Intro: Greatest Common Divisors I

Daniel Kane

Department of Computer Science and Engineering University of California, San Diego

Data Structures and Algorithms Algorithmic Toolbox

Learning Objectives

- Define greatest common divisors.
- Compute greatest common divisors inefficiently.

- Put fraction $\frac{a}{b}$ in simplest form.
- Divide numerator and denominator by d, to get $\frac{a/d}{b/d}$.

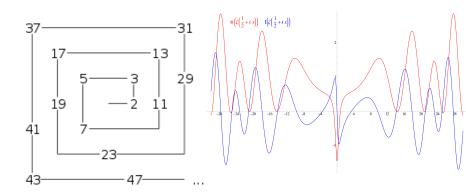
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Definition

For integers, a and b, their greatest common divisor or gcd(a, b) is the largest integer d so that d divides both a and b.

Number Theory



Cryptography



Computation

Compute GCD

Input: Integers $a, b \ge 0$.

Output: gcd(a, b).

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Output: gcd(a, b).

Run on large numbers like

gcd(3918848, 1653264).

Naive Algorithm

Function NaiveGCD(a, b)

```
best \leftarrow 0
```

for d from 1 to a+b: if d|a and d|b:

 $best \leftarrow d$

return best

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best \leftarrow 0
for d from 1 to a + b:
    if d|a and d|b:
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- \blacksquare Runtime approximately a + b.
- Very slow for 20 digit numbers.

Intro: Greatest Common Divisors II

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Data Structures and Algorithms Algorithmic Toolbox

Learning Objectives

- Implement the Euclidean Algorithm.
- Approximate the runtime.

Definition

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Compute GCD

Input: Integers $a, b \ge 0$.

Output: gcd(a, b).

Key Lemma

Lemma

Let a' be the remainder when a is divided by b, then

$$\gcd(a,b)=\gcd(a',b)=\gcd(b,a').$$

Proof

Proof (sketch)

- lacksquare a = a' + bq for some q
- d divides a and b if and only if it divides
 a' and b

Euclidean Algorithm

Function EuclidGCD(a, b) if b = 0:

return a $a' \leftarrow$ the remainder when a is divided by breturn EuclidGCD(b, a')

Euclidean Algorithm

Function EuclidGCD(a, b)

 $\begin{array}{l} \text{if } b = 0: \\ \text{return } a \\ a' \leftarrow \text{the remainder when } a \text{ is} \\ \text{divided by } b \\ \text{return EuclidGCD}(b, a') \end{array}$

Produces correct result by Lemma.

gcd(3918848, 1653264)

```
gcd(3918848, 1653264)
= gcd(1653264, 612320)
```

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```

Runtime

- Each step reduces the size of numbers by about a factor of 2.
- Takes about log(ab) steps.

Runtime

- Each step reduces the size of numbers by about a factor of 2.
- Takes about log(ab) steps.
- GCDs of 100 digit numbers takes about 600 steps.
- Each step a single division.

Summary

- Naive algorithm is too slow.
- The correct algorithm is much better.
- Finding the correct algorithm requires knowing something interesting about the problem.