### CS660: Intro to Database Systems

# Class 26: More Recovery & NoSQL

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https://bu-disc.github.io/CS660/

## **Course Evaluation**

12:30-12:45 course evaluation

https://tinyurl.com/CS660-F23-CourseEval

if the above does not work:

https://go.blueja.io/inAWTDZkT0CDuuMCUTba5g



Recovery

## **EXAMPLE 1**

# Example 1

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, prevLSN=NULL
15	<i>Update</i> , T1, P2, prevLSN=10
20	Commit, T1, prevLSN=15
	CRASH



### **Active Transaction Table**

[xact\_id, status, lastLSN]

### Dirty Page Table

[page\_id, recLSN]

### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, prevLSN=NULL
15	<i>Update</i> , T1, P2, prevLSN=10
20	Commit, T1, prevLSN=15



Active Transaction Table [xact id, status, lastLSN]

Dirty Page Table [page\_id, recLSN]

Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, prevLSN=NULL
15	<i>Update</i> , T1, P2, prevLSN=10
20	Commit, T1, prevLSN=15



Active Transaction Table T1, running, 10

Dirty Page Table P1, 10

Master Record:

LOG
Begin Checkpoint
End Checkpoint (empty ATT, DPT)
<i>Update</i> , T1, P1, prevLSN=NULL
<i>Update</i> , T1, P2, prevLSN=10
Commit, T1, prevLSN=15



Active Transaction Table T1, running, 10 15

Dirty Page Table

P1, 10 P2, 15

Master Record:

LOG
Begin Checkpoint
End Checkpoint (empty ATT, DPT)
<i>Update</i> , T1, P1, prevLSN=NULL
<i>Update</i> , T1, P2, prevLSN=10
Commit, T1, prevLSN=15

Master Record:

last checkpoint at LSN 00



**Active Transaction Table** 

T1, running, 15 committing, 20

### **Dirty Page Table**

P1, 10

P2, 15

Analysis phase done!

- → need to REDO from oldest recLSN (10)
- → T1 is committed (need to add End)
- → nothing to UNDO

LOG
Begin Checkpoint
End Checkpoint (empty ATT, DPT)
<i>Update</i> , T1, P1, prevLSN=NULL
<i>Update</i> , T1, P2, prevLSN=10
Commit, T1, prevLSN=15



#### **Active Transaction Table**

T1, running, 15 committing, 20

### Dirty Page Table

P1, 10 P2, 15

Start from LSN 10:

Read P1
if (pageLSN(P1)<10):
 redo LSN 10
 set pageLSN(P1)=10

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, prevLSN=NULL
15	<i>Update</i> , T1, P2, prevLSN=10
20	Commit, T1, prevLSN=15



#### **Active Transaction Table**

T1, running, 15 committing, 20

### Dirty Page Table

P1, 10 P2, 15

Move to LSN 15:

Read P2
if (pageLSN(P2)<15):
 redo LSN 15
 set pageLSN(P2)=15

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, prevLSN=NULL
15	<i>Update</i> , T1, P2, prevLSN=10
20	Commit, T1, prevLSN=15



### **Active Transaction Table**

T1, running, 15 committing, 20

### **Dirty Page Table**

P1, 10

P2, 15

Move to LSN 20: Prepare End Record

LOG
Begin Checkpoint
End Checkpoint (empty ATT, DPT)
<i>Update</i> , T1, P1, prevLSN=NULL
<i>Update</i> , T1, P2, prevLSN=10
Commit, T1, prevLSN=15
End, T1, prevLSN=20



### **Active Transaction Table**

T1, committing, 20

### **Dirty Page Table**

P1, 10

P2, 15

#### Redo phase done!

- → nothing to UNDO
- → Recovery is completed!

Add End to LSN 25: Remove T1 from ATT

Recovery

## **EXAMPLE 2**

# Example 2

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
	CRASH



Active Transaction Table

[xact\_id, status, lastLSN]

Dirty Page Table

[page\_id, recLSN]

### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



Active Transaction Table

[xact\_id, status, lastLSN]

Dirty Page Table

[page\_id, recLSN]

#### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20

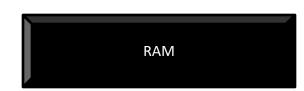


Active Transaction Table T1, running, 10

Dirty Page Table P1, 10

### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



Active Transaction Table T1, running, 15

Dirty Page Table

P1, 10 P2, 15

#### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, running, 15 T2, running, 20

### Dirty Page Table

P1, 10 P2, 15 P3, 20

#### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, running, 15 committing, 25 T2, running, 20

