

# Mobile Robots

## Assembly and Control

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Mathias Hudoba de Badyn

Advanced Topics in Control  
May 16, 2022

**ETH** zürich

AUTOMATIC  
CONTROL  
LABORATORY **ifa**

## Announcements

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1. Homework 3 is extended, Vahid is coming at 17:45 to answer questions

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2. Project ideas should be finalized

Recap from last week

Princeton Series in APPLIED MATHEMATICS

# Graph Theoretic Methods in Multiagent Networks



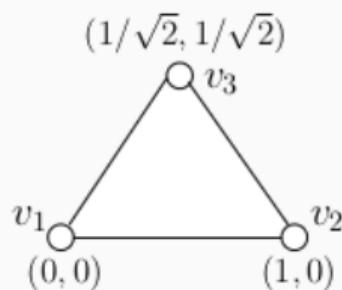
Mehran Mesbahi  
and Magnus Egerstedt

# Recap of Rigidity

## Desired Interagent distances

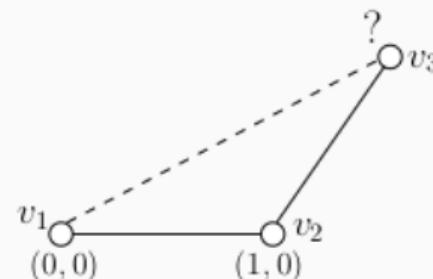
$$D = \{d_{ij} \in \mathbb{R} : d_{ij} > 0; i, j = 1, \dots, n, i \neq j\}$$

$D$  must be **feasible**: there exist points  $\Xi := \{\xi_1, \dots, \xi_n \mid \xi_i \in \mathbb{R}^p\}$  such that  $\|\xi_i - \xi_j\| = d_{ij}$ ,



$$D = \{d_{12} = d_{13} = d_{23} = 1\}$$

(a) Feasible formation



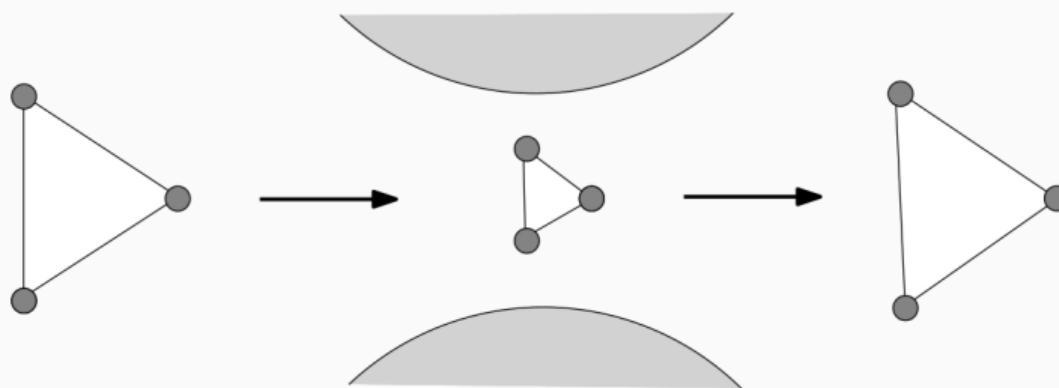
$$D = \{d_{12} = d_{23} = 1, d_{13} = 3\}$$

(b) Infeasible formation

## Recap of Rigidity

We refer to the **scale invariant** formation  $D$  as any set of distances  $D'$  such that  $D' = \alpha D$  for any  $\alpha \in \mathbb{R}_+$ .

Used in 'cluttered' environments where strict interagent distances are not required.



