

UE Unix/C L3 Informatique

\$whoami



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- Teaching: ISTIC mainly L3info
- Research: CAIRN team, INRIA-IRISA
- Research Interests:
 - Computer Architecture, Embedded Systems, Mixed-critical Systems, Fault Tolerance, Low energy design, Task Scheduling and Allocation, Memory Management, Design Space Exploration



Organisation 2020-2021



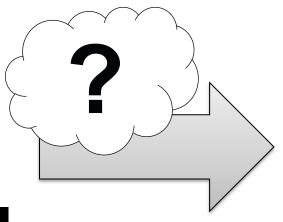
- CM: Angeliki Kritikakou
 - 8 main courses
 - Interactive course: Kahoot Tests
- TD: Angeliki Kritikakou
 - 2 exercise courses
- TP: A. Kritikakou, J. C. Engel, A. Maddi
 - Final lab exam:
 - Independent
 - ★ 03 December 2020 (Last Lab)
 - **X** Duration: 1h (slot 14:30 or 15:30)
- Note: 1/2 Theoretical exam + 1/2 Labs
 - Multiple choice questions

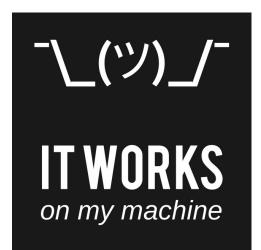


Introduction

Creation...From a text file to an executable FNINE







foo.c

```
include <stdio.h>
int main()
{
   int i = 17;
   int j = 21;
   int k = i+j;
   printf("i+j=%d\n", k);
   return 0;
}
1,1 Top
```

```
$ ./foo
i+j=38
$
```

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L5

Problem-oriented language level

L5: Problem-oriented language

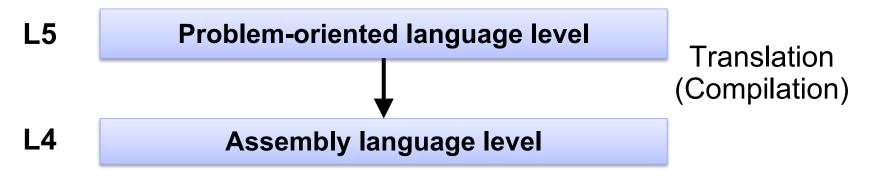


File foo.c is written with the syntax of C language

```
include <stdio.h>
int main()
{
  int i = 17;
  int j = 21;
  int k = i+j;
  printf("i+j=%d\n", k);
  return 0;
}
```

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L4: Assembly Language



1.2-11

All

foo.s:

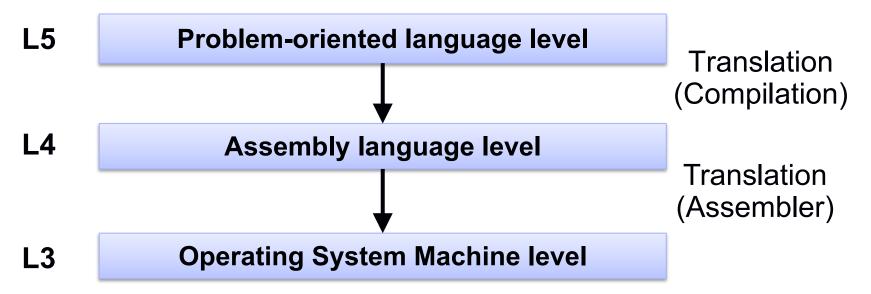
Result of compilation

\$ gcc —S foo.c

```
File Edit View Search Terminal Help
         .file
                    "foo.c"
          .section
                    .rodata
.LC0:
                    "i+j=%d\n"
          .string
          .text
          .globl
                    main
                    main, @function
          .type
nain:
.LFB0:
          .cfi startproc
          pushq
                    %г Бр
          .cfi def cfa offset 16
          .cfi_offset 6, -16
         movq
                    %rsp, %rbp
          .cfi_def_cfa_register 6
         suba
                    $16, %rsp
         movl
                    $17, -12(%rbp)
                    $21, -8(%rbp)
         movl
                    -8(%rbp), %eax
         movl
                    -12(%rbp), %edx
         movl
                    %edx, %eax
         addl
                    %eax, -4(%rbp)
         movl
         movl
                    $.LCO, %eax
                    -4(%rbp), %edx
         movl
                    %edx, %esi
         movl
                    %rax, %rdi
         movq
                    $0, %eax
         movl
         call
                    printf
         movl
                    $0. %eax
          .cfi_def_cfa 7, 8
         ret
          .cfi_endproc
.LFE0:
          .size
                    main, .-main
          .ident
                    "GCC: (Ubuntu/Linaro 4.6.3-1ubuntu5) 4.6.3"
                    .note.GNU-stack,"",@progbits
          sectron
```

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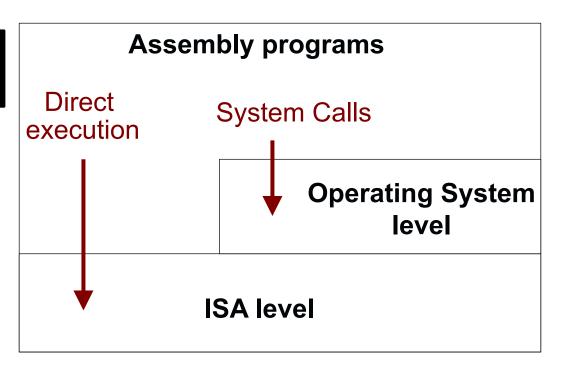


L3: Operating system



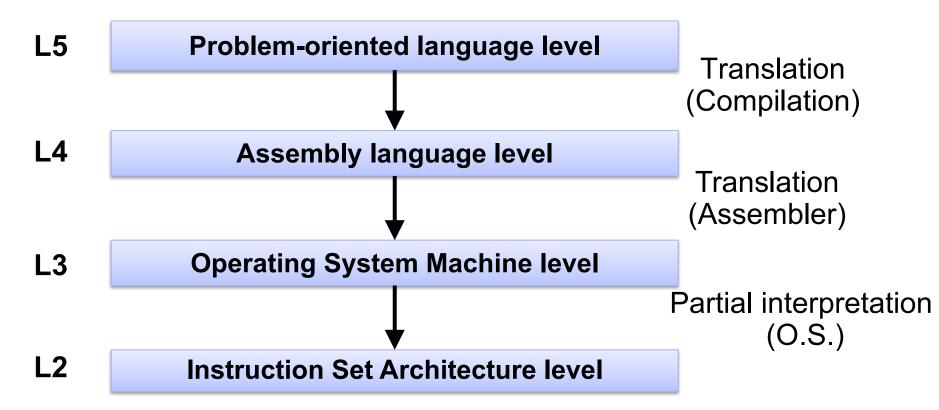
- foo: Result of assembler
- Executable:
 - Operating System
 - Instruction Set Architecture (ISA)

\$ gcc -c foo.c -o foo



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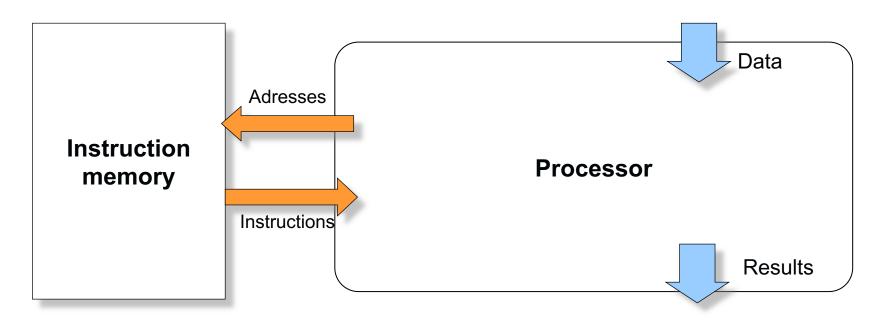


L2: Instruction Set Architecture



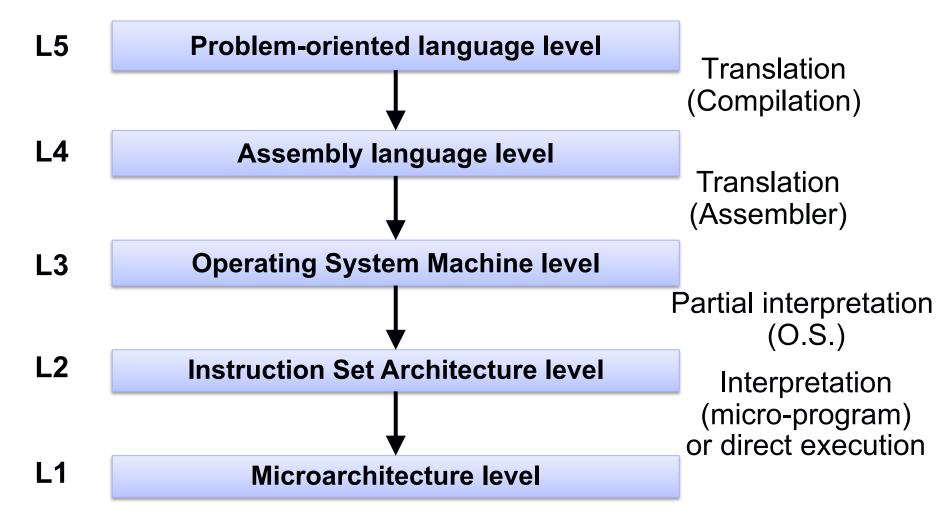
foo:

- stored in the instruction memory
- consists of several machine code instructions
- depend on the ISA of the used processor (X86, ARM, SPARC, NIOS2...)



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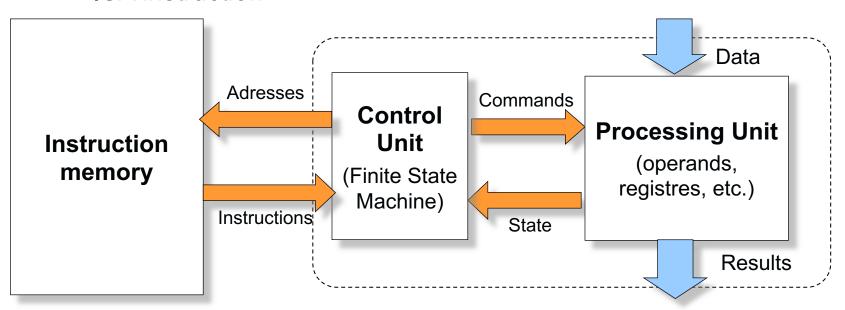




L1: Micro architecture

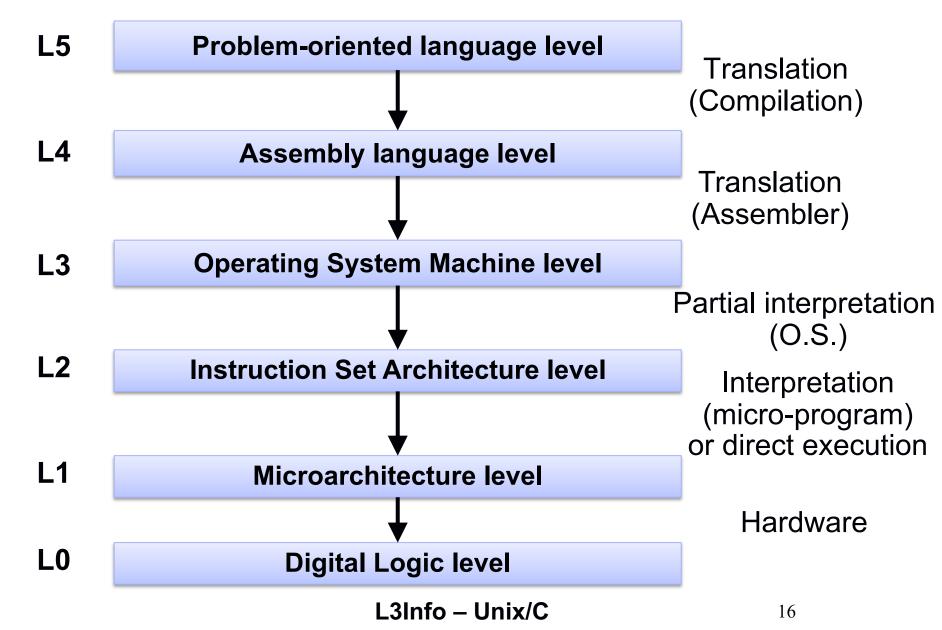


- Each ISA instruction is implemented by a set of microoperations (micro-ops)
- Micro-ops:
 - Created by the Control Unit
 - Describe how the Processing Unit has to behave to execute the ISA instruction



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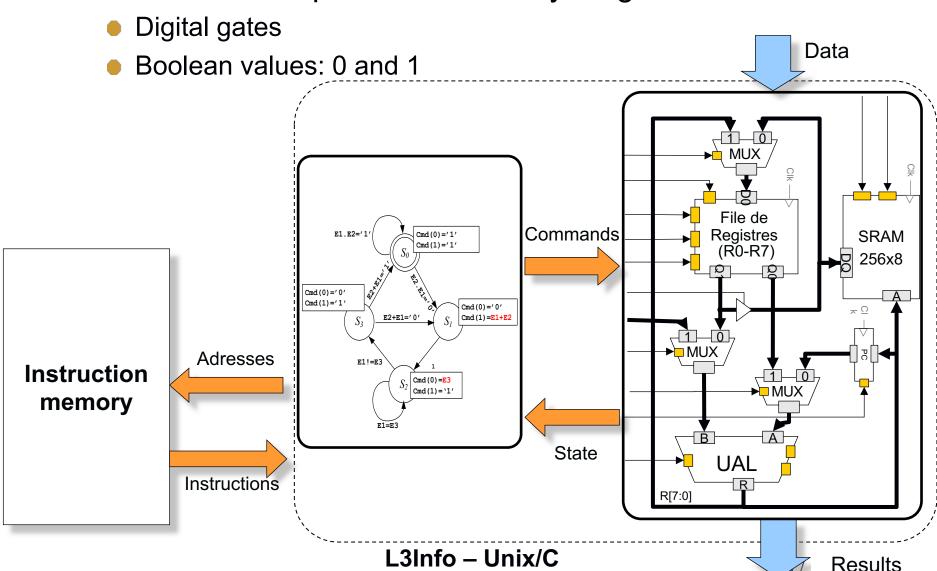




L0: Digital Logic

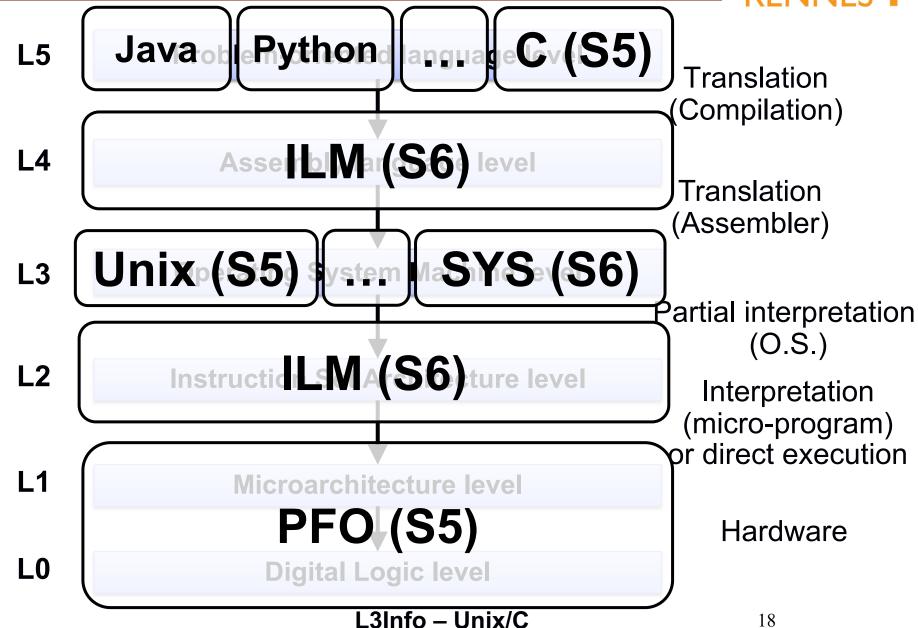


Hardware is responsible for everything!



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Outline



- Part 1: Linux and Shell scripting:
 - 2 CM, 1 TD, 1 TP

- Part 2: C Language
 - 6 CM, 1 TD, 8 TP

Kahoot time! (2 Q)

- ① Goto:
 - Insert:
- 3 Insert:

https://kahoot.it/

PIN number

a nickname



Part A: Linux and Shell scripting

Operating system



- Manages:
 - Process: creation, execution, termination, ...
 - Memory: isolation, paging, virtual memory, ...
 - Files: creation, access, permissions, ...
 - Peripherals: usage, common interface, ...
- Provides a higher layer abstraction to interact with hardware peripherals
 - System calls
- Focus: Utilisation of the Operating System
 - Linux

Linux



- First prototype in 1991 (post to "comp.os.minix"):
 - Linus Torvalds (Finland student age 21):

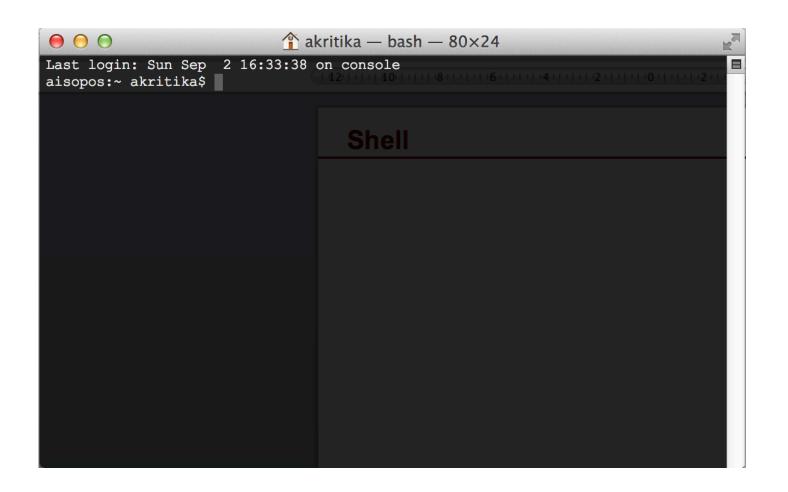
"Hello everybody out there using minix –

I am doing a (free) operating system (just a hobby, won't be big and professional like gnu) for 386 (486) AT clones..."

- Naming:
 - Linus: "Freax" → "Free", "Freak", "x" (from Unix)
 - Ari Lemmke (server admin): Renamed to "Linux"
- Free and open source operating system
 - > 18.000.000 lines of source code

Interface with the system: The Shell





Shell



- Command Line Interface (CLI)
 - Input: The commands you want to execute (use the keyboard)
 - Output: Shell passes the commands to the OS, which executes them, and provides the result (usually to the screen)
 - Scripts: Combination of commands
 - Variables
- Interaction with the Shell: Terminal emulators
 - Konsole, Xterm, Gnome-terminal...
- Several shells are available
 - sh (Bourne shell), bash (Bourne shell), csh, zsh, ksh ...
- Each user has a preselected shell
 - Selection in the file: /etc/passwd
 - Command to change the shell: chsh
 L3Info Unix/C

Directories



- Home Directory (\$HOME)
 - Each user has a home directory connected with his name
 - Name: 2363415
 - Home directory: /private/student/5/25/2363415 or ~/2363415
- Working Directory (\$PWD)
 - Current directory
 - When connected: Working directory = Home directory
- Current directory: .
- Parent directory: ...
- Root directory: /
- User home directory: ~

Paths



- Absolute
 - From root directory /:/private/student/5/25/2363415
- Relative
 - From working directory:
 - Assume the following working directory: /private/student/5/25/2363415
 - You can access a file in another directory /private/student/5/25/2363414 by:

../2363414

Wildcard characters



 Special characters which can be used in combination with commands

- One and only one character: ?
 - Ex: ls /etc/rc.????
- From 0 to many characters: *
 - Ex: ls /etc/rc.*
- Specific characters : [...]
 - Ex: 1s [abc]oo.c
 aoo.c, boo.c, coo.c

How to treat them as normal (by the Shell) VERSITÉ DE SHENNES

- Preserve the literal value of next character: \
- Preserve the literal value of a character sequence:
 '...' (ticks)
- As before, except for the characters \$ and `: "..." (quotes)
- Command replacement: Execution of the command and replace with what it has been produced as output
 `command` (back ticks)
 \$ (command)

Useful information



- tab key:
 - One tab key: automatically fills the word you are started typing, if exists
 - Two tab keys: Prints all the existing words in that start with the letters that you have typed
- clear command: Cleans the screen in the shell
- Up arrow key: Shows the previous command
- history command: Show all previous commands
- !!: Executes the previous command
- Do not use special characters in the filenames, that include also the space character and the hyphen character (-).
- Put as default language ENGLISH, the French translation in not complete ... (LANG=EN)

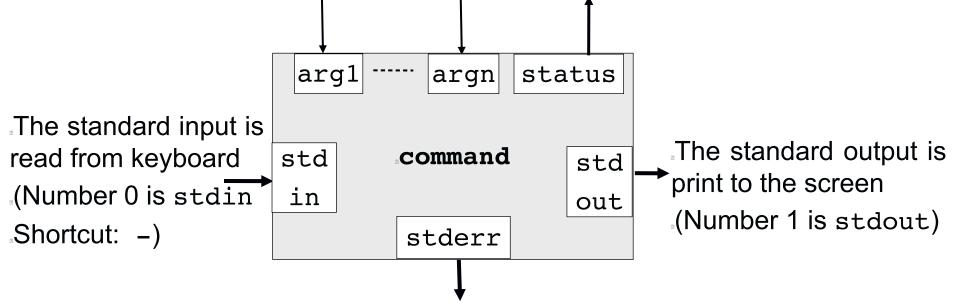
Structure of a command



command arg1 arg2 ... argn

The arguments of a command are the parameters passed during its call

The exit status is the value returned after the execution of the command to inform about normal or not execution



