Higher Order Programming: Folds

Concepts of Programming Languages Lecture 8

Outline

- » Look at one more common HOF in detail: fold_left (and fold_right)
- >> Look at HOFs on data types other than lists

Practice Problem

Implement the function

val smallest_prime_factor : int -> int

so that **smallest_prime_factor n** is the smallest prime factor of **n** if **n > 1**

Use this to define the predicate **p** such that **List.filter p l** returns the elements of **l** which are the product of two distinct primes

Recap

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- 3. passed as arguments to another function
- Note. Types are not first-class values

In OCaml, functions are **first-class values**They can be:

- 1. returned by another function
- 2. given names with <u>let-definitions</u>
- 3. passed as arguments to another function

Note. Types are not first-class values

Recall: Functions as Parameters

```
# let apply f x = f x;;
val apply : ('a -> 'b) -> 'a -> 'b = <fun>
# apply add_five 10;;
- : int = 15
```

This is very interesting in OCaml...

This allows us to create new functions which are parametrized by old ones

Recall: Functions as Parameters

```
# let apply f x = f x;;
val apply : ('a -> 'b) -> 'a -> 'b = <fun>
# apply add_five 10;;
- : int = 15
```

This is very interesting in OCaml...

This allows us to create new functions which are parametrized by old ones

```
let rec fact n =
   match n with
   | 0 -> 1
   | n -> n * fact (n - 1)

let rec sum n =
   match n with
   | 0 -> 0
   | n -> n + sum (n - 1)
```

Some functions cannot be polymorphic

But can we still abstract the core functionality?

```
let rec fact n =
   match n with
   | 0 -> 1
   | n -> n * fact (n - 1)

let rec sum n =
   match n with
   | 0 -> 0
   | n -> n + sum (n - 1)
```

Some functions cannot be polymorphic

But can we still abstract the core functionality?

```
let rec accum f n start =
  let rec go n =
    match n with
    | 0 -> start
    | n -> f n (go (n - 1))
  in go n
```

In order to generalize this function, we need to be able to take the operation as a parameter

Now we have a single function which we can reuse elsewhere

```
let rec accum f n start =
  let rec go n =
    match n with
    | 0 -> start
    | n -> f n (go (n - 1))
  in go n
```

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Now we have a single function which we can reuse elsewhere

Recall: Definition of Map

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» If the list is empty there is nothing to do

Recall: Definition of Map

- » If the list is empty there is nothing to do
- » If the list is nonempty, we apply f to its
 first element, and recurse

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

» If the list is empty there is nothing to do

```
let rec filter p l =
  match l with
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  | x :: xs ->
    (if p x then [x] else []) @ filter p xs
```

- » If the list is empty there is nothing to do
- » If the first element satisfies our predicate we keep
 it and recurse

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

- » If the list is empty there is nothing to do
- » If the first element satisfies our predicate we keep
 it and recurse
- >> Otherwise, we drop it and recurse

Folds

map

```
transform each element (keep every element)
```

```
map
    transform each element (keep every
    element)

filter    keep some elements based on a
    predicate
```

```
map
    transform each element (keep every
    element)

filter    keep some elements based on a
    predicate

fold    combine elements via an accumulation
    function
```

```
let rec concat ls =
  match ls with
  | [] -> []
  | xs :: xss -> xs @ concat xss
```

```
let map f l =
  let rec go l =
    match l with
    | [] -> []
    | x :: xs -> (f x) :: go xs
in go l
```

```
let rec sum l =
  match l with
  [] -> 0
  | x :: xs -> x + sum xs
  base
```

```
let rec rev l =
  match l with
  | [] -> []
  | x :: xs -> rev xs @ [x]
  base
```

```
let rec sum l =
  match l with
  [] -> 0
  | x :: xs -> x + sum xs
  base rec. call
```

```
let rec rev l =
  match l with
  | [] -> []
  | x :: xs -> rev xs @ [x]
  base rec. call
```

```
let map f l =
  let rec go l =
    match l with
    | [] -> []
    | x :: xs -> (f x) :: go xs
in go l base
    rec. call
    combine
```

