Bidirectional Typing Rules for Mini

Aidan Ewart

November 11, 2020

Basically the same as [1], except bidirectional. The reason we still have rigidity on checking is that certain operations (like APP) check non-rigidly.

$$\Gamma \vdash t \Rightarrow^m \tau$$

$$\frac{V_{\text{AR}}}{v : \forall \bar{a}.\tau \in \Gamma} \qquad \qquad \frac{A_{\text{PP}}}{\Gamma \vdash v \Rightarrow^{m} [\bar{a} \mapsto \bar{v}]\tau} \qquad \qquad \frac{\Gamma \vdash f \Rightarrow^{w} \tau_{1} \to \tau_{2}}{\Gamma \vdash (f \ x) \Rightarrow^{m} \tau_{2}}$$

$$\frac{\text{Abs-Infer}}{\Gamma, x :^m \tau_1 \vdash t \Leftarrow^m \tau_2} \frac{\Gamma, x :^m \tau_1 \vdash t \Leftarrow^m \tau_2}{\Gamma \vdash (\text{lam } x \ t) \Rightarrow^m \tau_1 \rightarrow \tau_2}$$

Let-Infer

$$\frac{\Gamma \vdash u \upharpoonright^{m_1}}{\Gamma \vdash u \Rightarrow^w \tau_1} \frac{\Gamma \vdash u \upharpoonright^{m_1}}{\bar{a} = \operatorname{ftv}(\tau_1) - \operatorname{ftv}(\Gamma)} \frac{\Gamma, x :^{m_1} \forall \bar{a}.\tau_1 \vdash t \Rightarrow^{m_2} \tau_2}{\Gamma \vdash (\operatorname{let} (x \ u) \ t) \Rightarrow^{m_2} \tau_2}$$

$$\frac{ \begin{array}{c} \text{Fix-}\overline{\text{Infer}} \\ \Gamma \vdash \overline{(x \ u)} \mapsto \Delta & \Gamma, \Delta \vdash t \Rightarrow^{m_2} \tau_2 \\ \hline \Gamma \vdash (\text{fix} \ \overline{(x \ u)} \ t) \Rightarrow^{m_2} \tau_2 \end{array} \qquad \frac{ \begin{array}{c} A_{\text{NN}} \\ \overline{a} = \text{ftv}(\tau) - \text{ftv}(\Gamma) & \Gamma, \overline{a} \vdash x \Leftarrow^r \tau \\ \hline \Gamma \vdash (x \colon : \tau) \Rightarrow^m \tau \end{array} }$$

CONS

$$C: \forall \bar{a}.\bar{\tau_1} \to \bar{\tau_2} \in \Gamma$$
 $\theta = [\overline{a \mapsto v}] \text{ s.t. } \overline{\Gamma \vdash \theta(v) \Leftarrow^r \tau_1}$
 $\Gamma \vdash C\bar{v} \Rightarrow^m \theta(\tau_2)$

PRIMOP
$$\frac{O: \bar{\tau}_1 \to \tau_2}{\Gamma \vdash O\bar{v} \Rightarrow^m \tau_2} \qquad \theta = [\overline{a} \mapsto \overline{v}] \text{ s.t. } \Gamma \vdash \theta(v) \Leftarrow^r \tau_1$$

$$\Gamma \vdash t \Leftarrow^m \tau$$

$$\frac{\text{Abs-Check}}{\Gamma, x :^m \tau_1 \vdash t \Leftarrow^m \tau_2} \frac{\Gamma, x :^m \tau_1 \vdash t \Leftarrow^m \tau_2}{\Gamma \vdash (\texttt{lam } x \ t) \Leftarrow^m \tau_1 \rightarrow \tau_2}$$

Let-Check

$$\frac{\Gamma \vdash u \upharpoonright^{m_1}}{\Gamma \vdash u \Rightarrow^w \tau_1} \frac{\bar{a} = \operatorname{ftv}(\tau_1) - \operatorname{ftv}(\Gamma)}{\Gamma \vdash (\operatorname{let} (x \ u) \ t) \Leftarrow^{m_2} \tau_2}$$

FIX-CHECK
$$\frac{\Gamma \vdash \overline{(x\ u)} \mapsto \Delta \qquad \Gamma, \Delta \vdash t \Leftarrow^{m_2} \tau_2}{\Gamma \vdash (\text{fix } \overline{(x\ u)} \ t) \Leftarrow^{m_2} \tau_2}$$

MATCH-CHECK

$$\frac{\Gamma \vdash u \mid^{m_p} \quad \Gamma \vdash u \Rightarrow^w \tau_p \quad \Gamma \vdash \overline{p,t} \Leftarrow^{\langle m_p,m_t \rangle} \tau_p, \tau_t}{\Gamma \vdash (\mathsf{match} \ u \ \overline{(p \rightarrow t)}) \Leftarrow^{m_t} \tau_t} \qquad \frac{\Gamma \vdash t \Rightarrow^m \tau}{\Gamma \vdash t \Leftarrow^m \tau}$$

