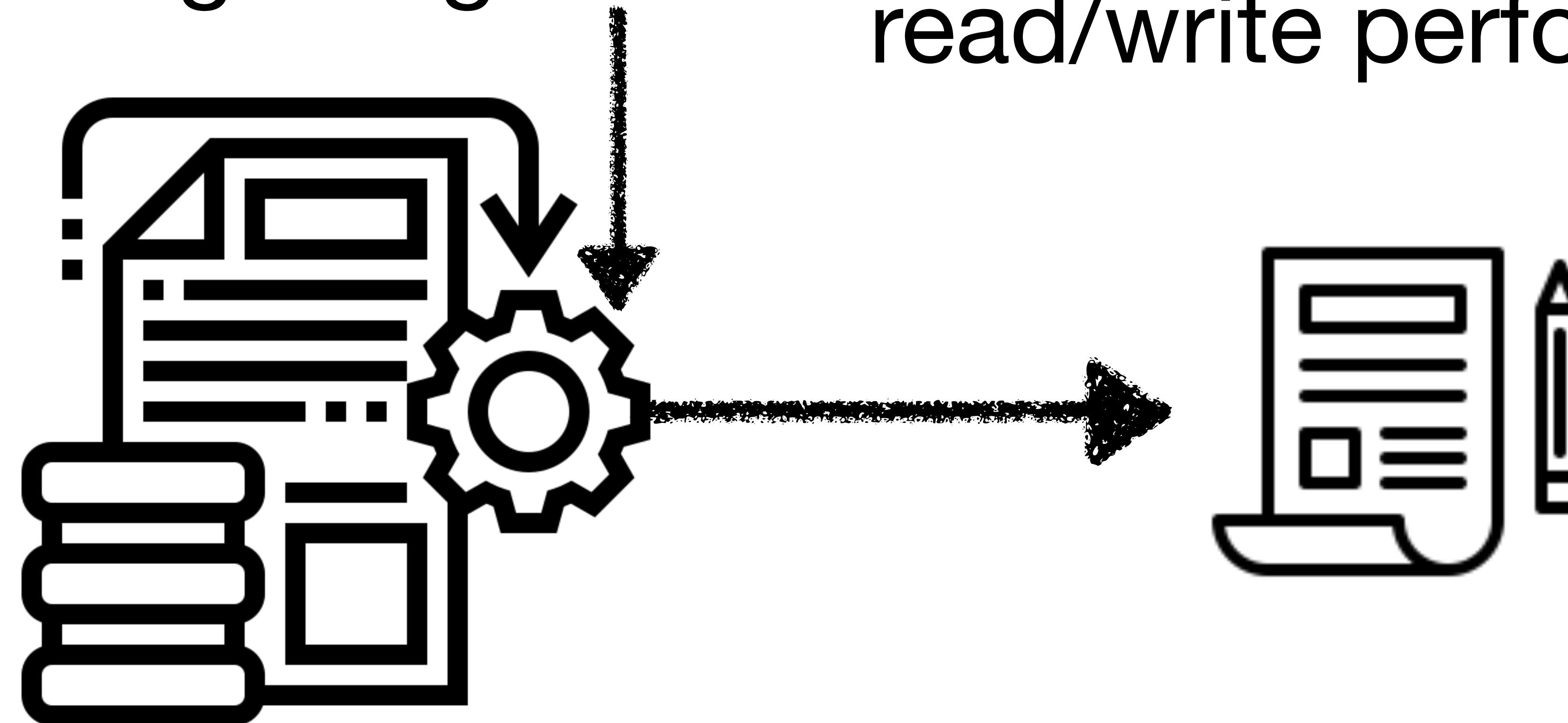


# Cosine: A Cloud-Cost Optimized Self-Designing Key-Value Storage Engine

Subarna Chatterjee, Meena Jagadeesan,  
Wilson Qin, Stratos Idreos

Data Systems Laboratory (DASLab)  
Harvard University

storage-engines



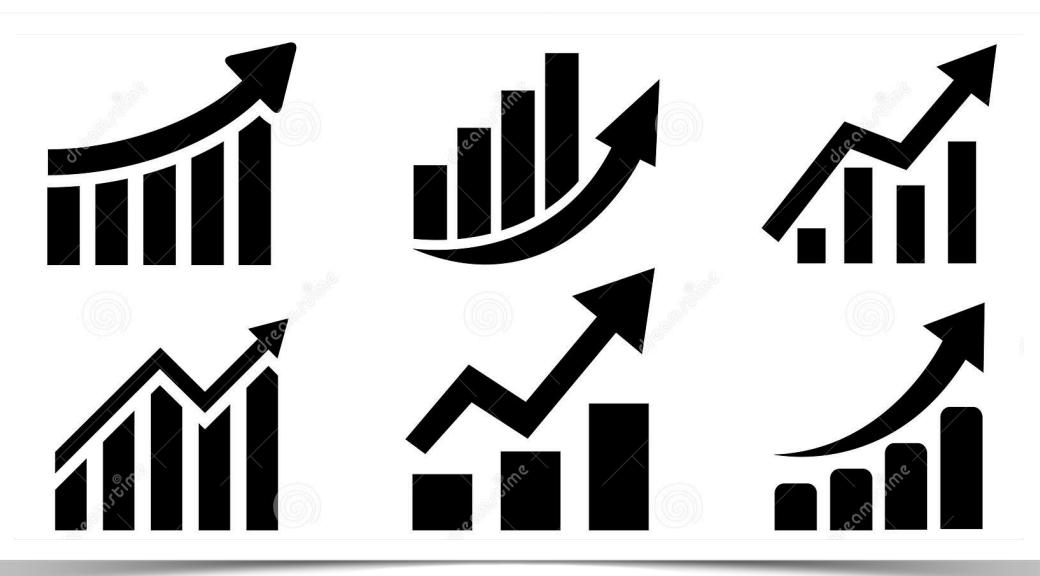
Data-systems

read/write performance

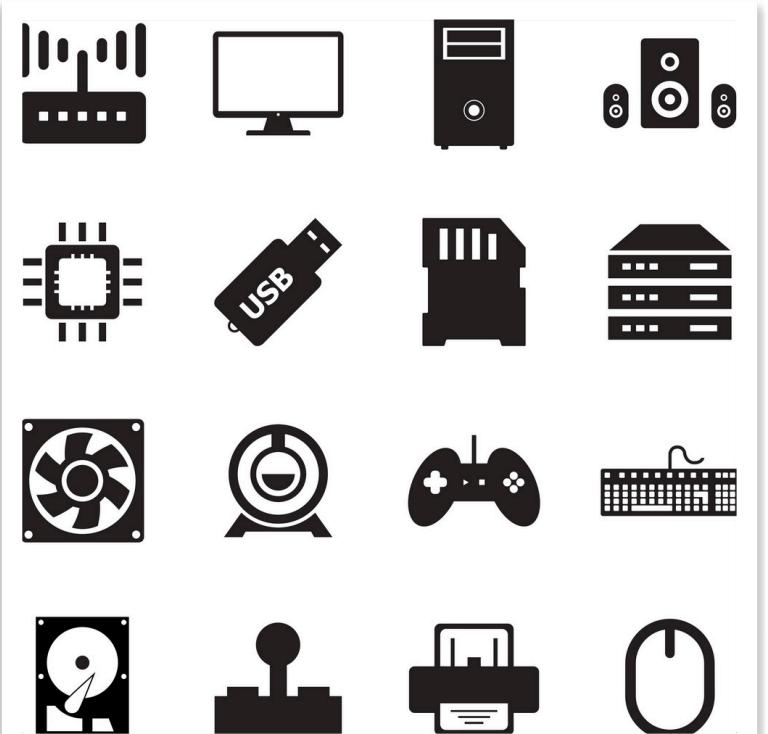


# CONTEXT

data



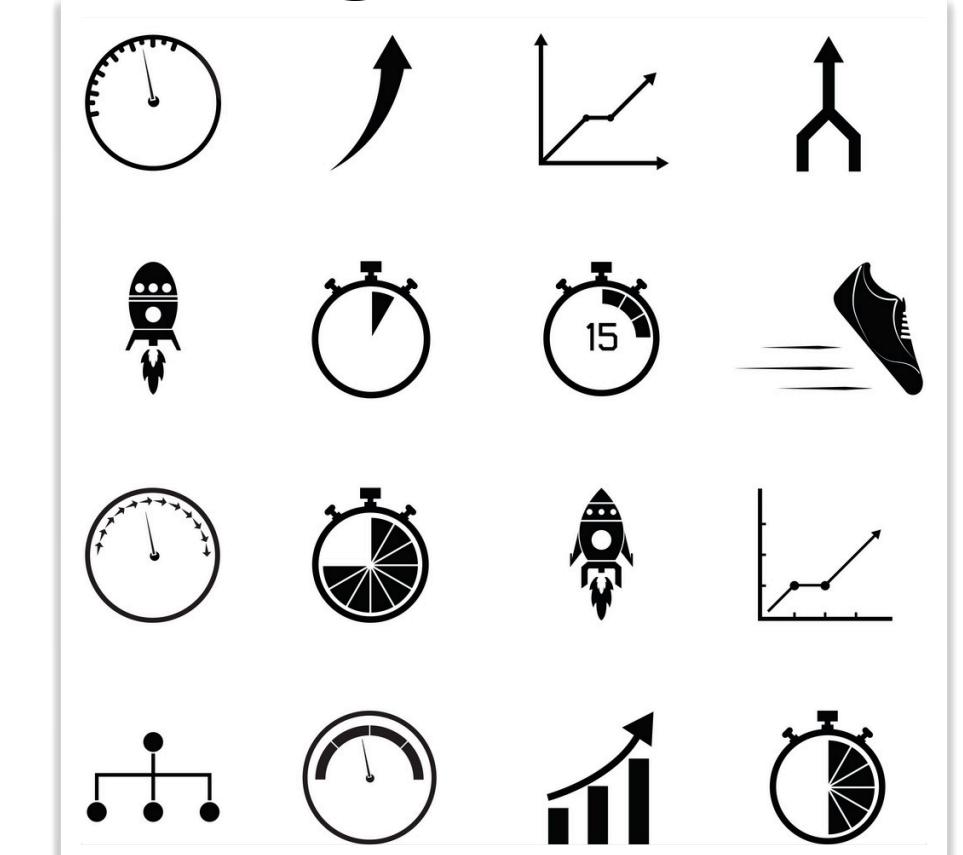
hardware



applications



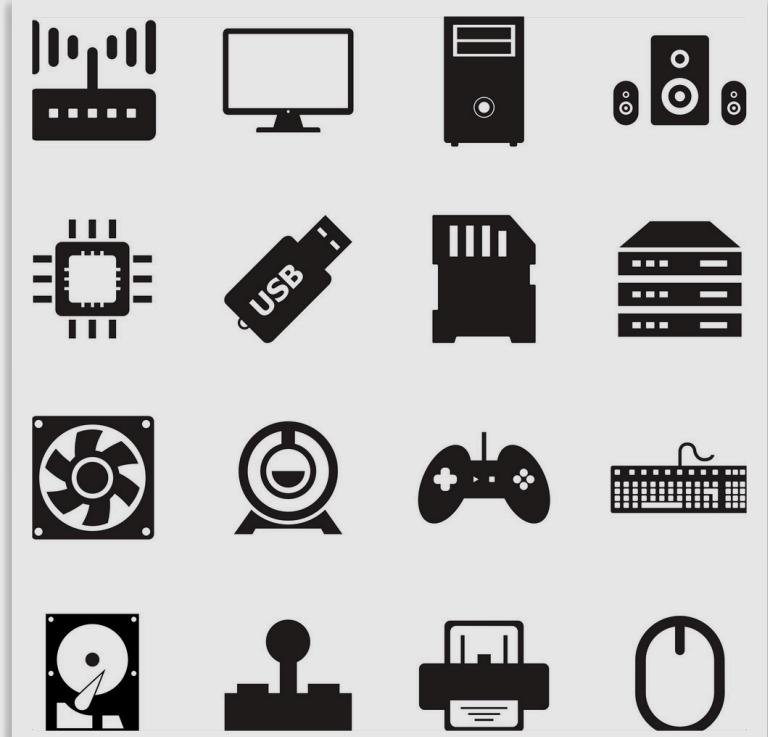
performance  
goals



# The CONTEXT keeps changing ...



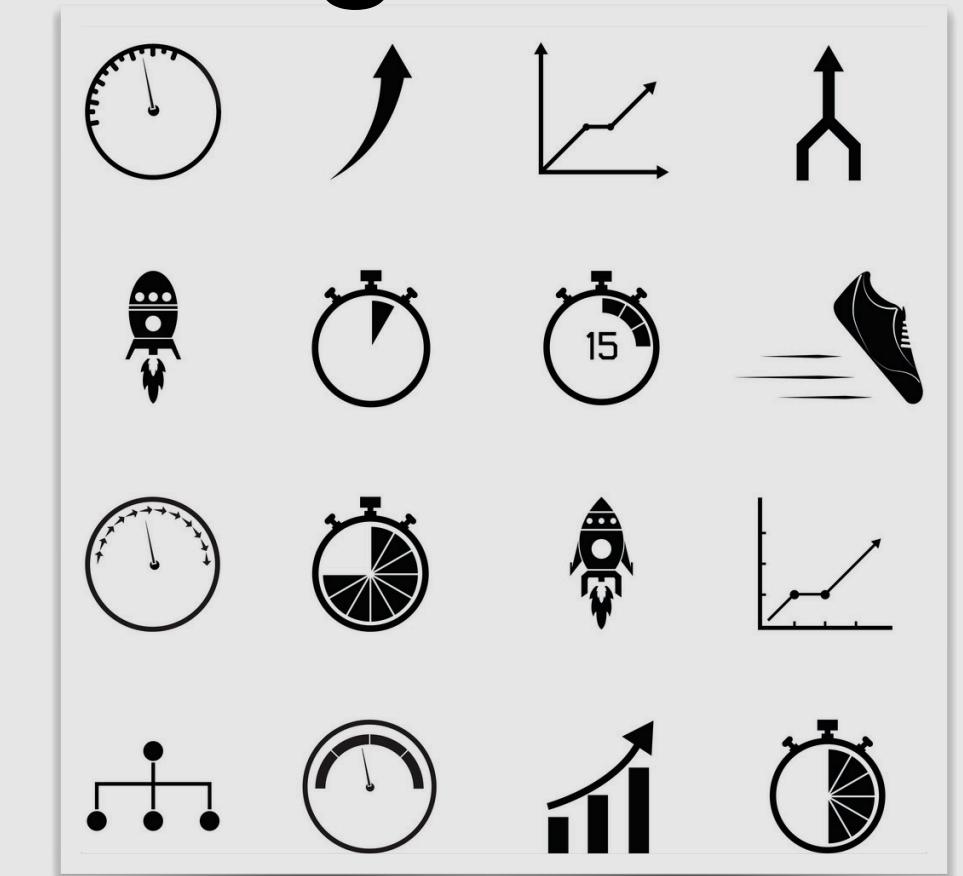
hardware



applications



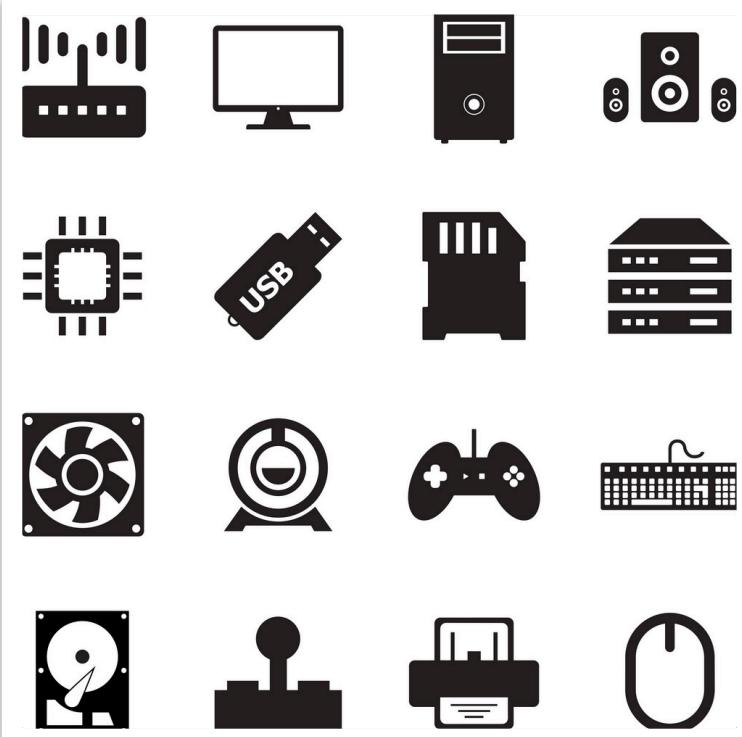
performance  
goals



# The CONTEXT keeps changing ...



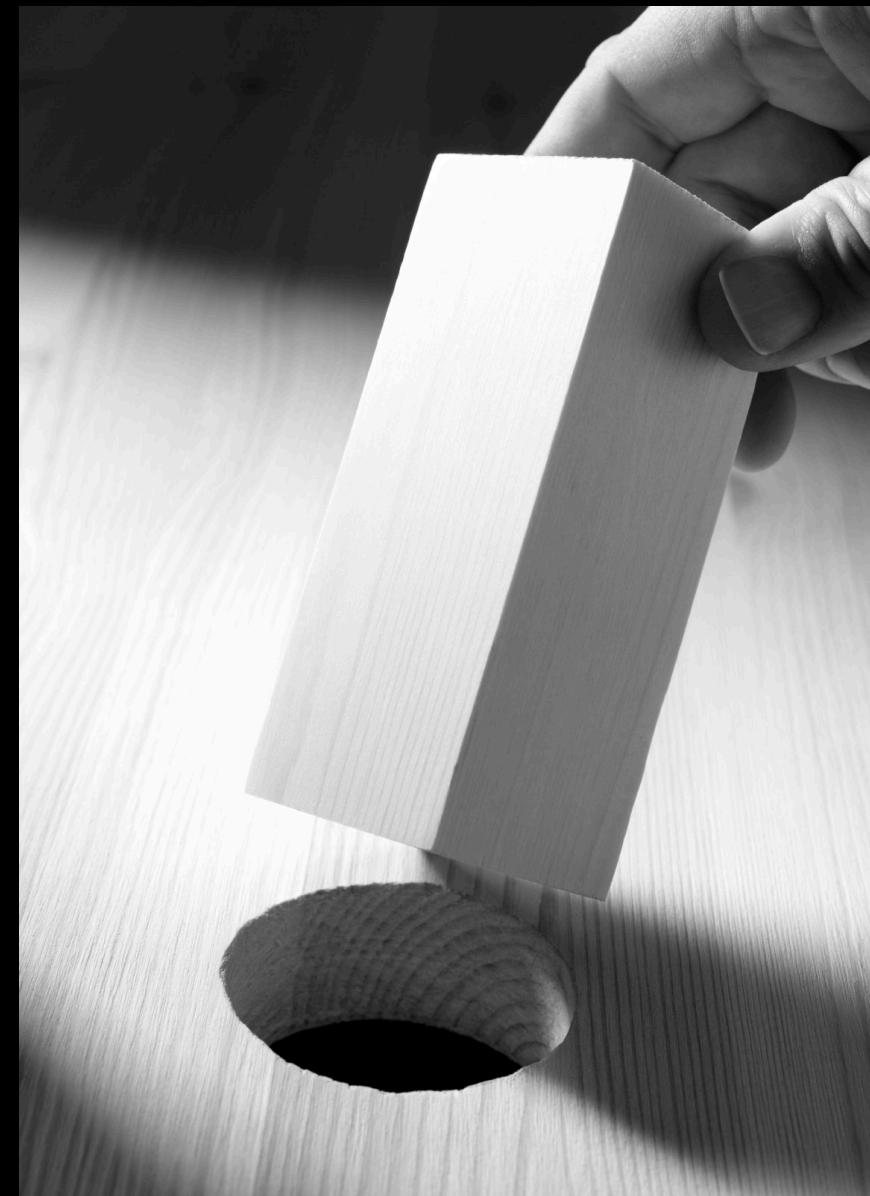
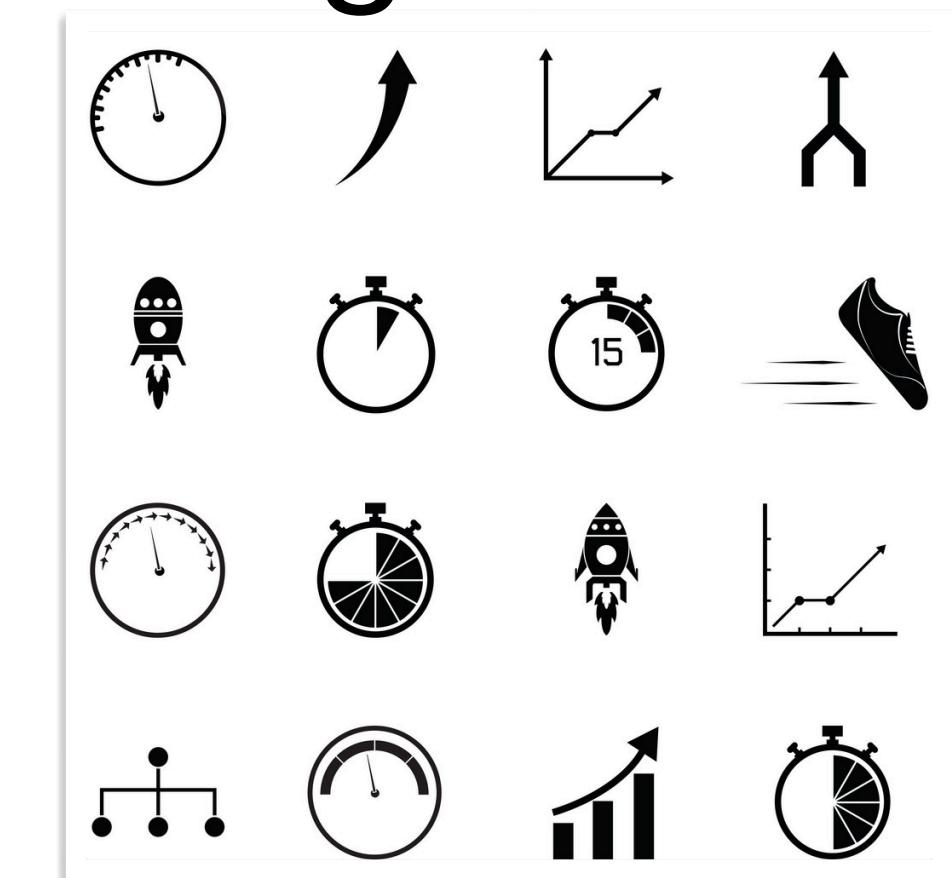
hardware



applications



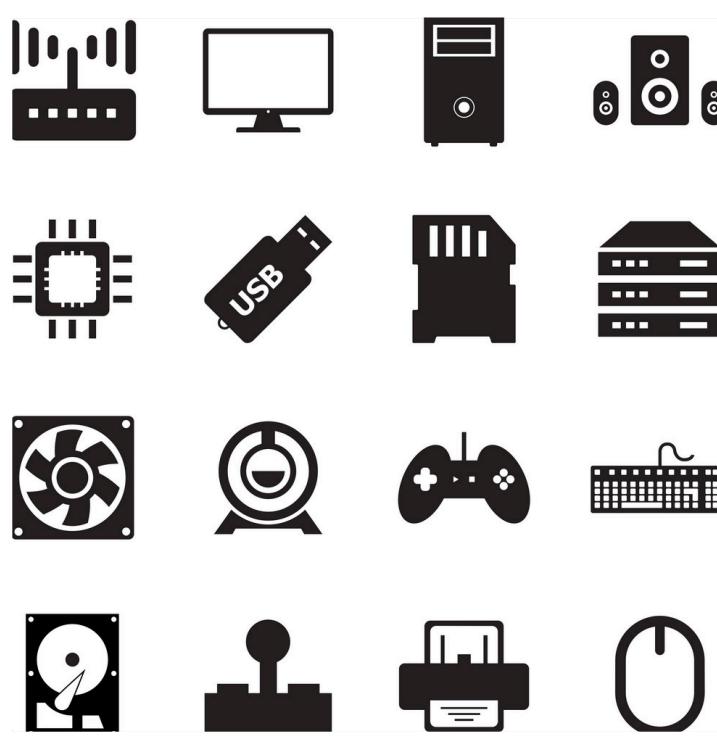
performance  
goals



# The CONTEXT keeps changing ...



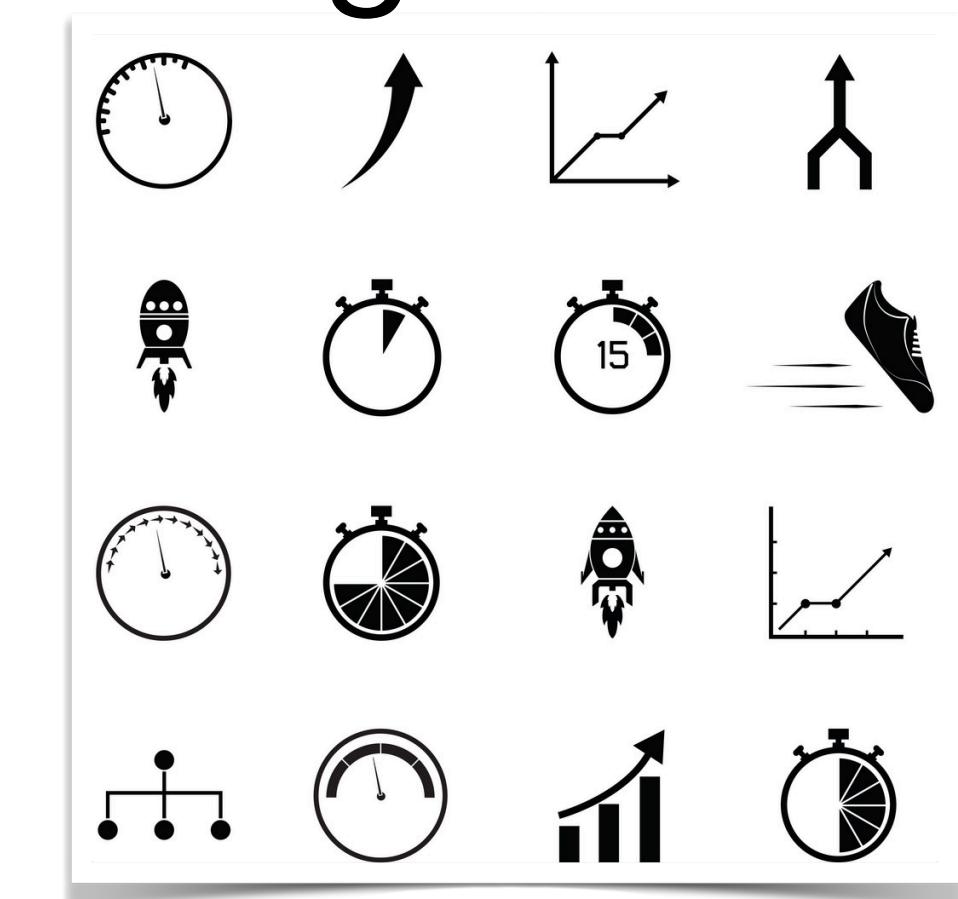
# hardware



# applications



# performance goals

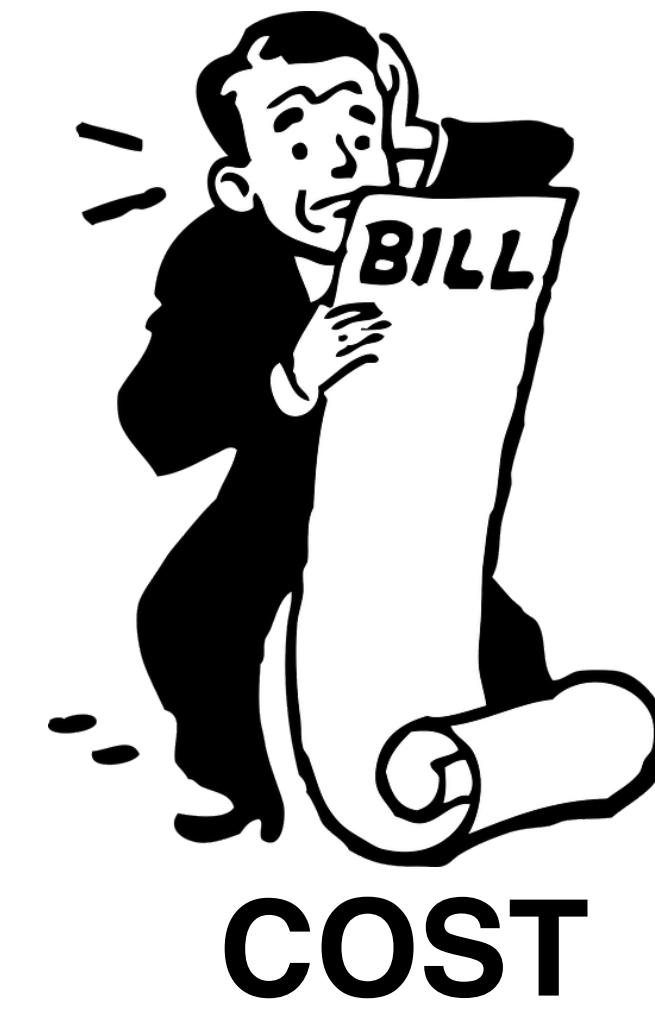


