

EuroHack15: Introduction to OpenACC

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Contents

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- What is OpenACC?
- How do GPUs work?
 - Will my code run well on a GPU using OpenACC?
- What does OpenACC look like?
- How do I use it?
 - The basic concepts
 - The basic directives
 - Advanced topics will be covered later in this tutorial
 - And in other training courses, like <u>this one</u>.
- Plus a few hints, tips, tricks and gotchas along the way
 - Not all guaranteed to be relevant, useful (or even true).

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Accelerator programming

- CRAY
- Why do we need a new GPU programming model?
- Aren't there enough ways to drive a GPU already?
 - CUDA (incl. NVIDIA CUDA-C & PGI CUDA-Fortran)
 - OpenCL
- All are quite low-level and closely coupled to the GPU
 - User needs to rewrite kernels in specialist language:
 - Hard to write and debug
 - Hard to optimise for specific GPU
 - Hard to port to new accelerator
 - Multiple versions of kernels in codebase
 - Hard to add new functionality
- Aside: OpenMP4accel is another directive-based model
 - Currently lacks some features of OpenACC
 - Compiler support also currently immeature
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Directive-based programming



Directives provide a high-level alternative

+ Based on original source code (Fortran, C, C++)

- + Easier to maintain/port/extend code
- + Users with OpenMP experience find it a familiar programming model
- + Compiler handles repetitive coding (cudaMalloc, cudaMemcpy...)
- + Compiler handles default scheduling; user tunes only where needed

Possible performance sacrifice

- Important to quantify this
- Can then tune the compiler
- Small performance sacrifice is acceptable
 - trading-off portability and productivity

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- after all, who handcodes in assembler for CPUs these days?

OpenACC Directives for Accelerators

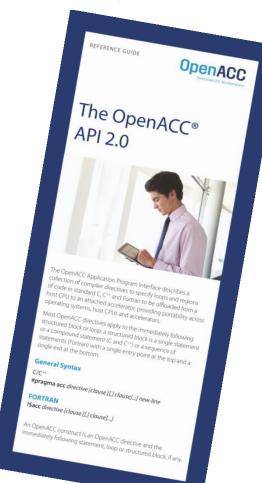


A common directive programming model for today's GPUs

- Announced at SC11 conference
- Offers portability between compilers
 - Drawn up by: NVIDIA, Cray, PGI, CAPS
 - Multiple compilers offer:
 - portability, debugging, permanence
- Works for Fortran, C, C++
 - Standard available at <u>openacc.org</u>
 - Initially implementations targeted at NVIDIA GPUs

Compiler support: all now complete

- Cray CCE: complete OpenACC 2.0 in v8.2
- PGI Accelerator: v12.6 onwards
- gcc: work started in late 2013
- Various other compilers in development









Accelerator directives



Modify original source code with directives

- Non-executable statements (comments, pragmas)
 - Can be ignored by non-accelerating compiler
 - CCE -hnoacc also suppresses compilation
- Sentinel: acc
 - C/C++: preceded by #pragma
 - Structured block {...} avoids need for end directives
 - Fortran: preceded by !\$ (or c\$ for FORTRAN77)
 - Usually paired with !\$acc end * directive
 - Directives can be capitalized
- Continuation to extra lines allowed
 - C/C++: \ (at end of line to be continued)
 - Fortran:
 - Fixed form: c\$acc& or !\$acc& on continuation line
 - Free form: & at end of line to be continued
 - continuation lines can start with either !\$acc or !\$acc&

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```
// C/C++ example
#pragma acc *
{structured block}
```

```
! Fortran example
!$acc *
<structured block>
!$acc end *
```

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Conditional compilation



- In theory, OpenACC code should be identical to CPU
 - only difference are the directives (i.e. comments)
- In practise, you may need slightly different code
 - For example, to cope with:
 - calls to OpenACC runtime API functions
 - where you need to recode for OpenACC
 - such as for performance reasons
 - you should try to minimise this
 - usually better OpenACC code is better CPU code
- CPP macro defined to allow conditional compilation

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- _OPENACC == yyyymm
 - Version 1.0: 201111
 - Version 2.0: 201306

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A first example

- Execute loop nest on GPU
- Compiler does the work:
 - Data movement
 - determines data use in loopnest
 - at start and end of loopnest:
 - allocates/frees GPU memory
 - moves data to/from GPU
 - Synchronisation

#pragma acc parallel loop
for (j=0; j<N; j++) {
 for (i=1; i<N-1; i++) {
 c[j][i] = a[j][i] + b[j][i]
 }
}
write-only</pre>
read-only

- Loop "schedule": spreading loop iterations over threads on GPU
 - OpenACC will "partition" (workshare) more than one loop in a loopnest
 - compare this to OpenMP, which only partitions the outer loop
- Caching (e.g. explicit use GPU shared memory for reused data)
 - automatic caching can be important
- User can tune all default behavior with optional clauses on directives

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Accelerator kernels



- We call a loopnest that will execute on the GPU a "kernel"
 - this language is similar to CUDA
 - the loop iterations will be divided up and executed in parallel
- We have choice of two directives to create a kernel
 - parallel loop or kernels loop
 - both generate an accelerator computational task from a loopnest
 - also known as a "kernel"
 - the language is confusing
- Why are there two and what's the difference?
 - You can use either
 - or both, in different parts of the code
 - This tutorial concentrates on using the parallel loop directive

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A first full OpenACC program: "Hello World"