The standard error is by default printed to the screen and it is used by the commands to print the error messages (Number 2 is stderr)

L3Info – Unix/C

Shell Commands



- Directory management
- File management
- Redirections
- Process management
- Shell programming: Scripts
- This course does NOT give an exhaustive list of commands and a complete command description → Use:
 - documentation
 - man or info command
 - apropos command

Directory management



- Print the full pathname of the working directory: pwd
- Change directories: cd
 - cd or cd ~: Go to home directory
 - cd path: Go to directory indicated by path
- Creation/Delete directories
 - mkdir dirname: Creates a directory named dirname
 - rmdir dirname: Deletes the directory named dirname (empty)
 - rm -r dirname: Deletes the directory named dirname recursively (subdirectories) (not empty)
 - NO TRASHBIN!
- Rename: mv dirsource dirdest
- Hierarchical copy (subdirectories): cp -r dirsource dirdest
 L3Info - Unix/C
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Listing



- 1s: Lists the files in the working directory
- ls dirname: Lists the content of directory dirname

```
angie@achilleas:~/Entertainment$ ls
movies Music photos
```

- Some parameters passed as arguments:
 - -1: Use a long listing format (details)

```
angie@achilleas:~/Entertainment$ ls -l
total 12
drwxrwxr-x 8 angie angie 4096 Φεβ 1 2015 movies
drwxr-xr-x 3 angie angie 4096 Φεβ 1 2015 Music
drwxrwxr-x 5 angie angie 4096 Οκτ 20 2015 photos
```

-a: Not ignore entries starting with . (hidden folders/files)

```
angie@achilleas:~/Entertainment$ ls -al total 28 drwxrwxr-x 5 angie angie 4096 Δεκ 4 2014 . drwxr-xr-x 48 angie angie 12288 Μάι 30 08:53 .. drwxrwxr-x 8 angie angie 4096 Φεβ 1 2015 movies drwxr-xr-x 3 angie angie 4096 Φεβ 1 2015 Music drwxrwxr-x 5 angie angie 4096 Οκτ 20 2015 photos
```

File management



- Create file by:
 - Update access & modification time: touch filename
- Delete file: rm filename
- Rename: mv filesource filedest
- Copy:
 - cp filesource filedest
 - cp filesource dirdest
- Remove non-directory suffix part from a filename: dirname
- Remove the directory part suffix from a filename: basename
- Count the number of words of a file: wc —w filename

File management



- Print the contents of a file:
 - more filename: Prints page by page the content of file filename on the stdout (Q: Exit from more)
 - less filename: Similar to more, but it allows to go up/down to the pages (Q: Exit from less)
 - head —n 42 filename: Print the first 42 lines of filename
 - tail —n 42 filename: Print the last 42 lines of filename
- Concatenate files (create and view files):
 - ocat filename1 ... filenameN (print context of files at stdout)
 - cat: Print at stdout whatever is tapped at stdin (cntrl+D indicates EOF)
 - : Uses what is tapped with the keyboard instead of an existing file
- Difference between two files:
 - diff filename1 filename2

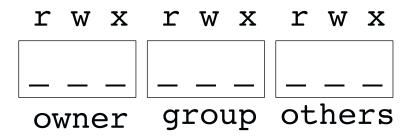
Directory/File management: Examples UNIVERSITÉ DE RENNES



Kahoot time! (4 Q)

Directory/File permissions





- Change permissions: chmod chmod [who] op permissions [,op permissions]* name
- Who:
 - Owner (u)
 - Group (g)
 - Others (o)
 - Nothing (all)

- Op:
 - Add (+)
 - Remove (-)
 - Define (=)

- Permissions:
 - Read or List (r)
 - Write (w)
 - Execute or Access(x)

Directory/File permissions: Examples UNIVERSITÉ DE RENNES



Kahoot time! (3 Q)

File management: Search files



Search files based on criteria (print path):

find directory [criteria]

- The search is performed to directory (recursive to subdirectories)
- Apply a command at the found files: option —exec cmd {} \;
 - Apply the command cmd in each found file (it replaces the {})
 - ★ End of the command: \;
- Some search criteria:
 - -name filename: based on the filename (-iname: ignore case)
 - -user username: based on the owner
 - -size number[c|k|M]: based on size [B, KB, MB]
 - -type [d|f|1]: based on the type [directory, file, link]
 - -mtime x: based on time of modification (less than x days)
 -
- Criteria can be combined: -o -a -not
 L3Info Unix/C

File management: Search files



Kahoot time! (2 Q)

File management: Search in files



- Print all lines of the files of the directory that contain a word:
 grep string directory
 - The string can be expressed as a regular expression
 - Whatever character : ..
 - Repetition of pattern: *
 - ★ Interval of characters: [x-y]
 - Return value: 0 not found, else !=0
- Examples:
 - grep -r test ./
 - grep "m.*t" filename
 - grep "toto[1-3]" filename

Redirections



- Redirection of standard input (stdin):
 - Use the content of a file as input to a command (instead of keyboard)
 - x command < filename</pre>
- Redirection of standard output (stdout)
 - Print the output of a command into a file (instead of screen)
 - x command > filename (overwrite old content of file)
 - x command >> filename (append to the content of file)
- Pipe: redirect standard output to standard input of commands
 - Use the output of a command as input to another command
 - A subshell is created to execute command_2
 - x command 1 | command 2

Redirections: Examples



command



command > res.txt



command < data.txt</pre>



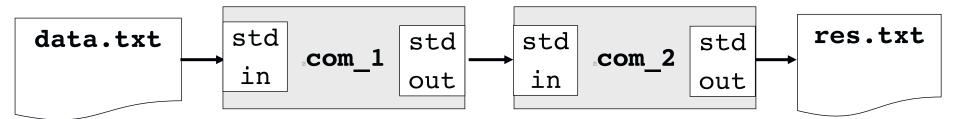
Redirections: Examples



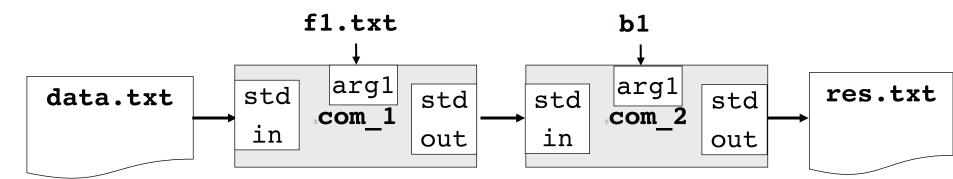
command < data.txt > res.txt



com_1 < data.txt | com_2 > res.txt



com_1 f1.txt < data.txt | com_2 b1 > res.txt



Redirections: Examples



Kahoot time! (3 Q)

Command implementation



Builtin

- Implemented in the code of the shell
- They affect the internal state of the shell (ex. cd)

External

- Implemented as a separate code
- Shell interprets the command as a request to load and run the program that implements the command
- Stored in /bin

Builtin Commands



Command	Description	Example
cd	Change directory	cd
declare	Define variable	declare myvar
echo	Print to standard output (stdout)	echo hello
exec	Replace bash with another process	exec 1s
exit	Quit bash	exit
export	Define global variable	export myvar=1
history	Print the history of commands	history
kill	Send signal to a process	kill 1121
let	Compute arithmetic operation	let myvar=3+5

Builtin Commands



Command	Description	Example
local	Define local variable	local myvar=5
pwd	Print working directory	pwd
read	Read from standard input (stdin)	read myvar
readonly	Locks a variable	readonly myvar
return	End of a function and return a value	return 1
set	Print variables	set
shift	Move parameters	shift 2
test	Test a condition	test —d temp
trap	Follow a signal	trap "echo signal" 3

Process management: Execution



- Process: The running instance of a program, e.g.:
 - Execution of a command
 - Execution of your program
- Foreground execution: Interactive processes
 - Connected to a shell terminal
 - User can send input
- Background execution: Non-Interactive or automated processes
 - Not connected to a terminal
 - User cannot send input until the process is moved to the foreground

Process management: Changing grounds RENNES

- Run in Background: &:
 - Ex: gedit &
- From FG to BG:
 - Ex: geditPress CTRL + Z (puts process in idle state)bg
- From BG to FG:
 - Ex: fg %jobid
 (jobid is not PID. It is a unique number given by the shell, e.g. when process goes to BG)

Process management: Active processes UNIVERSITÉ DE RENNES

- Print the catalogue of the running processes: ps
 - -a: Select all processes associated with a terminal.
 - **PID:** Unique identity of the process
 - **TTY:** Console created the process
 - TIME: Total execution time (sec)
 - CMD: Command executed

```
angie@achilleas:~/Downloads$ ps
PID TTY TIME CMD
12190 pts/1 00:00:00 bash
14414 pts/1 00:00:00 vi
14447 pts/1 00:00:00 evince
14454 pts/1 00:00:00 ps
```

Process management: Control execution RENNES

- End a running process
 - Global method:
 - x kill PID
 - Foreground execution in current shell (through SIGNALS)
 - x CTRL + c (Interrupt)
 - x CTRL + D (EOF)
 - ★ CTRL + \ (Quit)
- Pause a process
 - CTRL + z
- Resume
 - Foreground (fg)
 - Background (bg)

Shell Variables

- RENNES 1
- Initial values: defined in system and user configuration file
- Use uppercase by convention
- Definition of variables based on the type (bash shell)
 - String variables: MYVAR="value" (NO SPACES around =)
 - Integer variables: declare —i MYVAR
 - Constant variables: readonly DATA=value
 - Array variables: declare —a MYARRAY MYARRAY[0]="one"; MYARRAY[1]=5
- Use of variables: \$
 - E.g. \$MYVAR
- Lifetime of variables:
 - End of session termination of shell
 - Deleted by user: unset VARNAME

Shell Variables: Scope



- Global: Public variable seen by all subshells
 - Set export attitude for the variable: export
 - Print global variables: env, printenv
- Local variables: Seen by current shell
 - Print list of local and global variables: set

Shell Variables: Scope example



```
angie@achilleas:~$ myvar="hello"
angie@achilleas:~$ set | grep myvar
myvar=hello
angie@achilleas:~$ echo "$$"
12190
angie@achilleas:~$ bash
angie@achilleas:~$ echo "$$"
15752
angie@achilleas:~$ set | grep myvar
angie@achilleas:~$ env | grep myvar
angie@achilleas:~$ exit
exit
angie@achilleas:~$ export myvar="hello"
angie@achilleas:~$ bash
angie@achilleas:~$ echo "$$"
15825
angie@achilleas:~$ set | grep myvar
myvar=hello
angie@achilleas:~$
```

Predefined shell variables



Variable	Description
USER	Name of user account
HOME	Personal folder of user
TERM	Type of terminal
SHELL	Name of shell
PATH	List of directories with executable commands
MANPATH	List of directories with help information for the commands
PWD	Current directory
OLDPWD	Previous current directory
HOSTNAME	Name of system

Scripts: Structure



- Program consisting of commands:
 - Use an editor:

```
x vi (vim), emacs, sublime, gedit, ...
```

- 1st line defines the shell to be used:
 - #!/bin/bash (for bash shell)
- Write program commands
- Commands can be separated by:
 - x a new line
 - Semicolon: ;
- Split a command to 2 or more lines: \ (backslash)
- For comments: #
- Save file with extension .sh
- The script lines (commands) are interpreted and NOT compiled

Scripts: Execution



No need for permissions:

bash path_to_script/myscript.sh

- bash command is executed
- ./script.sh is the command-line argument.
- Provide eXecute permissions and execute:

chmod +x myscript.sh

- ./myscript.sh
 - Shell forks itself and uses a system call (e.g. execve) to make the OS execute the file in the forked process.
 - OS checks the file's permissions
 - Program loader looks at the file to find how to execute it
 (#! for scripts and then which interpreter to use: /bin/bash).
- New process: attention to the variables and functions

Command-line arguments



- Arguments can be passed along with the script execution:
 - ./myscript.sh arg1 ... argN
- The arguments are seen as specific variables in the script code:
 - \$0: Name of the script
 - \$1: 1st argument
 - \$N: Nth argument
 - \$#: Number of arguments
 - \$*: All the arguments passed to the script, as a string
 - \$@: Behaves like \$* except that when the arguments are quoted, they are broken up properly if there are spaces in them
- Maximum length of arguments is defined by operating system and measured in Kilobytes.
 - getconf ARG_MAX

Variables



- Initialization: nomvar=expresion
 - ATTENTION! No space around =
 - expression:
 - × Value
 - **X** Reference to a variable
 - * `command` (replaced with the result to the stdout)
- Use: \$nomvar
 - Print current directory

```
y=`pwd`
echo $y
```

Print the first two arguments of a script

Read input



- We can ask the user to provide some value(s) to the variables
 - read var-name1 var-name2 ... var-nameN
 - Reads a line from the stin and affects the first value to varname1, the second value to var-name2, etc.
 - Returns 0, if the read was successful
- We can print a message to the user:
 - read -p "Enter value:" var
- If variable is not specified, then the default is used:
 - \$REPLY

Print to output



- Basic print to stdout: echo
 - Strings
 - Variables
 - Wildcards characters
 - Can have several lines

```
angie@achilleas:~/Entertainment$ echo hello
hello
angie@achilleas:~/Entertainment$ echo *
movies Music photos
angie@achilleas:~/Entertainment$ echo print '@' "user"
print @ user
angie@achilleas:~/Entertainment$ echo 'hell
> o there'
hell
o there
```

Arithmetic operations



- No need to define the variables as integers
- Arithmetic operations with integers
 - \$((...)): Portability among shells
 - ((...))

- let (without spaces as it is an initialisation)
- expr (with spaces in arguments)

Condition expressions: test



- A condition expression is:
 - a set of commands that are executed
- The test command exits with the exit status determined by the condition_expression

```
test condition_expression
```

```
[ condition_expression ]
```

[[condition_expression]]

Similar to single brackets, but more powerful

```
(( arithmetic condition_expression ))
( command )
```

```
L3Info – Unix/C
```

SPACES ARE

REQUIRED !!

Exit status



Syntax
\$?

ATTENTION!!

- Function
 - If the command ends correctly (TRUE): exits with value 0 (exit 0)
 - If there was an error (FALSE): exits with value 1 (exit 1)

```
$ true ; echo $?
$ 0
$ false ; echo $?
$ 1
```

- Commands that do nothing except return an exit status of:
 - true: Return an exit status of zero
 - false: Return an exit status of one

Arithmetic operations



Kahoot time! (4 Q)

Condition expressions: test operators UNIVERSITÉ REN

Operator	Description
-gt	Greater than
-ge	Greater than or equal
-lt	Less than
-le	Less than or equal
-eq or =	Equal (-eq for integers and = for strings)
-ne or !=	Not equal (-ne for integers and != for strings)
-n str	Non-zero size string
-z str	Zero size string
-d dirname	dirname is a directory
-s filename	filename size is not zero
-f filename	filename exists
-r filename	filename exists and we have read access
-w filename	filename exists and we have write access
-x filename	filename exists and we have execute access

Boolean control operators



- Syntax (AND): -a or &&
- command1 && command2
- Function:
 - command2 is executed if, and only if, command1 returns an exit status of zero.
- Syntax (OR): -o or | | command2
- Function:
 - command2 is executed if, and only if, command1 returns a nonzero exit status.
- The exit status of AND and OR lists is the exit status of the last command executed in the list
- Syntax (NOT): !

Conditions



Kahoot time! (4 Q)

AND and OR control operators: Example RENNES

```
$ true | echo "echo executed"
$ false | echo "echo executed"
echo executed
$ true && echo "echo executed"
echo executed
$ false && echo "echo executed"
$
```

Conditions: if, elif, else



Syntax:

```
if [ condition_expression_1 ] ; then
    1<sup>st</sup> set of commands
elif [ condition_expression_2 ]
then
    2<sup>nd</sup> set of commands
else
    3<sup>rd</sup> set of commands
fi
```

ALL 3 SPACES ARE REQUIRED!!

- Function
 - The if checks if the condition_expression_1 returns TRUE
 x 1st set of commands is executed
 - Else, if the condition_expression_2 returns TRUE
 x 2nd set of commands is executed
 - Otherwise, 3rd set of commands is executed
 L3Info Unix/C

Conditions: if: Examples



```
#!/bin/bash
read -p "Enter a filename: " filename
if [ ! -w "$filename" ]; then
  echo "File $filename is not writable"
  exit 1
elif [! -r "$filename"]; then
  echo "File $filename is not readable"
  exit 1
else
  echo "File $filename is writable and readable"
fi
```

Conditions: if: Examples



```
#!/bin/bash
TMPFILE="diff.out"
diff $1 $2 > $TMPFILE
if [ ! -s "$TMPFILE" ]; then
        echo "Files are the same"
else
        more $TMPFILE
fi
if [ -f "$TMPFILE" ]; then
        rm $TMPFILE
fi
```

Conditions: case



Syntax:

Function

- If the value of **variable** is **pattern 1**, the 1st set of commands is executed, if it is **pattern 2**, the 2nd set of commands is executed etc.
- Otherwise (*) → Nth set of commands

Conditions: case: Example



```
#!/bin/bash
read -p "Enter a command: " command
case $command in
all | ALL )
        echo "Display all files..."
        ls -la ;;
list | LIST)
        echo "Display all non-hidden files ... "
        ls -11 ;;
* )
        echo "Invalid choice"
esac
```

Iterations: for



Syntax:

```
for variable in list_of_values
do
   set of commands
done
```

Function

- In each iteration variable takes a value from the list_of_values
 * The set of commands is execute for this value
- list of values can be:

File names:

x command output: \$(ls /tmp/*)

L3Info - Unix/C

filename1 filename2

Iterations: for: Examples



```
#!/bin/bash
for file in f1 f2 f3
do
    cp $file $1/$file
done
```

```
#!/bin/bash
for i in 6 3 1 2
do
echo $i
done | sort -n
```

Iterations: while



Syntax:

```
while
    [ condition_expression ]
do
    set of commands
done
```

- Function
 - As long as the condition_expression returns TRUE
 The set of commands is executed
 - Otherwise, the iteration ends

Iterations: while: Examples



Iterations: until



Syntax:

```
until
   [ condition_expression ]
do
   set of commands
done
```

- Function
 - Execute 1st set of commands
 - As long as the condition_expression commands is NOT true (while consition_expression is false)
 - The set_of_commands is executed
 - Otherwise the iteration ends

Iterations: until: Examples



Controlling loop commands



- Stop with current loop iteration: break
- Continue with the next iteration: continue
- Terminate script: exit

Controlling loop commands: Examples RENNES

```
#!/bin/bash
for x in 1 2 3 4 5 6 7 8 9
do
      echo -n "x is $x. "
      if [ "$x" -eq "4" ]; then
            echo
            continue
      fi
      echo "Most numbers get here, except 4."
      if [ "$x" -eq "7" ]; then
            break
      fi
done
echo "We have finished the loop now."
echo "At the end, x is now $x."
```

Controlling loop commands: Example RENNES

```
x is 1. Most numbers get here, except 4.
x is 2. Most numbers get here, except 4.
x is 3. Most numbers get here, except 4.
x is 4.
x is 5. Most numbers get here, except 4.
x is 6. Most numbers get here, except 4.
x is 7. Most numbers get here, except 4.
We have finished the loop now.
At the end, x is now 7.
```

Functions



Syntax

```
function name[()]
{
   list of commands;
[return]}
```

Function

- Functions have to be defined in the beginning of the script
- They may have or not arguments
- Arguments and return values can be of any type
- The variables defined inside the function are GLOBAL. They have to be defined LOCAL, if we want to reduce their scope

Functions: Example



```
#!/bin/bash
Outside="a global variable"
function mine() {
        local inside="this is local"
        echo $outside
        echo $inside
        outside="a global with new value"
}
echo $outside
mine
echo $outside
echo $inside
```

```
a global variable
a global variable
this is local
a global with new value
```



TD1: Linux and Shell Exercises

Ex1: Wildcard characters

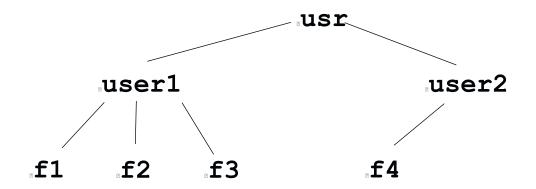


- List the entries inside the folder /usr/bin whose name starts with the letter *m*.
- List the entries inside the folder /usr/bin whose name starts with the letter *m* and it has exactly 3 characters
- List the entries inside the folder /usr/bin whose name starts with the letter *m* and it has at least 3 characters
- List the entries inside the folder /usr/bin whose name starts with the letter *m* and it has an extension (postfix after .)

Ex2: Redirections



Given the following file hierarchy

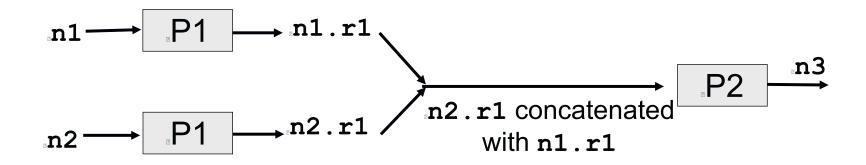


- user2 is connected at his home folder
- Use the command cat:
 - Create a file **f5** with the following content5 2 3 -1
 - Create a file f6 by concatenating f5 after f4
 - Add the file f1 of user1 at the file f6
 - Copy the file f2 of user1 as a new file f7 of user2

Ex3: Redirections



Create script with the name execution called with 3 arguments n1 n2 n3 (which are the filenames) performing the following manipulations



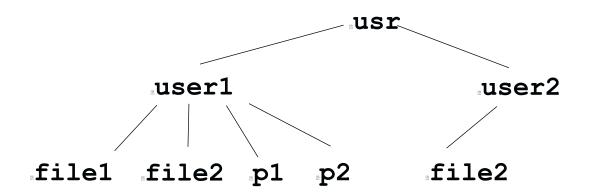
■ P1 and P2 are programs that read from the stdin and print their result to stdout

Ex3: Redirections



What happens when the following instruction is executed in the home directory of user1?

execution file1 ../user2/file2 fileN



Write an interactive version of the script execution that asks from the user to provide n1, n2 and n3, instead of passing them as arguments

Ex4: Script arguments



Create a script with the name argument using the command expr that prints one by one the arguments passed during the script call.

