# Higher Order Programming: Folds

**Concepts of Programming Languages Lecture 8** 

#### Outline

- » Look at one more common HOF in detail: fold\_left (and fold\_right)
- >> Look at HOFs on data types other than lists

#### Practice Problem

Implement the function

val smallest\_prime\_factor : int -> int

so that **smallest\_prime\_factor n** is the smallest prime factor of **n** if **n > 1** 

Use this to define the predicate **p** such that **List.filter p l** returns the elements of **l** which are the product of two distinct primes

# Recap

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- 3. passed as arguments to another function
- Note. Types are not first-class values

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- 3. passed as arguments to another function

Note. Types are not first-class values

#### Recall: Functions as Parameters

```
# let apply f x = f x;;
val apply : ('a -> 'b) -> 'a -> 'b = <fun>
# apply add_five 10;;
- : int = 15
```

This is very interesting in OCaml...

This allows us to create new functions which are parametrized by old ones

#### Recall: Functions as Parameters

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# let apply f x = f x;;
val apply : ('a -> 'b) -> 'a -> 'b = <fun>
# apply add_five 10;;
- : int = 15
```

This is very interesting in OCaml...

This allows us to create new functions which are parametrized by old ones

```
let rec fact n =
   match n with
   | 0 -> 1
   | n -> n * fact (n - 1)

let rec sum n =
   match n with
   | 0 -> 0
   | n -> n + sum (n - 1)
```

Some functions cannot be polymorphic

But can we still abstract the core functionality?

```
let rec fact n =
   match n with
   | 0 -> 1
   | n -> n * fact (n - 1)

let rec sum n =
   match n with
   | 0 -> 0
   | n -> n + sum (n - 1)
```

Some functions cannot be polymorphic

But can we still abstract the core functionality?

```
let rec accum f n start =
  let rec go n =
    match n with
    | 0 -> start
    | n -> f n (go (n - 1))
  in go n
```

In order to generalize this function, we need to be able to take the operation as a parameter

Now we have a single function which we can reuse elsewhere

```
let rec accum f n start =
  let rec go n =
    match n with
    | 0 -> start
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  in go n
```

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#### Recall: Definition of Map

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» If the list is empty there is nothing to do

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- » If the list is empty there is nothing to do
- » If the list is nonempty, we apply f to its
  first element, and recurse

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

» If the list is empty there is nothing to do

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
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```

- » If the list is empty there is nothing to do
- » If the first element satisfies our predicate we keep
  it and recurse

```
let rec filter p l =
  match l with
  | [] -> []
  | x :: xs ->
     (if p x then [x] else []) @ filter p xs
```

- » If the list is empty there is nothing to do
- » If the first element satisfies our predicate we keep
  it and recurse
- >> Otherwise, we drop it and recurse

## Folds

map

```
transform each element (keep every element)
```

```
map
    transform each element (keep every
    element)

filter    keep some elements based on a
    predicate
```

```
map
    transform each element (keep every
    element)

filter    keep some elements based on a
    predicate

fold    combine elements via an accumulation
    function
```

```
let rec concat ls =
  match ls with
  | [] -> []
  | xs :: xss -> xs @ concat xss
```

```
let map f l =
  let rec go l =
    match l with
    | [] -> []
    | x :: xs -> (f x) :: go xs
in go l
```

```
let rec sum l =
  match l with
  [] -> 0
  | x :: xs -> x + sum xs
  base
```

```
let rec rev l =
  match l with
  | [] -> []
  | x :: xs -> rev xs @ [x]
  base
```

```
let rec sum l =
  match l with
  [] -> 0
  | x :: xs -> x + sum xs
  base rec. call
```

```
let rec rev l =
  match l with
  | [] -> []
  | x :: xs -> rev xs @ [x]
  base rec. call
```

```
let map f l =
  let rec go l =
    match l with
    | [] -> []
    | x :: xs -> (f x) :: go xs
in go l base
    rec. call
    combine
```

## Fold as Specialized Pattern Matching