Fold as Specialized Pattern Matching

```
let rec sum l =
  let base = 0 in
  match l with
  | [] -> base
  | x :: xs -> x + sum xs
```

```
let rec sum l =
  let base = 0 in
  match l with
  | [] -> base
  | x :: xs -> x + sum xs
```

```
let rec sum l =
  let base = 0 in
  let op = (+) in
  match l with
  | [] -> base
  | x :: xs -> op x (sum xs)
```

```
let rec sum l =
  let base = 0 in
  let op = (+) in

match l with
  | [] -> base
  | x :: xs -> op x (sum xs)
```

```
let sum l =
  let base = 0 in
  let op = (+) in
  let rec go l =
    match l with
    [] -> base
x::xs-> op x (go xs)
  in go l
```

```
let sum l =
  let base = 0 in
  let op = (+) in
 let rec go l =
    match l with
    [] -> base
    x :: xs \rightarrow op x (go xs)
  in go l
                    fold right
```

```
let sum l =
  let base = 0 in
  let op = (+) in
  List.fold_right op l base
```

```
let sum l = List.fold_right (+) l 0
```

```
let sum l = List.fold_right (+) l 0
```

We get a one-liner for **sum** (and a whole lot of other functions)

Folds are very nice for "iterating" over a list

```
1 :: (2 :: (3 :: (4 :: (5 :: (6 :: (7 :: []))))))
\downarrow sum = fold_right (+) l 0
1 + (2 + (3 + (4 + (5 + (6 + (7 + 0))))))
```

We can think of fold_right as "replacing" :: with + and [] with 0

```
1 :: (2 :: (3 :: (4 :: (5 :: (6 :: (7 :: []))))))
\downarrow \text{ prod} = \text{fold\_right (*) l 1}
1 * (2 * (3 * (4 * (5 * (6 * (7 * 1))))))
```

We can think of $fold_right$ as "replacing" :: with * and [] with 1

```
[1] :: ([2] :: ([3] :: ([4] :: ([5] :: ([6] :: ([7] :: []))))))

concat = fold_right (@) l []

[1] @ ([2] @ ([3] @ ([4] @ ([5] @ ([6] @ ([7] @ []))))))
```

We can think of **fold_right** as "replacing" :: with @ and [] with []

We can think of fold_right as "replacing" :: with op and [] with base

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

>> On empty, return the base element

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

- >> On empty, return the base element
- » On nonempty, recurse on the tail and apply op to the head and the result

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l

Is this tail recursive?
```

- >> On empty, return the base element
- » On nonempty, recurse on the tail and apply op to the head and the result

Understanding Check

```
Write filter using List.fold_right
Write append (@) using List.fold_right
     append (10:020:30:11) (4:15:16:1)
              W W U

1:: 2::3::(4::5::6::[])
```

demo

(tail recursive fold attempt)

Tail-Recursive Fold Attempt

```
let fold_right_tr op l base =
  let rec go l acc =
    match l with
  | [] -> acc
  | x :: xs -> go xs (op acc x)
in go l base
```

Can you see what's wrong with this definition?



```
fold_right (+) [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
```

```
fold_right (+) [1;2;3] 0 ===

1 + fold_right (+) [2;3] 0 ===

1 + (2 + fold_right (+) [3] 0) ===

1 + (2 + (3 + fold_right (+) [] 0)) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
1 + 5
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
go [] (((0 + 1) + 2) + 3) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
go [] (((0 + 1) + 2) + 3) ====
((0 + 1) + 2) + 3
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
1 + 5
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
go [3] ((0 + 1) + 2) ===
go [] (((0 + 1) + 2) + 3) ===
((0 + 1) + 2) + 3 ===
(1 + 2) + 3 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
go [3] ((0 + 1) + 2) ===
go [] (((0 + 1) + 2) + 3) ===
((0 + 1) + 2) + 3 ===
(1 + 2) + 3 ===
3 + 3 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1)  ===
go [3] ((0 + 1) + 2)  ===
go [] (((0 + 1) + 2) + 3)  ===
((0 + 1) + 2) + 3  ===
(1 + 2) + 3  ===
6
```