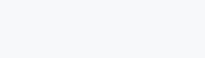
# Bottleneck: Sub-Optimal Data Systems

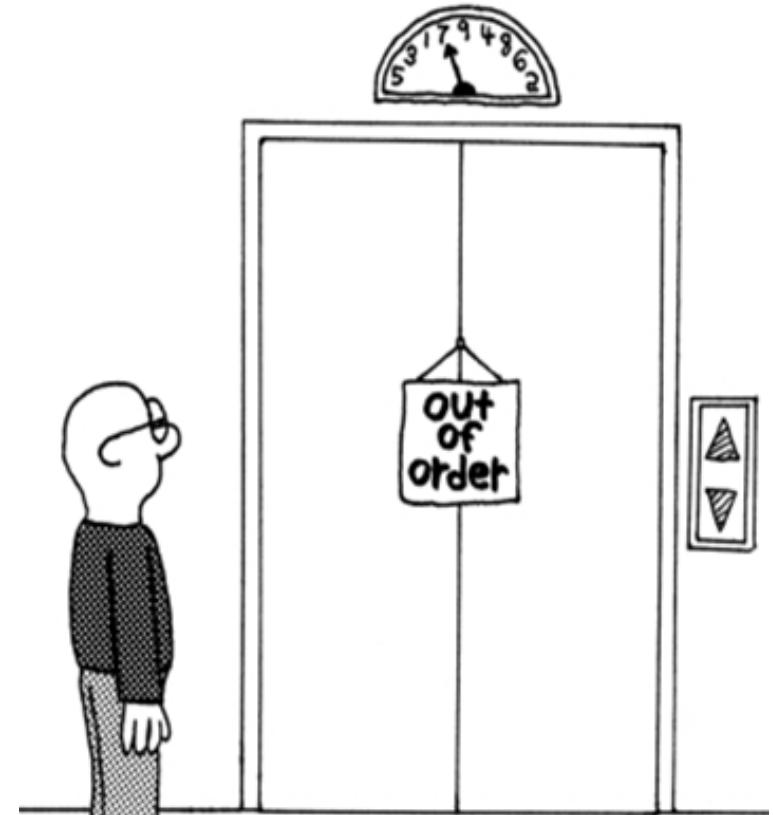
# Bottleneck: Sub-Optimal Data Systems



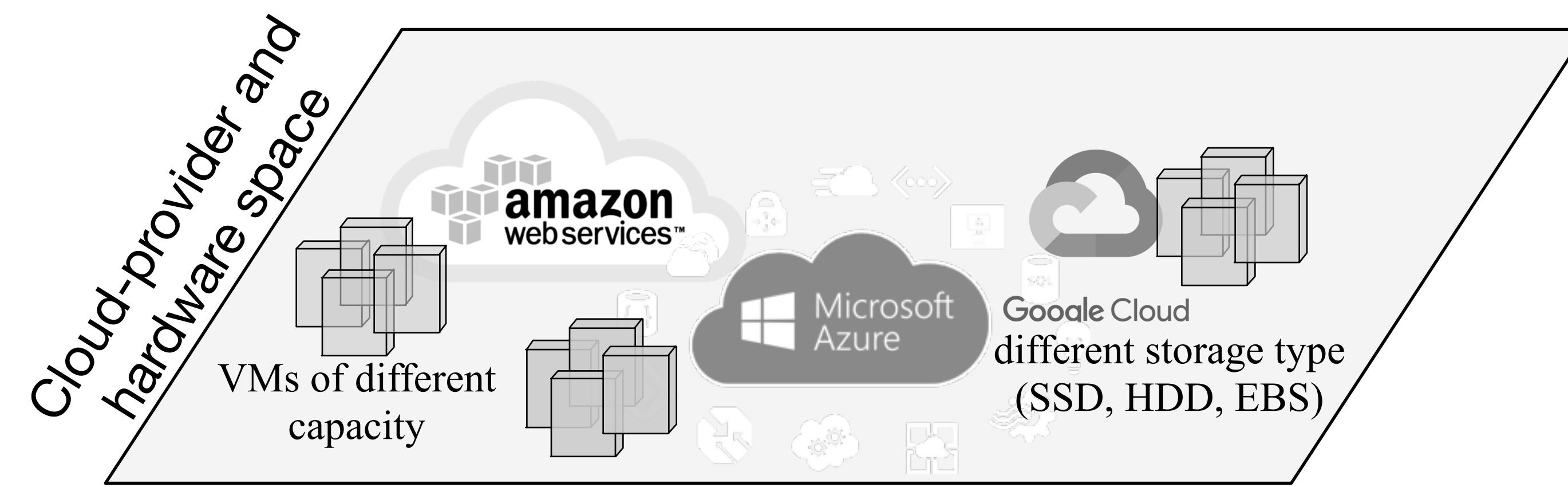
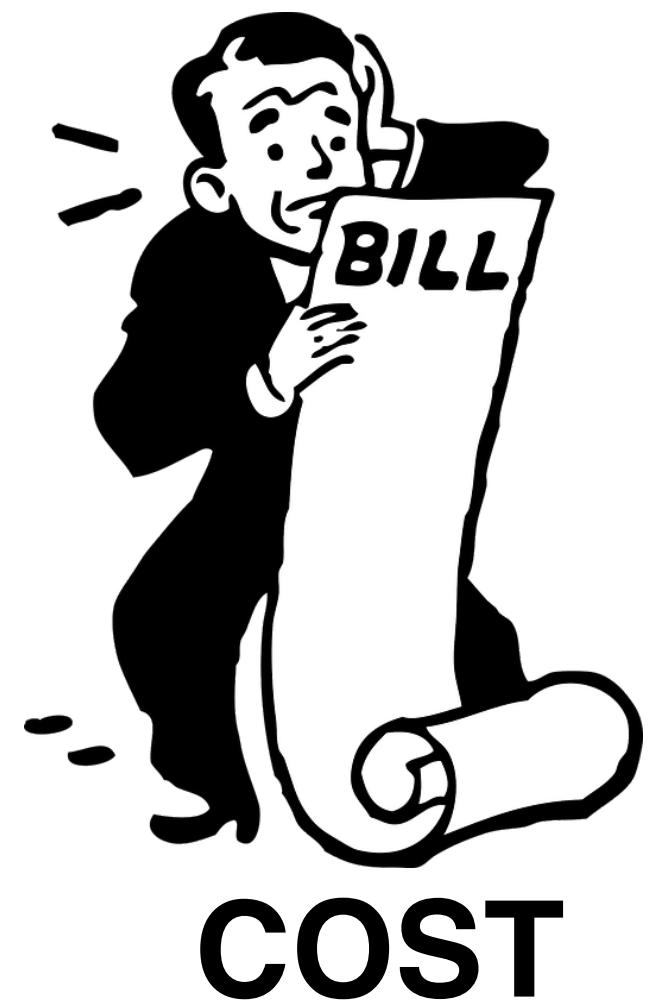
**PERFORMANCE**

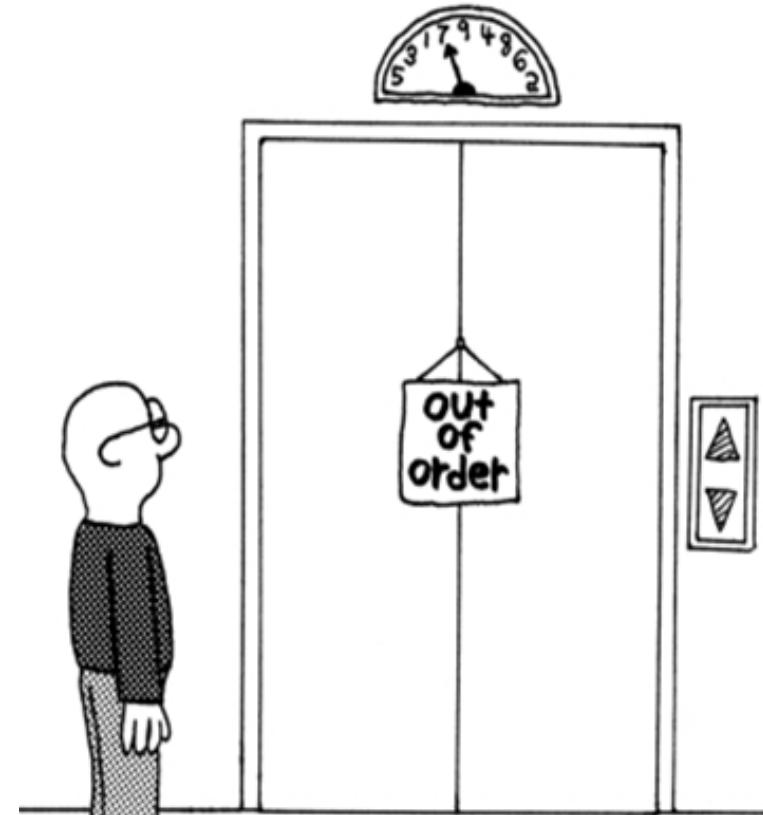


-  Ryan Booth @that1guy\_15 · Mar 4  
Oh fun! Another **high cloud bill**. This is exactly how I wanted to spend the rest of my week...
-  Matt Getty @aspen · Feb 16  
The **Cloud** is as **high** in the sky as the **bill** from AWS\*  
\* for a tremendous amount of workloads
-  CiscoEvents @CiscoEvents · Jan 28  
Is your **cloud bill** too **high**? Learn how you can control **cloud costs** with CloudCenter Suite.
-  Translucent Computing @translucentcomp · Jan 20  
**#Kubernetes** **cloud** costs are getting out of hand. Working with many clients and different **cloud** service providers, in more than 70% of cases, we see the VM cluster nodes are underutilized, which leads to a **high cloud bill**.