### **Dirty Page Table**

P1, 10

P2, 15

P3, 20

#### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, <del>20</del> 30

### **Dirty Page Table**

P1, 10 P2, 15

P3, 20

## Analysis phase done!

- → need to REDO from oldest recLSN (10)
- → T1 is committed (need to add End)
- → need to undo T2: ToUndo={30}

#### Master Record:

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

P1, 10 P2, 15 P3, 20

Read P1 if (pageLSN(P1)<10):

Start from LSN 10:

redo LSN 10 (set value to ZZZ) set pageLSN(P1)=10

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

P1, 10 P2, 15 P3, 20

Read P2

Move to LSN 15:

```
if (pageLSN(P2)<15):
redo LSN 15 (set value to XXX)
set pageLSN(P2)=15
```

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

P1, 10 P2, 15 P3, 20

Read P3

Move to LSN 20:

```
if (pageLSN(P3)<20):
redo LSN 20 (set value to PPP)
set pageLSN(P3)=20
```

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### **Dirty Page Table**

P1, 10

P2, 15

P3, 20

Move to LSN 25: Prepare End record

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### **Dirty Page Table**

P1, 10 P2, 15 P3, 20

Move to LSN 30:

```
Read P1
if (pageLSN(P1)<30):
    redo LSN 30 (set value to TTT)
    set pageLSN(P1)=30
```

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
35	End, T1, prevLSN=25



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

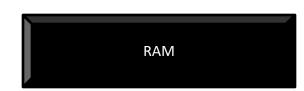
P1, 10 P2, 15 P3, 20

Redo phase done!

→ need to UNDO T2: ToUndo={30}

Add End to 35: Remove T1 from ATT

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
35	End, T1, prevLSN=25



### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

P1, 10 P2, 15

P3, 20 ToUndo={30}

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
35	End, T1, prevLSN=25
40	CLR, T2, Undo LSN 30, P1 (ZZZ), undoNextLSN=20



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### **Dirty Page Table**

P1, 10
P2, 15
P3, 20
ToUndo={30}

#### Undo LSN 30:

write CLR for P1; set ZZZ to P1
ToUndo ← 20

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
35	End, T1, prevLSN=25
40	CLR, T2 Undo LSN 30, P1 (ZZZ), undoNextLSN=20
45	CLR, T2 Undo LSN 20, P3 (LLL), undoNextLSN=NULL



#### **Active Transaction Table**

T1, committing, 25 T2, running, 30

### Dirty Page Table

P1, 10
P2, 15
P3, 20
ToUndo={20}

#### Undo LSN 20:

write CLR for P3; set LLL to P3 ToUndo  $\leftarrow \emptyset$ 

LSN	LOG
00	Begin Checkpoint
05	End Checkpoint (empty ATT, DPT)
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
20	<i>Update</i> , T2, P3, old=LLL, new=PPP, prevLSN=NULL
25	Commit, T1, prevLSN=15
30	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=20
35	End, T1, prevLSN=25
40	CLR, T2 Undo LSN 30, P1 (ZZZ), undoNextLSN=20
45	CLR, T2 Undo LSN 20, P3 (LLL), undoNextLSN=NULL
50	End, T2, prevLSN=45



#### **Active Transaction Table**

T1, committing, 25

### **Dirty Page Table**

P1, 10

P2, 15

P3, 20 ToUndo={ $\emptyset$ }

Write End for T2 & Remove T2 from ATT

Recovery completed!

Recovery

## **EXAMPLE 3**

# Example 3

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
	CRASH

#### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

[xact\_id, status, lastLSN]

### Dirty Page Table

[page\_id, recLSN]

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL 15 *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 20 Begin Checkpoint 25 End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 30 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 *Abort*, T2, prevLSN=40 50 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

#### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

[xact\_id, status, lastLSN]

### Dirty Page Table

[page id, recLSN]

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, running, 10 T2, running, 15

### **Dirty Page Table**

P1, 10 P3, 15

load ATT & DPT from checkpoint!

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 15 20 Begin Checkpoint 25 End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] 30 *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 *Abort*, T2, prevLSN=40 50 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

#### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, running, <del>10</del> 30 T2, running, 15

### **Dirty Page Table**

P1, 10 P3, 15

P2, 30

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL 15 *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 20 Begin Checkpoint 25 End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] 30 *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 Abort, T2, prevLSN=40 50 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, running, 30 committing, 35 T2, running, 15

### **Dirty Page Table**

P1, 10

P3, 15

P2, 30

# Example 3: Analysis Phase

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL 15 *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 20 Begin Checkpoint 25 End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] 30 *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 Abort, T2, prevLSN=40 50 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

T1, committing, 35 T2, running, <del>15</del> 40

### **Dirty Page Table**

P1, 10

P3, 15

P2, 30

# Example 3: Analysis Phase

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL 15 *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 20 Begin Checkpoint 25 End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] 30 *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 *Abort*, T2, prevLSN=40 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15 50