```
let rec sum l =
  let base = 0 in
  match l with
  | [] -> base
  | x :: xs -> x + sum xs
```

```
let rec sum l =
  let base = 0 in
  match l with
  | [] -> base
  | x :: xs -> x + sum xs
```

```
let rec sum l =
  let base = 0 in
  let op = (+) in
  match l with
  | [] -> base
  | x :: xs -> op x (sum xs)
```

```
let rec sum l =
  let base = 0 in
  let op = (+) in

match l with
  | [] -> base
  | x :: xs -> op x (sum xs)
```

```
let sum l =
  let base = 0 in
  let op = (+) in
  let rec go l =
    match l with
    [] -> base
x::xs-> op x (go xs)
  in go l
```

```
let sum l =
  let base = 0 in
  let op = (+) in
 let rec go l =
    match l with
    [] -> base
    x :: xs \rightarrow op x (go xs)
  in go l
                    fold right
```

```
let sum l =
  let base = 0 in
  let op = (+) in
  List.fold_right op l base
```

```
let sum l = List.fold_right (+) l 0
```

```
let sum l = List.fold_right (+) l 0
```

We get a one-liner for **sum** (and a whole lot of other functions)

Folds are very nice for "iterating" over a list

```
1 :: (2 :: (3 :: (4 :: (5 :: (6 :: (7 :: []))))))
\downarrow sum = fold_right (+) l 0
1 + (2 + (3 + (4 + (5 + (6 + (7 + 0))))))
```

We can think of fold\_right as "replacing" :: with + and [] with 0

```
1 :: (2 :: (3 :: (4 :: (5 :: (6 :: (7 :: []))))))
\downarrow \text{ prod} = \text{fold\_right (*) l 1}
1 * (2 * (3 * (4 * (5 * (6 * (7 * 1))))))
```

We can think of  $fold_right$  as "replacing" :: with \* and [] with 1

```
[1] :: ([2] :: ([3] :: ([4] :: ([5] :: ([6] :: ([7] :: []))))))

concat = fold_right (@) l []

[1] @ ([2] @ ([3] @ ([4] @ ([5] @ ([6] @ ([7] @ []))))))
```

We can think of **fold\_right** as "replacing" :: with @ and [] with []

We can think of fold\_right as "replacing" :: with op and [] with base

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

>> On empty, return the base element

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l
```

- >> On empty, return the base element
- » On nonempty, recurse on the tail and apply op to the head and the result

```
let fold_right op l base =
  let rec go l =
    match l with
    | [] -> base
    | x :: xs -> op x (go xs)
  in go l

Is this tail recursive?
```

- >> On empty, return the base element
- » On nonempty, recurse on the tail and apply op to the head and the result

### Understanding Check

Write filter using List.fold\_right

Write append (@) using List.fold\_right

$$[1; 2; 3] @ [4; 5; 6] = [1; 2; 3; 4; 5; 6]$$

$$[0; 2; 3; 4; 5; 6] = [1; 2; 3; 4; 5; 6]$$

$$[1; 2; 3] @ [4; 5; 6] = [1; 2; 3; 4; 5; 6]$$

$$[1; 2; 3] @ [4; 5; 6] = [1; 2; 3; 4; 5; 6]$$

# demo

(tail recursive fold attempt)

### Tail-Recursive Fold Attempt

```
let fold_right_tr op l base =
  let rec go l acc =
    match l with
  | [] -> acc
  | x :: xs -> go xs (op acc x)
in go l base
```