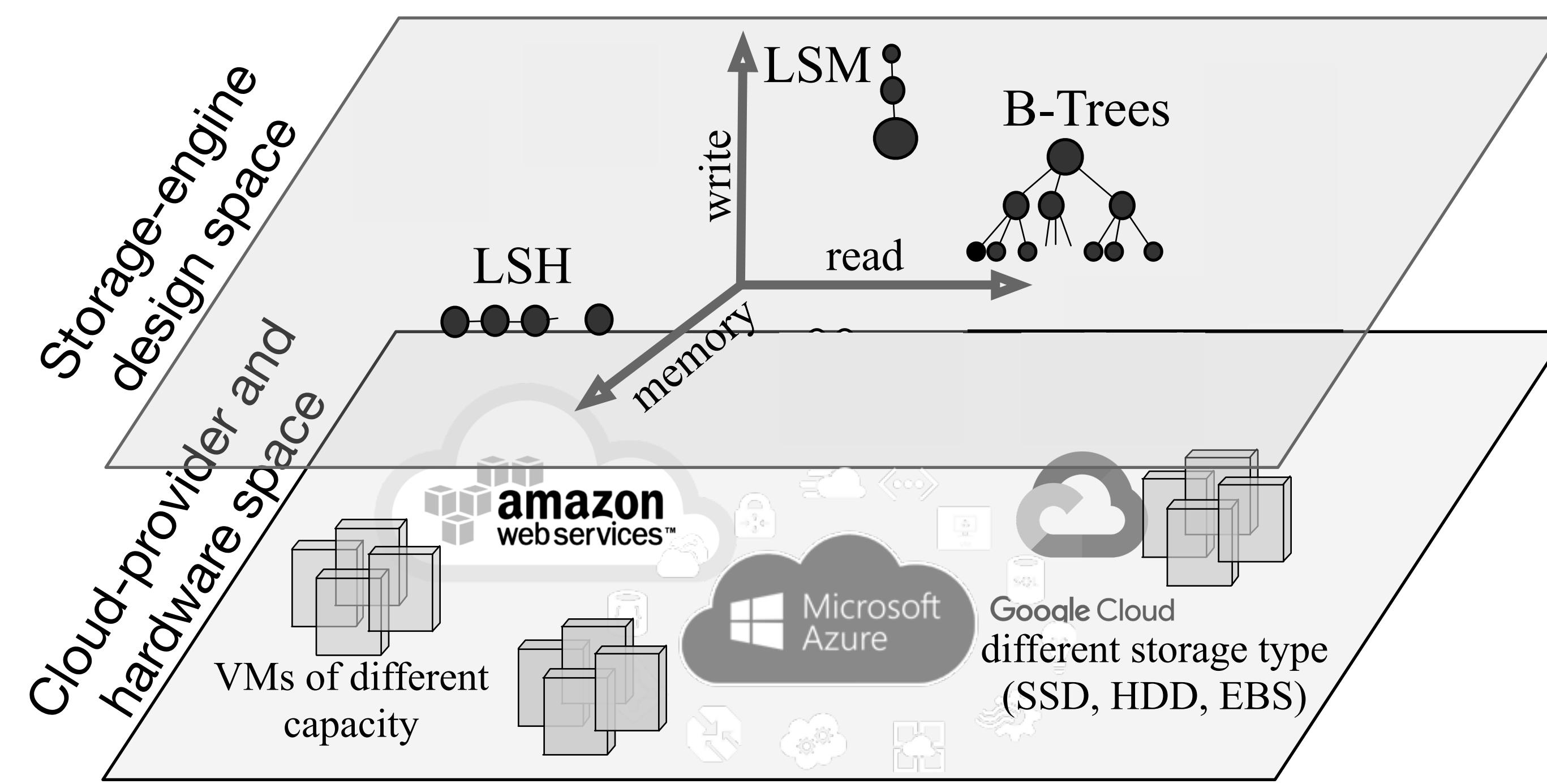
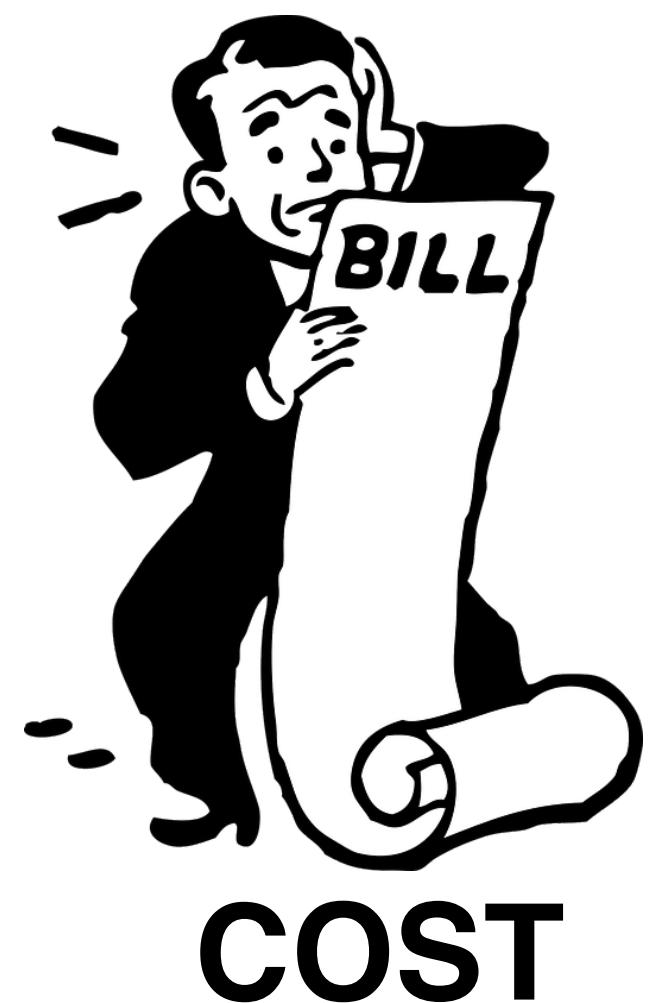