For instance, the call:

% argument one two three

Should print in the stdout:

```
3 arguments :
   argument 1 - one
   argument 2 - two
   argument 3 - three
```

Ex5: Search



- Create a script with the name mysearch passing two arguments during the call, i.e. motif and filename. The program prints either "filename has motif" or "filename does not have motif"
- Create a script mysearch2 filename that ask the user the motif to be searched inside the filename until the motif end is entered by the user
- Create a script mysearch3 f1 .. fn that performs the same thing as mysearch2, but now it searches inside the files given as arguments.
- /dev/null = trash

Ex6: Find



Describe the effect of the following instructions:

```
find ./ -name "tp2" -exec {} < f_in > {}.out \;
find /var/log -exec grep "www.athabasca.com" {} \;
find . -name "rc.conf" -exec chmod o+r {} \;
find /usr/local -type f -name "*.html".
```



Part B: C Language

Outline



- Introduction
- Structure of C program and compilation
- Basic components of C language
- Pointers and Arrays
- Functions and parameters
- Dynamic memory allocation
- Structures, Lists and Hash tables
- Strings
- I/O



CM 3: Introduction, Structure and Tools

Why C?



- JAVA is an evolved language
 - High abstraction away from the physical objects
 - Useful when we want to be abstracted from the hardware
 - Software engineering
- C is closer the hardware
 - Relatively controls what happens to the memory
 - Manipulation of the physical objects of the hardware
 - System engineering
- Main operating systems are written in C
 - Unix, GNU/Linux, etc.

What this course covers?



- It is NOT a yet another programming course with extended presentation of the language syntax,
 - Differences with JAVA, similar syntaxes but still very different!
 - Mainly "tricky" parts of C language!
- Being a master of C force is long and difficult
 - Requires practise, practise and practise
- But, there is a lot of documentation!
 - The functions of libraries are documented
 - x Example:man fprintf
 - Google is your best friend...not to copy paste, but read and understand!

C language



- Developed by <u>Thompson</u> & <u>Ritchie</u> in 1972
- Originally developed: 90-95% of the Unix kernel



- It has been standardized by the American National Standards Institute (ANSI) in 1989 (C ANSI)
- Loosely/weakly typed language (so, warnings are important !)



Structure of C program and compilation

C program's structure



- A program is saved into one or multiple compilation C files
- One compilation file consist of a source file or a header file
 - Source file .c → code
 - Header file .h → declarations

Communication among .c files



- code2.c can use the functions and the global variables defined:
 - locally in this compilation file code2.c
 - in another compilation file code1.c (code1.c exports them)
- code1.c defines the functions and the global variables:
 - exported by code1.c to be used by another file code2.c
 - used locally in code1.c

code2.c

code1.c

Definition of variable b;

Use of variable b;

Use of variable a;

Definition of variable **a**;
Use of variable **a**;

Communication among .c files



- The definition of the global variables and constants is by default local:
- To be "seen" by file code2.c → the global variable has to be exported:
 - declarations are placed in the header file code1.h
 x extern: For the global variable to be seen by other .c files
 - definitions are in the source file code1.c
 - code2.c includes the header code1.h of code1.c through the directive #include

Directive #include



Syntax:

```
#include <filename.h> (1)
#include "filename.h" (2)
```

Description

 (1) and (2): Include the file called filename which is stored in the folder

```
/usr/include
```

- Only (2): Search the file called filename in the working directory
 - x if the file is saved in another directory you have to give the path "path/filename"

Functionality:

• Includes (before compilation) the content of filename. This text is treated as it was part of the current file .c

Communication among .c files: Example RENNES

```
#include "code1.h" /* Declaration made available here */ code1.c
/* Variable defined and initialized here */
int global_variable = 1; /* Definition checked against declaration */
int increment(void) { return global_variable++; }
```

```
extern int global_variable; /* Declaration of the variable */ code1.h
```

```
#include "code1.h"
#include <stdio.h>

void use_it(void)
{
    printf("Global variable: %d\n", global_variable);
}
```

Communication among .c files: Example RENNES

```
extern void use_it(void);
extern int increment(void);
```

```
prog.c
#include "code1.h"
#include "prog.h"
#include <stdio.h>
int main(void)
    use it();
    global variable += 10;
    use it();
    increment();
    use it();
    return 0;
```

A. Kritikakou

Communication among .c files: Example RENNES

```
#include "code1.h" /* Declaration made available here */ code1.c
#include "prog.h"

/* Variable defined here */
int global_variable = 1; /* Definition checked against declaration */
int increment(void) { return global_variable++; }
```

```
extern int global_variable; /* Declaration of the variable */ code1.h
```

```
#include "code1.h"
#include "prog.h"
#include <stdio.h>

void use_it(void)
{
    printf("Global variable: %d\n", global_variable++);
}
```

General rules



- For any given global variable
 - exactly 1 header file declares it
 - exactly 1 source file defines it and preferably initializes it too.
 (otherwise undefined behavior → anything could happen!)
 - The source file that defines the variable also includes the header (ensuring definition and declaration are consistent)
- A source file
 - never contains extern declarations of variables
 - always include the (sole) header that declares them
- A header file only contains extern declarations of variables
 - never static
 (each source file makes its own version of the global variable)

Structure of .c file



- Description of the functionality and the use (comments ②)
- The directives for the pre-processor
 - Include files, Macros, Conditional compilation, ...
- Global variables used only by this file
 - if not initialized, they are by default 0
- The functions that act on the variables
 - Function prototype
 - Function body
- No order is imposed to the definitions, BUT any element has to be declared BEFORE its utilisation

Structure of .h file



- Declaration of whatever is allowed to be "seen" outside file .c
 - Functions
 - Global variables
 - Data structures
 - ...
- It is the source file .c interface to outside world

Structure of .c file: Example

```
int ppcm(int x, int y)
int a=x;
 * main.c
                                                          int b=y;
                                                          while (a!=b)
 while (a>b) b=b+y;
                                  Pre-processor directives
#include <stdlib.h>
                                                             while (a<b) a=a+x:
                                  - Include files
#include <stdio.h>
                                                           return a;
                                  - Macros
#define PI 3.14
int variableGlobale = 3;
                                   Global variable defined in this file
extern int ppcm (int a, int b);
                                  Function declaration: is defined in another file
void printINT(int val)
                                  Function defined in this file
    printf("%d\n",val);
                                   (can be exported)
int main(int argc, char *argv[])
                                                           The main function is
    int variableLocale = ppcm(2,variableGlobale);
                                                           the starting point
    printINT(variableLocale);
    return EXIT_SUCCESS;
                                                           Only 1 main function!
                                                                    18
```

* DDCM.C

Same example split in 3 files

```
UNIVERSITÉ DE RENIVES
```

```
* main.c
#include <stdlib.h>
#include <stdio.h>
#include "ppcm.h"
#define PI 3.14
int variableGlobale = 3:
//extern int ppcm (int a, int b);
void printINT(int val)
   printf("%d\n",val);
int main(int argc, char *argv[])
   int variableLocale = ppcm(2,variableGlobale);
   printINT(variableLocale);
   return EXIT_SUCCESS;
```

```
/****************
* ppcm.h
*******************
#ifndef PPCM_H
#define PPCM_H
int ppcm(int x, int y);
#endif
```

```
/*********************************

* ppcm.c

******************************

int ppcm(int x, int y)

{
    int a=x;
    int b=y;
    while (a!=b)
    {
        while (a>b) b=b+y;
        while (a<b) a=a+x;
    }
    return a;
}
```

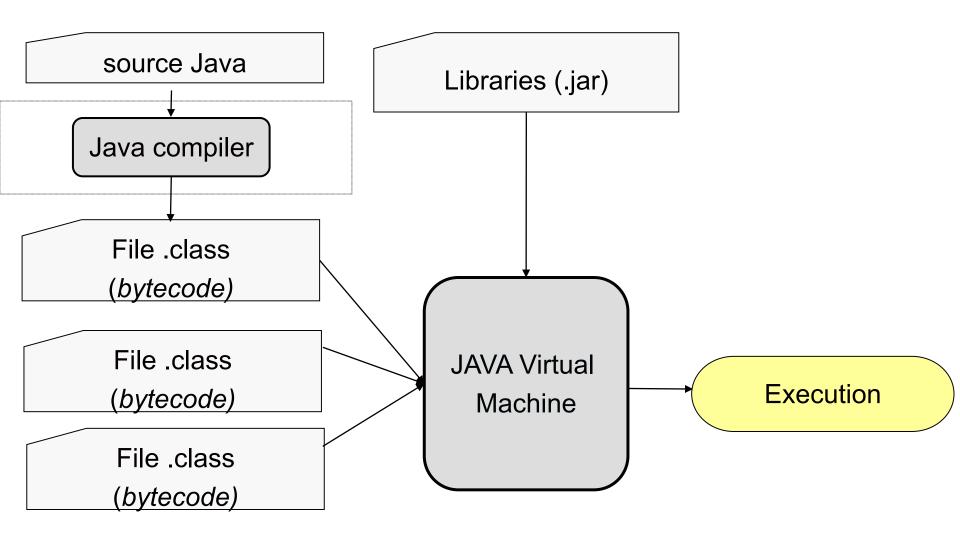
JAVA compilation



- The .java files have the code that will be translated into the files .class by the compiler javac
 - The program is split into classes
- The files.class consist of the bytecode java: a lower level language closer to the machine language (but not a machine language)
- The bytecode can be executed in whatever machine that has a Java Virtual Machine. The JVM interprets the bytecode to machine language (portable, but slow)
 - The link among classes is performed during the execution (Dynamic link)

JAVA compilation





A. Kritikakou

L3Info - C/Unix

C compilation



- The call of the C compiler is associated with 3 tools
- 1st step: The pre-processor (CPP):
 - Handles the directives (#include, #define etc.)
- 2nd step: C compiler and assembler (GCC):
 - C compiler performs one-to-one translation of a compilation file (.c and .h) into an assembly file (.s).
 - ★ This step is hidden by default (use gcc —S to see it)
 - Then, the assembler generates an object file (.o) with
 - * the machine instructions of the functions defined in this .c file
 - information on the functions used in this .c file, but defined in another file
- 3rd step: The linker (LD/GCC):
 - Makes the link (resolves symbols) among the different program's objects (static linking) to one executable binary

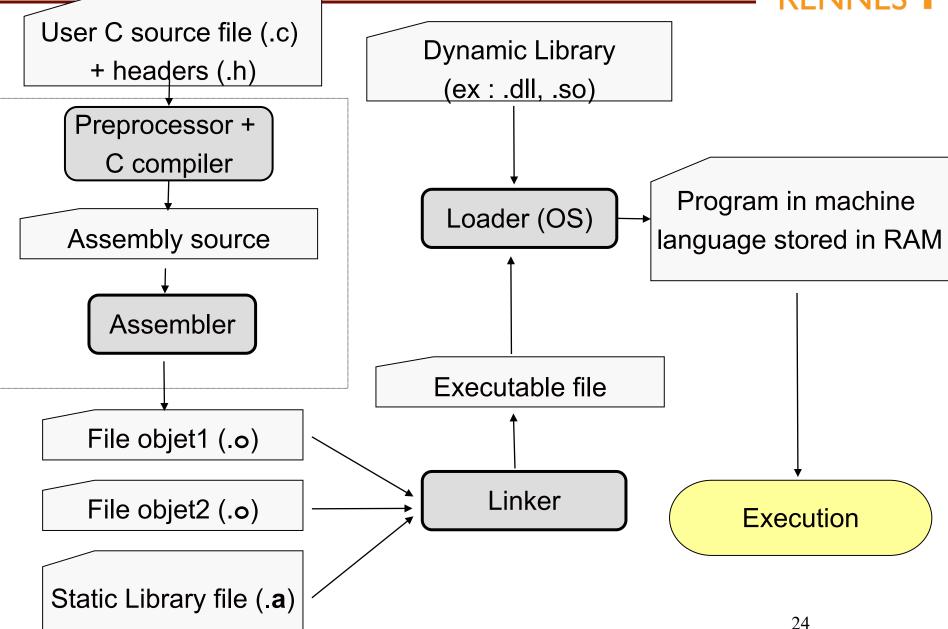
C compilation



- The generated binary depends on
 - Architecture/processorAND
 - operating system
 - → It is not portable!
- Possible to use "cross compilation"
 - Use a machine to create an binary for another machine
 - -march
 - x native
 - × i386

C compilation





C compilation: GCC



- GNU (GNU is Not Unix!) Compiler Collection is free software:
 - Design to target several machines
 - Exist by default in Linux
- Syntax (see man gcc):

```
gcc [-c] [-g] [-I dir] [-o nom] f1 ... fn
```

- Options
 - -c: Produce an object file .o from .c file by stopping before linking
 - -I dir: adds dir in the include path (used header files .h)
 - -o nom: Specifies the name of output file. By default:
 - object files have the same name as their code files
 - Executable file: a.out
 - -g: Inserts the information needed by the debugger

GCC: One compilation file



Produce an executable called module from the file module.c

```
gcc module.c -o module
```

```
statesta de al contratado de la contratada del contratada de la contratada de la contratada de la contratada
            * main.c
           * le classique Hello World
           #include <stdio.h>
#include <stdlib.h>
int main(void)
                                              printf("Hello World\n");
                                               return EXIT_SUCCESS;
```

```
bash-3.2$ ls
main.c
bash-3.2$ gcc main.c -o hello
bash-3.2$ ls
hello main.c
bash-3.2$ ./hello
Hello World
bash-3.2$ ■
```

GCC: Several compilation files



Produce an object file called module.o from a code file module.c (-c: stop before linking)

```
gcc -o module.o -c module.c
```

Produce an executable file named prog by linking its object(s) files

```
gcc -o prog module1.o ... moduleN.o
```

C compilation: The tool make



Principle

- Use to automate the compilation (large) projects, which consists of several files
- Handles file dependencies (e.g. #includes)
- Allow to automatize the update of the files
- Functionality: Create a file called Makefile that
 - Defines the dependencies between targets
 - Defines the dependencies between each target and the corresponding files
 - Provides the commands required to be executed to create the target based on the dependencies (files)



A C program consist of 3 files:

prog.c, code1.c, code2.c

Graph with file dependencies

Final Target:

prog

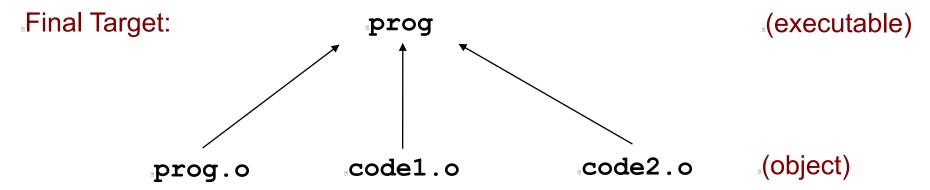
(executable)



A C program consist of 3 files:

prog.c, code1.c, code2.c

Graph with file dependencies



To produce the executable file **prog**, the following command has to be executed:

gcc -o prog prog.o code1.o code2.o

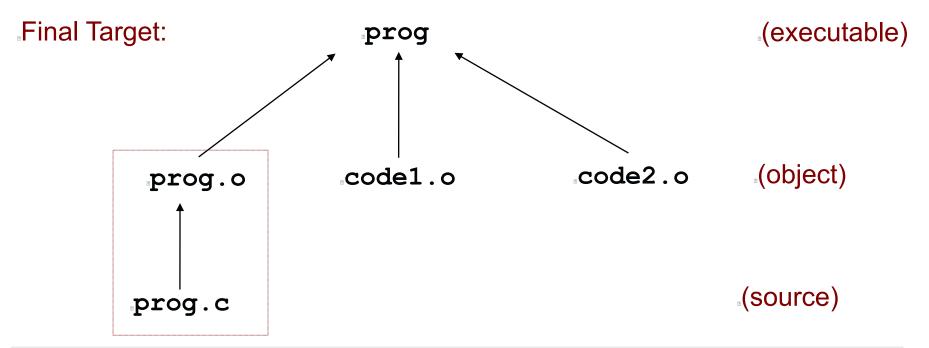
The file prog depends on the object files prog.o, code1.o and code2.o



A C program consist of 3 files:

prog.c, code1.c, code2.c

Graph with file dependencies



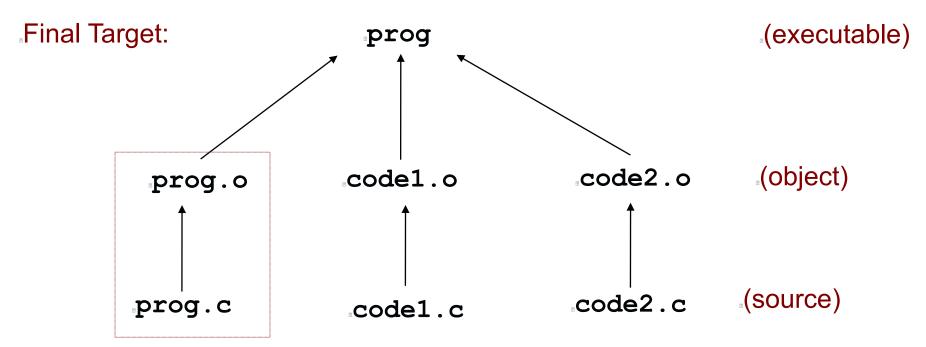
To produce the object file prog.o, we have to execute gcc -c prog.c So, the file prog.o depends on the source file prog.c



A C program consist of 3 files:

prog.c, code1.c, code2.c

Graph with file dependencies



The graph of file dependencies provides the definition of the targets, dependencies and commands of the makefile

Structure of the Makefile



Format of a rule

target: dep_file1 dep_file2 ...
(TAB) command

Application of rule:

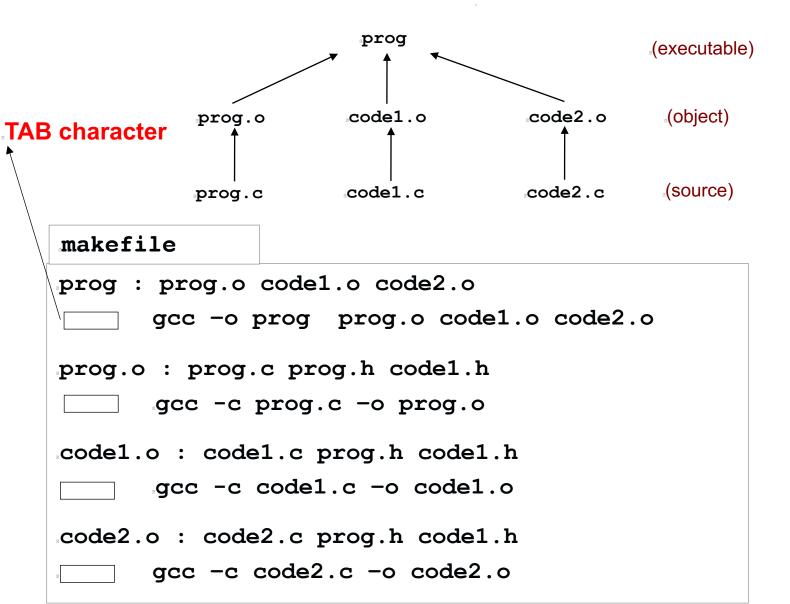
 If the date of the last modification in one of the dependent files dep_file1, dep_file2, etc. is more recent that the date of the file target, the tool make executes the command

Use: command make

- The tool executes the first target of the file makefile of the working directory
- make label: Tool executes the rule associated with the target called label

Structure of the Makefile: Example

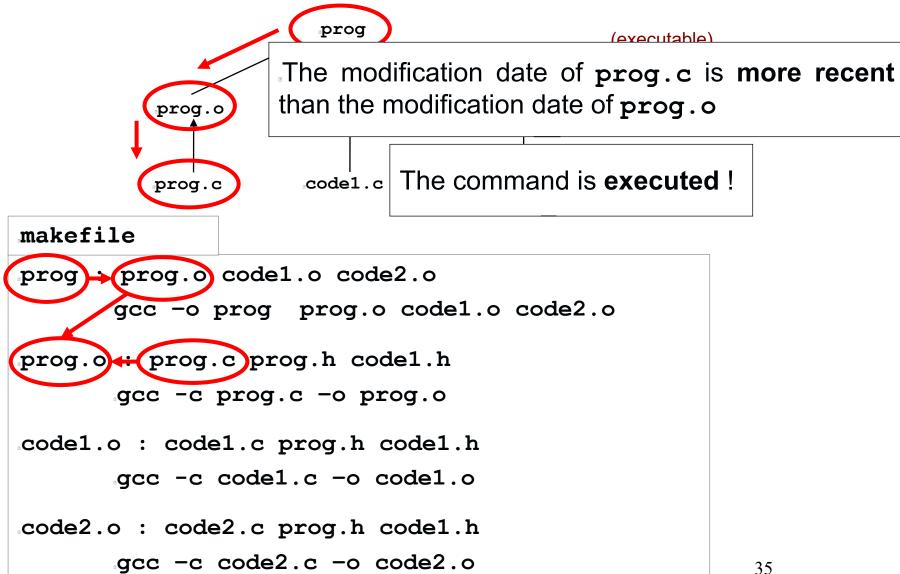




Structure of the Makefile: Example



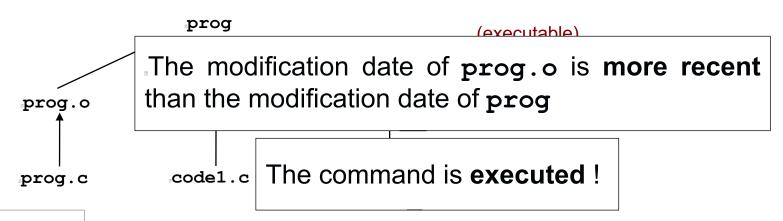
We modify prog.c, what happens when we execute make?