CHECK-INFER
$$\frac{\Gamma \vdash t \Rightarrow^m \tau}{\Gamma \vdash t \Leftarrow^m \tau}$$

$$\Gamma \vdash p, t \Rightarrow^{\langle m_t, \tau_p \rangle} \tau_t$$

$$\frac{P\text{Con-Infer}}{C: \forall \bar{a}.\bar{\tau_1} \rightarrow P\bar{\tau_v} \in \Gamma} \frac{\Gamma, \overline{v:w\ \tau_v} \vdash t \Rightarrow^{m_t} \tau_t}{\Gamma \vdash C\bar{v}, t \Rightarrow^{\langle m_t, P\bar{\tau_v} \rangle} \tau_t} \qquad \frac{P\text{Lit-Infer}}{L: \tau_p \in \Gamma} \frac{\Gamma \vdash t \Rightarrow^{m_t} \tau_t}{\Gamma \vdash L, t \Rightarrow^{\langle m_t, \tau_p \rangle} \tau_t}$$

PLIT-INFER
$$L: \tau_p \in \Gamma \qquad \Gamma \vdash t \Rightarrow^{m_t} \tau_t$$

$$\Gamma \vdash L, t \Rightarrow^{\langle m_t, \tau_p \rangle} \tau_t$$

$$\frac{\Gamma \vdash t \Rightarrow^{m_t} \tau_t}{\Gamma \vdash_{\cdot} t \Rightarrow^{\langle m_t, \tau_p \rangle} \tau_t}$$

$$\Gamma \vdash p, t \Leftarrow^{\langle m_p, m_t \rangle} \tau_p, \tau_t$$

$$\frac{C : \forall \bar{a}.\bar{\tau_1} \to P\bar{\tau_2} \in \Gamma \qquad \Gamma, \overline{v : w \tau_v} \vdash t \Leftarrow^m \tau_t}{\Gamma \vdash C\bar{v}, t \Leftarrow^{\langle w, m \rangle} P\bar{\tau_v}, \tau_t}$$

PCon-R

$$\frac{C: \forall \bar{a}.\bar{\tau_1} \to P\bar{\tau_2} \in \Gamma \qquad \theta = \mathrm{mgu}(\bar{\tau_p}, \bar{\tau_2}) \qquad \Gamma, \overline{v:^r \theta(\tau_v)} \vdash \theta(t) \Leftarrow^m \theta(\tau_t)}{\Gamma \vdash C\bar{v}, t \Leftarrow^{\langle r, m \rangle} P\bar{\tau_v}, \tau_t}$$

$$\frac{\text{PLIT-CHECK}}{L: \tau_p \in \Gamma} \frac{\Gamma \vdash t \Leftarrow^{m_t} \tau_t}{\Gamma \vdash L, t \Leftarrow^{\langle m_p, m_t \rangle} \tau_p, \tau_t} \qquad \frac{\text{PWILD-CHECK}}{\Gamma \vdash t \Leftarrow^{m_t} \tau_t} \frac{\Gamma \vdash t \Leftarrow^{m_t} \tau_t}{\Gamma \vdash \cdot, t \Leftarrow^{\langle m_p, m_t \rangle} \tau_p, \tau_t}$$

$$\frac{\Gamma \vdash t \Leftarrow^{m_t} \tau_t}{\Gamma \vdash \cdot \cdot \cdot t \Leftarrow^{\langle m_p, m_t \rangle} \tau_p, \tau_t}$$

$$\Gamma \vdash \overline{(x \ u)} \mapsto \Delta$$

FIXDEFS

$$\frac{a_1 = \operatorname{ftv}(\tau_3) - \operatorname{ftv}(\Gamma)}{a_2 = \operatorname{ftv}(\tau_1) - \operatorname{ftv}(\Gamma)} \frac{\Gamma, x_1 :^{m_1} \forall a_1.\tau_3, \overline{u_2}, \overline{u_2} :^{m_2} \tau_1}{\Delta = x_1 :^{m_1} \forall a_1.\tau_3, \overline{x_2} :^{m_2} \tau_1} \vdash \overline{u_1 \Leftarrow^w \forall a_1.\tau_3, \overline{u_2} \Leftarrow^w \tau_1} \\ \frac{a_2 = \operatorname{ftv}(\tau_1) - \operatorname{ftv}(\Gamma)}{\Gamma \vdash \overline{(x_1 \ (u_1 : : \tau_3)), \overline{(x_2 \ u_2)}} \mapsto \Delta}$$

$$\begin{array}{c|c} \hline \Gamma \vdash t \mid^m \\ \hline \\ SCR-VAR \\ \underline{v \mid^m \in \Gamma} \\ \overline{\Gamma \vdash v \mid^m} \end{array} \quad \begin{array}{c|c} SCR-APP \\ \underline{\Gamma \vdash f \mid^{m_1}} \quad \Gamma \vdash x \mid^{m_2} \\ \hline \Gamma \vdash (f \ x) \mid^{\lfloor m_1, m_2 \rfloor} \end{array} \quad \begin{array}{c|c} SCR-LET \\ \underline{\Gamma \vdash u \mid^{m_1}} \quad \Gamma, x \mid^{m_1} \vdash t \mid^{m_2} \\ \hline \Gamma \vdash (let \ (x \ u) \ t) \mid^{m_2} \\ \hline \\ SCR-ABS \\ \underline{\Gamma, x \mid^{w} \vdash t \mid^{m_1}} \\ \overline{\Gamma \vdash (lam \ x \ t) \mid^{m_1}} \end{array} \quad \begin{array}{c|c} SCR-ANN \\ \overline{\Gamma \vdash (x : : \tau) \mid^{\tau}} \end{array} \quad \begin{array}{c|c} SCR-OTHER \\ \overline{\Gamma \vdash t \mid^{w}} \\ \hline \end{array}$$

References

[1] Simon Peyton Jones, Dimitrios Vytiniotis, and Stephanie Weirich. Simple unification-based type inference for GADTs. Jan. 2005.