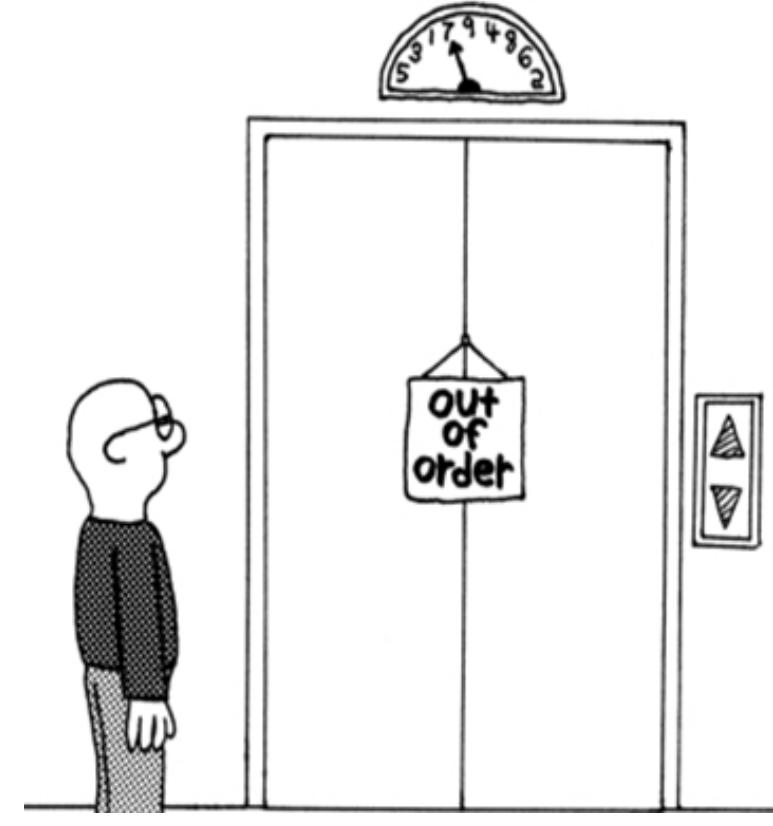
## PERFORMANCE



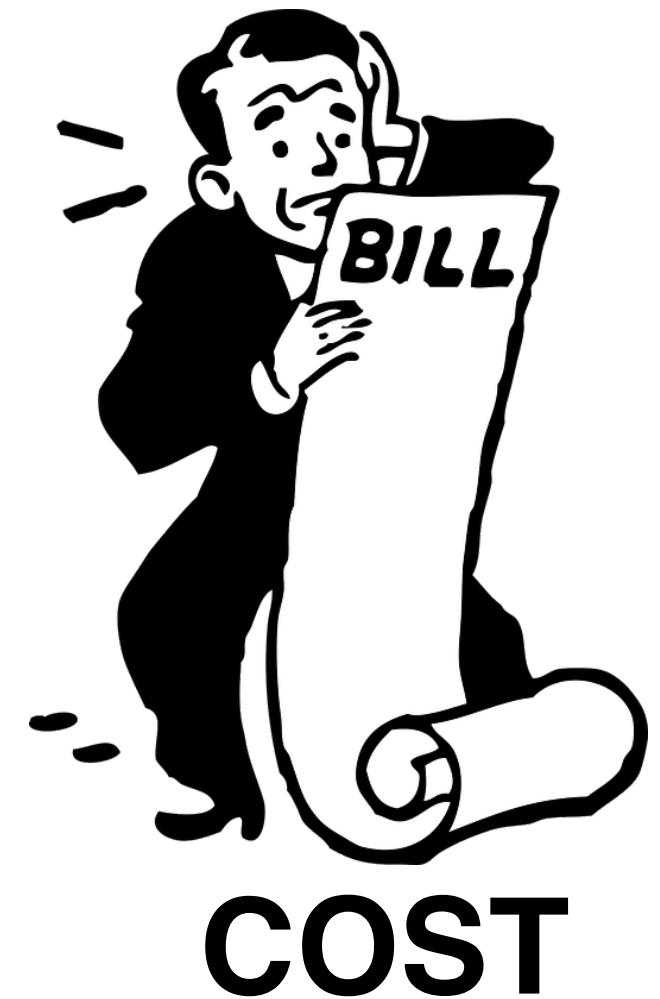


## PERFORMANCE

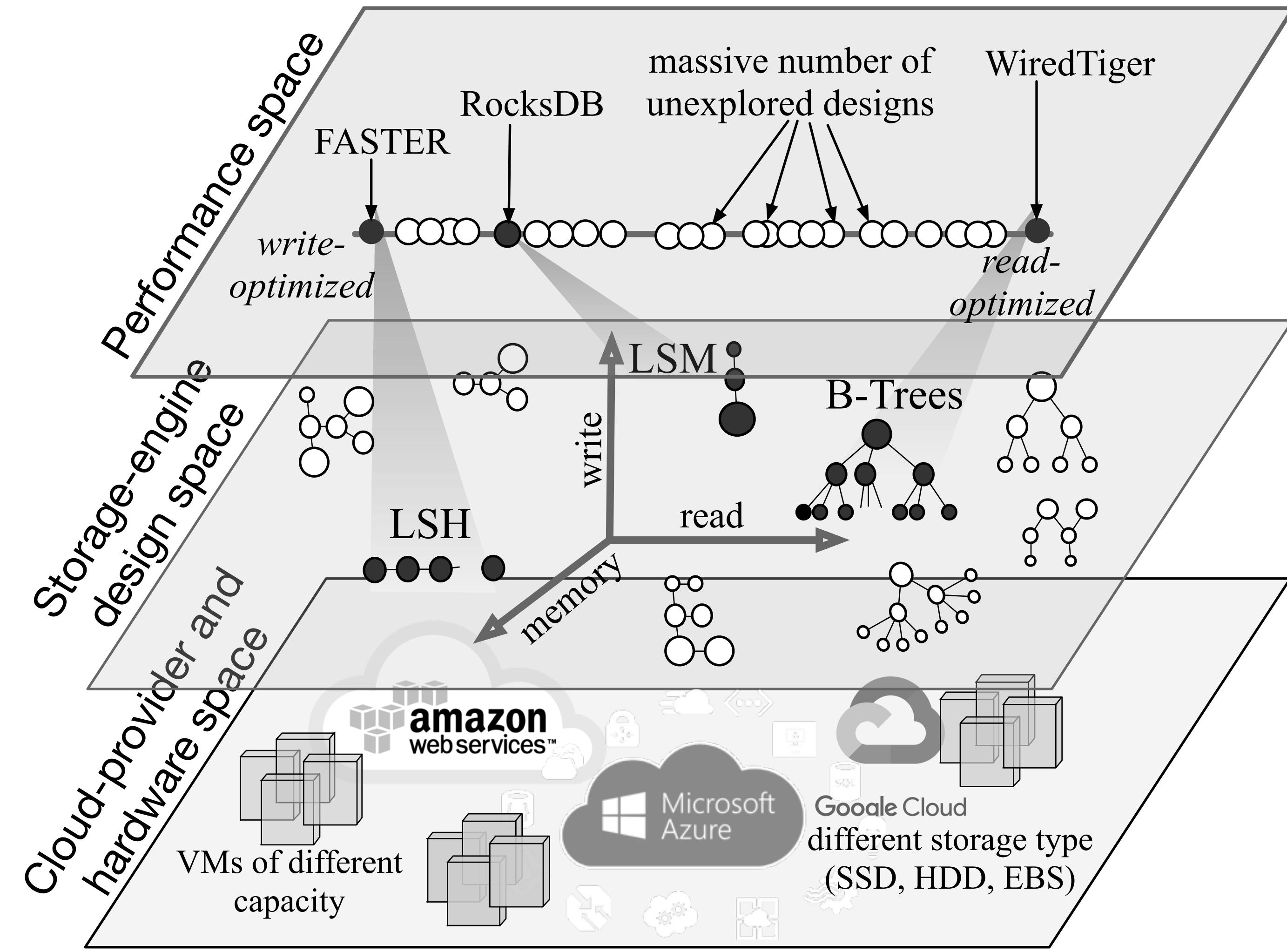




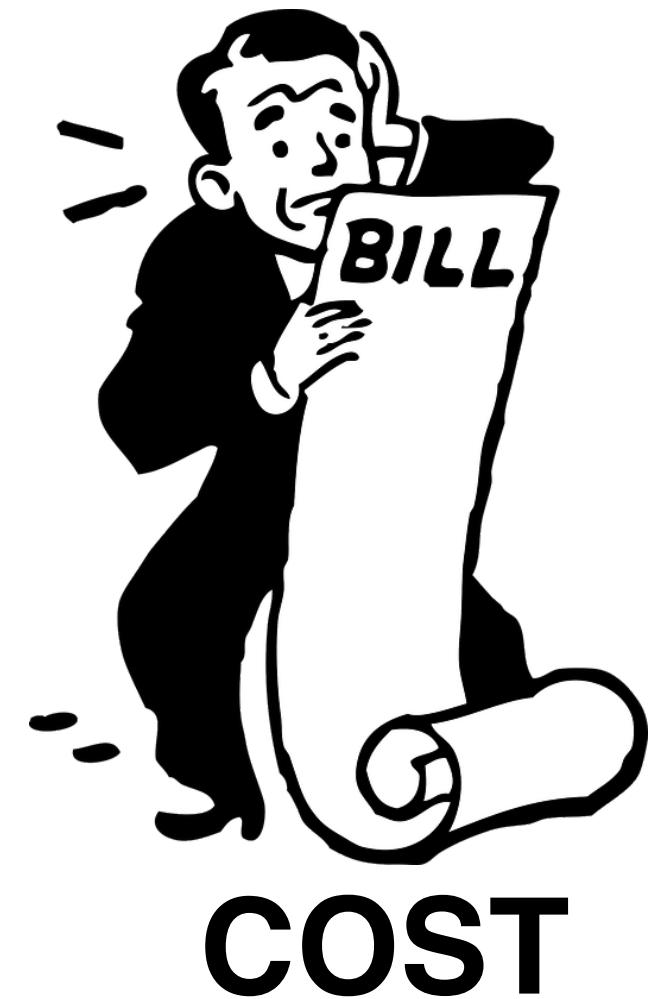
## PERFORMANCE



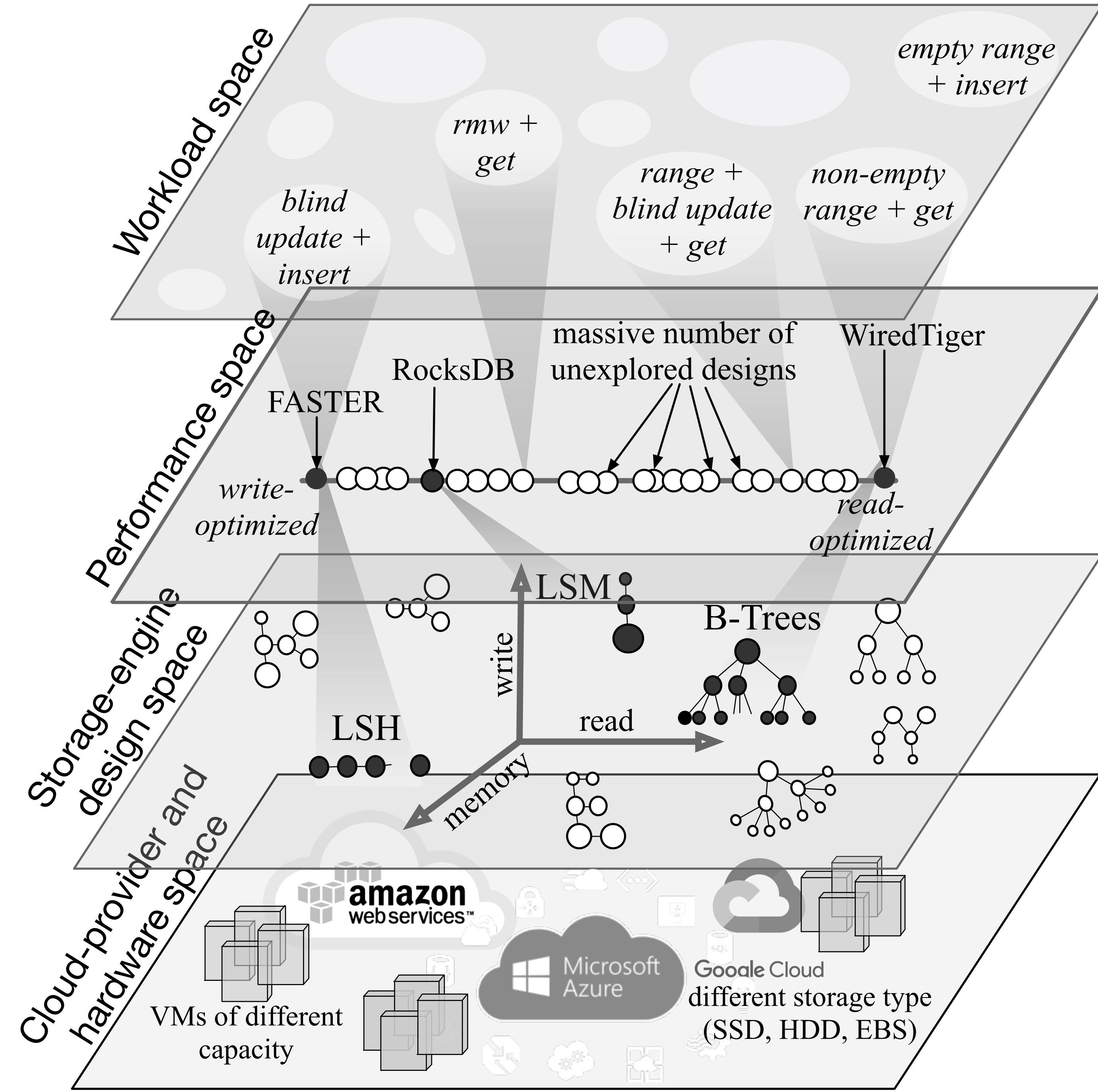
## COST



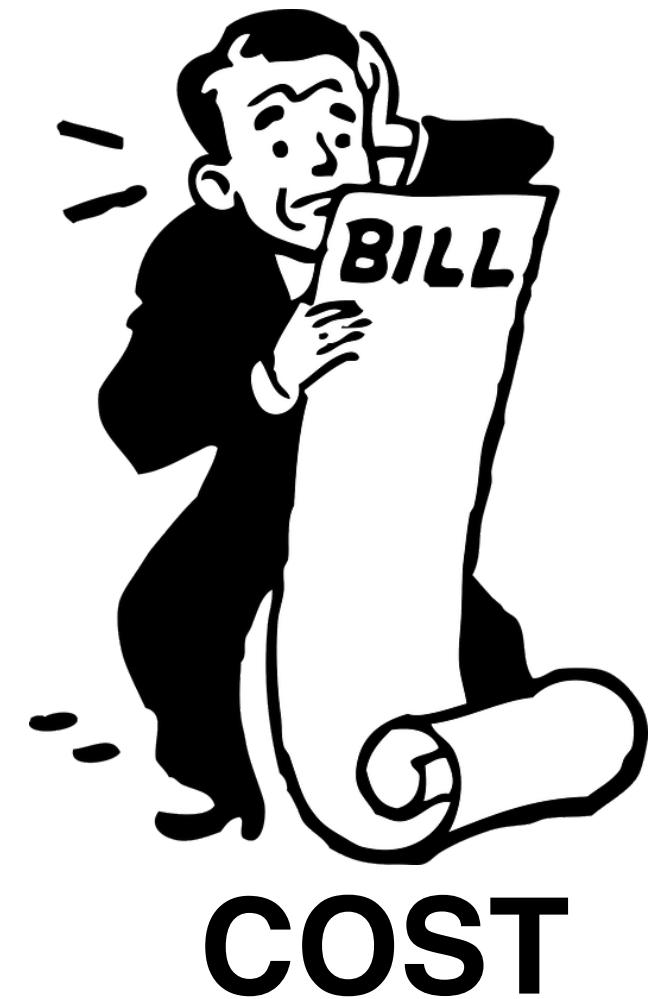
# PERFORMANCE



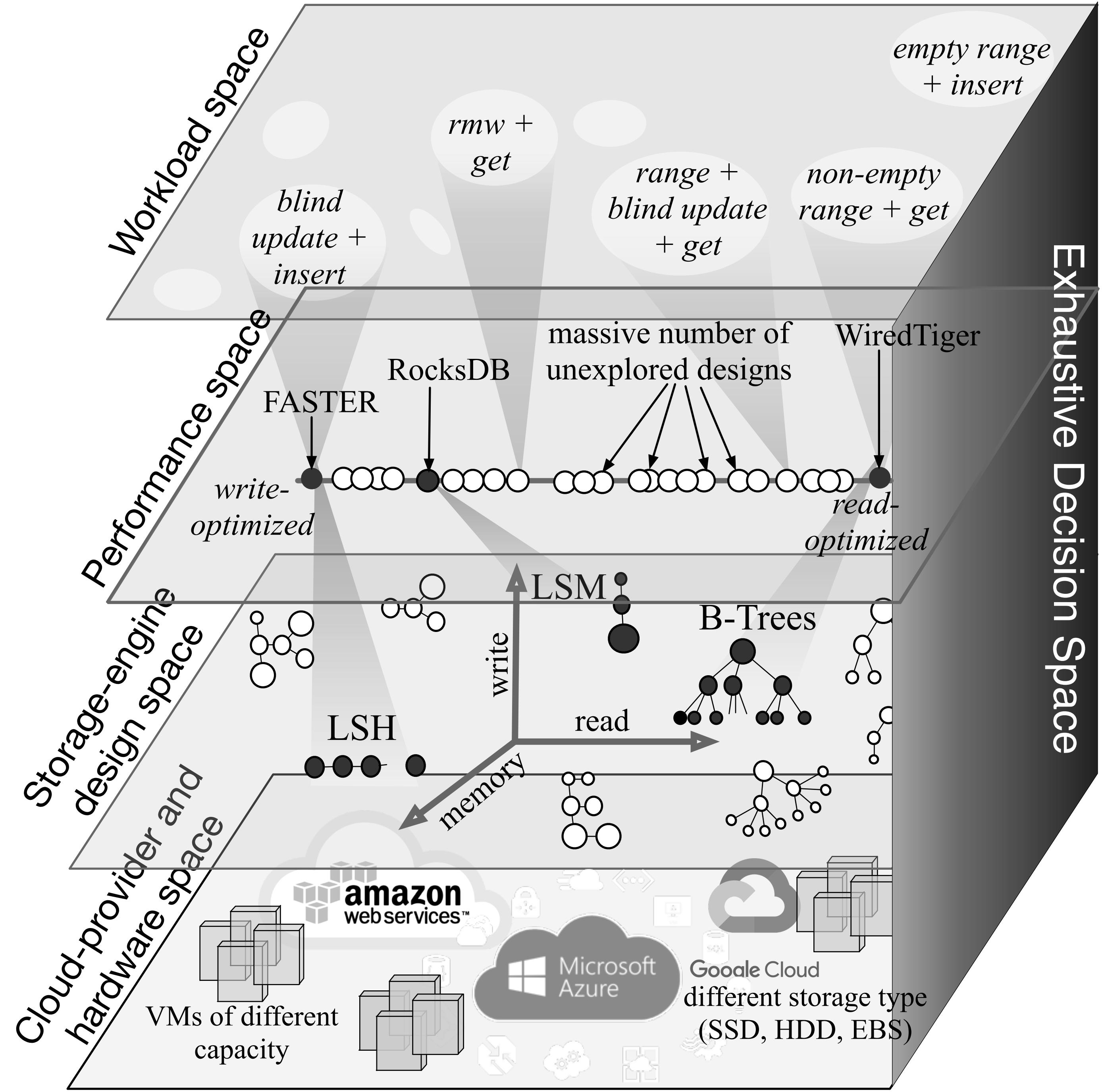
COST

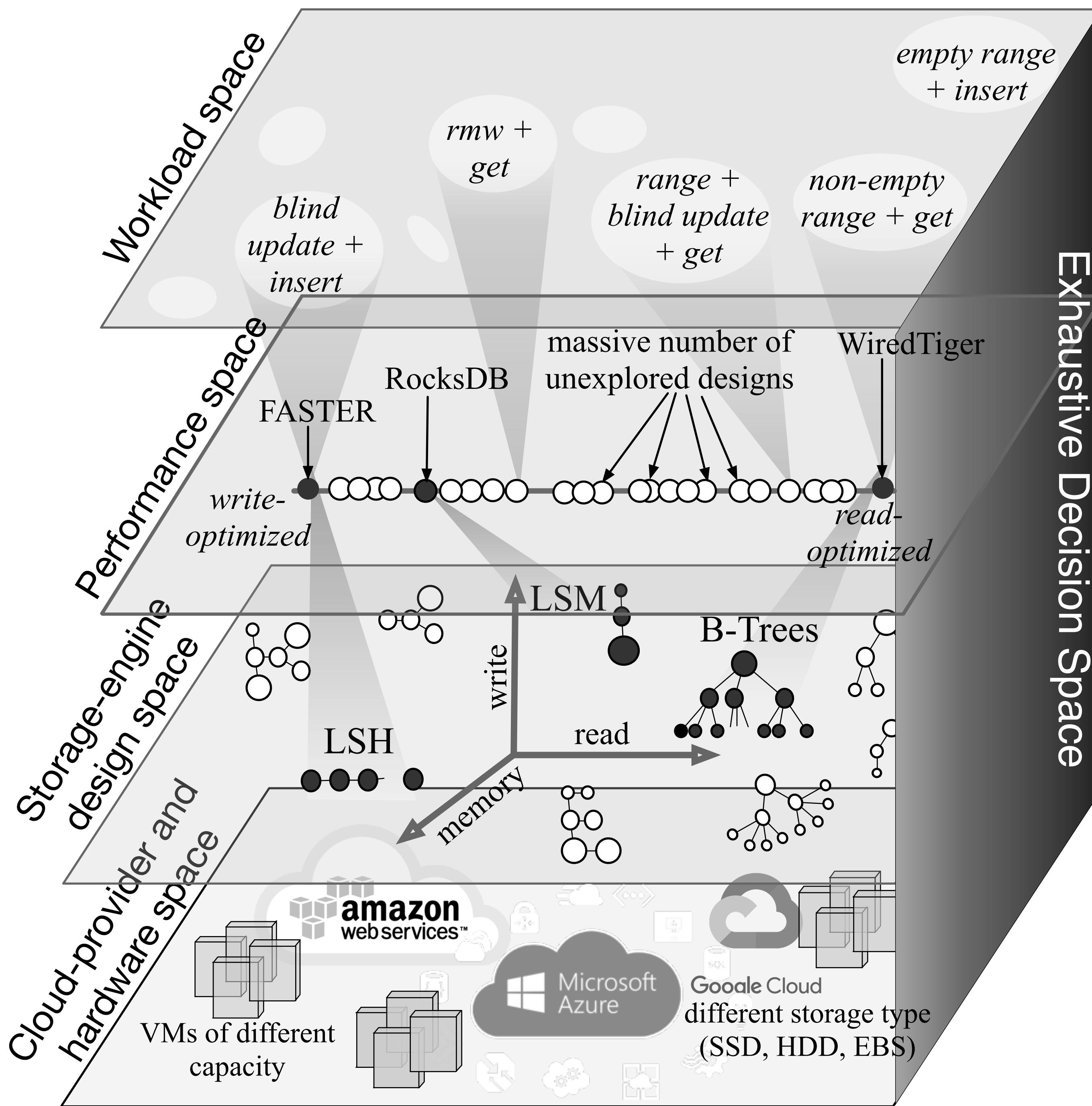


# PERFORMANCE



COST

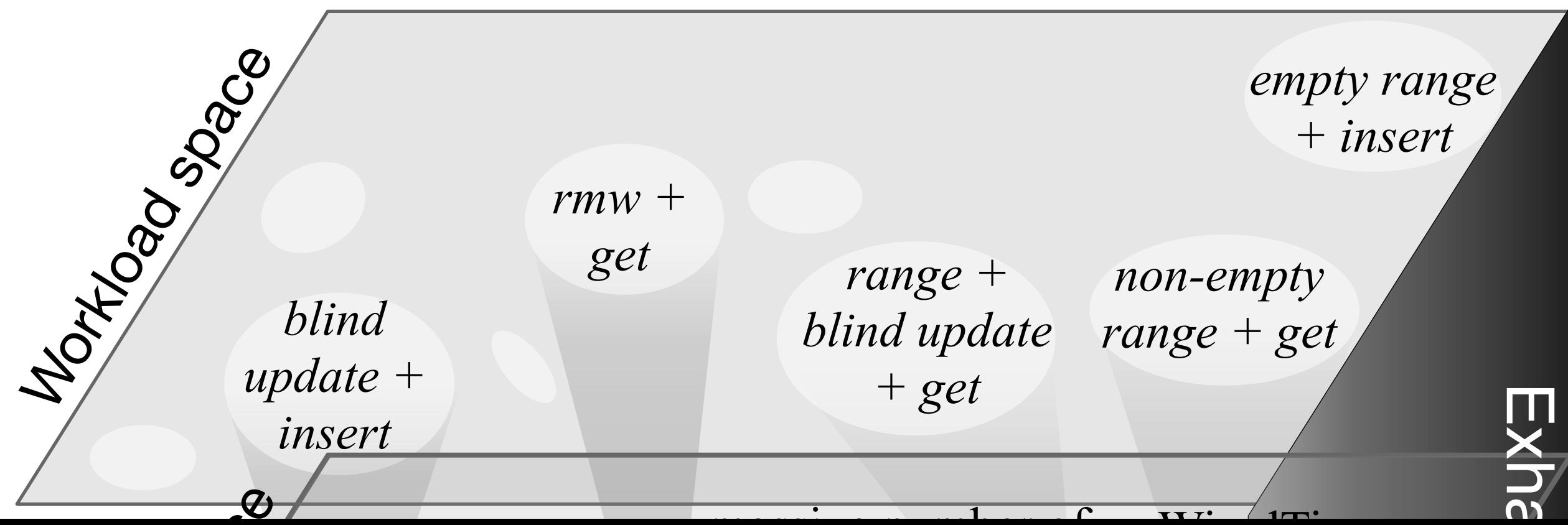




**Complex**

**Vast**  
**( $10^{35}$  possibilities)**

**Manual**  
**(Limited exploration of systems)**



**GOAL:** To create the perfect data system  
tailored for each context



Gilad David Maayan

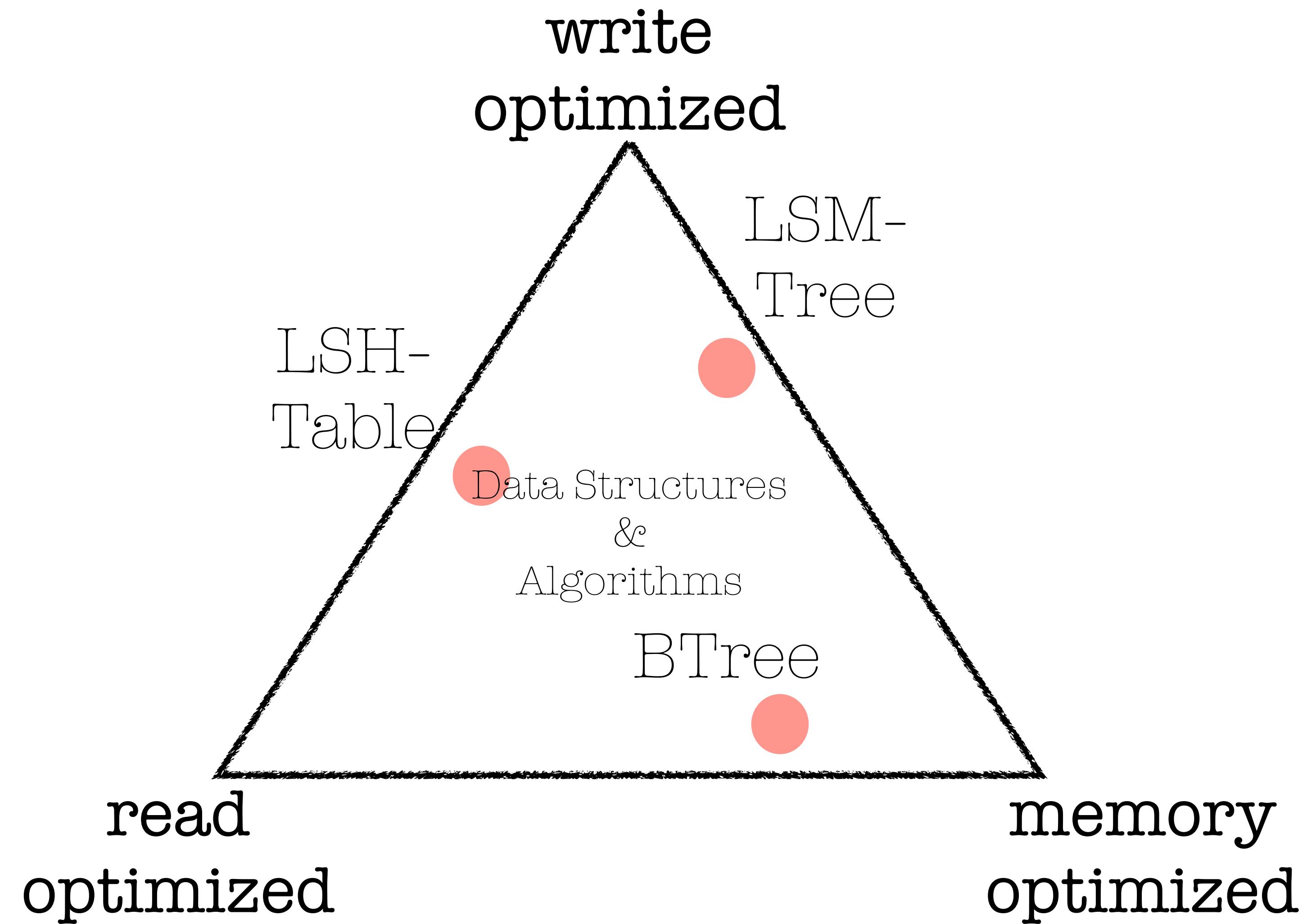
Posted on Sep 25, 2020 • Updated on Dec 10, 2020