Can you see what's wrong with this definition?



```
fold_right (+) [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
```

```
fold_right (+) [1;2;3] 0 ===

1 + fold_right (+) [2;3] 0 ===

1 + (2 + fold_right (+) [3] 0) ===

1 + (2 + (3 + fold_right (+) [] 0)) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
1 + 5
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
go [] (((0 + 1) + 2) + 3) ====
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 + 1) ====
go [3] ((0 + 1) + 2) ====
go [] (((0 + 1) + 2) + 3) ====
((0 + 1) + 2) + 3
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ===
1 + 5
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
go [3] ((0 + 1) + 2) ===
go [] (((0 + 1) + 2) + 3) ===
((0 + 1) + 2) + 3 ===
(1 + 2) + 3 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1) ===
go [3] ((0 + 1) + 2) ===
go [] (((0 + 1) + 2) + 3) ===
((0 + 1) + 2) + 3 ===
(1 + 2) + 3 ===
3 + 3 ===
```

```
fold_right (+) [1;2;3] 0 ===
1 + fold_right (+) [2;3] 0 ===
1 + (2 + fold_right (+) [3] 0) ===
1 + (2 + (3 + fold_right (+) [] 0)) ===
1 + (2 + (3 + 0)) ===
1 + (2 + 3) ====
6
```

```
fold_right_tr (+) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 + 1)  ===
go [3] ((0 + 1) + 2)  ===
go [] (((0 + 1) + 2) + 3)  ===
((0 + 1) + 2) + 3  ===
(1 + 2) + 3  ===
6
```



```
fold_right (-) [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
1 - (-1)
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 - 1) ====
go [3] ((0 - 1) - 2) ====
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ====
go [1;2;3] 0 ====
go [2;3] (0 - 1) ====
go [3] ((0 - 1) - 2) ====
go [] (((0 - 1) - 2) - 3) ====
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3 ===
(-3) - 3
```

```
fold_right (-) [1;2;3] 0 ===
1 - fold_right (-) [2;3] 0 ===
1 - (2 - fold_right (-) [3] 0) ===
1 - (2 - (3 - fold_right (-) [] 0)) ===
1 - (2 - (3 - 0)) ===
1 - (2 - 3) ===
2 ===
```

```
fold_right_tr (-) [1;2;3] 0 ===
go [1;2;3] 0 ===
go [2;3] (0 - 1) ===
go [3] ((0 - 1) - 2) ===
go [] (((0 - 1) - 2) - 3) ===
((0 - 1) - 2) - 3 ===
((-1) - 2) - 3 ===
(-3) - 3 ===
-6
```

$$1 - (2 - (3 - 0))$$

$$((0-1)-2)-3$$

#### Changing parentheses is fine for (+) but not for (-)

# Associativity

Definition: A binary operation  $\square: A \times A \to A$  is associative if it satisfies  $a\square(b\square c) = (a\square b)\square c$  for any  $a,b,c \in A$ 

Example: Addition and multiplication are associative, whereas subtraction and division are not

#### Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

#### Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

Folding left is just our incorrect tail recursive right folding (with a change in the order of arguments)

#### Definition of Fold Left

```
let fold_left op base l =
  let rec go l acc =
    match l with
    | [] -> acc
    | x :: xs -> go xs (op acc x)
  in go l base
```

Folding left is just our incorrect tail recursive right folding (with a change in the order of arguments)

```
fold_left is a left-associative fold
fold_right is a right-associative fold
```

```
1:: (2:: (3:: (4:: [])))
fold_left op base l
                   op 1 (op 2 (op 3 (op 4 base)))
op (op (op (op base 1) 2) 3) 4
```

```
1:: (2:: (3:: (4:: [])))
fold_left op base l
                    op 1 (op 2 (op 3 (op 4 base)))
op (op (op base 1) 2) 3) 4
```

# Tail-Recursive Fold Right

```
let fold_right_tr op l base =
  List.fold_left
    (fun x y -> op y x)
    base
    (List.rev l)
```

We can write fold\_right in terms of fold left by reversing the list and "reversing" the operation

**Challenge:** Write a tail-recursive fold right without reversing the list

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

$$1 - r (2 - r (3 - r (4 - r 0)))$$

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

$$1 - r (2 - r (3 - r (4 - r 0)))$$

$$= 1 - r (2 - r (3 - r (0 - 4)))$$

```
Let x - r y := y - x, subtraction with the arguments flipped
```

```
1 - r (2 - r (3 - r (4 - r 0)))
= 1 - r (2 - r (3 - r (0 - 4)))
= 1 - r (2 - r ((0 - 4) - 3))
```

