DYNAMIC PROGRAMMING

Introduction

SIMPLE RECURSIVE ALGORITHMS I

- A simple recursive algorithm:
 - Solves the base cases directly
 - Recurs with a simpler subproblem
 - Does some extra work to convert the solution to the simpler subproblem into a solution to the given problem
- I call these "simple" because several of the other algorithm types are inherently recursive

DIVIDE AND CONQUER

- A divide and conquer algorithm consists of two parts:
 - Divide the problem into smaller subproblems of the same type, and solve these subproblems recursively
 - Combine the solutions to the subproblems into a solution to the original problem
- Traditionally, an algorithm is only called "divide and conquer" if it contains at least two recursive calls

GREEDY ALGORITHMS

- An optimization problem is one in which you want to find, not just *a* solution, but the *best* solution
- A "greedy algorithm" sometimes works well for optimization problems
- A greedy algorithm works in phases: At each phase:
 - You take the best you can get right now, without regard for future consequences
 - You hope that by choosing a *local* optimum at each step, you will end up at a *global* optimum

EXAMPLE: COUNTING MONEY

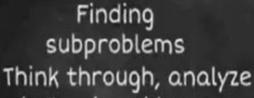
- Suppose you want to count out a certain amount of money, using the fewest possible bills and coins
- A greedy algorithm would do this would be: At each step, take the largest possible bill or coin that does not overshoot
 - Example: To make \$6.39, you can choose:
 - o a \$5 bill
 - a \$1 bill, to make \$6
 - \circ a 25¢ coin, to make \$6.25
 - \circ A 10¢ coin, to make \$6.35
 - o four 1¢ coins, to make \$6.39
- For US money, the greedy algorithm always gives the optimum solution

Dynamic Programming

- Dynamic Programming(DP) is a method for solving complex problems by breaking them down into simpler sub problems".
- Also refer to as "Smart Recursion" or "Intelligent brute-force"
- For example DP is like building a wall brick by brick. Here each brick is a sub-problem and building a wall is a complex problem.
- In DP we will not repeat a sub problem we memorize the result

TO WRITE A DYNAMIC PROGRAMMING SOLUTION





what subproblems are needed to be solved in order to complete the main problem.



Recursion

Write a recursive solution to the given problem such that each call to the function complete one subproblem



Memoization

Store the results of each subproblem to be used later whenever the same subproblem occur again.

DYNAMIC PROGRAMMING ALGORITHMS

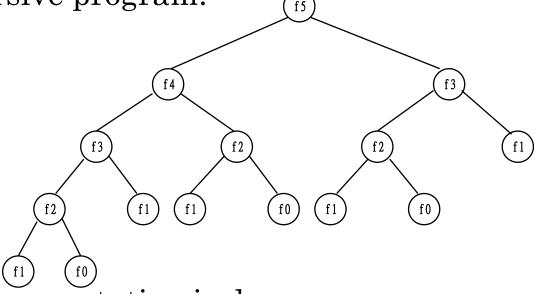
- A dynamic programming algorithm remembers past results and uses them to find new results
- Dynamic programming is generally used for optimization problems
 - Multiple solutions exist, need to find the "best" one
 - Requires "optimal substructure" and "overlapping subproblems"
 - Optimal substructure: Optimal solution contains optimal solutions to subproblems
 - Overlapping subproblems: Solutions to subproblems can be stored and reused in a bottom-up fashion
- This differs from Divide and Conquer, where subproblems generally need not overlap

FIBONACCI SEQUENCE

• Fibonacci sequence: 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

$$\begin{aligned} \mathbf{F_i} &= i & \text{if } i \leq 1 \\ \mathbf{F_i} &= \mathbf{F_{i-1}} + \mathbf{F_{i-2}} & \text{if } i \geq 2 \end{aligned}$$

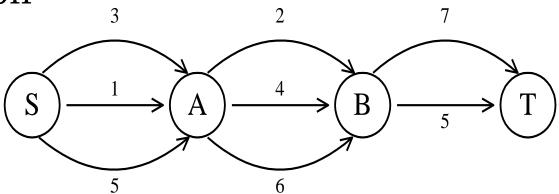
• Solved by a recursive program:



- Much replicated computation is done.
- It should be solved by a simple loop.