## AWS to Azure: Making the Move

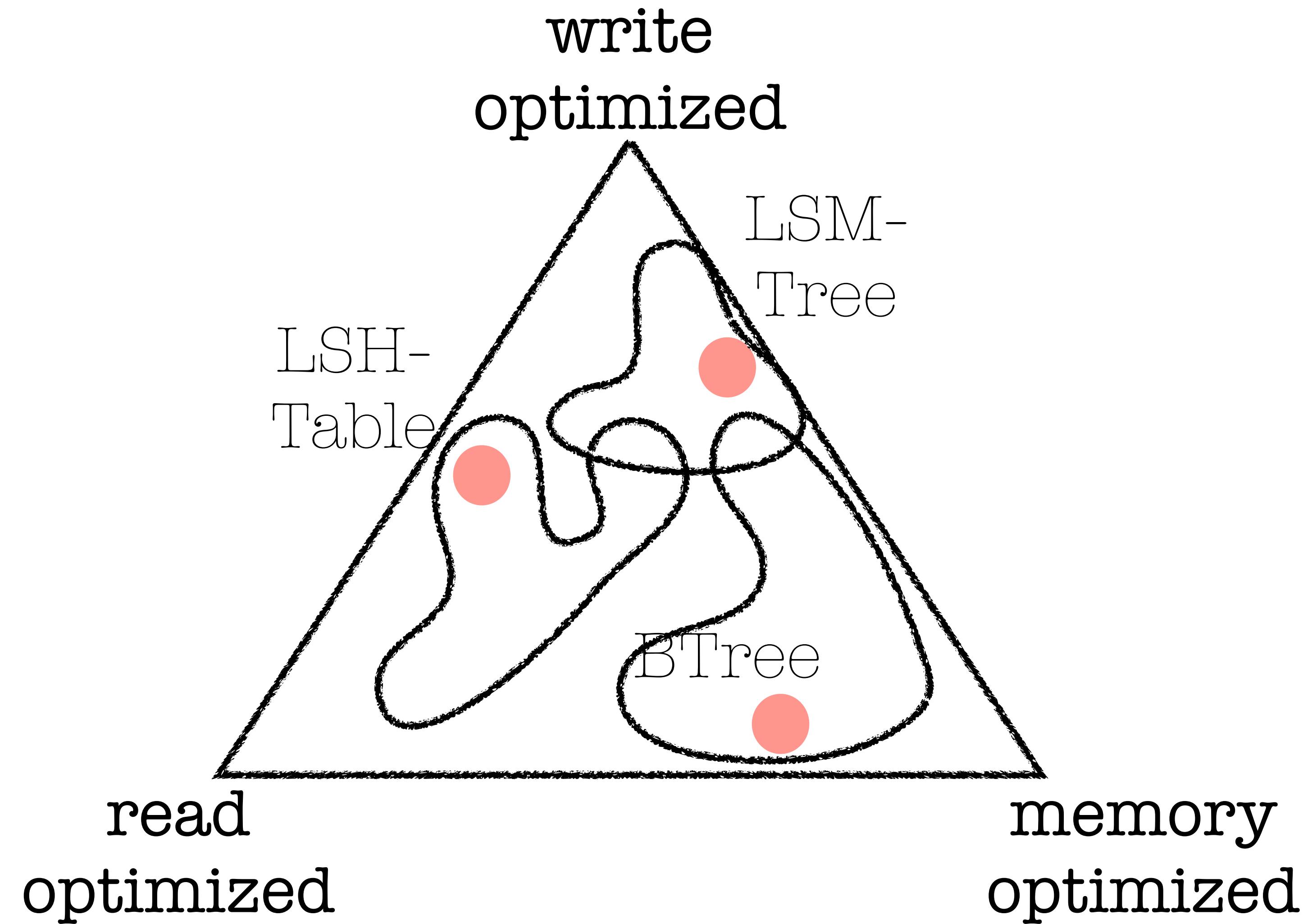
#azure #aws

Both Amazon Web Services and Microsoft Azure are considered top [cloud computing companies](#). However, there are certain aspects unique to each.

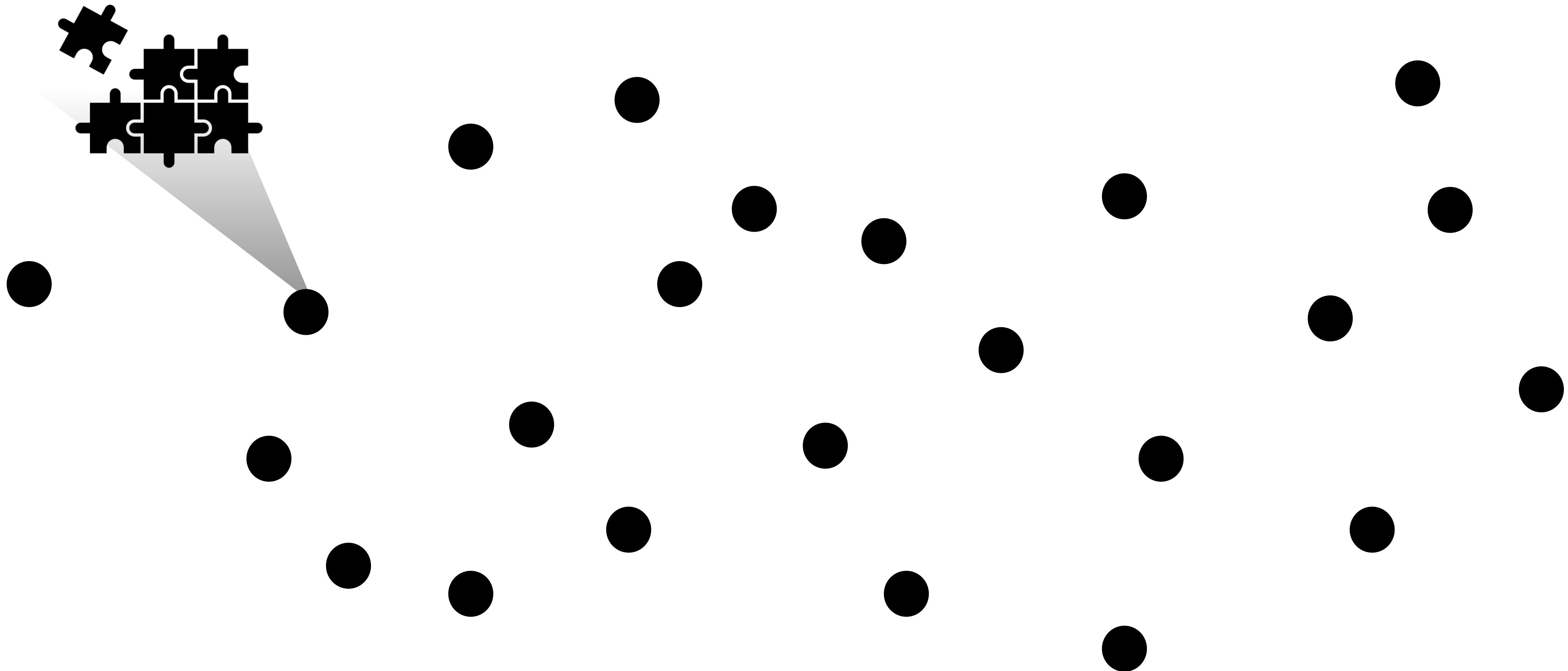
# Intuition



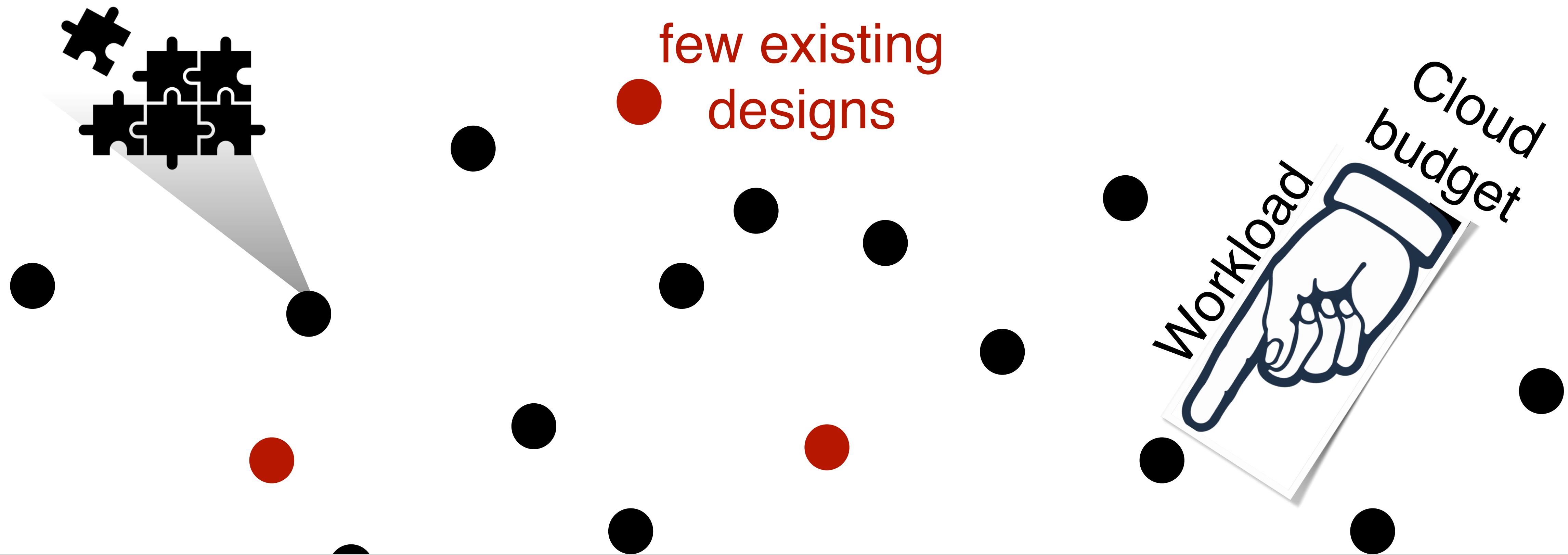
# Intuition



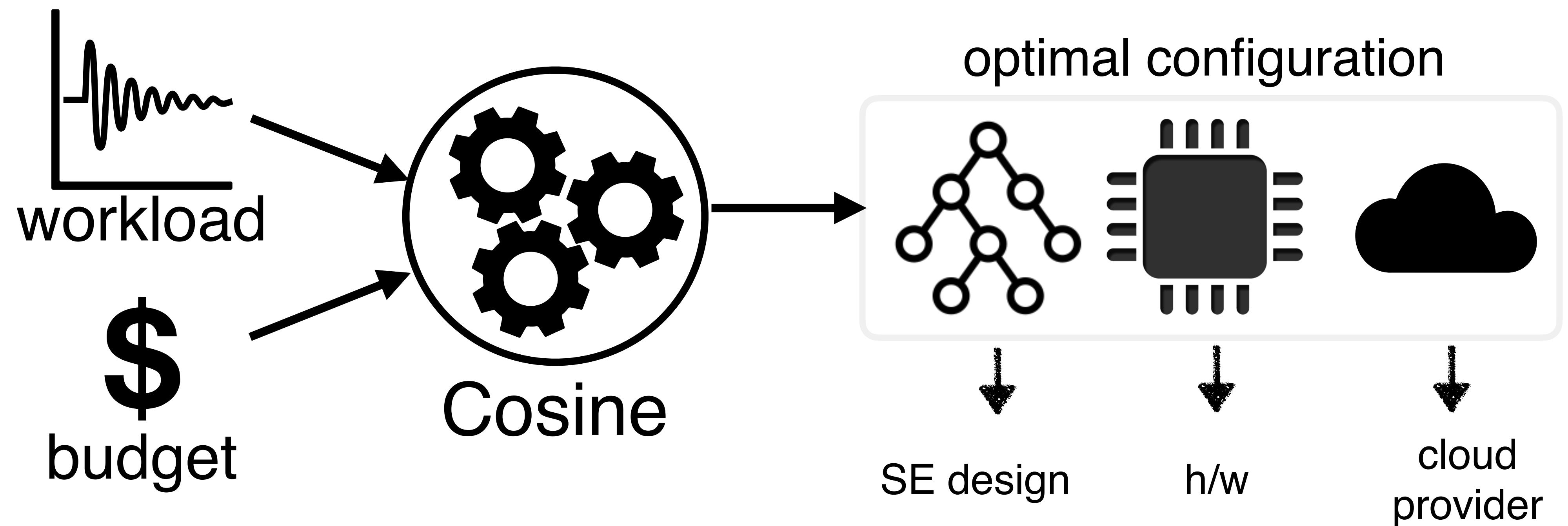
# THE KEY INTUITION



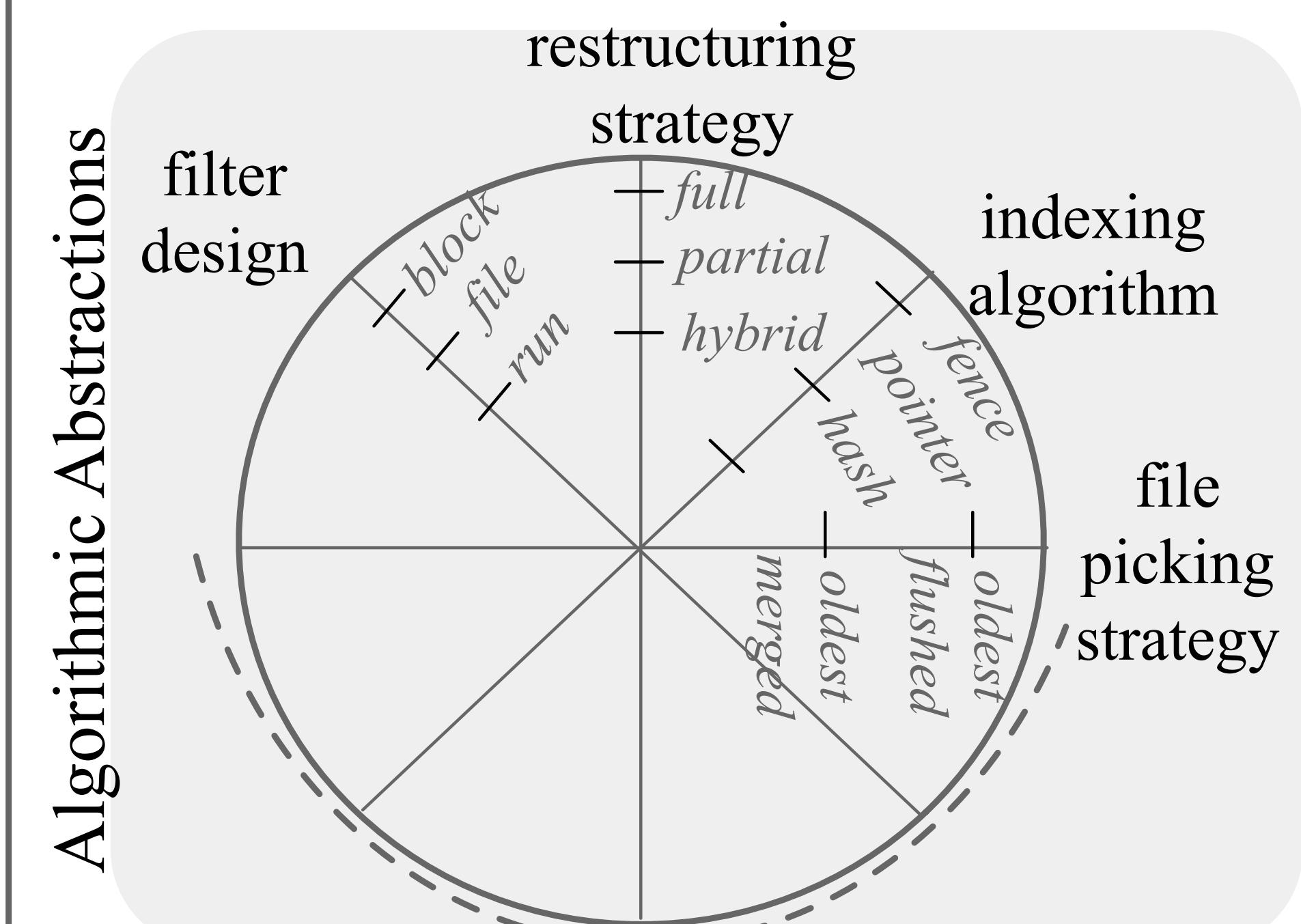
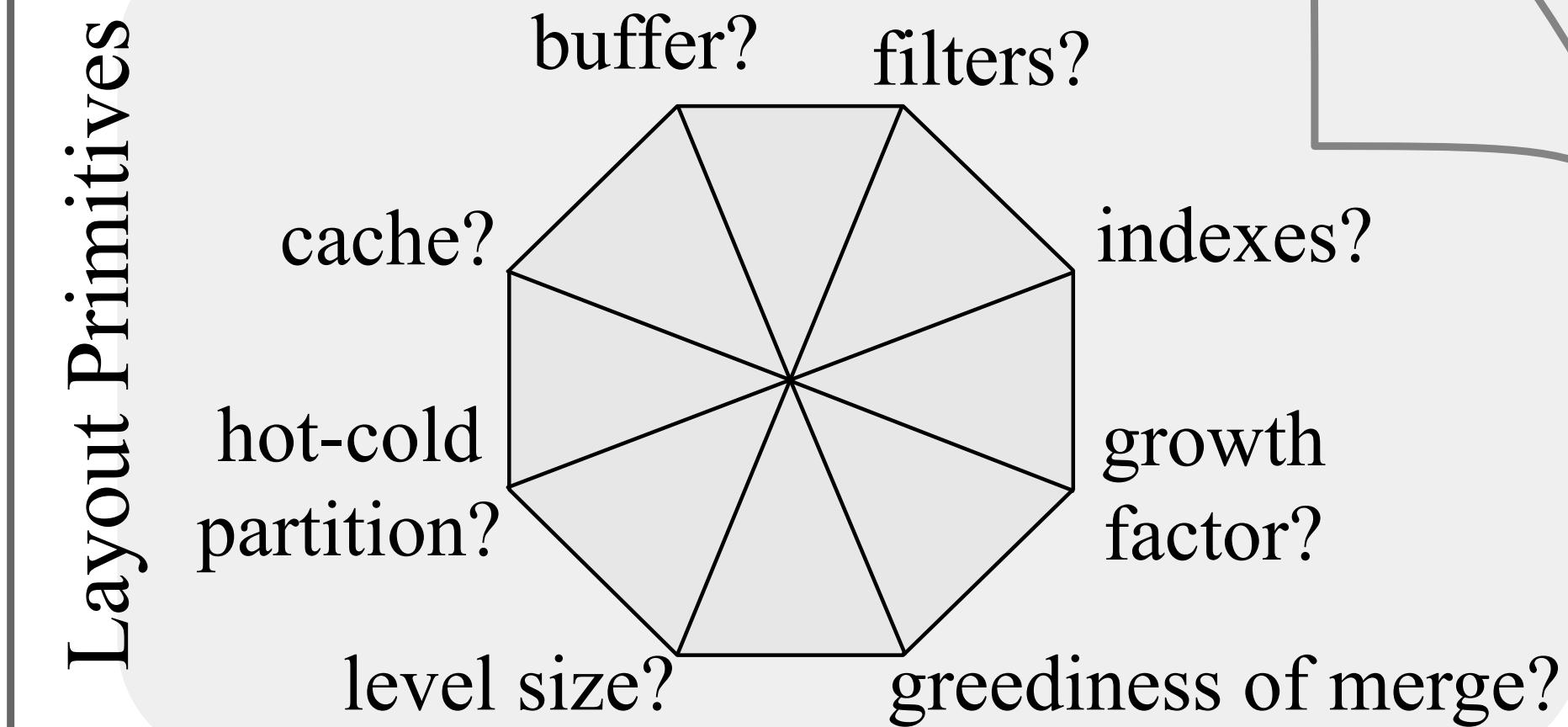
# THE KEY INTUITION