Structure of the Makefile: Example UNIVERSITÉ



■ We modify **prog.c**, what happens when we execute **make**?



makefile prog : prog.o code1.o code2.o gcc -o prog prog.o code1.o code2.o prog.o : prog.c prog.h code1.h gcc -c prog.c -o prog.o code1.o : code1.c prog.h code1.h gcc -c code1.c -o code1.o code2.o : code2.c prog.h code1.h gcc -c code2.c -o code2.o

Structure of the Makefile: Example UNIVERSITÉ DI



- Add a clean target: Remove object files, executable
- Use: make clean

```
makefile
prog : prog.o code1.o code2.o
      qcc -o prog prog.o code1.o code2.o
prog.o : prog.c prog.h code1.h
      gcc -c prog.c -o prog.o
code1.o : code1.c prog.h code1.h
      gcc -c code1.c -o code1.o
code2.o : code2.c prog.h code1.h
      gcc -c code2.c -o code2.o
clean :
      rm *.o proq
```

Structure of the Makefile: Example 2 UNIVERSITÉ DE CONTROLLE CONTR



```
hello: hello.o main.o

gcc -o hello hello.o main.o

hello.o: hello.c hello.h

gcc -c hello.c

main: main.c hello.h

gcc -c main.c

clean:

rm *.o
```

Structure of the Makefile: Example 2 UNIVERSITÉ DE



```
CC=qcc
CFLAGS=-q -Wall -I.
DEP=hello.h
OBJ=hello.o main.o
EXEC=hello
all: $(EXEC)
$ (EXEC) : $ (OBJ)
        $(CC) -o $@ $^
%.o: %.c $ (DEP)
        $(CC) -c $(CFLAGS) $< -o $@
.PHONY: clean mrproper
clean:
        rm *.0
mrproper: clean
        rm $(EXEC)
```

Common compilation error messages



Example of an error message

```
bash-3.2$ gcc main.c ppcm.c -o main
main.c:19: error: expected '=', ',', ';', 'asm' or '__attribute__' before 'main'
bash-3.2$ |
```

- Set of errors!
 - Always start correcting the first one (they can be dependent)!!

```
bash-3.2$ gcc main.c ppcm.c -o main
main.c: In function 'main':
main.c:21: error: 'a' undeclared (first use in this function)
main.c:21: error: (Each undeclared identifier is reported only once
main.c:21: error: for each function it appears in.)
main.c:22: error: expected ';' before 'int'
main.c:23: error: 'variableLocale' undeclared (first use in this function)
bash-3.2$
```

Debugging



- Usual bugs
 - Infinite loops
 - Loop boundaries
 - Non-Initialized variables
 - Segmentation faults
 - Algorithmic faults
 - ...
- Tools to support the correction of the code
 - The debugger is a software that helps the developer in the analysis of the bugs of a program
 - Examples : GDB, valgrind ...

Tool GDB



- With the GDB, we can
 - Start the execution of the program
 - Stop the execution of the program based on some conditions
 - Examine what happened when the program stopped
 - Execute the program step-by-step
 - Perform modifications during the execution of the program
- Use: Make the link between the program and the debugger
 - Add the information of debug during compilation
 - ✗ Option -g and -ggdb

Tool GDB: Commands



- Start execution: run
- Show the stack of the call: backtrace
- Break Point : break (help breakpoints)
- Continue execution : cont
- Execute one instruction : step
- Display the variables
 - Display nom_variable
 - watch nom_variable (allows to monitor)
 - info locals (local variables)
- **...**
- GDB Quick Reference Sheet !!!

Tool Valgrind



- With Valgrind we can identify memory leaks
 - Example of a memory leak

```
/*************
* test.c
    **************/
#include <stdlib.h>
int main()
{
    int a=2;
    int* p = (int*) malloc(1024);
    p=&a;
    return 0;
}
```

Use: valgrind --tool=memcheck ./test



...before going in more details about the C syntax...

Code style is very important!



- Understanding the code written by someone else is a hard job!
- Use a universal language (readable by everyone)
 - Πιστεύω ότι κανένας δεν μπορεί να καταλάβει αυτή τη γραμμή χωρίς να χρησιμοποιήσει μεταφραστή
- Decide a code style and be CONSISTENT
- Comments
 - Documentation: Programmers that want to use the code
 - Implementation: Programmers that maintain the code, i.e. they have to understand your implementation
- Names that mean something
 - int z Vs int counter



CM 4:

Basic components of C language: Primitive data types, Libraries and Control structures

Instructions & Comments



- Instruction is any expression finishing with;
- Block of instructions:

```
{
list of instructions
}
```

- Comments
 - /* this is a comment */
 - // this is one line comment
- The language C does not support disambiguation
 - The keywords of the language are reserved and cannot be used for naming variables and functions.

Value representation



- Hexadecimal numbers begin with 0x:
 - 0x23 is 23₁₆ (which is 35 in the base of 10)
- Octal numbers begin with 0:
 - 023 is 23₈ (which is 19 in the base of 10)
- No prefix for decimal numbers:
 - 23 is 23₁₀
- Negative representation operator
 - -0x23 is the negation of 0x23

Data object definition



Syntax:

```
type identifier;
```

- Function:
 - Reserve in the memory a part equal to the size required to store a data object of the type type
 - Associate the starting address of this memory part with the identifier

Primitive data types: Type specifiers



- Same as in JAVA:
 - Integral
 - x char, short, int, long
 - Real and complex
 - float, double
- With additional specifiers for the representation:
 - unsigned: always 0 or positive
 - Representation: Binary
 - signed: positive or negative:
 - × Representation: implementation depended
 - 1's complement
 - 2's complement (usual one)
 - sign and magnitude
- Example: unsigned short int

Primitive data types: Size



- The exact data type size is NOT part of the C standard
 - int can be 16 or 32 bits...
- The standard only says that the range of a larger type shall be at least as big as the smaller type.
 - char ≤ short ≤ long ≤ long long
 - char: The minimum supported data type
- Compiler decides based on the target platform architecture! (though, there is a minimum allowed width)
- sizeof() can give the actual size in bytes
- To give a specific storage size use the YintX_t type
 - Y = u for unsigned
 - \bullet x = 8, 16, 32 etc for the number of bits to be used

Primitive data types: Min values

RENNES 1			
Туре	Storage size	Value range	
char	usually signed	(but not standard)	
unsigned char	1 byte	0 to 2 ⁸ – 1	
signed char	1 byte	$-(2^7)$ to $2^7 - 1$	
int	2 bytes	$-(2^{15})$ to $2^{15} - 1$	
	(usually 4 bytes)		
unsigned int	2 bytes	0 to 2 ¹⁶ – 1	
	(usually 4 bytes)		
short	2 bytes	$-(2^{15})$ to $2^{15} - 1$	
unsigned short	2 bytes	0 to 2 ¹⁶ – 1	
long	4 bytes	$-(2^{31})$ to $2^{31} - 1$	
unsigned long	4 bytes	0 to 2 ¹⁶ – 1	
long long	8 bytes	$-(2^{63})$ to $2^{63} - 1$	
unsigned long long	8 bytes	0 to 2 ⁶⁴ – 1	
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Primitive data types: Float



Туре	Storage size	Value range	Precision
float	4 byte	1.2E-38 to 3.4E+38	single
double	8 byte	2.3E-308 to 1.7E+308	double
long double	10 byte	3.4E-4932 to 1.1E+4932	extended

Defining primitive data objects: Examples RENNES

Definitions

- int i;
- float f1, f2;
- unsigned char c1, c2;
- What is value of these identifiers?
- Whatever is already stored in the memory just reserved... (except global and static variables)
 - Definitions with initialization
 - int i = 3;
 - float f1 = 1.2, f2 = 0.3;
 - unsigned char c1 = 'A', c2 = 65;

INITIALIZE ALL VARIABLES!!

Value of ASCII character

Operators



- Perform operations on data objects
- They have priorities in the order they are applied:
 - Check the priority order
 - USE PARENTHESIS
- Operator types
 - Arithmetic
 - Relational
 - Logical
 - Bitwise
 - Assignment
 -

Arithmetic operators

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		A = 10, B = 20
Operator	Description	Example
+	Adds two operands.	A + B = 30
_	Subtracts second operand from the first.	A - B = -10
*	Multiplies both operands.	A * B = 200
1	Divides numerator by de-numerator.	B / A = 2
%	Remainder after integer division.	B % A = 0
!!!	(assignment at the same time).	
++	Increases the integer value by one	A++ = 11
	Decreases the integer value by one.	A = 9

- Postfix (Post-increment): Do the operation and return old value
 - A++ returns 10 and then does A=11
- Prefix (Pre-increment): Do the operation and return new value
 - ++A does =11 and then returns A

Relational operators

		A = 10, B = 20
Op.	Description	Example
==	Checks if the values of two operands are equal.	
	If yes, the condition becomes true.	(A == B) is not true.
!=	Checks if the values of two operands are not equal	
	different. If yes, the condition becomes true.	(A != B) is true.
>	Checks if the value of left operand is greater than the	
	value of right operand. If yes, it becomes true.	(A > B) is not true.
<	Checks if the value of left operand is less than the	
	value of right operand. If yes, it becomes true.	(A < B) is true.
>=	Checks if the value of left operand is greater than or	
	equal to the value of right operand.	$(A \ge B)$ is not true.
<=	Checks if the value of left operand is less than or	
	equal to the value of right operand.	(A <= B) is true

Any integer value can be considered as Boolean

- ★ A FALSE expression is coded as a zero integer
- ★ A TRUE expression is coded as any non-zero integers value.

Logical operators



■ The operands are TRUE or FALSE

		A = 1, B = 0
Op.	Description	Example
&&	Logical AND: If both the operands are non-zero, then the condition becomes true.	(A && B) is false.
	Does not evaluate 2 nd operand, if the first is false	
II	Logical OR: If any of the two operands is non-zero, then the condition becomes true.	(A B) is true.
	Does not evaluate 2 nd operand, if the first is true	!(A && B) is true.
!	Logical NOT: It is used to reverse the logical state of its operand. If a condition is true, then Logical NOT operator will make it false.	

Bitwise operators



■ The operands are integral (short, int, unsigned, etc)

	A = 0x3 (0011) B = 0x2 (001)		
Op.	Description	Example	
&	Binary AND: Copies a bit to the result, if it exists in both operands.	(A & B) = 0x2 (0010)	
I	Binary OR: Copies a bit if it exists in either operand.	$(A \mid B) = 0x3$ (0011)	
٨	Binary XOR: Copies the bit if it is set only in one operand but not both.	$(A ^ B) = 0x1$ (0001)	
~	Binary Ones Complement: Has the effect of 'flipping' bits.	$(\sim A) = 0xC$ (1100)	
<<	Binary Left Shift : The left operands value is moved left by the number of bits specified by the right operand.	(A << 2) = 0xC (1100)	
^	Binary Right Shift: The left operands value is moved right by the number of bits specified by the right operand.	(A>>2) = 0x0 (0000)	

Assignment operators



	A =	: 10, B = 20
Op.	Description	Example
=	Simple assignment operator. Assigns values from	C = A + B
	right side operands to left side operand	C = 30
+=		
-=		
*=		
\=	In can be combined with the other arithmetic	A 1 – 1
%=	operators (+, -, *, %) and bitwise operators (&, , ^, <<,>>).	A += 1 A = 11
& =	It performs the arithmetic operation with right	A - 11
=	operand and the left operand and assign the result	
^=	to the left operand.	
<<=		
>>=		

Common error: Assignment Vs Testing UNIVERSITÉ DE



- Assignment: nom_variable = expression ;
 - x=3;
 - It is an expression that stores the value 3 to the variable x
- Test of equality: expression1 == expression2 ;
 - x == 3;
 - It returns 1 if the stored value in x is equal to 3, otherwise 0

ATTENTION:

= and == are different!



Kahoot time! (6 Q)

So, what happens if we mix types in an operation?

a (operator) b type a ≠ type b

- Automatic IMPLICIT and SILENT conversion!
- Both operands (a, b) must be data objects of the same type
- Converted to the common higher rank of the operands based on the type width:

PROMOTION

- Integer promotions:
 - signed > unsigned
 - small integer types: char, short, int16, int8 etc (BOTH signed and unsigned) are converted to
 - x int: IF the range of int can represent ALL values of the original
 type
 - x else: unsigned int

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And during assignment?



....Still automatic IMPLICIT and SILENT conversion!

Converted to the type of data object that stores the result:

Higher rank: PROMOTION

Lower rank: **DEMOTION**

Primitive data types: Implicit conversion UNIVERSITÉ DE long double double float unsigned long long long long **DEMOTION PROMOTION** unsigned long long unsigned int Small integer types short char

1 Byte

2 Bytes

66

Explicit conversion: Cast operator



Syntax (type) identifier;

- Creates a copy of the data object and changes its type.
- The data object of the identifier is NOT modified

With the Implicit or explicit conversion UNIVERSITÉ DE



- The type changes →
- So, the storage size and/or the representation changes →
- So, the value CAN change

And, when we cast different representation RENNES

- Casting float to int: It changes representation!
 - float is represented in IEEE 754 single precision
 - int is usually 2's Complement

```
float f = 3.9;
int val = (int) f;
```

- Value gets truncated!
 - ★ 3.9 is truncated to 3.0 and stored as 3 in 2's complement
- Casting int to float: It changes representation!

```
int f = 3 ;
float val = (float) f ;
```

Value to convert to float (thus, 3 is converted to 3.0)

Type conversion: Examples



- int x;
- Implicit conversion: x = 1.5 + 1.6;
 - 1.5 → double
 - 1.6 \rightarrow double
 - 1.5 + 1.6 = 3.1 double addition
 - \bullet x \rightarrow int
 - 3.1 → 3 DEMOTION to match with int

$$x = 3$$

- Explicit conversion: x = (int) 1.5 + (int) 1.6;
 - 1.5 \rightarrow double
 - (int) 1.5 \rightarrow 1 DEMOTION
 - 1.6 \rightarrow double
 - (int) 1.6 \rightarrow 1 DEMOTION
 - 1 + 1 = 2 integer addition
 - \bullet x \rightarrow int

$$\times x = 2$$

Type conversion: Examples



- float x;
- Implicit conversion:

$$x = 3/5;$$

- 3 → int
- \bullet 5 \rightarrow int
- 3/5 = 0 integer division
- \bullet x \rightarrow float
- 0.0 PROMOTION to match with float

$$x = 0.0$$

- Explicit conversion
 - \bullet 3 \rightarrow int
 - (float) $3 \rightarrow 3.0$ PROMOTION
 - Promotion 5 → 5.0 PROMOTION
 - \bullet 3.0/5.0 = 0 float division
 - \bullet x \rightarrow float

$$\times x = 0.6$$

x = (float) 3/5;

Type conversion: Usual problems



- Casting signed to unsigned:
 - Negative value is converted to a high positive value
 - Especially occurs during function calls or return values
- Casting a smaller signed int type to a larger unsigned int type
 - Sign extension: The sign bit is used in the unused bits, i.e. the Most Significant Bis (MSB) of the larger type
 - char and short are signed!
 - Negative value is converted to a high positive value



Type conversion: Usual problems



- Casting a larger type to a smaller one (DEMOTION):
 - Truncation: Cuts off the MSB (overflow occurs).

```
int f = 0x12345678 ;
short int g;
g = f;
```

Comparisons, especially with unsigned, char and short!

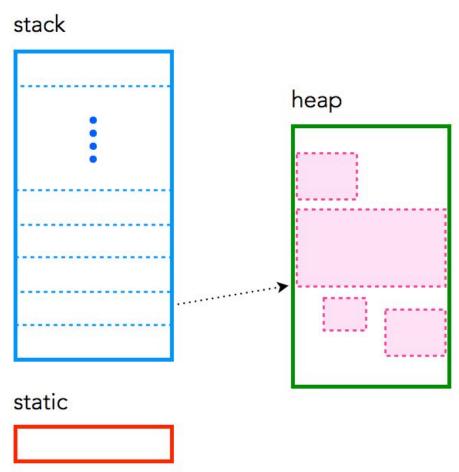
```
unsigned int a = 0;
if (a > -1)
    printf("TRUE\n");
else
    printf("FALSE\n");
    SECURITY ISSUES
```

Variable allocation and lifetime



Static memory:

- Variable Lifetime: complete program
- Stores:
 - ✗ Global variables: data objects declared outside of all function blocks
 - X Variables defined as static
- Global and static variables are initialized to zero by default by the compiler
- Allocation at compile/link time



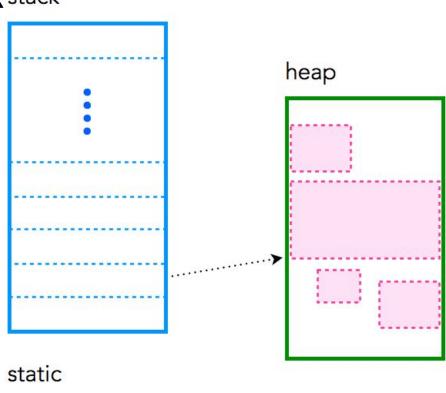
Variable allocation and lifetime





Stack:

- Variable Lifetime: specific block stack
- Stores
 - **X Local variables**: data objects declared within a block
 - X Variables defined as auto (by default the local ones)
 - X Variables defined as
 register, but the compiler
 couldn't store them to
 registers
- Allocated in an
 - Automatic way
 - Continuous way
- Allocation during execution



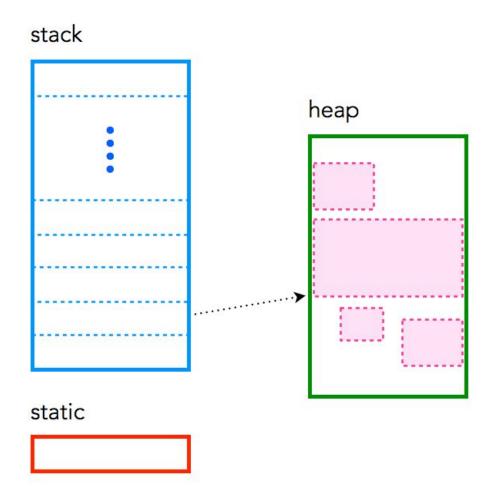
Variable allocation and lifetime





Heap:

- Variable lifetime: Managed by the program
- Large pool of memory
- Not allocated in contiguous order
- Dynamic allocation
- Allocation during execution



Data objects: Qualifiers



- Express additional information about a value
 - Ensure the correct use of the data object
- const: The value of the data object does not change once it has been initialized
- volatile: The compiler does not optimize, the data object is always accessed from the memory
- restrict: The only way to access a data object is through the pointer that points to it

Object-like Macros



Syntax: Through pre-processor directive #define #define IDENTIFIER value

Function:

- Gives a symbolic name to a data object
- Convention: Symbolic names are given in UPPERCASE

Examples

- #define PI 3.1416
- #define NOM "foo"

Other reserved constants



- Escape sequences:
 - The character '\n' is a new line (on Unix systems)
 - The character '\b' is a backspace
 - The character '\t' is a tab space
 - The character '\0' is null character

- The double quotes "..." defines a string of characters
 - There is NO string type in C
 - String = array of characters that ends with the null character (\0)
- The simple quote '.' defines a character



Basic components of C language: Libraries

Libraries



- In C, there are a lot of functions already implemented into libraries
 - Include system calls
 - Mathematical computations
 - ...
- The programmer can and should use them!
 - Add corresponding directive to the head of the compilation file
 - #include <stdio.h>
 - #include <stdlib.h>
 - #include <math.h>
 - #include <string.h>
 - ...

Two basic I/O functions!



We have two main functions for read and write to the standard input and ouptut:

```
Write: int printf(format, e<sub>1</sub>, e<sub>2</sub>,..., e<sub>n</sub>);
```

Read: int scanf(format, a₁, a₂, ..., a_n);

To use them we need to add the directive

```
#include <stdio.h>
```

Formated write: printf(...)



Syntax

• int printf(format, e₁, e₂,..., e_n);

Description

- Writes (displays) the parameters e₁, e₂, ..., e_n to the standard output based on the parameter format
- The format consists of I/O expressions and, potentially, text.
- Each I/O expression describes how the writing of one parameter e_i
 will be done using the character %
- It can have a variable number of parameters, but
 The number of e_i has to correspond to the number of I/O expressions specified in the format

Result

Number of displayed characters (value < 0 in case of an error)

Formated read: scanf(...)



Syntax

int scanf(format, a₁, a₂, ..., a_n);

Description

- Reads from the standard input (keyboard) the values and interprets them based on the parameter format
- It assigns each read value to the variables, whose addresses are given by the parameters a₁, a₂, ..., a_n
 - Address of a variable is noted by the operand & nom_variable
- It can have a variable number of parameters, but
 - ★ The number of a_i has to correspond to the number of I/O expressions specified in the **format**

Result

Returns the number of elements correctly assigned

Format of I/O expressions



- %d: implies that the I/O is a signed integer (int)
- %c : implies that the I/O is an ASCII character (char)
- %f: implies that the I/O is a float (float).
- %1f: implies that the I/O is a double (double).
- %x: implies that the I/O is a unsigned integer to be printed using the hexadecimal representation
- %s: implies that the I/O is a string
-
- More info: man format or google printf
- Remarks
 - The format can also have normal text (whatever does not have the % character)

Example



```
int main()
    int i;
    float f;
    printf("Entrez un entier :");
    scanf("%d",&i);
    printf("Entrez un nombre réel :");
    scanf("%f",&f);
    printf("Vous avez saisi %d et %f \n", i, f);
```

What happens if we use a wrong format?NIVERSITÉ DE RENNES

- Usually we have only a compiler warning
- Question: What prints the following program?

```
#include <stdio.h>
void main()
{
    int a = 1654 ;
    printf("%s\n",a);
}
```

- Answer: Segmentation fault (core dumped) or unexpected behaviour
 - Tries to print a string of characters stored in the memory starting at the address 1654 and its end is indicated by the \0 character.
 - X At the end, the program will try to access an illegal address
 - Very common error!