THE SHORTEST PATH

 To find a shortest path in a multi-stage graph

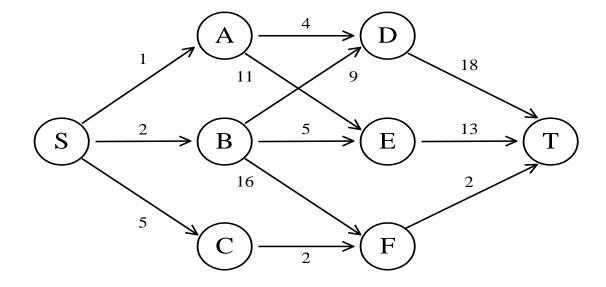


• Apply the greedy method: the shortest path from S to T:

$$1 + 2 + 5 = 8$$

THE SHORTEST PATH IN MULTISTAGE GRAPHS

oe.g.



- The greedy method can not be applied to this case: (S, A, D, T) 1+4+18=23.
- The real shortest path is:

$$(S, C, F, T)$$
 $5+2+2=9$.

Introduction

- Dynamic Programming is an algorithm design technique for *optimization problems*: often minimizing or maximizing.
- Solves problems by combining the solutions to subproblems that contain common sub-sub-problems.

- DP can be applied when the solution of a problem includes solutions to subproblems
- We need to find a recursive formula for the solution
- We can recursively solve subproblems, starting from the trivial case, and save their solutions in memory
- In the end we'll get the solution of the whole problem

STEPS

- Steps to Designing a Dynamic Programming Algorithm
- 1. Characterize optimal sub-structure
- 2. Recursively define the value of an optimal solution
- 3. Compute the value bottom up
- 4. (if needed) Construct an optimal solution

DIFF. B/W DYNAMIC PROGRAMMING AND DIVIDE & CONQUER:

- o Divide-and-conquer algorithms split a problem into separate subproblems, solve the subproblems, and combine the results for a solution to the original problem.
 - Example: Quicksort, Mergesort, Binary search
- Divide-and-conquer algorithms can be thought of as top-down algorithms

- Dynamic Programming split a problem into subproblems, some of which are common, solve the subproblems, and combine the results for a solution to the original problem.
 - Example: Matrix Chain Multiplication, Longest Common Subsequence
- Dynamic programming can be thought of as bottom-up

DIFF. B/W DYNAMIC PROGRAMMING AND DIVIDE & CONQUER (CONT...):

- In divide and conquer, subproblems are independent.
- Divide & Conquer solutions are simple as compared to Dynamic programming.
- Divide & Conquer can be used for any kind of problems.
- Only one decision sequence is ever generated

- In Dynamic Programming , subproblems are not independent.
- Dynamic programming solutions can often be quite complex and tricky.
- Dynamic programming is generally used for Optimization Problems.
- Many decision sequences may be generated.

PRINCIPLE OF OPTIMALITY

- Principle of optimality: Suppose that in solving a problem, we have to make a sequence of decisions $D_1, D_2, ..., D_n$. If this sequence is optimal, then the last k decisions, 1 < k < n must be optimal.
- e.g. the shortest path problem

 If i, i₁, i₂, ..., j is a shortest path from i to j, then i₁, i₂, ..., j must be a shortest path from i₁ to j
- o In summary, if a problem can be described by a multistage graph, then it can be solved by dynamic programming.

DYNAMIC PROGRAMMING

- Forward approach and backward approach:
 - Note that if the recurrence relations are formulated using the forward approach then the relations are solved backwards. i.e., beginning with the last decision
 - On the other hand if the relations are formulated using the backward approach, they are solved forwards.
- To solve a problem by using dynamic programming:
 - Find out the recurrence relations.
 - Represent the problem by a multistage graph.

- Multistage graph
- Computing a binomial coefficient
- Matrix-chain multiplication
- Longest Common Subsequence
- o 0/1 Knapsack
- The Traveling Salesperson Problem
- Warshall's algorithm for transitive closure
- Floyd's algorithm for all-pairs shortest paths