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

```
1 - r (2 - r (3 - r (4 - r 0)))
= 1 - r (2 - r (3 - r (0 - 4)))
= 1 - r (2 - r ((0 - 4) - 3))
= 1 - r (((0 - 4) - 3) - 2)
```

```
Let x - r y := y - x, subtraction with
the arguments flipped
```

```
1 - r (2 - r (3 - r (4 - r 0)))
= 1 - r (2 - r (3 - r (0 - 4)))
= 1 - r (2 - r ((0 - 4) - 3))
= 1 - r (((0 - 4) - 3) - 2)
= (((0 - 4) - 3) - 2) - 1
```

```
let rec all bs =
  match bs with
    | [] -> true
    | false :: _ -> false
    | true :: t -> all t

let all = List.fold_left (&&) true
```

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fold\_left has to traverse the entire list, it can't short-circuit

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Which is better?

**fold\_left** has to traverse the entire list, it can't short-circuit But the fold code is shorter and arguably clearer...

» For associative operations, use fold\_left

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- » For associative operations, use fold\_left
- The types are difficult to remember, let the compiler remind you
- » Don't use folds for everything, but also don't
   use pattern matching for everything. Think
   about the use case

# Understanding Check

```
(fun acc x -> insert le x acc)
                   let rec insert le v l =
                    match l with
                      if le v x
                      then v:: l
                      else x :: insert le v l
In terms of fold_left implement the function
val sort : ('a -> 'a -> bool) ->/'a list -> 'a list
so that sort le l is the list l in sorted order
according to le
let sort le l = List. fold-left [] l
```

foldl op [] [2;3;1] fold op (op (2:3] 1 fold l op (op [] 2) [3;1] op [2] (3;1] op (op (23 3) 1;7;37 foldl op [7:3] (1)

# Beyond Lists

### Mappable Data

A lot of data types hold uniform kinds of data which can then be mapped over

Formally, these are called Functors

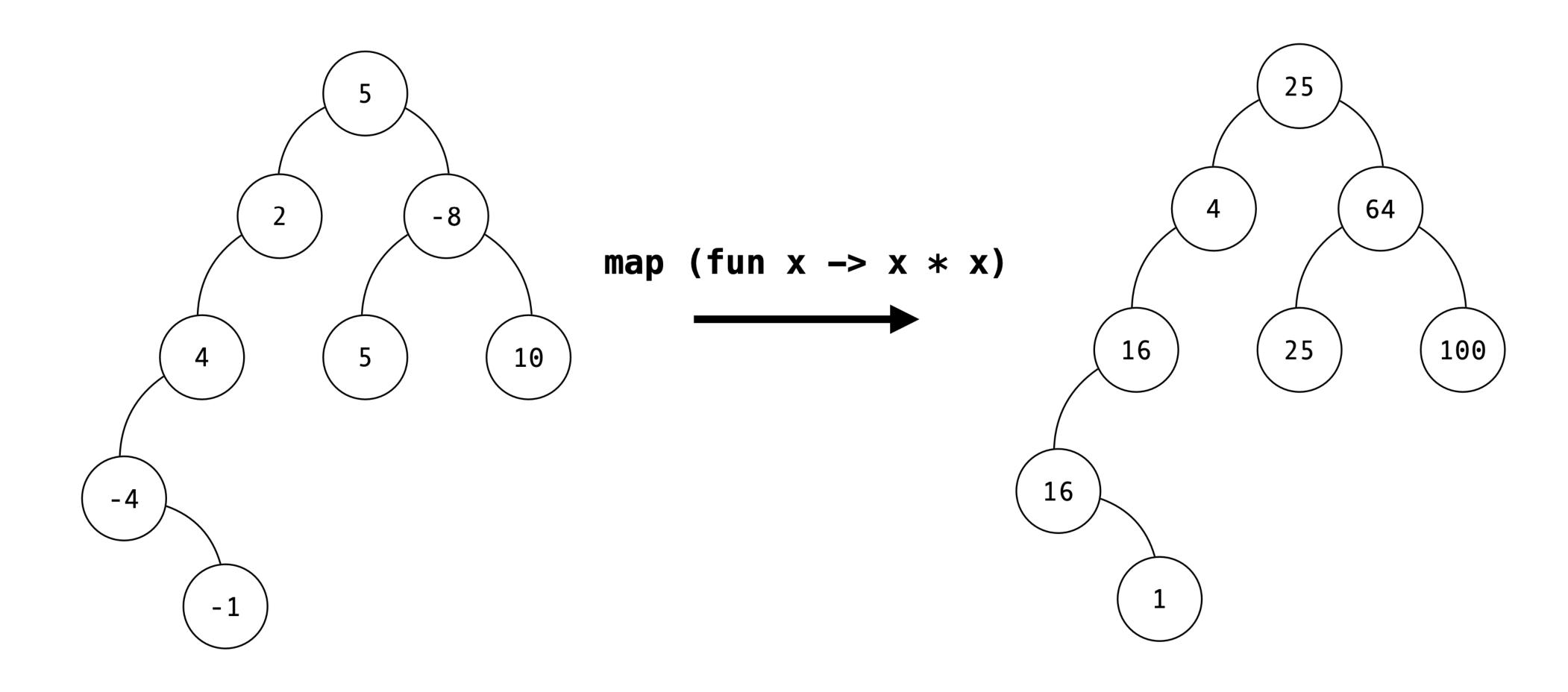
#### Trees

```
type 'a tree =
    | Leaf
    | Node of 'a * 'a tree * 'a tree

let map f t =
    let rec go t =
        match t with
    | Leaf -> Leaf
        | Node (x, l, r) -> Node (f x, go l, go r)
    in go t
```