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, running, 40 aborting, 45

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30

# Example 3: Analysis Phase

#### LSN LOG 10 *Update*, T1, P1, old=YYY, new=ZZZ, prevLSN=NULL 15 *Update*, T2, P3, old=UUU, new=VVV, prevLSN=NULL 20 Begin Checkpoint End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}] 25 30 *Update*, T1, P2, old=WWW, new=XXX, prevLSN=10 35 *Commit*, T1, prevLSN=30 40 *Update*, T2, P1, old=ZZZ, new=TTT, prevLSN=15 45 *Abort*, T2, prevLSN=40 50 CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### **Master Record:**

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 45 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30

#### Analysis phase done!

- → need to REDO from oldest recLSN (10)
- → T1 is committed (need to add End)
- → need to (continue) undo T2: ToUndo={50}

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

### **Dirty Page Table**

P1, 10 P3, 15 P2, 30

Start from LSN 10:

Read P1
if (pageLSN(P1)<10):
 redo LSN 10 (set value to ZZZ)
 set pageLSN(P1)=10

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

### **Dirty Page Table**

P1, 10 P3, 15 P2, 30

Read P3

Move to LSN 15:

if (pageLSN(P3)<15):
redo LSN 15 (set value to VVV)
set pageLSN(P3)=15

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

## **Dirty Page Table**

P1, 10 P3, 15 P2, 30

Move to LSN 30:

Read P2
if (pageLSN(P2)<30):
 redo LSN 30 (set value to XXX)
 set pageLSN(P2)=30

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

## **Dirty Page Table**

P1, 10 P3, 15

P2, 30

Move to LSN 35: Prepare End Record

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

#### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

### **Dirty Page Table**

P1, 10 P3, 15 P2, 30

Read P1 (if it has been flushed at chkpt, check pageLSN in storage; it should be 10) redo LSN 40 (set value to TTT) set pageLSN(P1)=40

Move to LSN 40:

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15

#### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

### **Dirty Page Table**

P1, 10 P3, 15 P2, 30

Read P1 (if it has been flushed, check pageLSN in storage; it should be 40) redo LSN 50 (set value to ZZZ) set pageLSN(P1)=50

Move to LSN 50:

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35

#### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30

Add End to 55: Remove T1 from ATT

Redo phase done!

→ need to UNDO T2: ToUndo={50}

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35

### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30 ToUndo={50}

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35

### Master Record:

last checkpoint at LSN 20



### **Active Transaction Table**

T1, committing, 35 T2, aborting, 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30 ToUndo={50}

Undo LSN 50:

ToUndo ← 15

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35
T2, aborting, 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30

ToUndo={<del>50</del> 15}

Undo LSN 50:

ToUndo ← 15

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35
60	CLR, T2 Undo LSN 15, P3 (UUU), prevLSN=50, undoNext=NULL

### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35
T2, aborting, 50

## **Dirty Page Table**

P1, 10

P3, 15

P2, 30

ToUndo={15}

#### Undo LSN 15:

write CLR for P3; set UUU to P3 ToUndo  $\leftarrow \emptyset$ 

LSN	LOG
10	<i>Update</i> , T1, P1, old=YYY, new=ZZZ, prevLSN=NULL
15	<i>Update</i> , T2, P3, old=UUU, new=VVV, prevLSN=NULL
20	Begin Checkpoint
25	End Checkpoint: ATT=[{T1,r,10}, {T2,r,15}] DPT=[{P1,10}, {P3,15}]
30	<i>Update</i> , T1, P2, old=WWW, new=XXX, prevLSN=10
35	Commit, T1, prevLSN=30
40	<i>Update</i> , T2, P1, old=ZZZ, new=TTT, prevLSN=15
45	Abort, T2, prevLSN=40
50	CLR, T2 Undo LSN 40, P1 (ZZZ), prevLSN=45, undoNextLSN=15
55	End, T1, prevLSN=35
60	CLR, T2 Undo LSN 15, P3 (UUU), prevLSN=50, undoNext=NULL
65	End, T2, prevLSN=60

#### Master Record:

last checkpoint at LSN 20



#### **Active Transaction Table**

T1, committing, 35
T2, aborting, 50

### **Dirty Page Table**

P1, 10

P3, 15

P2, 30

ToUndo={Ø}

Write End for T2 & Remove T2 from ATT

Recovery completed!

## Summary

Checkpointing: A quick way to limit the amount of log to scan on recovery.

## Recovery works in 3 phases:

Analysis: Forward from checkpoint.

Redo: Forward from oldest recLSN.

Undo: Backward from end to first LSN of oldest Xact alive at crash.