```
fold_right (-) [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
1 - (-1)
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 - 1) ====
go [3] ((0 - 1) - 2) ====
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 - 1) ====
go [3] ((0 - 1) - 2) ====
go [] (((0 - 1) - 2) - 3) ====
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3 ===
(-3) - 3
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3 ===
(-3) - 3 ===
-6
```

$$1 - (2 - (3 - 0))$$

$$((0-1)-2)-3$$

Changing parentheses is fine for (+) but not for (-)

Associativity

Definition: A binary operation $\square: A \times A \to A$ is associative if it satisfies $a\square(b\square c) = (a\square b)\square c$ for any $a,b,c \in A$

Example: Addition and multiplication are associative, whereas subtraction and division are not

Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

Folding left is just our incorrect tail recursive right folding (with a change in the order of arguments)

Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

Folding left is just our incorrect tail recursive right folding (with a change in the order of arguments)

```
fold_left is a left-associative fold
fold_right is a right-associative fold
```

```
1:: (2:: (3:: (4:: [])))
fold_left op base l
                   op 1 (op 2 (op 3 (op 4 base)))
op (op (op (op base 1) 2) 3) 4
```

```
1:: (2:: (3:: (4:: [])))
fold_left op base l
                    op 1 (op 2 (op 3 (op 4 base)))
op (op (op base 1) 2) 3) 4
```

Tail-Recursive Fold Right

```
let fold_right_tr op l base =
  List.fold_left
    (fun x y -> op y x)
    base
    (List.rev l)
```

We can write fold_right in terms of fold left by reversing the list and "reversing" the operation

Challenge: Write a tail-recursive fold right without reversing the list

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

$$1 - r (2 - r (3 - r (4 - r 0)))$$

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

$$1 - r (2 - r (3 - r (4 - r 0)))$$

$$= 1 - r (2 - r (3 - r (0 - 4)))$$

```
Let x - r y := y - x, subtraction with the arguments flipped
```

```
1 - r (2 - r (3 - r (4 - r 0)))
= 1 - r (2 - r (3 - r (0 - 4)))
= 1 - r (2 - r ((0 - 4) - 3))
```

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