```
PROGRAM main
  INTEGER :: a(N)
  <stuff>
!$acc parallel loop
 D0 i = 1,N
  a(i) = i
  ENDDO
!$acc end parallel loop
!$acc parallel loop
 DO i = 1, N
  a(i) = 2*a(i)
  ENDDO
!$acc end parallel loop
  <stuff>
END PROGRAM main
```

- Two accelerator parallel regions
 - Compiler creates two kernels
 - Loop iterations automatically divided across GPU threads
 - First kernel initialises array
 - Compiler will determine a is write-only
 - Second kernel updates array
 - Compiler will determine a is read-write
 - Breaking parallel region=barrier
 - No barrier directive (global or within SM)

- Note:
 - Code can still be compiled for the CPU

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Data scoping

- Codes process data, using other data to do this
 - all this data is held in structures, such as arrays or scalars
- In a serial code (or pure MPI), there are no complications
- In a thread-parallel code (OpenACC, OpenMP etc.)
 - Things are more complicated:
 - Some data will be the same for each thread (e.g. the main data array)
 - The threads can (and usually should) share a single copy of this data
 - Some data will be different (e.g. loop index values)
 - Each thread will need it's own private copy of this data
- Data scoping arranges this. It is done:
 - automatically (by the compiler) or explicitly (by the programmer)
- If the data scoping is incorrect, we get:
 - incorrect (and inconsistent) answers ("race conditions"), and/or
 - a memory footprint that is too large to run

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Understanding data scoping

- Data scoping ensures the right answer
 - We want the same answer when executing in parallel as when serially

Declare variables in parallel region to be shared or private

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- shared
 - all loop iterations process the same version of the variable
 - variable could be a scalar or an array
 - a and b are shared arrays in this example

private

- each loop iteration uses the variable separately
- again, variable could be a scalar or an array
- t is a private scalar in this example
- loop index variables (like i) are also private
- firstprivate: a variation on private
 - each thread's copy set to initial value
 - loop limits (like N) should be firstprivate

```
for (i=0; i<N; i++) {
    t = a[i];
    t++;
    b[i] = 2*t;
}</pre>
```

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Data scoping in OpenACC (and OpenMP)



- In OpenMP, we have exactly these data clauses
 - shared, private, firstprivate

In OpenACC

- private, firstprivate are just the same
- shared variables are more complicated in OpenACC
 - because we also need to think about data movements to/from GPU
- We sub-classify shared variables by how they are used on the GPU:
- copyin: a shared variable that is used read-only by the GPU
- copyout: a shared variable that is used write-only
- copy: a shared variable that is used read-write

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a shared variable that is a temporary scratch space create: (although there is still an unused copy on the host in this case)

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Data scoping with OpenACC



parallel regions:

- scalars and loop index variables are private by default
- arrays are shared by default
 - the compiler chooses which shared-type: copyin, copyout, etc.
- explicit data clauses over-ride automatic scoping decisions
- You can also add the default(none) clause

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then you have to do everything explicitly (or you get a compiler error)

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A more-explicit first version