## Recap of Rigidity

Interagent distances

$$D = \{d_{ij} = d_{ji} \geq 0, i, j = 1, \dots, n, i \neq j\}$$

<u>formation</u>	<u>specification</u>	<u>interpretation</u>
scale invariant	$D$	$\ x_i - x_j\  = \alpha d_{ij}$ for some $\alpha > 0$
rigid	$D$	$\ x_i - x_j\  = d_{ij}$
translational invariant	$\Xi$	$x_i = \xi_i + \tau$ for some $\tau \in \mathbf{R}^p$

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where

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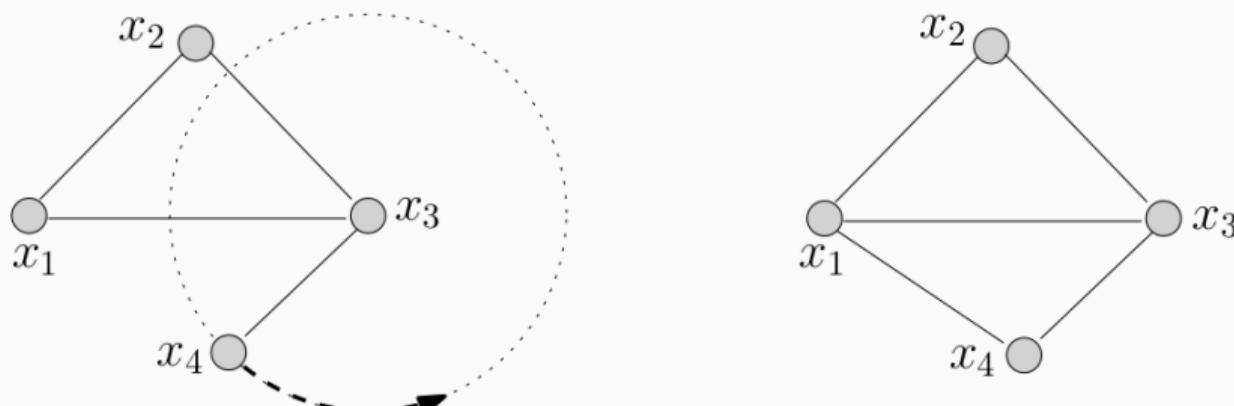
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- is called **edge consistent** if  $\|x_i(t) - x_j(t)\|$  is constant for all  $ij \in \mathcal{E}, t \in \mathbb{R}_+$ , or
- **rigid**, if  $\|x_i(t) - x_j(t)\|$  is constant for all  $i \neq j$ .

## Definition

A framework  $(\Xi, \mathcal{G}_f)$  is called **rigid** if and only if all edge-consistent trajectories of the framework are rigid trajectories.

Otherwise, the framework is called flexible.



**Left:** Flexible framework with edge-consistent trajectories.

**Right:** Rigid framework

## Definition

A framework  $\mathcal{G}(\Xi)$  is **infinitesimally rigid** if  $R(\mathcal{G}(\Xi))u = 0$  for all  $u$  satisfying

$$(u_i(t) - u_j(t))^T (x_i(t) - x_j(t)) = 0 \text{ for all } ij \in \mathcal{E}$$

Some properties of  $R(\mathcal{G}(\Xi))$ :

- Has  $|\mathcal{E}|$  rows
- Has  $pn$  columns ( $p$  is the state dimension)
- Is an algebraic representation of a graph plus the points  $\Xi$

## Theorem

If  $\mathcal{G}$  is rigid (i.e., there exists at least one  $\Xi$  such that  $\mathcal{G}(\Xi)$  is infinitesimally rigid), then the set of all  $\Xi$  such that  $\mathcal{G}(\Xi)$  is infinitesimally rigid is a dense, open subset of  $\mathbb{R}^{pn}$ .

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2. Feasible configurations can be approximated arbitrarily well (remember what dense + open means)

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2. Shape-based control (controlling to target  $\Xi$ )
3. Relative state-based control (controlling to  $\|x_i - x_j\| = d_{ij}$ )

## Formation Specification via Relative States

---

$$z(t) = \begin{bmatrix} x_1(t) - x_2(t) \\ x_2(t) - x_3(t) \end{bmatrix} \in \mathbb{R}^4$$

This specifies **all** relative formation states, including

$$x_1(t) - x_3(t) = [x_1(t) - x_2(t)] + [x_2(t) - x_3(t)]$$

## Formation Specification via Relative States

We can also set a formation center  $x_o(t)$  and control the states,

$$z(t) = \begin{bmatrix} x_o(t) - x_1(t) \\ x_1(t) - x_2(t) \\ x_2(t) - x_3(t) \end{bmatrix} \in \mathbb{R}^6$$