Upon Undo, write CLRs.

Redo "repeats history": Simplifies the logic!

## CS660: Intro to Database Systems

Class 26: NoSQL

Instructor: Manos Athanassoulis

https://bu-disc.github.io/CS660/

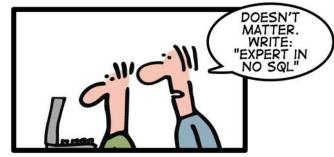
## What is NoSQL?

from "Geek and Poke"

### HOW TO WRITE A CV







Leverage the NoSQL boom

## What is NoSQL?

An emerging "movement" around non-relational software for Big Data

Roots are in the Google, Amazon, Facebook homegrown software stacks







A NoSQL system provides a mechanism for storage and retrieval of data that uses looser consistency models than traditional relational databases in order to achieve horizontal scaling and higher availability.

NoSQL comes from "Not SQL" or "Not only SQL" to emphasize that some NoSQL systems allow SQL-like queries.

## **NoSQL Stores**

offer an easy to program storage model

simplification of relational

two attributes (a key and a value) value has variable size

## NoSQL features

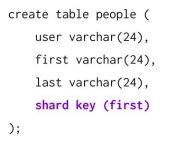
## Scalability is crucial!

load increased rapidly for many applications

## Large servers are expensive

## Solution: use clusters of small commodity machines

- need to shard the data (maybe use replication)
- cheap (usually open source!)
- cloud-based storage



people people\_0 people\_1 people\_2 people\_3 first first last first last nancy nick jsmith iohn smith tholmes tom

image from: singlestore.com

## NoSQL features

Sometimes not a well-defined schema

### Allow for semi-structured data

- still need to provide ways to query efficiently (use of index methods)
- need to express specific types of queries easily

# Scalability

Often cited as the main reason for moving from DB technology to NoSQL

DB Position: there is no reason a parallel DBMS cannot scale to 1000's of nodes

NoSQL Position: a) Prove it; b) it will cost too much anyway

## Flavors of NoSQL

## Four main types:

key-value stores
document databases
column-family (aka big-table) stores
graph databases

Here we will talk more about "Document" databases (MongoDB)

# **Key-Value Stores**

There are many systems like that:







Simple data model: key/value pairs the DBMS *does not attempt to interpret* the value

Queries are limited to query by key

- get/put/update/delete a key/value pair
- iterate over key/value pairs

## **Document Databases**

Examples include:





Special type of key/value that *value is a document*.

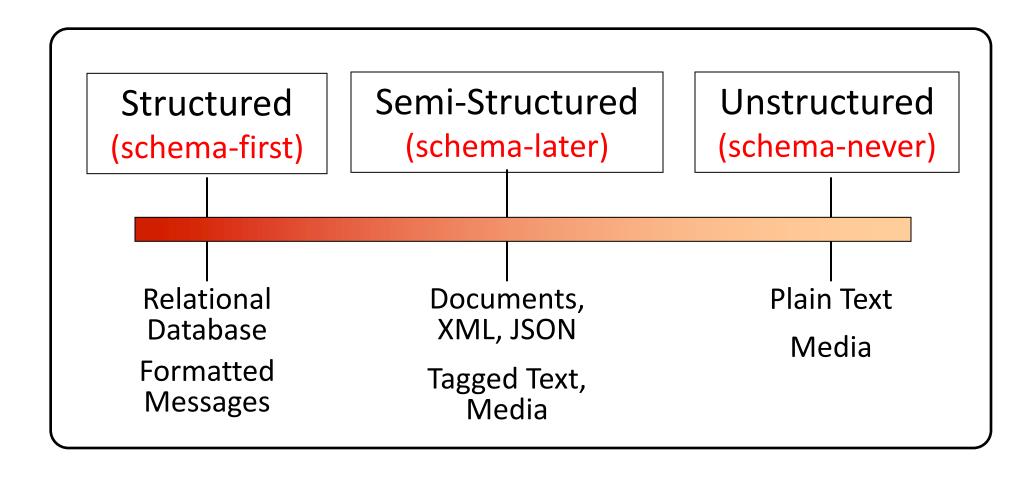
- use some sort of semi-structured data model: XML/JSON
- the *value can be examined* and used by the system (unlike in key/data stores)

Queries based on key (as in key/value stores), but also on the document (value).

Here again, there is support for *sharding* and *replication*.

the sharding can be based on values within the document

# The Structure Spectrum



# Recap: NoSQL

simplification of relational: {key, value}

unique key variable-size value

## Key/Value stores:

- get/put/update/delete a key/value pair
- iterate over key/value pairs

## <u>Document</u> stores (value is a *semi-structured* document ):

- use some sort of semi-structured data model: XML/JSON
- the value can be examined and used by the system (unlike in key/data stores)
- queries based on key (as in key/value stores), but also on the document (value).