```
1 - r (2 - r (3 - r (4 - r 0)))
= 1 - r (2 - r (3 - r (0 - 4)))
= 1 - r (2 - r ((0 - 4) - 3))
= 1 - r (((0 - 4) - 3) - 2)
```

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= (((0 - 4) - 3) - 2) - 1
```

```
let rec all bs =
  match bs with
    | [] -> true
    | false :: _ -> false
    | true :: t -> all t

let all = List.fold_left (&&) true
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Which is better?

fold_left has to traverse the entire list, it can't short-circuit But the fold code is shorter and arguably clearer...

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- » For associative operations, use fold_left
- The types are difficult to remember, let the compiler remind you
- » Don't use folds for everything, but also don't
 use pattern matching for everything. Think
 about the use case

Understanding Check List fold left

let sort le l=

In terms of fold_left implement the function

val sort : ('a -> 'a -> bool) -> 'a list -> 'a list

so that **sort le l** is the list **l** in sorted order according to **le**

fl op [] [
$$z;1;3$$
] \Rightarrow

fl op (op [] z) [$1;3$] \Rightarrow

fl op [1] [$1;3$] \Rightarrow

fl op [1] [$1;3$] \Rightarrow

fl op (op [1] 1) [3] \Rightarrow

fl op [$1;z$] [3] \Rightarrow

fl op (op ($1;z$] 3) [1] \Rightarrow

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Beyond Lists

Mappable Data

A lot of data types hold uniform kinds of data which can then be mapped over

Formally, these are called Functors

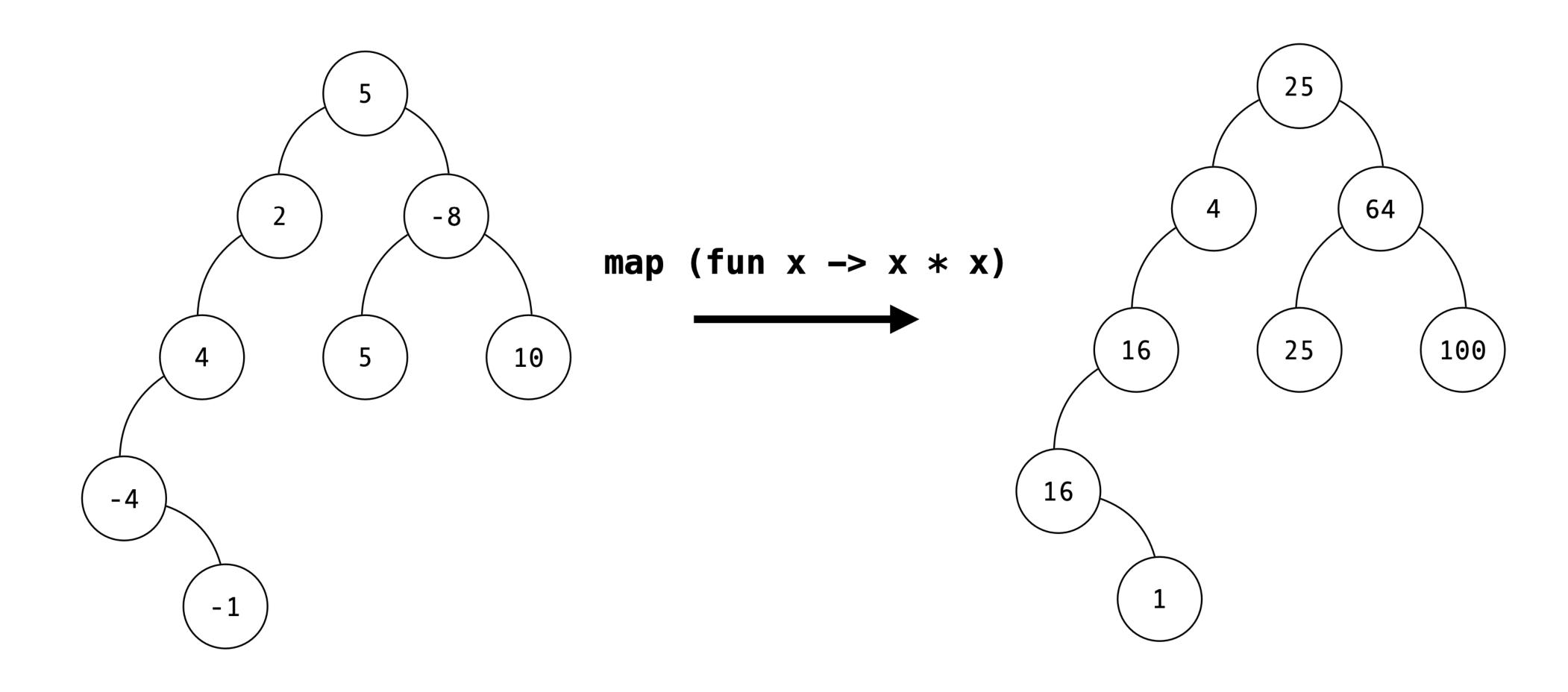
Trees

```
type 'a tree =
    | Leaf
    | Node of 'a * 'a tree * 'a tree

let map f t =
    let rec go t =
        match t with
    | Leaf -> Leaf
        | Node (x, l, r) -> Node (f x, go l, go r)
    in go t
```

Mapping over a tree maintains the structure but recursively updates values with **f**

The Picture



Options

On None, leave the None
On Some x, apply f to x