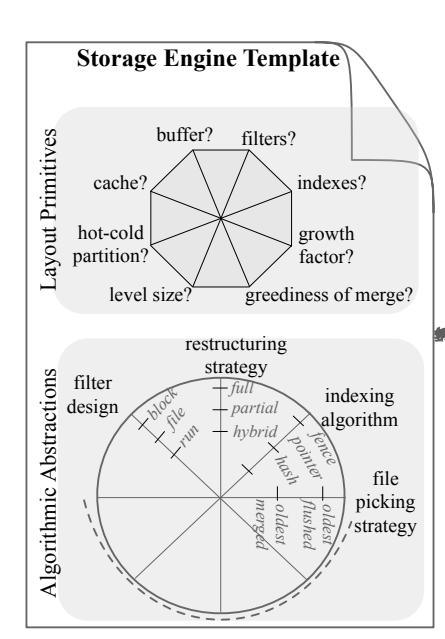


WHAT IF WE COULD REASON ABOUT THE MASSIVE DESIGN SPACE OF STORAGE ENGINES?



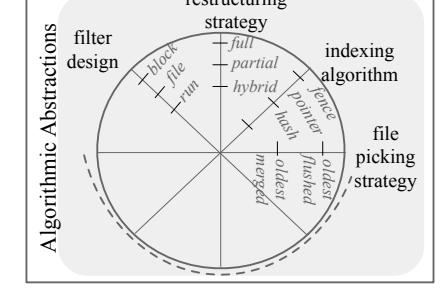
# Storage Engine Template





# MEMORY

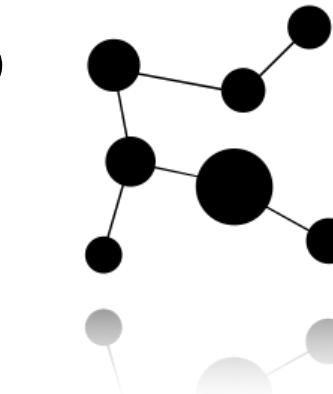
# DISK



# MEMORY

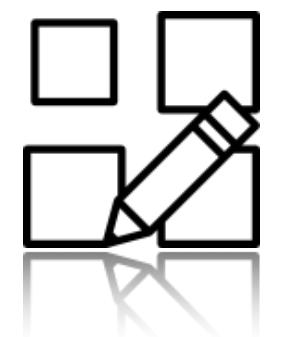
# DISK

storage pattern?  
*{flat logs, hierarchical}*



level size?  $\leftarrow \square \rightarrow$   
 $[1, \dots, L]$

greediness of merge?  
*[1 (high), ..., T (low)]*



file size?  
MB ... GB

hot-cold partition?

# MEMORY

buffer?



[1..M]

Bloom? Cuckoo?



...

[1..M]

filters?

indexes?



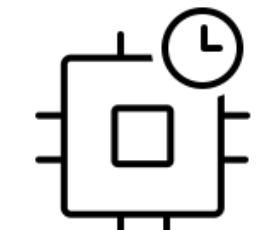
hash table?



zone map?

...

cache?

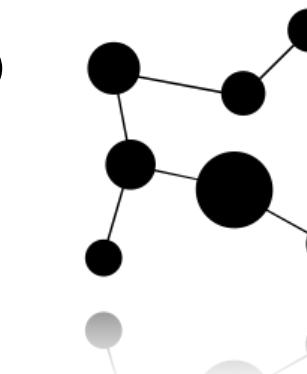


[1..M]



# DISK

storage pattern?  
{flat logs, hierarchical}

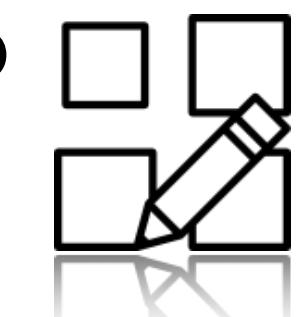


level size? ↵

[1, ..., L]

file size?  
MB ... GB

hot-cold partition?



greediness of merge?  
[1 (high), ..., T (low)]

# Superstructure

MEMORY

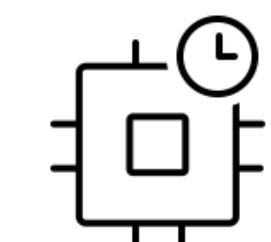
buffer?



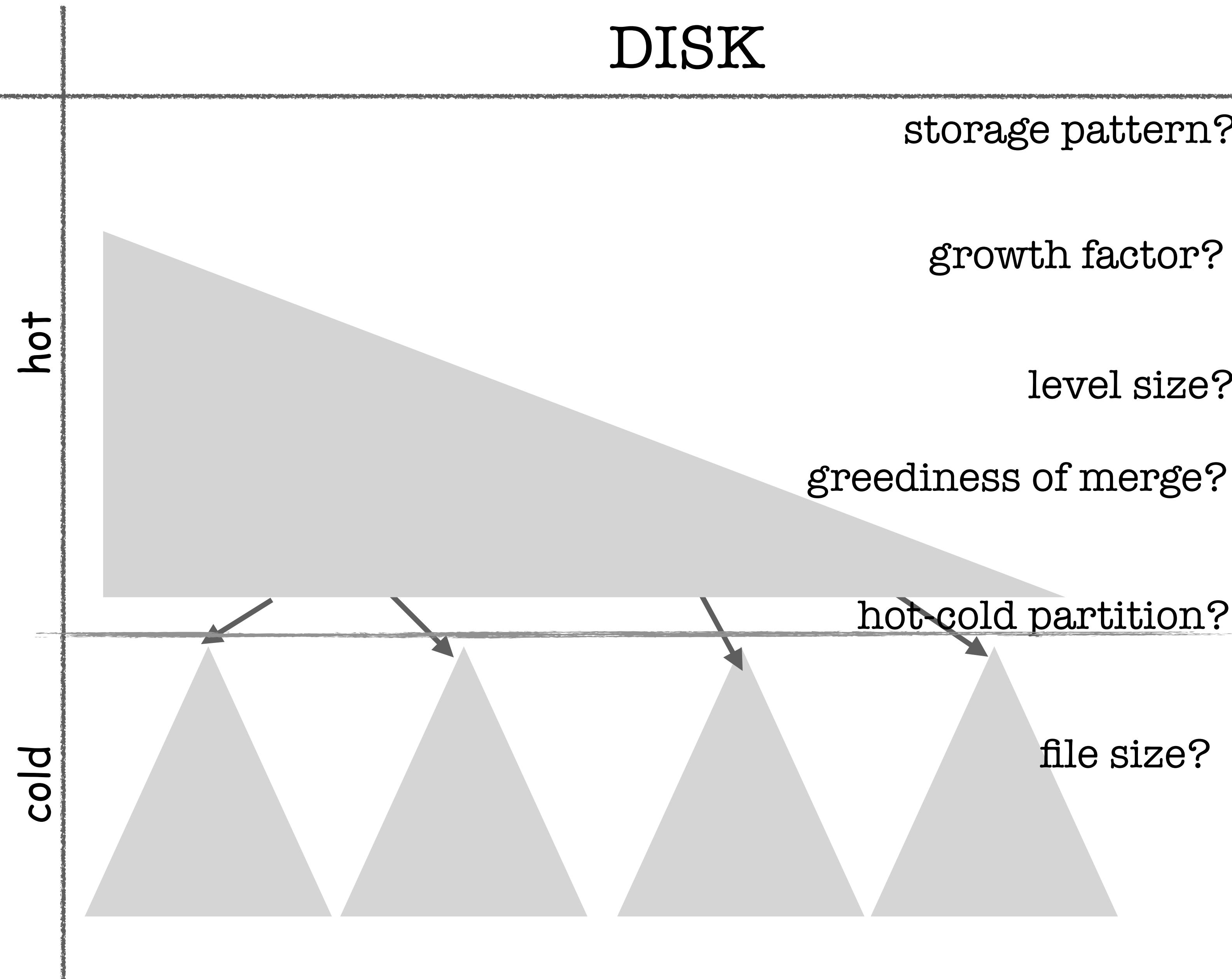
filters?



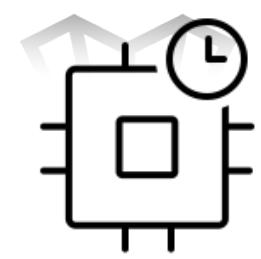
indexes?



cache?

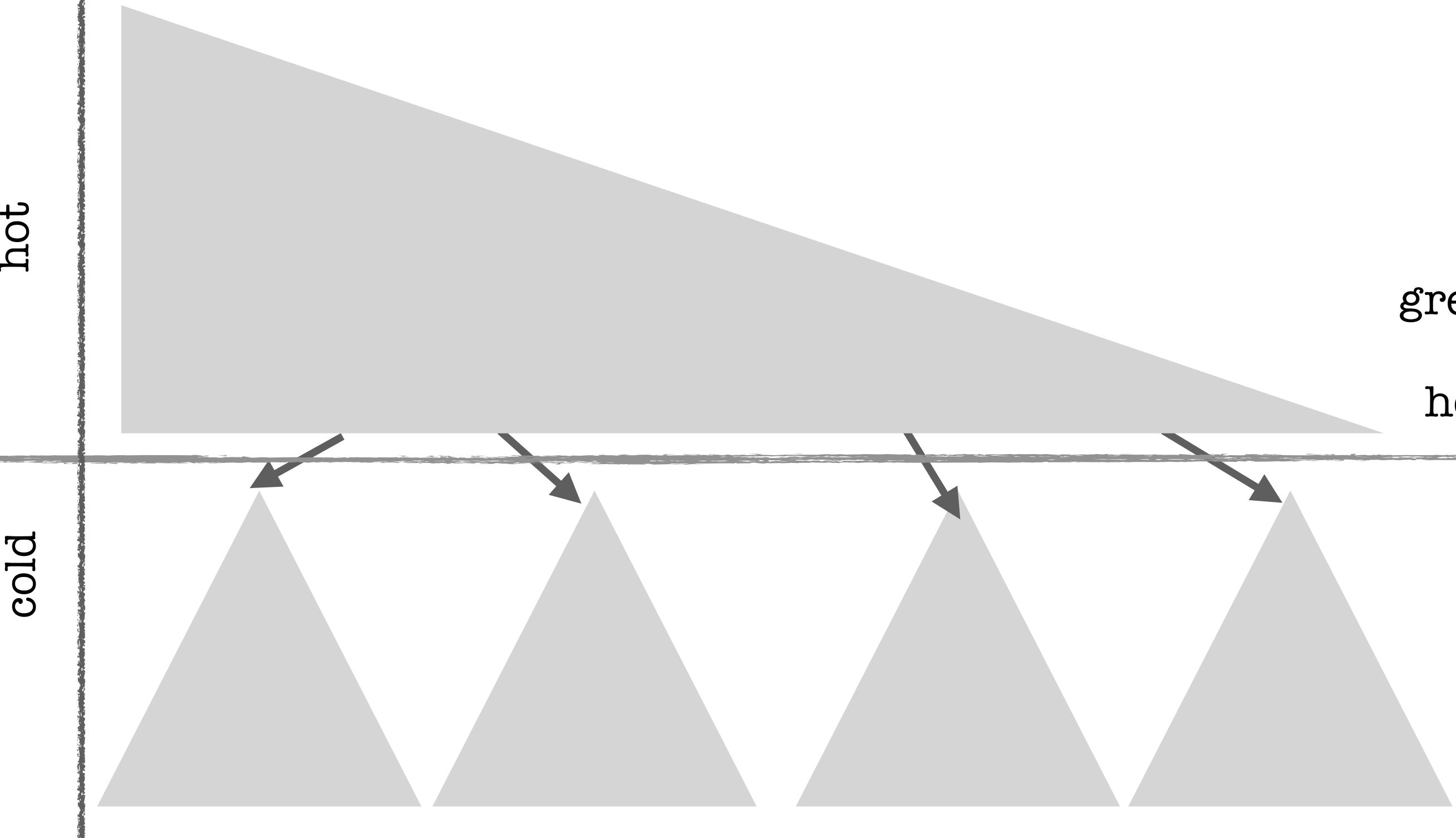


## MEMORY

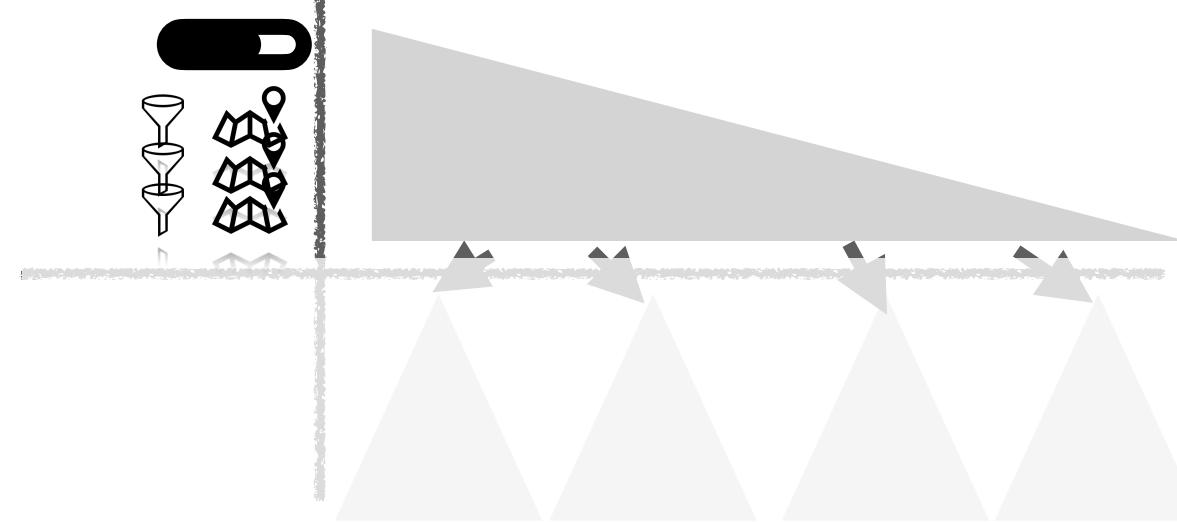
- buffer? 
- filters? 
- indexes? 
- cache? 

## DISK

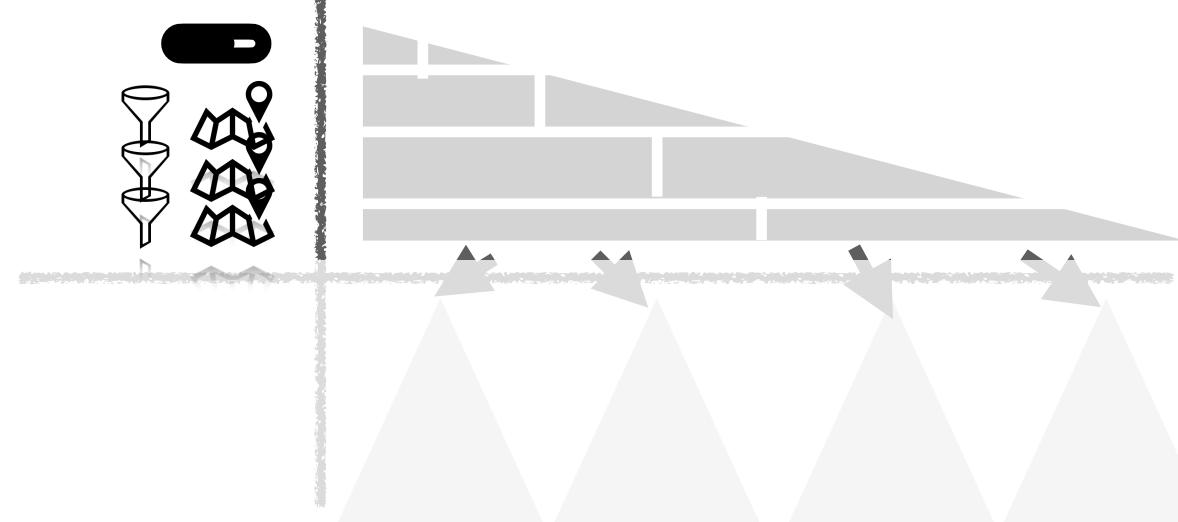
- storage pattern?
- growth factor?
- level size?
- greediness of merge?
- hot-cold partition?
- file size?