Type conversion: Inbuilt typecast functions Find RENNES

- atof() : convert string to float
- atoi() : convert string to integer
- atol() : convert string to long
- itoa(): convert integer to string
- ltoa(): convert long to string



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Basic components of C language: Control structures

Control structures



- Similar to the ones in JAVA and Scripts
 - NO BOOLEAN type!

Example

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```
int i=7;
while(i)
  i--;
                            Both ways are equivalent, BUT
                            the second one is more readable!
int i=7;
while(i!=0)
i--;
```

Conditions: If



```
if (expression) {
    list of instructions1
}

if (expression) {
    list of instructions1
} else{
    list of instructions2
}
```

Conditions: ternary instruction



```
expression1 ? expression2 : expression3;
```

- The value of a ternary instruction (like a multiplexer?) is determined as following:
 - expression1 is evaluated.
 - If it is true, then expression2 is evaluated and becomes the result of the entire expression.
 - If it is false, then expression3 is evaluated and its value becomes the result of the expression.

Conditions: switch - case



```
switch(expression) {
  case constant-expression1:
         list of instructions1;
                                         (optional)
         break;
  case constant-expression2:
         list of instruction2;
         break ;
                                         (optional)
  default:
         list of instructions3;
```

Loops



Syntax while while (expression) { list of instructions; Syntax do-while **do** { list of instructions } while (expression) ; Syntax for for(expression1; expression2; expression3) { list of instructions; expression1: initialization expression2: condition of termination expression3: step

Loops: Control statements



- Change execution from its normal sequence
- Instruction break;
 - Terminates the loop or switch statement
 - Transfers execution to the statement immediately following the END of the loop or switch.
- Instruction continue;
 - Causes the loop to skip the remaining of its BODY
 - Immediately retests the condition expression before reiterating.
- Instruction goto label;
 - NEVER USE goto ! (except special cases)

Other types of objects



- Derived data types:
 - Pointer type
 - Array type
 - Structure type
 - Union type
- Enumerated types:
 - Arithmetic types used to define variables
 - Variables can only assign certain discrete integer values.
- Function type
- Void type



CM 5: Pointers and Arrays

Pointer type



```
type * identifier;
```

- Function:
 - The pointer type, *, is applied over a specific type of an object
 - The type can be anything: integer, character, float, pointer, structure, function etc.
 - The identifier points (refers) to an objet of that type
 - **X** It contains the **address** of the object that it is pointing to

Pointer



Useful

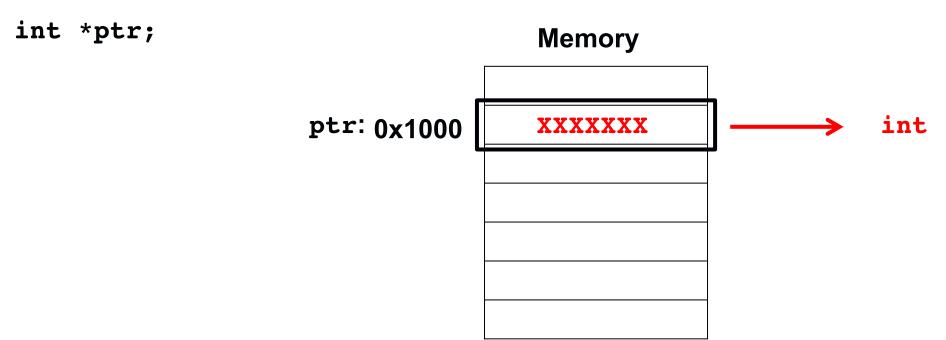
- Mechanism of indexing
- List structures
- Passing arguments to a function
- Dynamic allocation

Often associated with **NULL**

- NULL is a value of pointer that points to nothing (nothing is an invalid address)
- Attention CAPITAL letters!!

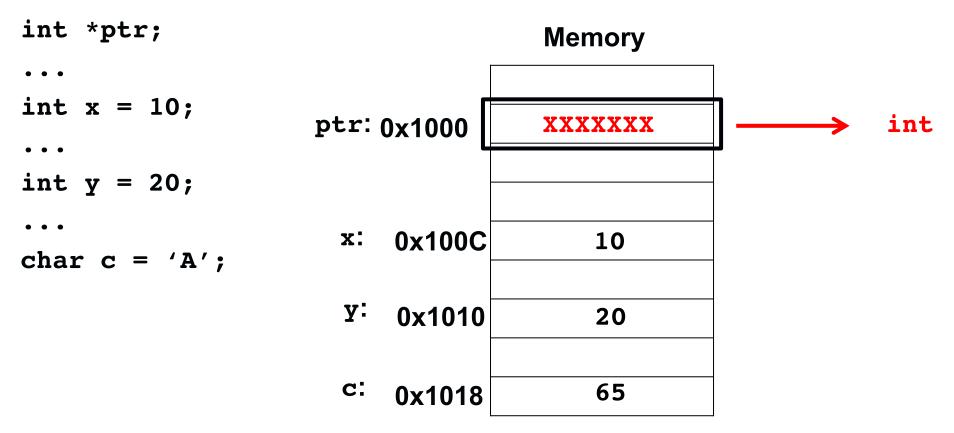
Pointer: Definition





Pointer: Definition

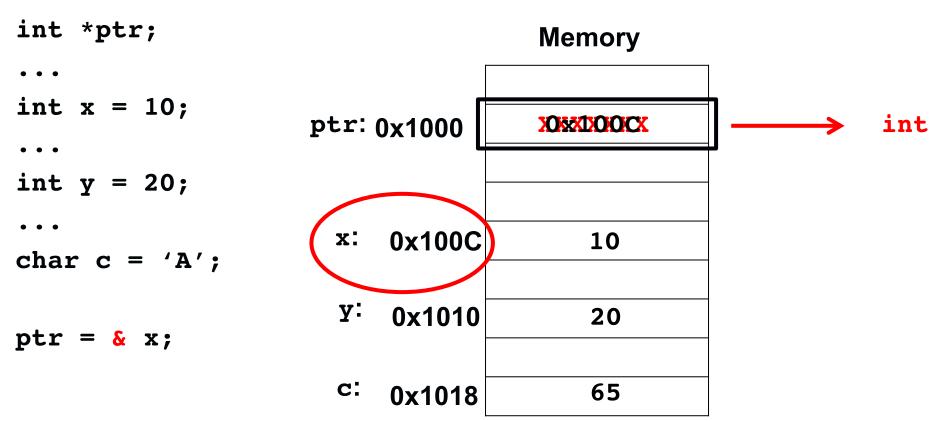




- The object ptr is declared as a pointer to integers (int*)
 - It can contain the address of x or to y (both are int), but not to c
 (as it is a char)
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Pointer Operations: Address reception UNIVERSITÉ DE RENN





Address reception (noted as &) returns the address of an object

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Pointer Operations: Address reception UNIVERSITÉ DE RENN



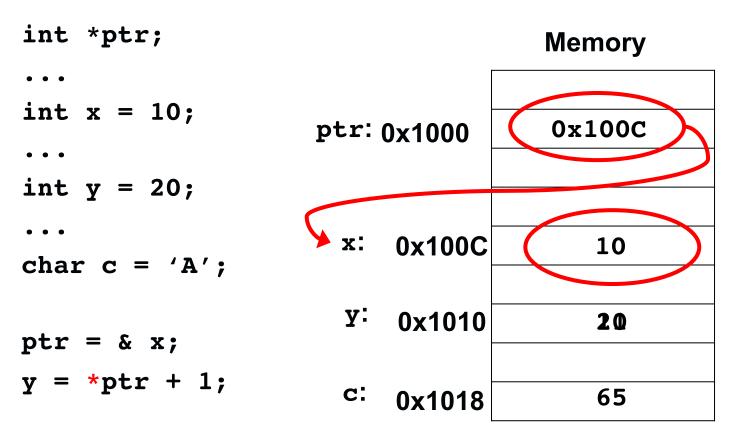
```
int *ptr;
                                     Memory
int x = 10;
                     ptr: 0x1000
                                     0x100C
int y = 20;
                          0x100C
                                         10
char c = 'A';
                       y:
                          0x1010
                                         20
ptr = & x;
                       c:
                                         65
                          0x1018
```

Address reception (noted as &) returns the address of an object

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Pointer Operations: Dereference/Accessing NES



■ The dereference (noted as *) returns the value stored in the memory where the pointer points to

Pointer Operations: Remarks



ATTENTION:

Pointer definition !=

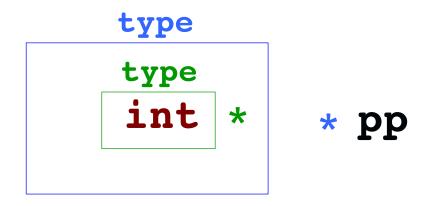
Pointer dereference !=

Multiplication operation

Pointer to pointer



- The object, where a pointer points to, can be also a pointer
- In this case, the type of the object is type*, i.e. pointer type
- int **pp; means that pp is a pointer that points to a pointer that points to integers



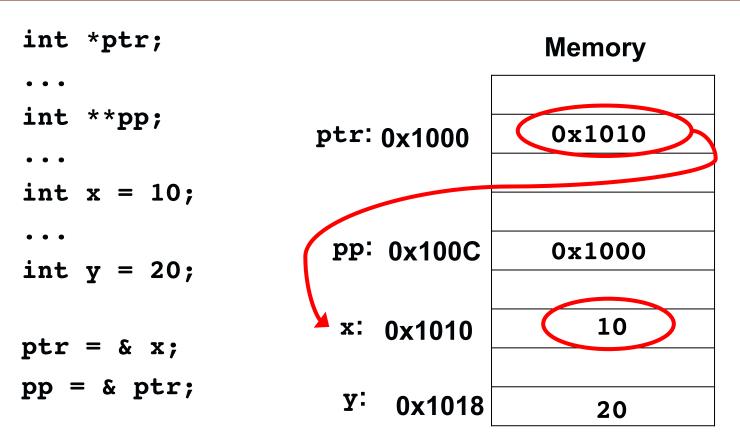


<pre>int *ptr;</pre>		Memory		
<pre>int **pp; int x = 10;</pre>	ptr: 0x1000	xxxx		int
<pre>int y = 20;</pre>	PP: 0x100C	xxxx		int*
	x: 0x1010	10		
	y: 0x1018	20		



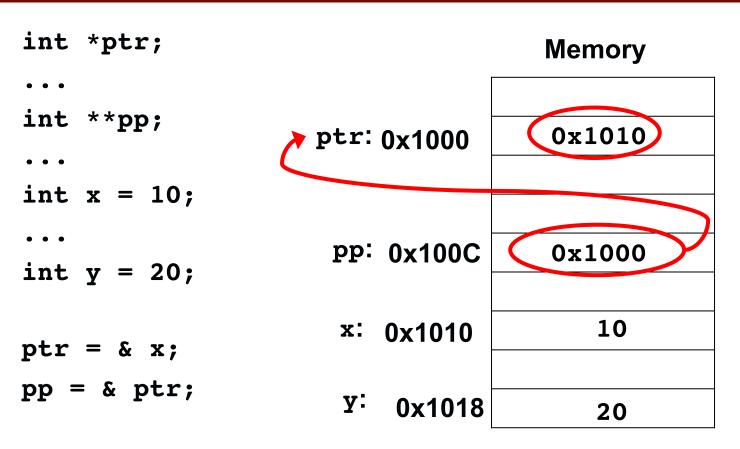
<pre>int *ptr;</pre>		Memory		
<pre>int **pp;</pre>				
•••	ptr: 0x1000	0x1010		
int $x = 10;$				
int y = 20;	pp: 0x100C	xxxx		int*
ptr = & x;	x: 0x1010	10		
	у: 0x1018	20		





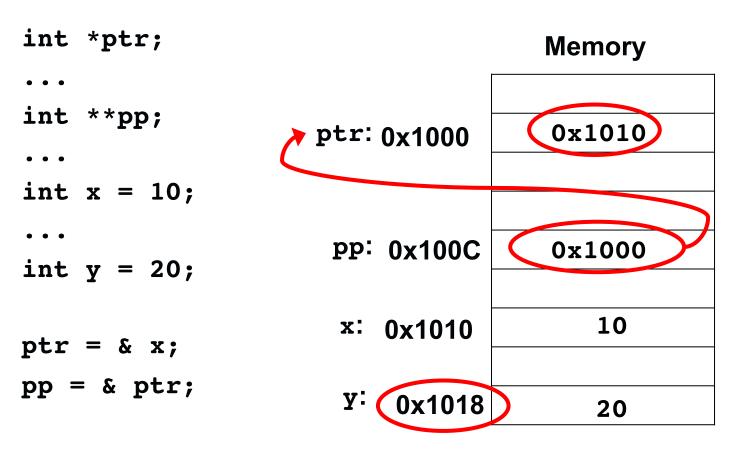
■ What is the result of *ptr?





■ What is the result of *pp?



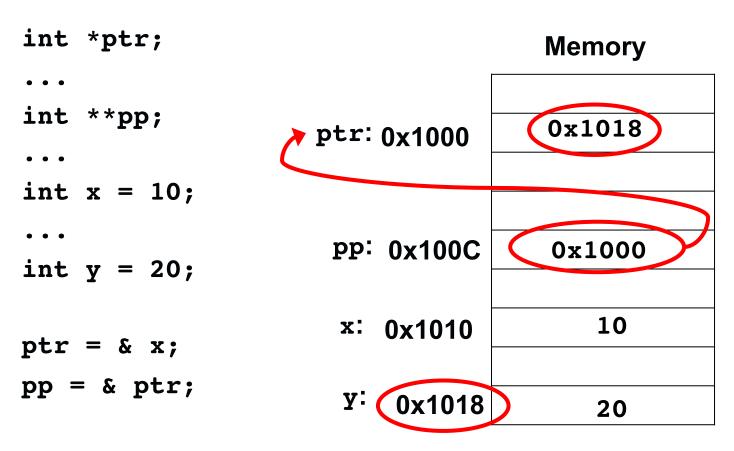


- What is the result of *pp = &y?
- The instruction *pp=&y modifies the content of the object pointed by pp (which is ptr) by storing the address of the object y

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- What is the result of *pp = &y?
- The instruction *pp=&y modifies the content of the object pointed by pp (which is ptr) by storing the address of the object y

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<pre>int *ptr;</pre>		Memory
• • •		
<pre>int **pp;</pre>	ptr: 0x1000	0x1018
• • •		
int $x = 10;$		
···	pp: 0x100C	0x1000
int $y = 20;$		
ptr = & x;	x: 0x1010	10
-		
pp = & ptr;	y: 0x1018	20

What is now the result of *ptr?

Let's...



Kahoot time !! Q1 – Q6

Pointer arithmetic



- A pointer is NOT ONLY an address!
- It provides information about
 - What type is the object stored in the address where the pointer points to
 - And thus, what is the size of the object.
 - x Data type's size can be given by the sizeof() operator.
 sizeof(int), sizeof(short)
 - The return value depends on the compiler and the target architecture!

Pointer arithmetic operators



type *ptr;int x;

Operator	Description
ptr + x	Adds sizeof(type) * x to the pointer value
ptr - x	Subtracts sizeof(type) * x to the pointer value ptr
ptr ++	Increases the pointer value by sizeof(type)
ptr	Decreases the pointer value by sizeof(type)

Example (assuming sizeof(int) == 4).

<pre>int *ptr;</pre>	ptr:	0x1000	
	ptr+1	0x1004	
	ptr+2	0x1008	

Pointer arithmetic: Problems!



- You can make a pointer point anywhere, the compiler will NOT complain. The following are valid:
 - ptr + 10000
 - ptr 10000
 - char *ptr = (char *) 0x0000ffff
- SECURITY ISSUES
- However, problems occur when you attempt to access an invalid memory address by dereferencing the pointer
- Usually, the operating system forbids access outside of the program's allocated memory
 - If the program can access the memory location, the data is not valid
 - If the program cannot access, it results to segmentation fault
 - ★ The operating system sends the SEGFAULT signal to the process that caused the memory violation, the process core dumps and terminates

Array type



```
type identifier[SIZE];
```

- Function:
 - The identifier is an array with a size equal to SIZE
 - Each array element is a object of type type
 - The array is statically allocated. During definition
 - X The SIZE must be is a constant
 - * The memory to store **SIZE** elements of type **type** is reserved
 - **X** The name of the array (**identifier**) is a **constant pointer** that points to the address of the first element of the array
 - The index of the array is an integer
 - Values from 0 to SIZE-1
- It can be initialized during its definition (SIZE=elements) int array[]={1,3,-2,19,-79};

Array type: Definition



int arr[10];

Memory

arr ——>arr[0]:0x1000 arr[1]:0x1004

. . .

arr[9]:0x1028

wemory	
XXXXXXX	
xxxxxx	
XXXXXX	
XXXXXX	
XXXXXX	
XXXXXXX	
XXXXXXX	
XXXXXX	
XXXXXX	
XXXXXX	

arr == & arr[0]

*arr ==

arr[0]

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Arrays and Pointers

- UNIVERSITÉ DE RENNES 1
- The notion of array is very close to the notion of pointer
- arr[i] == *(arr+i), so they indicate the same element!
- arr+i = arr + sizeof(type)*i

Memory

0

int arr[10];

 $arr \longrightarrow arr[0]:0x1000$

arr[1]:0x1004

arr[2]:0x1008

 $arr+3 \rightarrow arr[3] 0x100C$

arr is a constant pointer

 \rightarrow

arr++ is not allowed

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*(arr+3)

arr[9]:0x1028

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1
2
3
4
5
6
7
8
9

Arrays and Pointers

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We can use a pointer to point to the array

```
int arr[10];
int *ptr = arr;
                          ptr

arr[0]:0x1000

                        ptr++ --- arr[1]:0x1004
                        ptr++ \longrightarrow arr[2]:0x1008
                        ptr++ ---> arr[3]0x100C
                                   arr[9]:0x1028
```

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Memory 0 3 5 6 8 9

Arrays and Pointers



- What is the difference?
- When we define an array:
 - We reserve one memory block equal to the size required to store the values of all the array elements
 - The array name is a constant pointer, its value cannot be modified
- When we define a pointer:
 - we do NOT reserve the memory where we will store the values of the array elements,
 - We reserve ONLY the memory, where we will store the pointer
 - The pointer is a variable, its value can be modified

An array of pointers



- The type of the elements of an array can be also pointers, that is of type type*
- Example:

```
int *arr[10];
int **ptr = arr;
```

- An element in the array arr (e.g. arr[0]) is of type int*
- The constant pointer arr is the address of the first element of the array
- Hence, the ptr is a pointer to a pointer of integers
 ptr+i = ptr + sizeof(int*) * i

An array of characters: String



- There is no string type in C
- A string is an array of characters finishing with '\0' char string[SIZE]
- The SIZE must be large enough to hold this additional byte!
 - Compiler puts the '\0' character at the first empty element of the array
 - No automatic array bound checking!
 - ✗ If number of characters and '\0'> array maximum size
 - UNDEFINED BEHAVIOUR
- Responsibility of the programmer:
 - content of the string <= size reserved in the array

This is the source of A LOT of ERRORS!

SECURITY ISSUES

String initialisation



$$c \longrightarrow c[0]: 0x100C$$

Memory

b \0	С	đ
	С	đ
\0		

String: Remarks



Static declaration

```
char str1[]="hello";
char str2[]="world";
```

- We CANNOT assign a string to another (str2 = str1;). They are constant pointers!
 - We have to copy one to the other!
- Once the string is initialized, it CANNOT be assigned to another set of characters

```
str1 = "bye";
```

Arrays and Pointers: Attention!



Example:

```
int arr[10] ;
int *ptr1, *ptr2 ;
ptr1 = arr + 3 ; //ptr1 == & arr[3]
ptr2 = ptr1 - 2 ; //ptr1 == & arr[1] ;
```

- Pointer ptr1 points the fourth element of the array
- Pointer ptr2 points two elements back from ptr1.
- Still, within the memory block of the array arr.

Arrays and Pointers: Attention!



Example:

```
int arr[10];
int *ptr1, *ptr2;
ptr1 = arr + 30; //ptr1 == & arr[30]
ptr2 = ptr1 - 2; //ptr1 == & arr[28]
```

- Now, pointer ptr1 points the outside the memory block of the array!
- There is no compiler error, the compiler still computes the address!
- There is no core dump, if we are accessing the address space dedicated to the program!
- The data that we will read through *ptr1 are not valid!

Let's...



Kahoot time !! Q7 – Q11

2D Array: Array of Arrays



Syntax:

```
type identifier[ROWS][COLS] ;
```

- A 2D array named identifier is defined as
 - a number of **ROWS**
 - Each row is an 1D array of cols elements of type type
- When initializing during definition:
 - It is always required to give the COLS dimension
 - The number of **ROWS** is NOT always required
 - **x** it can be calculated based on the number of columns and the initialization.

Array of Arrays: Initialization



```
int arr[3][2] = {
                        { 1 , 4 },
                        { 5 , 2 },
                        { 6 , 5 }
                  };
\blacksquare int arr[3][2] = { 1, 4, 5, 2, 6, 5 };
int arr[3][2] = { { 1 }, { 5, 2 }, { 6 }};