## MongoDB (An example of a Document Database)

Data are organized in *collections*. A <u>collection</u> stores a *set of documents*.

Collection (like table) and document (like record)

- BUT each document can have different attributes even in the same collection
- Semi-structured schema!

Only requirement: every document should have an "\_id" field

– humongous => Mongo

# Example MongoDB

```
"_id":ObjectId("4efa8d2b7d284dad101e4bc9"),
 "Last Name": " Cousteau",
 "First Name": "Jacques-Yves",
 "Date of Birth": "06-1-1910" },
"_id": ObjectId("4efa8d2b7d284dad101e4bc7"),
 "Last Name": "PELLERIN",
 "First Name": "Franck",
 "Date of Birth": "09-19-1983",
 "Address": "1 chemin des Loges",
"City": "VERSAILLES" }
```

# Example Document Database: MongoDB

Key features include:

JSON-style documents

actually, uses BSON (JSON's binary format)

replication for high availability

auto-sharding for scalability

key & document-based queries

can create an index on any attribute for faster reads

under the hood, a simple key-value store called WiredTiger! design based on LSM-trees/B-Trees

# MongoDB Terminology

```
relational term <==> MongoDB equivalent
database <==> database
table <==> collection
row <==> document
attributes <==> fields (field-name:value pairs)
primary key <==> the id field, which is the key associated with
the document
```

## **JSON**

### JSON is an alternative data model for semi-structured data

JavaScript Object Notation

### Built on two key structures:

- an *object*, which is a sequence of name/value pairs {"\_id": "1000", "name": "Sanders Theatre", "capacity": 1000 }
- an array of values [ "123", "222", "333" ]

### A *value* can be:

- an atomic value: string, number, true, false, null
- an object
- an array

# The \_id Field

## Every MongoDB document must have an \_id field.

its value must be unique within the collection acts as the primary key of the collection it is the key in the key/value pair

## If you create a document without an \_id field:

MongoDB adds the field for you assigns it a unique BSON (binary JSON) ObjectID example from the MongoDB shell:

```
> db.test.save({ rating: "PG-13" })
> db.test.find() { "_id" :ObjectId("528bf38ce6d3df97b49a0569"), "rating" : "PG-13" }
```

Note: quoting field names is optional (see rating above)

# Capturing Relationships in MongoDB

## Two options:

1. store references to other documents using their \_id values

2. embed documents within other documents

# Example relationships

```
" id":ObjectId("52ffc33cd85242f436000001"),
   "name": "Tom Benzamin ",
   "contact": "987654321",
   "dob": "01-01-1991"
   " id":ObjectId("52ffc4a5d85242602e000000"),
  "building": "22 A, Indiana Apt",
  "pincode": 123456,
  "city": "Los Angeles",
  "state": "California"
Here an example of reference-based relationship
     " id":ObjectId("52ffc33cd85242f436000001"),
     "contact": "987654321",
    "dob": "01-01-1991",
    "name": "Tom Benzamin",
     "address ids": [
      ObjectId("52ffc4a5d85242602e000000")
```

```
And, here is an example of embedded relationship:
      " id":ObjectId("52ffc33cd85242f436000001"),
      "contact": "987654321".
      "dob": "01-01-1991",
      "name": "Tom Benzamin",
      "address": [
         "building": "22 A, Indiana Apt",
         "pincode": 123456,
         "city": "Los Angeles",
         "state": "California"
         "building": "170 A, Acropolis Apt",
         "pincode": 456789,
         "city": "Chicago",
         "state": "Illinois"
```

### **CRUD**

#### Create

```
> db.collection.insert( <document> )
> db.collection.save( <document> )
> db.collection.update( <query>, <update>, { upsert: true } )
Read
> db.collection.find( <query>, <projection> )
Update
> db.collection.update( <query>, <update>, <options> )
Delete
> db.collection.remove( <query>, <justOne> )
```

## Queries in MongoDB

Each query can only access a single collection of documents. Use a method called

```
> db.collection.find(<selection>, , projection>)
```

### **Example**: find the titles of all R-rated movies:

```
> db.movies.find({ rating: 'R' }, { title: 1 })
```

## Projection

Specify the name of the fields that you want in the output with 1 (0 hides the value)

#### Example:

```
> db.movies.find({},{"title":1,_id:0})
(will report the title but not the id)
```

## Selection

You can specify the condition on the corresponding attributes using the find:

```
> db.movies.find({ rating: "R", year: 2000 }, { title: 1, runtime: 1 })
Operators for other types of comparisons:
```

