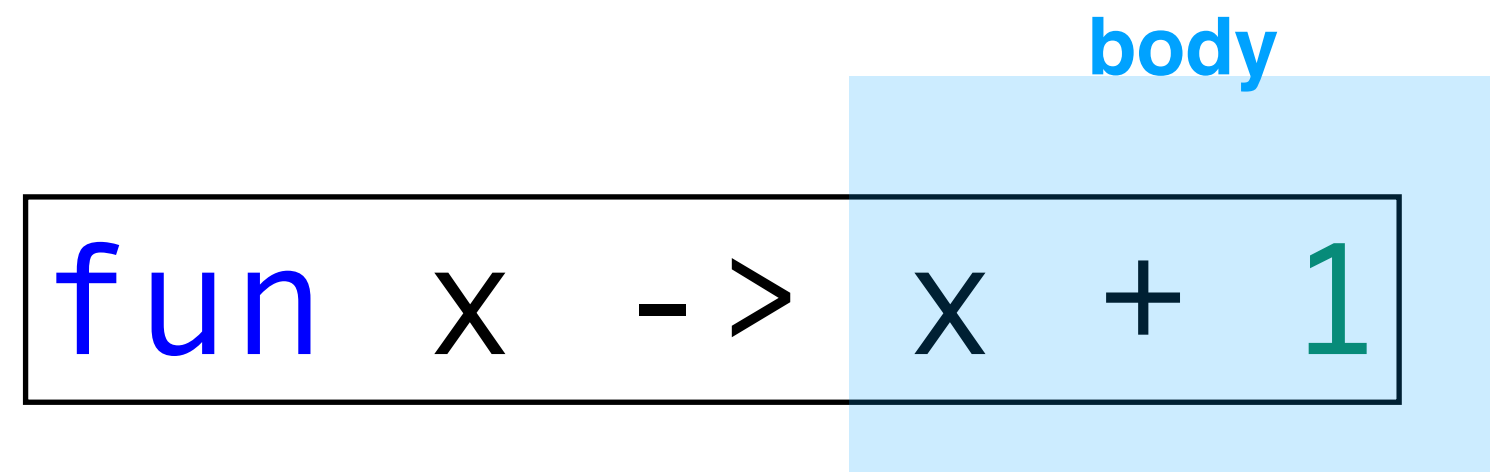
demo

(simple use of ifs)

Third (and Most Important) Abstraction of the Day

Functions

Functions (Informal)



fun x -> x + 1

The diagram shows the function definition `fun x -> x + 1`. The text is enclosed in a black rectangular box. A light blue rectangular box highlights the expression `x + 1` on the right side of the definition. Above the light blue box, the word `body` is written in blue text.

Syntax: fun VARIABLE -> EXPRESSION

Typing: assume VARIABLE has type **T1**; compute the type of EXPRESSION; suppose it is **T2**; type of function is **T1 -> T2**

Semantics: A function is a value; evaluates to itself

Functions are also Expressions

`fun x -> 3 + x`

body

`fun y -> 2. * . x`

body

`fun x y z -> if x then y else z`

body

`fun x -> if x > 0 then x else -x`

body

In OCaml, we can define *anonymous functions*, which are just **functions without names**

We Can Give them Names using Let

let f = fun x -> 3 + x

body

let g = fun y -> 2. * . x

body

let h = fun x y z -> if x then y else z

body

let abs = fun x -> if x > 0 then x else -x

body

Another Way to Define Functions

```
let abs = fun x -> if x > 0 then x else -x
```

```
let h = fun x y z -> if x then y else z
```

Another Way to Define Functions

```
let abs = fun x -> if x > 0 then x else -x
```

```
let abs x = if x > 0 then x else -x
```

```
let h = fun x y z -> if x then y else z
```

Another Way to Define Functions

```
let abs = fun x -> if x > 0 then x else -x
```

```
let abs x = if x > 0 then x else -x
```

```
let h = fun x y z -> if x then y else z
```

```
let h x y z = if x then y else z
```

Another Way: Curried Functions

```
let h = fun x y z -> if x then y else z
```

```
let h = fun x -> fun y -> fun z -> if x then y else z
```

Another way of thinking about functions:

The only kind of function we have is *single argument*

This seems restrictive, but ultimately it doesn't affect us at all

We can *simulate* multi-argument functions with nested functions. This is called **Currying** after Haskell Curry