```
let mkMatrix (vals : 'a list list) : 'a matrix option = ...
let transpose (mx : 'a matrix) : 'a matrix = ...
let vals = ...
let a = Option.map transpose (mkMatrix vals)
```

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This is a very common pattern for working with options if we want to "keep computing" as long as the option still holds a value

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Map allows us to "lift" non-option functions to option functions

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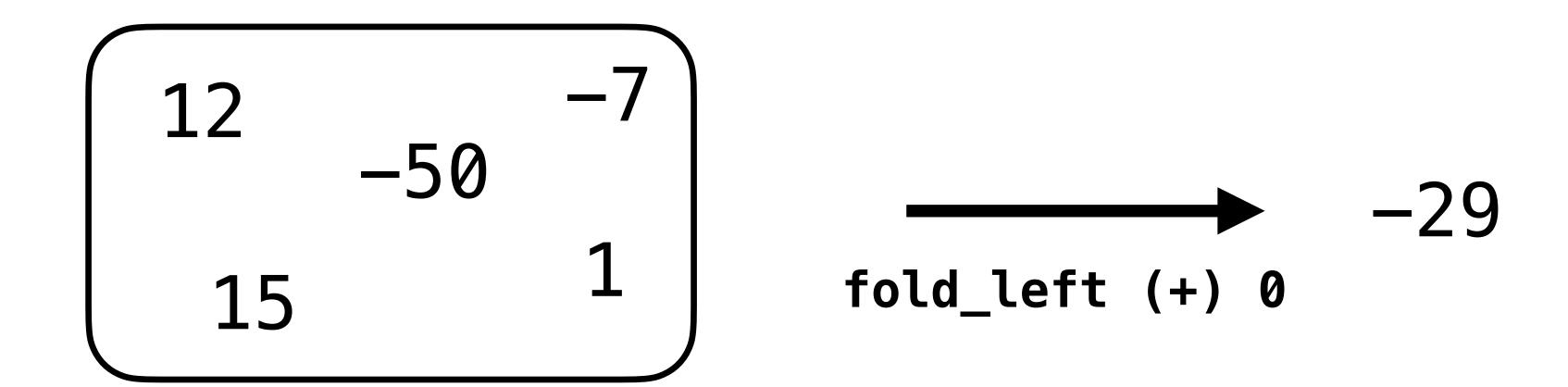
Map allows us to "lift" non-option functions to option functions

We can avoid pattern matching explicitly on options

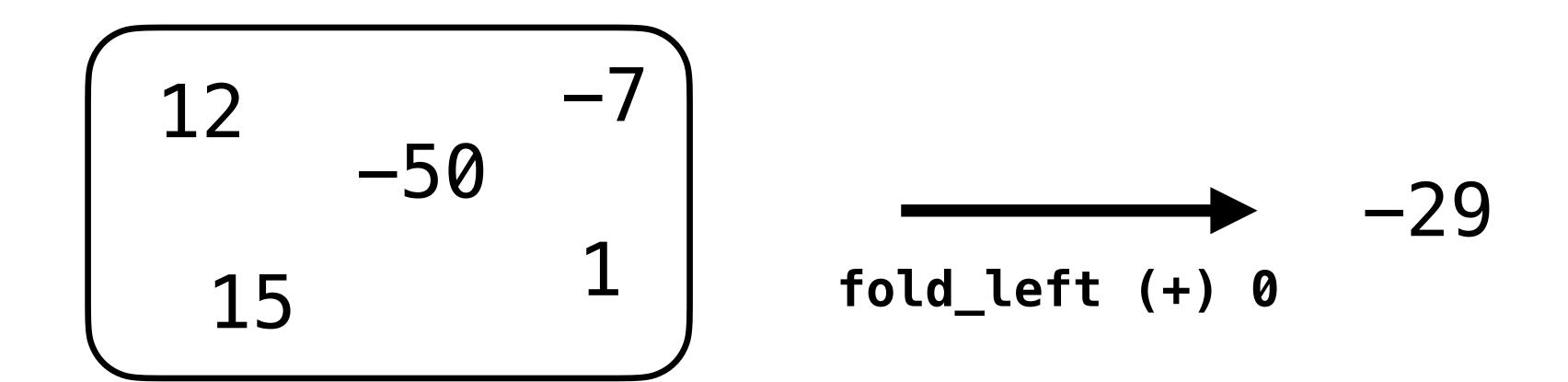
demo

(option mapping)

Foldable Data

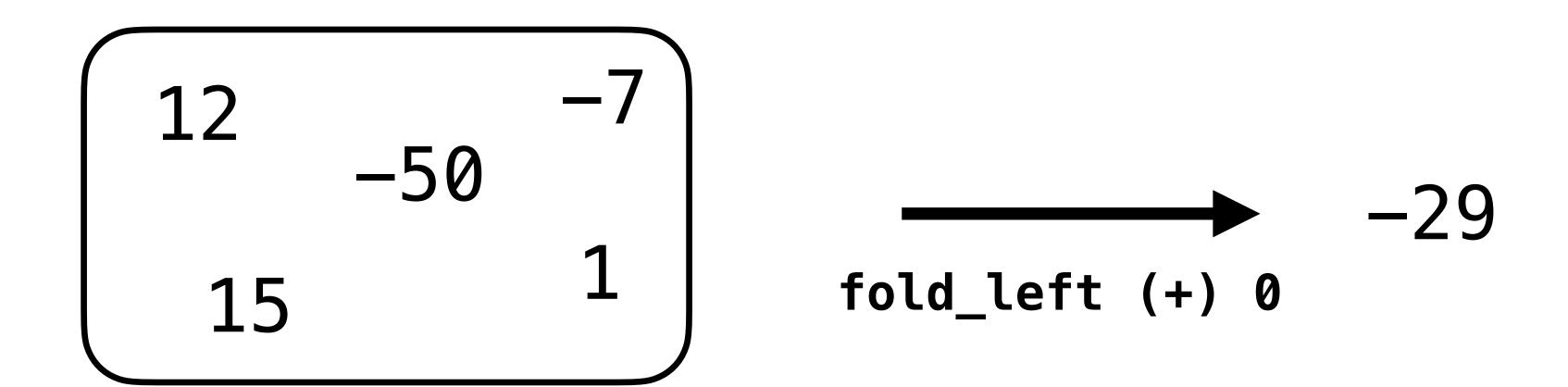