MongoDB	SQL equivalent
\$gt, \$gte	>, >=
\$lt, \$lte	<, <=
\$ne	!=

**Example**: find the names of movies with an earnings <= 200000

```
> db.movies.find({ earnings: { $1te: 200000 }})
```

For logical operators \$and, \$or, \$nor

use an array of conditions and apply the logical operator among the array conditions:

```
> db.movies.find({ $or: [ { rating: "R" }, { rating: "PG-13" } ] })
```

## Demo

https://mongoplayground.net/

### Other Structure Issues

#### **NoSQL**

- a) Tables are unnatural
- b) "joins" are evil
- c) need to be able to "grep" my data

#### DB

- a) Tables are a natural/neutral structure
- b) data independence lets you precompute joins under the covers
- c) this is a price of all the DBMS goodness you get

This is an Old Debate - Object-oriented databases, XML DBs, Hierarchical, ...

## Fault Tolerance

DBs: coarse-grained FT – if trouble, restart transaction

- Fewer, Better nodes, so failures are rare
- Transactions allow you to kill a job and easily restart it

NoSQL: Massive amounts of cheap HW, failures are the norm and massive data means long running jobs

- So must be able to do mini-recoveries
- This causes some overhead (file writes)

CS660: Intro to Database Systems

# Database Systems and Beyond

Instructor: Manos Athanassoulis

https://bu-disc.github.io/CS660/

## Database Systems

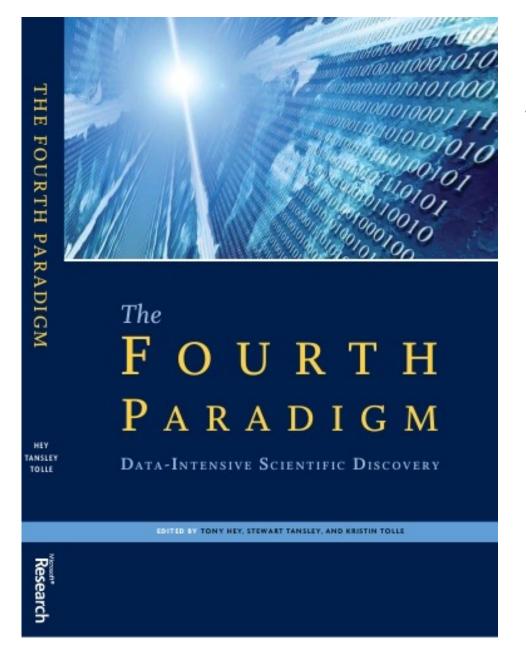
we spent a whole semester on Database Systems what is next?

what can we do with data?

data-driven science

data-driven discovery

data-driven governance



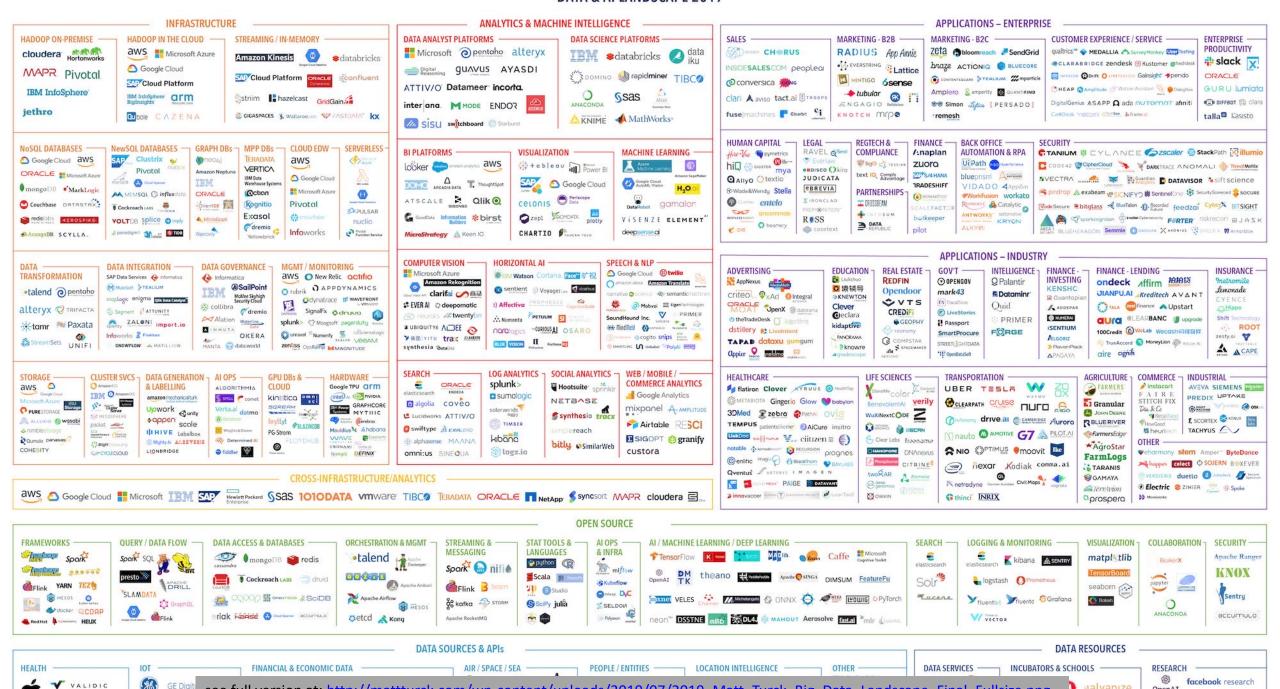
"Experimental, theoretical, and computational science are all being affected by the data deluge, and a fourth, 'data-intensive' science paradigm is emerging.