Curried Functions Return Functions

```
let f = fun x -> fun y -> fun z -> x + y + z
```

We should think of the above function as something which takes an input and returns **another function**

In other words, we *partially apply* the function

Application

```
(fun x -> fun y -> x + y + 1) 3 2
```

Application is done by *juxtaposition* which means we put the arguments next to the function

Application is *left-associative*, which means we pass arguments from left to right

Application (Informally)

```
(fun x -> fun y -> x + y + 1) 3 2
```

Application (Informally)

```
(fun x -> fun y -> x + y + 1) 3 2
```

Syntax: FUNCTION-EXPR ARG-EXPR

Application (Informally)

```
(fun x -> fun y -> x + y + 1) 3 2
```

Syntax: FUNCTION-EXPR ARG-EXPR

Typing: *compute type of* FUNCTION-EXPR; *say it is*
T1 *->* **T2**; *compute type of* ARG-EXPR; *it must be* **T1**;
then the type of expression is **T2**

Application (Informally)

```
(fun x -> fun y -> x + y + 1) 3 2
```

Syntax: FUNCTION-EXPR ARG-EXPR

Typing: *compute type of FUNCTION-EXPR; say it is T1 -> T2; compute type of ARG-EXPR; it must be T1; then the type of expression is T2*

Semantics: *Evaluate ARG-EXPR; substitute its value into the body of FUNCTION-EXPR; evaluate the result*

Application (Example)

Application (Example)

```
(fun x -> fun y -> x + y + 1) 3 2
```

Application (Example)

`(fun x -> fun y -> x + y + 1) 3 2`

is the same as

`(fun x -> (fun y -> x + y + 1)) 3 2`

Application (Example)

`(fun x -> fun y -> x + y + 1) 3 2`

is the same as

`(fun x -> (fun y -> x + y + 1)) 3 2`

is the same as

`((fun x -> (fun y -> x + y + 1)) 3) 2`

Application (Example)

(fun x -> fun y -> x + y + 1) 3 2

is the same as

(fun x -> (fun y -> x + y + 1)) 3 2

is the same as

((fun x -> (fun y -> x + y + 1)) 3) 2

Application (Example)

(fun x -> fun y -> x + y + 1) 3 2

is the same as

(fun x -> (fun y -> x + y + 1)) 3 2

is the same as

((fun x -> (fun y -> x + y + 1)) 3) 2

evaluates to

(fun y -> 3 + y + 1) 2

Application (Example)

(fun x -> fun y -> x + y + 1) 3 2

is the same as

(fun x -> (fun y -> x + y + 1)) 3 2

is the same as

((fun x -> (fun y -> x + y + 1)) 3) 2

evaluates to

(fun y -> 3 + y + 1) 2

Application (Example)

(fun x -> fun y -> x + y + 1) 3 2

is the same as

(fun x -> (fun y -> x + y + 1)) 3 2

is the same as

((fun x -> (fun y -> x + y + 1)) 3) 2

evaluates to

(fun y -> 3 + y + 1) 2

evaluates to

3 + 2 + 1

Application (Example)

`(fun x -> fun y -> x + y + 1) 3 2`

is the same as

`(fun x -> (fun y -> x + y + 1)) 3 2`

is the same as

`((fun x -> (fun y -> x + y + 1)) 3) 2`

evaluates to

`(fun y -> 3 + y + 1) 2`

evaluates to

`3 + 2 + 1`

is the same as

`(3 + 2) + 1`

Application (Example)

`(fun x -> fun y -> x + y + 1) 3 2`

is the same as

`(fun x -> (fun y -> x + y + 1)) 3 2`

is the same as

`((fun x -> (fun y -> x + y + 1)) 3) 2`

evaluates to

`(fun y -> 3 + y + 1) 2`

evaluates to

`3 + 2 + 1`

is the same as

`(3 + 2) + 1`

evaluates to

`5 + 1`

Application (Example)

(fun x -> fun y -> x + y + 1) 3 2

is the same as

(fun x -> (fun y -> x + y + 1)) 3 2

is the same as

((fun x -> (fun y -> x + y + 1)) 3) 2

evaluates to

(fun y -> 3 + y + 1) 2

evaluates to

3 + 2 + 1

is the same as

(3 + 2) + 1

evaluates to

5 + 1

evaluates to

6

demo

(anonymous and curried functions)

Summary

OCaml is a language of **expressions**, everything is an expression

OCaml has everything we need to do basic programming