Foldable Data



There are also a lot of data types which hold uniform data that we might want to fold over

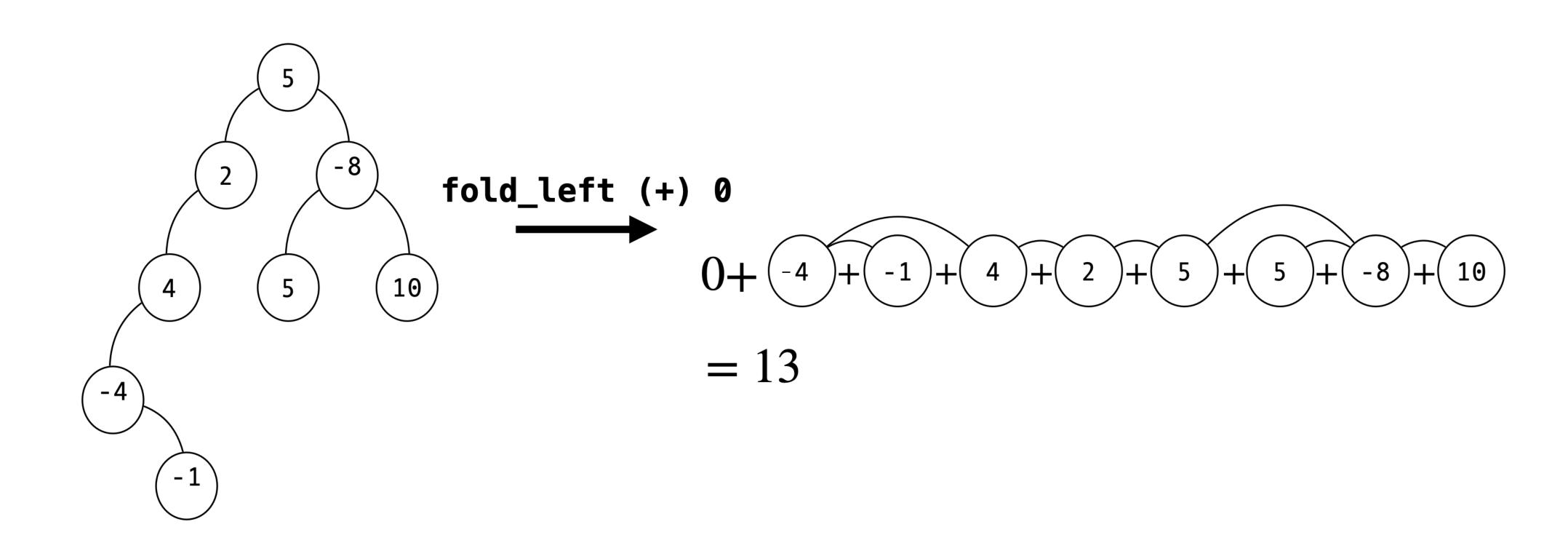
Foldable Data



There are also a lot of data types which hold uniform data that we might want to fold over

We have to deal with associativity and the order that elements are processed

Trees (The Picture)



```
let fold_left op base t =
  let rec go acc t=
  match t with
  | Leaf -> acc
  | Node (x, l, r) -> go (op (go acc l) x) r
  in go base t
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Understanding Check

Implement fold_left for ntrees

Summary

Folds are used to **combine** data with an accumulation function

The order that we combine things matters if the accumulation function is not associative

We can map and fold (and even filter) more than just lists