```
PROGRAM main
  INTEGER :: a(N)
 <stuff>
!$acc parallel loop copyout(a)
 D0 i = 1,N
  a(i) = i
  ENDDO
!$acc end parallel loop
!$acc parallel loop copy(a)
 DO i = 1,N
  a(i) = 2*a(i)
  ENDDO
!$acc end parallel loop
  <stuff>
END PROGRAM main
```

- We could choose to make the data movements explicit
 - maybe because we want to
 - maybe also use default(none) clause
 - or maybe compiler is overcautious

- Note:
 - Array a is needlessly moved from/to GPU between kernels

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- You could call this "data sloshing"
- This will have a big impact on performance

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OpenACC data regions



Data regions allow data to remain on the accelerator

- e.g. for processing by multiple accelerator kernels
- specified arrays only move at start/end of data region

Data regions only label a region of code

- they do not define or start any sort of parallel execution
- just specify GPU memory allocation and data transfers
- can contain host code, nested data regions and/or device kernels

Be careful:

- Inside data region we have two copies of each of the specified arrays
- These only synchronise at the start/end of the data region
 - and only following the directions of the explicit data clauses
- Otherwise, you have two separate arrays in two separate memory spaces
- Unified memory systems will have only one copy!

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Defining OpenACC data regions

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- Two ways to define data regions:
 - Structured data regions:
 - Fortran: !\$acc data [data-clauses] ... !\$acc end data
 - C/C++: #pragma acc data [date-clauses] {...}
 - Unstructured data regions (new in OpenACC v2):
 - Fortran: !\$acc enter data [data-clauses] ... !\$acc exit data [data-clauses]
 - C/C++: #pragma enter data [data-clauses] ... #pragma exit data [data-clauses]
- For most "procedural code", use structured data regions
- Unstructured data regions
 - Useful for more "Object Oriented" coding styles, e.g.
 - Separate constructor/destructor methods in C++
 - Separate subroutines for malloc (or allocate) and free (or deallocate)
- A data region with no data clauses is pointless
 - that is, it is redundant (and does nothing)

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A second version

```
PROGRAM main
  INTEGER :: a(N)
  <stuff>
!$acc data copyout(a)
!$acc parallel loop
 DO i = 1,N
  a(i) = i
 ENDDO
!$acc end parallel loop
!$acc parallel loop
 DO i = 1,N
   a(i) = 2*a(i)
  ENDDO
!$acc end parallel loop
!$acc end data
  <stuff>
END PROGRAM main
```

- Now added a data region
 - Specified arrays only moved at boundaries of data region
 - Unspecified arrays moved by each kernel
 - No compiler-determined movements for data regions
- Data region can contain host code and accelerator regions
- Copies of arrays independent
- No automatic synchronisation within data region
 - User-directed synchronisation possible with update directive

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Data scoping with OpenACC (2)

parallel regions:

- scalars and loop index variables are firstprivate by default
- arrays are shared by default
 - the compiler chooses which shared-type: copyin, copyout, etc.
- explicit data clauses over-ride automatic scoping decisions
 - You can also add the default(none) clause
 - then you have to do everything explicitly (or you get a compiler error)

data regions:

- only shared-type scoping clauses are allowed
- there is NO default/automatic scoping
- un-scoped variables on data regions
 - will be scoped at each of the enclosed parallel regions
 - automatically, unless the programmer does this explicitly
 - this probably leads to unwanted data-sloshing of large arrays
- Using data region scoping in enclosed parallel regions:
 - same routine: omit scoping clauses on enclosed parallel directives
 - different routine: use present clause on enclosed parallel directives

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Sharing GPU data between subprograms

```
PROGRAM main
   INTEGER :: a(N)
   <stuff>
!$acc data copyout(a)
!$acc parallel loop
   DO i = 1,N
    a(i) = i
   ENDDO
!$acc end parallel loop
   CALL double_array(a)
!$acc end data
   <stuff>
END PROGRAM main
```

```
SUBROUTINE double_array(b)
  INTEGER :: b(N)
!$acc parallel loop present(b)
  DO i = 1,N
   b(i) = double_scalar(b(i))
  ENDDO
!$acc end parallel loop
END SUBROUTINE double_array
```

```
INTEGER FUNCTION double_scalar(c)
  INTEGER :: c
  double_scalar = 2*c
END FUNCTION double_scalar
```

- present clause uses GPU version of b without data copy
 - Original calltree structure of program can be preserved
- One kernel is now in subroutine (maybe in separate file)
 - OpenACC 1.0: function calls inside parallel regions required inlining
 - OpenACC 2.0: compilers support calls and nested parallelism

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Reduction variables



- Reduction variables are a special case of shared variables
 - where we will need to combine values across loop iterations
 - e.g. sum, max, min, logical-and etc. acting on a shared array
- We need to tell the compiler to treat this appropriately
 - Use the reduction clause for this (added to parallel loop directive)
 - same expression in OpenACC as in OpenMP
 - Examples:
 - sum: use clause reduction(+:t)
 - Note sum could involve adding and/or subtracting
 - max: use clause reduction(max:u)