Mapping over a tree maintains the structure but recursively updates values with **f** 

### The Picture



### Options

On None, leave the None
On Some x, apply f to x

```
let mkMatrix (vals : 'a list list) : 'a matrix option = ...
let transpose (mx : 'a matrix) : 'a matrix = ...
let vals = ...
let a = Option.map transpose (mkMatrix vals)
```

```
let mkMatrix (vals : 'a list list) : 'a matrix option = ...
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This is a very common pattern for working with options if we want to "keep computing" as long as the option still holds a value

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Map allows us to "lift" non-option functions to option functions

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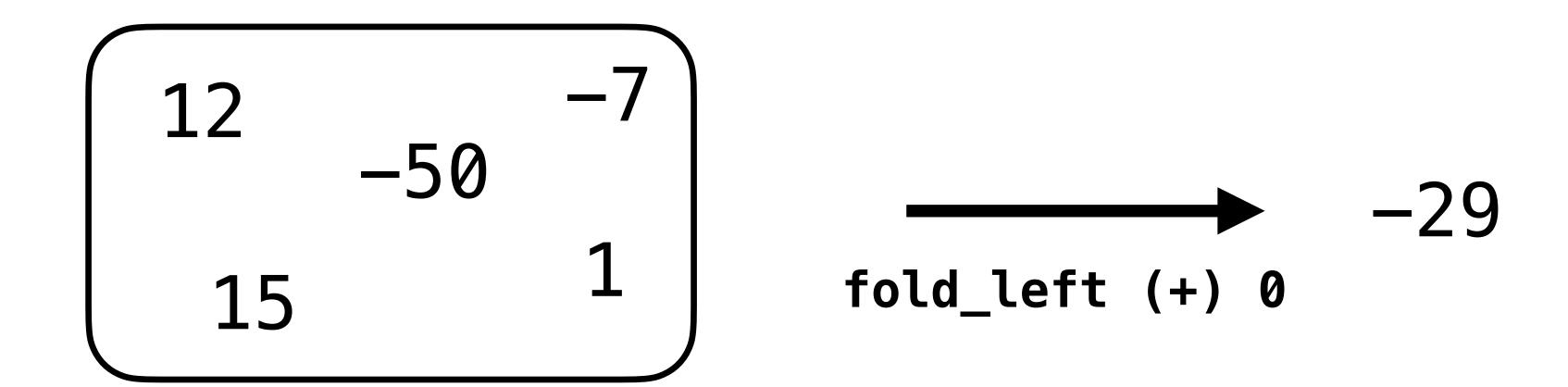
Map allows us to "lift" non-option functions to option functions

We can avoid pattern matching explicitly on options

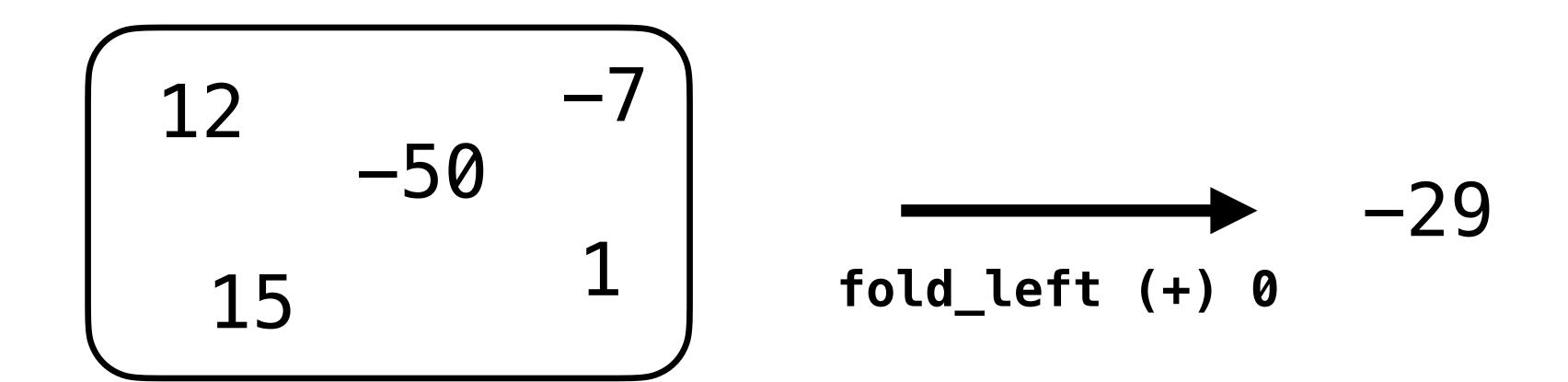
# demo

(option mapping)

### Foldable Data

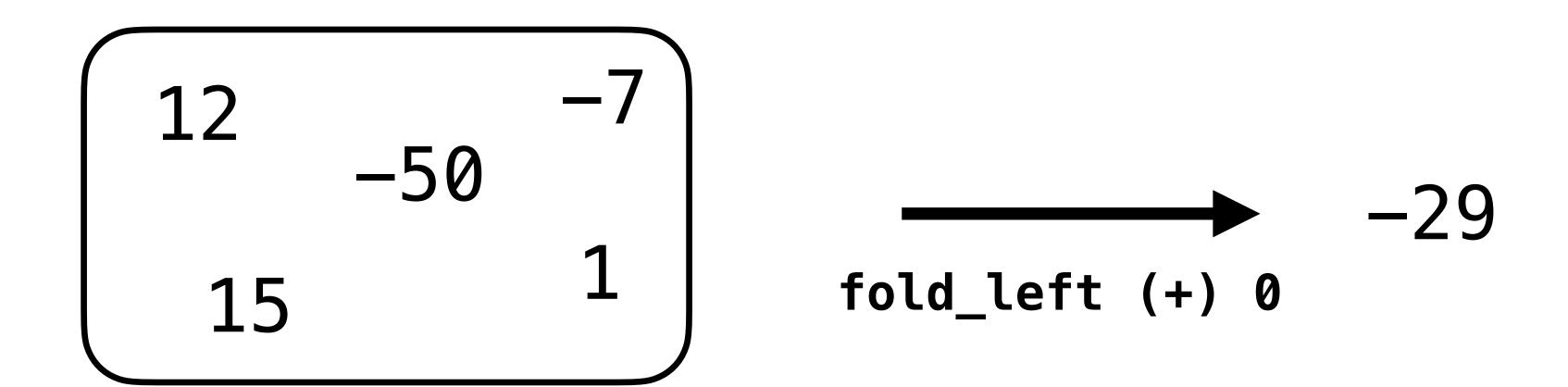


#### Foldable Data



There are also a lot of data types which hold uniform data that we might want to fold over

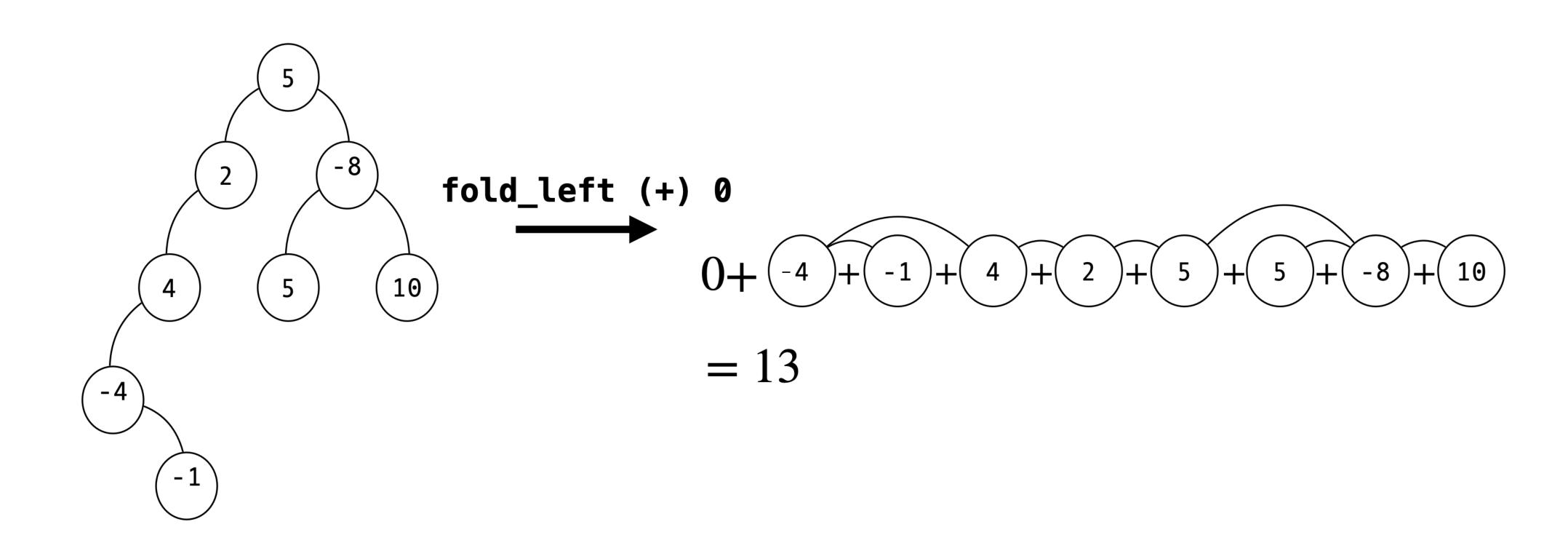
#### Foldable Data



There are also a lot of data types which hold uniform data that we might want to fold over

We have to deal with associativity and the order that elements are processed

# Trees (The Picture)



```
let fold_left op base t =
  let rec go acc t=
  match t with
  | Leaf -> acc
  | Node (x, l, r) -> go (op (go acc l) x) r
  in go base t
```

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### Understanding Check

Implement fold\_left for ntrees

### Summary

Folds are used to **combine** data with an accumulation function

The order that we combine things matters if the accumulation function is not associative

We can map and fold (and even filter) more than just lists