    Not initialised elements will get 0

\blacksquare int arr[3][2] = {{ 0 }};
```

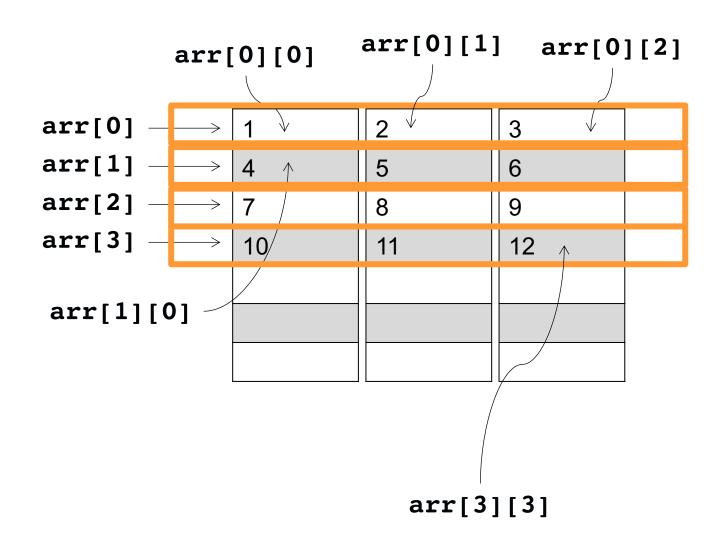
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All elements are initialised to 0

Array of Arrays



int arr[4][3] ;



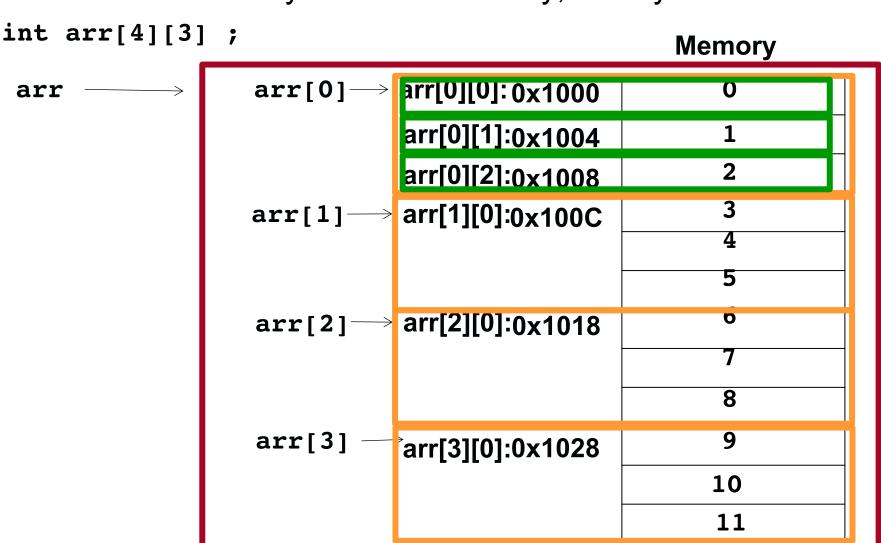
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Array of Arrays



It is continuously stored in memory, row by row



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Array of Arrays: Remarks!



- The name of the array is a constant pointer the points to the first element of the array
 - 2 levels of hierarchy
 - × Row
 - Columns -> elements
 - first level of hierarchy is the row, so it points to the first raw!

```
x arr == &arr[0]
```

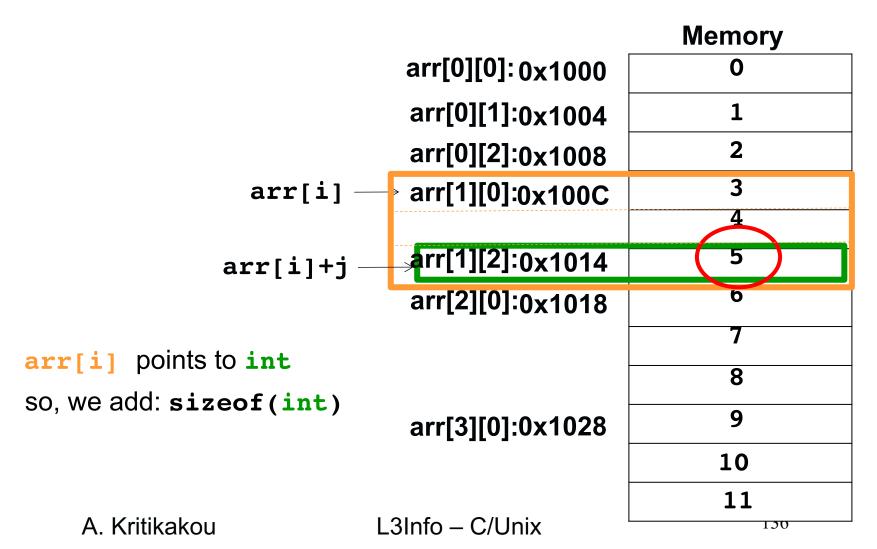
- The name of the array is NOT a pointer to pointer to type,(type**)
 - This is NOT OK with the pointer arithmetic!
- The name of the 1D arrays, i.e. the rows, point to the first element of the 1D array

Array of Arrays: Pointer arithmetic



■ arr[i][j] = *(arr[i] + j)

int arr[4][3];



Array of Arrays: Pointer arithmetic



int arr[4][3];

arr ———————————————————————————————————
arr+i —> so, we add the size of the array: sizeof(int)*COLS *(arr+i)+j
<pre>*(arr+i) is arr[i] i.e. 1D array of int so, we add: sizeof(int)</pre>

	Memory	
arr[0][0]:0x1000	0	
arr[0][1]:0x1004	1	
arr[0][2]:0x1008	2	
arr[1][0]:0x100C	3	
	4	
arr[1][2]:0x1014	(5)	
		4
arr[2][0]:0x1018	0	.
	7	
	7	
	7	
arr[2][0]:0x1018	7 8	
arr[2][0]:0x1018	7 8 9	

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Array of Arrays: Pointer arithmetic



arr points to 1D array of int

*arr+i*COLS

Manually add the size of the array instead of compiler

(*arr)+i*COLS+j

	Memory
arr[0][0]:0x1000	0
arr[0][1]:0x1004	1
arr[0][2]:0x1008	2
arr[1][0]:0x100C	3
	4
arr[1][2]:0x1014	5
arr[2][0]:0x1018	6
	7
	8
arr[3][0]:0x1028	9
	10
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Pointer to an array



```
int arr[ROWS][COLS];
int (*ptr) [COLS];
ptr = arr;
```

- (*ptr) is a pointer to an 1D array of cols elements. The parentheses are NOT optional!
 - Without the parentheses, ptr becomes an array of cols pointers,
 not a pointer to an array of cols ints
- Doing pointer arithmetic, we can compute correctly the address of ptr[row][col]:



Kahoot time! Q12 – Q19



CM 6: Functions

Function body



Syntax

```
type identifier (type1 parameter1, type2 parameter2, ...)
{
     ...
     return expression;
}
```

- Function
 - Subroutine to performs a specific job with a unique name (identifier)
 - The functions are defined at the same level as main
 - Can be used (called) using its name several times from different positions
 - Information
 - can be passed to the function through the function parameters
 - can be returned to the position where the function was called through the return statement
 - The type of the function equals to the type of the value returned
 - X If no function type is defined, then it is by default int
 - X If no return value, then the function type is void

Return value



- Syntax: return expression;
- The possible types that can be returned by the evaluation of the expression are
 - The primitive data types
 - The structures
 - The pointers
 - Attention: We CANNOT return arrays! We can return a pointer char*,
 int*, ...
 - passed as argument during the function call
 - dynamically allocated inside the function

Exit value



Syntax: exit(int);

- Directive:
 - #include <stdlib.h>
- Function:
 - Exits the program
- Parameters
 - 0 : usually means your program completed successfully (EXIT_SUCCESS)
 - Non zero: Different values that can be used as error codes (EXIT FAILURE)

Function call



Syntax

```
variable = identifier(argument1, argument2, ...)
```

- Function
 - When a function call occurs, the function parameters:
 - x are created (LOCAL variables)
 - x initialized with the value of the passed argument
 - parameter1 = argument1
 - parameter1 is a copy of argument1 manipulated inside the function
 - * the original argument is unchanged.
 - The function MUST be at least declared before it is called
 - × Put the function body before the first function call
 - Put the function declaration before any function body

Function prototype



Syntax

```
type identifier (type1 parameter1, type2 parameter2 ,...);
```

- The function declaration is the function prototype
- The complete bloc of the function is replaced by a semicolon;
- The name of the parameters can be omitted, but not their type

Function call: Remarks



- The type of the arguments must be the same as the type of the function parameters
- When they are different
 - Automatic and implicit conversion!
 - ★ char and short (signed, unsigned) → int
 - \times float \rightarrow double
 - Implicit conversion rules are overridden by function prototype
- The **type** of the **variable**, where the return value is assigned, must be the same as the **return type** of the function

Function: Examples



Function bodies

- No Arguments and No return: void proc1() { ... }
- Arguments and No return: void proc2(int n, float x) { ... }
- No Arguments and return: int proc3() { ... }
- Arguments and return: int proc4(int n, float x) { ... }

Function calls

```
float r ;
int a ;
```

- No Arguments and No return: proc1();
- Arguments and No return: proc2(a, r); or proc2(42, 3.14);
- No Arguments and return: a=proc3();
- Arguments and return: a=proc4(a,r);

Variable allocation in memory



Static:

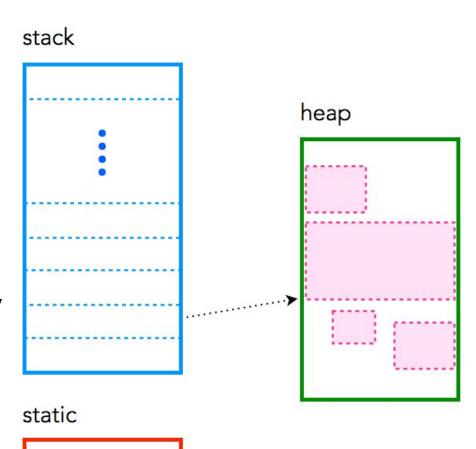
- global and static variable storage
- permanent for the entire run of the program.
- Allocation at compile/link time

Stack:

- local variable storage
- automatic, continuous memory
- Allocation during execution

Heap:

- dynamic storage
- large pool of memory, not allocated in contiguous order
- Allocation during execution



Stack



- Each time a function call occurs, the stack grows to store the local variables of the function
- Each time a function returns, the stack shrinks and removes the function local variables
- So, stack variables ONLY exist while the function that created them is running
- No need to manage the memory yourself, variables are allocated and freed automatically
- The stack has size limits

What is visible inside a function?



- The parameters of the function
- The local variables (the ones defined inside the function bloc)
- The global variables of the compilation file
- The imported global variables (specified by extern)
- The other functions declared before this function
 - inside the compilation file
 - external



Kahoot time! Q1- Q5

Passing arguments: By value



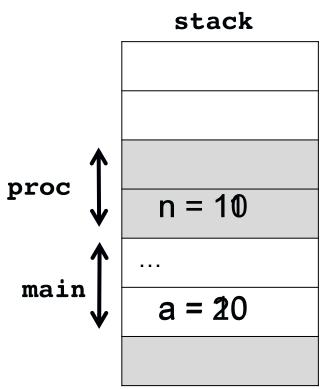
Function parameters: LOCAL variables initialized with the value of the passed argument

- Pass by Value: pass the data object itself
- The parameter is initialized with the actual value of the data object
- When a function returns:
 - The space reserved for the local variables of the function is erased!
 - With that, any modification occurred on the local variables

Pass by value



```
void proc(int n) {
         n++;
}
void main {
    int a=10;
    proc(a);
    a = a + 10;
    ...
}
```



Passing arguments: by Reference



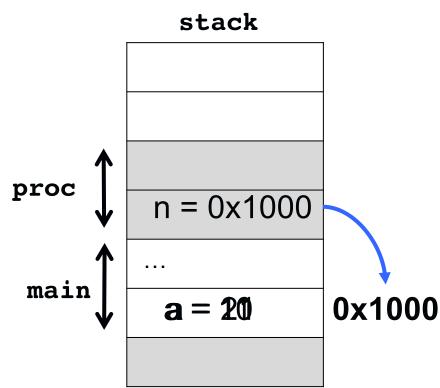
Function parameters: LOCAL variables initialized with the value of the passed argument

- Pass by Reference: pass the pointer of the data object
- The parameter is initialized with the address of the data object
- When a function returns:
 - The space reserved for the local variables of the function is erased!
 - But... the modifications now occurred at addresses belonging to the function that made the call!

Pass by reference



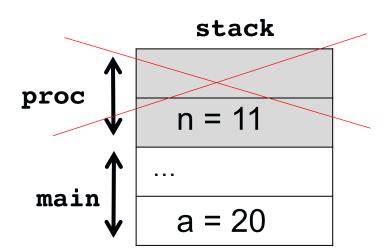
```
void proc(int *n) {
          (*n)++;
}
void main{
        int a=10;
        proc(&a);
        a = a + 10;
        ...
}
```



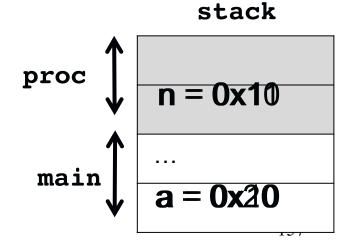
Pass by value: Pointer type



```
void proc(int n) {
    n++;
}
void main{
    int a=10;
    proc(a);
    a = a + 10;
...
}
```



```
void proc(int *n) {
    n++;
}
void main{
    int *a=0x10;
    proc(a);
    a = a + 0x10;
...
}
```



Pass by reference: Pointer



```
void proc(int *n){
        (*n)++;
 }
 void main{
        int a=10;
        proc(&a);
        a = a + 10;
            stack
proc
        n = 0x1000
 main
           a = 21
                          L3Info - C/Unix
```

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```
void proc(int **n){
           (*n)++;
    void main{
           int *a=0x10;
           proc(&a);
           a = a + 0x10;
            stack
proc
         n = 0x1000
main
                     0x1000
```

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Kahoot time! Q6- Q10

Passing arguments: Array



When an array is a function parameter, the compiler translates it as a pointer

```
void foo( int arr[]) { ... } \rightarrow void foo( int *arr) { ... }
```

- So, inside the function we cannot know the size of the array! All we have is a pointer ...
 - If you try to compute the array's size using sizeof(arr), you
 just get the size of a pointer,
 - x e.g. 4 in a code compiled for a machine with 32 bits as addresses
- It's important to pass the array size as a parameter.

```
void foo( int arr[], int size ) { ... }
```

Passing arguments: Array



- The parameter equals to the address of the 1^{er} element (arr=&arr[0])
- Since the parameter is the address of the array, we can modify the array elements
- The result is visible outside of the function

Example

Function to initialize the n values of an 1D array arr
void init_vector (int arr [], int n)
{
 int i;
 for (i=0;i<n;i++) { arr[i]=0; }
}</pre>

Passing arguments: 2D Arrays



- Array of arrays:
 - int arr[][COLS]
- Pointer to an array:
 - int (*arr) [COLS]

BUT NOT double pointer!

- int **arr
- compiler error because it cannot compute the size!
- compiler may NOT prevent you, BUT you may not have the expected results in some situations!

Passing arguments: Arrays



Kahoot time! Q13-Q15

The main function



- An executable can be called with a list of arguments
- ./prog arg1 arg2 ... argn
- These arguments are the parameters of the function main
- The prototype of main function:

```
int main(int argc, char* argv[])
```

- argc: size of argv (number of passed arguments + 1)
- argv: is a table of strings
 - x argv[0] is a char* pointing to the name of the program
 - **x** argv[1] is a char* pointing to the 1st argument
 - x argv[arc-1] is a char* pointing to the last argument

Remark

• If we want to use an argument as int, we have to convert it! int atoi(char* ch) of stdlib.h

Example



```
#include<stdio.h>
int main(int argc, char* argv[])
{
   int i;
   for (i=1;i<argc;i++)
   {
      printf("l'argument n°%d vaut %s ",i,argv[i]);
   }
}</pre>
```

What does prints execution of the command: ./prog foo 12 bar ?



Library functions: Inputs/Outputs, strings and files

Input / Output functions



- Several functions to manipulate I/O.
 - To use them we have to add their directive to the head of the compilation file using #include
 - We have to provide a format that indicates the type of the objects over whom we want to perform an I/O operation
- We have seen 2 basic I/O functions:
 - Write to the standard output (printf)
 - Read from the standard entry (scanf)

Read/Write a character



Read a character:

```
#include <stdio.h>
int getchar()
```

- Returns one character read from the standard input as an unsigned char cast to an int
- Returns EOF when the file ends or in case of an error

```
Write a character :
```

```
#include <stdio.h>
int putchar(int c)
```

Writes a character (an unsigned char) specified by the argument
 to the standard output

Example:

String manipulation



#include<string.h>

Read and write strings:

scanf(...)/printf(...),
gets(...)/puts(...),

fgets(...)/fputs(...), ...

- strlen(...), ...
- Comparison of strings: strcmp(...), ...
- Copy of strings: strcpy(...), ...
- Formatted output to a string: sprintf(...), ...
- **...**

Length of string :

String declaration: Remarks



Static declaration

```
char str1[]="hello";
char str2[]="world";
```

- We CANNOT assign a string to another (str2 = str1;). They are constant pointers!
 - We have to copy one to the other!
- Once the string is initialized, it CANNOT be assigned to another set of characters

```
str1 = "bye";
```

Basic I/O limitations with strings



- Basic I/O are NOT capable verifying bounds
- Example:

```
char name = "Hello"
scanf("%s", name );
printf("%s", name );
```



- lacksquare scanf("%s", name); ightarrow %s is dangerous!
 - name is a pointer, the starting address of the string
 - NO verification on how long is the information to be written!
 - Analogy: Give a pen to someone and point where to start writing
 But then have no control! he may continue writing on the desk...
 - Solution:

```
x char name[10];
x scanf("%9s", name);
```

Basic I/O limitations with strings



- scanf ends after a space:
 - receiving multiword strings in one variable
- Example

```
char name = "Hello John"
scanf("%s", name);
printf("%s", name);
```

- Solution?
 - gets(name);
 - puts(name);

Read/Write a string



Syntax:

```
char *gets(char *str)
```

Function:

- Reads a line from the stdin and stores it in the string pointed by str (stops when \n is read or EOF)
- It does NOT include the ending newline character \n in the resulting string
- It does NOT allow to specify a maximum size for name
 - **× DANGEROUS**

Returns:

- str on success
- NULL
 - × Error
 - **×** when **EOF** occurs, while no characters have been read.
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Read/Write a string



Syntax:

```
char *fgets (char *str, int size, FILE* file);
```

- Function:
 - Reads size 1 characters from input stream file and stores it in the string pointed by str
 - Stops whichever happens first:
 - x size 1 characters have been read
 - X A newline occurs, \n
 - The end-of-file is reached, EOF
 - It includes new line character `\n` pressed by the user
 - A terminating null character `\o'is automatically appended after the characters.
- Returns:
 - str on success
 - NULL :Error or when EOF occurs, while no characters have been read.
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Read/Write a string



- int sscanf (const char *str, const char * format, ...);
 - Like scanf but reads data from string pointed by str as the format specifies

And much other functions, check libraries documentation

Length of a string: strlen()



Syntax:

```
size_t strlen(const char *str);
```

- Function:
 - Counts the number of characters of the string pointed by str until it finds the terminating null character (`\0')
 - the terminating null character (`\0') is NOT INCLUDED in the length of the string
 - size_t is unsigned integer type of at least 16 bit
- Returns: the length of the string
- Example

```
char ch[50]="how long am I?";
printf("%d\n",strlen(ch));
```

Concatenate strings: strcat()



Syntax:

```
char *strcat(char *dest, const char *src) ;
```

Function

- appends the string pointed to by src to the end of the string pointed to by dest.
- the string pointed to by dest SHOULD be LARGE enough to hold both strings AND the terminating character
- The strings pointed by src and dest must NOT overlap!

Return value:

 The function strcat () returns normally a pointer to the destination string dest.

Copy strings: strcpy()



Syntax:

```
char *strcpy(char *dest, char *src);
```

Function

- The function strcpy() copies the string pointed by src (including) the character $' \setminus 0'$) into the string pointed by **dest**.
- The strings pointed by src and dest must NOT overlap!
- SECURITYISSUES It does **NOT** verify the **SIZE** of the strings before copying
- No knowledge of how large src is!
- src can be larger than dest

Return:

The function strcpy() returns normally a pointer to the destination string dest.

Solution?

strncpy(...)