## Leveled LSM



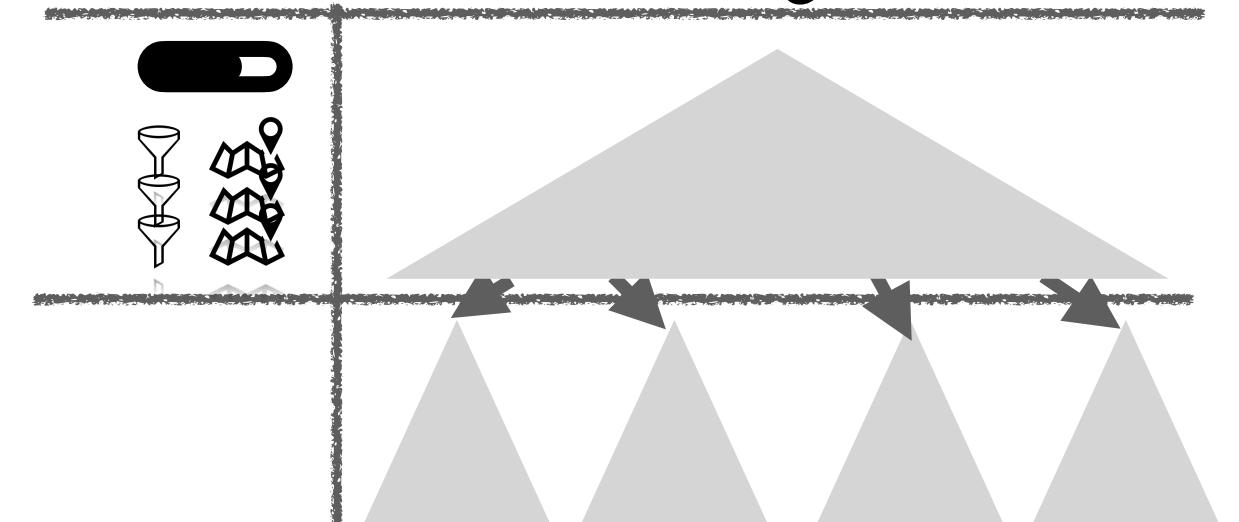
## Tiered LSM



## BTree

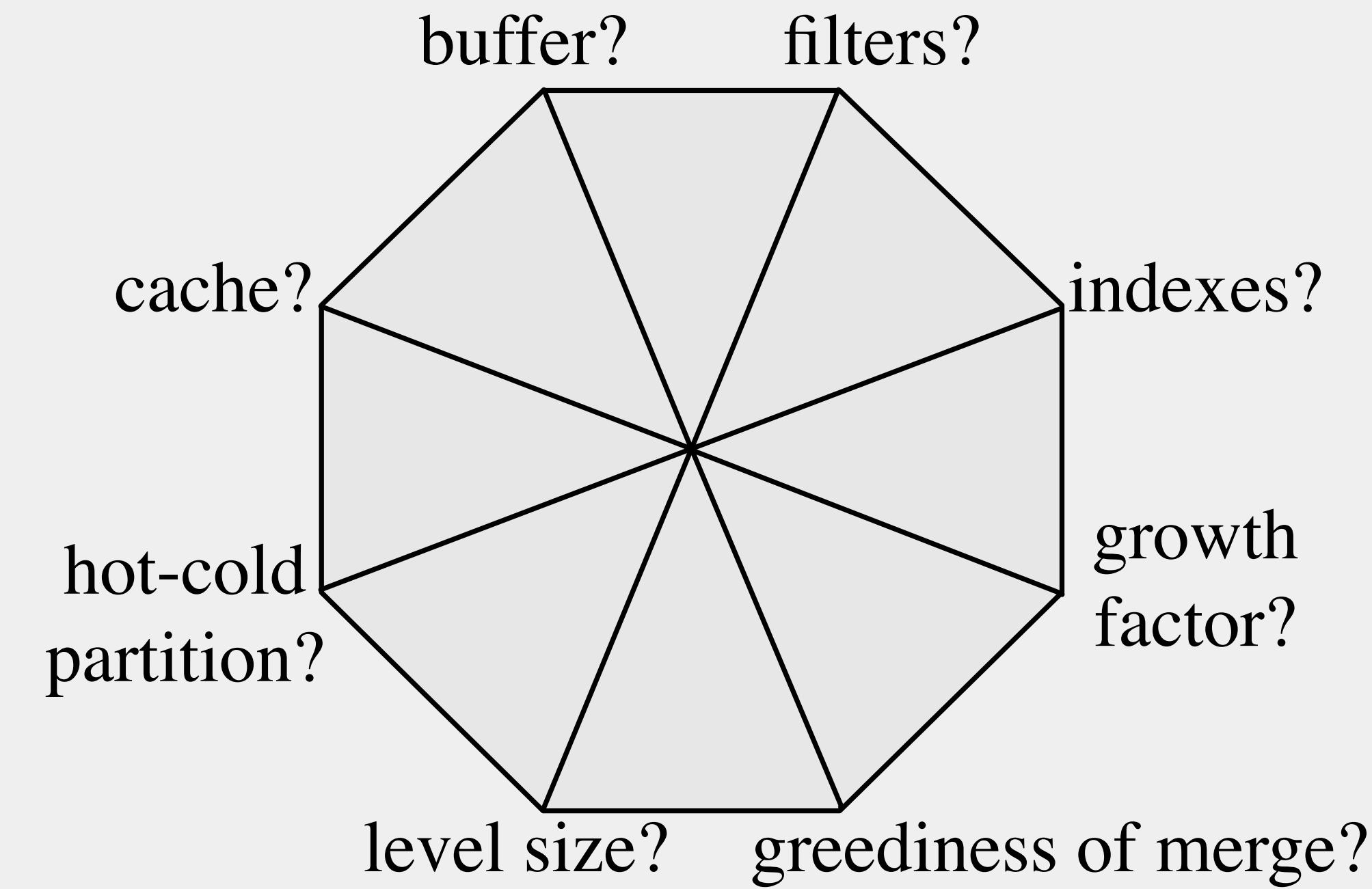


## New layout



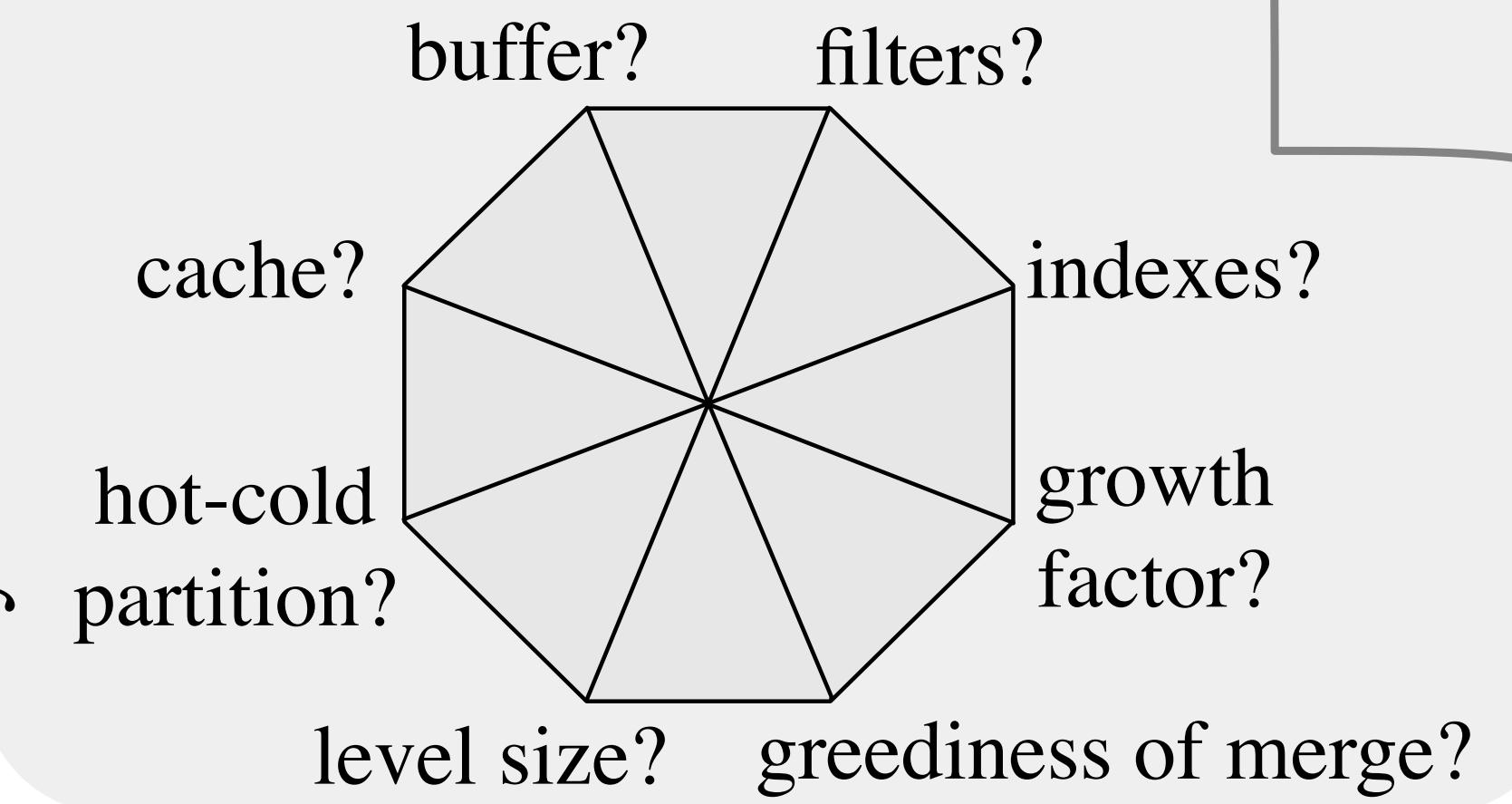
# Storage Engine Template

Layout Primitives

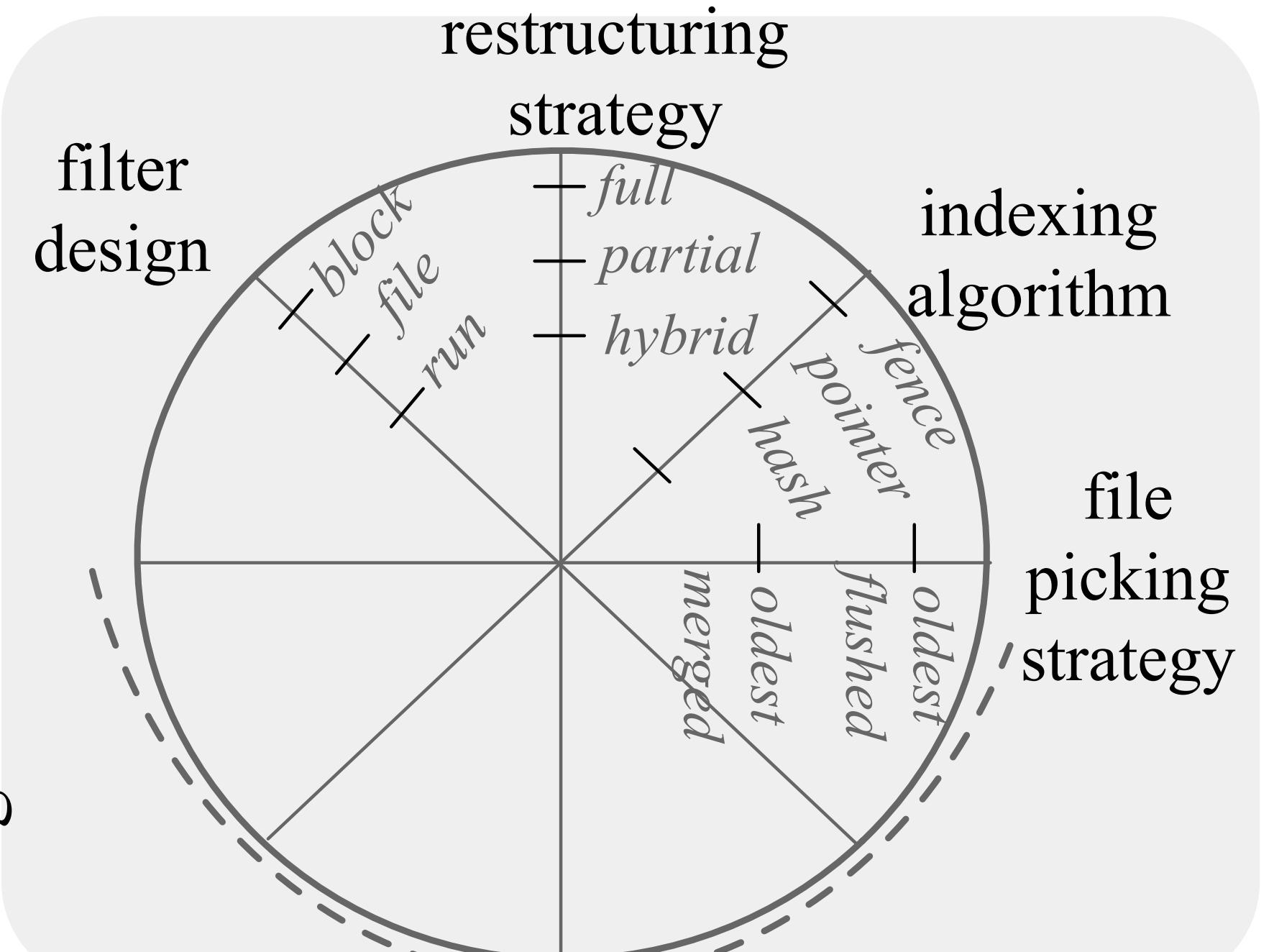


# Storage Engine Template

Layout Primitives

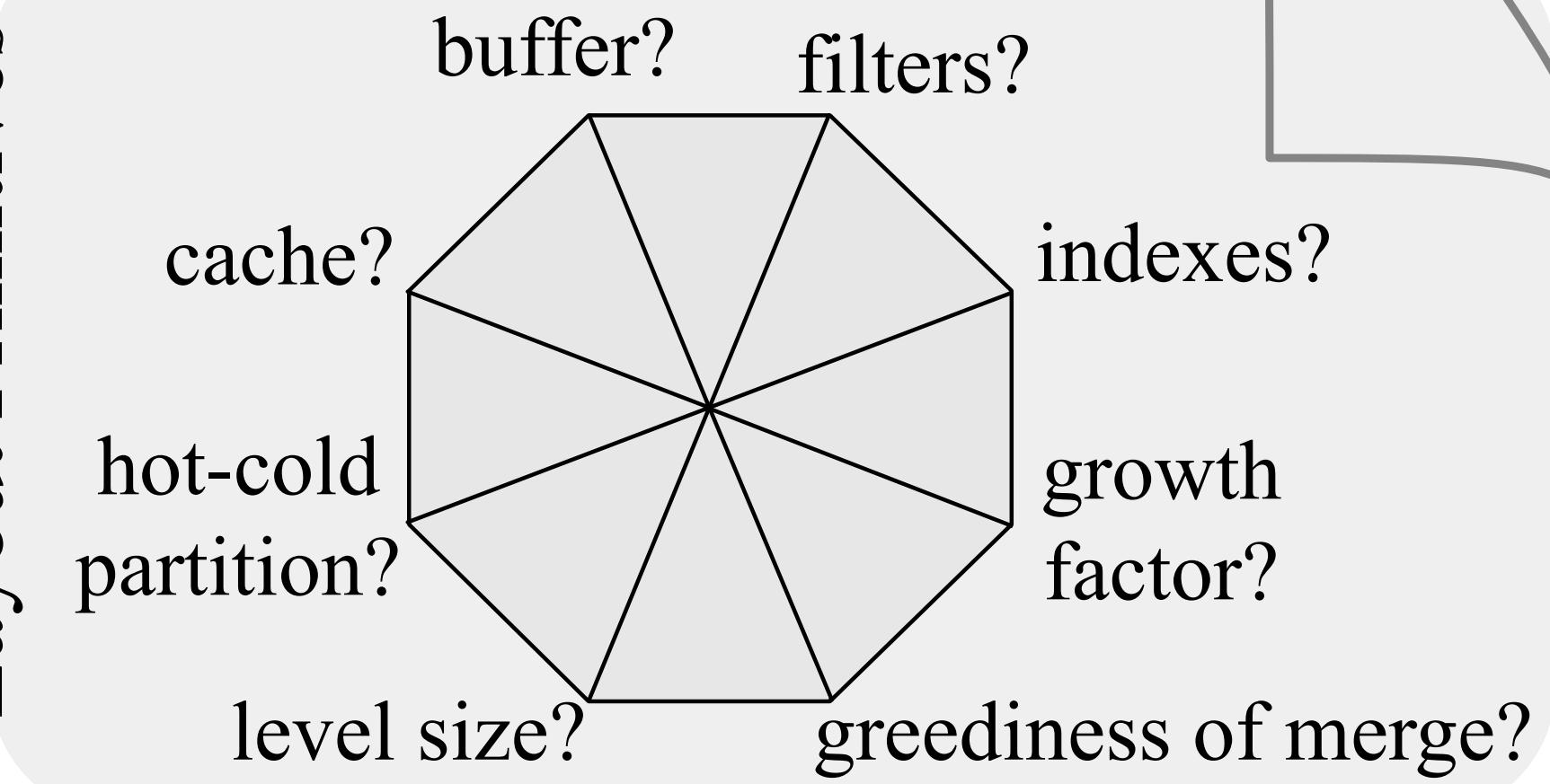


Algorithmic Abstractions

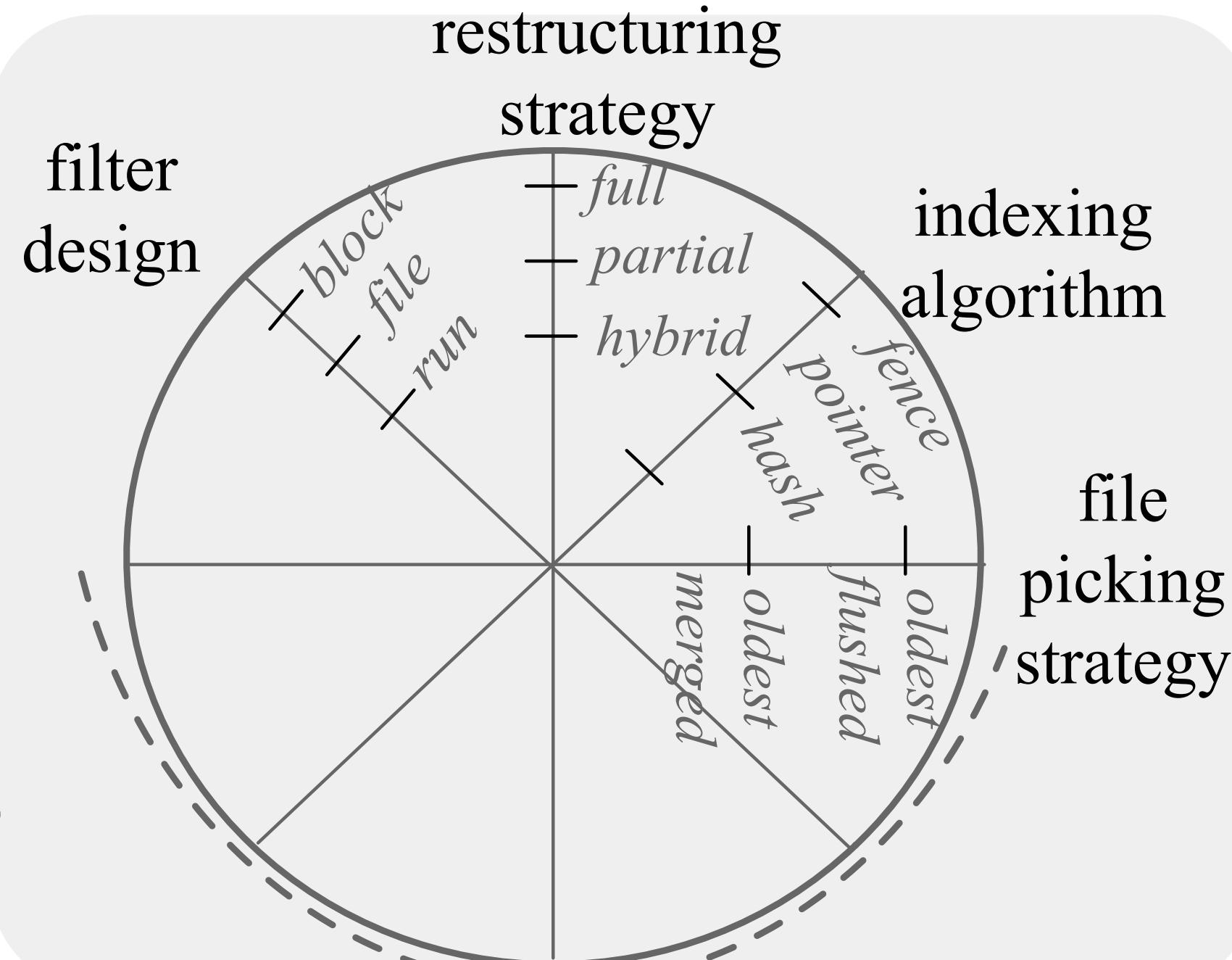


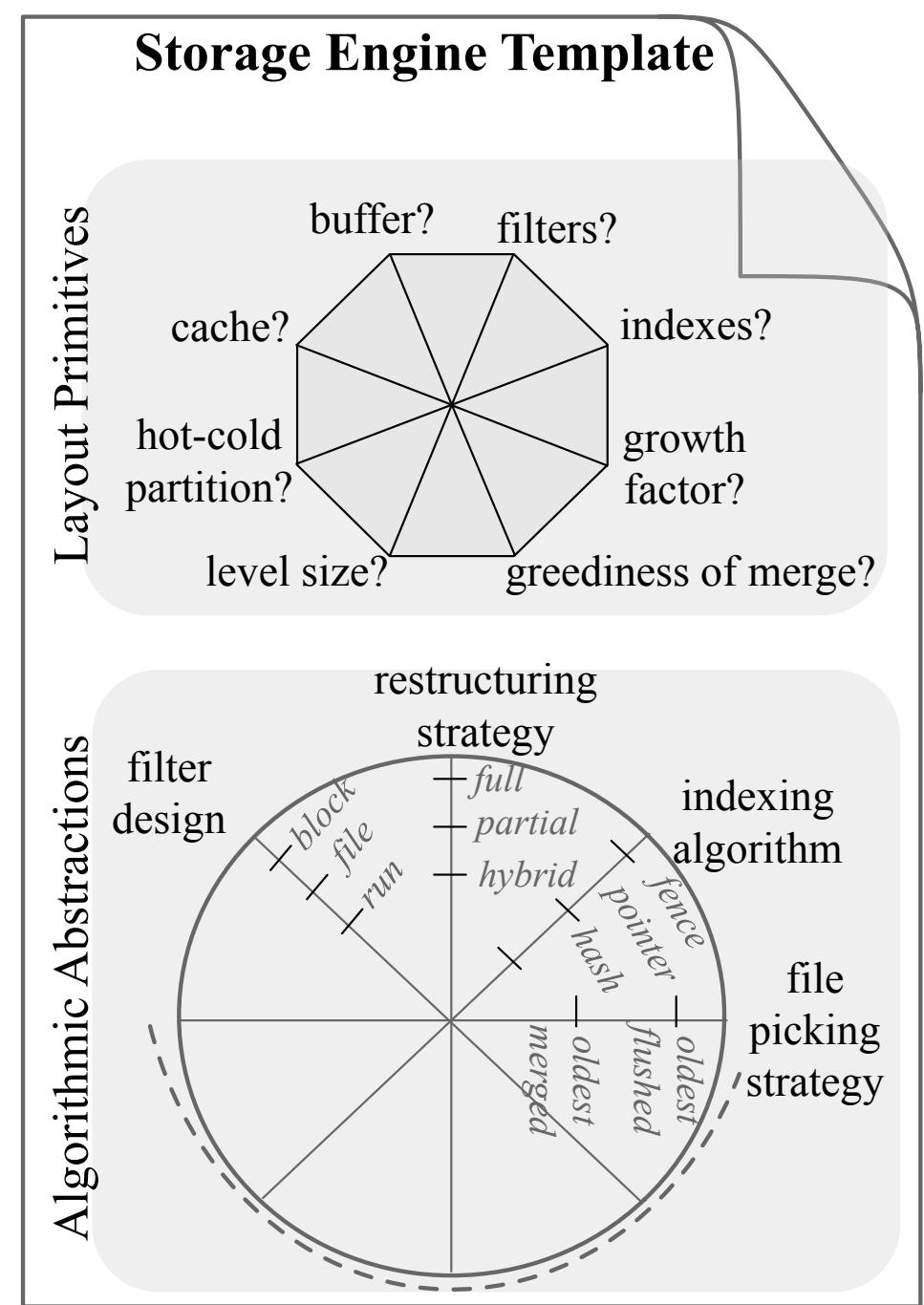
# Storage Engine Template

Layout Primitives



Algorithmic Abstractions

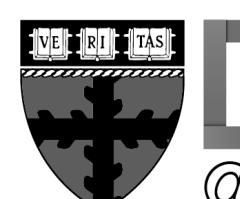




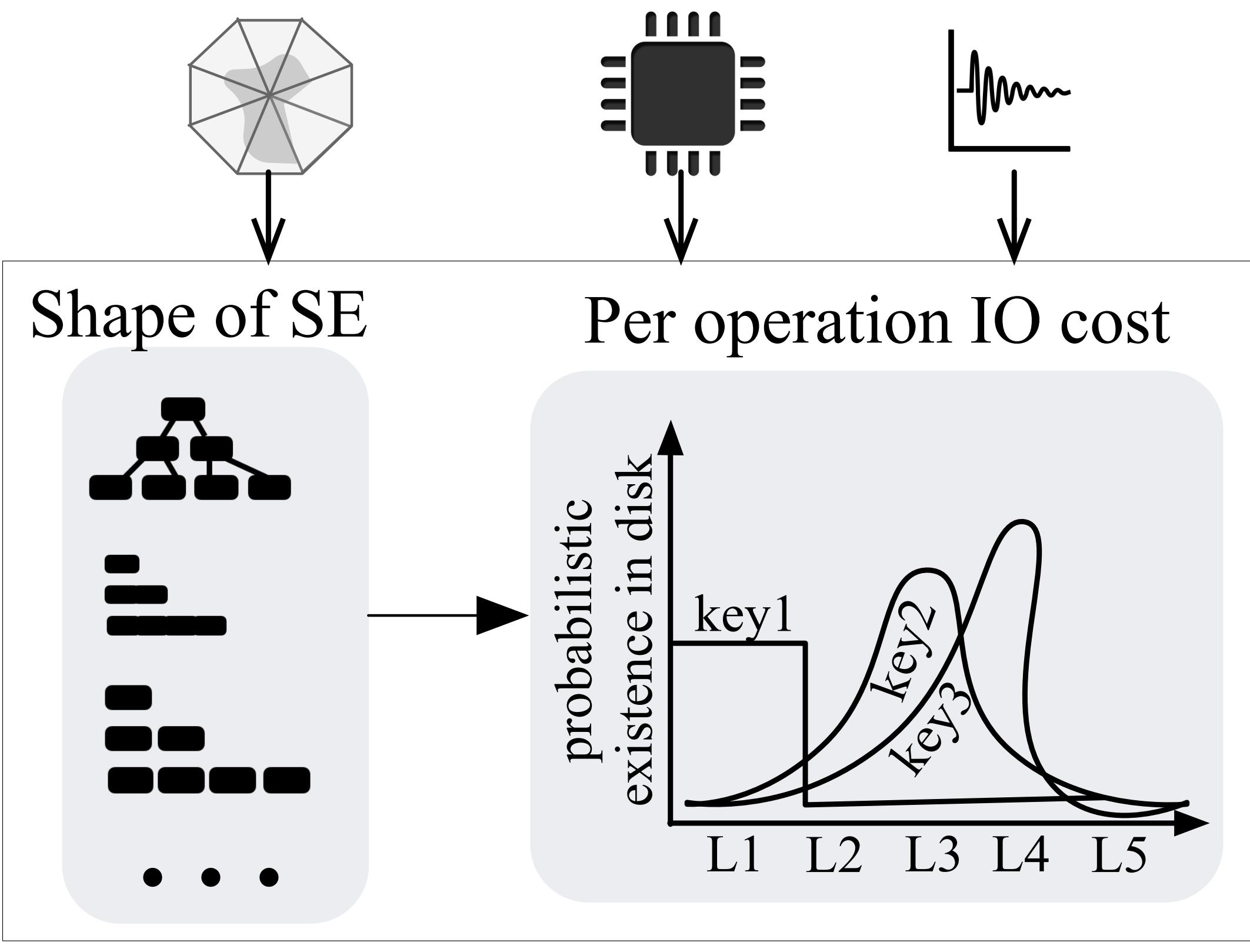
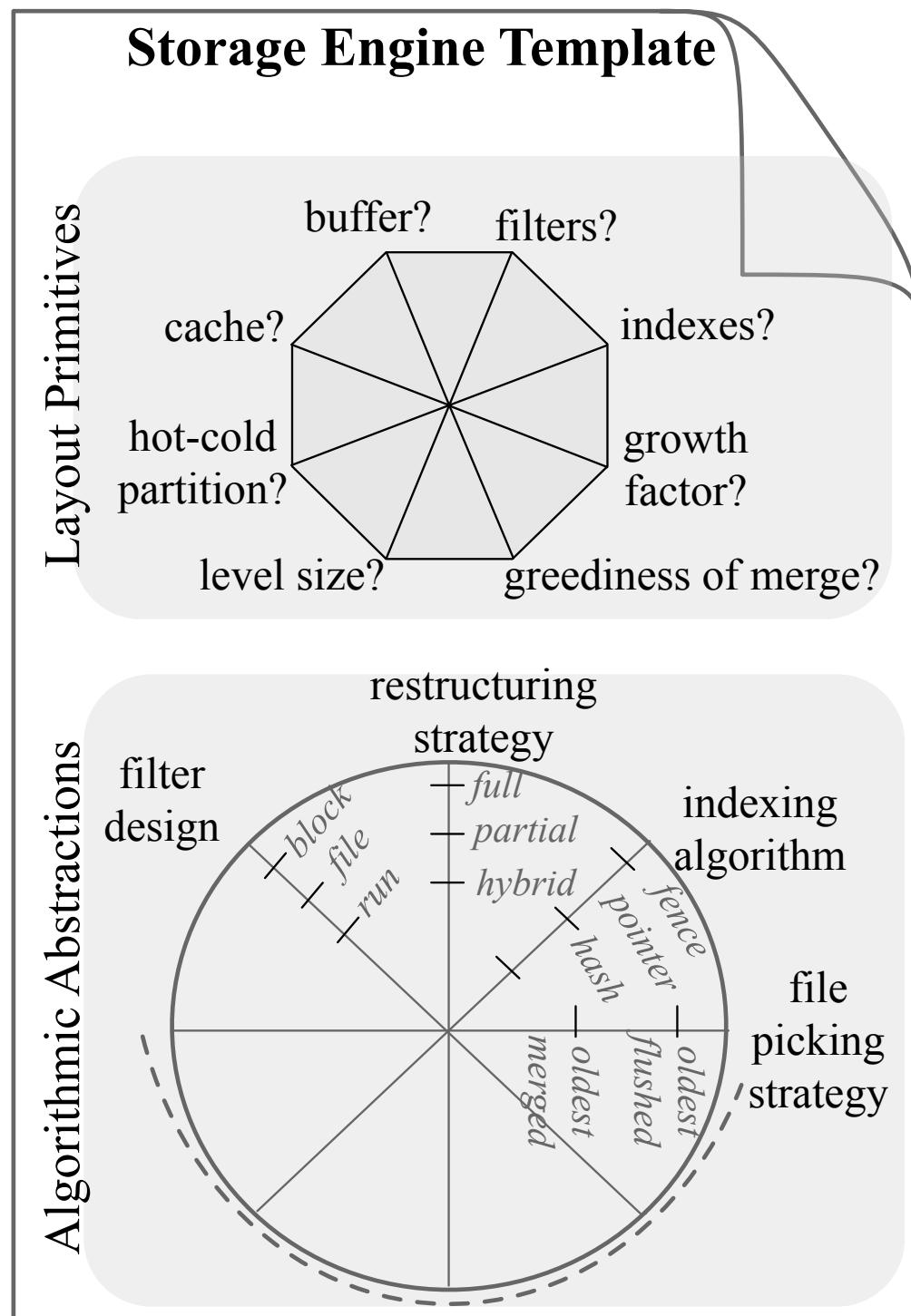
↑ LAYOUT PRIMITIVES → ALGORITHMIC ABSTRACTIONS ↓

Design and hardware specification  
*initialized by search through engine design space*

Design Abstractions of Template		Type/Domain	Example templates for diverse data structures			
			RocksDB variants	WiredTiger variants	FASTER variants	A new design
1.	<b>Key size:</b> Denotes the size of keys in the workload.	unsigned int			auto-configured from the sample workload	
2.	<b>Value size:</b> Denotes the size of values in the workload. All values are accepted as variable-length strings.	string/slice <i>max size set to 1 GB</i>			auto-configured from the sample workload	
3.	<b>Size ratio (T):</b> The maximum number of entries in a block (e.g. growth factor in LSM trees or fanout of B-trees).	unsigned integer   function (func)	[2,.. 32]	[32, 64, 128, 256, ..]	[1000, 1001, ...] (T is large)	2
4.	<b>Runs per hot level (K):</b> At what capacity hot levels are compacted. Rule: should be less than size ratio.	unsigned int	[1.. T]		[T-1]	7
5.	<b>Runs per cold level (Z):</b> At what capacity cold levels are compacted. Rule: should be less than size ratio.	unsigned int	[1.. T]	[1]		32
6.	<b>Logical block size (B):</b> Number of consecutive disk blocks.	unsigned int			[2048, 4096, ...]	
7.	<b>Buffer capacity (<math>M_B</math>):</b> Denotes the amount of memory allocated to in-memory buffer/memtables. Configurable w.r.t file size.	64-bit floating point   function (func)	[64 MB, 128 MB, ...]	[1 MB, 2 MB, ...]	[64 MB, 128 MB, ...]	h/w dependent
