After doing this tutorial you should be able to:

- 1. use the recursive definition to tell if something is a Boolean formula or not,
- 2. evaluate a given Boolean formula for a given assignment,
- 3. tell if a given Boolean formula is satisfiable/valid,
- 4. model facts and reasoning in Boolean logic.

**Problem 1.** Are the following expressions propositional logic formulas (according to the definition of the syntax, the conventions, as well as the extended definition of the syntax, given in the lecture)?

- 1.  $(p \land q)$
- **2.** (*p*)
- 3.  $p \wedge q$
- 4.  $\neg(p \land q)$
- 5.  $(p \oplus q)$
- 6.  $(p \rightarrow q)$
- 7.  $(p \rightarrow (q \lor r))$
- 8.  $\neg\neg\neg\neg\neg p$

**Problem 2.** A formula *G* is **valid** if every assignment satisfies *G*. Valid formulas capture "logical truths". Give some examples of valid formulas.

**Problem 3.** This is Exercise 7.10 in Russel and Norvig's *Artificial Intelligence* (2010).

Decide whether each of the following sentences is valid, unsatisfiable, or neither. Verify your answer using truth-tables (or any method you like).

- 1.  $Smoke \rightarrow Smoke$
- 2.  $Smoke \rightarrow Fire$
- 3.  $(Smoke \rightarrow Fire) \rightarrow (\neg Smoke \rightarrow \neg Fire)$
- 4.  $Smoke \lor Fire \lor \neg Fire$
- 5.  $((Smoke \land Heat) \rightarrow Fire) \leftrightarrow ((Smoke \rightarrow Fire) \lor (Heat \rightarrow Fire))$
- 6.  $(Smoke \rightarrow Fire) \rightarrow ((Smoke \land Heat) \rightarrow Fire)$
- 7.  $Big \lor Dumb \lor (Big \rightarrow Dumb)$

**Problem 4.** Suppose you had code that solved the satisfiability problem. Show how you could call that code to solve the validity problem, i.e., to test if a given formula is valid or not.

**Problem 5.**[Exam 2022] Is it true that if the formula  $F \vee G$  is valid then either F is valid or G is valid? Give a short explanation/justification of your answer.

**Problem 6.**[Assignment problem in 2021] Each of the following statements about propositional logic is false. For each statement, show why the statement is false with a counterexample using only the variables p, q (you don't need to use both variables).

- 1. If *F* is valid and *G* is satisfiable then  $(F \rightarrow G)$  is valid.
- 2. If *F* is satisfiable and *G* is satisfiable then  $(F \land G)$  is satisfiable.
- 3. If *F* is valid then  $\neg F$  is satisfiable.
- 4. If *F* is satisfiable then  $\neg F$  is satisfiable.

**Problem 7.** Write a recursive function SubForms(G) that returns the set of subformulas of G (i.e. in a similar style to the way that the truth value function tv is defined in the lecture). Do this for the basic syntax, and for the extended syntax.

**Problem 8.** Intuitively, the size of a propositional formula F is the number of symbols in F, ignoring parentheses. So, e.g.,  $|(p \lor \neg p)| = 4$ .

- 1. Give a recursive procedure defining the size of F, denoted |F|.
- 2. Show by induction that the number of subformulas of F, denoted  $|\mathbf{SubForms}(F)|$ , is at most |F|. Hint: use induction on the size of F.