The goal is to have a world in which all of the science literature is online, all of the science data is online, and they interoperate with each other. Lots of new tools are needed to make this happen."

Faster Innovation through Data-Intensive Approaches

Need for Innovation in Data Management!

#### **DATA & AI LANDSCAPE 2019**

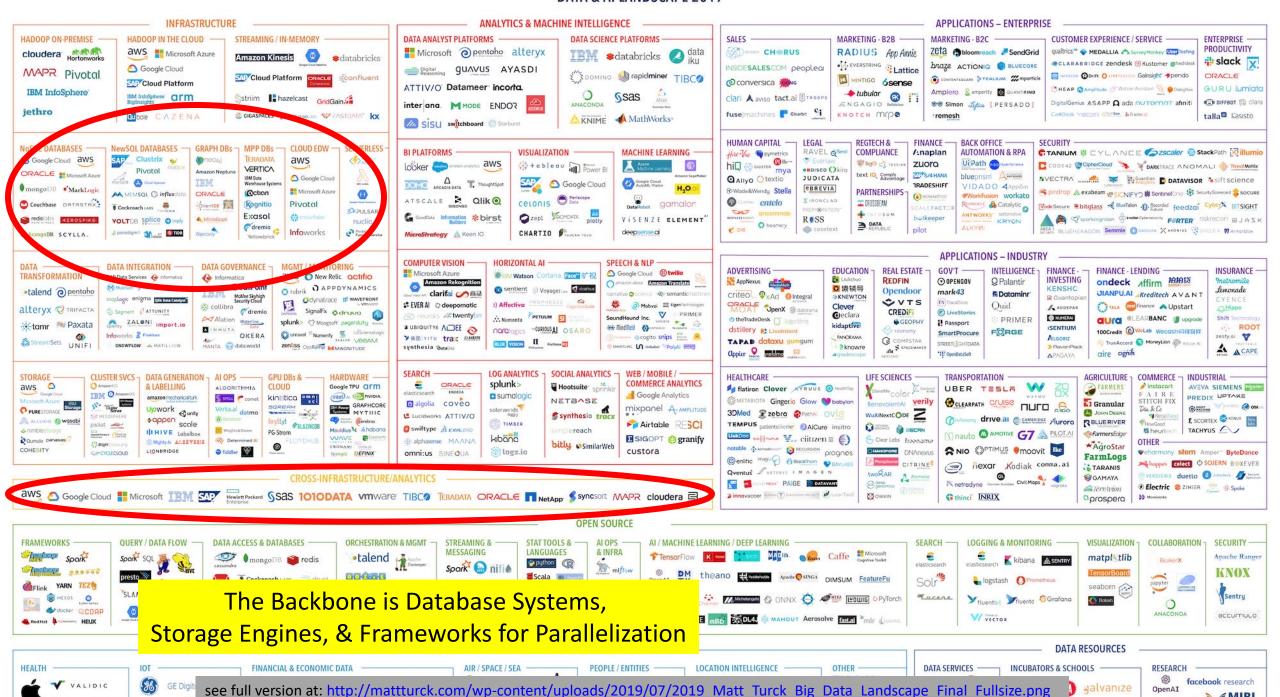


see full version at: http://mattturck.com/wp-content/uploads/2019/07/2019 Matt Turck Big Data Landscape Final Fullsize.png

OpenAI

MIRI MIRI

#### DATA & AI LANDSCAPE 2019



MIRI

### increase throughput by parallelization

"scale-up"
use more powerful machines (>#CPUs, >RAM)