Copy strings: strcpy():Example



```
#include <stdio.h>
#include <string.h>
void print(int i){
    char text1[10] = "Hippolyte";
    char text2[5] = "Ares";
    printf("The initial content of text1 is :%s \n", text1)
    if (i==0){ strcpy(text1,text2) ;
        printf("text1 has now :%s \n", text1);}
    else{
        strcpy(text2,text1);
        printf("text2 has now :%s \n", text2);}
                 void main(){
                     print(∅);
                     print(1);
                                                   179
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```

Copy strings: strncpy()



Syntax:

```
char *strncpy(char *dest, char *src, size_t n)
```

Function:

- Copies UP TO to n characters from the string pointed by src to dest.
- Length of src < n → the remaining part of dest is padded with null bytes.
- Length of src > n → no termination character at dest

Return:

 The function strncpy() returns normally a pointer to the destination string dest.

Compare strings: strcmp()



Syntax:

```
char strcmp (char *str1, char *str2);
```

- Function
 - The function strcmp() compares two strings of pointed by str1 and by str2.
- Return: The function strcmp() returns an integer:
 - if Return value < 0 then it indicates str1 is less than str2.
 - if Return value > 0 then it indicates str2 is less than str1.
 - if Return value = 0 then it indicates str1 is equal to str2.

Compare strings: strncmp()



Syntax:

```
int strncmp(const char *str1, const char *str2, size_t n)
```

- Function
 - The function strncmp() compares at most the first n number of characters between two strings of pointed by str1 and by str2.
- Return value: The function strncmp() returns an integer:
 - if Return value < 0 then it indicates str1 is less than str2.
 - if Return value > 0 then it indicates str2 is less than str1.
 - if Return value = 0 then it indicates str1 is equal to str2.

Compare strings: Example



```
#include<stdio.h>
#include<string.h>
void main(){
    char nom1[100] :
    char nom2[100] :
    int res;
    strcpy ( nom1, "C programming at L3 INFO - ISTIC" );
    strcpy ( nom2, "C programming is fun" );
    res = strcmp(nom1, nom2);
    if (res == 0) {
        printf("%s and %s are identical\n", nom1, nom2);
    } else if (res<0) {</pre>
        printf("%s is less than %s \n", nom1, nom2);
    } else {
       printf("%s is less than %s \n", nom2, nom1);
                                                     183
```

Compare strings: Example



```
#include <stdio.h>
#include <string.h>
int main(){
    char s1[20];
    fgets(s1, 20, stdin);
    strcmp(s1, "login");
    if (strcmp(s1, "login") == 0)
        printf("s2 = \"login\"\n");
    else
        printf("s2 != \"login\"\n");
    return 0;
```



Kahoot time! Q1 – Q10

I/O and files



- In C, the I/O are implemented using files
- There are a lot of predefined files
 - The standard input stdin
 - The standard output stdout
 - The standard error stderr
- We have already seen the basic formatted I/Os
 - printf (...): write to the file of the stdout
 - scanf (...) : read from the file stdin
- More global I/Os exist
 - Non formatted
 - Based on real files

File declaration



- The access to the files is performed through logical files
 - We have to use the directive #include <stdio.h>
 - We define them as type FILE
 - We use pointer to access them FILE*
- Example of declaring a logical file
 - FILE* file;
- Remark:
 - A logical file of a program has to be linked with an existing physical file in the file system
 - The link is created through the operations of opening and closing the logical files

File open: fopen



Syntax

```
FILE* fopen (char *name, char *mode );
```

- Function
 - Open a physical file called name with the access rights defined by the string mode:
 - "r" open for read
 - "w" open for write/create
 - "a" open for write at the end of the file
- Parameters
 - name : external name of the file to be opened
 - mode : mode and right of opening.
- Return value
 - A pointer over the logical file is returned. In case of an error the return pointer is NULL.

File close: fclose



Syntax

```
int fclose(FILE *file);
```

- Function
 - Close the logical file pointed out by the pointer passed as argument
- Return value
 - The function return 0 if the file is correctly closed
 - It returns the constant EOF, in case of an error
- ATTENTION: it is obligatory to close an opened file in mode write. Otherwise the content is NOT stored!

Read from a file



- int getc (FILE * file);
 - Read a character from the file pointed by file
 - Advances the position indicator for the file.
 - Returns:
 - * the character read as an unsigned char cast to an int
 - **X** the constant **EOF** in case of an error.
 - getchar() is equal to getc(stdin)
- int fscanf(FILE *file, char *format,
 list_args);
 - Works as scanf(), but the read is performed from the file pointed by file instead of the standard input

Read from a file



- int fread(void *buf, int size, int nb, FILE
 *file);
 - Read the packets of data stored to the file pointed by file and write them to a buffer pointed by the pointer buf
 - The parameter size gives the size of each packet
 - The parameter nb indicates the number of packets to write
 - Return the number of packets read from the file, 0 if we have reached the end of the file
 - Attention : the buffer pointer by buf should be large enough to store all the packets read!

Write to a file



- int putc (int c, FILE * file);
 - Write a character c (converted in unsigned char) in the file pointed by file
 - Advances the position indicator for the stream
 - Returns:
 - x returns the character written as an unsigned char cast to an int
 - * the constant value EOF in case of an error
 - putchar() is equivalent to putc(c, stdout).
- int fprintf(FILE *file, char *format, list_args);
 - Works like printf(), but the writing is performed into the file pointed by file instead of the standard output

Write to a file



- int fwrite(void *buf, int size, int nb, FILE
 *file);
 - Write the packest of data stored to the buffer pointed by the pointer
 buf in the file pointed by file
 - The parameter size gives the size of the packet
 - The parameter nb indicates the number of packets to write
 - Return the number of packets written to the file

Example



```
#include <stdio.h>
#include <string.h> /* strlen() */
#include <stdlib.h> /* exit() */
int main() {
    FILE *file;
    char *msg = "123456789012345";
    int n;
    file = fopen("tmp/test_write", "w");
    if (file!= NULL) {
        n = strlen(msg);
        int res = fwrite(msg, sizeof(char), n, file);
        if (res!=n) {
            fprintf(stderr,"Error during writing\n");
            exit(2);
        }
        fprintf(stdout,"Write successful\n");
        fclose(file);
    }else{
        printf("Couldn't open file for writing\n");
    return 0;
```



Dynamic memory allocation

Static & Automatic memory allocation UNIVERSITÉ



Static:

- allocated when the program starts
- Variable lifetime is the entire program

Automatic:

- allocated during execution (function calls)
- managed by the OS and the compiler
- Variable lifetime is the block where they are defined
- The variable size is fixed
- Usually stored on the stack

Limitations & inefficiencies



- Not always possible to know the exact number of array elements. May depend:
 - Parameters of program execution
 - Information provided by the user during execution
- Limited control over the lifetime of the variables
- Worst case: Stack overflow
 - Example: Recursive functions

Dynamic memory allocation



- The memory to store the program variables is allocated during the execution of the program
- Stored in the heap (large free pool of memory)
- The programmer controls:
 - the exact size (instantiated during execution)
 - Lifetime of memory locations

Worst-case: Memory leaks!

Problems: Dangling pointers!

. . .

Memory Leak



- Allocated blocks in the heap that are now unused and unreferenced!
- Improper use of dynamic memory → causes the memory consumption to be increased with the time
- In JAVA, garbage collector takes care of the not removed reserved parts in the heap!
- In C, the programmer is responsible for that!

For **EACH** allocation → We NEED a DEALLOCATION

Dangling pointer



- An object is deleted or de-allocated
- If the value of the pointer is not modified
 - the pointer still points to the memory location of the deleted/deallocated memory.

Dynamic memory: Functions



- Directive:
 - #include <stdlib.h>
- Allocation
 - malloc()
 - calloc()
 - realloc()
- Deallocation:
 - free()

Dynamic allocation: malloc()



Syntax:

```
void *malloc(size_t size);
```

Function

- Reserves a block with a size equal to size bytes in the heap
 x size_t is unsigned integer type of at least 16 bit
- The content of the memory block it is NOT initialized
- To calculate the required memory size to store an object we should use the function sizeof(type)
 - ★ The actual size of data object is compiler and target machine depended

Return value:

- pointer of void type with the address of the first reserved byte
 (void*)
- We have to explicitly cast the pointer to the type of data we want to store(int*, char*, ...)

Dynamic de-allocation: free(...)



Syntax:

```
void free(void *pointer);
```

- Function
 - free() frees the memory block pointed out by the pointer, which has been reserved by a call, such as malloc(), calloc() ou realloc()
 - If the pointer has not been obtained by one of these functions, or it has been already freed by free(), the behavior is UNDEFINED
 - If the pointer ptr is NULL, nothing happens
- No return value
- The C programmer is responsible for setting free the memory, when it is no more needed

Dynamic allocation: Example



```
#include<stdio.h>
#include<stdlib.h>
int main()
    int* p;
    int i;
    int taille = 100;
    p = (int*)(malloc (taille * sizeof(int)));
    if (p == NULL)
       printf("Allocation impossible:");
       exit(EXIT_FAILURE);
    for (i=0; i<taille;i++){ p[i]=0; }
    free(p);
```

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1D Array with size elements



```
type *arr = (type *)malloc(size * sizeof(type));
int *arr=(int *)malloc(4*sizeof(int));
```

	Heap
0x1000	xxx
0x1004	xxx
0x1008	xxx
0x100C	xxx
	xxx
0x1014	xxx
0x1018	xxx
	xxx
	xxx
0x1028	xxx
	xxx
	xxx

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1D Array with size characters -> string UNIVERSITÉ DE CA



char *arr=(char *)malloc(4*sizeof(char));

	Heap
0x1000	xxx
0x1004	xxx
0x1008	xxx
0x100C	xxx
	xxx
0x1014	xxx
0x1018	xxx
0x1028	XXX
	xxx
	xxx
	xxx
	xxx

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String declaration: Remarks



Dynamic declaration

```
char* s=(char*)malloc(sizeof(char) * SIZE);
s = "hello";
```

• We can assign a **char pointer** to another **char** pointer

```
char *p; p=s;
```

We can assign another set of characters

```
s="bye";
```

2D Array: Single pointer



XXX

```
type *arr= (type *)malloc(rows * cols * sizeof(type));
 int *arr=(int *)malloc(4*3*sizeof(int));
                                                 Heap
                                      0x1000
                                                  XXX
                                      0x1004
                                                  XXX
                                      0x1008
                                                  XXX
                                                  XXX
                                      0x100C
                                                  XXX
                                                  XXX
                                      0x1014
                                                  XXX
                                      0x1018
                                                  XXX
                                                  XXX
                                                  XXX
                                      0x1028
                                                  XXX
```

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2D Array: Array of pointers



```
type *arr[ROWS];
for (i = 0 ; i < ROWS; i++) {
  arr[i]=(type *)malloc(cols*sizeof(type));
                                                      Heap
                                         0x1000
                                                       XXX
                                         0x1004
                                                       XXX
                                         0x1008
                                                       XXX
                  Stack
                                                       XXX
                 arr[0]
     arr
                                         0x1010
                                                       XXX
                 arr[1]
                                                       XXX
                 arr[2]
                 arr[3]
                                         0x1028
                                                       XXX
                                                       XXX
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                                                       XXX
```

2D Array: Pointer to pointer



```
type ** arr;
arr=(type **)malloc(rows*sizeof(type*));
for (i = 0 ; i < rows ; i++) {
  arr[i]=(type *)malloc(cols*sizeof(type));
                                                     Heap
                                        0x1000
                                                     XXX
                                        0x1004
                                                     XXX
                                        0x1008
                                                     XXX
                 Stack
                                                    arr[0]
                                        0x1010
                                                    arr[1]
                   arr
                                                    arx[2]
                                                    axx[x3]
                                        0x1028
                                                     XXX
                                                     XXX
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                          L3Info – C/Unix
                                                     XXX
```

Dynamic Memory Allocation



Kahoot time! Q1 – Q7



User-defined data types: Enum, typedef, struct

Enumerated type declaration & definition NIVERSITÉ DE RENNES

Syntax:

```
enum identifier {element1,...}
```

- Function:
 - The enumeration creates a data type that groups a set of elements that behave as integral constants
 - Declares a new enumeration type with the name enum identifier
- A variable is defined as enum identifier, it can be assigned one of the elements
- Example:
 - enum day{mon, tue, wen, thu, fri, sat, sun};
 - **x** mon is coded by 0, tue by 1, ..., sun par 6
 - This is a new type called enum day
 - x enum day today;
 - x today=wen;

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Structure type declaration



Syntax

```
struct identifier {
   member definition;
   ...
};
```

NO space is reserved!

It only creates a new type that describes the members of the group!

Function:

- A structure creates a data type that groups items of possibly different types into a single type.
- Declares a new structure type named struct identifier
- The members of the structure can be integral, float, arrays, pointers.
- Example:

```
struct address{
    char street[100];
    int number;
};
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```

Structure type declaration



- The structure type declaration CANNOT be initialized!
- This is a new data type declaration, not a structure variable definition!
- Only when the structure variable is defined!
 - The corresponding memory is reserved
 - Now we can initialize the members of the structure

Structure variable definition & initialisation RENNES

Separately from the structure declaration:

```
struct address MyAddress = {"Paul Bert", 12};
```

During structure declaration:

```
struct address{
     char street[100];
     int number;
} MyAddress = {"Paul Bert", 12};
```

Space is reserved!

Equal to the one required to store the structure members aligned

Structure padding



- Architecture of a computer processor is made in such a way so it can read 1 word from memory at a time
 - e.g. 4 byte in 32 bit processor
- To use this processor advantage, the data are aligned in memory in words
 - e.g. multiples of 4 bytes in 32-bit processor
- One or more memory addresses may LEFT EMPTY to achieve this alignment!

Structure padding: Example



```
struct structure1
                                   struct structure1
        int id1;
                                           int id1;
        int id2;
                                           char c1;
        char c1;
                                           int id2;
        char c2;
                                           char c2;
        float ratio;
                                           float ratio;
};
                                   };
0x1000
                      4 bytes
                                     0x1000
             id1
                                                  id1
0x1004
                      4 bytes
                                     0x1004
             id2
                                              c1
0x1008
                                     0x1008
                                                  id2
                      1 byte 1 byte
            c2 rat
        c1
0x100C
                                     0x100C
                      4 bytes
                                              c2
          io atio
0x1010
                                     0x1010
                                                 ratio
```

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Structure variable: Access members UNIVERSITÉ DI



```
struct address{
   char street[100];
   int number;
};
```

Through a structure variable: .

```
struct address MyAddress;
strcpy(MyAddress.street, "Paul Bert");
MyAddress.number=12;
```

■ Through a pointer to the structure variable: →

```
struct address MyAddress;
struct address *MyAddressPtr = &MyAddress;
strcpy(MyAddressPtr → street, "Paul Bert");
MyAddressPtr → number=12;
```

Definition of a new type



Syntax

```
typedef type IDENTIFIER;
```

- Examples
 - typedef int LENGHT;
 - typedef enum {false, true} BOOLEAN;
 - typedef struct date date;
 - Use of typedef to avoid repetition of keyword struct

- Remark
 - There is no identifier for the enum in the definition of the type
 BOOLEEN

Structure type: Examples



```
struct date /* déclaration du type struct date */
{
    short j
    short m ;
    int a ;
};
```

```
/* déclaration du type struct UNE_DATE */
typedef enum {lun, mar, mer, jeu, ven, sam, dim} JOUR ;
typedef struct
{
    JOUR j ;
    struct date d ;
} UNE_DATE ;
```

Access struct members



Variable structure

```
UNE_DATE foo = { mer, {3 ,9 ,2014 } } ;
foo.j = mer ;
foo.d.a = 2014 ;
```

Structure pointer

```
UNE_DATE* queljour = &foo;
queljour->j = foo.j;
queljour->d.a = 2014;
```



Kahoot time! Q8 - Q12

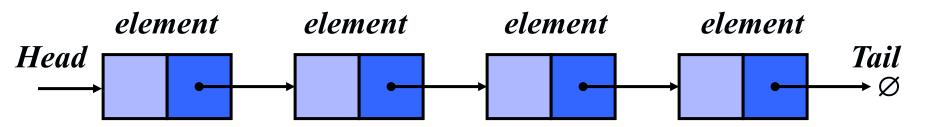


Data structures: Lists and Hash tables

Linked lists Vs arrays



- Linear data structure (like arrays)
- Linked list elements NOT stored in contiguous location
 - are linked using pointers



Why Linked list?



Array limitations:

- Fixed size: We must know the upper limit on the number of elements in advance
- Insert/Delete an element: Expensive
 - Create room by shifting existing elements!
 - x Ex.: Add new element 101 in an array with sorted ids
 id[]=[100, 103, 105, 200]

List Advantages:

- Dynamic size
- Ease of insertion/deletion

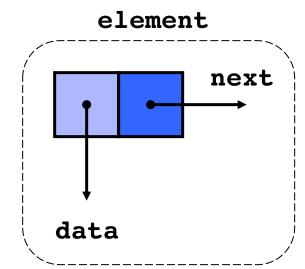
List Drawbacks:

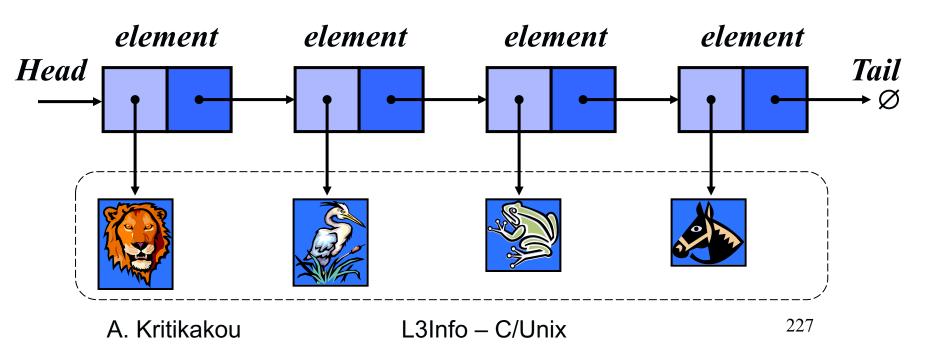
- Random access is not allowed
 - Access elements sequentially
- Extra memory space required for the pointer to each element

Linked list



- Each element has
 - data
 - next: Link to the next element
- In C, such an element can be represented using structures





Linked list: Example

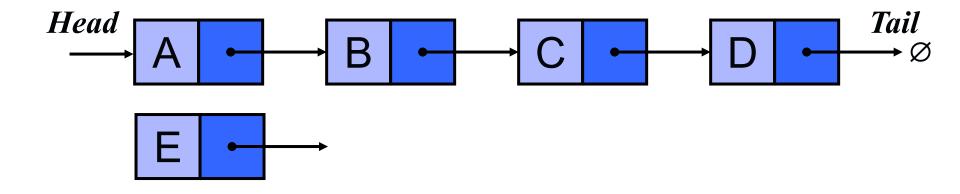


```
struct element {
   int data;
   struct element *next;
};
int main(){
  struct element *head = NULL;
  head=(struct element*)malloc(sizeof(struct element);
  head->data=1;
  head->next=NULL;
  return 0;
```

Operations: Insertion



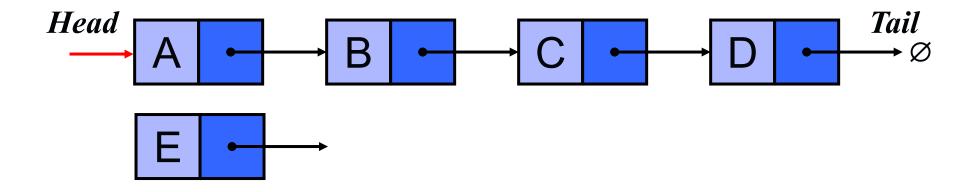
At the head:



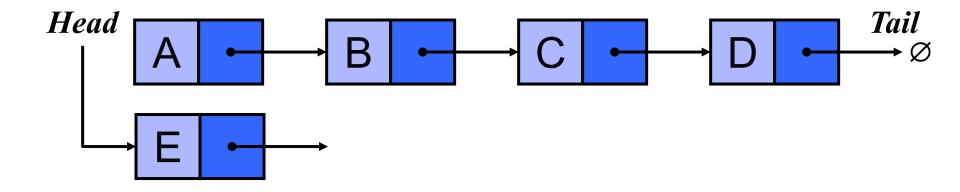
Operations: Insertion



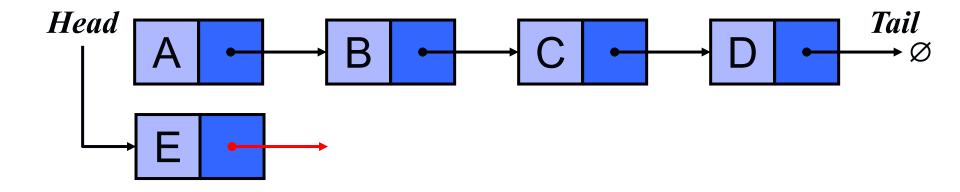
At the head:



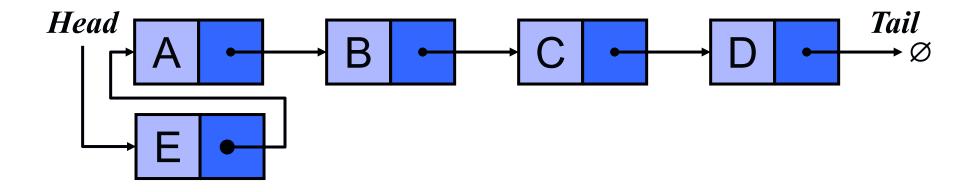




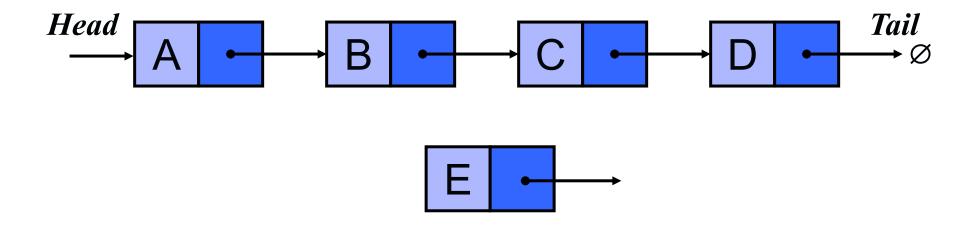




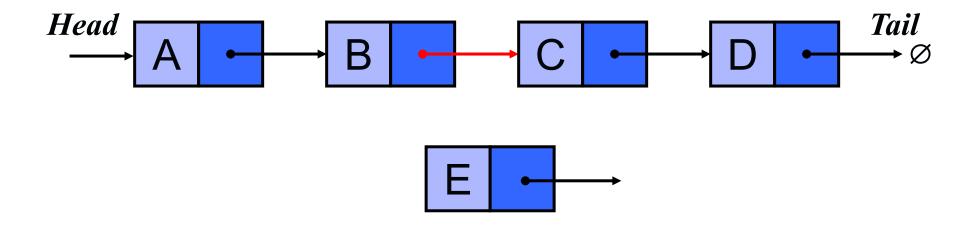




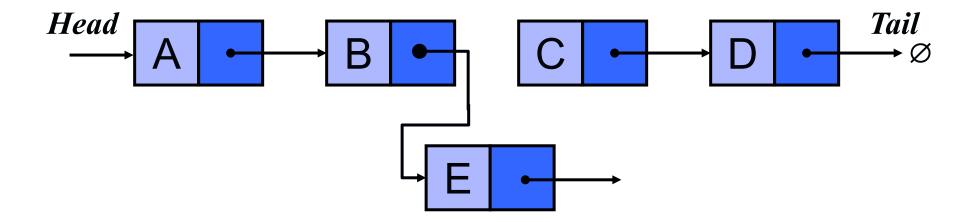




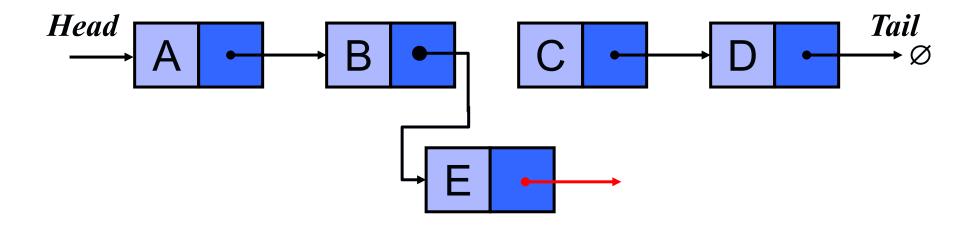




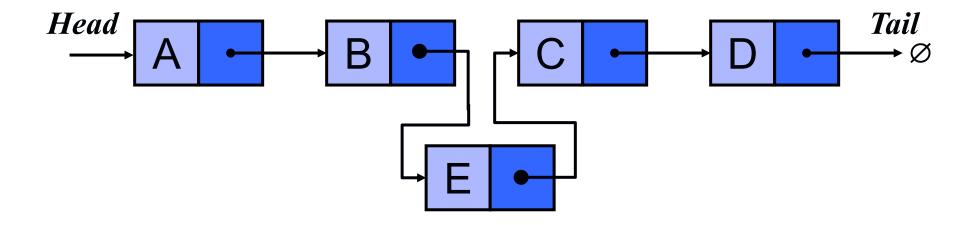




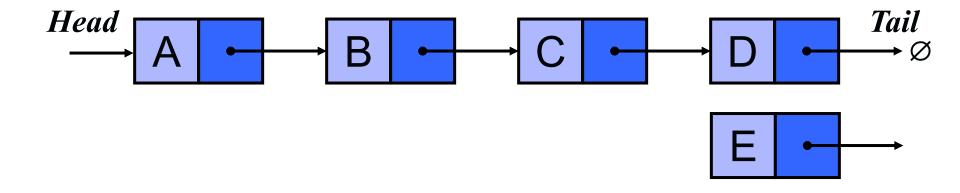




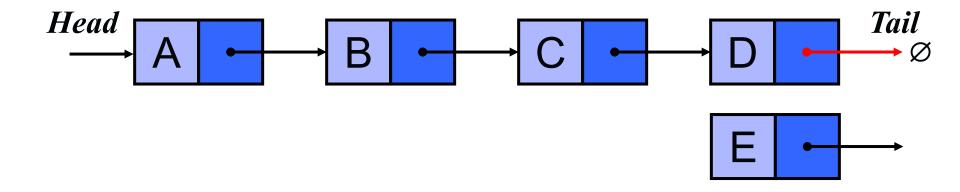




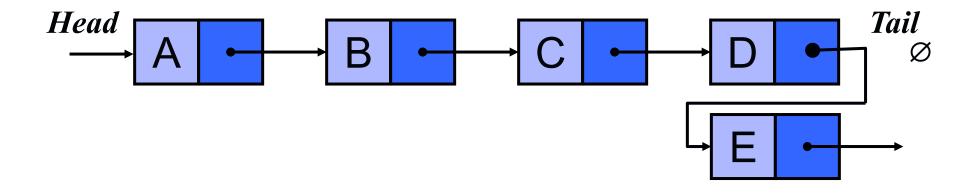




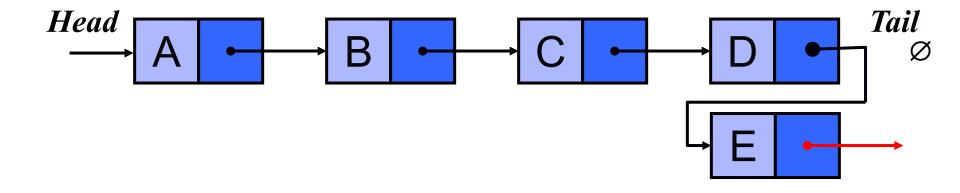




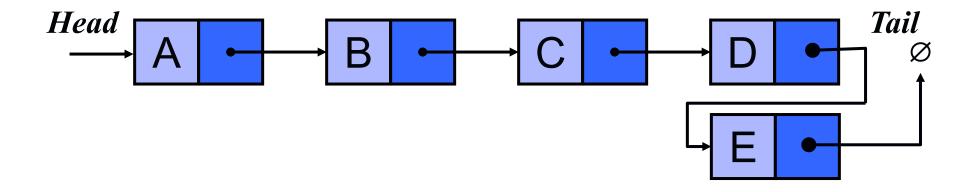






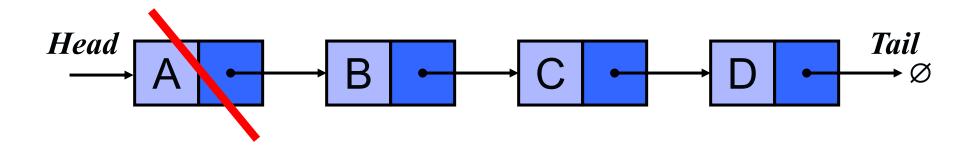






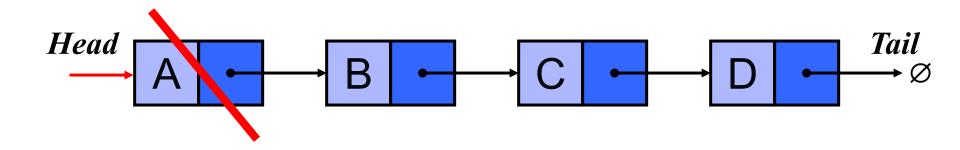
Operations: Deletion





Operations: Deletion

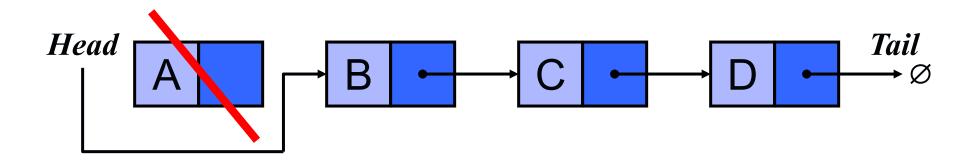




Operations: Deletion

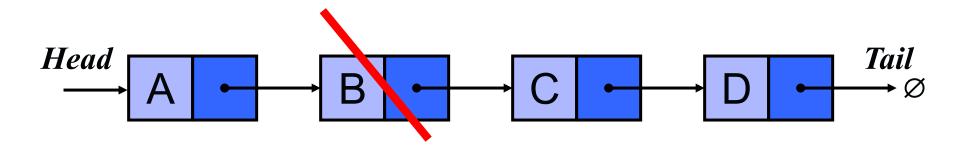


At the head:



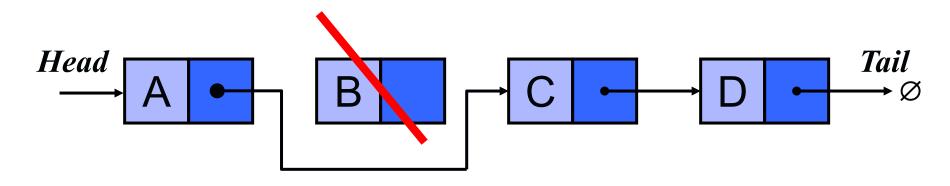
ATTENTION: Free the memory space used for element A





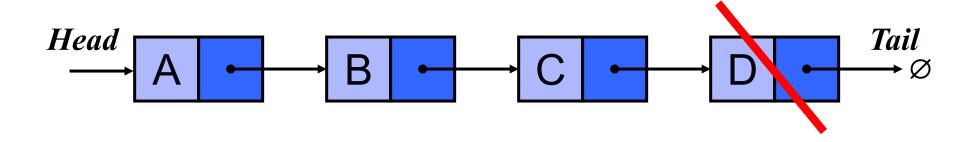


In the middle:



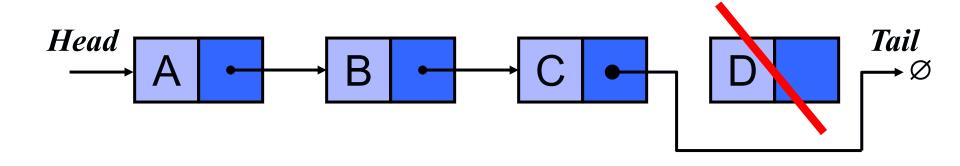
ATTENTION: Free the memory space used for element B







At the tail:



ATTENTION: Free the memory space used for element D

Doubly linked list



- Each element has
 - data
 - prev: Link to the previous element
 - next: Link to the next element

prev next data

Advantages

- Traversed both in forward and backward direction
- Delete operation is more efficient, previous pointer is given
 - ✗ Singly list: We have to traverse the list to find it

Disadvantages:

- Extra memory space for storing the previous pointer
- All operations have to maintain this extra pointer.

Hash Table



- An array of a fixed size hsize
- An element is associated with an integer value called hash value
 - a number within the range 0 to hsize 1
 - depicts in which cell of the hash table the element will be stored (key)
 - calculated using a Hash Function
 - Maps a data set of arbitrary size to a data set of fixed size
 - unique value

0	
1	
2	John
3	Mary
4 5	
5	
6	Helen
7	
8	Nick
9	

- Fast access of the data:
 - Search in the index indicated by the key

Hash Table: Some of its use



- Caches: Data tables used to speed up the data
- Object representation: Dynamic languages (Perl, Python, JavaScript) use hash tables to implement objects
- Algorithms to make computing faster
- Database indexing: In disk-based data tables and database indexes

Hash Function: Simple examples



- Integers:
 - Return key MOD hsize
- Strings:
 - We add up the ASCII values of the chars, and return the modulo operation

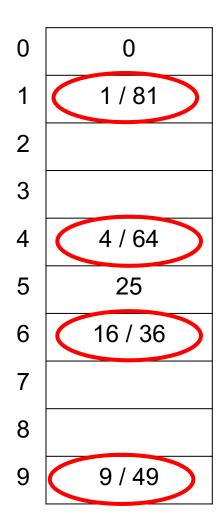
Collisions



When more than one element have the same hash value!

Example:

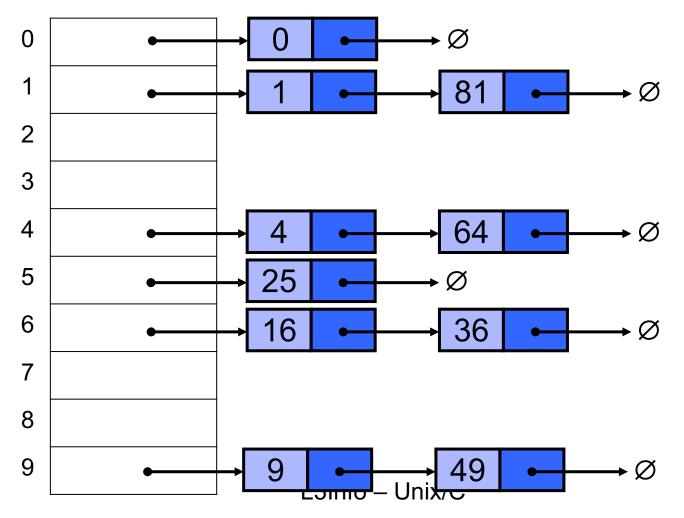
- hsize = 10
- Elements: the first 10 perfect squares
- 0, 1, 4, 9, 16, 25, 36, 49, 64, 81
- hash function: hash(x)=x mod 10



Collisions: Solution



- Separate chaining (open hashing):
 - Each element of the hash table is a linked list
 - Collision: Store both elements in the same linked list



Problems



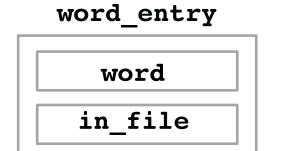
- The distribution of the values may NOT be even!
- Example
 - hsize = 10007
 - Variable: strings of 8 ASCII characters
 - The hash function: adds up the ASCII values of the chars and return the modulo operation
 - → it can only assume values 0 1016, which is 127*8



- Create a dictionary and search engine:
 - Read files word by word
 - Add to the dictionary one word found in a file and the number of times found in the file
 - Search for a word in the dictionary
- Use hash table and linked lists to store the words with the same key



```
struct word_entry
{
   char word[SIZE];
   int in_file;
   int times;
   struct word_entry *next;
};
typedef struct word_entry word_entry;
```



next

times

- The structure word_entry has as elements
 - a word of size SIZE
 - an integer in_file which is the identifier of the file where the word has been found
 - an integer times which shows how many times the word exist in this file
 - a pointer *next to a structure word_entry



```
struct word_list
{
    struct word_entry *first_word;
    struct word_entry *last_word;
};
typedef struct word_list word_list;

word_list
first_word

first_word

last_word

last_word

// Property word_list;
```

- The structure word_list has as elements
 - a pointer *first_word to a structure word_entry
 - a pointer *last_word to a structure word_entry



```
struct hash_table
{
word_list *htable;
int hsize;
};
typedef struct hash_table hash_table;
hash_table
hash_table
```

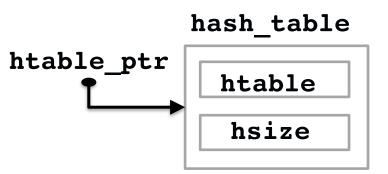
- The hash_table structure contains
 - a pointer htable to a type word_list
 A pointer to a pointer of the structure word_entry
 - the fixed size hsize
- We have finished the hash table declaration?



```
hash_table *htable_ptr;
htable_ptr = (hash_table*) malloc (sizeof(hash_table));
```

- Hash table allocation:
 - malloc reserves a memory space equal to the structure hash_table
 - htable_ptr will point to a structure containing an integer and a pointer to a word_list
- Hash table size initialisation htable_ptr->hsize = PRIME_NUMBER;
- We have finished the hash table declaration?







- Hast table pointer initialisation:
 - We need to allocate space for the linked lists.
 - We use an array of structures word_list equal to the table size

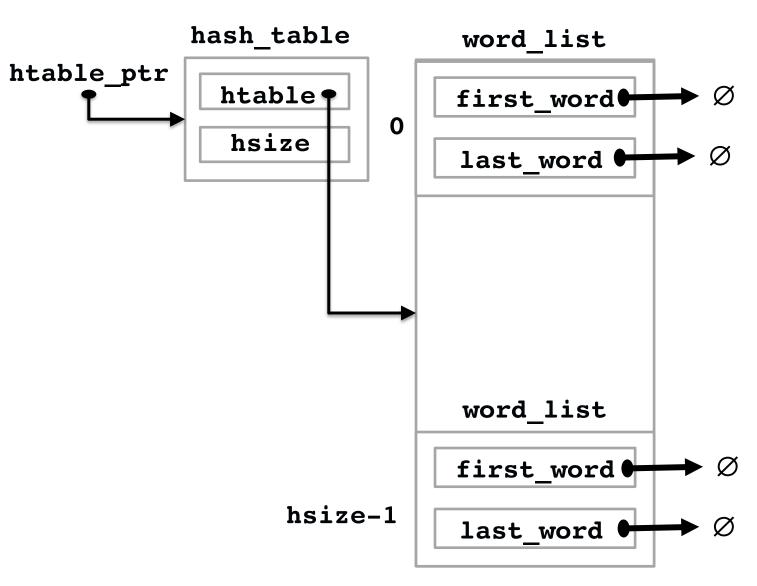
```
htable_ptr->htable =
  (word_list*) malloc(sizeof(word_list) * htable_ptr->hsize );
```

We need to initialize the pointers of the structures word_list to
 NULL

```
for(i = 0; i < htable_ptr->hsize ; i++ ){
  htable_ptr->htable[i].first_word = NULL;
  htable_ptr->htable[i].last_word = NULL;
}
```



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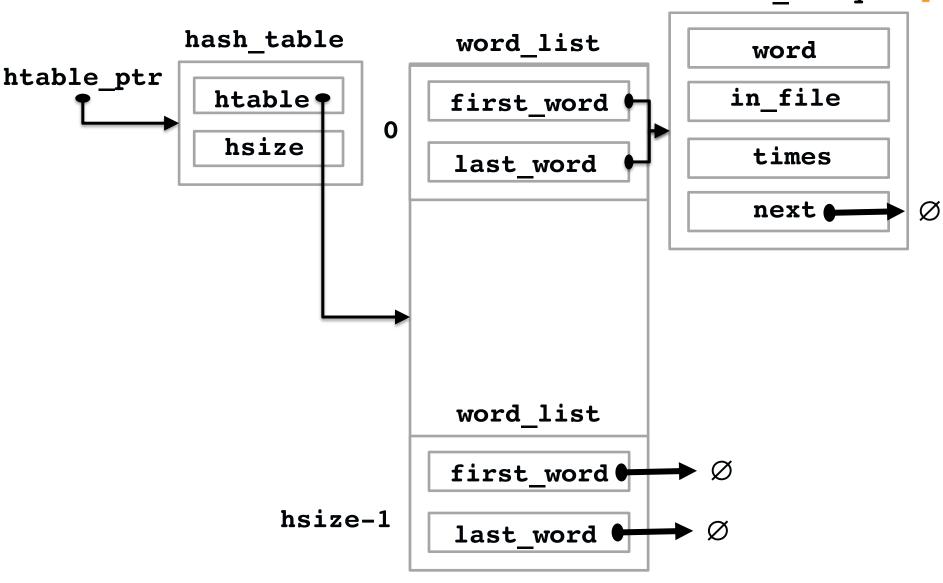
A. Kritikakou L3Info – C/Unix



- When a new word is added:
 - We need to allocate space for the word
 - Connected to the list based on the hash key

```
htable_ptr->htable[key].first_word =
          (word_entry*)malloc(sizeof(word_entry));
```





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TD2: C exercices

Exercice 1: Basic I/O



What prints the following program?

```
#include<stdio.h>
int main(){
  char nom[20] = "Balthazar";
  double solde = 1000.25;
   int age = 21;
  printf("Hello %s, you have %d years\n", nom, age);
  printf("Your age is written %x in hex\n", age);
  printf("Your name start with the letter %c\n", nom[0]);
  printf("Your bank account balance is %lf\n", solde);
  printf("The address of the table name is %p\n", nom);
```

Exercice 2: Pointers



Assume the following instructions

```
int x=0;
int y=0;
int* p=NULL;
p = &x;
x=2;
y=4;
*p = y;
x = 5;
y = *p;
p = &y;
```

Provide the values of x, y and p

Exercice 3: Pointers



Give the output of the program

```
#include <stdio.h>
void echange(int * a, int b) {
       int sauve; sauve = *a; *a = b; b = sauve;
void echange2(int * a, int ** b) {
       int sauve; sauve = *a; *a = **b; **b = sauve;
}
void echange3(int * a,int * b) {
       int sauve; sauve = *a; *a = *b; *b = sauve;
}
int main(){
       int a, b, d; int* c = &d; a = 1; b = 2; (*c) = 3;
       echange(\&a,b); printf("a= \&d, b= \&d\n",a,b);
       echange2(&a,&c); printf("a= %d, c= %d\n",a,*c);
       echange3(&a,&b); printf("a= %d, b= %d\n",a,b);
       echange3(&b,c); printf("b= %d, c= %d\n",b,*c);
       echange(\&a, \&b); printf("a= \&d, b= \&d\n",a,b);
       return 0;
```

Exercice 4: Strings



Write a function called reverse that reverses the characters of a string, without declaring a second array. The prototype of the function is the following:

```
int reverse(char* str) ;
```

Exercice 5: Strings and pointers



Give the output of the following program

```
#include <stdio.h>
#include <string.h>

char* mess="hello world\n"

int main() {
    strcpy(mess+4,mess) ;
    return 0;
}
```

Exercice 6: Linked lists



We have the structure list_elem_t for the list elements typedef struct s list {

The head of list is pointing by the pointer list_head list elem t* list head;

Write a function

list_elem_t* create_element(int val), which creates a new list element and initializes its value equal to the passed argument val and the next pointer to NULL

Exercice 6: Linked lists



Assume the function insert_head, which inserts a new element at the head of the list that has been passed as parameter

```
int insert_head(list_elem_t* 1, int val) {
    list_elem_t * nouveau = create_element(val);
    nouveau->next = 1;
    return 0;
}
```

- Is it correct?
 - What happens when insert_head(list_head, 1) is executed?
 - Correct the function!

Exercice 6: Linked lists



Write a function:

int insert_tail(list_elem_t** 1, int val), which creates a new list element at the tail of the list and return 0 in case of success and -1 in case of failure

Exercice 7: Files



- Write a main function that
 - Reads the name of a file (name1) from the keyboard
 - Reads the name of a file (name2) from the keyboard
 - Copies character by character the content of the file name1 to the file name2 and replace all the 'X' with 'Y'.
 - If the copy is not possible, prints an error message