8.	<b>Indexes (<math>M_{FP}</math>):</b> Amount of memory allocated to indexes (fence pointers/hashtables).	64-bit floating point   function (func)	memory to cover L	memory for first level	memory for hash table	h/w dependent
9.	<b>Bloom filter memory (<math>M_{BF}</math>):</b> Denotes the bits/entry assigned to Bloom filters.	64-bit float   func(FPR)	10 bits/key			func(FPR)
10.	<b>Bloom filter design:</b> Denotes the granularity of Bloom filters, e.g., one Bloom filter instance per block or per file or per run. The default is file.	block   file   run	file			file
11.	<b>Compaction/Restructuring algorithm:</b> Full does level-to-level compaction; partial is file-to-file; and hybrid uses both full and partial at separate levels.	partial   full   hybrid	full, partial	partial	partial	hybrid
12.	<b>Run strategy:</b> Denotes which run to be picked for compaction (only for partial/hybrid compaction).	first   last_full   fullest	first, fullest, last_full		first	fullest
13.	<b>File picking strategy:</b> Denotes which file to be picked for compaction (for partial/hybrid compaction). For LSM-trees we set default to dense_fp as it empirically works the best. B-trees pick the first file found to be full. LSH-table restructures at the granularity of runs.	oldest_merged   oldest_flushed   dense_fp   sparse_fp   choose_first	dense_fp	choose_first		dense_fp (hot), choose_first (cold)
14.	<b>Merge threshold:</b> If a level is more than x% full, a compaction is triggered.	64-bit floating point	[0.7..1]	0.5		0.75
15.	<b>Full compaction levels:</b> Denotes how many levels will have full compaction (only for hybrid compaction). The default is set to 2.	unsigned integer   function (func)	[1..L]			L-Y (from optimal config)
16.	<b>No. of CPUs:</b> Number of available cores to use in a VM.	unsigned int			Use all available cores	
17.	<b>No of threads:</b> Denotes how many threads are used to process the workload.	unsigned int			Use 1 thread per CPU core	

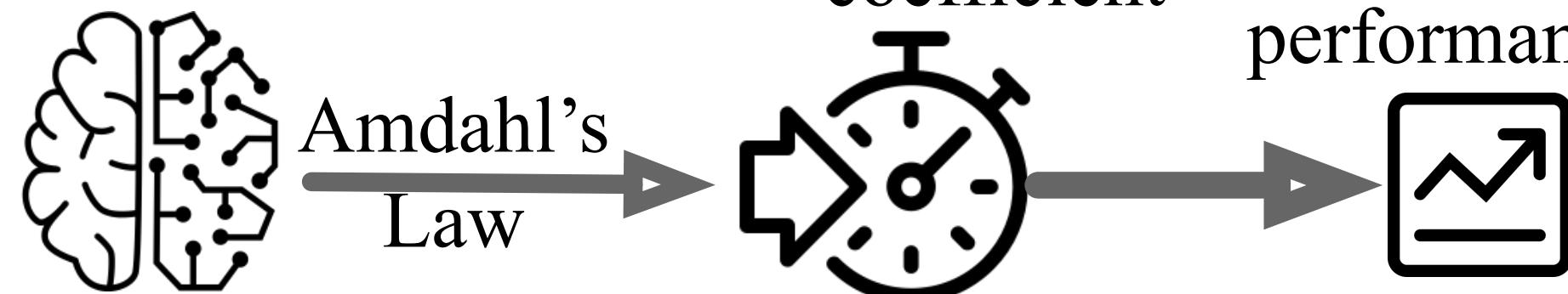


# Distribution-Aware I/O Model



# Concurrency-Aware CPU Model

Empirically learn proportion of parallelizable component



Get speedup coefficient

Get actual end-to-end performance