## Formation Specification via Relative States – Incidence Matrix

$$z(t) = \begin{bmatrix} x_0(t) - x_1(t) \\ x_1(t) - x_2(t) \\ x_2(t) - x_3(t) \end{bmatrix} \in \mathbb{R}^6$$

$$x(t) = \begin{bmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \end{bmatrix}$$
$$D(\mathcal{D}) = \begin{bmatrix} 1 & 0 \\ -1 & 1 \\ 0 & -1 \end{bmatrix} \otimes I_2$$

## Formation Specification via Relative States – RSS

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This is called the **relative state specification (RSS)**.

## Formation Specification via Relative States – RSS Example

## Equivalence/Converting between RSS's

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### Proposition

Let  $\mathcal{D}_j$  be a directed spanning tree RSS and let  $\mathcal{D}_d$  be an arbitrary RSS. Then,

$$T_{dj} = D(\mathcal{D}_d)^T D(\mathcal{D}_j) [D(\mathcal{D}_j)^T D(\mathcal{D}_j)]^{-1}$$

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$$T_{dj} = D(\mathcal{D}_d)^T D(\mathcal{D}_j) [D(\mathcal{D}_j)^T D(\mathcal{D}_j)]^{-1}$$

This is the **projection** of the edge states of  $\mathcal{D}_j$  onto those of  $\mathcal{D}_d$ .

shape-based control

Recall:

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$$(R1): \|x_i(t) - x_j(t)\| \rightarrow d_{ij} \text{ for all } ij \in E_f$$

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The goal of formation control:

(R1):  $\|x_i(t) - x_j(t)\| \rightarrow d_{ij}$  for all  $ij \in E_f$

(R2): If the interaction graph  $\mathcal{G}(t)$  is dynamic, then it should converge to  $E_f \subseteq \mathcal{E}$  for all  $t \geq T$ , for finite  $T$ .

the static graph case

## Shape-based control – The static graph case

---

Suppose (R2) is satisfied: all pairs of target relative states  $ij \in E_j$  are also communicating, in that  $i$  and  $j$  exchange information ( $ij \in \mathcal{E}$ )

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By reaching agreement over the  $\tau_i$ 's, each agent would satisfy  $x_i - \xi_i = \tau$ , and so  $\|x_i - x_j\| = d_{ij}$  as required by (R1)

We can let each agent  $i$  perform consensus on  $\tau_i$  with its neighbours in  $E_f$ :

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Since we assume that (R2) is satisfied,  $E_f \subseteq \mathcal{E}$ , and so this is a feasible communication strategy.

Note that  $\dot{\tau}_i = \dot{x}_i$ , and

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Therefore, we can write

$$\dot{x}_i(t) = \dot{\tau}_i(t) = - \sum_{j \in N_f(i)} (\tau_i(t) - \tau_j(t)) = - \sum_{j \in N_f(t)} (x_i(t) - x_j(t)) - (\xi_i - \xi_j)$$

### Theorem

Consider the protocol,

$$\dot{x}_i(t) = - \sum_{j \in N_f(t)} (x_i(t) - x_j(t)) - (\xi_i - \xi_j)$$

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$$x_i(t) - \xi_i \rightarrow \tau, \quad \forall i \in \mathcal{N}.$$

as  $t \rightarrow \infty$ .

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(Way harder!)

Suppose  $\mathcal{G}$  is dynamic, i.e. the edge set is a function of time:  $\mathcal{E} = \mathcal{E}(t)$ .

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### Corollary

Given a connected, target formation graph  $\mathcal{G}_f = (\mathcal{V}, E_f)$ , the above protocol will asymptotically drive all agents to a constant displacement of the target positions if for all  $t \geq 0, E_f \subseteq \mathcal{E}(t)$ .

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There are many results on graphs that vary randomly over time, but requires complicated machinery to analyze

relative state-based control

Let's start with a feedback control design procedure.

Suppose each agent  $i \in \mathcal{N}$  is a single integrator  $\dot{x}_i = u_i$ .