"scale-out" use more machines

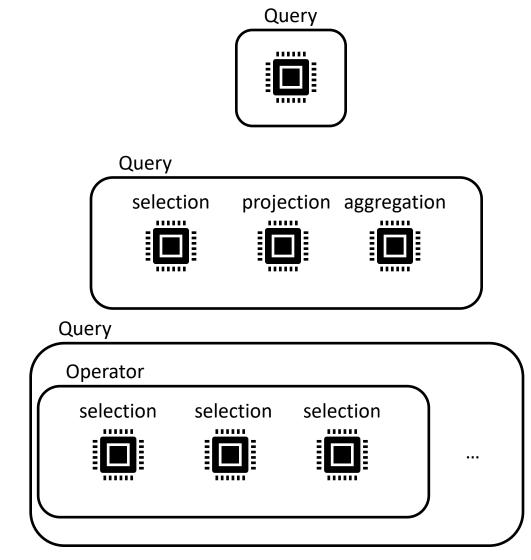
# Scale Up Execution

how to use more cores (threads)?

inter-query parallelism each query runs on one processor

inter-operator parallelism
each query runs on multiple processors
an operator runs on one processor

intra-operator parallelism an operator runs on multiple processors



## Scale Up Storage

needs more disks!

how to distribute data?

block partition hash partition range partition

how to distribute data accesses?

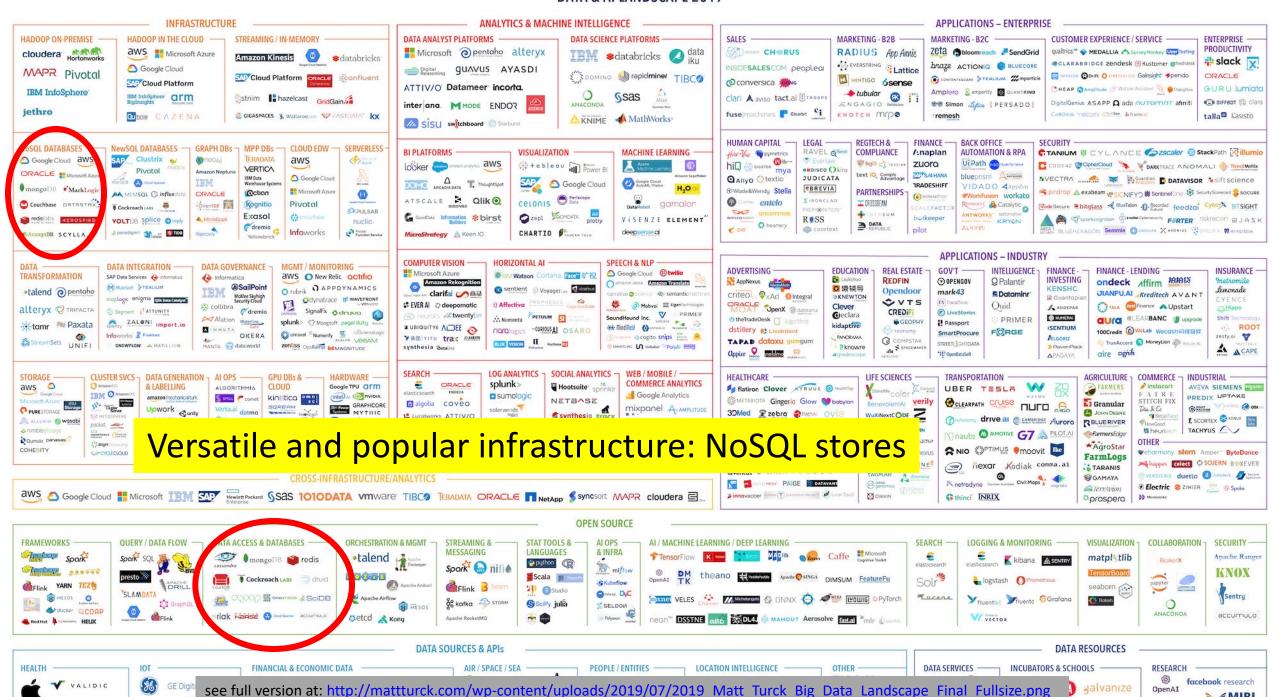
## Scale Out

similar questions across machines

new bottlenecks?

move data across machines: network!

#### DATA & AI LANDSCAPE 2019



MIRI.

diving into the internals of modern data systems

cutting-edge designs / research projects / engineering projects

CS 561: Data Systems Architectures
Spring 2024

# A path in data science & data engineering

- (1) strong data systems skills
  - (i) coding skills
  - (ii) system architecture insights performance tradeoffs

(2) application domain knowledge

(3) statistics, machine learning, math tools

## **Open Discussion**

Questions?
for NoSQL
for DBMS
for next semester (CS561!)
for life after college (Academia vs. Industry vs. ?)

Next: Review and questions for final