```
D0 i = 1, N
   t = t + a(i) - b(i)
   u = MAX(u,a(i))
ENDDO
```

- Note: OpenACC only allows reductions of scalars
 - not of array elements
 - advice:
 - try rewriting to use a temporary scalar in the loopnest for the reduction

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Data scoping gotchas: OpenACC vs. OpenMP



- In OpenACC parallel regions:
 - scalars <u>and</u> loop index variables are firstprivate by default
- Compare this to OpenMP parallel regions:
 - loop index variables are private by default, but scalars are shared
- Be careful of this, especially:

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- if you program (separately) using the two programming models, or
- if you are translating an OpenMP code to OpenACC

Minimising data movements



- This is the single-biggest OpenACC optimisation for GPUs
- There are three techniques:
 - Keep data on the GPU as long as possible
 - use data regions and port all enclosed loopnests (as we have seen)
 - Only move arrays when you need to
 - using the update directive
 - Only move the data you need to move
 - using array sections

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The update directive



Data regions keep data on device

- can span multiple compute kernels and serial (host) code
- create <u>copies</u> of data arrays on device for duration of data region
- host, device copies only synchronised at start/end of data region
 - as requested by explicit data clauses

You can synchronise copies manually within a data region

- for instance:
 - to copy a halo buffer back to the host for communication
 - to copy values of an array to the CPU for checking or printing

You do this using the update directive, for instance:

- update host(a) copies entire array a from device to host
 - OpenACC 2.0: can use self instead of host
- update device(a) copies entire array a from host to device

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Array sections



- array section notation allows this, using ":" notation
- syntax differs slightly between languages
 - Fortran uses start:end, so first N elements is a(1:N)
 - C/C++ uses start:<u>length</u>, so first N elements is b[0:N]
 - Advice: be careful when switching languages!
 - Use profiler, CRAY_ACC_DEBUG commentary to see how much data moved

Sections allowed in data clauses and with update

For multi-dimensional arrays

- specified sections must be a contiguous block of memory
- can only specify one incomplete section on slowest-moving index

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- Fortran: slowest index is right-most
 - so a(1:N,2:N-1) is allowed, but a(2:N-1,1:N) is not
- C/C++: slowest index is left-most
 - so b[1:N-2][0:N] is allowed, but b[0:N][1:N-2] is not

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Unshaped pointers



"Unshaped pointer" compiler/runtime errors

- Fortran arrays have a complicated data structure
 - includes a descriptor that contains information about size and shape
 - the compiler therefore knows how much data to transfer
- C/C++: arrays are often just pointers
 - especially if the arrays were dynamically allocated or passed by reference
 - How many bytes should be transferred here: copy(c)?
 - So you usually need to be more explicit
 - You need to put in data clauses
 - And specify the slicing (even if this is the whole array): copy(c[0:N])

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Sharing data between kernels



- Must ensure that the runtime uses the shared data already present
- If the kernel is in the same routine as the data region
 - just don't mention those bits of data in the parallel/kernels clauses
- If the kernel is in a different routine to the data region
 - On the parallel or kernels directive:
 - Specify the relevant data with present clause
 - instead of other shared clauses (copy, copyin, copyout, create)
 - don't rely on automatic scoping for shared data in this case
- If an array is declared present, but is not on accelerator
 - you get a runtime error and the program crashes
 - This is usually what you would like to happen
 - rather than running to completion with the wrong answer
 - But there are other ways to do things...

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Directives in summary

Compute regions

created using parallel loop or kernels loop directives

Data regions

created using data or enter/exit data directives

Data clauses are applied to:

- accelerated loopnests: parallel and kernels directives
 - here they over-ride relevant parts of the automatic compiler analysis
 - you can switch off all automatic scoping with default(none) clause (in v2)
- data regions: data directive (plus enter/exit data in OpenACC v2)
 - There is no automatic scoping in data regions (arrays or scalars)
- Shared clauses (copy, copyin, copyout, create)
 - supply list of scalars, arrays (or array sections)
- Private clauses (private, firstprivate, reduction)
 - only apply to accelerated loopnests (parallel and kernels directives)
- present clause (used for nested data/compute regions)

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And take a breath...



- You now know everything you need to start accelerating
 - You can successfully port a lot of codes just knowing this much
 - The performance at this stage isn't bad, either
 - you can often beat the CPU version of the code running across all the cores
- So what is the rest of OpenACC (and this tutorial) for?
 - Some codes require more functionality to port
 - OpenACC also has a lot of performance tuning options
- The emphasis in this introduction has been on
 - explaining data scoping and using data regions
- Why?
 - because optimising data movements is far more important than tuning
 - minimising data transfers typically speeds up GPU execution by 10x-100x
 - performance tuning maybe gains you 2x-3x

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and you can't start to get this until you first stop data-sloshing