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The formation error at time  $t$  is defined as,

$$e(t) = z_{\text{ref}} - z(t).$$

We can write the dynamics of the error as,

$$\dot{e}(t) = -D(\mathcal{D})^T u(t).$$

*Proof:*

A common feedback design is to integrate proportional to the error:

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Then, we can write,

$$\dot{e} = -kL_e(\mathcal{D})e(t)$$

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The edge Laplacian is positive-definite for a spanning tree, and so

$$\lim_{t \rightarrow \infty} e(t) = 0,$$

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and so the  $n$  single integrators achieve  $z_{\text{ref}}$  if  $k > 0$  as  $t \rightarrow \infty$ .

## Relative state-based control – Dynamics

---

What do the dynamics of  $x$  look like under this control?

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## Lemma

Under the control law  $u(t) = kD(\mathcal{D})e(t)$ , the state dynamics are,

$$\dot{x}(t) = -kL(\tilde{\mathcal{D}})x(t) + kD(\mathcal{D})z_{\text{ref}},$$

where  $\tilde{\mathcal{D}}$  is the undirected version of the spanning tree.

*Proof:*

Why does the Laplacian use the undirected version of  $\mathcal{D}$ ?

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relative state-based control –  
double-integrator case

## Relative state-based control – Double Integrator

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We further assume that  $\ddot{z}_{\text{ref}} = 0$ , i.e. the reference velocity does not change with time.

## Relative state-based control – Double Integrator

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Again, we set the error as  $e(t) = z_{\text{ref}} - D(\mathcal{D})^T x(t)$ .

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We can show that,

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*Proof:*

Define the (proportional-derivative) feedback controller,

$$u(t) = k \begin{bmatrix} D(\mathcal{D}) & D(\mathcal{D}) \end{bmatrix} \begin{bmatrix} e(t) \\ \dot{e}(t) \end{bmatrix}.$$

Define the (proportional-derivative) feedback controller,

$$u(t) = k \begin{bmatrix} D(\mathcal{D}) & D(\mathcal{D}) \end{bmatrix} \begin{bmatrix} e(t) \\ \dot{e}(t) \end{bmatrix}.$$

Then, the closed-loop system is

$$\begin{bmatrix} \dot{e}(t) \\ \ddot{e}(t) \end{bmatrix} = \begin{bmatrix} 0 & I \\ -kL_e(\mathcal{D}) & -kL_e(\mathcal{D}) \end{bmatrix} \begin{bmatrix} e(t) \\ \dot{e}(t) \end{bmatrix}$$

## Relative state-based control – Closed-loop system

*Proof:*

## Stability of the closed-loop system

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## Stability of the closed-loop system

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*continued:*

## Theorem

$$\begin{bmatrix} \dot{x}(t) \\ \ddot{x}(t) \end{bmatrix} = \begin{bmatrix} 0 & I \\ -kL(\tilde{\mathcal{D}}) & -kL(\tilde{\mathcal{D}}) \end{bmatrix} \begin{bmatrix} x(t) \\ \dot{x}(t) \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ kD(\mathcal{D}) & kD(\mathcal{D}) \end{bmatrix} \begin{bmatrix} z_{\text{ref}}(t) \\ \dot{z}_{\text{ref}} \end{bmatrix}$$

# Conclusion

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1. Shape-based control (Take points  $\Xi$  and control towards them)

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1. Shape-based control (Take points  $\Xi$  and control towards them)
2. Relative state-based control (Define relative states  $x_i - x_j$  and control them so  $\|x_i - x_j\| = d_{ij}$ .)
3. Single/double integrator versions

## Homework: Extensions of these ideas

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1. Internal dynamics:  $\dot{x}_i = ax(t) + bu(t)$

## Homework: Extensions of these ideas

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1. Internal dynamics:  $\dot{x}_i = ax(t) + bu(t)$
2. Unicycle control:  $r_i(t) = x_i(t) + jy_i(t)$

## Reading Assignment

- Mesbahi & Egerstedt 6.2-6.6

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- Mesbahi & Egerstedt 6.2-6.6
- Mesbahi & Egerstedt 7.1-7.4

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## Exercise Session

Coverage Control with **Alberto**

